

# A FAMILY OF 4-POINTS DYADIC HIGH RESOLUTION SUBDIVISION SCHEMES

**ABSTRACT.** By using temporary placeholders on a dense grid, we generalize the 4-point dyadic cubic Deslauriers-Dubuc scheme. Interpolated values require 2 steps to stabilize as they are first interpolated on a coarse scale through a tetradic filter and then on a finer scale using a dyadic filter. The interpolants are  $C^1$  and can be chosen to reproduce polynomials of degree 4. These generalized interpolatory subdivision schemes have minimal support and no additional memory requirement. This work has applications in CAGD and wavelet theory.

## 1. INTRODUCTION

Interpolatory subdivision schemes interpolate a discrete set of data points in a local manner, that is, the value of the interpolation function at a given point depends on a small number of nearby data points. The classical dyadic algorithm introduced by Deslauriers and Dubuc [4, 2] finds the midpoint values by fitting a Lagrange polynomial through the  $2N$  closest data points. By repeating this algorithm again and again, each time doubling the number of data points or nodes, we eventually have a dense set of data points and we can determine uniquely a smooth interpolation function. They also inherit their order of approximation from the corresponding Lagrange interpolation scheme. Because interpolatory subdivision schemes relate data points from one scale to the data points at another scale, it is not surprising that they are a key ingredient in the construction of compactly supported wavelets [1, 3].

More recently, Merrien [9, 10, 5] introduced Hermite subdivision schemes. Since Merrien subdivision schemes use Hermite nodes, they have twice the approximation order and better regularity for a given support. For example 2-points Hermite schemes are differentiable and can have up to cubic order whereas the corresponding Deslauriers-Dubuc scheme (the linear spline) isn't differentiable and only has linear order.

Since Deslauriers-Dubuc schemes have important applications, it is tempting to add extra nodes to Deslauriers and Dubuc schemes as an attempt to improve them to get "high resolution" schemes. One might hope to retain the useful properties such as the approximation order and the regularity while making the scheme more local for example. Doubling the number of nodes however is costly (effectively doubling the memory requirements), however, since a dyadic subdivision scheme doubles its memory usage at each step, we can choose to use right away this upcoming extra storage space without any cost. In effect, we can simply make use of the memory we will allocate later in any case. Therefore, we can double (or more) the number of nodes. These new placeholders can then be used to record a first coarse scale guess (using a tetradic filter) which we can later combine with a finer scale interpolation (using a dyadic filter). As a special case, we may choose to ignore the coarse scale estimate, in which case our approach amounts to a Deslauriers-Dubuc scheme; we can also use this approach to reproduce polynomials of degree 4 by a Richardson extrapolation approach. This paper shows that by summing up the tetradic (coarse) interpolation recorded in placeholders and dyadic (fine) interpolations, we get a range of smooth ( $C^1$ ) high resolutions schemes reproducing cubic polynomials.

## 2. SUBDIVISION SCHEMES

Interpolatory subdivision schemes were first introduced by Deslauriers and Dubuc (quote). Let  $b > 1$  be an integer, given two integers  $k, j$ , the number  $x_{j,k} = k/b^j$  is said to be  $b$ -adic (of depth  $j$ ). For a fixed  $j$ , the  $b$ -adic numbers form a regularly-spaced set of nodes. Given some corresponding data  $\{y_{J,k}\}_{k \in \mathbb{N}}$  on the dyadic numbers of depth  $J$ , we want to build a smooth function  $f$  such that  $f(x_{J,k}) = y_{J,k} \forall k \in \mathbb{N}$ . Starting with this initial data  $(y_{J,k})$  and using the linear formula

$$(2.1) \quad y_{j+1,l} = \sum_{k \in \mathbb{N}} \gamma_{bk-l} y_{j,k}$$

for some constant array  $\gamma$ , we get values  $y_{j,k}$  for any  $j > J$  and since  $b$ -adic numbers form a dense set of the real numbers, there is at most one continuous function such that  $f(x_{j,k}) = y_{j,k}$  for all  $k, j > J$ .

