QuadraticResidue . .

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```

# 0 Basic

#### OA .vimrc

#### **OB PBDS**

```
OC pragma
#pragma GCC optimize("Ofast,unroll-loops")
#pragma GCC target("avx, avx2, sse, sse2
     ,sse3,ssse3,sse4,popcnt,abm,mmx,fma,tune=native")
// chrono
     ::steady_clock::now().time_since_epoch().count()
     LambdaCompare
 0D
auto cmp = [](int x, int y) { return x < y; };</pre>
std::set<int, decltype(cmp)> st(cmp);
      Graph
 1A 2SAT/SCC
struct SAT { // O-base
   int low[N], dfn[N], bln[N], n, Time, nScc;
   bool instack[N], istrue[N];
   stack<int> st;
   vector<int> G[N], SCC[N];
   void init(int _n) {
     n = _n; // assert(n * 2 <= N);
     for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
   void add_edge(int a, int b) { G[a].emplace_back(b); }
   int rv(int a) {
     if (a >= n) return a - n;
     return a + n;
   void add_clause(int a, int b) {
     add_edge(rv(a), b), add_edge(rv(b), a);
   void dfs(int u) {
     dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
       do {
         tmp = st.top(), st.pop();
         instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != υ);
       ++nScc;
     }
   bool solve() {
     Time = nScc = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       if (!dfn[i]) dfs(i);
     for (int i =
         0; i < n + n; ++i) SCC[bln[i]].emplace_back(i);
     for (int i = 0; i < n; ++i) {</pre>
       if (bln[i] == bln[i + n]) return false;
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true;
  }
};
 1B BCC Vertex
```

```
int n, m, dfn[N], low[N], is_cut[N], nbcc = 0, t = 0;
vector<int> g[N], bcc[N], G[2 * N];
stack<int> st;
void tarjan(int p, int lp) {
  dfn[p] = low[p] = ++t;
  st.push(p);
  for (auto i : g[p]) {
    if (!dfn[i]) {
```

```
tarjan(i, p);
       low[p] = min(low[p], low[i]);
       if (dfn[p] <= low[i]) {</pre>
         nbcc++
         is_cut[p] = 1;
         for (int x = 0; x != i; st.pop()) {
           x = st.top();
           bcc[nbcc].push_back(x);
         bcc[nbcc].push_back(p);
    } else low[p] = min(low[p], dfn[i]);
}
void build() { // [n+1,n+nbcc] cycle, [1,n] vertex
  for (int i = 1; i <= nbcc; i++) {</pre>
     for (auto j : bcc[i]) {
       G[i + n].push_back(j);
       G[j].push_back(i + n);
  }
| }
```

#### **1C** VirtualTree

```
// requires DFS io, lca, is_child
vector<int> tre[N];
bool cmp(int a, int b){ return in[a] < in[b]; }</pre>
void add_edge(int a, int b){
  tre[a].emplace_back(b);
  tre[b].emplace_back(a);
void virtual_tree(vector<int> arr, int k){
  vector<int> sta:
  sort(arr.begin(), arr.end(), cmp);
  for (int i = 1; i < k; i++)
    arr.emplace_back(lca(arr[i], arr[i - 1]));
  sort(arr.begin(), arr.end(), cmp);
  arr.resize
       (unique(arr.begin(), arr.end()) - arr.begin());
  for (auto i : arr){
    while (!sta.empty
         () && !is_child(sta.back(), i)) sta.pop_back();
     if (!sta.empty()) add_edge(sta.back(), i);
     sta.push_back(i);
| }
```

### MinimumMeanCycle

```
let dp[i][j] = min length from 1 to j exactly i edges
ans = min (dp[n + 1][u] - dp[i][u]) / (n + 1 - i) */
```

# 1E MaximumCliqueDyn

```
struct MaxClique { // fast when N <= 100</pre>
 bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
   n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
 void add_edge(int u, int v) {
   G[v][v] = G[v][v] = 1;
 void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(all(r)
          , [&](int x, int y) { return d[x] > d[y]; });
   }
   vector<int> c(r.size());
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
```

```
while ((cs[k] & G[p]).any()) ++k;
       if (k > rgt) cs[++rgt + 1].reset();
       cs[k][p] = 1;
       if (k < lft) r[tp++] = p;
    for (int k = lft; k <= rgt; ++k)</pre>
       for (int p = cs[k]._Find_first
           (); p < N; p = cs[k]._Find_next(p))
         r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<</pre>
       int> &r, vector<int> &c, int l, bitset<N> mask) {
     while (!r.empty()) {
       int p = r.back();
       r.pop_back(), mask[p] = 0;
       if (q + c.back() <= ans) return;</pre>
       cur[q++] = p;
       vector<int> nr;
       for (int i : r) if (G[p][i]) nr.emplace_back(i);
       if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
       else if (q > ans) ans = q, copy_n(cur, q, sol);
       c.pop_back(), --q;
    }
  }
  int solve() {
    vector<int> r(n);
     ans = q = 0, iota(all(r), 0);
     pre_dfs(r, 0, bitset<N>(string(n, '1')));
     return ans;
|};
```

```
1F MinimumSteinerTree
struct SteinerTree { // O-base
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcst[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) {</pre>
      fill_n(dst[i], n, INF);
      dst[i][i] = vcst[i] = 0;
  }
  void chmin(int &x, int val) {
    x = min(x, val);
  void add_edge(int vi, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)</pre>
      for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)</pre>
          chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
    shortest_path();
    int t = ter.size(), full = (1 << t) - 1;</pre>
    for (int i = 0; i <= full; ++i)</pre>
      fill_n(dp[i], n, INF);
    copy_n(vcst, n, dp[0]);
    for (int msk = 1; msk <= full; ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
        for (int i = 0; i < n; ++i)</pre>
           dp[msk
               ][i] = vcst[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
        for (int sub = (
             msk - 1) \& msk; sub; sub = (sub - 1) \& msk)
           chmin(dp[msk][i],
               dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
      for (int i = 0; i < n; ++i) {</pre>
```

ll in[N];

```
tdst[i] = INF;
                                                               void init() { E.clear(); }
         for (int j = 0; j < n; ++j)</pre>
                                                               void add_edge(int u, int v, ll w) {
           chmin(tdst[i], dp[msk][j] + dst[j][i]);
                                                                 if (u != v) E.emplace_back(edge{u, v, w});
      copy_n(tdst, n, dp[msk]);
                                                               ll build(int root, int n) {
                                                                 ll ans = 0;
                                                                 for (;;) {
    return *min_element(dp[full], dp[full] + n);
                                                                   fill_n(in, n, INF);
}; // O(V 3^T + V^2 2^T)
                                                                   for (int i = 0; i < (int)E.size(); ++i)</pre>
                                                                     if (E[i].u != E[i].v && E[i].w < in[E[i].v])</pre>
1G DominatorTree
                                                                       pe[E[i].v] = i, in[E[i].v] = E[i].w;
                                                                   for (int u = 0; u < n; ++u) // no solution
struct DominatorTree { // 1-base
                                                                     if (u != root && in[u] == INF) return -INF;
  vector<int> G[N], rG[N];
                                                                   int cntnode = 0:
  int n, pa[N], dfn[N], id[N], Time;
                                                                   fill_n(id, n, -1), fill_n(vis, n, -1);
  int semi[N], idom[N], best[N];
                                                                   for (int u = 0; u < n; ++u) {
  vector<int> tree[N]; // dominator_tree
                                                                     if (u != root) ans += in[u];
  void init(int _n) {
                                                                     int v = v;
    n = _n;
                                                                     while (vis[v] != u && !~id[v] && v != root)
    for (int i = 1; i <= n; ++i)</pre>
                                                                       vis[v] = u, v = E[pe[v]].u;
      G[i].clear(), rG[i].clear();
                                                                     if (v != root && !~id[v]) {
  }
                                                                       for (int x = E[pe[v]].u; x != v;
  void add_edge(int u, int v) {
                                                                            x = E[pe[x]].u)
    G[u].emplace_back(v), rG[v].emplace_back(u);
                                                                         id[x] = cntnode;
                                                                       id[v] = cntnode++;
  void dfs(int u) {
    id[dfn[u] = ++Time] = u;
                                                                   }
    for (auto v : G[u])
                                                                   if (!cntnode) break; // no cycle
      if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
                                                                   for (int u = 0; u < n; ++u)
                                                                     if (!\sim id[v]) id[v] = cntnode++;
  int find(int y, int x) {
                                                                   for (int i = 0; i < (int)E.size(); ++i) {</pre>
    if (y <= x) return y;</pre>
                                                                     int v = E[i].v;
    int tmp = find(pa[y], x);
                                                                     E[i].v = id[E[i].v], E[i].v = id[E[i].v];
    if (semi[best[y]] > semi[best[pa[y]]])
                                                                     if (E[i].u != E[i].v) E[i].w -= in[v];
      best[y] = best[pa[y]];
    return pa[y] = tmp;
                                                                   n = cntnode, root = id[root];
  void tarjan(int root) {
                                                                 return ans;
    Time = 0;
                                                               }
    for (int i = 1; i <= n; ++i) {</pre>
                                                            1};
      dfn[i] = idom[i] = 0;
       tree[i].clear();
                                                             1I
                                                                 DMST
      best[i] = semi[i] = i;
                                                            #define rep(i, a, b) for (int i = a; i < (b); ++i)</pre>
    dfs(root);
                                                             #define sz(x) (int)(x).size()
    for (int i = Time; i > 1; --i) {
                                                            typedef vector<int> vi;
      int u = id[i];
                                                            struct RollbackUF {
      for (auto v : rG[u])
                                                               vi e;
                                                               vector<pii> st;
         if (v = dfn[v]) {
                                                               RollbackUF(int n) : e(n, -1) {}
          find(v, i);
                                                               int size(int x) { return -e[find(x)]; }
           semi[i] = min(semi[i], semi[best[v]]);
                                                               int find(int x) { return e[x] < 0 ? x : find(e[x]); }</pre>
                                                               int time() { return sz(st); }
      tree[semi[i]].emplace_back(i);
                                                               void rollback(int t) {
      for (auto v : tree[pa[i]]) {
                                                                 for (int i = time(); i-- > t;)
         find(v, pa[i]);
                                                                   e[st[i].first] = st[i].second;
         idom[v] =
                                                                 st.resize(t);
           semi[best[v]] == pa[i] ? pa[i] : best[v];
      }
                                                               bool join(int a, int b) {
      tree[pa[i]].clear();
                                                                 a = find(a), b = find(b);
                                                                 if (a == b) return false;
    for (int i = 2; i <= Time; ++i) {</pre>
                                                                 if (e[a] > e[b]) swap(a, b);
       if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
                                                                 st.push_back({a, e[a]});
       tree[id[idom[i]]].emplace_back(id[i]);
                                                                 st.push_back({b, e[b]});
                                                                 e[a] += e[b];
  }
                                                                 e[b] = a;
|};
                                                                 return true;
1H DMST(slow)
                                                              }
                                                            };
struct DMST { // O(VE)
                                                            struct Edge {
  struct edge {
                                                               int a, b;
    int u, v;
                                                               ll w;
    ll w;
                                                            };
  };
                                                            struct Node { // lazy skew heap node
  vector<edge> E; // O-base
                                                               Edge key;
  int pe[N], id[N], vis[N];
                                                               Node *l, *r;
```

ll delta;

