6.100uadraticResidue .

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```

Basic 1

1.1 .vimrc

```
set ru nu cin cul sc so=3 ts=4 sw=4 bs=2 ls=2 mouse=a
ino {<CR> {<CR>}<C-o>0
ino jj <esc>
ino jk <esc>
map <F7> :w<CR>:!g++
    "%" -std=c++17 -Wall -Wextra -Wshadow -Wconversion
      -fsanitize=address,undefined -q && ./a.out<CR>
ca Hash w !cpp -dD -P -fpreprocessed
     \| tr -d "[:space:]" \| md5sum \| cut -c-6
```

1.2 PBDS

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
tree<int, null_type, less<int>, rb_tree_tag
     tree_order_statistics_node_update> bst;
// order_of_key(n): # of elements <= n
// find_by_order(n): 0-indexed
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
__gnu_pbds::priority_queue
    <int, greater<int>, thin_heap_tag> pq;
```

1.3 pargma

```
#pragma GCC optimize("Ofast,unroll-loops")
#pragma GCC target("avx,avx2,sse,sse2
    ,sse3,ssse3,sse4,popcnt,abm,mmx,fma,tune=native")
// chrono
    ::steady_clock::now().time_since_epoch().count()
    Graph
2.1 2SAT/SCC
  int low[N], dfn[N], bln[N], n, Time, nScc;
```

```
struct SAT { // O-base
   bool instack[N], istrue[N];
   stack<int> st;
   vector<int> G[N], SCC[N];
   void init(int _n) {
     n = _n; // assert(n * 2 <= N);
     for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
   void add_edge(int a, int b) { G[a].emplace_back(b); }
   int rv(int a) {
     if (a >= n) return a - n;
     return a + n;
   void add_clause(int a, int b) {
     add_edge(rv(a), b), add_edge(rv(b), a);
   void dfs(int u) {
     dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
       do {
         tmp = st.top(), st.pop();
         instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != u);
       ++nScc;
     }
   bool solve() {
     Time = nScc = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       if (!dfn[i]) dfs(i);
     for (int i =
         0; i < n + n; ++i) SCC[bln[i]].emplace_back(i);
     for (int i = 0; i < n; ++i) {</pre>
       if (bln[i] == bln[i + n]) return false;
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true;
   }
|};
```

2.2 BCC Vertex

```
int n, m, dfn[N], low[N], is_cut[N], nbcc = 0, t = 0;
vector<int> g[N], bcc[N], G[2 * N];
stack<int> st;
void tarjan(int p, int lp) {
  dfn[p] = low[p] = ++t;
  st.push(p);
  for (auto i : g[p]) {
    if (!dfn[i]) {
      tarjan(i, p);
      low[p] = min(low[p], low[i]);
      if (dfn[p] <= low[i]) {</pre>
        nbcc++;
```

```
is_cut[p] = 1;
    for (int x = 0; x != i; st.pop()) {
        x = st.top();
        bcc[nbcc].push_back(x);
    }
    bcc[nbcc].push_back(p);
    }
    lelse low[p] = min(low[p], dfn[i]);
}

void build() { // [n+1,n+nbcc] cycle, [1,n] vertex
    for (int i = 1; i <= nbcc; i++) {
        for (auto j : bcc[i]) {
            G[i + n].push_back(j);
            G[j].push_back(i + n);
        }
    }
}</pre>
```

2.3 MinimumMeanCycle

```
|/* 0(V^3)
|let dp[i][j] = min length from 1 to j exactly i edges
|ans = min (dp[n + 1][u] - dp[i][u]) / (n + 1 - i) */
```

2.4 MaximumCliqueDyn

```
struct MaxClique { // fast when N <= 100</pre>
 bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
   n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
 }
 void add_edge(int u, int v) {
   G[u][v] = G[v][u] = 1;
 void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(all(r)
          , [&](int x, int y) { return d[x] > d[y]; });
   }
    vector<int> c(r.size());
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first
          (); p < N; p = cs[k]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
   dfs(r, c, l + 1, mask);
 }
 void dfs(vector<</pre>
      int> &r, vector<int> &c, int l, bitset<N> mask) {
   while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.emplace_back(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
    }
 }
 int solve() {
   vector<int> r(n);
   ans = q = 0, iota(all(r), 0);
```