A subdivision scheme is interpolatory and will satisfy  $f(x_{J,k}) = y_{J,k}$  if  $\gamma_{bk} = 0$  except when  $k = 0$ . We say that a subdivision scheme is stationary if the array  $\gamma$  is constant (doesn't depend on  $j$ ). An interpolatory subdivision scheme is said to be  $2N$ -points if  $\gamma_l = 0$  for  $|l| > Nb$ . The interpolation function  $f$  computed from a  $2N$ -points  $b$ -adic

scheme with initial data  $y_{0,0} = 1$  and  $y_{0,k} = 0$  for all  $k \neq 0$  is said to be the fundamental function and has a compact support of  $[-(Nb-1)/(b-1), (Nb-1)/(b-1)]$  or  $[1-2N, 2N-1]$  when  $b = 2$ . Hence as  $N$  increases the support of the fundamental function increases.

For each  $N = 1, 2, 3, \dots$  there exists a corresponding an interpolatory Deslauriers-Dubuc subdivision scheme and they are built from the midpoint evaluation of Lagrange polynomial of degree  $2N-1$  on  $2N$  points. For  $b = 2$  (dyadic case), the 4-points Deslauriers-Dubuc scheme can be defined by the array  $\gamma^{DD2}$  given by  $\gamma_0^{DD2} = 1, \gamma_1^{DD2} = \gamma_{-1}^{DD2} = -9/16, \gamma_3^{DD2} = \gamma_{-3}^{DD2} = -1/16$  with  $\gamma_k^{DD2} = 0$  otherwise; for  $b = 4$  (tetradic case), the scheme is defined by  $\gamma^{DD4}$  given by  $\gamma_{2k}^{DD4} = \gamma_k^{DD2} \forall k \in \mathbb{N}$ ,  $\gamma_{-1}^{DD4} = \gamma_1^{DD4} = 105/128, \gamma_{-3}^{DD4} = \gamma_3^{DD4} = 35/128, \gamma_{-5}^{DD4} = \gamma_5^{DD4} = -7/128, \gamma_{-7}^{DD4} = \gamma_7^{DD4} = -5/128$ , with  $\gamma_k^{DD2} = 0$  otherwise. Because 4-points Deslauriers-Dubuc schemes are derived from cubic Lagrange polynomials, they reproduces cubic polynomials, that is, if the initial data  $y_{j,k}$  satisfies  $y_{j,k} = p(x_{j,k}) \forall k \in \mathbb{N}$  for some cubic polynomial  $p$  then the interpolation function  $f$  is this same cubic polynomial  $f = p$ . Hence the two cases presented above ( $\gamma^{DD2}$  and  $\gamma^{DD4}$ ) reproduce cubic polynomials. It can also be shown also that they both give differentiable ( $C^1$ ) interpolation functions.

We are interested in measuring how well a given subdivision scheme can approximation functions. One such measure is given by the approximation order of the scheme [8, definition 2]. We say that a subdivision scheme has approximation order  $p$  if given given any smooth function  $g \in C^p([0, 1])$ , the interpolation function  $f$  satisfying  $f(x_{j,k}) = g(x_{j,k}) \forall k \in \mathbb{N}$  is such that  $\|f - g\|_{L^\infty([0, 1])} \leq C/2^{jp}$  for a constant  $C$  independent of  $j$ .

For a continuous subdivision scheme reproducing polynomials of degree  $p$ , it is sufficient for the scheme to converge to a continuous function to have approximation order  $p+1$  [8]. Specifically, this means that 4-points Deslauriers-Dubuc schemes have approximation order 4.

### 3. HIGH RESOLUTION SUBDIVISION SCHEMES

**3.1. Definitions.** An more general alternative to equation 2.1 is given by the linear equation

$$y_{j+1,l} = \sum_{m=1}^M \sum_{k \in \mathbb{N}} \gamma_{Nbkm+m-1-l}^{(m)} y_{j,Nk+m-1}$$

where  $\gamma^{(1)}, \dots, \gamma^{(M)}$  are constant arrays (independent from  $j$ ). Because this new formula uses  $M$  times the usual number of nodes (see equation 2.1), we say that it is a “high resolution subdivision scheme” however it can still be said to be  $b$ -adic because the number of nodes is increasing by  $b$ . As a special case, when  $M = 1$ , we have a usual subdivision scheme.