auto update = [&](int u) {

```
void prop() {
                                                                for (X[u] = 1; C[u][X[u]]; X[u]++);
    key.w += delta;
    if (l) l->delta += delta;
                                                              auto color = [&](int u, int v, int c) {
    if (r) r->delta += delta;
                                                                int p = G[u][v];
    delta = 0;
                                                                G[\upsilon][v] = G[v][\upsilon] = c;
                                                                C[u][c] = v;
 Edge top() {
                                                                C[v][c] = u;
    prop();
                                                                C[v][p] = C[v][p] = 0;
    return key;
                                                                if (p) X[u] = X[v] = p;
 }
                                                                else update(u), update(v);
};
                                                                return p;
Node *merge(Node *a, Node *b) {
                                                              };
 if (!a || !b) return a ?: b;
                                                              auto flip = [&](int u, int c1, int c2) {
 a->prop(), b->prop();
                                                                int p = C[u][c1];
 if (a->key.w > b->key.w) swap(a, b);
                                                                swap(C[u][c1], C[u][c2]);
  swap(a->l, (a->r = merge(b, a->r)));
                                                                if (p) G[u][p] = G[p][u] = c2;
  return a;
                                                                if (!C[v][c1]) X[v] = c1;
                                                                if (!C[u][c2]) X[u] = c2;
void pop(Node *&a) {
                                                                return p;
 a->prop();
                                                              }:
 a = merge(a->l, a->r);
                                                              for (int i = 1; i <= n; i++) X[i] = 1;</pre>
                                                              for (int t = 0; t < (int)E.size(); t++) {</pre>
                                                                int u = E[t].first, v0 = E[t].second, v = v0,
pair<ll, vi> dmst(int n, int r, vector<Edge> &g) {
                                                                    c0 = X[u], c = c0, d;
 RollbackUF uf(n);
                                                                vector<pii> L;
  vector<Node *> heap(n);
                                                                int vst[n] = {};
 for (Edge e : g)
                                                                while (!G[u][v0]) {
   heap[e.b] = merge(heap[e.b], new Node{e});
                                                                  L.emplace_back(v, d = X[v]);
 ll res = 0;
                                                                  if (!C[v][c])
 vi seen(n, -1), path(n), par(n);
                                                                     for (a = (int)L.size() - 1; a >= 0; a--)
  seen[r] = r;
                                                                      c = color(u, L[a].first, c);
  vector<Edge> Q(n), in(n, {-1, -1}), comp;
                                                                  else if (!C[u][d])
  deque<tuple<int, int, vector<Edge>>> cycs;
                                                                     for (a = (int)L.size() - 1; a >= 0; a--)
  rep(s, 0, n) {
                                                                      color(u, L[a].first, L[a].second);
    int u = s, qi = 0, w;
                                                                  else if (vst[d]) break;
    while (seen[u] < 0) {
                                                                  else vst[d] = 1, v = C[u][d];
      if (!heap[u]) return {-1, {}};
      Edge e = heap[u]->top();
                                                                if (!G[u][v0]) {
      heap[u]->delta -= e.w, pop(heap[u]);
                                                                  for (; v; v = flip(v, c, d), swap(c, d));
      Q[qi] = e, path[qi++] = u, seen[u] = s;
                                                                  if (C[u][c0]) {
      res += e.w, u = uf.find(e.a);
                                                                    for (a = (int)L.size() - 2;
      if (seen[u] == s) { /// found cycle, contract
                                                                         a >= 0 \&\& L[a].second != c; a--)
        Node *cyc = 0;
        int end = qi, time = uf.time();
                                                                     for (; a >= 0; a--)
        do cyc = merge(cyc, heap[w = path[--qi]]);
                                                                      color(u, L[a].first, L[a].second);
        while (uf.join(u, w));
                                                                  } else t--;
        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
        cycs.push_front(\{u, time, \{\&Q[qi], \&Q[end]\}\}\);
                                                              }
      }
                                                            }
   }
                                                           |} // namespace Vizing
    rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
                                                            1K MinimumCliqueCover
 for (auto &[u, t, cmp] : cycs) {
                                                            struct CliqueCover { // O-base, O(n2^n)
   // restore sol (optional)
                                                              int co[1 << N], n, E[N];</pre>
    uf.rollback(t);
                                                              int dp[1 << N];</pre>
    Edge inEdge = in[u];
                                                              void init(int _n) {
    for (auto &e : cmp) in[uf.find(e.b)] = e;
                                                                n = _n, fill_n(dp, 1 << n, 0);
    in[uf.find(inEdge.b)] = inEdge;
                                                                fill_n(E, n, 0), fill_n(co, 1 << n, 0);
 rep(i, 0, n) par[i] = in[i].a;
                                                              void add_edge(int u, int v) {
  return {res, par};
                                                                E[v] \mid = 1 << v, E[v] \mid = 1 << v;
                                                              int solve() {
1J VizingTheorem
                                                                for (int i = 0; i < n; ++i)</pre>
namespace Vizing { // Edge coloring
                                                                  co[1 << i] = E[i] | (1 << i);
                    // G: coloring adjM
                                                                co[0] = (1 << n) - 1;
int C[N][N], G[N][N];
                                                                dp[0] = (n \& 1) * 2 - 1;
void clear(int n) {
                                                                for (int i = 1; i < (1 << n); ++i) {
 for (int i = 0; i <= n; i++) {</pre>
                                                                  int t = i & -i;
    for (int j = 0; j <= n; j++) C[i][j] = G[i][j] = 0;</pre>
                                                                  dp[i] = -dp[i ^ t];
                                                                  co[i] = co[i ^ t] & co[t];
 }
void solve(vector<pii> &E, int n) {
                                                                for (int i = 0; i < (1 << n); ++i)</pre>
 int X[n] = {}, a;
                                                                  co[i] = (co[i] \& i) == i;
```

fwt(co, 1 << n, 1); // needs FWHT</pre>

```
for (int ans = 1; ans < n; ++ans) {
   int sum = 0; // probabilistic
   for (int i = 0; i < (1 << n); ++i)
      sum += (dp[i] *= co[i]);
   if (sum) return ans;
   }
  return n;
}
</pre>
```

# 1L CountMaximalClique

```
struct BronKerbosch { // 1-base
   int n, a[N], g[N][N];
   int S, all[N][N], some[N][N], none[N][N];
   void init(int _n) {
    n = _n;
     for (int i = 1; i <= n; ++i)</pre>
       for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
     g[v][v] = g[v][v] = 1;
  }
  void dfs(int d, int an, int sn, int nn) {
     if (S > 1000) return; // pruning
     if (sn == 0 && nn == 0) ++S;
     int u = some[d][0];
     for (int i = 0; i < sn; ++i) {</pre>
       int v = some[d][i];
       if (g[v][v]) continue;
int tsn = 0, tnn = 0;
       copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)</pre>
         if (g[v][some[d][j]])
           some[d + 1][tsn++] = some[d][j];
       for (int j = 0; j < nn; ++j)</pre>
         if (g[v][none[d][j]])
           none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
       some[d][i] = 0, none[d][nn++] = v;
     }
  }
  int solve() {
    iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
     return S;
|};
```

#### 1M Theorems

 $|\max$  independent edge  $\mathsf{set}| = |V| - |\min$  edge cover  $|\max$  independent  $\mathsf{set}| = |V| - |\min$  vertex cover |

# 2 Flow-Matching

# 2A HopcroftKarp

```
struct HopcroftKarp
     { // O-based, return btoa to get matching
 bool dfs(int a, int L, vector<vector<int>> &g,
    vector<int> &btoa, vector<int> &A,
    vector<int> &B) {
    if (A[a] != L) return 0;
   A[a] = -1;
    for (int b : g[a])
     if (B[b] == L + 1) {
        B[b] = 0;
        if (btoa[b] == -1 ||
          dfs(btoa[b], L + 1, g, btoa, A, B))
          return btoa[b] = a, 1;
   return 0;
 }
 int solve(vector<vector<int>> &g, int m) {
    int res = 0;
    vector<int> btoa(m, -1), A(g.size()),
     B(btoa.size()), cur, next;
```

```
for (;;) {
       fill(all(A), 0), fill(all(B), 0);
       cur.clear();
       for (int a : btoa)
         if (a != -1) A[a] = -1;
       for (int a = 0; a < (int)g.size(); a++)</pre>
         if (A[a] == 0) cur.push_back(a);
       for (int lay = 1;; lay++) {
         bool islast = 0;
         next.clear();
         for (int a : cur)
           for (int b : g[a]) {
             if (btoa[b] == -1) {
               B[b] = lay;
               islast = 1;
             } else if (btoa[b] != a && !B[b]) {
               B[b] = lay;
               next.push_back(btoa[b]);
           }
         if (islast) break;
         if (next.empty()) return res;
         for (int a : next) A[a] = lay;
         cur.swap(next);
       for (int a = 0; a < (int)g.size(); a++)</pre>
         res += dfs(a, 0, g, btoa, A, B);
  }
};
2B
     KM
struct KM { // O-base, maximum matching
  11 w[N][N], hl[N], hr[N], slk[N];
   int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
```

```
bool vl[N], vr[N];
void init(int _n) {
  n = _n;
  for (int i = 0; i < n; ++i)</pre>
    fill_n(w[i], n, -INF);
void add_edge(int a, int b, ll wei) {
 w[a][b] = wei;
bool Check(int x) {
  if (vl[x] = 1, \sim fl[x])
    return vr[qu[qr++] = fl[x]] = 1;
  while (\sim x) swap(x, fr[fl[x] = pre[x]]);
  return 0;
void bfs(int s) {
  fill_n(slk
      , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
  ql = qr = 0, qu[qr++] = s, vr[s] = 1;
  for (ll d;;) {
    while (ql < qr)</pre>
      for (int x = 0, y = qu[ql++]; x < n; ++x)
        if (!vl[x] && slk
             [x] >= (d = hl[x] + hr[y] - w[x][y])) {
          if (pre[x] = y, d) slk[x] = d;
          else if (!Check(x)) return;
    d = INF;
    for (int x = 0; x < n; ++x)
      if (!vl[x] \&\& d > slk[x]) d = slk[x];
    for (int x = 0; x < n; ++x) {
      if (vl[x]) hl[x] += d;
      else slk[x] -= d;
      if (vr[x]) hr[x] -= d;
    for (int x = 0; x < n; ++x)
      if (!vl[x] && !slk[x] && !Check(x)) return;
  }
}
ll solve() {
```

# 2C MCMF

```
struct MinCostMaxFlow { // O-base
  struct Edge {
    tl from, to, cap, flow, cost, rev;
  } *past[N];
  vector<Edge> G[N];
  int inq[N], n, s, t;
 ll dis[N], up[N], pot[N];
 bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 \&\& dis[u] > d) {
        dis[v] = d, vp[v] = cap, past[v] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
    };
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[v] = 0;
      for (auto &e : G[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
          e.to, d2, min(up[u], e.cap - e.flow), &e);
      }
    }
    return dis[t] != INF;
 }
 bool Dijkstra() {
    fill_n(dis, n, INF);
    priority_queue<pll, vector<pll>, greater<pll>>> pq;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
        pq.push(pll(d, u));
      }
    };
    relax(s, 0, INF, 0);
    while (!pq.empty()) {
      auto [d, u] = pq.top();
      pq.pop();
      if (dis[v] != d) continue;
      for (auto &e : G[u]) {
        ll d2 = dis[u] + e.cost + pot[u] - pot[e.to];
        relax(
          e.to, d2, min(up[u], e.cap - e.flow), &e);
      }
    }
    return dis[t] != INF;
  void solve(int _s, int _t, ll &flow, ll &cost,
    bool neg = true) {
    s = _s, t = _t, flow = 0, cost = 0;
    if (neg) BellmanFord(), copy_n(dis, n, pot);
    // do BellmanFord() if time isn't tight
    for (; Dijkstra(); copy_n(dis, n, pot)) {
      for (int i = 0; i < n; ++i)</pre>
        dis[i] += pot[i] - pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], G[e.to][e.rev].flow -= up[t];
```

```
}
}
void init(int _n) {
    n = _n, fill_n(pot, n, 0);
    for (int i = 0; i < n; ++i) G[i].clear();
}
void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].emplace_back(
    Edge{a, b, cap, 0, cost, (int)G[b].size()});
    G[b].emplace_back(
    Edge{b, a, 0, 0, -cost, (int)G[a].size() - 1});
};</pre>
```

# 2D GeneralGraphMatching

```
struct Matching { // O-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
    }
  }
  void Blossom(int x, int y, int l) {
    for (; Find(x) != l; x = pre[y]) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
       for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
    }
  bool Bfs(int r) {
    iota(all(fa), 0); fill(all(s), -1);
    q = queue<int>(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : G[x])
        if (s[v] == -1) {
           if (pre[u] = x, s[u] = 1, match[u] == n) {
             for (int a = u, b = x, last;
                 b != n; a = last, b = pre[a])
               last =
                   match[b], match[b] = a, match[a] = b;
            return true;
           q.push(match[u]); s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int l = LCA(u, x);
           Blossom(x, u, l); Blossom(u, x, l);
        }
    }
    return false;
  Matching(int _n): n(_n), fa(n + 1), s(n + 1), vis
       (n + 1), pre(n + 1, n), match(n + 1, n), G(n) {}
  void add_edge(int u, int v)
  { G[u].emplace_back(v), G[v].emplace_back(u); }
  int solve()
    int ans = 0;
    for (int x = 0; x < n; ++x)
      if (match[x] == n) ans += Bfs(x);
    return ans;
  } // match[x] == n means not matched
|};
```

# 2E MaxWeightMaching

```
#define rep(i, l, r) for (int i = (l); i <= (r); ++i)
struct WeightGraph { // 1-based, note int!
    struct edge {
    int u, v, w;
    };</pre>
```

```
int n, nx;
vector<int> lab;
vector<vector<edge>> g;
vector<int> slack, match, st, pa, S, vis;
vector<vector<int>> flo, flo_from;
queue<int> q;
WeightGraph(int n_)
  : n(n_{-}), nx(n * 2), lab(nx + 1),
    g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1),
    flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
  match = st = pa = S = vis = slack;
  rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
int ED(edge e) {
  return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
}
void update_slack(int u, int x, int &s) {
 if (!s || ED(g[v][x]) < ED(g[s][x])) s = u;</pre>
}
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)
    if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
      update_slack(u, x, slack[x]);
void q_push(int x) {
  if (x \le n) q.push(x);
  else
    for (int y : flo[x]) q_push(y);
void set_st(int x, int b) {
  st[x] = b;
  if(x > n)
    for (int y : flo[x]) set_st(y, b);
vector<int> split_flo(auto &f, int xr) {
  auto it = find(all(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
    reverse(1 + all(f)), it = f.end() - pr;
  auto res = vector(f.begin(), it);
  return f.erase(f.begin(), it), res;
}
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  int xr = flo_from[u][g[u][v].u];
  auto &f = flo[u], z = split_flo(f, xr);
  rep(i, 0, (int)z.size() - 1)
    set_match(z[i], z[i ^ 1]);
  set_match(xr, v);
  f.insert(f.end(), all(z));
void augment(int u, int v) {
  for (;;) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
  }
}
int lca(int u, int v) {
  static int t = 0;
  ++t;
  for (++t; u || v; swap(u, v))
    if (u) {
      if (vis[u] == t) return u;
      vis[u] = t;
      u = st[match[u]];
      if (u) u = st[pa[u]];
  return 0;
}
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + all(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0;
```

```
match[b] = match[o];
  vector<int> f = {o};
  for (int x = u, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  reverse(1 + all(f));
  for (int x = v, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  flo[b] = f;
  set_st(b, b);
  for (int x = 1; x <= nx; ++x)
    g[b][x].w = g[x][b].w = 0;
  fill(all(flo_from[b]), 0);
  for (int xs : flo[b]) {
    for (int x = 1; x <= nx; ++x)
      if (g[b][x].w == 0 ||
        ED(g[xs][x]) < ED(g[b][x])
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)</pre>
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  }
  set_slack(b);
}
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) {
      xs = x;
      continue;
    pa[xs] = g[x][xs].u;
    S[xs] = 1, S[x] = 0;
    slack[xs] = 0;
    set_slack(x);
    q_push(x);
    xs = -1;
  for (int x : flo[b])
    if (x == xr) S[x] = 1, pa[x] = pa[b];
    else S[x] = -1, set_slack(x);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
    int nu = st[match[v]];
    pa[v] = e.u;
    S[v] = 1;
    slack[v] = slack[nu] = 0;
    S[nu] = 0;
    q_push(nu);
  } else if (S[v] == 0) {
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
  return false;
}
bool matching() {
  fill(all(S), -1), fill(all(slack), 0);
  q = queue<int>();
  for (int x = 1; x <= nx; ++x)</pre>
    if (st[x] == x \&\& !match[x])
      pa[x] = 0, S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
    while (q.size()) {
      int u = q.front();
      q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (ED(q[v][v]) != 0)
            update_slack(u, st[v], slack[st[v]]);
          else if (on_found_edge(g[u][v]))
```

```
return true:
           }
       int d = INF;
       for (int b = n + 1; b <= nx; ++b)</pre>
         if (st[b] == b && S[b] == 1)
           d = min(d, lab[b] / 2);
       for (int x = 1; x <= nx; ++x)
         if (int s = slack[x];
             st[x] == x \&\& s \&\& S[x] <= 0)
           d = min(d, ED(g[s][x]) / (S[x] + 2));
       for (int u = 1; u <= n; ++u)</pre>
         if (S[st[u]] == 1) lab[u] += d;
         else if (S[st[u]] == 0) {
           if (lab[u] <= d) return false;</pre>
           lab[u] -= d;
         }
       rep(b, n + 1, nx) if (st[b] == b && S[b] >= 0)
         lab[b] += d * (2 - 4 * S[b]);
       for (int x = 1; x <= nx; ++x)
         if (int s = slack[x]; st[x] == x && s &&
             st[s] != x \&\& ED(g[s][x]) == 0)
           if (on_found_edge(g[s][x])) return true;
       for (int b = n + 1; b <= nx; ++b)
         if (st[b] == b && S[b] == 1 && lab[b] == 0)
           expand_blossom(b);
     return false;
  }
  pair<ll, int> solve() {
    fill(all(match), 0);
     rep(u, 0, n) st[u] = u, flo[u].clear();
     int w_max = 0;
     rep(u, 1, n) rep(v, 1, n) {
       flo_from[u][v] = (u == v ? u : 0);
       w_{max} = max(w_{max}, g[u][v].w);
    fill(all(lab), w_max);
     int n_matches = 0;
     tot_weight = 0;
     while (matching()) ++n_matches;
     rep(u, 1, n) if (match[u] \&\& match[u] < u)
       tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
  }
  void add_edge(int u, int v, int w) {
     g[v][v].w = g[v][v].w = w;
|};
```

