6.100uadraticResidue

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```

1 Basic

1.1 .vimrc

1.2 PBDS

```
// Tree and fast PQ
#include <bits/extc++.h>
using namespace __gnu_pbds;

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
tree<int, null_type, less<int>, rb_tree_tag
    , tree_order_statistics_node_update> bst;

// order_of_key(n): # of elements <= n
// find_by_order(n): 0-indexed

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
__gnu_pbds::priority_queue
    <int, greater<int>, thin_heap_tag> pq;
```

1.3 pragma

```
2.1 2SAT/SCC
struct SAT { // O-base
   int low[N], dfn[N], bln[N], n, Time, nScc;
   bool instack[N], istrue[N];
   stack<int> st;
   vector<int> G[N], SCC[N];
   void init(int _n) {
     n = _n; // assert(n * 2 <= N);
     for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
   void add_edge(int a, int b) { G[a].emplace_back(b); }
   int rv(int a) {
     if (a >= n) return a - n;
     return a + n;
   void add_clause(int a, int b) {
     add_edge(rv(a), b), add_edge(rv(b), a);
   void dfs(int υ) {
     dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
       do {
         tmp = st.top(), st.pop();
         instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != u);
       ++nScc;
     }
  }
   bool solve() {
     Time = nScc = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       if (!dfn[i]) dfs(i);
     for (int i =
         0; i < n + n; ++i) SCC[bln[i]].emplace_back(i);
     for (int i = 0; i < n; ++i) {</pre>
       if (bln[i] == bln[i + n]) return false;
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true;
  }
};
```

2.2 BCC Vertex

```
int n, m, dfn[N], low[N], is_cut[N], nbcc = 0, t = 0;
vector<int> g[N], bcc[N], G[2 * N];
stack<int> st;
void tarjan(int p, int lp) {
   dfn[p] = low[p] = ++t;
   st.push(p);
   for (auto i : g[p]) {
      if (!dfn[i]) {
        tarjan(i, p);
   }
}
```

```
low[p] = min(low[p], low[i]);
      if (dfn[p] <= low[i]) {
        nbcc++;
        is_cut[p] = 1;
        for (int x = 0; x != i; st.pop()) {
           x = st.top();
          bcc[nbcc].push_back(x);
        bcc[nbcc].push_back(p);
    } else low[p] = min(low[p], dfn[i]);
  }
void build() { // [n+1,n+nbcc] cycle, [1,n] vertex
  for (int i = 1; i <= nbcc; i++) {</pre>
    for (auto j : bcc[i]) {
      G[i + n].push_back(j);
      G[j].push_back(i + n);
  }
į }
```

2.3 MinimumMeanCycle

```
|/* 0(V^3)
|let dp[i][j] = min length from 1 to j exactly i edges
|ans = min (dp[n + 1][v] - dp[i][v]) / (n + 1 - i) */
```

2.4 MaximumCliqueDyn

```
struct MaxClique { // fast when N <= 100</pre>
  bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
   n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
 }
  void add_edge(int u, int v) {
    G[v][v] = G[v][v] = 1;
  void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(all(r)
          , [&](int x, int y) { return d[x] > d[y]; });
    vector<int> c(r.size());
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rqt) cs[++rqt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first
          (); p < N; p = cs[k]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
 }
  void dfs(vector<</pre>
      int> &r, vector<int> &c, int l, bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.emplace_back(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
 }
```

```
int solve() {
    vector<int> r(n);
     ans = q = 0, iota(all(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
     return ans;
|};
2.5 DMST(slow)
struct DMST { // O(VE)
  struct edge {
     int u, v;
     ll w;
  };
  vector<edge> E; // 0-base
  int pe[N], id[N], vis[N];
  ll in[N];
  void init() { E.clear(); }
  void add_edge(int u, int v, ll w) {
    if (u != v) E.emplace_back(edge{u, v, w});
  ll build(int root, int n) {
     ll ans = 0;
     for (;;) {
      fill_n(in, n, INF);
       for (int i = 0; i < E.size(); ++i)</pre>
         if (E[i].u != E[i].v && E[i].w < in[E[i].v])</pre>
           pe[E[i].v] = i, in[E[i].v] = E[i].w;
       for (int u = 0; u < n; ++u) // no solution</pre>
        if (u != root && in[u] == INF) return -INF;
       int cntnode = 0;
      fill_n(id, n, -1), fill_n(vis, n, -1);
      for (int u = 0; u < n; ++u) {
         if (u != root) ans += in[u];
        int v = u;
        while (vis[v] != u && !~id[v] && v != root)
           vis[v] = u, v = E[pe[v]].u;
         if (v != root && !~id[v]) {
           for (int x = E[pe[v]].u; x != v;
                x = E[pe[x]].u)
             id[x] = cntnode;
           id[v] = cntnode++;
        }
      ļ
      if (!cntnode) break; // no cycle
       for (int u = 0; u < n; ++u)
        if (!~id[u]) id[u] = cntnode++;
```

2.6 DMST

return ans;

}

}

|};

```
#define rep(i, a, b) for (int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef vector<int> vi;
struct RollbackUF {
  vi e;
  vector<pii> st;
  RollbackUF(int n) : e(n, -1) {}
  int size(int x) { return -e[find(x)]; }
  int find(int x) { return e[x] < 0 ? x : find(e[x]); }</pre>
  int time() { return sz(st); }
  void rollback(int t) {
    for (int i = time(); i-- > t;)
      e[st[i].first] = st[i].second;
    st.resize(t);
  }
```

for (int i = 0; i < E.size(); ++i) {</pre>

n = cntnode, root = id[root];

E[i].v = id[E[i].v], E[i].v = id[E[i].v];

if (E[i].u != E[i].v) E[i].w -= in[v];

int v = E[i].v;

```
bool join(int a, int b) {
                                                                in[uf.find(inEdge.b)] = inEdge;
    a = find(a), b = find(b);
    if (a == b) return false;
                                                              rep(i, 0, n) par[i] = in[i].a;
    if (e[a] > e[b]) swap(a, b);
                                                              return {res, par};
    st.push_back({a, e[a]});
                                                           |}
    st.push_back({b, e[b]});
                                                            2.7 VizingTheorem
   e[a] += e[b];
    e[b] = a;
                                                           | namespace Vizing { // Edge coloring
    return true;
                                                                                // G: coloring adjM
 }
                                                            int C[N][N], G[N][N];
};
                                                            void clear(int n) {
struct Edge {
                                                              for (int i = 0; i <= n; i++) {</pre>
 int a, b;
                                                                for (int j = 0; j <= n; j++) C[i][j] = G[i][j] = 0;</pre>
  ll w;
                                                            }
struct Node { /// lazy skew heap node
                                                            void solve(vector<pii> &E, int n, int m) {
  Edge key;
                                                              int X[n] = {}, a;
  Node *l, *r;
  ll delta;
                                                              auto update = [&](int u) {
  void prop() {
                                                                for (X[u] = 1; C[u][X[u]]; X[u]++);
    key.w += delta;
    if (l) l->delta += delta;
                                                              auto color = [&](int u, int v, int c) {
    if (r) r->delta += delta;
                                                                int p = G[u][v];
    delta = 0;
                                                                G[\upsilon][v] = G[v][\upsilon] = c;
 }
                                                                C[u][c] = v;
 Edge top() {
                                                                C[v][c] = u;
    prop():
                                                                C[\upsilon][p] = C[\upsilon][p] = 0;
    return key;
                                                                if (p) X[u] = X[v] = p;
                                                                else update(u), update(v);
};
                                                                return p;
Node *merge(Node *a, Node *b) {
 if (!a || !b) return a ?: b;
                                                              auto flip = [&](int u, int c1, int c2) {
  a->prop(), b->prop();
                                                                int p = C[u][c1];
 if (a->key.w > b->key.w) swap(a, b);
                                                                swap(C[u][c1], C[u][c2]);
  swap(a->l, (a->r = merge(b, a->r)));
                                                                if (p) G[u][p] = G[p][u] = c2;
  return a;
                                                                if (!C[u][c1]) X[u] = c1;
                                                                if (!C[u][c2]) X[u] = c2;
void pop(Node *&a) {
                                                                return p;
 a->prop();
 a = merge(a->l, a->r);
                                                              for (int i = 1; i <= n; i++) X[i] = 1;</pre>
                                                              for (int t = 0; t < E.size(); t++) {</pre>
                                                                int u = E[t].first, v0 = E[t].second, v = v0,
pair<ll, vi> dmst(int n, int r, vector<Edge> &g) {
                                                                    c0 = X[v], c = c0, d;
 RollbackUF uf(n);
                                                                vector<pii> L;
  vector<Node *> heap(n);
                                                                int vst[n] = {};
 for (Edge e : g)
                                                                while (!G[u][v0]) {
   heap[e.b] = merge(heap[e.b], new Node{e});
                                                                  L.emplace_back(v, d = X[v]);
                                                                  if (!C[v][c])
 vi seen(n, -1), path(n), par(n);
                                                                    for (a = (int)L.size() - 1; a >= 0; a--)
  seen[r] = r:
                                                                      c = color(u, L[a].first, c);
  vector<Edge> Q(n), in(n, {-1, -1}), comp;
                                                                  else if (!C[u][d])
  deque<tuple<int, int, vector<Edge>>> cycs;
                                                                    for (a = (int)L.size() - 1; a >= 0; a--)
  rep(s, 0, n) {
                                                                      color(u, L[a].first, L[a].second);
    int u = s, qi = 0, w;
                                                                  else if (vst[d]) break;
    while (seen[u] < 0) {
                                                                  else vst[d] = 1, v = C[u][d];
      if (!heap[u]) return {-1, {}};
      Edge e = heap[u]->top();
                                                                if (!G[u][v0]) {
      heap[u]->delta -= e.w, pop(heap[u]);
                                                                  for (; v; v = flip(v, c, d), swap(c, d));
      Q[qi] = e, path[qi++] = v, seen[v] = s;
                                                                  if (C[u][c0]) {
      res += e.w, u = uf.find(e.a);
                                                                    for (a = (int)L.size() - 2;
      if (seen[u] == s) { /// found cycle, contract
                                                                         a >= 0 && L[a].second != c; a--)
        Node *cyc = 0;
        int end = qi, time = uf.time();
                                                                    for (; a >= 0; a--)
        do cyc = merge(cyc, heap[w = path[--qi]]);
                                                                      color(u, L[a].first, L[a].second);
        while (uf.join(u, w));
                                                                  } else t--;
        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
                                                                }
        cycs.push_front({u, time, {&Q[qi], &Q[end]}});
                                                              }
                                                           |} // namespace Vizing
    rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
                                                            2.8 MinimumCliqueCover
 for (auto &[u, t, comp] :
                                                            struct CliqueCover { // O-base, O(n2^n)
    cycs) { // restore sol (optional)
                                                              int co[1 << N], n, E[N];</pre>
    uf.rollback(t);
                                                              int dp[1 << N];</pre>
    Edge inEdge = in[v];
                                                              void init(int _n) {
    for (auto &e : comp) in[uf.find(e.b)] = e;
                                                                n = _n, fill_n(dp, 1 << n, 0);
```

```
fill_n(E, n, 0), fill_n(co, 1 << n, 0);
  void add_edge(int u, int v) {
    E[u] \mid = 1 << v, E[v] \mid = 1 << u;
  int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;
     dp[0] = (n \& 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {
       int t = i & -i;
       dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)
       co[i] = (co[i] \& i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {</pre>
       int sum = 0; // probabilistic
       for (int i = 0; i < (1 << n); ++i)</pre>
         sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n;
  }
|};
```

2.9 CountMaximalClique

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
    g[v][v] = g[v][v] = 1;
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
       if (g[u][v]) continue;
      int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
      all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)</pre>
         if (g[v][some[d][j]])
           some[d + 1][tsn++] = some[d][j];
      for (int j = 0; j < nn; ++j)</pre>
         if (g[v][none[d][j]])
           none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
       some[d][i] = 0, none[d][nn++] = v;
    }
  }
  int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
|};
```

2.10 Theorems

 $|\max \text{ independent edge set}| = |V| - |\min \text{ edge cover}| \\ |\max \text{ independent set}| = |V| - |\min \text{ vertex cover}|$

3 Flow-Matching

3.1 HopcroftKarp