In what follows, we set  $b = N = 2$  and consider the equation

$$(3.1) \quad y_{j+1,l} = \sum_{k \in \mathbb{N}} \gamma_{4k-l}^{(1)} y_{j,2k} + \gamma_{4k+1-l}^{(2)} y_{j,2k+1}$$

with

$$(3.2) \quad \gamma_k^{(1)} = \gamma_k^{DD4} - \alpha \frac{((-1)^k + 1)}{2} \gamma_{k/2}^{DD2} + \alpha \delta_{k,0} \forall k \in \mathbb{N}$$

$$(3.3) \quad \gamma_0^{(2)} = \alpha, \gamma_k^{(2)} = 0 \text{ otherwise}$$

for some parameter  $\alpha \in \mathbb{R}$ . As we will see, this choice is made so that the scheme can reproduce cubic polynomials for all  $\alpha$ . The odd terms  $\{y_{j,2k+1}\}_{k \in \mathbb{N}}$  will be referred as placeholders because their assigned value will change in general. In the simplest case,  $\alpha = 0$ , equation 3.1 becomes

$$(3.4) \quad y_{j+1,l} = \sum_{k \in \mathbb{N}} \gamma_{4k-l}^{DD4} y_{j,2k}.$$

Because  $\gamma^{(2)} = 0$  in this case, we can see that the placeholders are effectively ignored. Indeed, we observe that this last equation discards odd terms at each step:  $y_{j+1,l}$  depends only on terms of the form  $y_{j,2k}$  (even terms) and not at all on the odd terms  $y_{j,2k+1}$ . Hence, we can replace equation 3.4 by

$$y_{j+1,2l} = \sum_{k \in \mathbb{N}} \gamma_{4k-2l}^{DD4} y_{j,2k}$$

but because  $\gamma_{2k}^{DD4} = \gamma_k^{DD2}$ , this last equation becomes  $y_{j+1,2l} = \gamma_{2k-l}^{DD2} y_{j,2k}$  and if we define  $\tilde{y}_{j,k} = y_{j,2k}$  then

$$(3.5) \quad \tilde{y}_{j+1,2l} = \sum_{k \in \mathbb{N}} \gamma_{2k-l}^{DD2} \tilde{y}_{j,2k}$$

which we recognize as the cubic Deslauriers-Dubuc scheme.

**Proposition 3.1.** *For  $\alpha = 0$ , the high resolution scheme given by equations 3.1, 3.2, and 3.3 is equivalent to the 4-points dyadic Deslauriers-Dubuc subdivision scheme (discarding the odd nodes or placeholders in the first iteration).*

In general, since  $\gamma_{2k}^{DD4} = \gamma_k^{DD2}$ , we can rewrite equation 3.1 for even and odd terms. Firstly, setting  $l = 2s$  ( $l$  even), we have

$$\begin{aligned} y_{j+1,2s} &= \sum_{k \in \mathbb{N}} \gamma_{4k-2s}^{(1)} y_{j,2k} + \gamma_{4k+1-2s}^{(2)} y_{j,2k+1} \\ &= \sum_{k \in \mathbb{N}} (\gamma_{4k-2s}^{DD4} - \alpha \gamma_{2k-s}^{DD2} + \alpha \delta_{4k,2s}) y_{j,2k} + \alpha \delta_{4k+1,s} y_{j,2k+1} \\ &= \sum_{k \in \mathbb{N}} ((1-\alpha) \gamma_{2k-s}^{DD2} + \alpha \delta_{2k,s}) y_{j,2k} + \delta_{2k+1,s+1} \alpha y_{j,2k+1} \end{aligned}$$

so that when  $s$  is even ( $l = 2s = 4r$ ), we have the interpolatory condition

$$(3.6) \quad y_{j+1,4r} = y_{j,2r}$$

otherwise, when  $s$  is odd ( $l = 2s = 4r + 2$ )

$$(3.7) \quad y_{j+1,4r+2} = \alpha y_{j,2r+1} + (1-\alpha) \sum_{k \in \mathbb{N}} \gamma_{2k-2r-1}^{DD2} y_{j,2k}.$$