### GlobalMinCut

```
struct StoerWagner { // O(V^3), is it O(VE + V log V)?
 int vst[N], edge[N][N], wei[N];
  void init(int n) {
    for (int i = 0; i < n; ++i) fill_n(edge[i], n, 0);</pre>
 void addEdge(int u, int v, int w) {
    edge[u][v] += w;
    edge[v][u] += w;
 int search(int &s, int &t, int n) {
    fill_n(vst, n, 0), fill_n(wei, n, 0);
    s = t = -1;
    int mx, cur;
    for (int j = 0; j < n; ++j) {</pre>
      mx = -1, cur = 0;
      for (int i = 0; i < n; ++i)</pre>
        if (wei[i] > mx) cur = i, mx = wei[i];
      vst[cur] = 1, wei[cur] = -1;
      s = t;
      t = cur;
      for (int i = 0; i < n; ++i)</pre>
        if (!vst[i]) wei[i] += edge[cur][i];
    return mx;
```

```
int solve(int n) {
     int res = INF;
     for (int x, y; n > 1; n--) {
       res = min(res, search(x, y, n));
       for (int i = 0; i < n; ++i)</pre>
         edge[i][x] = (edge[x][i] += edge[y][i]);
       for (int i = 0; i < n; ++i) {</pre>
         edge[y][i] = edge[n - 1][i];
         edge[i][y] = edge[i][n - 1];
       } // edge[y][y] = 0;
     }
     return res;
  }
|} sw;
```

# 2G BoundedFlow(Dinic)

```
struct BoundedFlow { // O-base
  struct edge { // note int!
    int to, cap, flow, rev;
  };
  vector<edge> G[N];
  int n, s, t, dis[N], cur[N], cnt[N];
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n + 2; ++i)</pre>
      G[i].clear(), cnt[i] = 0;
  void add_edge(int u, int v, int lcap, int rcap) {
    cnt[u] -= lcap, cnt[v] += lcap;
    G[u].emplace_back(
      edge{v, rcap, lcap, (int)G[v].size()});
    G[v].emplace_back(
      edge{u, 0, 0, (int)G[u].size() - 1});
  void add_edge(int u, int v, int cap) {
    G[u].emplace_back(
      edge{v, cap, 0, (int)G[v].size()});
    G[v].emplace_back(
      edge{u, 0, 0, (int)G[u].size() - 1});
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < (int)G[u].size(); ++i) {</pre>
      edge &e = G[u][i];
      if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
        int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df, G[e.to][e.rev].flow -= df;
          return df;
      }
    }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n + 3, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (edge &e : G[u])
        if (!~dis[e.to] && e.flow != e.cap)
          q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
```

```
}
  bool solve() {
     int sum = 0;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
         add_edge(n + 1, i, cnt[i]), sum += cnt[i];
       else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
    for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
         G[n + 1].pop_back(), G[i].pop_back();
       else if (cnt[i] < 0)</pre>
         G[i].pop_back(), G[n + 2].pop_back();
    return sum != -1;
  }
  int solve(int _s, int _t) {
    add_edge(_t, _s, INF);
     if (!solve()) return -1; // invalid flow
     int x = G[_t].back().flow;
    return G[_t].pop_back(), G[_s].pop_back(), x;
|};
```

# 2H GomoryHuTree

```
| BoundedFlow Dinic;
| int g[N];
| void add_edge(int u, int v, int w); // TODO
| void GomoryHu(int n) { // O-base
| fill_n(g, n, 0);
| for (int i = 1; i < n; ++i) {
| Dinic.init(n);
| // build the graph
| add_edge(i, g[i], Dinic.maxflow(i, g[i]));
| for (int j = i + 1; j <= n; ++j)
| if (g[j] == g[i] && ~Dinic.dis[j])
| g[j] = i;
| }
```

#### 2I MinCostCirculation

```
struct MinCostCirculation { // O-base
 struct Edge {
    tl from, to, cap, fcap, flow, cost, rev;
 } *past[N];
 vector<Edge> G[N];
 ll dis[N], inq[N], n;
 void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, Edge *e) {
     if (dis[v] > d) {
       dis[u] = d, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
   };
   relax(s, 0, 0);
   while (!q.empty()) {
     int u = q.front();
     q.pop(), inq[v] = 0;
     for (auto &e : G[u])
        if (e.cap > e.flow)
          relax(e.to, dis[u] + e.cost, &e);
 }
 void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
   if (dis[cur.from] + cur.cost < 0) {</pre>
     ++cur.flow, --G[cur.to][cur.rev].flow;
     for (int
           i = cur.from; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        ++e.flow, --G[e.to][e.rev].flow;
```

```
++cur.cap;
   }
   void solve(int mxlg) {
     for (int b = mxlg; b >= 0; --b) {
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           e.cap *= 2, e.flow *= 2;
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
    }
  }
   void init(int _n) { n = _n;
     for (int i = 0; i < n; ++i) G[i].clear();</pre>
   void add_edge(ll a, ll b, ll cap, ll cost) {
     G[a].emplace_back(Edge{a, b,
          0, cap, 0, cost, (ll)G[b].size() + (a == b));
     G[b].emplace_back(Edge
         {b, a, 0, 0, 0, -cost, (ll)G[a].size() - 1});
  }
|} mcmf; // O(VE * ElogC)
```

# 2J FlowModelsBuilding

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - For each edge (x,y,l,u), connect x→y with capacity u-l.
     For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - of outgoing lower bounds. 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - there's no solution. Otherwise, f' is the answer. 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x,y) \in M$ ,  $x \rightarrow y$  otherwise.
  - 2. DFS from unmatched vertices in X. 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x,y,c), connect  $x \rightarrow y$  with (cost,cap) = (c,1) if c>0, otherwise connect  $y \rightarrow x$  with (cost,cap) = (-c,1)
  - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost,cap)=(0,d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \to T$  with  $(cost \, con) = (0 d(v))$
  - $\begin{array}{c} (cost, cap) = (0, -d(v)) \\ \text{6. Flow from } S \text{ to } T\text{, the answer is the cost of the flow} \\ C+K \end{array}$
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let  ${\cal K}$  be the sum of all weights
  - 3. Connect source  $s \rightarrow v$ ,  $v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow  $f\!<\!K|V|$
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v).
  - 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on  $G^\prime\,.$  • Project selection problem

```
1. If p_v > 0, create edge (s,v) with capacity p_v; otherwise, create edge (v,t) with capacity -p_v.

2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.

3. The mincut is equivalent to the maximum profit of a subset of projects.

Dual of minimum cost maximum flow

1. Capacity c_{uv}, Flow f_{uv}, Cost w_{uv}, Required Flow difference for vertex b_u.

2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

\min \sum_{uv} w_{uv} f_{uv} inline void Update(int x) { Tree[x].lazy ^= 1; swap(ls(x), rs(x)); inline void PushDown(int x) if (!Tree[x].lazy) return if (ls(x)) Update(ls(x)); if (rs(x)) Update(rs(x)); Tree[x].lazy = 0; line void Rotate(int x) { inline void Rotate(int x
```

 $p_u \ge 0$ 

# 3 Data Struture

 $\sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_u$ 

# 3A Treap

```
mt19937 rd(1);
#define sz(t) ((t) == 0 ? 0 : (t)->size)
struct Treap {
  int pri, size;
  Treap *l, *r;
  Treap(ll val = 0)
    : pri(rd()), size(1), l(0), r(0) {};
  void push();
  void pull() { size = 1 + sz(l) + sz(r); }
};
void spilt(int k, Treap *rt, Treap *&a, Treap *&b) {
  if (!rt) return a = b = 0, void();
  rt->push();
  int lsz = 1 + sz(rt->l);
  if (k >= lsz)
    a = rt, spilt(k - lsz, a->r, a->r, b), a->pull();
  else b = rt, spilt(k, b->l, a, b->l), b->pull();
Treap *merge(Treap *l, Treap *r) {
  if (!l) return r;
  if (!r) return l;
  if (l->pri < r->pri) {
    l->push(), l->r = merge(l->r, r), l->pull();
    return l;
  } else {
    r->push(), r->l = merge(l, r->l), r->pull();
    return r;
  }
}
```

#### 3B LinkCutTree

```
#define ls(x) Tree[x].son[0]
#define rs(x) Tree[x].son[1]
#define fa(x) Tree[x].fa
struct node {
  int son[2], Min, id, fa, lazy;
} Tree[N];
int n, m, q, w[N], Min;
struct Node {
  int u, v, w;
} a[N];
inline bool IsRoot(int x) {
  return (ls(fa(x)) == x \mid\mid rs(fa(x)) == x) ? false
                                               : true;
inline void PushUp(int x) {
  Tree[x].Min = w[x], Tree[x].id = x;
  if (ls(x) && Tree[ls(x)].Min < Tree[x].Min) {</pre>
    Tree[x].Min = Tree[ls(x)].Min;
    Tree[x].id = Tree[ls(x)].id;
  if (rs(x) && Tree[rs(x)].Min < Tree[x].Min) {</pre>
    Tree[x].Min = Tree[rs(x)].Min;
    Tree[x].id = Tree[rs(x)].id;
| }
```

```
Tree[x].lazy ^= 1;
  swap(ls(x), rs(x));
inline void PushDown(int x) {
  if (!Tree[x].lazy) return;
  if (ls(x)) Update(ls(x));
  if (rs(x)) Update(rs(x));
  Tree[x].lazy = 0;
inline void Rotate(int x) {
  int y = fa(x), z = fa(y), k = rs(y) == x,
      w = Tree[x].son[!k];
  if (!IsRoot(y)) Tree[z].son[rs(z) == y] = x;
  fa(x) = z, fa(y) = x;
  if (w) fa(w) = y;
  Tree[x].son[!k] = y, Tree[y].son[k] = w;
  PushUp(y);
}
inline void Splay(int x) {
  stack<int> Stack;
  int y = x, z;
  Stack.push(y);
  while (!IsRoot(y)) Stack.push(y = fa(y));
  while (!Stack.empty())
    PushDown(Stack.top()), Stack.pop();
  while (!IsRoot(x)) {
    y = fa(x), z = fa(y);
    if (!IsRoot(y))
      Rotate((ls(y) == x) ^( (ls(z) == y) ? x : y);
    Rotate(x);
  PushUp(x);
}
inline void Access(int root) {
  for (int x = 0; root; x = root, root = fa(root))
    Splay(root), rs(root) = x, PushUp(root);
inline void MakeRoot(int x) {
  Access(x), Splay(x), Update(x);
}
inline int FindRoot(int x) {
  Access(x), Splay(x);
  while (ls(x)) x = ls(x);
  return Splay(x), x;
}
inline void Link(int u, int v) {
  MakeRoot(u);
  if (FindRoot(v) != u) fa(u) = v;
}
inline void Cut(int u, int v) {
  MakeRoot(u);
  if (FindRoot(v) != u || fa(v) != u || ls(v)) return;
  fa(v) = rs(u) = 0;
}
inline void Split(int u, int v) {
  MakeRoot(u), Access(v), Splay(v);
inline bool Check(int u, int v) {
  return MakeRoot(u), FindRoot(v) == u;
inline int LCA(int root, int u, int v) {
  MakeRoot(root), Access(u), Access(v), Splay(u);
  if (!fa(u)) {
    Access(u), Splay(v);
    return fa(v);
  return fa(u);
}
/* ETT
每次進入節點和走邊都放入一次共 3n - 2
node(u) 表示進入節點 u 放入 treap 的位置
edge(u, v) 表示 u -> v 的邊放入 treap 的位置 (push v)
Makeroot u
 L1 = [begin, node(u) - 1], L2 = [node(u), end]
```

```
National Taiwan University | RngBased
   -> L2 + L1
                                                                 sa.resize(n), cnt.resize(n);
                                                                 rk.resize(n), tmp.resize(n);
Insert u, v :
                                                                 iota(all(sa), 0);
  Tu \rightarrow L1 = [begin, node(u) - 1], L2 = [node(u), end]
                                                                 sort(all(sa),
   Tv \rightarrow L3 = [begin, node(v) - 1], L4 = [node(v), end]
  -> L2 + L1 + edge(u, v) + L4 + L3 + edge(v, u)
                                                                 rk[0] = 0;
                                                                 for (int i = 1; i < n; i++)</pre>
Delect u, v
                                                                   rk[sa[i]] =
  maybe need swap u, v
   T -> L1 + edge(u, v) + L2 + edge(v, u) + L3
   -> L1 + L3, L2
                                                                   fill(all(cnt), 0);
      String
4A
     KMP
int KMP(string s, string t) {
                                                                       add(sa[i], -k);
  t = " "s + t; // consistency with ACa
                                                                   sa.swap(tmp);
  int ans = 0;
                                                                    tmp[sa[0]] = 0;
  vector<int> f(t.size(), 0);
  f[0] = -1;
  for (int i = 1, j = -1; i < (int)t.size(); i++) {</pre>
    while (j \ge 0 \&\& t[j + 1] != t[i]) j = f[j];
     f[i] = ++j;
  }
                                                                   rk.swap(tmp);
  for (int i = 0, j = 0; i < (int)s.size(); i++) {</pre>
                                                                 }
    while (j \ge 0 \&\& t[j + 1] != s[i]) j = f[j];
                                                               }
     if (++j + 1 == (int)t.size()) ans++, j = f[j];
                                                               void LCP(string s) {
  return ans;
                                                                 lcp.resize(n);
| }
                                                                 for (int i = 0; i < n; i++)</pre>
4B Z
                                                                   else {
                                                                      if (k) k--;
int Z[N];
void z(string s) {
                                                                     int j = sa[rk[i] - 1];
                                                                     while (
  for (int i = 1, mx = 0; i < (int)s.size(); i++) {</pre>
    if (i < Z[mx] + mx)
                                                                        k++
       Z[i] = min(Z[mx] - i + mx, Z[i - mx]);
                                                                     lcp[rk[i]] = k;
     while (
       Z[i] +
                                                               }
            i < (int)s.size() && s[i + Z[i]] == s[Z[i]])
                                                            |};
       Z[i]++;
     if (Z[i] + i > Z[mx] + mx) mx = i;
                                                                  SAIS
                                                             4E
}
                                                             auto sais(const auto &s) {
    Manacher
                                                               if (n == 1) return vector{0};
                                                               vector<int> c(z);
int man[N]; // len: man[i] - 1
                                                               for (int x : s) ++c[x];
void manacher(string s) { // uses 2|s|+1
                                                               partial_sum(all(c), begin(c));
  string t;
                                                               vector<int> sa(n);
  for (int i = 0; i < (int)s.size(); i++) {</pre>
                                                               auto I = views::iota(0, n);
     t.push_back('$');
                                                               vector<bool> t(n, true);
     t.push_back(s[i]);
                                                                 t[i] =
  t.push_back('$');
  int mx = 1:
                                                               auto is_lms = views::filter(
  for (int i = 0; i < (int)t.size(); i++) {</pre>
    man[i] = 1;
                                                               auto induce = [&] {
    man[i] = min(man[mx] + mx - i, man[2 * mx - i]);
                                                                 for (auto x = c; int y : sa)
     while (man[i] + i < (int)t.size() && i >= man[i] &&
                                                                    if (y--)
       t[i + man[i]] == t[i - man[i]])
       man[i]++;
     if (i + man[i] > mx + man[mx]) mx = i;
                                                                   if (y--)
  }
| }
                                                               };
```