```
struct HopcroftKarp
      { // O-based, return btoa to get matching
  bool dfs(int a, int L, vector<vector<int>> &g,
     vector<int> &btoa, vector<int> &A,
     vector<int> &B) {
     if (A[a] != L) return 0;
     A[a] = -1;
     for (int b : g[a])
       if (B[b] == L + 1) {
         B[b] = 0;
         if (btoa[b] == -1 ||
           dfs(btoa[b], L + 1, g, btoa, A, B))
           return btoa[b] = a, 1;
    return 0;
  }
  int solve(vector<vector<int>> &q, int m) {
    int res = 0:
     vector<int> btoa(m, -1), A(g.size()),
       B(btoa.size()), cur, next;
     for (;;) {
       fill(all(A), 0), fill(all(B), 0);
       cur.clear();
       for (int a : btoa)
         if (a != -1) A[a] = -1;
       for (int a = 0; a < g.size(); a++)</pre>
         if (A[a] == 0) cur.push_back(a);
       for (int lay = 1;; lay++) {
         bool islast = 0;
         next.clear();
         for (int a : cur)
           for (int b : g[a]) {
             if (btoa[b] == -1) {
               B[b] = lay;
               islast = 1:
             } else if (btoa[b] != a && !B[b]) {
               B[b] = lay;
               next.push_back(btoa[b]);
          }
         if (islast) break;
         if (next.empty()) return res;
         for (int a : next) A[a] = lay;
         cur.swap(next);
       for (int a = 0; a < g.size(); a++)</pre>
         res += dfs(a, 0, g, btoa, A, B);
    }
  }
|};
3.2 KM
struct KM { // O-base
  ll w[N][N], hl[N], hr[N], slk[N];
  int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
  bool vl[N], vr[N];
  void init(int _n) {
    n = n;
     for (int i = 0; i < n; ++i)</pre>
       fill_n(w[i], n, -INF);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool Check(int x) {
     if (vl[x] = 1, \sim fl[x])
       return vr[qu[qr++] = fl[x]] = 1;
     while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    fill_n(slk
          n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
     ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (ll d;;) {
```

```
while (ql < qr)</pre>
         for (int x = 0, y = qu[ql++]; x < n; ++x)
           if (!vl[x] && slk
               [x] >= (d = hl[x] + hr[y] - w[x][y])) {
             if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
         if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
         if (vl[x]) hl[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !Check(x)) return;
    }
  }
  ll solve() {
    fill_n(fl
          n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
     for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(w[i], w[i] + n);
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 res = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
|};
```

3.3 MCMF

```
struct MinCostMaxFlow { // O-base
  struct Edge {
   ll from, to, cap, flow, cost, rev;
 } *past[N];
 vector<Edge> G[N];
  int inq[N], n, s, t;
 ll dis[N], up[N], pot[N];
 bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
     if (cap > 0 && dis[u] > d) {
       dis[u] = d, up[u] = cap, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
   };
    relax(s, 0, INF, 0);
   while (!q.empty()) {
     int u = q.front();
     q.pop(), inq[v] = 0;
     for (auto &e : G[u]) {
       11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
        relax(
          e.to, d2, min(up[u], e.cap - e.flow), &e);
     }
   }
   return dis[t] != INF;
 }
 bool Dijkstra() {
    fill_n(dis, n, INF);
    priority_queue<pll, vector<pll>, greater<pll>>> pq;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
     if (cap > 0 && dis[u] > d) {
        dis[v] = d, up[v] = cap, past[v] = e;
        pq.push(pll(d, u));
     }
   };
    relax(s, 0, INF, 0);
    while (!pq.empty()) {
     auto [d, u] = pq.top();
     pq.pop();
     if (dis[v] != d) continue;
     for (auto &e : G[u]) {
```

```
ll d2 = dis[v] + e.cost + pot[v] - pot[e.to];
        relax(
           e.to, d2, min(up[u], e.cap - e.flow), &e);
    }
    return dis[t] != INF;
  void solve(int _s, int _t, ll &flow, ll &cost,
    bool neg = true) {
     s = _s, t = _t, flow = 0, cost = 0;
     if (neg) BellmanFord(), copy_n(dis, n, pot);
     // do BellmanFord() if time isn't tight
     for (; Dijkstra(); copy_n(dis, n, pot)) {
      for (int i = 0; i < n; ++i)</pre>
        dis[i] += pot[i] - pot[s];
       flow += up[t], cost += up[t] * dis[t];
       for (int i = t; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         e.flow += up[t], G[e.to][e.rev].flow -= up[t];
    }
  }
  void init(int _n) {
    n = _n, fill_n(pot, n, 0);
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].emplace_back(
      Edge{a, b, cap, 0, cost, (int)G[b].size()});
     G[b].emplace_back(
       Edge{b, a, 0, 0, -cost, (int)G[a].size() - 1});
};
```

3.4 GeneralGraphMatching

}

```
struct Matching { // O-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
    }
  }
  void Blossom(int x, int y, int l) {
    for (; Find(x) != l; x = pre[y]) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
    }
  }
  bool Bfs(int r) {
    iota(all(fa), 0); fill(all(s), -1);
    q = queue<int>(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : G[x])
        if (s[u] == -1) {
          if (pre[u] = x, s[u] = 1, match[u] == n) {
            for (int a = u, b = x, last;
                b != n; a = last, b = pre[a])
                  match[b], match[b] = a, match[a] = b;
            return true;
          }
          q.push(match[u]); s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int l = LCA(u, x);
          Blossom(x, u, l); Blossom(u, x, l);
```

```
3.5 MaxWeightMaching
#define rep(i, l, r) for (int i = (l); i <= (r); ++i)
struct WeightGraph { // 1-based
  struct edge {
   int u, v, w;
  int n, nx;
  vector<int> lab;
  vector<vector<edge>> q;
 vector<int> slack, match, st, pa, S, vis;
 vector<vector<int>> flo, flo_from;
  queue<int> q;
 WeightGraph(int n_)
    : n(n_{-}), nx(n * 2), lab(nx + 1),
      g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1),
      flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
    match = st = pa = S = vis = slack;
   rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
 int ED(edge e) {
   return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
 void update_slack(int u, int x, int &s) {
    if (!s || ED(g[u][x]) < ED(g[s][x])) s = u;</pre>
 void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)
     if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
        update_slack(u, x, slack[x]);
 }
 void q_push(int x) {
    if (x \le n) q.push(x);
   else
      for (int y : flo[x]) q_push(y);
 void set_st(int x, int b) {
    st[x] = b;
    if(x > n)
     for (int y : flo[x]) set_st(y, b);
 vector<int> split_flo(auto &f, int xr) {
    auto it = find(ALL(f), xr);
    if (auto pr = it - f.begin(); pr % 2 == 1)
     reverse(1 + ALL(f)), it = f.end() - pr;
    auto res = vector(f.begin(), it);
    return f.erase(f.begin(), it), res;
 void set_match(int u, int v) {
   match[u] = g[u][v].v;
    if (u <= n) return;</pre>
    int xr = flo_from[u][g[u][v].u];
   auto &f = flo[u], z = split_flo(f, xr);
    rep(i, 0, (int)z.size() - 1)
     set_match(z[i], z[i ^ 1]);
    set_match(xr, v);
    f.insert(f.end(), all(z));
 }
 void augment(int u, int v) {
    for (;;) {
     int xnv = st[match[u]];
     set_match(u, v);
```

```
if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
int lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v))
    if (u) {
      if (vis[u] == t) return u;
      vis[u] = t;
      u = st[match[u]];
      if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + all(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0;
  match[b] = match[o];
  vector<int> f = {o};
  for (int x = u, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  reverse(1 + all(f));
  for (int x = v, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  flo[b] = f;
  set_st(b, b);
  for (int x = 1; x <= nx; ++x)
    g[b][x].w = g[x][b].w = 0;
  fill(all(flo_from[b]), 0);
  for (int xs : flo[b]) {
    for (int x = 1; x <= nx; ++x)
      if (g[b][x].w == 0 | |
        ED(g[xs][x]) < ED(g[b][x])
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) {
      xs = x;
      continue;
    pa[xs] = g[x][xs].u;
    S[xs] = 1, S[x] = 0;
    slack[xs] = 0;
    set_slack(x);
    q_push(x);
    xs = -1;
  for (int x : flo[b])
    if (x == xr) S[x] = 1, pa[x] = pa[b];
    else S[x] = -1, set_slack(x);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
    int nu = st[match[v]];
    pa[v] = e.u;
    S[v] = 1;
    slack[v] = slack[nu] = 0;
    S[nu] = 0;
    q_push(nu);
  } else if (S[v] == 0) {
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
```

void init(int n) {

for (int i = 0; i < n; ++i) fill_n(edge[i], n, 0);</pre>

```
return false;
                                                               void addEdge(int u, int v, int w) {
  }
                                                                 edge[u][v] += w;
  bool matching() {
                                                                 edge[v][u] += w;
    fill(all(S), -1), fill(all(slack), 0);
    q = queue<int>();
                                                               int search(int &s, int &t, int n) {
    for (int x = 1; x <= nx; ++x)
                                                                 fill_n(vst, n, 0), fill_n(wei, n, 0);
      if (st[x] == x \&\& !match[x])
                                                                 s = t = -1
                                                                 int mx, cur;
        pa[x] = 0, S[x] = 0, q_push(x);
                                                                 for (int j = 0; j < n; ++j) {</pre>
    if (q.empty()) return false;
                                                                   mx = -1, cur = 0;
    for (;;) {
                                                                   for (int i = 0; i < n; ++i)</pre>
      while (q.size()) {
                                                                     if (wei[i] > mx) cur = i, mx = wei[i];
        int u = q.front();
                                                                   vst[cur] = 1, wei[cur] = -1;
        q.pop();
        if (S[st[v]] == 1) continue;
                                                                   t = cur;
        for (int v = 1; v <= n; ++v)</pre>
                                                                   for (int i = 0; i < n; ++i)</pre>
           if (g[u][v].w > 0 && st[u] != st[v]) {
                                                                     if (!vst[i]) wei[i] += edge[cur][i];
             if (ED(g[u][v]) != 0)
               update_slack(u, st[v], slack[st[v]]);
                                                                 return mx;
             else if (on_found_edge(g[u][v]))
                                                               }
               return true;
                                                               int solve(int n) {
                                                                 int res = INF;
                                                                 for (int x, y; n > 1; n--) {
      int d = INF;
                                                                   res = min(res, search(x, y, n));
      for (int b = n + 1; b <= nx; ++b)</pre>
                                                                   for (int i = 0; i < n; ++i)</pre>
        if (st[b] == b && S[b] == 1)
                                                                     edge[i][x] = (edge[x][i] += edge[y][i]);
           d = min(d, lab[b] / 2);
                                                                   for (int i = 0; i < n; ++i) {</pre>
      for (int x = 1; x <= nx; ++x)</pre>
                                                                     edge[y][i] = edge[n - 1][i];
        if (int s = slack[x];
                                                                     edge[i][y] = edge[i][n - 1];
             st[x] == x && s && s[x] <= 0)
                                                                   } // edge[y][y] = 0;
           d = min(d, ED(g[s][x]) / (S[x] + 2));
      for (int u = 1; u <= n; ++u)</pre>
                                                                 return res;
        if (S[st[u]] == 1) lab[u] += d;
                                                               }
        else if (S[st[u]] == 0) {
                                                            |} sw;
          if (lab[v] <= d) return false;</pre>
           lab[v] -= d;
                                                             3.7 BoundedFlow(Dinic)
        }
                                                            | struct BoundedFlow { // O-base
      rep(b, n + 1, nx) if (st[b] == b \&\& S[b] >= 0)
        lab[b] += d * (2 - 4 * S[b]);
                                                               struct edge {
                                                                 int to, cap, flow, rev;
      for (int x = 1; x <= nx; ++x)
        if (int s = slack[x]; st[x] == x \&\& s \&\&
                                                               vector<edge> G[N];
             st[s] != x \&\& ED(g[s][x]) == 0)
                                                               int n, s, t, dis[N], cur[N], cnt[N];
           if (on_found_edge(g[s][x])) return true;
                                                               void init(int _n) {
      for (int b = n + 1; b <= nx; ++b)
                                                                 n = _n;
        if (st[b] == b && S[b] == 1 && lab[b] == 0)
                                                                 for (int i = 0; i < n + 2; ++i)</pre>
           expand_blossom(b);
                                                                   G[i].clear(), cnt[i] = 0;
    return false;
                                                               void add_edge(int u, int v, int lcap, int rcap) {
  }
                                                                 cnt[u] -= lcap, cnt[v] += lcap;
  pair<ll, int> solve() {
                                                                 G[u].emplace_back
    fill(all(match), 0);
                                                                     (edge{v, rcap, lcap, (int)G[v].size()});
    rep(u, 0, n) st[u] = u, flo[u].clear();
    int w_max = 0;
                                                                 G[v].emplace_back
                                                                     (edge{u, 0, 0, (int)G[u].size() - 1});
    rep(u, 1, n) rep(v, 1, n) {
      flo_from[u][v] = (u == v ? u : 0);
                                                               void add_edge(int u, int v, int cap) {
      w_{max} = max(w_{max}, q[u][v].w);
                                                                 G[u].emplace_back
    }
    fill(all(lab), w_max);
                                                                     (edge{v, cap, 0, (int)G[v].size()});
    int n_matches = 0;
                                                                 G[v].emplace_back
    tot_weight = 0;
                                                                     (edge{u, 0, 0, (int)G[u].size() - 1});
    while (matching()) ++n_matches;
    rep(u, 1, n) if (match[u] \&\& match[u] < u)
                                                               int dfs(int u, int cap) {
      tot_weight += g[u][match[u]].w;
                                                                 if (u == t || !cap) return cap;
    return make_pair(tot_weight, n_matches);
                                                                 for (int &i = cur[v]; i < G[v].size(); ++i) {</pre>
                                                                   edge &e = G[v][i];
  void add_edge(int u, int v, int w) {
                                                                   if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
    g[v][v].w = g[v][v].w = w;
                                                                     int df = dfs(e.to, min(e.cap - e.flow, cap));
  }
                                                                     if (df) {
∣}:
                                                                       e.flow += df, G[e.to][e.rev].flow -= df;
                                                                       return df;
3.6
      GlobalMinCut
                                                                     }
                                                                   }
struct StoerWagner { // O(V^3), is it O(VE + V log V)?
                                                                 }
  int vst[N], edge[N][N], wei[N];
                                                                 dis[v] = -1;
```

return 0;