Secondly, if  $l$  is odd ( $l = 2s + 1$ ), we have

$$\begin{aligned} y_{j+1,2s+1} &= \sum_{k \in \mathbb{N}} \gamma_{4k-2s-1}^{(1)} y_{j,2k} + \gamma_{4k-2s-1}^{(2)} y_{j,2k+1} \\ (3.8) \quad &= \sum_{k \in \mathbb{N}} \gamma_{4k-2s-1}^{DD4} y_{j,2k}. \end{aligned}$$

Equations 3.6, 3.7, and 3.8 can be used to describe the chosen high resolution schemes: while equation 3.6 is the interpolatory condition, equation 3.8 fills the placeholders with tetradic (coarse scale) interpolated values whereas equation 3.7 combines the value stored in the placeholder with the newly available interpolated value (fine scale) given by the summation term which we recognize from the dyadic Deslauriers-Dubuc interpolation.

**3.2. Reproduced polynomials.** Assume that for some  $j$ ,  $y_{j,k} = p_3(x_{j,k})$  for some cubic polynomial  $p_3$ . We then have that

$$\sum_{k \in \mathbb{N}} \gamma_{2k-2r-1}^{DD2} y_{j,2k} = y_{j,2r+1} = p_3(x_{j,2r+1})$$

and thus equation 3.7 becomes  $y_{j+1,4r+2} = p_3(x_{j,2r+1})$ . Moreover, equation 3.8 implies  $y_{j+1,2s+1} = p_3(x_{j+1,2s+1})$ . We can conclude that  $y_{j+1,k} = p_3(x_{j+1,k})$  if  $y_{j,k} = p_3(x_{j,k})$ . For practical implementations of a high subdivision scheme, it is necessary to first apply a one-step subdivision scheme since in general, we don't have placeholder values precomputed. This can be achieved with a dyadic Deslauriers-Dubuc filter. Let  $\{y_{j,k}\}_k$  be some initial data. As a first step, we apply equation

$$(3.9) \quad y_{j+1,2l} = \sum_{k \in \mathbb{N}} \gamma_{2k-l}^{DD2} y_{j,2k}$$

followed by equation 3.1 with  $j = j + 1$ , and so on. By induction, we get the following lemma.

**Lemma 3.2.** *The high resolution scheme given by equations 3.1, 3.2, and 3.3 (using 3.9 as an initialization step) reproduces cubic polynomials.*

We can also get a stronger result by choosing  $\alpha$ . Let  $p_4(x) = a_4 x^4 + p_3(x)$  where  $p_3$  is some cubic polynomial. Suppose that for some  $j$ ,  $y_{j,2k} = p_4(x_{j,2k})$  (and accordingly  $y_{j-1,2k} = p_4(x_{j-1,2k})$ ). Substituting equation 3.8 into 3.7, we get

$$\begin{aligned} y_{j+1,4r+2} &= \alpha y_{j,2r+1} + (1-\alpha) \sum_{k \in \mathbb{N}} \gamma_{2k-2r-1}^{DD2} y_{j,2k} \\ (3.10) \quad &= \alpha \sum_{k \in \mathbb{N}} \gamma_{4k-2r-1}^{DD4} y_{j-1,2k} + (1-\alpha) \sum_{k \in \mathbb{N}} \gamma_{2k-2r-1}^{DD2} y_{j,2k}. \end{aligned}$$

We can compute each of the sums explicitly

$$\begin{aligned} \sum_{k \in \mathbb{N}} \gamma_{2k-2r-1}^{DD2} y_{j,2k} &= p_3(x_{j,2r+1}) + a_4 \sum_{k \in \mathbb{N}} \gamma_{2k-2r-1}^{DD2} (x_{j,2k})^4 \\ &= p_4(x_{j,2r+1}) - \frac{9a_4}{16 \times 2^4} \end{aligned}$$

and

$$\begin{aligned} \sum_{k \in \mathbb{N}} \gamma_{4k-2r-1}^{DD4} y_{j-1,2k} &= p_3(x_{j,2r+1}) + a_4 \sum_{k \in \mathbb{N}} \gamma_{4k-2s-1}^{DD4} (x_{j-1,2k})^4 \\ &= p_4(x_{j,2r+1}) - \frac{105a_4}{16 \times 2^{4j}}. \end{aligned}$$