# SuffixArray

```
struct SuffixArray {
#define add(x, k) (x + k + n) % n
  vector<int> sa, cnt, rk, tmp, lcp;
  // sa: order, rk[i]: pos of s[i..],
 // lcp[i]: LCP of sa[i], sa[i-1]
 void SA(string s) { // remember to append '\1'
   int n = (int)s.size();
```

```
[&](int i, int j) { return s[i] < s[j]; });
      rk[sa[i - 1]] + (s[sa[i - 1]] != s[sa[i]]);
  for (int k = 1; k <= n; k <<= 1) {</pre>
    for (int i = 0; i < n; i++)</pre>
      cnt[rk[add(sa[i], -k)]]++;
    for (int i = 1; i < n; i++) cnt[i] += cnt[i - 1];</pre>
    for (int i = n - 1; i >= 0; i--)
      tmp[--cnt[rk[add(sa[i], -k)]]] =
    for (int i = 1; i < n; i++)</pre>
      tmp[sa[i]] = tmp[sa[i - 1]] +
        (rk[sa[i - 1]] != rk[sa[i]] ||
          rk[add(sa[i - 1], k)] !=
            rk[add(sa[i], k)]);
 int n = (int)s.size(), k = 0;
    if (rk[i] == 0) lcp[rk[i]] = 0;
        \max(i, j) + k < n \&\& s[i + k] == s[j + k])
const int n = (int)s.size(), z = ranges::max(s) + 1;
for (int i = n - 2; i >= 0; --i)
    (s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  [&t](int x) { return x && t[x] && !t[x - 1]; });
      if (!t[y]) sa[x[s[y] - 1]++] = y;
  for (auto x = c; int y : sa | views::reverse)
      if (t[y]) sa[--x[s[y]]] = y;
vector<int> lms, q(n);
lms.reserve(n);
for (auto x = c; int i : I | is_lms)
 q[i] = (int)lms.size(),
 lms.emplace_back(sa[--x[s[i]]] = i);
induce();
vector<int> ns((int)lms.size());
for (int j = -1, nz = 0; int i : sa | is_lms) {
 if (j >= 0) {
```

```
int len = min({n - i, n - j, lms[q[i] + 1] - i});
      ns[q[i]] = nz += lexicographical_compare(
        begin(s) + j, begin(s) + j + len, begin(s) + i,
        begin(s) + i + len);
    j = i;
  }
  fill(all(sa), 0);
  auto nsa = sais(ns);
  for (auto x = c; int y : nsa | views::reverse)
    y = lms[y], sa[--x[s[y]]] = y;
  return induce(), sa;
// sa[i]: sa[i]-th suffix is the i-th lexicographically
// smallest suffix. hi[i]: LCP of suffix sa[i] and
// suffix sa[i - 1].
struct Suffix {
  int n:
  vector<int> sa, hi, ra;
  Suffix(const auto &_s, int _n)
    : n(_n), hi(n), ra(n) {
    vector<int> s(n + 1); // s[n] = 0;
    copy_n(_s, n, begin(s)); // _s shouldn't contain 0
    sa = sais(s);
    sa.erase(sa.begin());
    for (int i = 0; i < n; ++i) ra[sa[i]] = i;</pre>
    for (int i = 0, h = 0; i < n; ++i) {
      if (!ra[i]) {
        h = 0;
        continue;
      }
      for (int j = sa[ra[i] - 1];
           \max(i, j) + h < n \&\& s[i + h] == s[j + h];)
      hi[ra[i]] = h ? h-- : 0;
  }
|};
```

#### 4F ACAutomaton

```
#define sigma 26
#define base 'a'
struct AhoCorasick { // N: sum of length
  int ch[N][sigma] = \{\{\}\}, f[N] = \{-1\}, tag[N],
      mv[N][sigma], jump[N], cnt[N];
  int idx = 0, t = -1;
  vector<int> E[N], q;
  pii o[N];
  int insert(string &s) {
    int j = 0;
    for (int i = 0; i < (int)s.size(); i++) {</pre>
      if (!ch[j][s[i] - base])
        ch[j][s[i] - base] = ++idx;
      j = ch[j][s[i] - base];
    }
    tag[j] = 1;
    return j;
 }
 int next(int u, int c) {
   return u < 0 ? 0 : mv[u][c];</pre>
 }
 void dfs(int u) {
    o[v].F = ++t;
    for (auto v : E[u]) dfs(v);
    o[v].S = t;
 }
  void build() {
    int k = -1;
    q.emplace_back(0);
    while (++k < (int)q.size()) {</pre>
      int u = q[k];
      for (int v = 0; v < sigma; v++) {</pre>
        if (ch[u][v]) {
          f[ch[u][v]] = next(f[u], v);
          q.emplace_back(ch[u][v]);
```

```
mv[v][v] =
           (ch[u][v] ? ch[u][v] : next(f[u], v));
       if (u) jump[u] = (tag[f[u]] ? f[u] : jump[f[u]]);
     }
     reverse(q.begin(), q.end());
     for (int i = 1; i <= idx; i++)</pre>
       E[f[i]].emplace_back(i);
     dfs(0);
   void match(string &s) {
     fill(cnt, cnt + idx + 1, 0);
     for (int i = 0, j = 0; i < (int)s.size(); i++)</pre>
       cnt[j = next(j, s[i] - base)]++;
     for (int i : q)
       if (f[i] > 0) cnt[f[i]] += cnt[i];
  }
|} ac;
```

#### 4G MinRotation

```
int mincyc(string s) {
  int n = (int)s.size();
  s = s + s;
  int i = 0, ans = 0;
  while (i < n) {
    ans = i;
    int j = i + 1, k = i;
    while (j < 2 * n && s[j] >= s[k]) {
        k = (s[j] > s[k] ? i : k + 1);
        ++j;
    }
    while (i <= k) i += j - k;
}
return ans;
</pre>
```

#### 4H ExtSAM

```
#define CNUM 26
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
    fill_n(next[tot], CNUM, 0);
    len[tot] = cnt[tot] = link[tot] = 0;
    return tot++;
  void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len
        [p] + 1 == len[q]) return link[cur] = q, cur;
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[
          clone][i] = len[next[q][i]] ? next[q][i] : 0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
  void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
```

```
cnt[cur = nxt] += 1;
    }
  }
  void build() {
    queue<int> q;
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
           q.push(insertSAM(cur, i));
    vector<int> lc(tot);
    for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
    partial_sum(all(lc), lc.begin());
         = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
  }
  void solve() {
    for (int i = tot - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
|};
```

#### 4I PalindromeTree

```
struct PalindromicTree {
  struct node {
    int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
                   // pal. suf.
    node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
    }
  };
  vector<node> St:
  vector<char> s;
  int last, n;
  PalindromicTree() : St(2), last(1), n(0) {
    St[0].fail = 1, St[1].len = -1, s.emplace_back(-1);
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.emplace_back(0), St.emplace_back(-1);
    St[0].fail = 1, s.emplace_back(-1);
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
      x = St[x].fail;
    return x;
  inline void add(int c) {
    s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = (int)St.size();
      St.emplace_back(St[cur].len + 2);
      St[now].fail =
        St[get_fail(St[cur].fail)].next[c];
      St[cur].next[c] = now;
      St[now].num = St[St[now].fail].num + 1;
    last = St[cur].next[c], ++St[last].cnt;
  inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
      St[i->fail].cnt += i->cnt;
  }
  inline int size() { // The number of diff. pal.
    return (int)St.size() - 2;
};
```

# 5 Number Theory

#### 5A Primes

#### 5B ExtGCD

```
// beware of negative numbers!
void extgcd(ll a, ll b, ll c, ll &x, ll &y) {
  if (b == 0) x = c / a, y = 0;
  else {
    extgcd(b, a % b, c, y, x);
    y -= x * (a / b);
  }
} // |x| <= b/2, |y| <= a/2</pre>
```

#### 5C FloorCeil

```
|int floor(int a, int b)
|{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
|int ceil(int a, int b)
|{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

#### 5D FloorSum

Computes

$$f(a,b,c,n) = \sum_{i=0}^{n} \left\lfloor \frac{a \cdot i + b}{m} \right\rfloor$$

Furthermore, Let  $m = \left| \frac{an+b}{c} \right|$ :

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \left\lfloor \frac{ai+b}{c} \right\rfloor \\ &= \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c,c-b-1,a,m-1)) \\ -h(c,c-b-1,a,m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) = & \sum_{i=0}^{n} \left\lfloor \frac{ai+b}{c} \right\rfloor^2 \\ = & \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \\ + \left\lfloor \frac{a}{c} \right\rfloor \cdot \left\lfloor \frac{b}{c} \right\rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \left\lfloor \frac{a}{c} \right\rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \left\lfloor \frac{b}{c} \right\rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ - 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

```
| Il floorsum(Il A, Il B, Il C, Il N) {
    if (A == 0) return (N + 1) * (B / C);
    if (A > C || B > C)
        return (N + 1) * (B / C) +
        N * (N + 1) / 2 * (A / C) +
        floorsum(A % C, B % C, C, N);
    ll M = (A * N + B) / C;
    return N * M - floorsum(C, C - B - 1, A, M - 1);
    l}
```

#### 5E MillerRabin

```
if ((a = a % n) == 0) return 1;
if (n % 2 == 0) return n == 2;
ll tmp = (n - 1) / ((n - 1) & (1 - n));
ll t = __lg(((n - 1) & (1 - n))), x = 1;
for (; tmp; tmp >>= 1, a = mul(a, a, n))
if (tmp & 1) x = mul(x, a, n);
if (x == 1 || x == n - 1) return 1;
while (--t)
if ((x = mul(x, x, n)) == n - 1) return 1;
return 0;
}
```

### 5F PollardRho

```
|map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2
      == 0) return PollardRho(n / 2), ++cnt[2], void();
  11 x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return;
    }
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
    d = gcd(abs(x - y), n);
| }
```

#### 5G Fraction

```
struct fraction {
  ll n, d;
  fraction
       (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
    ll t = gcd(n, d);
    n /= t, d /= t;
    if (d < 0) n = -n, d = -d;
  fraction operator-() const
  { return fraction(-n, d); }
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator-(const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;
  }
|};
```

# 5H ChineseRemainder

```
| Il solve(Il x1, Il m1, Il x2, Il m2) {
| ll g = gcd(m1, m2);
| if ((x2 - x1) % g) return -1; // no sol
| m1 /= g; m2 /= g;
| ll x, y;
| extgcd(m1, m2, __gcd(m1, m2), x, y);
| ll cm = m1 * m2 * g;
| ll res = x * (x2 - x1) * m1 + x1;
| // be careful with overflow
| return (res % lcm + lcm) % lcm;
| }
```

# 5I Factorial $\mathsf{Mod}p^k$