```
bool bfs() {
     fill_n(dis, n + 3, -1);
     queue<int> q;
     q.push(s), dis[s] = 0;
     while (!q.empty()) {
       int u = q.front();
       q.pop();
       for (edge &e : G[u])
         if (!~dis[e.to] && e.flow != e.cap)
           q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
int flow = 0, df;
     while (bfs()) {
       fill_n(cur, n + 3, 0);
       while ((df = dfs(s, INF))) flow += df;
    return flow;
  }
  bool solve() {
     int sum = 0;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
         add_edge(n + 1, i, cnt[i]), sum += cnt[i];
       else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);</pre>
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
         G[n + 1].pop_back(), G[i].pop_back();
       else if (cnt[i] < 0)
         G[i].pop_back(), G[n + 2].pop_back();
     return sum != -1;
  }
  int solve(int _s, int _t) {
     add_edge(_t, _s, INF);
     if (!solve()) return -1; // invalid flow
     int x = G[_t].back().flow;
     return G[_t].pop_back(), G[_s].pop_back(), x;
  }
|};
```

3.8 GomoryHuTree

```
| MaxFlow Dinic;
| int g[N];
| void GomoryHu(int n) { // O-base |
| fill_n(g, n, 0);
| for (int i = 1; i < n; ++i) {
| Dinic.reset();
| add_edge(i, g[i], Dinic.maxflow(i, g[i]));
| for (int j = i + 1; j <= n; ++j)
| if (g[j] == g[i] && ~Dinic.dis[j])
| g[j] = i;
| }
| }</pre>
```

3.9 MinCostCirculation

```
while (!q.empty()) {
       int u = q.front();
       q.pop(), inq[v] = 0;
       for (auto &e : G[u])
         if (e.cap > e.flow)
           relax(e.to, dis[u] + e.cost, &e);
    }
  void try_edge(Edge &cur) {
     if (cur.cap > cur.flow) return ++cur.cap, void();
     BellmanFord(cur.to);
     if (dis[cur.from] + cur.cost < 0) {</pre>
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
            i = cur.from; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         ++e.flow, --G[e.to][e.rev].flow;
       }
     ++cur.cap;
  }
  void solve(int mxlq) {
     for (int b = mxlg; b >= 0; --b) {
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           e.cap *= 2, e.flow *= 2;
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
    }
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].emplace_back(Edge{a, b,
          0, cap, 0, cost, (ll)G[b].size() + (a == b));
     G[b].emplace_back(Edge
         {b, a, 0, 0, 0, -cost, (ll)G[a].size() - 1});
|} mcmf; // O(VE * ElogC)
```

3.10 FlowModelsBuilding

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \rightarrow y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - of outgoing lower bounds. 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching ${\cal M}$ on bipartite graph (X,Y)
 - 1. Redirect every edge: $y \rightarrow x$ if $(x,y) \in M$, $x \rightarrow y$ otherwise.
 - 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited. • Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x,y,c), connect $x \to y$ with (cost,cap) = (c,1) if c>0, otherwise connect $y \to x$ with (cost,cap) = (-c,1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0 , connect $S\to v$ with (cost, cap)=(0, d(v))

```
5. For each vertex v with d(v) < 0, connect v \to T with
     (cost, cap) = (0, -d(v))
     Flow from S to T, the answer is the cost of the flow
     C+K

    Maximum density induced subgraph

  1. Binary search on answer, suppose we're checking
     answer T
  2. Construct a max flow model, let K be the sum of all
     weights
  3. Connect source s \rightarrow v, v \in G with capacity K
  4. For each edge (u,v,w) in G, connect u \rightarrow v and v \rightarrow u with
     capacity w
  5. For v \in G, connect it with sink v \to t with capacity
     K+2T-(\sum_{e\in E(v)}w(e))-2w(v)
  6. T is a valid answer if the maximum flow f < K|V|
 Minimum weight edge cover
  1. For each v \in V create a copy v', and connect u' \to v'
     with weight w(u,v).
  2. Connect v \rightarrow v' with weight 2\mu(v), where \mu(v) is the cost
     of the cheapest edge incident to v.
  3. Find the minimum weight perfect matching on G'.
 Project selection problem
  1. If p_v > 0, create edge (s,v) with capacity p_v; otherwise,
     create edge (v,t) with capacity -p_v .
  2. Create edge (u,v) with capacity w with w being the
     cost of choosing u without choosing v.
  The mincut is equivalent to the maximum profit of a
     subset of projects.
Dual of minimum cost maximum flow
 1. Capacity c_{uv}, Flow f_{uv}, Cost w_{uv}, Required Flow difference for vertex b_u.
  2. If all w_{uv} are integers, then optimal solution can
     happen when all p_u are integers.
          \min \sum w_{uv} f_{uv}
                         \min \! \sum \! b_u p_u \! + \! \sum \! c_{uv} \! \max (0, \! p_v \! - \! p_u \! - \! w_{uv})
                                                          p_u > 0
```

4 Data Struture

4.1 LinkCutTree

```
#define ls(x) Tree[x].son[0]
#define rs(x) Tree[x].son[1]
#define fa(x) Tree[x].fa
struct node {
 int son[2], Min, id, fa, lazy;
} Tree[N];
int n, m, q, w[N], Min;
struct Node {
 int u, v, w;
} a[N];
inline bool IsRoot(int x) {
  return (ls(fa(x)) == x \mid\mid rs(fa(x)) == x) ? false
inline void PushUp(int x) {
 Tree[x].Min = w[x], Tree[x].id = x;
  if (ls(x) && Tree[ls(x)].Min < Tree[x].Min) {</pre>
    Tree[x].Min = Tree[ls(x)].Min;
    Tree[x].id = Tree[ls(x)].id;
 }
 if (rs(x) \&\& Tree[rs(x)].Min < Tree[x].Min) {
    Tree[x].Min = Tree[rs(x)].Min;
    Tree[x].id = Tree[rs(x)].id;
inline void Update(int x) {
 Tree[x].lazy ^= 1;
  swap(ls(x), rs(x));
inline void PushDown(int x) {
 if (!Tree[x].lazy) return;
  if (ls(x)) Update(ls(x));
  if (rs(x)) Update(rs(x));
 Tree[x].lazy = 0;
inline void Rotate(int x) {
 int y = fa(x), z = fa(y), k = rs(y) == x,
```

```
w = Tree[x].son[!k];
  if (!IsRoot(y)) Tree[z].son[rs(z) == y] = x;
  fa(x) = z, fa(y) = x;
  if (w) fa(w) = y;
  Tree[x].son[!k] = y, Tree[y].son[k] = w;
  PushUp(y);
inline void Splay(int x) {
  stack<int> Stack;
  int y = x, z;
  Stack.push(y);
  while (!IsRoot(y)) Stack.push(y = fa(y));
  while (!Stack.empty())
    PushDown(Stack.top()), Stack.pop();
  while (!IsRoot(x)) {
    y = fa(x), z = fa(y);
    if (!IsRoot(y))
      Rotate((ls(y) == x) ^(ls(z) == y) ? x : y);
    Rotate(x);
  PushUp(x);
inline void Access(int root) {
  for (int x = 0; root; x = root, root = fa(root))
    Splay(root), rs(root) = x, PushUp(root);
inline void MakeRoot(int x) {
 Access(x), Splay(x), Update(x);
inline int FindRoot(int x) {
  Access(x), Splay(x);
  while (ls(x)) x = ls(x);
  return Splay(x), x;
inline void Link(int u, int v) {
  MakeRoot(u);
  if (FindRoot(v) != u) fa(u) = v;
inline void Cut(int u, int v) {
  MakeRoot(u):
  if (FindRoot(v) != u || fa(v) != u || ls(v)) return;
  fa(v) = rs(u) = 0;
}
inline void Split(int u, int v) {
  MakeRoot(u), Access(v), Splay(v);
inline bool Check(int u, int v) {
  return MakeRoot(u), FindRoot(v) == u;
}
inline int LCA(int root, int u, int v) {
  MakeRoot(root), Access(u), Access(v), Splay(u);
  if (!fa(u)) {
    Access(u), Splay(v);
    return fa(v);
  return fa(u);
}
/* ETT
每次進入節點和走邊都放入一次共 3n - 2
node(u) 表示進入節點 u 放入 treap 的位置
edge(u, v) 表示 u -> v 的邊放入 treap 的位置 (push v)
Makeroot u
 L1 = [begin, node(u) - 1], L2 = [node(u), end]
  -> L2 + L1
Insert u. v :
  Tu \rightarrow L1 = [begin, node(u) - 1], L2 = [node(u), end]
  Tv -> L3 = [begin, node(v) - 1], L4 = [node(v), end]
  -> L2 + L1 + edge(u, v) + L4 + L3 + edge(v, u)
Delect u, v :
  maybe need swap u, v
  T -> L1 + edge(u, v) + L2 + edge(v, u) + L3
  -> L1 + L3, L2
```

5 String 5.1 KMP

```
int KMP(string s, string t) {
  t = " "s + t; // consistency with ACa
  int n = t.size(), ans = 0;
   vector<int> f(t.size(), 0);
   f[0] = -1;
  for (int i = 1, j = -1; i < t.size(); i++) {</pre>
     while (j \ge 0 \&\& t[j + 1] != t[i])
       j = f[j];
     f[i] = ++j;
  }
   for (int i = 0, j = 0; i < s.size(); i++) {</pre>
     while (j \ge 0 \&\& t[j + 1] != s[i])
       j = f[j];
     if (++j + 1 == t.size()) ans++, j = f[j];
   return ans;
| }
 5.2 Z
int Z[N];
 void z(string s) {
```

```
for (int i = 1, mx = 0; i < s.size(); i++) {</pre>
    if (i < Z[mx] + mx)
      Z[i] = min(Z[mx] - i + mx, Z[i - mx]);
    while (
      Z[i] + i < s.size() \&\& s[i + Z[i]] == s[Z[i]])
      Z[i]++;
    if (Z[i] + i > Z[mx] + mx) mx = i;
  }
}
```

5.3 Manacher

```
int man[N]; // len: man[i] - 1
void manacher(string s) { // uses 2|s|+1
  string t;
  for (int i = 0; i < (int)s.size(); i++) {</pre>
    t.push_back('$');
    t.push_back(s[i]);
  }
  t.push_back('$');
  int mx = 1;
  for (int i = 0; i < (int)t.size(); i++) {</pre>
    man[i] = 1;
    man[i] = min(man[mx] + mx - i, man[2 * mx - i]);
    while (man[i] + i < (int)t.size() && i >= man[i] &&
      t[i + man[i]] == t[i - man[i]])
      man[i]++;
    if (i + man[i] > mx + man[mx]) mx = i;
|}
```

5.4 SuffixArray

```
struct SuffixArray {
#define add(x, k) (x + k + n) % n
 vector<int> sa, cnt, rk, tmp, lcp;
 // sa: order, rk[i]: pos of s[i..],
 // lcp[i]: LCP of sa[i], sa[i-1]
 void SA(string s) { // remember to append '\1'
    int n = s.size();
    sa.resize(n), cnt.resize(n);
    rk.resize(n), tmp.resize(n);
    iota(all(sa), 0);
    sort(all(sa),
      [&](int i, int j) { return s[i] < s[j]; });</pre>
    rk[0] = 0;
    for (int i = 1; i < n; i++)</pre>
      rk[sa[i]] =
        rk[sa[i - 1]] + (s[sa[i - 1]] != s[sa[i]]);
    for (int k = 1; k <= n; k <<= 1) {</pre>
      fill(all(cnt), 0);
```