```
pre_dfs(r, 0, bitset<N>(string(n, '1')));
     return ans;
|};
2.5 DMST(slow)
| struct zhu_liu { // O(VE)
   struct edge {
     int u, v;
     ll w;
   vector<edge> E; // 0-base
   int pe[N], id[N], vis[N];
   ll in[N];
   void init() { E.clear(); }
   void add_edge(int u, int v, ll w) {
     if (u != v) E.emplace_back(edge{u, v, w});
   ll build(int root, int n) {
     ll ans = 0;
     for (;;) {
       fill_n(in, n, INF);
       for (int i = 0; i < E.size(); ++i)</pre>
         if (E[i].u != E[i].v && E[i].w < in[E[i].v])</pre>
           pe[E[i].v] = i, in[E[i].v] = E[i].w;
       for (int u = 0; u < n; ++u) // no solution
         if (u != root && in[u] == INF) return -INF;
       int cntnode = 0;
       fill_n(id, n, -1), fill_n(vis, n, -1);
       for (int u = 0; u < n; ++u) {
         if (u != root) ans += in[u];
         int v = u;
         while (vis[v] != u && !~id[v] && v != root)
           vis[v] = u, v = E[pe[v]].u;
         if (v != root && !~id[v]) {
           for (int x = E[pe[v]].u; x != v;
                x = E[pe[x]].u)
             id[x] = cntnode;
           id[v] = cntnode++;
         }
       }
       if (!cntnode) break; // no cycle
       for (int u = 0; u < n; ++u)
         if (!~id[u]) id[u] = cntnode++;
       for (int i = 0; i < E.size(); ++i) {</pre>
         int v = E[i].v;
         E[i].v = id[E[i].v], E[i].v = id[E[i].v];
         if (E[i].u != E[i].v) E[i].w -= in[v];
       }
      n = cntnode, root = id[root];
     }
     return ans;
|};
2.6 DMST
#define rep(i, a, b) for (int i = a; i < (b); ++i)</pre>
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef vector<int> vi;
struct RollbackUF {
   vi e;
   vector<pii> st;
   RollbackUF(int n) : e(n, -1) {}
   int size(int x) { return -e[find(x)]; }
   int find(int x) { return e[x] < 0 ? x : find(e[x]); }</pre>
   int time() { return sz(st); }
   void rollback(int t) {
     for (int i = time(); i-- > t;)
       e[st[i].first] = st[i].second;
     st.resize(t);
  }
```

bool join(int a, int b) {

a = find(a), b = find(b);
if (a == b) return false;

```
if (e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]});
                                                           }
    st.push_back({b, e[b]});
    e[a] += e[b];
    e[b] = a;
    return true;
};
struct Edge {
  int a, b;
 ll w;
struct Node { /// lazy skew heap node
  Edge key;
  Node *l, *r;
  ll delta;
  void prop() {
    key.w += delta;
    if (l) l->delta += delta;
    if (r) r->delta += delta;
    delta = 0;
 Edge top() {
    prop();
    return key;
};
Node *merge(Node *a, Node *b) {
 if (!a || !b) return a ?: b;
 a->prop(), b->prop();
 if (a->key.w > b->key.w) swap(a, b);
  swap(a->l, (a->r = merge(b, a->r)));
  return a;
void pop(Node *&a) {
 a->prop();
 a = merge(a->l, a->r);
pair<ll, vi> dmst(int n, int r, vector<Edge> &g) {
 RollbackUF uf(n);
  vector<Node *> heap(n);
  for (Edge e : g)
   heap[e.b] = merge(heap[e.b], new Node{e});
  11 res = 0;
  vi seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector<Edge> Q(n), in(n, {-1, -1}), comp;
  deque<tuple<int, int, vector<Edge>>> cycs;
  rep(s, 0, n) {
    int u = s, qi = 0, w;
    while (seen[u] < 0) {</pre>
      if (!heap[u]) return {-1, {}};
      Edge e = heap[u]->top();
      heap[u]->delta -= e.w, pop(heap[u]);
      Q[qi] = e, path[qi++] = u, seen[u] = s;
      res += e.w, u = uf.find(e.a);
      if (seen[u] == s) { /// found cycle, contract
        Node *cyc = 0;
        int end = qi, time = uf.time();
        do cyc = merge(cyc, heap[w = path[--qi]]);
        while (uf.join(u, w));
        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
        cycs.push_front(\{u, time, \{\&Q[qi], \&Q[end]\}\}\);
      }
    }
    rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
  for (auto &[u, t, comp] :
   cycs) { // restore sol (optional)
    uf.rollback(t);
    Edge inEdge = in[u];
    for (auto &e : comp) in[uf.find(e.b)] = e;
    in[uf.find(inEdge.b)] = inEdge;
 }
 rep(i, 0, n) par[i] = in[i].a;
```

```
2.7 VizingTheorem
```

return {res, par};