```
// 0(p^k + log^2 n), pk = p^k
|ll prod[MAXP];
|ll fac_no_p(ll n, ll p, ll pk) {
    prod[0] = 1;
    for (int i = 1; i <= pk; ++i)
        if (i % p) prod[i] = prod[i - 1] * i % pk;
        else prod[i] = prod[i - 1];
        ll rt = 1;
        for (; n; n /= p) {
            rt = rt * mpow(prod[pk], n / pk, pk) % pk;
            rt = rt * prod[n % pk] % pk;
        }
        return rt;
} // (n! without factor p) % p^k</pre>
```

# 5J QuadraticResidue

```
|// Berlekamp-Rabin, log^2(p)
ll trial(ll y, ll z, ll m) {
  ll a0 = 1, a1 = 0, b0 = z, b1 = 1, p = (m - 1) / 2;
  while (p) {
     if (p & 1)
       tie(a0, a1) =
         make_pair((a1 * b1 % m * y + a0 * b0) % m,
           (a0 * b1 + a1 * b0) % m);
     tie(b0, b1) =
       make_pair((b1 * b1 % m * y + b0 * b0) % m,
         (2 * b0 * b1) % m);
    p >>= 1;
  if (a1) return inv(a1, m);
  return -1;
}
mt19937 rd(49);
ll psqrt(ll y, ll p) {
  if (fpow(y, (p - 1) / 2, p) != 1) return -1;
  for (int i = 0; i < 30; i++) {</pre>
    ll z = rd() \% p;
    if (z * z % p == y) return z;
    11 x = trial(y, z, p);
    if (x == -1) continue;
    return x;
  return -1;
}
```

# 5K MeisselLehmer

```
if (n <= 1) return 0;
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector<int> smalls(v + 1), skip(v + 1), roughs(s);
  vector<ll> larges(s);
  for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
  for (int i = 0; i < s; ++i) {</pre>
    roughs[i] = 2 * i + 1;
    larges[i] = (n / (2 * i + 1) + 1) / 2;
  for (int p = 3; p <= v; ++p) {
    if (smalls[p] > smalls[p - 1]) {
      int q = p * p;
      ++pc;
      if (1LL * q * q > n) break;
      skip[p] = 1;
      for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
      int ns = 0;
      for (int k = 0; k < s; ++k) {
        int i = roughs[k];
        if (skip[i]) continue;
        11 d = 1LL * i * p;
        larges[ns] = larges[k] - (d <= v ? larges</pre>
            [smalls[d] - pc] : smalls[n / d]) + pc;
        roughs[ns++] = i;
      }
      s = ns;
      for (int j = v / p; j >= p; --j) {
```

```
smalls[j] - pc, e = min(j * p + p, v + 1);
      for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
  }
}
for (int k = 1; k < s; ++k) {</pre>
  const ll m = n / roughs[k];
  ll t = larges[k] - (pc + k - 1);
  for (int l = 1; l < k; ++l) {</pre>
    int p = roughs[l];
    if (1LL * p * p > m) break;
    t -= smalls[m / p] - (pc + l - 1);
  larges[0] -= t;
}
return larges[0];
```

# DiscreteLog

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered_map<int, int> p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
     y = 1LL * y * x % m;
     b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {</pre>
     s = 1LL * s * b % m;
     if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {</pre>
     if (s == y) return i;
     s = 1LL * s * x % m;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
|}
```

#### 5M Theorems

· Cramer's Rule

Vandermonde's Identity

$$C(n\!+\!m,\!k)\!=\!\sum_{i=0}^k\!C(n,\!i)C(m,\!k\!-\!i)$$

Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}\!=\!d(i)$ ,  $L_{ij}\!=\!-c$  where c is the number of

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\mathsf{det}( ilde{L}_{rr})|$  .
- Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$  , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1,d_2,...,d_n$  for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees. Let  $T_{n,k}$  be the number of labeled forests on n vertices
  - with k components, such that vertex 1,2,...,k belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### Erdős-Gallai Theorem

A sequence of nonnegative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+\cdots+d_n$  is even and  $\sum_{i=1}^{n}d_{i}\!\leq\!k(k\!-\!1)\!+\sum_{i=1}^{n}$  min $(d_{i},\!k)$  holds for every  $1\!\leq\!k\!\leq\!n$  .

i=k+1Gale-Ryser Theorem

A pair of sequences of nonnegative integers  $a_1 \ge \cdots \ge a_n$ and  $b_1, \ldots, b_n$  is bigraphic (degree sequence of bipartie

graph) if and only if 
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$ 

holds for every  $1 \le k \le n$ .

Fulkerson-Chen-Anstee Theorem

A sequence  $(a_1,b_1),...,(a_n,b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic (in, out degree of a di-

rected graph) if and only if 
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and  $\sum_{i=1}^k a_i \leq$ 

$$\sum_{i=1}^k \min(b_i, k-1) + \sum_{i=k+1}^n \min(b_i, k)$$
 holds for every  $1 \leq k \leq n$ .

Möbius Inversion Formula

- $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
- $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Lagrange Multiplier

  - Optimize  $f(x_1,...,x_n)$  when k constraints  $g_i(x_1,...,x_n)=0$ . Lagrangian function  $\mathcal{L}(x_1,\ \dots,\ x_n,\ \lambda_1,\ \dots,\ \lambda_k)=0$ - Lagrangian function  $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)-\sum_{i=1}^k\lambda_ig_i(x_1,\ldots,x_n)$ .
    - The solution corresponding to the original con-
  - strained optimization is always a saddle point of the Lagrangian function.

#### 5N Estimation

### 50 Numbers

• Bernoulli numbers

$$\begin{split} &B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ &\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!} . \\ &S_{m}(n) = \sum_{j=0}^{n} k^m = \frac{1}{n!} \sum_{j=0}^{m} {m+1 \choose j} B_n^{\pm} n^{m+1-k} \end{split}$$

 $S_m(n)=\sum_{k=1}^n k^m=\frac{1}{m+1}\sum_{k=0}^m {m+1\choose k}B_k^+n^{m+1-k}$  • Stirling numbers of the second kind Partitions of ndistinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$
 
$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} \binom{k}{i} i^n$$
 
$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

$$\begin{array}{c} {\displaystyle \mathop{\text{catalan numbers}}_{k=1}^{k=1}} \\ {\displaystyle \mathop{\text{c}}_{n}^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}} \\ {\displaystyle \mathop{\text{c}}_{n}^{(k)}(x) = 1 + x[\mathop{\text{c}}_{n}^{(k)}(x)]^{k}} \end{array}$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element.  $k \ j$ :s s.t.  $\pi(j)\!>\!\pi(j+1)\text{, }k+1\text{ }j\text{:s s.t. }\pi(j)\!\geq\! j\text{, }k\text{ }j\text{:s s.t. }\pi(j)\!>\! j\text{.}$ 

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$
  
 $E(n,0) = E(n,n-1) = 1$ 

 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$ 

#### GeneratingFunctions

- Ordinary Generating Function  $A(x) = \sum_{i>0} a_i x^i$ 
  - $A(rx) \Rightarrow r^n a_n$

  - $A(x)+B(x)\Rightarrow a_n+b_n$   $A(x)B(x)\Rightarrow \sum_{i=0}^n a_ib_{n-i}$
  - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $xA(x)' \Rightarrow na_n$
  - $\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function  $A(x) = \sum_{i \ge 0} \frac{a_i}{i!} x_i$

```
- A(x)+B(x) \Rightarrow a_n+b_n
     -A^{(k)}(x) \Rightarrow a_{n+k}
      - A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}
     - A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}
       - xA(x) \Rightarrow na_n
• Special Generating Function
      - (1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i
      -\frac{1}{(1-x)^n} = \sum_{i>0} \binom{i}{n-1} x^i
      - S_k = \sum_{x=1}^n x^k: S = \sum_{p=0}^\infty x^p = \frac{e^x - e^{x(n+1)}}{1 - e^x}
```

# Linear Algebra 6A GuassianElimination

```
struct matrix { // m variables, n equations
   int n, m;
  fraction A[N][N + 1], sol[N];
  int solve() { //-1: inconsistent, >= 0: rank
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !A[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
       for (int j = 0; j < n; ++j) {</pre>
         if (i == j) continue;
         fraction tmp = -A[j][piv] / A[i][piv];
         for (int k = 0; k <= m; ++k)</pre>
           A[j][k] = tmp * A[i][k] + A[j][k];
       }
     int rank = 0;
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !A[i][piv].n) ++piv;</pre>
       if (piv == m && A[i][m].n) return -1;
       else if (piv < m)</pre>
         ++rank, sol[piv] = A[i][m] / A[i][piv];
     return rank;
  }
|};
```

# BerlekampMassey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(output.size() + 1), me, he;
   for (int f = 0, i = 1; i <= output.size(); ++i) {</pre>
     for (int j = 0; j < me.size(); ++j)</pre>
      d[i] += output[i - j - 2] * me[j];
     if ((d[i] -= output[i - 1]) == 0) continue;
     if (me.empty()) {
       me.resize(f = i);
       continue;
    vector<T> o(i - f - 1);
    T k = -d[i] / d[f];
     o.emplace_back(-k);
     for (T x : he) o.emplace_back(x * k);
     o.resize(max(o.size(), me.size()));
     for (int j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
     if (i - f + (int
         )he.size() >= (int)me.size()) he = me, f = i;
    me = o;
  }
   return me;
| }
```

# Simplex

Standard form: maximize  $\mathbf{c}^T \mathbf{x}$  subject to  $A \mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq \mathbf{0}$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $ar{x}_i\!=\!0$  or  $\sum_{j=1}^m\!A_{ji}ar{y}_j\!=\!c_i$  holds and for all  $i\!\in\![1,m]$  either  $ar{y}_i\!=\!0$ or  $\sum_{j=1}^{n} A_{ij} \bar{x}_j = b_j$  holds.

```
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
```

 $3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j$ 

```
1. In case of minimization, let c_i' = -c_i
     • \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
```

```
• \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
4. If x_i has no lower bound, replace x_i with x_i - x_i'
// n variable, m constraints, M >= n + 2m
struct simplex {
  const double inf = 1 / .0, eps = 1e-9;
   int n, m, k, var[N], inv[N], art[N];
   double A[M][N], B[M], x[N];
   void init(int _n) { n = _n, m = 0; }
   void equation(vector<double> a, double b) {
     for (int i = 0; i < n; i++) A[m][i] = a[i];</pre>
     B[m] = b, var[m] = n + m, ++m;
   void pivot(int r, int c, double bx) {
     for (int i = 0; i <= m + 1; i++)</pre>
       if (i != r && abs(A[i][c]) > eps) {
         x[var[i]] = bx * A[i][c] / A[i][var[i]];
         double f = A[i][c] / A[r][c];
         for (int j = 0; j <= n + m + k; j++)</pre>
           A[i][j] -= A[r][j] * f;
         B[i] -= B[r] * f;
       }
   double phase(int p) {
     while (true) {
       int in = (int)(min_element(A[m + p],
         A[m + p] + n + m + k + 1) - A[m + p]);
       if (A[m + p][in] >= -eps) break;
       double bx = inf;
       int piv = -1;
       for (int i = 0; i < m; i++)</pre>
         if (A[i][in] > eps && B[i] / A[i][in] <= bx)</pre>
           piv = i, bx = B[i] / A[i][in];
       if (piv == -1) return inf;
       int out = var[piv];
       pivot(piv, in, bx);
       x[out] = 0, x[in] = bx, var[piv] = in;
     return x[n + m];
  }
   double solve(vector<double> c) {
     auto invert = [&](int r) {
       for (int j = 0; j <= n + m; j++) A[r][j] *= -1;</pre>
       B[r] *= -1;
     k = 1;
     for (int i = 0; i < n; i++) A[m][i] = -c[i];</pre>
     fill(A[m + 1], A[m + 1] + N, 0.0);
     for (int i = 0; i <= m + 1; i++)</pre>
       fill(A[i] + n, A[i] + n + m + 2, 0.0),
         var[i] = n + i, A[i][n + i] = 1;
     for (int i = 0; i < m; i++) {</pre>
       if (B[i] < 0) {
         ++k;
         for (int j = 0; j <= n + m; j++)
           A[m + 1][j] += A[i][j];
         invert(i);
         var[i] = n + m + k, A[i][var[i]] = 1,
         art[var[i]] = n + i;
       x[var[i]] = B[i];
     phase(1);
     if (*max_element(
           x + (n + m + 2), x + (n + m + k + 1)) > eps)
       return .0 / .0;
     for (int i = 0; i <= m; i++)</pre>
       if (var[i] > n + m)
         var[i] = art[var[i]], invert(i);
     return phase(0);
|} lp;
```

# **Polynomials** 7A NTT (FFT)