```
for (int i = 0; i < n; i++)</pre>
         cnt[rk[add(sa[i], -k)]]++;
       for (int i = 1; i < n; i++) cnt[i] += cnt[i - 1];</pre>
       for (int i = n - 1; i >= 0; i--)
         tmp[--cnt[rk[add(sa[i], -k)]]] =
           add(sa[i], -k);
       sa.swap(tmp);
       tmp[sa[0]] = 0;
       for (int i = 1; i < n; i++)</pre>
         tmp[sa[i]] = tmp[sa[i - 1]] +
           (rk[sa[i - 1]] != rk[sa[i]] ||
             rk[add(sa[i - 1], k)] !=
               rk[add(sa[i], k)]);
       rk.swap(tmp);
  }
  void LCP(string s) {
     int n = s.size(), k = 0;
     lcp.resize(n);
     for (int i = 0; i < n; i++)</pre>
       if (rk[i] == 0) lcp[rk[i]] = 0;
       else {
         if (k) k--;
         int j = sa[rk[i] - 1];
         while (
           \max(i, j) + k < n \&\& s[i + k] == s[j + k])
         lcp[rk[i]] = k;
  }
};
```

5.5 SAIS

```
auto sais(const auto &s) {
  const int n = SZ(s), z = ranges::max(s) + 1;
  if (n == 1) return vector{0};
  vector<int> c(z); for (int x : s) ++c[x];
  partial_sum(ALL(c), begin(c));
  vector<int> sa(n); auto I = views::iota(0, n);
  vector<bool> t(n, true);
  for (int i = n - 2; i >= 0; --i)
    t[i] = (
         s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  auto is_lms = views::filter([&t](int x) {
    return x && t[x] && !t[x - 1];
  auto induce = [&] {
    for (auto x = c; int y : sa)
      if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
    for (auto x = c; int y : sa | views::reverse)
      if (y--) if (t[y]) sa[--x[s[y]]] = y;
  vector<int> lms, q(n); lms.reserve(n);
  for (auto x = c; int i : I | is_lms)
    q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
  induce(); vector<int> ns(SZ(lms));
  for (int j = -1, nz = 0; int i : sa | is_lms) {
    if (j >= 0) {
      int len = min({n - i, n - j, lms[q[i] + 1] - i});
       ns[q[i]] = nz += lexicographical_compare(
           begin(s) + j, begin(s) + j + len,
           begin(s) + i, begin(s) + i + len);
    j = i;
  }
  fill(ALL(sa), 0); auto nsa = sais(ns);
  for (auto x = c; int y : nsa | views::reverse)
    y = lms[y], sa[--x[s[y]]] = y;
  return induce(), sa;
}
// sa[i]: sa[i]-th suffix
      is the i-th lexicographically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
  int n; vector<int> sa, hi, ra;
```

5.6 ACAutomaton

```
#define sigma 26
#define base 'a'
struct AhoCorasick { // N: sum of length
  int ch[N][sigma] = {{}}, f[N] = {-1}, tag[N],
  mv[N][sigma], jump[N], cnt[N];
int idx = 0, t = -1;
  vector<int> E[N], q;
  pii o[N];
  int insert(string &s, int t) {
    int j = 0;
    for (int i = 0; i < (int)s.size(); i++) {</pre>
      if (!ch[j][s[i] - base])
        ch[j][s[i] - base] = ++idx;
      j = ch[j][s[i] - base];
    tag[j] = 1;
    return j;
  int next(int u, int c) {
    return u < 0 ? 0 : mv[u][c];</pre>
  }
  void dfs(int u) {
    o[u].F = ++t;
    for (auto v : E[u]) dfs(v);
    o[u].S = t;
  void build() {
    int k = -1;
    q.emplace_back(0);
    while (++k < q.size()) {</pre>
      int u = q[k];
      for (int v = 0; v < sigma; v++) {</pre>
        if (ch[u][v]) {
          f[ch[u][v]] = next(f[u], v);
           q.emplace_back(ch[u][v]);
        }
        mv[u][v] =
           (ch[u][v] ? ch[u][v] : next(f[u], v));
      }
      if (u) jump[u] = (tag[f[u]] ? f[u] : jump[f[u]]);
    reverse(q.begin(), q.end());
    for (int i = 1; i <= idx; i++)</pre>
      E[f[i]].emplace_back(i);
    dfs(0);
  }
  void match(string &s) {
    fill(cnt, cnt + idx + 1, 0);
    for (int i = 0, j = 0; i < (int)s.size(); i++)</pre>
      cnt[j = next(j, s[i] - base)]++;
    for (int i : q)
      if (f[i] > 0) cnt[f[i]] += cnt[i];
  }
} ac;
```

5.7 MinRotation

```
int mincyc(string s) {
  int n = s.size();
  s = s + s;
```

```
int i = 0, ans = 0;
while (i < n) {
    ans = i;
    int j = i + 1, k = i;
    while (j < s.size() && s[j] >= s[k]) {
        k = (s[j] > s[k] ? i : k + 1);
        ++j;
    }
    while (i <= k) i += j - k;
}
return ans;
</pre>
```

```
5.8 ExtSAM
#define CNUM 26
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
    fill_n(next[tot], CNUM, 0);
    len[tot] = cnt[tot] = link[tot] = 0;
    return tot++;
  void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len
        [p] + 1 == len[q]) return link[cur] = q, cur;
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[
          clone][i] = len[next[q][i]] ? next[q][i] : 0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
  void insert(const string &s) {
    int cur = 0:
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
    }
  }
  void build() {
    queue<int> q;
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    vector<int> lc(tot);
    for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
    partial_sum(all(lc), lc.begin());
    for (int i
        = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
  void solve() {
    for (int i = tot - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
```

|};

5.9 PalindromeTree

```
struct PalindromicTree {
  struct node {
    int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
                  // pal. suf.
    node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
    }
  };
  vector<node> St;
  vector<char> s;
  int last, n;
  PalindromicTree() : St(2), last(1), n(0) {
    St[0].fail = 1, St[1].len = -1, s.emplace_back(-1);
  }
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.emplace_back(0), St.emplace_back(-1);
    St[0].fail = 1, s.emplace_back(-1);
  }
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
      x = St[x].fail;
    return x;
  }
  inline void add(int c) {
    s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = St.size();
      St.emplace_back(St[cur].len + 2);
      St[now].fail =
        St[get_fail(St[cur].fail)].next[c];
      St[cur].next[c] = now;
      St[now].num = St[St[now].fail].num + 1;
    }
    last = St[cur].next[c], ++St[last].cnt;
  }
  inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
      St[i->fail].cnt += i->cnt;
  }
  inline int size() { // The number of diff. pal.
    return (int)St.size() - 2;
};
```

6 Number Theory

6.1 Primes

12721 13331 14341 75577 123457 222557 556679 999983 1097774749 1076767633 100102021 999997771 1001010013 1000512343 987654361 999991231 999888733 98789101 987777733 999991921 1010101333 1010102101 10000000000039 100000000000037 2305843009213693951 4611686018427387847 9223372036854775783 18446744073709551557

6.2 ExtGCD

```
|// beware of negative numbers!
|void extgcd(ll a, ll b, ll c, ll &x, ll &y) {
| if (b == 0) x = c / a, y = 0;
| else {
| extgcd(b, a % b, c, y, x);
| y -= x * (a / b);
| }
|} // |x| <= b/2, |y| <= a/2</pre>
```

6.3 FloorCeil

```
| int floor(int a, int b)
|{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
| int ceil(int a, int b)
|{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

6.4 FloorSum

Computes

$$f(a,b,c,n) = \sum_{i=0}^{n} \left\lfloor \frac{a \cdot i + b}{m} \right\rfloor$$

Furthermore, Let $m = \left| \frac{an+b}{c} \right|$:

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \left\lfloor \frac{ai+b}{c} \right\rfloor \\ &= \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \mod c, b \mod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c - b - 1, a, m - 1) \\ -h(c, c - b - 1, a, m - 1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \left\lfloor \frac{ai+b}{c} \right\rfloor^2 \\ &= \begin{cases} \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \\ &+ \left\lfloor \frac{a}{c} \right\rfloor \cdot \left\lfloor \frac{b}{c} \right\rfloor \cdot n(n+1) \\ &+ h(a \mod c, b \mod c, c, n) \\ &+ 2 \left\lfloor \frac{a}{c} \right\rfloor \cdot g(a \mod c, b \mod c, c, n) \\ &+ 2 \left\lfloor \frac{b}{c} \right\rfloor \cdot f(a \mod c, b \mod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ &- 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

```
| Il floorsum(ll A, ll B, ll C, ll N) {
    if (A == 0) return (N + 1) * (B / C);
    if (A > C || B > C)
        return (N + 1) * (B / C) +
        N * (N + 1) / 2 * (A / C) +
        floorsum(A % C, B % C, C, N);
    ll M = (A * N + B) / C;
    return N * M - floorsum(C, C - B - 1, A, M - 1);
}
```

6.5 MillerRabin

```
// n < 4,759,123,141 3 : 2, 7, 61
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 4,759,123,141
// n < 3,474,749,660,383 6 : primes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool Miller_Rabin(ll a, ll n) {
  if ((a = a % n) == 0) return 1;
  if (n % 2 == 0) return n == 2;
  11 tmp = (n - 1) / ((n - 1) & (1 - n));
  for (; tmp; tmp >>= 1, a = mul(a, a, n))
    if (tmp \& 1) x = mul(x, a, n);
  if (x == 1 || x == n - 1) return 1;
  while (--t)
    if ((x = mul(x, x, n)) == n - 1) return 1;
  return 0;
```

6.6 PollardRho

}

return -1;

if (a1) return inv(a1, m);

```
d = gcd(abs(x - y), n);
                                                           | }
  }
                                                           mt19937 rd(49);
|}
                                                           ll psqrt(ll y, ll p) {
                                                             if (fpow(y, (p - 1) / 2, p) != 1) return -1;
6.7 Fraction
                                                             for (int i = 0; i < 30; i++) {
                                                               ll z = rd() \% p;
struct fraction {
                                                               if (z * z % p == y) return z;
   ll n, d;
                                                               ll x = trial(y, z, p);
   fraction
                                                               if (x == -1) continue;
       (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
                                                               return x;
    ll t = gcd(n, d);
    n /= t, d /= t;
                                                             return -1;
    if (d < 0) n = -n, d = -d;
                                                          }
  }
  fraction operator-() const
                                                           6.11 MeisselLehmer
  { return fraction(-n, d); }
                                                          | ll PrimeCount(ll n) \{ // n \sim 10^13 => < 2s \}
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
                                                             if (n <= 1) return 0;</pre>
  fraction operator-(const fraction &b) const
                                                             int v = sqrt(n), s = (v + 1) / 2, pc = 0;
                                                             vector<int> smalls(v + 1), skip(v + 1), roughs(s);
  { return fraction(n * b.d - b.n * d, d * b.d); }
                                                             vector<ll> larges(s);
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
                                                             for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
  fraction operator/(const fraction &b) const
                                                             for (int i = 0; i < s; ++i) {
  { return fraction(n * b.d, d * b.n); }
                                                               roughs[i] = 2 * i + 1;
                                                               larges[i] = (n / (2 * i + 1) + 1) / 2;
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;
                                                             for (int p = 3; p <= v; ++p) {
  }
                                                               if (smalls[p] > smalls[p - 1]) {
                                                                 int q = p * p;
|};
                                                                 ++pc;
6.8
      ChineseRemainder
                                                                 if (1LL * q * q > n) break;
                                                                 skip[p] = 1;
for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
  ll g = gcd(m1, m2);
                                                                 int ns = 0;
  if ((x2 - x1) % g) return -1; // no sol
                                                                 for (int k = 0; k < s; ++k) {
  m1 /= g; m2 /= g;
                                                                   int i = roughs[k];
  ll x, y;
                                                                   if (skip[i]) continue;
  extgcd(m1, m2,
                   _gcd(m1, m2), x, y);
                                                                   11 d = 1LL * i * p;
  ll lcm = m1 * m2 * g;
                                                                   larges[ns] = larges[k] - (d <= v ? larges</pre>
  ll res = x * (x2 - x1) * m1 + x1;
                                                                        [smalls[d] - pc] : smalls[n / d]) + pc;
  // be careful with overflow
                                                                   roughs[ns++] = i;
  return (res % lcm + lcm) % lcm;
                                                                 }
|}
                                                                 s = ns;
                                                                 for (int j = v / p; j >= p; --j) {
6.9 Factorial\mathsf{Mod}p^k
                                                                         smalls[j] - pc, e = min(j * p + p, v + 1);
|// O(p^k + log^2 n), pk = p^k
                                                                   for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
11 prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
                                                               }
  prod[0] = 1;
                                                             }
  for (int i = 1; i <= pk; ++i)</pre>
                                                             for (int k = 1; k < s; ++k) {
    if (i % p) prod[i] = prod[i - 1] * i % pk;
                                                               const ll m = n / roughs[k];
    else prod[i] = prod[i - 1];
                                                               ll t = larges[k] - (pc + k - 1);
  11 rt = 1;
                                                               for (int l = 1; l < k; ++l) {</pre>
  for (; n; n /= p) {
                                                                 int p = roughs[l];
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
                                                                 if (1LL * p * p > m) break;
    rt = rt * prod[n % pk] % pk;
                                                                 t -= smalls[m / p] - (pc + l - 1);
  }
  return rt;
                                                               larges[0] -= t;
} // (n! without factor p) % p^k
6.10 QuadraticResidue
                                                             return larges[0];
                                                          }
// Berlekamp-Rabin, log^2(p)
                                                           6.12 DiscreteLog
ll trial(ll y, ll z, ll m) {
  11 a0 = 1, a1 = 0, b0 = z, b1 = 1, p = (m - 1) / 2;
                                                           int DiscreteLog(int s, int x, int y, int m) {
  while (p) {
                                                             constexpr int kStep = 32000;
    if (p & 1)
                                                             unordered_map<int, int> p;
      tie(a0.a1) =
                                                             int b = 1;
        make_pair((a1 * b1 % m * y + a0 * b0) % m,
                                                             for (int i = 0; i < kStep; ++i) {</pre>
           (a0 * b1 + a1 * b0) % m);
                                                               p[y] = i;
    tie(b0, b1) =
                                                                y = 1LL * y * x % m;
       make_pair((b1 * b1 % m * y + b0 * b0) % m,
                                                               b = 1LL * b * x % m;
                                                             }
        (2 * b0 * b1) % m);
    p >>= 1;
                                                             for (int i = 0; i < m + 10; i += kStep) {</pre>
```