```
namespace Vizing { // Edge coloring
                     // G: coloring adjM
int C[N][N], G[N][N];
void clear(int n) {
  for (int i = 0; i <= n; i++) {</pre>
    for (int j = 0; j <= n; j++) C[i][j] = G[i][j] = 0;</pre>
}
void solve(vector<pii> &E, int n, int m) {
  int X[n] = {}, a;
  auto update = [&](int u) {
    for (X[u] = 1; C[u][X[u]]; X[u]++);
  auto color = [&](int u, int v, int c) {
     int p = G[u][v];
     G[u][v] = G[v][u] = c;
    C[u][c] = v;
    C[v][c] = u;
    C[\upsilon][p] = C[\upsilon][p] = 0;
     if (p) X[u] = X[v] = p;
     else update(u), update(v);
    return p;
  };
  auto flip = [&](int u, int c1, int c2) {
     int p = C[u][c1];
     swap(C[u][c1], C[u][c2]);
     if (p) G[u][p] = G[p][u] = c2;
     if (!C[u][c1]) X[u] = c1;
     if (!C[u][c2]) X[u] = c2;
     return p;
  };
  for (int i = 1; i <= n; i++) X[i] = 1;</pre>
  for (int t = 0; t < E.size(); t++) {</pre>
     int u = E[t].first, v0 = E[t].second, v = v0,
         c0 = X[u], c = c0, d;
     vector<pii> L;
     int vst[n] = {};
     while (!G[u][v0]) {
       L.emplace_back(v, d = X[v]);
       if (!C[v][c])
         for (a = (int)L.size() - 1; a >= 0; a--)
           c = color(u, L[a].first, c);
       else if (!C[u][d])
         for (a = (int)L.size() - 1; a >= 0; a--)
           color(u, L[a].first, L[a].second);
       else if (vst[d]) break;
       else vst[d] = 1, v = C[v][d];
     if (!G[u][v0]) {
       for (; v; v = flip(v, c, d), swap(c, d));
       if (C[u][c0]) {
         for (a = (int)L.size() - 2;
              a >= 0 && L[a].second != c; a--)
         for (; a >= 0; a--)
           color(u, L[a].first, L[a].second);
       } else t--;
    }
|} // namespace Vizing
```

2.8 MinimumCliqueCover

```
| struct Clique_Cover { // 0-base, 0(n2^n) | int co[1 << N], n, E[N]; | int dp[1 << N]; | void init(int _n) { | n = _n, fill_n(dp, 1 << n, 0); | fill_n(E, n, 0), fill_n(co, 1 << n, 0); | }
```

```
void add_edge(int u, int v) {
    E[u] \mid = 1 << v, E[v] \mid = 1 << u;
  }
  int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;
     dp[0] = (n \& 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {</pre>
       int t = i & -i;
       dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)</pre>
       co[i] = (co[i] \& i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {</pre>
       int sum = 0; // probabilistic
       for (int i = 0; i < (1 << n); ++i)</pre>
         sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n;
  }
|};
```

2.9 CountMaximalClique

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
   int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    n = _n;
     for (int i = 1; i <= n; ++i)</pre>
       for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
     g[v][v] = g[v][v] = 1;
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
     if (sn == 0 && nn == 0) ++S;
     int u = some[d][0];
     for (int i = 0; i < sn; ++i) {</pre>
       int v = some[d][i];
       if (g[v][v]) continue;
       int tsn = 0, tnn = 0;
       copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)</pre>
         if (g[v][some[d][j]])
           some[d + 1][tsn++] = some[d][j];
       for (int j = 0; j < nn; ++j)</pre>
         if (g[v][none[d][j]])
           none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
       some[d][i] = 0, none[d][nn++] = v;
    }
  }
  int solve() {
    iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
     return S;
|};
```

2.10 Theorems

 $|\max \text{ independent edge set}| = |V| - |\min \text{ edge cover}| \\ |\max \text{ independent set}| = |V| - |\min \text{ vertex cover}| \\$

3 Flow-Matching

3.1 HopcroftKarp

```
| struct hopcroftKarp { // 0-based
| bool dfs(int a, int L, vector<vector<int>> &g,
| vector<int> &btoa, vector<int> &A,
| vector<int> &B) {
```

```
if (A[a] != L) return 0;
     A[a] = -1;
     for (int b : g[a])
       if (B[b] == L + 1) {
         B[b] = 0;
         if (btoa[b] == -1 ||
           dfs(btoa[b], L + 1, g, btoa, A, B))
           return btoa[b] = a, 1;
    return 0;
  }
  int solve(vector<vector<int>> &g, int m) {
    int res = 0;
     vector<int> btoa(-1, m), A(g.size()),
       B(btoa.size()), cur, next;
     for (;;) {
       fill(all(A), 0), fill(all(B), 0);
       cur.clear();
       for (int a : btoa)
         if (a != -1) A[a] = -1;
       for (int a = 0; a < g.size(); a++)</pre>
         if (A[a] == 0) cur.push_back(a);
       /// Find all layers using bfs.
       for (int lay = 1;; lay++) {
         bool islast = 0;
         next.clear();
         for (int a : cur)
           for (int b : g[a]) {
             if (btoa[b] == -1) {
               B[b] = lay;
               islast = 1:
             } else if (btoa[b] != a && !B[b]) {
               B[b] = lay;
               next.push_back(btoa[b]);
             }
          }
         if (islast) break;
         if (next.empty()) return res;
         for (int a : next) A[a] = lay;
         cur.swap(next);
       /// Use DFS to scan for augmenting paths.
      for (int a = 0; a < g.size(); a++)</pre>
         res += dfs(a, 0, g, btoa, A, B);
  }
|};
```