```
Form
                                     2^{16} + 1
                       65 537
                                     119 \cdot 2^{23} + 1
                  998 244 353
                                3
                                     1255 \cdot 2^{20} + 1
                1 315 962 881
                                3
                                     51 \cdot 2^{25} + 1
                1 711 276 033
                                29
                                     549755813881 \!\cdot\! 2^{24} \!+\! 1
   9 223 372 036 737 335 297
#define base ll // complex<double>
// const double PI = acosl(-1);
const ll mod = 998244353, g = 3;
base omega[4 * N], omega_[4 * N];
int rev[4 * N];
ll fpow(ll b, ll p);
ll inverse(ll a) { return fpow(a, mod - 2); }
void calcW(int n) {
 ll r = fpow(g, (mod - 1) / n), invr = inverse(r);
  omega[0] = omega_[0] = 1;
  for (int i = 1; i < n; i++) {</pre>
    omega[i] = omega[i - 1] * r % mod;
    omega_[i] = omega_[i - 1] * invr % mod;
 // double arg = 2.0 * PI / n;
  // for (int i = 0; i < n; i++)
 // {
 //
      omega[i] = base(cos(i * arg), sin(i * arg));
 //
       omega_[i] = base(cos(-i * arg), sin(-i * arg));
  // }
void calcrev(int n) {
  int k = __lg(n);
  for (int i = 0; i < n; i++) rev[i] = 0;</pre>
  for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < k; j++)</pre>
      if (i & (1 << j)) rev[i] ^= 1 << (k - j - 1);</pre>
vector<base> NTT(vector<base> poly, bool inv) {
  base *w = (inv ? omega_ : omega);
  int n = (int)poly.size();
  for (int i = 0; i < n; i++)</pre>
    if (rev[i] > i) swap(poly[i], poly[rev[i]]);
  for (int len = 1; len < n; len <<= 1) {</pre>
    int arg = n / len / 2;
    for (int i = 0; i < n; i += 2 * len)</pre>
      for (int j = 0; j < len; j++) {</pre>
        base odd =
          w[j * arg] * poly[i + j + len] % mod;
        poly[i + j + len] =
           (poly[i + j] - odd + mod) % mod;
        poly[i + j] = (poly[i + j] + odd) \% mod;
      }
 }
  if (inv)
    for (auto &a : poly) a = a * inverse(n) % mod;
  return poly;
vector<base> mul(vector<base> f, vector<base> g) {
 int sz = 1 << (__lg(f.size() + g.size() - 1) + 1);</pre>
  f.resize(sz), g.resize(sz);
 calcrev(sz);
 calcW(sz);
 f = NTT(f, 0), g = NTT(g, 0);
  for (int i = 0; i < sz; i++)</pre>
    f[i] = f[i] * g[i] % mod;
  return NTT(f, 1);
```

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
    for (int i = 0; i < n; i += L)</pre>
      for (int j = i; j < i + (L >> 1); ++j)
        a[j + (L >> 1)] += a[j] * op;
const int P = 21; // power of max N
int f[
    P][1 << P], g[P][1 << P], h[P][1 << P], ct[1 << P];
void
    subset_convolution(int *a, int *b, int *c, int L) {
  // c_k = \sum_{i = 0} a_i * b_j
  int n = 1 << L;
  for (int i = 1; i < n; ++i)
    ct[i] = ct[i \& (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)</pre>
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)
    for (int j = 0; j <= i; ++j)</pre>
      for (int x = 0; x < n; ++x)
        h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i)</pre>
    fwt(h[i], n, -1);
  for (int i = 0; i < n; ++i)</pre>
    c[i] = h[ct[i]][i];
```

# 7C PolynomialOperations

```
|#define poly vector<ll>
poly inv(poly A) {
  A.resize(1 << (__lg(A.size() - 1) + 1));
  poly B = {inverse(A[0])};
  for (int n = 1; n < (int)A.size(); n <<= 1) {</pre>
     poly pA(A.begin(), A.begin() + 2 * n);
    calcrev(4 * n), calcW(4 * n);
    pA.resize(4 * n), B.resize(4 * n);
     pA = NTT(pA, 0);
     B = NTT(B, 0);
     for (int i = 0; i < 4 * n; i++)</pre>
       B[i] =
         ((B[i] * 2 - pA[i] * B[i] % mod * B[i]) % mod +
           mod) %
         mod:
     B = NTT(B, 1);
    B.resize(2 * n);
  return B;
}
pair<poly, poly> div(poly A, poly B) {
  if (A.size() < B.size()) return make_pair(poly(), A);</pre>
  int n = A.size(), m = B.size();
  poly revA = A, invrevB = B;
  reverse(all(revA)), reverse(all(invrevB));
  revA.resize(n - m + 1);
  invrevB.resize(n - m + 1);
  invrevB = inv(invrevB);
  poly Q = mul(revA, invrevB);
  Q.resize(n - m + 1);
  reverse(all(Q));
  poly R = mul(Q, B);
  R.resize(m - 1);
  for (int i = 0; i < m - 1; i++)</pre>
    R[i] = (A[i] - R[i] + mod) \% mod;
  return make_pair(Q, R);
}
poly modulo(poly A, poly B) { return div(A, B).S; }
| ll fast_kitamasa(ll k, poly A, poly C) {
  int n = A.size();
  C.emplace_back(mod - 1);
```

```
poly Q, R = \{0, 1\}, F = \{1\};
   R = modulo(R, C);
  for (; k; k >>= 1) {
    if (k & 1) F = modulo(mul(F, R), C);
     R = modulo(mul(R, R), C);
     k >>= 1;
   11 ans = 0;
   for (int i = 0; i < F.size(); i++)</pre>
     ans = (ans + A[i] * F[i]) % mod;
   return ans;
}
 vector<ll> fpow(vector<ll> f, ll p, ll m) {
  int b = 0;
  while (b < f.size() && f[b] == 0) b++;</pre>
  f = vector<ll>(f.begin() + b, f.end());
  int n = f.size();
  f.emplace_back(0);
  vector<ll> q(min(m, b * p), 0);
   q.emplace_back(fpow(f[0], p));
   for (int k = 0; q.size() < m; k++) {</pre>
    11 res = 0;
     for (int i = 0; i < min(n, k + 1); i++)</pre>
       res = (res +
               p * (i + 1) % mod * f[i + 1] % mod *
                 q[k - i + b * p]) %
         mod:
     for (int i = 1; i < min(n, k + 1); i++)</pre>
       res = (res -
               f[i] * (k - i + 1) % mod *
                 q[k - i + 1 + b * p]) %
         mod:
     res = (res < 0 ? res + mod : res) *
       inv(f[0] * (k + 1) % mod) % mod;
     q.emplace_back(res);
   return q;
| }
```

#### 7D NewtonMethod+MiscGF

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)\!=\!0\pmod{x^{2^k}}$ , then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$

- $A^{-1}$ :  $B_{k+1} = B_k(2 AB_k) \mod x^{2^{k+1}}$
- $\ln A$ :  $(\ln A)' = \frac{A'}{A}$
- $\exp A$ :  $B_{k+1} = B_k(1 + A \ln B_k) \mod x^{2^{k+1}}$
- $\sqrt{A}$ :  $B_{k+1} = \frac{1}{2}(B_k + AB_k^{-1}) \mod x^{2^{k+1}}$

# 8 Geometry

#### 8A Basic

```
typedef pair<pdd, pdd> Line;
struct Cir{ pdd 0; double R; };
const double pi = acos(-1);
const double eps = 1e-8;
pll operator+(pll a, pll b)
{ return pll(a.F + b.F, a.S + b.S); }
pll operator-(pll a, pll b)
{ return pll(a.F - b.F, a.S - b.S); }
pll operator-(pll a)
{ return pll(-a.F, -a.S); }
pll operator*(pll a, ll b)
{ return pll(a.F * b, a.S * b); }
pdd operator/(pll a, double b)
{ return pdd(a.F / b, a.S / b); }
ll dot(pll a, pll b)
{ return a.F * b.F + a.S * b.S; }
| ll cross(pll a, pll b)
```

```
{ return a.F * b.S - a.S * b.F; }
ll abs2(pll a)
{ return dot(a, a); }
double abs(pll a)
{ return sqrt(dot(a, a)); }
int sign(ll a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pll a, pll b, pll c)
{ return sign(cross(b - a, c - a)); }
bool collinearity(pll p1, pll p2, pll p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pll a, pll b, pll c) {
   return collinearity
        (a, b, c) \&\& sign(dot(a - c, b - c)) <= 0;
bool seg_strict_intersect
     (pdd p1, pdd p2, pdd p3, pdd p4) {
   int a123 = ori(p1, p2, p3);
   int a124 = ori(p1, p2, p4);
   int a341 = ori(p3, p4, p1);
   int a342 = ori(p3, p4, p2);
return a123 * a124 < 0 && a341 * a342 < 0;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
   int a123 = ori(p1, p2, p3);
   int a124 = ori(p1, p2, p4);
   int a341 = ori(p3, p4, p1);
   int a342 = ori(p3, p4, p2);
   if (a123 == 0 && a124 == 0)
     return btw(p1, p2, p3) || btw(p1, p2, p4) ||
btw(p3, p4, p1) || btw(p3, p4, p2);
   return a123 * a124 <= 0 && a341 * a342 <= 0;
}
\textbf{pdd} \hspace{0.1cm} \textbf{intersect(pdd} \hspace{0.1cm} \textbf{p1,} \hspace{0.1cm} \textbf{pdd} \hspace{0.1cm} \textbf{p2,} \hspace{0.1cm} \textbf{pdd} \hspace{0.1cm} \textbf{p3,} \hspace{0.1cm} \textbf{pdd} \hspace{0.1cm} \textbf{p4)} \hspace{0.1cm} \{
   double a123 = cross(p2 - p1, p3 - p1);
   double a124 = cross(p2 - p1, p4 - p1);
   return (p4
        * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
pdd orth(pdd p1)
{ return pdd(-p1.S, p1.F); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
     p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + orth(p2 - p1
     ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
     (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) {
   pdd dp = p1 - p0
        , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
   return q0 + pdd(
        cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
|} // from line p0--p1 to q0--q1, apply to r
     ConvexHull
void hull(vector<pll> &dots) { // n=1 => ans = {}
```

### 8C SortByAngle

```
| bool down(pll k) {
    return sign(k.S) < 0 ||
        (sign(k.S) == 0 && sign(k.F) < 0);
    }
    int cmp(pll a, pll b, bool same = true) {
```

```
int A = down(a), B = down(b);
if (A != B) return A < B;
if (sign(cross(a, b)) == 0)
  return same ? abs2(a) < abs2(b) : -1;
return sign(cross(a, b)) > 0;
}
```

#### **8D Formulas**

Rotation

$$M(\theta) = \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix}$$

90 degree: (x,y) = (Y-y,x)

· Pick's theorem

For simple integer-coordinate polygon,

$$A = \mathsf{B} + \frac{I}{2} - 1$$

Where A is the area; B,I is #lattice points in the interior, on the boundary.

· Spherical Cap

- A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :  $\arcsin(a/r)$ .
- Volume  $= \pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos\theta)(1 \cos\theta)^2/3$ .
- Area  $=2\pi rh=\pi(a^2+h^2)=2\pi r^2(1-\cos\theta)$  .
- · Nearest points of two skew lines
  - Line 1: ${m v}_1\!=\!{m p}_1\!+\!t_1{m d}_1$
  - Line 2: $m{v}_2 = m{p}_2 + t_2 m{d}_2$
  - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
  - $\boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}$
  - $\boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$
  - $c_1 = p_1 + \frac{(p_2 p_1) \cdot n_2}{d_1 \cdot n_2} d_1$
  - $c_2 = p_2 + \frac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2$

# 8E TriangleHearts

```
pdd excenter(
  pdd p0, pdd p1, pdd p2) { // radius = abs(center)
p1 = p1 - p0, p2 = p2 - p0;
  auto [x1, y1] = p1;
  auto [x2, y2] = p2;
  double m = 2. * cross(p1, p2);
  pdd center = pdd((x1 * x1 * y2 - x2 * x2 * y1 +
                      y1 * y2 * (y1 - y2)),
                  (x1 * x2 * (x2 - x1) - y1 * y1 * x2 +
                    x1 * y2 * y2)) /
  return center + p0;
pdd incenter(
  pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
  double a = abs(p2 - p3), b = abs(p1 - p3),
         c = abs(p1 - p2);
  double s = a + b + c;
  return (p1 * a + p2 * b + p3 * c) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3) {
  return (p1 + p2 + p3) / 3;
pdd orthcenter(pdd p1, pdd p2, pdd p3) {
  return masscenter(p1, p2, p3) * 3
    excenter(p1, p2, p3) * 2;
```

#### 8F PointSegmentDist

```
|double PointSegDist(pdd q0, pdd q1, pdd p) {
| if (abs(q0 - q1) <= eps) return abs(q0 - p);
| if (dot(q1 - q0,
| p - q0) >= -eps && dot(q0 - q1, p - q1) >= -eps)
| return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
| return min(abs(p - q0), abs(p - q1));
|}
```

#### 8G PointInCircle

```
// return q'
    s relation with circumcircle of tri(p[0],p[1],p[2])
bool in_cc(const array<pll, 3> p, pll q) {
    __int128 det = 0;
    for (int i = 0; i < 3; ++i)
        det += __int128(abs2(p[i]) - abs2(q)) *
            cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

#### 8H PointInConvex

# 8I PointTangentConvex

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
/* bool pred(int a, int b);
f(0) \sim f(n-1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - l > 1) {
     int m = (l + r) / 2;
     if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
     else l = m;
   return pred(l, r % n) ? l : r % n;
}
pii get_tangent(vector<pll> &C, pll p) {
   auto gao = [&](int s) {
     return cyc_tsearch((int)C.size(), [&](int x, int y)
     { return ori(p, C[x], C[y]) == s; });
  return pii(gao(1), gao(-1));
|} // return (a, b), ori(p, C[a], C[b]) >= 0
```