s = 1LL * s * b % m;

if (p.find(s) != p.end()) return i + kStep - p[s];

```
return -1;
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {
    if (s == y) return i;
    s = 1LL * s * x % m;
  }
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
```

6.13 Theorems

· Cramer's Rule

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}\!=\!d(i)$, $L_{ij}\!=\!-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(ilde{L}_{11})|$.
- The number of directed spanning tree rooted at \boldsymbol{r} in G is $|\det(\tilde{L}_{rr})|$.

Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence $d_1,d_2,...,d_n$ for each labeled vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex 1,2,...,k belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

Erdős-Gallai Theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+\cdots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k) \text{ holds for every } 1 \leq k \leq n \text{.}$$

Gale-Ryser Theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1,\ldots,b_n is bigraphic (degree seqence of bipartie

graph) if and only if
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$

holds for every $1 \le k \le n$. Fulkerson-Chen-Anstee Theorem

A sequence $(a_1,b_1),...,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq \cdots \geq a_n$ is digraphic (in, out degree of a di-

rected graph) if and only if
$$\sum\limits_{i=1}^n a_i = \sum\limits_{i=1}^n b_i$$
 and $\sum\limits_{i=1}^k a_i \leq$

$$\sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n\text{.}$$

- Möbius Inversion Formula
 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
 - $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Lagrange Multiplier
 - Optimize $f(x_1,...,x_n)$ when k constraints $g_i(x_1,...,x_n)\!=\!0$.
 - Lagrangian function $\mathcal{L}(x_1,\ \dots,\ x_n,\ \lambda_1,\ \dots,\ \lambda_k)=f(x_1,\dots,x_n)-\sum_{i=1}^k\lambda_ig_i(x_1,\dots,x_n)$. The solution corresponding to the original con-
 - strained optimization is always a saddle point of the Lagrangian function.

6.14 Estimation

- Estimation

6.15 Numbers

• Bernoulli numbers
$$B_0-1, B_1^{\pm}=\pm \frac{1}{2}, B_2=\frac{1}{6}, B_3=0$$

$$\sum_{j=0}^m {m+1 \choose j} B_j=0, \text{ EGF is } B(x)=\frac{x}{e^x-1}=\sum_{n=0}^\infty B_n \frac{x^n}{n!} \cdot S_m(n)=\sum_{k=1}^n k^m=\frac{1}{m+1}\sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

Stirling numbers of the second kind Partitions of ndistinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
 • Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

6.16 GeneratingFunctions

- Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$
 - $A(rx) \Rightarrow r^n a_n$
 - $A(x) + B(x) \Rightarrow a_n + b_n$
 - $A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x)' \Rightarrow na_n$
 - $\frac{A(x)}{1-x}$ $\Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i \ge 0} \frac{a_i}{i!} x_i$
 - $A(x)+B(x) \Rightarrow a_n+b_n$

 - -A(x) = A(x) =
- Special Generating Function
- $(1+x)^n = \sum_{i>0} \binom{n}{i} x^i$
- $-\frac{1}{(1-x)^n} = \sum_{i>0} {i \choose n-1} x^i$
- $S_k = \sum_{x=1}^n x^k$: $S = \sum_{p=0}^\infty x^p = \frac{e^x e^{x(n+1)}}{1 e^x}$

Linear Algebra

7.1 GuassianElimination

```
struct matrix { // m variables, n equations
  int n, m;
  fraction A[N][N + 1], sol[N];
  int solve() { //-1: inconsistent, >= 0: rank
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
      while (piv < m && !A[i][piv].n) ++piv;</pre>
      if (piv == m) continue;
      for (int j = 0; j < n; ++j) {</pre>
        if (i == j) continue;
        fraction tmp = -A[j][piv] / A[i][piv];
        for (int k = 0; k <= m; ++k)</pre>
          A[j][k] = tmp * A[i][k] + A[j][k];
      }
    int rank = 0;
```

```
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     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !A[i][piv].n) ++piv;</pre>
       if (piv == m && A[i][m].n) return -1;
       else if (piv < m)</pre>
         ++rank, sol[piv] = A[i][m] / A[i][piv];
     return rank;
  }
|};
7.2
       BerlekampMassey
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(output.size() + 1), me, he;
  for (int f = 0, i = 1; i <= output.size(); ++i) {</pre>
     for (int j = 0; j < me.size(); ++j)</pre>
       d[i] += output[i - j - 2] * me[j];
     if ((d[i] -= output[i - 1]) == 0) continue;
     if (me.empty()) {
       me.resize(f = i);
       continue;
    vector<T> o(i - f - 1);
    T k = -d[i] / d[f];
    o.emplace_back(-k);
    for (T x : he) o.emplace_back(x * k);
```

7.3 Simplex

me = o;

return me;

}

|}

if (i - f + (**int**

Standard form: maximize $\mathbf{c}^T\mathbf{x}$ subject to $A\mathbf{x} \leq \mathbf{b}$ and $\mathbf{x} \geq 0$. Dual LP: minimize $\mathbf{b}^T\mathbf{y}$ subject to $A^T\mathbf{y} \geq \mathbf{c}$ and $\mathbf{y} \geq \mathbf{0}$. $\bar{\mathbf{x}}$ and $\bar{\mathbf{y}}$ are optimal if and only if for all $i \in [1,n]$, either $\bar{x}_i = 0$ or $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$ holds and for all $i \in [1,m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$ holds.

1. In case of minimization, let $c_i' = -c_i$

for (int j = 0; j < me.size(); ++j) o[j] += me[j];</pre>

)he.size()) >= (int)me.size()) he = me, f = i;

2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \rightarrow \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$ 3. $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ • $\sum_{i \le n} A_{ji} x_i \le b_i$

o.resize(max(o.size(), me.size()));

• $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$ • $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j$

4. If x_i has no lower bound, replace x_i with $x_i - x_i'$

```
// n variable, m constraints, M >= n + 2m
struct simplex {
 const double inf = 1 / .0, eps = 1e-9;
 int n, m, k, var[N], inv[N], art[N];
 double A[M][N], B[M], x[N];
 void init(int _n) { n = _n, m = 0; }
 void equation(vector<double> a, double b) {
    for (int i = 0; i < n; i++) A[m][i] = a[i];</pre>
    B[m] = b, var[m] = n + m, ++m;
 void pivot(int r, int c, double bx) {
    for (int i = 0; i <= m + 1; i++)
      if (i != r && abs(A[i][c]) > eps) {
        x[var[i]] -= bx * A[i][c] / A[i][var[i]];
        double f = A[i][c] / A[r][c];
        for (int j = 0; j <= n + m + k; j++)</pre>
          A[i][j] -= A[r][j] * f;
        B[i] -= B[r] * f;
      }
 }
  double phase(int p) {
    while (true) {
      int in = min_element(
                 A[m + p], A[m + p] + n + m + k + 1) -
        A[m + p];
      if (A[m + p][in] >= -eps) break;
      double bx = inf;
```

```
int piv = -1;
       for (int i = 0; i < m; i++)</pre>
         if (A[i][in] > eps && B[i] / A[i][in] <= bx)</pre>
           piv = i, bx = B[i] / A[i][in];
       if (piv == -1) return inf;
       int out = var[piv];
       pivot(piv, in, bx);
       x[out] = 0, x[in] = bx, var[piv] = in;
     return x[n + m];
   double solve(vector<double> c) {
     auto invert = [&](int r) {
       for (int j = 0; j <= n + m; j++) A[r][j] *= -1;</pre>
       B[r] *= -1;
     };
     k = 1:
     for (int i = 0; i < n; i++) A[m][i] = -c[i];</pre>
     fill(A[m + 1], A[m + 1] + N, 0.0);
     for (int i = 0; i <= m + 1; i++)</pre>
       fill(A[i] + n, A[i] + n + m + 2, 0.0),
         var[i] = n + i, A[i][n + i] = 1;
     for (int i = 0; i < m; i++) {
       if (B[i] < 0) {
         for (int j = 0; j <= n + m; j++)</pre>
           A[m + 1][j] += A[i][j];
         invert(i);
         var[i] = n + m + k, A[i][var[i]] = 1,
         art[var[i]] = n + i;
       x[var[i]] = B[i];
     phase(1):
     if (*max_element(
           x + (n + m + 2), x + (n + m + k + 1)) > eps)
       return .0 / .0;
     for (int i = 0; i <= m; i++)</pre>
       if (var[i] > n + m)
         var[i] = art[var[i]], invert(i);
     k = 0;
     return phase(0);
} lp;
```