Hence, by substituting these two results (fine and coarse scales interpolation) and setting  $\alpha = -3/35$ , we get

$$y_{j+1,4r+2} = p_4(x_{j,2r+1}) = p_4(x_{j+1,4r+2})$$

since for  $\alpha = -3/35$ ,  $105\alpha + 9\alpha = 0$ . Therefore, the scheme reproduces polynomials of degree 4. Of course, this result assumes that we initialize the data so that  $y_{j,2k+1} = p_4(x_{j,2k+1}) - \frac{105a_4}{16 \times 2^{4j}}$  and  $y_{j,2k} = p_4(x_{j,2k})$  for all  $k \in \mathbb{N}$ . We can get this result naturally by having  $y_{j-1,k} = p_4(x_{j,k})$  and applying first the high resolution scheme (equation 3.1) with  $\alpha = 1$  so that equations 3.6 and 3.7 will guarantee  $y_{j,2k} = p_4(x_{j,k})$  whereas equation 3.8 will initialize the placeholders properly.

**Proposition 3.3.** *For any given  $j$ , if  $y_{j-1,k} = p_4(x_{j-1,k})$  where  $p_4$  is a polynomial of degree 4 then applying the high resolution scheme given by equations 3.1, 3.2, and 3.3 first with  $\alpha = 1$  as an initialization step and then with  $\alpha = -3/35$  will guarantee that  $y_{j',2k} = p_4(x_{j',2k})$  for  $\forall k \in \mathbb{N}$  and all  $j' \geq j-1$ .*

**3.3. Sufficient conditions for regularity.** To study the regularity of high resolution schemes, it is convenient rewrite formula 3.1 in terms of (trigonometric) polynomials. Given some data  $y_{j,k}$ , define  $P^j(z) = \sum_{k \in \mathbb{N}} y_{j,k} z^k$ . If  $P_2(z) = \sum_{k \in \mathbb{N}} \gamma_k^{DD2} z^k$ , then equation 3.9, can be rewritten  $P^{j+1}(z) = P_2(z)P^j(z^2)$ . Similarly, if  $P_4(z) = \sum_{k \in \mathbb{N}} \gamma_k^{DD4} z^k$ , then the tetradic subdivision scheme is given by  $P^{j+1}(z) = P_4(z)P^j(z^2)$ . It can be shown that we can rewrite the general equation for high resolution subdivision schemes as

$$P^{j+1}(z) = \sum_{i=1}^M \Gamma_i(z) P^j(e^{2\pi i/b} z^b).$$

where  $\Gamma_i$  must be Laurent polynomials. Using symbols, equation 3.1 becomes

$$\begin{aligned} P^{j+1}(z) &= \{P_4(z) - \alpha P_2(z^2) + \alpha\} \left( \frac{P^j(z^2) + P^j(-z^2)}{2} \right) + \alpha \left( \frac{P^j(z^2) - P^j(-z^2)}{2} \right) \\ &= \left\{ \frac{P_4(z) - \alpha P_2(z^2)}{2} + \alpha \right\} P^j(z^2) + \frac{P_4(z) - \alpha P_2(z^2)}{2} P^j(-z^2) \\ (3.11) \quad &= \Gamma_1(z) P^j(z^2) + \Gamma_2(z) P^j(-z^2) \end{aligned}$$

Because  $\gamma_{2k}^{DD4} = \gamma_k^{DD2} \forall k \in \mathbb{N}$ , we observe that

$$P_4(z) - \alpha P_2(z^2) = \frac{P_4(z) - P_4(-z)}{2} + (1 - \alpha) \frac{P_4(z) + P_4(-z)}{2}$$

and thus

$$\begin{aligned} \Gamma_1(z) &= \Gamma_2(z) + \alpha \\ \Gamma_2(z) &= \frac{P_4(z) - P_4(-z)}{4} + (1 - \alpha) \frac{P_4(z) + P_4(-z)}{4}. \end{aligned}$$