# 8J CircTangentCirc

```
vector<Line> go(Cir c1, Cir c2, int sign1) {
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = abs2(c1.0 - c2.0);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
    pdd n = pdd(v.F * c - sign2 * h * v.S,
      v.S * c + sign2 * h * v.F);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sign(p1.F - p2.F) == 0 and
      sign(p1.S - p2.S) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
    ret.emplace_back(Line(p1, p2));
```

```
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  return ret:
| }
8K LineCircleIntersect
vector<pdd> circleLine(pdd c, double r, pdd a, pdd b) {
  pdd p
                                                             if (a > r) {
        = a + (b - a) * dot(c - a, b - a) / abs2(b - a);
   double s = cross
       (b - a, c - a), h2 = r * r - s * s / abs2(b - a);
  if (h2 < 0) return {};
  if (h2 == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return {p - h, p + h};
      LineConvexIntersect
int cyc_tsearch(int n, auto pred); // ref: TanPointHull
int TangentDir(vector<pll>> &C, pll dir) {
   return cyc_tsearch((int)C.size(), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
                                                             double S = 0;
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
                                                                 ori(
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = (int)C.size();
  if (cmpL(A) < 0 \mid \mid cmpL(B) > 0)
    return pii(-1, -1); // no collision
                                                           80
  auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r;) {
      int m = ((l + r + (l < r ? 0 : n)) / 2) % n;</pre>
                                                             return sign
       (cmpL(m) == cmpL(t) ? l : r) = m;
    }
    return (l + !cmpL(r)) % n;
  };
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.F == res.S) // touching the corner i
    return pii(res.F, -1);
  if (!cmpL(res.F) &&
                                                                 vector
    !cmpL(res.S)) // along side i, i+1
    switch ((res.F - res.S + n + 1) % n) {
    case 0: return pii(res.F, res.F);
    case 2: return pii(res.S, res.S);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned in the same order as the line hits the
  convex */
  return res;
} // convex cut: (r, l]
                                                                     }
      CircIntersectCirc
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
  pdd o1 = a.0, o2 = b.0;
   double r1 =
        a.R, r2 = b.R, d2 = abs2(o1 - o2), d = sqrt(d2);
  if(d < max
                                                                   }
       (r1, r2) - min(r1, r2) \mid \mid d > r1 + r2) return 0;
  pdd \ U = (o1 + o2) * 0.5
       + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
  double A = sqrt((r1 + r2 + d) *
        (r1 - r2 + d) * (r1 + r2 - d) * (-r1 + r2 + d));
        = pdd(o1.S - o2.S, -o1.F + o2.F) * A / (2 * d2);
  p1 = v + v, p2 = v - v;
   return 1;
|}
                                                               }
8N
      PolyIntersectCirc
                                                          |}
// Divides into multiple triangle, and sum up
```

const double PI = acos(-1);

double \_area(pdd pa, pdd pb, double r) {

```
if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  if (abs(pb) < eps) return 0;</pre>
  double S, h, theta;
  double a = abs(pb), b = abs(pa), c = abs(pb - pa);
  double cosB = dot(pb, pb - pa) / a / c,
         B = acos(cosB);
  double cosC = dot(pa, pb) / a / b, C = acos(cosC);
    S = (C / 2) * r * r;
    h = a * b * sin(C) / c;
    if (h < r && B < PI / 2)
      S = (acos(h / r) * r * r -
        h * sqrt(r * r - h * h));
  } else if (b > r) {
    theta = PI - B - asin(sin(B) / r * a);
    S = .5 * a * r * sin(theta) +
      (C - theta) / 2 * r * r;
  } else S = .5 * sin(C) * a * b;
double area_poly_circle(const vector<pdd> poly,
  const pdd &0, const double r) {
  for (int i = 0; i < (int)poly.size(); ++i)</pre>
    S += _area(poly[i] - 0,
           poly[(i + 1) % (int)poly.size()] - 0, r) *
        0, poly[i], poly[(i + 1) % (int)poly.size()]);
  return fabs(S);
    PolyUnion
double rat(pll a, pll b) {
      (b.F) ? (double)a.F / b.F : (double)a.S / b.S;
} // all poly. should be ccw
double polyUnion(vector<vector<pll>>> &poly) {
  double res = 0;
  for (auto &p : poly)
    for (int a = 0; a < (int)p.size(); ++a) {</pre>
      pll A = p[a], B = p[(a + 1) % (int)p.size()];
          <pair<double, int>> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < (int)q.size(); ++b) {</pre>
          pll C = q[b], D = q[(b + 1) \% (int)q.size()];
          int sc = ori(A, B, C), sd = ori(A, B, D);
          if (sc != sd && min(sc, sd) < 0) {</pre>
            double sa = cross(D
                  - C, A - C), sb = cross(D - C, B - C);
            segs.emplace_back
                 (sa / (sa - sb), sign(sc - sd));
          if (!sc && !sd &&
              &q < &p && sign(dot(B - A, D - C)) > 0) {
            segs.emplace_back(rat(C - A, B - A), 1);
            segs.emplace_back(rat(D - A, B - A), -1);
      sort(all(segs));
      for (auto &s : segs) s.F = clamp(s.F, 0.0, 1.0);
      double sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < (int)segs.size(); ++j) {</pre>
        if (!cnt) sum += segs[j].F - segs[j - 1].F;
        cnt += segs[j].S;
      res += cross(A, B) * sum;
  return res / 2;
```

# **8P MinkowskiSum**

```
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vector<pll> Minkowski
     (vector<pll> A, vector<pll> B) \{ // |A|, |B| >= 3 \}
   hull(A), hull(B);
  vector<pll> C(1, A[0] + B[0]), s1, s2;
  for (int i = 0; i < A.size(); ++i)</pre>
     s1.emplace_back(A[(i + 1) % A.size()] - A[i]);
  for (int i = 0; i < B.size(); i++)</pre>
     s2.emplace_back(B[(i + 1) % B.size()] - B[i]);
  for (int i = 0, j = 0; i < A.size() || j < B.size();)</pre>
    if (j >= B.size()
          || (i < A.size() \&\& cross(s1[i], s2[i]) >= 0))
      C.emplace_back(B[j % B.size()] + A[i++]);
     else
      C.emplace_back(A[i % A.size()] + B[j++]);
  return hull(C), C;
|}
8Q MinMaxEnclosingRect
const double qi = acos(-1) / 2 * 3;
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
  hull(dots):
   double Max = 0, Min = INF, deg;
  int n = (int)dots.size();
   dots.emplace_back(dots[0]);
  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
     pll nw = vec(i + 1);
     while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
```

while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))

while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))</pre>

Min = min(Min, (double)(dot(nw, vec(r)) - dot

(nw, vec(l))) \* cross(nw, vec(u)) / abs2(nw));

, l), vec(u) / abs(diff(r, l)) / abs(vec(u)));

(r, l)) \* abs(vec(u)) \* sin(deg) \* sin(deg));

# MinEnclosingCircle

u = (u + 1) % n;

r = (r + 1) % n;

l = (l + 1) % n;

deg = acos(dot(diff(r

deg = (qi - deg) / 2;

return pdd(Min, Max);

}

| }

Max = max(Max, abs(diff))

if (!i) l = (r + 1) % n;

```
pdd Minimum_Enclosing_Circle
    (vector<pdd> dots, double &r) {
  pdd cent;
  random_shuffle(all(dots));
  cent = dots[0], r = 0;
  for (int i = 1; i < (int)dots.size(); ++i)</pre>
    if (abs(dots[i] - cent) > r) {
      cent = dots[i], r = 0;
      for (int j = 0; j < i; ++j)</pre>
        if (abs(dots[j] - cent) > r) {
          cent = (dots[i] + dots[j]) / 2;
          r = abs(dots[i] - cent);
          for(int k = 0; k < j; ++k)</pre>
            if(abs(dots[k] - cent) > r)
                    excenter(dots[i], dots[j], dots[k]),
               r = abs(cent - dots[i]);
        }
  return cent;
```

#### 8S CircleCover

```
// N ~= 1000
struct CircleCover {
  int C;
  Cir c[N];
```

```
bool g[N][N], overlap[N][N];
   // Area[i] : area covered by at least i circles
   double Area[ N ];
   void init(int _c){ C = _c;}
   struct Teve {
     pdd p; double ang; int add;
     Teve() {}
     Teve(pdd _a
         , double _b, int _c):p(_a), ang(_b), add(_c){}
     bool operator<(const Teve &a)const
     {return ang < a.ang;}
   eve[N * 2];
   // strict: x = 0, otherwise x = -1
   bool disjuct(Cir &a, Cir &b, int x)
   \{ return \ sign(abs(a.0 - b.0) - a.R - b.R) > x; \}
   bool contain(Cir &a, Cir &b, int x)
   {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
   bool contain(int i, int j) {
     /* c[j] is non-strictly in c[i]. */
     return (sign
         (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
         R) == 0 \& i < j) && contain(c[i], c[j], -1);
   void solve(){
     fill_n(Area, C + 2, 0);
     for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         overlap[i][j] = contain(i, j);
     for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
             disjuct(c[i], c[j], -1));
     for(int i = 0; i < C; ++i){</pre>
       int E = 0, cnt = 1;
       for(int j = 0; j < C; ++j)</pre>
         if(j != i && overlap[j][i])
           ++cnt;
       for(int j = 0; j < C; ++j)</pre>
         if(i != j && g[i][j]) {
           pdd aa, bb;
           CCinter(c[i], c[j], aa, bb);
           double A =
                 atan2(aa.S - c[i].O.S, aa.F - c[i].O.F);
           double B =
                 atan2(bb.S - c[i].O.S, bb.F - c[i].O.F);
           eve[E++] = Teve
                (bb, B, 1), eve[E++] = Teve(aa, A, -1);
           if(B > A) ++cnt;
         ን
       if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
       else{
         sort(eve, eve + E);
         eve[E] = eve[0];
         for(int j = 0; j < E; ++j){</pre>
           cnt += eve[j].add;
           Arealcnt
                ] += cross(eve[j].p, eve[j + 1].p) * .5;
           double theta = eve[j + 1].ang - eve[j].ang;
           if (theta < 0) theta += 2. * pi;</pre>
           Area[cnt] += (theta
                 - sin(theta)) * c[i].R * c[i].R * .5;
         }
  }
|};
      LineCmp
```

```
struct lineCmp { // coordinates should be even!
  bool operator()(Line l1, Line l2) const {
    int X :
      (\max(l1.F.F, l2.F.F) + \min(l1.S.F, l2.S.F)) / 2;
    ll p1 =
         (X - l1.F.F) * l1.S.S + (l1.S.F - X) * l1.F.S,
       p2 =
         (X - 12.F.F) * 12.S.S + (12.S.F - X) * 12.F.S,
```

```
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       q1 = (l1.S.F - l1.F.F), q2 = (l2.S.F - l2.F.F);
    if (q1 == 0) p1 = l1.F.S + l1.S.S, q1 = 2;
    if (q2 == 0) p2 = l2.F.S + l2.S.S, q2 = 2;
    // for query a point: ask make_pair(P, P)
    if (l1.F == l2.F || l2.F == l2.S) l1 = l2;
    return make_tuple((__int128)(p1 * q2), l1) <</pre>
      make_tuple((\_int128)(p2 * q1), l2);
|};
8U
     Trapezoidalization
template<class T>
struct SweepLine {
  struct cmp {
    cmp(const SweepLine &_swp): swp(_swp) {}
    bool operator()(int a, int b) const {
```

```
if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
      return swp.slope_cmp(a, b);
    return swp.get_y(a) + swp.eps < swp.get_y(b);</pre>
  const SweepLine &swp;
 _cmp;
T curTime, eps, curQ;
vector<Line> base;
multiset<int, cmp> sweep;
multiset<pair<T, int>> event;
vector<typename multiset<int, cmp>::iterator> its;
    <typename multiset<pair<T, int>>::iterator> eits;
bool slope_cmp(int a, int b) const {
  assert(a != -1);
  if (b == -1) return 0:
  return sign(cross(base
      [a].S - base[a].F, base[b].S - base[b].F)) < 0;
T get_y(int idx) const {
  if (idx == -1) return curQ;
  Line l = base[idx];
  if (l.F.F == l.S.F) return l.S.S;
  return ((curTime - l.F.F) * l.S.S
      + (l.S.F - curTime) * l.F.S) / (l.S.F - l.F.F);
}
void insert(int idx) {
  its[idx] = sweep.insert(idx);
  if (its[idx] != sweep.begin())
    update_event(*prev(its[idx]));
  update_event(idx);
  event.emplace
      (base[idx].S.F, idx + 2 * (int)base.size());
void erase(int idx) {
  assert(eits[idx] == event.end());
  auto p = sweep.erase(its[idx]);
  its[idx] = sweep.end();
  if (p != sweep.begin())
    update_event(*prev(p));
void update_event(int idx) {
  if (eits[idx] != event.end())
    event.erase(eits[idx]);
  eits[idx] = event.end();
  auto nxt = next(its[idx]);
  if (nxt ==
       sweep.end() || !slope_cmp(idx, *nxt)) return;
  auto t = intersect(base[idx].
      F, base[idx].S, base[*nxt].F, base[*nxt].S).F;
  if (t + eps < curTime || t</pre>
       >= min(base[idx].S.F, base[*nxt].S.F)) return;
  eits[idx
      ] = event.emplace(t, idx + (int)base.size());
}
void swp(int idx) {
  assert(eits[idx] != event.end());
  eits[idx] = event.end();
  int nxt = *next(its[idx]);
  swap((int&)*its[idx], (int&)*its[nxt]);
```

```
swap(its[idx], its[nxt]);
     if (its[nxt] != sweep.begin())
       update_event(*prev(its[nxt]));
     update_event(idx);
   // only expected to call the functions below
   SweepLine(T t, T e, vector<Line> vec): _cmp
       (*this), curTime(t), eps(e), curQ(), base(vec),
        sweep(_cmp), event(), its((int)vec.size(), sweep
       .end()), eits((int)vec.size(), event.end()) {
     for (int i = 0; i < (int)base.size(); ++i) {</pre>
       auto &[p, q] = base[i];
       if (p > q) swap(p, q);
       if (p.F <= curTime && curTime <= q.F)</pre>
         insert(i);
       else if (curTime < p.F)</pre>
         event.emplace(p.F, i);
     }
  }
   void setTime(T t, bool ers = false) {
     assert(t >= curTime);
     while (!event.empty() && event.begin()->F <= t) {</pre>
       auto [et, idx] = *event.begin();
       int s = idx / (int)base.size();
       idx %= (int)base.size();
       if (abs(et - t) <= eps && s == 2 && !ers) break;</pre>
       curTime = et;
       event.erase(event.begin());
       if (s == 2) erase(idx);
       else if (s == 1) swp(idx);
       else insert(idx);
     curTime = t;
   T nextEvent() {
     if (event.empty()) return INF;
     return event.begin()->F;
   int lower_bound(T y) {
     curQ = y;
     auto p = sweep.lower_bound(-1);
     if (p == sweep.end()) return -1;
     return *p;
  }
|};
      HalfPlaneIntersect
8V
```

```
pll area_pair(Line a, Line b)
{ return pll(cross(a.S
      - a.F, b.F - a.F), cross(a.S - a.F, b.S - a.F)); }
bool isin(Line l0, Line l1, Line l2) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(l0, l2);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
  return (__int128)
        a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
|}
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(all(arr), [&](Line a, Line b) -> int {
    if (cmp(a.S - a.F, b.S - b.F, 0) != -1)
      return cmp(a.S - a.F, b.S - b.F, 0);
    return ori(a.F, a.S, b.S) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(
         dq.back().S - dq.back().F, p.S - p.F, 0) == -1)
       continue;
    while ((int)dq.size() >= 2
         && !isin(p, dq[(int)dq.size() - 2], dq.back()))
       dq.pop_back();
```

```
while
         ((int)dq.size() >= 2 \&\& !isin(p, dq[0], dq[1]))
       dq.pop_front();
    dq.emplace_back(p);
  while ((int)dq.size() >= 3 &&
        !isin(dq[0], dq[(int)dq.size() - 2], dq.back()))
     dq.pop_back();
  while ((int)
       dq.size() >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
    dq.pop_front();
  return vector<Line>(all(dq));
|}
```