8 Polynomials

8.1 NTT (FFT)

```
hoM
                                        Form
                         65 537
                                        2^{16}+1
                                        119 \cdot 2^{23} + 1
                    998 244 353
                                        1255 \!\cdot\! 2^{20} \!+\! 1
                 1 315 962 881
                  1 711 276 033
                                        51 \cdot 2^{25} + 1
                                   29
                                        549755813881 \!\cdot\! 2^{24} \!+\! 1
    9 223 372 036 737 335 297
|#define base ll // complex<double>
// const double PI = acosl(-1);
const ll mod = 998244353, g = 3;
base omega[4 * N], omega_[4 * N];
int rev[4 * N];
ll fpow(ll b, ll p);
ll inverse(ll a) { return fpow(a, mod - 2); }
void calcW(int n) {
  ll r = fpow(q, (mod - 1) / n), invr = inverse(r);
   omega[0] = omega_[0] = 1;
   for (int i = 1; i < n; i++) {</pre>
     omega[i] = omega[i - 1] * r % mod;
     omega_[i] = omega_[i - 1] * invr % mod;
  // double arg = 2.0 * PI / n;
  // for (int i = 0; i < n; i++)
   // {
```

```
omega[i] = base(cos(i * arg), sin(i * arg));
                                                                for (int i = 0; i < n; ++i)</pre>
        omega_[i] = base(cos(-i * arg), sin(-i * arg));
                                                                  c[i] = h[ct[i]][i];
  // }
                                                             }
                                                              8.3 PolynomialOperations
void calcrev(int n) {
                                                             |#define poly vector<ll>
  int k = __lg(n);
                                                              poly inv(poly A) {
  for (int i = 0; i < n; i++) rev[i] = 0;</pre>
                                                                A.resize(1 << (__lg(A.size() - 1) + 1));
  for (int i = 0; i < n; i++)</pre>
                                                                poly B = {inverse(A[0])};
     for (int j = 0; j < k; j++)</pre>
                                                                for (int n = 1; n < (int)A.size(); n <<= 1) {</pre>
       if (i & (1 << j)) rev[i] ^= 1 << (k - j - 1);</pre>
                                                                  poly pA(A.begin(), A.begin() + 2 * n);
                                                                  calcrev(4 * n), calcW(4 * n);
                                                                  pA.resize(4 * n), B.resize(4 * n);
vector<base> NTT(vector<base> poly, bool inv) {
                                                                  pA = NTT(pA, 0);
  base *w = (inv ? omega_ : omega);
                                                                  B = NTT(B, 0);
  int n = poly.size();
                                                                  for (int i = 0; i < 4 * n; i++)</pre>
  for (int i = 0; i < n; i++)</pre>
                                                                    B[i] =
     if (rev[i] > i) swap(poly[i], poly[rev[i]]);
                                                                      ((B[i] * 2 - pA[i] * B[i] % mod * B[i]) % mod +
                                                                        mod) %
  for (int len = 1; len < n; len <<= 1) {</pre>
                                                                      mod:
    int arg = n / len / 2;
                                                                  B = NTT(B, 1);
     for (int i = 0; i < n; i += 2 * len)</pre>
                                                                  B.resize(2 * n);
       for (int j = 0; j < len; j++) {</pre>
         base odd =
                                                                return B;
           w[j * arg] * poly[i + j + len] % mod;
                                                              }
         poly[i + j + len] =
                                                              pair<poly, poly> div(poly A, poly B) {
           (poly[i + j] - odd + mod) % mod;
                                                                if (A.size() < B.size()) return make_pair(poly(), A);</pre>
         poly[i + j] = (poly[i + j] + odd) % mod;
                                                                int n = A.size(), m = B.size();
                                                                poly revA = A, invrevB = B;
                                                                reverse(all(revA)), reverse(all(invrevB));
  if (inv)
                                                                revA.resize(n - m + 1);
     for (auto &a : poly) a = a * inverse(n) % mod;
                                                                invrevB.resize(n - m + 1);
   return poly;
                                                                invrevB = inv(invrevB);
                                                                poly Q = mul(revA, invrevB);
                                                                Q.resize(n - m + 1);
vector<base> mul(vector<base> f, vector<base> g) {
                                                                reverse(all(Q));
  int sz = 1 << (__lg(f.size() + g.size() - 1) + 1);</pre>
                                                                poly R = mul(Q, B);
  f.resize(sz), g.resize(sz);
                                                                R.resize(m - 1);
  calcrev(sz);
                                                                for (int i = 0; i < m - 1; i++)</pre>
  calcW(sz);
                                                                  R[i] = (A[i] - R[i] + mod) \% mod;
  f = NTT(f, 0), g = NTT(g, 0);
                                                                return make_pair(Q, R);
  for (int i = 0; i < sz; i++)</pre>
                                                              }
    f[i] = f[i] * g[i] % mod;
                                                              poly modulo(poly A, poly B) { return div(A, B).S; }
  return NTT(f, 1);
                                                              ll fast_kitamasa(ll k, poly A, poly C) {
į }
                                                                int n = A.size();
8.2 FHWT
                                                                C.emplace_back(mod - 1);
                                                                poly Q, R = \{0, 1\}, F = \{1\};
|/* x: a[j], y: a[j + (L >> 1)]
                                                                R = modulo(R, C);
or: (y += x * op), and: (x += y * op)
                                                                for (; k; k >>= 1) {
xor: (x, y = (x + y) * op, (x - y) * op)
                                                                  if (k & 1) F = modulo(mul(F, R), C);
invop: or, and, xor = -1, -1, 1/2 */
                                                                  R = modulo(mul(R, R), C);
void fwt(int *a, int n, int op) { //or
                                                                  k >>= 1;
  for (int L = 2; L <= n; L <<= 1)
     for (int i = 0; i < n; i += L)
                                                                ll ans = 0;
       for (int j = i; j < i + (L >> 1); ++j)
                                                                for (int i = 0; i < F.size(); i++)</pre>
         a[j + (L >> 1)] += a[j] * op;
                                                                  ans = (ans + A[i] * F[i]) % mod;
                                                                return ans;
const int N = 21;
                                                              }
int f[
                                                              vector<ll> fpow(vector<ll> f, ll p, ll m) {
     N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
                                                                int b = 0:
                                                                while (b < f.size() && f[b] == 0) b++;</pre>
     subset_convolution(int *a, int *b, int *c, int L) {
                                                                f = vector<ll>(f.begin() + b, f.end());
   // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
                                                                int n = f.size();
  int n = 1 << L;
                                                                f.emplace_back(0);
  for (int i = 1; i < n; ++i)</pre>
                                                                vector<ll> q(min(m, b * p), 0);
    ct[i] = ct[i \& (i - 1)] + 1;
                                                                q.emplace_back(fpow(f[0], p));
  for (int i = 0; i < n; ++i)</pre>
                                                                for (int k = 0; q.size() < m; k++) {</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
                                                                  11 res = 0;
  for (int i = 0; i <= L; ++i)</pre>
                                                                  for (int i = 0; i < min(n, k + 1); i++)</pre>
    fwt(f[i], n, 1), fwt(g[i], n, 1);
                                                                    res = (res +
  for (int i = 0; i <= L; ++i)</pre>
                                                                             p * (i + 1) % mod * f[i + 1] % mod *
     for (int j = 0; j <= i; ++j)</pre>
                                                                               q[k - i + b * p]) %
       for (int x = 0; x < n; ++x)
                                                                      mod;
         h[i][x] += f[j][x] * g[i - j][x];
                                                                  for (int i = 1; i < min(n, k + 1); i++)</pre>
  for (int i = 0; i <= L; ++i)</pre>
                                                                    res = (res -
    fwt(h[i], n, -1);
                                                                             f[i] * (k - i + 1) % mod *
```

```
q[k - i + 1 + b * p]) %
         mod:
     res = (res < 0 ? res + mod : res) *
       inv(f[0] * (k + 1) % mod) % mod;
     q.emplace_back(res);
  return q;
|}
```

8.4 NewtonMethod+MiscGF

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that $| \ \}$ // from line p0--p1 to q0--q1, apply to rF(P)=0 can be found iteratively. Denote by \mathcal{Q}_k the polynomial such that $F(Q_k) = 0 \pmod{x^{2^k}}$, then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$
 • $A^{-1}\!:\;B_{k+1}\!=\!B_k(2\!-\!AB_k)\mod x^{2^{k+1}}$

- $\ln A$: $(\ln A)' = \frac{A'}{A}$
- $\exp A$: $B_{k+1} = B_k(1 + A \ln B_k) \mod x^{2^{k+1}}$
- \sqrt{A} : $B_{k+1} = \frac{1}{2}(B_k + AB_k^{-1}) \mod x^{2^{k+1}}$

Geometry

9.1 Basic

```
typedef pair<pdd, pdd> Line;
struct Cir{ pdd 0; double R; };
const double eps = 1e-8;
pll operator+(pll a, pll b)
{ return pll(a.F + b.F, a.S + b.S); }
pll operator-(pll a, pll b)
{ return pll(a.F - b.F, a.S - b.S); }
pll operator-(pll a)
{ return pll(-a.F, -a.S); }
pll operator*(pll a, ll b)
{ return pll(a.F * b, a.S * b); }
pdd operator/(pll a, double b)
{ return pdd(a.F / b, a.S / b); }
ll dot(pll a, pll b)
{ return a.F * b.F + a.S * b.S; }
ll cross(pll a, pll b)
{ return a.F * b.S - a.S * b.F; }
ll abs2(pll a)
{ return dot(a, a); }
double abs(pll a)
{ return sqrt(dot(a, a)); }
int sign(ll a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pll a, pll b, pll c)
{ return sign(cross(b - a, c - a)); }
bool collinearity(pll p1, pll p2, pll p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pll a, pll b, pll c) {
  return collinearity
       (a, b, c) \&\& sign(dot(a - c, b - c)) <= 0;
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
     return btw(p1, p2, p3) || btw(p1, p2, p4) ||
  btw(p3, p4, p1) || btw(p3, p4, p2);
return a123 * a124 <= 0 && a341 * a342 <= 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
  double a124 = cross(p2 - p1, p4 - p1);
       * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
| }
```

```
pdd perp(pdd p1)
{ return pdd(-p1.S, p1.F); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
    ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
    (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) {
  pdd dp = p1 - p0
      , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  return q0 + pdd(
      cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
```

9.2 ConvexHull

```
void hull(vector<pll> &dots) { // n=1 => ans = {}
  sort(dots.begin(), dots.end());
  vector<pll> ans(1, dots[0]);
  for (int ct = 0; ct < 2; ++ct, reverse(all(dots)))</pre>
    for (int i = 1, t = ans.size()
         ; i < dots.size(); ans.emplace_back(dots[i++]))</pre>
       while (ans.size() > t &&
           ori(ans.end()[-2], ans.back(), dots[i]) <= 0)
         ans.pop_back();
  ans.pop_back(), ans.swap(dots);
}
```

9.3 SortByAngle

```
int cmp(pll a, pll b, bool same = true) {
#define is_neg(k) (
     sign(k.S) < 0 \mid | (sign(k.S) == 0 \&\& sign(k.F) < 0))
  int A = is_neg(a), B = is_neg(b);
  if (A != B)
     return A < B;
  if (sign(cross(a, b)) == 0)
    return same ? abs2(a) < abs2(b) : -1;
  return sign(cross(a, b)) > 0;
}
```

9.4 Formulas

Rotation

$$M(\theta) = \begin{bmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{bmatrix}$$

90 degree: (x,y) = (Y-y,x)

· Pick's theorem

For simple integer-coordinate polygon,

$$A = \mathsf{B} + \frac{I}{2} - 1$$

Where A is the area; B,I is #lattice points in the interior, on the boundary.

· Spherical Cap

- A portion of a sphere cut off by a plane. - r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : arcsin(a/r).
- Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos\theta) (1 \cos\theta)^2/3$. Volume
- Area $= 2\pi r \dot{h} = \pi (a^2 + \dot{h}^2) = 2\pi r^2 (1 \cos\theta)$.

Nearest points of two skew lines

```
- Line 1:{m v}_1\!=\!{m p}_1\!+\!t_1{m d}_1
- Line 2:{m v}_2\!=\!{m p}_2\!+\!t_2{m d}_2
 - \boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2
 - \boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}
- \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}
- c_1 = p_1 + \frac{(p_2 - p_1) \cdot n_2}{d_1 \cdot n_2} d_1
- c_2 = p_2 + \frac{(p_1 - p_2) \cdot n_1}{d_2 \cdot n_1} d_2
```

9.5 TriangleHearts

```
pdd excenter(
  pdd p0, pdd p1, pdd p2) { // radius = abs(center)
  p1 = p1 - p0, p2 = p2 - p0;
  auto [x1, y1] = p1;
  auto [x2, y2] = p2;
  double m = 2. * cross(p1, p2);
  pdd center = pdd((x1 * x1 * y2 - x2 * x2 * y1 +
                     y1 * y2 * (y1 - y2)),
    (x1 * x2 * (x2 - x1) - y1 * y1 * x2 +
      x1 * y2 * y2)) / m;
  return center + p0;
pdd incenter(
  pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
  double a = abs(p2 - p3), b = abs(p1 - p3),
        c = abs(p1 - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3) {
  return (p1 + p2 + p3) / 3;
pdd orthcenter(pdd p1, pdd p2, pdd p3) {
  return masscenter(p1, p2, p3) * 3 -
    excenter(p1, p2, p3) * 2;
```

9.6 PointSegmentDist

```
| double PointSegDist(pdd q0, pdd q1, pdd p) {
| if (abs(q0 - q1) <= eps) return abs(q0 - p);
| if (dot(q1 - q0,
| p - q0) >= -eps && dot(q0 - q1, p - q1) >= -eps)
| return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
| return min(abs(p - q0), abs(p - q1));
|}
```

9.7 PointInCircle

```
// return q'
s relation with circumcircle of tri(p[0],p[1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
    __int128 det = 0;
for (int i = 0; i < 3; ++i)
    det += __int128(abs2(p[i]) - abs2(q)) *
        cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

9.8 PointInConvex

9.9 PointTangentConvex

```
|/* The point should be strictly out of hull
| return arbitrary point on the tangent line */
|/* bool pred(int a, int b);
|f(0) ~ f(n - 1) is a cyclic-shift U-function
|return idx s.t. pred(x, idx) is false forall x*/
|int cyc_tsearch(int n, auto pred) {
```

```
if (n == 1) return 0;
int l = 0, r = n; bool rv = pred(1, 0);
while (r - l > 1) {
   int m = (l + r) / 2;
   if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
   else l = m;
}
return pred(l, r % n) ? l : r % n;
}
pii get_tangent(vector<pll> &C, pll p) {
   auto gao = [&](int s) {
     return cyc_tsearch((int)C.size(), [&](int x, int y)
     { return ori(p, C[x], C[y]) == s; });
};
return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

9.10 CircTangentCirc

```
vector<Line</pre>
    > go( const Cir& c1 , const Cir& c2 , int sign1 ){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = abs2(c1.0 - c2.0);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
    pdd n = pdd(v.F * c - sign2 * h * v.S,
      v.S * c + sign2 * h * v.F);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sign(p1.F - p2.F) == 0 and
        sign(p1.S - p2.S) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
    ret.emplace_back(Line(p1, p2));
  return ret;
```