When  $\alpha = 0$ ,  $\Gamma_1(z) = \Gamma_2(z) = \frac{P_4(z)}{2}$  and equation 3.1 becomes

$$P^{j+1}(z) = P_4(z) \left( \frac{P^j(z^2) + P^j(-z^2)}{2} \right)$$

and it can be shown to be equivalent to the Deslauriers-Dubuc dyadic scheme by writing

$$\begin{aligned} \frac{P^{j+1}(z) + P^{j+1}(-z)}{2} &= \left( \frac{P_4(z) - P_4(-z)}{2} \right) \left( \frac{P^j(z^2) + P^j(-z^2)}{2} \right) \\ &= P_2(z^2) \left( \frac{P^j(z^2) + P^j(-z^2)}{2} \right). \end{aligned}$$

which gives an alternate proof for proposition 3.1.

In order to study the regularity and stability of the chosen high resolution schemes, we need to find corresponding schemes for the (forward) finite differences

$$\frac{dy_{j,k}}{dx_{j,k}} = \frac{y_{j,k+1} - y_{j,k}}{1/2^j}$$

and of the correspondin higher order finite differences which can be defined recursively  $d^{(n)}y_{j,k} = d \left( d^{(n-1)}y_{j,k} \right) / (dx_{j,k})^n$ . Let

$$\begin{aligned} H_1^j(z) &= \sum_{k \in \mathbb{N}} 2^j (y_{j,k+1} - y_{j,k}) z^k \\ &= \sum_{k \in \mathbb{N}} 2^j y_{j,k} z^{k-1} - \sum_{k \in \mathbb{N}} 2^j y_{j,k} z^k \\ &= 2(1/z - 1)P^j(z) = 2(1 - z)P^j(z)/z, \end{aligned}$$

then  $P^j(z^2) = z^2 2^j H_1^j(z^2) / (1 - z^2)$ ,  $P^j(-z^2) = -z^2 2^j H_1^j(-z^2) / (1 + z^2)$ , and  $P^{j+1}(z) = z 2^{j+1} H_1^{j+1}(z) / (1 - z)$ . Substituting these equations into 3.11 gives

$$H_1^{j+1}(z) = \frac{2z(1-z)}{(1-z^2)} \Gamma_1(z) H_1^j(z^2) - \frac{2z(1-z)}{(1+z^2)} \Gamma_2(z) H_1^j(-z^2).$$

The higher order finite differences are given by

$$H_n^j(z) = \frac{2(1-z)}{z} H_{n-1}^j(z) = \left( \frac{2(1-z)}{z} \right)^n P^j(z)$$

where we let  $H_0(z) = P(z)$  and it can be seen that higher order finite differences are related through

$$(3.12) \quad H_n^{j+1}(z) = \left( \frac{2z}{1+z} \right)^n \Gamma_1(z) H_n^j(z^2) + \left( \frac{-2z(1-z)}{1+z^2} \right)^n \Gamma_2(z) H_n^j(-z^2).$$

$H_n$  is said to be equivalent to a high resolution subdivision scheme if  $\Gamma_1(z)/(1+z)$  and  $\Gamma_2(z)/(1+z^2)$  are Laurent polynomials. However, we have

$$P_4(z) = \frac{-(1+z)^4 (1+z^2)^4 (5z^2 - 12z + 5)}{128z^7}$$

and therefore  $\Gamma_2(z)/(1+z^2)^n$  is a Laurent polynomial for  $n = 1, 2, 3, 4$ . On the other hand, we can also show that  $\Gamma_1(z)/(1+z)^n$  is a Laurent polynomial for  $n = 1, 2, 3, 4$ . Therefore,  $H_n$  is the symbol of a high resolution subdivision scheme if  $n = 1, 2, 3, 4$ . While  $H_n$  is equivalent to the subdivision scheme of  $d^{(n)}y_{j,k} / (dx_{j,k})^n = 2^{jn} d^{(n)}y_{j,k}$ , we can define  $dH_n^j$  based on  $d \left( 2^{j(n-1)} d^{(n-1)}y_{j,k} \right)$  with, for example,  $dH_1^j(z) = (1-z)P^j(z)/z$  and more generally

$$(3.13) \quad dH_{n-1}^j(z) = \frac{(1-z)}{z} H_{n-1}^j(z) = \frac{2^{j(n-1)}(1-z)^n}{z^n} P^j(z).$$