# RotatingSweepLine

```
void rotatingSweepLine(vector<pii> &ps) {
  int n = (int)ps.size(), m = 0;
  vector<int> id(n), pos(n);
  vector<pii> line(n * (n - 1));
  for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++j)</pre>
       if (i != j) line[m++] = pii(i, j);
  sort(all(line), [&](pii a, pii b) {
     return cmp(ps[a.S] - ps[a.F], ps[b.S] - ps[b.F]);
  }); // cmp(): polar angle compare
  iota(all(id), 0);
  sort(all(id), [&](int a, int b) {
     if (ps[a].S != ps[b].S) return ps[a].S < ps[b].S;</pre>
     return ps[a] < ps[b];</pre>
  }); // initial order, since (1, 0) is the smallest
   for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
  for (int i = 0; i < m; ++i) {</pre>
    auto l = line[i];
     // do something
     tie(pos[l.F], pos[l.S], id[pos[l.F]], id[pos[l.S
         ]]) = make_tuple(pos[l.S], pos[l.F], l.S, l.F);
  }
| }
```

# 8X DelaunayTriangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle. */
struct Edge {
  int id; // oidx[id]
  list<Edge>::iterator twin;
  Edge(int _id = 0) : id(_id) {}
struct Delaunay { // O-base
 int n, oidx[N];
 list<Edge> head[N]; // result udir. graph
  pll p[N];
  void init(int _n, pll _p[]) {
    n = _n, iota(oidx, oidx + n, 0);
    for (int i = 0; i < n; ++i) head[i].clear();</pre>
    sort(oidx, oidx + n,
      [&](int a, int b) { return _p[a] < _p[b]; });</pre>
    for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];</pre>
    divide(0, n - 1);
 void addEdge(int u, int v) {
    head[u].push_front(Edge(v));
    head[v].push_front(Edge(u));
    head[u].begin()->twin = head[v].begin();
    head[v].begin()->twin = head[u].begin();
 }
 void divide(int l, int r) {
    if (l == r) return;
    if (l + 1 == r) return addEdge(l, l + 1);
    int mid = (l + r) >> 1, nw[2] = \{l, r\};
    divide(l, mid), divide(mid + 1, r);
    auto gao = [&](int t) {
      pll pt[2] = {p[nw[0]], p[nw[1]]};
```

```
for (auto it : head[nw[t]]) {
         int v = ori(pt[1], pt[0], p[it.id]);
         if (v > 0 ||
           (v == 0 \&\&
             abs2(pt[t ^ 1] - p[it.id]) <
               abs2(pt[1] - pt[0])))
          return nw[t] = it.id, true;
      return false;
     while (gao(0) || gao(1));
     addEdge(nw[0], nw[1]); // add tangent
     while (true) {
       pll pt[2] = {p[nw[0]], p[nw[1]]};
       int ch = -1, sd = 0;
      for (int t = 0; t < 2; ++t)
        for (auto it : head[nw[t]])
           if (ori(pt[0], pt[1], p[it.id]) > 0 &&
             (ch == -1 |
               in_cc({pt[0], pt[1], p[ch]}, p[it.id])))
             ch = it.id, sd = t;
      if (ch == -1) break; // upper common tangent
       for (auto it = head[nw[sd]].begin();
            it != head[nw[sd]].end();)
         if (seg_strict_intersect(
               pt[sd], p[it->id], pt[sd ^ 1], p[ch]))
           head[it->id].erase(it->twin),
             head[nw[sd]].erase(it++);
         else ++it;
      nw[sd] = ch, addEdge(nw[0], nw[1]);
  }
|} tool;
```

# 8Y VonoroiDiagram

```
// all coord. is even
      you may want to call halfPlaneInter after then
vector<vector<Line>> vec;
void build_voronoi_line(int n, pll *arr) {
  tool.init(n, arr); // Delaunay
  vec.clear(), vec.resize(n);
  for (int i = 0; i < n; ++i)</pre>
    for (auto e : tool.head[i]) {
      int u = tool.oidx[i], v = tool.oidx[e.id];
       pll m = (arr[v
           ] + arr[u]) / 2LL, d = perp(arr[v] - arr[u]);
       vec[u].emplace_back(Line(m, m + d));
}
```

### Misc

### **9A** MoAlgoWithModify

```
|// Mo's Algorithm With modification
// Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
  static const int blk = 2000;
  int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
    L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
  bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     if (RBid != q.RBid) return RBid < q.RBid;</pre>
     return T < q.T;</pre>
  }
};
void solve(vector<Query> query) {
  sort(all(query));
  int L=0, R=0, T=-1;
  for (auto q : query) { // TODO: fill in
     // while (T < q.T) addTime(L, R, ++T);
    // while (T > q.T) subTime(L, R, T--);
     // while (R < q.R) add(arr[++R]);
    // while (L > q.L) add(arr[--L]);
```

```
// while (R > q.R) sub(arr[R--]);
    // while (L < q.L) sub(arr[L++]);
    // answer query
|}
     MoAlgoOnTree
```

```
Mo's Algorithm On Tree
 Preprocess:
 1) I CA
 2) dfs with in[u] = dft++, out[u] = dft++
 3) ord[in[u]] = ord[out[u]] = u
 4) bitset<MAXN> inset
 */
 struct Query {
   int L, R, LBid, lca;
   Query(int u, int v) {
     int c = LCA(u, v);
     if (c == u || c == v)
       q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
     else if (out[u] < in[v])</pre>
       q.lca = c, q.L = out[v], q.R = in[v];
     else
       q.lca = c, q.L = out[v], q.R = in[u];
     q.Lid = q.L / blk;
  }
  bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     return R < q.R;</pre>
  }
 };
 void flip(int x) {
     if (inset[x]) sub(arr[x]); // TODO
     else add(arr[x]); // TODO
     inset[x] = ~inset[x];
 void solve(vector<Query> query) {
   sort(ALL(query));
   int L = 0, R = 0;
   for (auto q : query) {
     while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
     while (R > q.R) flip(ord[R--]);
     while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
  }
| }
```

# 9C MoAlgoAdvanced

- Mo's Algorithm With Addition Only
- Sort querys same as the normal Mo's algorithm.

  - For each query [l,r]:
     If l/blk = r/blk, brute-force.
  - If  $l/blk \neq curL/blk$ , initialize  $curL := (l/blk+1) \cdot blk$ ,  $curR := \{\}$
  - If r > curR, increase curR
- decrease  $\mathit{curL}$  to fit l, and then undo after answering
- Mo's Algorithm With Offline Second Time
  - Require: Changing answer  $\equiv$  adding f([l,r],r+1).
  - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1). – Part1: Answer all f([1,r],r+1) first.
  - Part2: Store  $curR \to R$  for curL (reduce the space to and then answer them by the second offline O(N), algorithm.
  - Note: You must do the above symmetrically for the left boundaries.

#### 9D HilbertCurve

```
ll hilbert(int n, int x, int y) {
 11 res = 0;
 for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y \& s) > 0;
    res += s * 111 * s * ((3 * rx) ^ ry);
```

```
if (ry == 0) {
       if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
  return res;
| \} // n = 2^k
```

#### 9E ManhattanMST

```
|#define p3i tuple<int, int, int>
struct DSU {
   vector<int> v;
   DSU(int n);
   int query(int u);
   void merge(int x, int y);
};
vector<p3i> manhattanMST(vector<pll> ps) {
   vector<int> id(ps.size());
   iota(id.begin(), id.end(), 0);
   vector<p3i> edges;
   for (int k = 0; k < 4; ++k) {
     sort(id.begin(), id.end(), [&](int i, int j) {
       return (ps[i] - ps[j]).F < (ps[j] - ps[i]).S;</pre>
     map<int, int> sweep;
     for (int i : id) {
       for (auto it = sweep.lower_bound(-ps[i].S);
            it != sweep.end(); sweep.erase(it++)) {
         int j = it->second;
         pll d = ps[i] - ps[j];
         if (d.S > d.F) break;
         edges.emplace_back(d.S + d.F, i, j);
       sweep[-ps[i].S] = i;
     for (auto &p : ps)
       if (k & 1) p.F = -p.F;
       else swap(p.F, p.S);
  return edges;
}
vector<int> MST(int n, const vector<p3i> &e) {
   vector<int> idx(e.size());
  iota(idx.begin(), idx.end(), 0);
sort(idx.begin(), idx.end(), [&](int i, int j) {
     return get<0>(e[i]) < get<0>(e[j]);
  });
   vector<int> r;
   DSU dsu(n);
   for (int o : idx) {
     const auto &[w, i, j] = e[o];
     if (dsu.query(i) == dsu.query(j)) continue;
     r.push_back(o);
     dsu.merge(i, j);
   return r;
```

# SternBrocotTree

- Construction: Root  $\frac{1}{1}$ , left/right neighbor  $\frac{0}{1}, \frac{1}{0}$ , each node is sum of last left/right neighbor:  $rac{a}{b},rac{c}{d}
  ightarrowrac{a+c}{b+d}$
- Property: Adjacent (mid-order DFS)  $\frac{a}{b}, \frac{c}{d} \Rightarrow bc ad = 1$ .
- Search known  $rac{p}{q}$ : keep L-R alternative. Each step can calcaulated  $\inf' O(1) \Rightarrow \text{total } O(\log C)$ .
- Search unknown  $\frac{p}{q}$ : keep L-R alternative. Each step can calcaulated in  $O(\log C)$  checks  $\Rightarrow$  total  $O(\log^2 C)$  checks.

#### 9G AllLCS

```
void all_lcs(string s, string t) { // O-base
   vector<int> h((int)t.size());
   iota(all(h), 0);
   for (int a = 0; a < (int)s.size(); ++a) {</pre>
     int v = -1;
     for (int c = 0; c < (int)t.size(); ++c)</pre>
       if (s[a] == t[c] || h[c] < v)
```

```
| swap(h[c], v);
| // LCS(s[0, a], t[b, c]) =
| // c - b + 1 - sum([h[i] >= b] | i <= c)
| // h[i] might become -1 !!
| }
|}
```

#### 9H MatroidIntersection

```
Start from S=\emptyset. In each iteration, let Y_1=\{x\not\in S\,|\,S\cup\{x\}\in I_1\} • Y_2=\{x\not\in S\,|\,S\cup\{x\}\in I_2\} If there exists x\in Y_1\cap Y_2, insert x into S. Otherwise for each x\in S, y\not\in S, create edges • x\to y if S-\{x\}\cup\{y\}\in I_1. • y\to x if S-\{x\}\cup\{y\}\in I_2. Find a shortest path (with RES) starting from a vertex in
```

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \not\in S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

# 9I SimulatedAnnealing

```
| double factor = 100000;
| const int base = 1e9; // remember to run ~ 10 times
| for (int it = 1; it <= 1000000; ++it) {
| // ans: answer, nw: current value
| if (exp(-(nw -
| ans) / factor) >= (double)(rd() % base) / base)
| ans = nw;
| factor *= 0.99995;
| }
```

#### 9J SMAWK

```
int opt[N]:
Il A(int x, int y); // target func
void smawk(vector<int> &r, vector<int> &c);
void interpolate(vector<int> &r, vector<int> &c) {
  int n = (int)r.size();
  vector<int> er;
 for (int i = 1; i < n; i += 2) er.emplace_back(r[i]);</pre>
  smawk(er, c);
 for (int i = 0, j = 0; j < c.size(); j++) {</pre>
    if (A(r[i], c[j]) < A(r[i], opt[r[i]]))</pre>
      opt[r[i]] = c[j];
    if (i + 2 < n \& c[j] == opt[r[i + 1]])
      j--, i += 2;
 }
void reduce(vector<int> &r, vector<int> &c) {
 int n = (int)r.size();
  vector<int> nc;
 for (int i : c) {
    int j = (int)nc.size();
      j \& A(r[j-1], nc[j-1]) > A(r[j-1], i))
      nc.pop_back(), j--;
    if (nc.size() < n) nc.emplace_back(i);</pre>
 }
  smawk(r, nc);
void smawk(vector<int> &r, vector<int> &c) {
 if (r.size() == 1 && c.size() == 1) opt[r[0]] = c[0];
  else if (r.size() >= c.size()) interpolate(r, c);
  else reduce(r, c);
```

# 9K Python

```
import math
math.isqrt(2) # integer sqrt
```

### 9L LineContainer

```
struct Line {
   mutable ll k, m, p;
   bool operator<(const Line &o) const {</pre>
     return k < o.k;
  bool operator<(ll x) const { return p < x; }</pre>
};
struct LineContainer : multiset<Line, less<>> {
   // (for doubles, use inf = 1/.0, div(a,b) = a/b)
   static const ll inf = LLONG_MAX;
   ll div(ll a, ll b) { // floored division
     return a / b - ((a ^ b) < 0 && a % b);
   bool isect(iterator x, iterator y) {
     if (y == end()) return x->p = inf, 0;
     if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
     else x->p = div(y->m - x->m, x->k - y->k);
     return x->p >= y->p;
   void add(ll k, ll m) {
     auto z = insert({k, m, 0}), y = z++, x = y;
     while (isect(y, z)) z = erase(z);
     if (x != begin() && isect(--x, y))
       isect(x, y = erase(y));
     while ((y = x) != begin() \&\& (--x)->p >= y->p)
       isect(x, erase(y));
  11 query(11 x) {
     assert(!empty());
     auto l = *lower_bound(x);
     return l.k * x + l.m;
};
```