9.11 LineCircleIntersect

9.12 LineConvexIntersect

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch((int)C.size(), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
  });
}
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = (int)C.size();
  if (cmpL(A) < 0 \mid | cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r; ) {
      int m = ((l + r + (l < r? 0 : n)) / 2) % n;
       (cmpL(m) == cmpL(t) ? l : r) = m;
    return (l + !cmpL(r)) % n;
  };
```

```
pii res = pii(gao(B, A), gao(A, B)); // (i, j)
if (res.F == res.S) // touching the corner i
  return pii(res.F, -1);
if (!
  cmpL(res.F) && !cmpL(res.S)) // along side i, i+1
  switch ((res.F - res.S + n + 1) % n) {
  case 0: return pii(res.F, res.F);
  case 2: return pii(res.S, res.S);
  }
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned
  in the same order as the line hits the convex */
  return res;
} // convex cut: (r, l]
```

9.13 CircIntersectCirc

9.14 PolyIntersectCirc

```
// Divides into multiple triangle, and sum up
const double PI = acos(-1);
double _area(pdd pa, pdd pb, double r) {
  if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  if (abs(pb) < eps) return 0;</pre>
  double S, h, theta;
  double a = abs(pb), b = abs(pa), c = abs(pb - pa);
  double cosB = dot(pb, pb - pa) / a / c,
         B = acos(cosB);
  double cosC = dot(pa, pb) / a / b, C = acos(cosC);
  if (a > r) {
    S = (C / 2) * r * r;
    h = a * b * sin(C) / c;
    if (h < r && B < PI / 2)
       S = (acos(h / r) * r * r -
        h * sqrt(r * r - h * h));
  } else if (b > r) {
    theta = PI - B - asin(sin(B) / r * a);
    S = .5 * a * r * sin(theta) +
      (C - theta) / 2 * r * r;
  } else S = .5 * sin(C) * a * b;
  return S;
double area_poly_circle(const vector<pdd> poly,
  const pdd &0, const double r) {
  double S = 0;
  for (int i = 0; i < (int)poly.size(); ++i)</pre>
    S += _area(poly[i] - 0,
            poly[(i + 1) % (int)poly.size()] - 0, r) *
      ori(
        0, poly[i], poly[(i + 1) % (int)poly.size()]);
  return fabs(S);
| }
```

9.15 MinkowskiSum

```
vector<pll> Minkowski
   (vector<pll> A, vector<pll> B) { // |A|, |B|>=3}
   hull(A), hull(B);
   vector<pll> C(1, A[0] + B[0]), s1, s2;
   for (int i = 0; i < A.size(); ++i)
   s1.emplace_back(A[(i + 1) % A.size()] - A[i]);</pre>
```

9.16 MinMaxEnclosingRect

```
const double INF = 1e18, qi = acos(-1) / 2 * 3;
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
   hull(dots):
   double Max = 0, Min = INF, deg;
   int n = (int)dots.size();
   dots.emplace_back(dots[0]);
   for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
     pll nw = vec(i + 1);
     while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
       u = (u + 1) \% n;
     while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
      r = (r + 1) \% n;
     if (!i) l = (r + 1) % n;
     while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))
       l = (l + 1) \% n;
     Min = min(Min, (double)(dot(nw, vec(r)) - dot
         (nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));
     deg = acos(dot(diff(r
         , l), vec(u)) / abs(diff(r, l)) / abs(vec(u)));
     deg = (qi - deg) / 2;
    Max = max(Max, abs(diff))
         (r, l) * abs(vec(u)) * sin(deq) * sin(deq));
  }
  return pdd(Min, Max);
}
```

9.17 MinEnclosingCircle

```
| pdd Minimum_Enclosing_Circle
     (vector<pdd> dots, double &r) {
   pdd cent;
   random_shuffle(all(dots));
   cent = dots[0], r = 0;
   for (int i = 1; i < (int)dots.size(); ++i)</pre>
     if (abs(dots[i] - cent) > r) {
       cent = dots[i], r = 0;
       for (int j = 0; j < i; ++j)</pre>
         if (abs(dots[j] - cent) > r) {
           cent = (dots[i] + dots[j]) / 2;
           r = abs(dots[i] - cent);
           for(int k = 0; k < j; ++k)</pre>
             if(abs(dots[k] - cent) > r)
               cent = excenter
                    (dots[i], dots[j], dots[k], r);
         }
  return cent:
}
```

9.18 CircleCover

```
const int N = 1021;
struct CircleCover {
  int C;
  Cir c[N];
  bool g[N][N], overlap[N][N];
  // Area[i] : area covered by at least i circles
  double Area[ N ];
  void init(int _C){ C = _C;}
  struct Teve {
    pdd p; double ang; int add;
    Teve() {}
```

if (q2 == 0) p2 = l2.F.S + l2.S.S, q2 = 2;
if (l1.F == l2.F || l2.F == l2.S) l1 = l2;

return make_tuple((__int128)(p1 * q2), l1) <
 make_tuple((__int128)(p2 * q1), l2);</pre>

```
Teve(pdd _a
         , double _b, int _c):p(_a), ang(_b), add(_c){}
    bool operator<(const Teve &a)const</pre>
    {return ang < a.ang;}
                                                                    Trapezoidalization
                                                            9.20
  eve[N * 2];
  // strict: x = 0, otherwise x = -1
                                                            |template<class T>
  bool disjuct(Cir &a, Cir &b, int x)
                                                            struct SweepLine {
  {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                              struct cmp {
  bool contain(Cir &a, Cir &b, int x)
                                                                 cmp(const SweepLine &_swp): swp(_swp) {}
  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
                                                                 bool operator()(int a, int b) const {
  bool contain(int i, int j) {
                                                                   if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
    /* c[j] is non-strictly in c[i]. */
                                                                    return swp.slope_cmp(a, b);
    return (sign
                                                                  return swp.qet_y(a) + swp.eps < swp.qet_y(b);</pre>
         (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
         R) == 0 \&\& i < j)) \&\& contain(c[i], c[j], -1);
                                                                const SweepLine &swp;
  }
                                                                 _cmp
                                                              T curTime, eps, curQ;
  void solve(){
                                                              vector<Line> base;
    fill_n(Area, C + 2, 0);
                                                              multiset<int, cmp> sweep;
    for(int i = 0; i < C; ++i)</pre>
                                                              multiset<pair<T, int>> event;
      for(int j = 0; j < C; ++j)</pre>
                                                              vector<typename multiset<int, cmp>::iterator> its;
        overlap[i][j] = contain(i, j);
    for(int i = 0; i < C; ++i)</pre>
                                                                   <typename multiset<pair<T, int>>::iterator> eits;
      for(int j = 0; j < C; ++j)</pre>
                                                              bool slope_cmp(int a, int b) const {
        g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                                                                 assert(a != -1);
             disjuct(c[i], c[j], -1));
                                                                 if (b == -1) return 0;
    for(int i = 0; i < C; ++i){</pre>
                                                                 return sign(cross(base
       int E = 0, cnt = 1;
                                                                     [a].Y - base[a].X, base[b].Y - base[b].X)) < 0;
      for(int j = 0; j < C; ++j)</pre>
        if(j != i && overlap[j][i])
                                                              T get_y(int idx) const {
           ++cnt;
                                                                 if (idx == -1) return curQ;
      for(int j = 0; j < C; ++j)</pre>
                                                                Line l = base[idx];
        if(i != j && g[i][j]) {
                                                                 if (l.X.X == l.Y.X) return l.Y.Y;
          pdd aa, bb;
                                                                 return ((curTime - l.X.X) * l.Y.Y
          CCinter(c[i], c[j], aa, bb);
                                                                     + (l.Y.X - curTime) * l.X.Y) / (l.Y.X - l.X.X);
          double A
                atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
                                                              void insert(int idx) {
          double B =
                                                                 its[idx] = sweep.insert(idx);
                atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
                                                                 if (its[idx] != sweep.begin())
          eve[E++] = Teve
                                                                   update_event(*prev(its[idx]));
               (bb, B, 1), eve[E++] = Teve(aa, A, -1);
                                                                 update_event(idx);
          if(B > A) ++cnt;
                                                                 event.emplace
        }
                                                                     (base[idx].Y.X, idx + 2 * (int)base.size());
      if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
      else{
                                                              void erase(int idx) {
        sort(eve, eve + E);
                                                                 assert(eits[idx] == event.end());
        eve[E] = eve[0];
                                                                 auto p = sweep.erase(its[idx]);
        for(int j = 0; j < E; ++j){</pre>
                                                                 its[idx] = sweep.end();
          cnt += eve[j].add;
                                                                 if (p != sweep.begin())
          Area[cnt
                                                                   update_event(*prev(p));
               ] += cross(eve[j].p, eve[j + 1].p) * .5;
          double theta = eve[j + 1].ang - eve[j].ang;
                                                              void update_event(int idx) {
          if (theta < 0) theta += 2. * pi;
                                                                 if (eits[idx] != event.end())
          Area[cnt] += (theta
                                                                  event.erase(eits[idx]);
                - sin(theta)) * c[i].R * c[i].R * .5;
                                                                 eits[idx] = event.end();
                                                                 auto nxt = next(its[idx]);
      }
                                                                 if (nxt ==
    }
                                                                      sweep.end() || !slope_cmp(idx, *nxt)) return;
  }
                                                                 auto t = intersect(base[idx].
|};
                                                                     X, base[idx].Y, base[*nxt].X, base[*nxt].Y).X;
9.19 LineCmp
                                                                 if (t + eps < curTime || t</pre>
                                                                      >= min(base[idx].Y.X, base[*nxt].Y.X)) return;
using Line = pair<pll, pll>;
                                                                 eits[idx
struct lineCmp { // coordinates should be even!
                                                                     ] = event.emplace(t, idx + (int)base.size());
  bool operator()(Line l1, Line l2) const {
                                                              void swp(int idx) {
       (\max(l1.F.F, l2.F.F) + \min(l1.S.F, l2.S.F)) / 2;
                                                                 assert(eits[idx] != event.end());
    ll p1 =
         (X - l1.F.F) * l1.S.S + (l1.S.F - X) * l1.F.S,
                                                                 eits[idx] = event.end();
        p2 =
                                                                 int nxt = *next(its[idx]);
         (X - 12.F.F) * 12.S.S + (12.S.F - X) * 12.F.S,
                                                                 swap((int&)*its[idx], (int&)*its[nxt]);
        q1 = (l1.S.F - l1.F.F), q2 = (l2.S.F - l2.F.F);
                                                                 swap(its[idx], its[nxt]);
    if (q1 == 0) p1 = l1.F.S + l1.S.S, q1 = 2;
                                                                 if (its[nxt] != sweep.begin())
```

update_event(*prev(its[nxt]));

// only expected to call the functions below

update_event(idx);

```
SweepLine(T t, T e, vector<Line> vec): _cmp
       (*this), curTime(t), eps(e), curQ(), base(vec),
        sweep(_cmp), event(), its((int)vec.size(), sweep
       .end()), eits((int)vec.size(), event.end()) {
    for (int i = 0; i < (int)base.size(); ++i) {</pre>
      auto &[p, q] = base[i];
      if (p > q) swap(p, q);
       if (p.X <= curTime && curTime <= q.X)</pre>
         insert(i);
       else if (curTime < p.X)</pre>
         event.emplace(p.X, i);
  }
  void setTime(T t, bool ers = false) {
    assert(t >= curTime);
    while (!event.empty() && event.begin()->X <= t) {</pre>
      auto [et, idx] = *event.begin();
      int s = idx / (int)base.size();
      idx %= (int)base.size();
      if (abs(et - t) <= eps && s == 2 && !ers) break;
curTime = et;</pre>
      event.erase(event.begin());
      if (s == 2) erase(idx);
      else if (s == 1) swp(idx);
      else insert(idx);
    curTime = t;
  }
  T nextEvent() {
    if (event.empty()) return INF;
    return event.begin()->X;
  int lower_bound(T y) {
    curQ = y;
    auto p = sweep.lower_bound(-1);
    if (p == sweep.end()) return -1;
    return *p:
|};
```

9.21 HalfPlaneIntersect

```
pll area_pair(Line a, Line b)
{ return pll(cross(a.S
     bool isin(Line l0, Line l1, Line l2) {
 // Check inter(l1, l2) strictly in l0
 auto [a02X, a02Y] = area_pair(l0, l2);
 auto [a12X, a12Y] = area_pair(l1, l2);
 if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
 return (__int128)
       a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
 sort(all(arr), [&](Line a, Line b) -> int {
   if (cmp(a.S - a.F, b.S - b.F, 0) != -1)
     return cmp(a.S - a.F, b.S - b.F, 0);
   return ori(a.F, a.S, b.S) < 0;</pre>
 });
 deque<Line> dq(1, arr[0]);
 for (auto p : arr) {
   if (cmp(
        dq.back().S - dq.back().F, p.S - p.F, 0) == -1)
      continue;
   while ((int)dq.size() >= 2
       && !isin(p, dq[(int)dq.size() - 2], dq.back()))
     dq.pop_back();
   while
        ((int)dq.size() \ge 2 \&\& !isin(p, dq[0], dq[1]))
     dq.pop_front();
   dq.emplace_back(p);
 }
 while ((int)dq.size() >= 3 &&
       !isin(dq[0], dq[(int)dq.size() - 2], dq.back())) | if(a.tri) a.tri->edge[a.side] = b;
```