Replacing  $H_{n-1}$  by  $dH_{n-1}$  in equation 3.12, we find

$$(3.14) \quad 2dH_{n-1}^{j+1}(z) = \left( \frac{2z}{1+z} \right)^n \Gamma_1(z) dH_{n-1}^j(z^2) + \left( \frac{-2z(1-z)}{1+z^2} \right)^n \Gamma_2(z) dH_{n-1}^j(-z^2).$$

$dH_{n-1}$  is said to be equivalent to a high resolution subdivision scheme if  $\Gamma_1(z)/(1+z)^n$  and  $\Gamma_2(z)/(1+z^2)^n$  are Laurent polynomials or if  $H_n$  is equivalent to a high resolution subdivision scheme which is true for  $n = 1, 2, 3, 4$ .

Using results from Dyn [6], we have the following theorem.

**Theorem 3.4.** (Dyn) *If  $dH_n$  as in equations 3.13 and 3.14 is equivalent to a high resolution subdivision scheme converging uniformly to zero for all initial data, then the corresponding scheme  $P$  as in equation 3.11 is  $C^n$ , that is, all interpolation functions  $f$  are  $C^n$*

*Proof.* See theorems 4.2 and 4.4 in [6] as they apply to this high resolution subdivision schemes as well. □

A sufficient condition for  $y_{j,k} \rightarrow 0$  uniformly as  $j \rightarrow \infty$  is that  $\max_{l=0,1} \{ \sum_{k \in \mathbb{N}} |\gamma_{2k-l}| \} < 1$  when  $y_{j+1,l} = \sum_{k \in \mathbb{N}} \gamma_{2k-l} y_{j,k}$ . Indeed, let

$$I_{j,l} = [2^{j+1} - 2 + 2^j(l-1) + l, 2^j(l+1) - 2^{j+1} + 2]$$

writing

$$M_{j,l} = \max \{ |y_{j,k}| : k \in I_{j,l} \}$$

then  $M_{j+1} \leq \max_{l=0,1} \{ \sum_{k \in \mathbb{N}} |\gamma_{2k-l}| \} M_j$ .

**Theorem 3.5.** For  $-7 < 32\alpha < 6$ , the high resolution subdivision scheme given by equation 3.11 are  $C^1$ .

*Proof.* We need to consider  $dH_1$  given by

$$\frac{dH_1^{j+1}(z)}{2} = \left( \frac{z}{1+z} \right)^2 \Gamma_1(z) dH_1^j(z^2) + \left( \frac{-z(1-z)}{1+z^2} \right)^2 \Gamma_2(z) dH_1^j(-z^2).$$

Let

$$dh1(z) = \left( \frac{z}{1+z} \right)^2 \Gamma_1(z) + \left( \frac{-z(1-z)}{1+z^2} \right)^2 \Gamma_2(z),$$

to show that the high resolution subdivision scheme corresponding to  $dH_1$  converges uniformly to zero, it is enough to

show that the sum of the absolute values of the coefficients of the odds and even terms,  $dh1_{even}(z) = \frac{dh1(z) + dh1(-z)}{2}$  and  $dh1_{odd}(z) = \frac{dh1(z) - dh1(-z)}{2}$ , are smaller than  $1/2$ . This is effectively equivalent to requiring that  $2^j |(y_{j,k+2} - y_{j,k+1}) - (y_{j,k+1} - y_{j,k})| = \lambda^j M \rightarrow 0 \forall k \in \mathbb{N}$  as  $j \rightarrow \infty$  ( $j > J$ ) where  $0 \leq \lambda < 1$  and  $M$  can be chosen to be  $2^J \max \{ |(y_{J,k+2} - y_{J,k+1}) - (y_{J,k+1} - y_{J,k})| \}$ .  $\square$

Proposition: cubic order of approx. for any  $\alpha$

Proposition: For  $\alpha$ , the fundamental function of the high resolution scheme has values of amplitude  $< 1$  for  $|x| > 1$  whereas the cubic Deslaurliers-Dubuc scheme reaches a maximum of  $1$ .

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