```
dq.pop_back();
while ((int)
    dq.size() >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
  dq.pop_front();
return vector<Line>(all(dq));
```

9.22 RotatingSweepLine

```
void rotatingSweepLine(vector<pii> &ps) {
   int n = (int)ps.size(), m = 0;
   vector<int> id(n), pos(n);
   vector<pii> line(n * (n - 1));
   for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++j)</pre>
       if (i != j) line[m++] = pii(i, j);
   sort(all(line), [&](pii a, pii b) {
     return cmp(ps[a.S] - ps[a.F], ps[b.S] - ps[b.F]);
   }); // cmp(): polar angle compare
   iota(all(id), 0);
   sort(all(id), [&](int a, int b) {
     if (ps[a].S != ps[b].S) return ps[a].S < ps[b].S;</pre>
     return ps[a] < ps[b];</pre>
   }); // initial order, since (1, 0) is the smallest
   for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
   for (int i = 0; i < m; ++i) {</pre>
     auto l = line[i];
     // do something
     tie(pos[l.F], pos[l.S], id[pos[l.F]], id[pos[l.S
         ]]) = make_tuple(pos[l.S], pos[l.F], l.S, l.F);
}
```

9.23 DelaunayTriangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
*/
const
     ll inf = MAXC * MAXC * 100; // lower_bound unknown
struct Tri:
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
};
struct Tri {
  pll p[3];
  Edge edge[3];
  Tri* chd[3];
  Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)
      if (ori(p[i], p[(i + 1) % 3], q) < 0)
        return 0:
    return 1;
  }
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
```

```
if(b.tri) b.tri->edge[b.side] = a;
                                                              tris = pool; triang.clear(); vst.clear();
                                                              random_shuffle(ps, ps + n);
                                                              Trig tri; // the triangulation structure
struct Trig { // Triangulation
                                                              for (int i = 0; i < n; ++i)</pre>
  Triq() {
    the root
                                                                 tri.add_point(ps[i]);
          = // Tri should at least contain all points
                                                              go(tri.the_root);
       new(tris++) Tri(pll(-inf, -inf),
            pll(inf + inf, -inf), pll(-inf, inf + inf));
                                                             9.24 VonoroiDiagram
  Tri* find(pll p) { return find(the_root, p); }
                                                            // all coord. is even
  void add_point(const
                                                                   you may want to call halfPlaneInter after then
        pll &p) { add_point(find(the_root, p), p); }
                                                            vector<vector<Line>> vec;
  Tri* the_root;
                                                            void build_voronoi_line(int n, pll *arr) {
  static Tri* find(Tri* root, const pll &p) {
                                                              tool.init(n, arr); // Delaunay
    while (1) {
                                                              vec.clear(), vec.resize(n);
       if (!root->has_chd())
                                                              for (int i = 0; i < n; ++i)</pre>
        return root;
                                                                for (auto e : tool.head[i]) {
       for (int i = 0; i < 3 && root->chd[i]; ++i)
                                                                   int u = tool.oidx[i], v = tool.oidx[e.id];
        if (root->chd[i]->contains(p)) {
                                                                   pll m = (arr[v
           root = root->chd[i];
                                                                       ] + arr[u]) / 2LL, d = perp(arr[v] - arr[u]);
           break;
                                                                   vec[u].emplace_back(Line(m, m + d));
    }
                                                            |}
    assert(0); // "point not found"
  }
                                                                   Misc
                                                             10
  void add_point(Tri* root, pll const& p) {
    Tri* t[3];
                                                            10.1 MoAlgoWithModify
     /* split it into three triangles */
    for (int i = 0; i < 3; ++i)
                                                            // Mo's Algorithm With modification
      t[i] = new(tris
                                                            // Block: N^{2/3}, Complexity: N^{5/3}
           ++) Tri(root->p[i], root->p[(i + 1) % 3], p);
                                                            struct Query {
                                                              const int blk = 2000;
    for (int i = 0; i < 3; ++i)
                                                              int L, R, LBid, RBid, T;
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
                                                              Query(int l, int r, int t):
    for (int i = 0; i < 3; ++i)
                                                                L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
                                                              bool operator<(const Query &q) const {</pre>
    for (int i = 0; i < 3; ++i)
                                                                 if (LBid != q.LBid) return LBid < q.LBid;</pre>
      root->chd[i] = t[i];
                                                                 if (RBid != q.RBid) return RBid < q.RBid;</pre>
    for (int i = 0; i < 3; ++i)
                                                                 return T < q.T;</pre>
      flip(t[i], 2);
                                                              }
                                                            };
  void flip(Tri* tri, int pi) {
                                                            void solve(vector<Query> query) {
    Tri* trj = tri->edge[pi].tri;
                                                              sort(all(query));
    int pj = tri->edge[pi].side;
                                                              int L=0, R=0, T=-1;
    if (!trj) return;
                                                              for (auto q : query) {
    if (!in_cc(tri->p
                                                                 while (T < q.T) addTime(L, R, ++T); // TODO
         [0], tri->p[1], tri->p[2], trj->p[pj])) return;
                                                                 while (T > q.T) subTime(L, R, T--); // TODO
     /* flip edge between tri,trj */
                                                                 while (R < q.R) add(arr[++R]); // TODO</pre>
    Tri* trk = new(tris++) Tri
                                                                 while (L > q.L) add(arr[--L]); // TODO
         (tri->p[(pi + 1) % 3], trj->p[pj], tri->p[pi]);
                                                                while (R > q.R) sub(arr[R--]); // TODO
    Tri* trl = new(tris++) Tri
                                                                while (L < q.L) sub(arr[L++]); // TODO</pre>
         (trj->p[(pj + 1) % 3], tri->p[pi], trj->p[pj]);
                                                                 // answer query
    edge(Edge(trk, 0), Edge(trl, 0));
                                                              }
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
                                                           }
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
                                                             10.2 MoAlgoOnTree
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd
         [0] = trk; tri->chd[1] = trl; tri->chd[2] = 0;
                                                            Mo's Algorithm On Tree
    trj->chd
                                                            Preprocess:
         [0] = trk; trj->chd[1] = trl; trj->chd[2] = 0;
                                                            1) LCA
    flip(trk, 1); flip(trk, 2);
                                                            2) dfs with in[u] = dft++, out[u] = dft++
    flip(trl, 1); flip(trl, 2);
                                                            3) ord[in[u]] = ord[out[u]] = u
  }
                                                            4) bitset<MAXN> inset
};
                                                            */
vector<Tri*> triang; // vector of all triangle
                                                            struct Query {
set<Tri*> vst;
                                                              int L, R, LBid, lca;
void go(Tri* now) { // store all tri into triang
                                                              Query(int u, int v) {
  if (vst.find(now) != vst.end())
                                                                 int c = LCA(u, v);
    return;
                                                                 if (c == u || c == v)
  vst.insert(now);
                                                                   q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
  if (!now->has_chd())
                                                                 else if (out[u] < in[v])</pre>
    return triang.emplace_back(now);
                                                                   q.lca = c, q.L = out[u], q.R = in[v];
  for (int i = 0; i < now->num_chd(); ++i)
                                                                 else
    go(now->chd[i]);
                                                                   q.lca = c, q.L = out[v], q.R = in[u];
                                                                 q.Lid = q.L / blk;
void build(int n, pll* ps) { // build triangulation
```

```
bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     return R < q.R;</pre>
   }
};
 void flip(int x) {
     if (inset[x]) sub(arr[x]); // TODO
     else add(arr[x]); // TODO
     inset[x] = ~inset[x];
 void solve(vector<Query> query) {
   sort(ALL(query));
   int L = 0, R = 0;
   for (auto q : query) {
     while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
     while (R > q.R) flip(ord[R--]);
     while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
   }
}
```

10.3 MoAlgoAdvanced

```
• Mo's Algorithm With Addition Only
```

- Sort querys same as the normal Mo's algorithm.
- For each query [l,r]:
- If l/blk = r/blk, brute-force.
- If $l/blk \neq curL/blk$, initialize $curL := (l/blk+1) \cdot blk, curR :=$ curL-1
- If r > curR, increase curR
- decrease $\widehat{cur}L$ to fit l, and then undo after answering
- Mo's Algorithm With Offline Second Time
 - Require: Changing answer \equiv adding f([l,r],r+1).
 - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l],r+1). Part1: Answer all f([1,r],r+1) first.

 - Part2: Store $curR \rightarrow R$ for curL (reduce the space to O(N)), and then answer them by the second offline
 - Note: You must do the above symmetrically for the left boundaries.

10.4 HilbertCurve

```
ll hilbert(int n, int x, int y) {
  11 \text{ res} = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y \& s) > 0;
    res += s * 111 * s * ((3 * rx) ^ ry);
    if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
    }
  }
  return res;
} // n = 2^k
```

10.5 SternBrocotTree

- Construction: Root $\frac{1}{1}$, left/right neighbor $\frac{0}{1}, \frac{1}{0}$, each node is sum of last left/right neighbor: $rac{a}{b},rac{c}{d}
 ightarrowrac{a+c}{b+d}$
- Property: Adjacent (mid-order DFS) $\frac{a}{b}, \frac{c}{d} \Rightarrow bc ad = 1$.
- Search known $\frac{p}{a}$: keep L-R alternative. Each step can calcaulated in $O(1) \Rightarrow$ total $O(\log C)$.
- Search unknown $rac{p}{q}$: keep L-R alternative. Each step can calcaulated in $O(\log C)$ checks \Rightarrow total $O(\log^2 C)$ checks.

10.6 AllLCS

```
void all_lcs(string s, string t) { // O-base
  vector<int> h((int)t.size());
  iota(all(h), 0);
  for (int a = 0; a < (int)s.size(); ++a) {</pre>
    int v = -1;
    for (int c = 0; c < (int)t.size(); ++c)</pre>
      if (s[a] == t[c] || h[c] < v)</pre>
        swap(h[c], v);
    // LCS(s[0, a], t[b, c]) =
```

```
// c - b + 1 - sum([h[i] >= b] | i <= c)
  // h[i] might become -1 !!
}
```

10.7 SimulatedAnnealing

```
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {</pre>
     // ans:
         answer, nw: current value, rnd(): mt19937 rnd()
     if (exp(-(nw - ans
         ) / factor) >= (double)(rnd() % base) / base)
         ans = nw;
     factor *= 0.99995;
}
```

10.8 SMAWK

}

```
int opt[N];
ll A(int x, int y); // target func
void smawk(vector<int> &r, vector<int> &c);
void interpolate(vector<int> &r, vector<int> &c) {
  int n = (int)r.size();
  vector<int> er;
  for (int i = 1; i < n; i += 2) er.emplace_back(r[i]);</pre>
  smawk(er, c);
  for (int i = 0, j = 0; j < c.size(); j++) {</pre>
     if (A(r[i], c[j]) < A(r[i], opt[r[i]]))</pre>
       opt[r[i]] = c[j];
     if (i + 2 < n && c[j] == opt[r[i + 1]])
       j--, i += 2;
}
void reduce(vector<int> &r, vector<int> &c) {
  int n = (int)r.size();
  vector<int> nc;
  for (int i : c) {
    int j = (int)nc.size();
     while (
       j \& A(r[j-1], nc[j-1]) > A(r[j-1], i))
       nc.pop_back(), j--;
     if (nc.size() < n) nc.emplace_back(i);</pre>
  }
  smawk(r, nc);
}
void smawk(vector<int> &r, vector<int> &c) {
  if (r.size() == 1 && c.size() == 1) opt[r[0]] = c[0];
  else if (r.size() >= c.size()) interpolate(r, c);
  else reduce(r, c);
1 }
```

10.9 Python

import math math.isqrt(2) # integer sqrt

10.10 LineContainer

```
struct Line {
  mutable ll k, m, p;
  bool operator<(const Line &o) const {</pre>
    return k < o.k;
  bool operator<(ll x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>>> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b)
  static const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b);
  bool isect(iterator x, iterator y) {
    if (y == end()) return x->p = inf, 0;
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
```

```
else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
}

void add(ll k, ll m) {
    auto z = insert({k, m, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y))
        isect(x, y = erase(y));
    while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
}

ll query(ll x) {
    assert(!empty());
    auto l = *lower_bound(x);
    return l.k * x + l.m;
}
};
```