3.2 KM

```
struct KM { // O-base
  ll w[N][N], hl[N], hr[N], slk[N];
  int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
  bool vl[N], vr[N];
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)</pre>
      fill_n(w[i], n, -INF);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool Check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    fill_n(slk
         , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (ll d;;) {
      while (ql < qr)
        for (int x = 0, y = qu[ql++]; x < n; ++x)
```

```
if (!vl[x] && slk
               [x] >= (d = hl[x] + hr[y] - w[x][y])) {
             if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
        if (vl[x]) hl[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!vl[x] && !slk[x] && !Check(x)) return;
    }
  }
  ll solve() {
    fill_n(fl
         , n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
    for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(w[i], w[i] + n);
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    ll res = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
  }
|};
```

3.3 MCMF

```
struct MinCostMaxFlow { // O-base
  struct Edge {
    ll from, to, cap, flow, cost, rev;
  } *past[N];
 vector<Edge> G[N];
 int inq[N], n, s, t;
 ll dis[N], up[N], pot[N];
 bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[v] = d, vp[v] = cap, past[v] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
   };
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[v] = 0;
      for (auto &e : G[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
        relax
            (e.to, d2, min(up[u], e.cap - e.flow), &e);
     }
    }
    return dis[t] != INF;
 }
  void solve(int _s
      , int _t, ll &flow, ll &cost, bool neg = true) {
    s = _s, t = _t, flow = 0, cost = 0;
    if (neg) BellmanFord(), copy_n(dis, n, pot);
    for (; BellmanFord(); copy_n(dis, n, pot)) {
      for (int
          i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], G[e.to][e.rev].flow -= up[t];
      }
   }
 }
  void init(int _n) {
    n = _n, fill_n(pot, n, 0);
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
```

3.4 GeneralGraphMatching

```
struct Matching { // O-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
   { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
  }
  void Blossom(int x, int y, int l) {
    for (; Find(x) != l; x = pre[y]) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
       for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
    }
  }
  bool Bfs(int r) {
    iota(all(fa), 0); fill(all(s), -1);
    q = queue<int>(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : G[x])
         if (s[u] == -1) {
           if (pre[u] = x, s[u] = 1, match[u] == n) {
             for (int a = u, b = x, last;
                b != n; a = last, b = pre[a])
               last =
                   match[b], match[b] = a, match[a] = b;
             return true;
          }
          q.push(match[u]); s[match[u]] = 0;
         } else if (!s[v] && Find(v) != Find(x)) {
          int l = LCA(u, x);
           Blossom(x, u, l); Blossom(u, x, l);
        }
    }
    return false;
  Matching(int_n): n(n), fa(n + 1), s(n + 1), vis
       (n + 1), pre(n + 1, n), match(n + 1, n), G(n) {}
  void add_edge(int u, int v)
  { G[u].emplace_back(v), G[v].emplace_back(u); }
  int solve() {
    int ans = 0;
     for (int x = 0; x < n; ++x)</pre>
      if (match[x] == n) ans += Bfs(x);
    return ans;
  } // match[x] == n means not matched
};
```

3.5 MaxWeightMaching

```
#define rep(i, l, r) for (int i = (l); i <= (r); ++i)
struct WeightGraph { // 1-based
    struct edge {
        int u, v, w;
    };
    int n, nx;
    vector<int> lab;
    vector<vector<edge>> g;
    vector<iint> slack, match, st, pa, S, vis;
    vector<vector<iint>> flo, flo_from;
    queue<iint> q;
```

```
WeightGraph(int n_)
  : n(n_{}), nx(n * 2), lab(nx + 1),
    g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1),
    flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
  match = st = pa = S = vis = slack;
  rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
int ED(edge e) {
  return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
void update_slack(int u, int x, int &s) {
  if (!s || ED(g[u][x]) < ED(g[s][x])) s = u;</pre>
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)</pre>
    if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
      update_slack(u, x, slack[x]);
void q_push(int x) {
  if (x \le n) q.push(x);
  else
    for (int y : flo[x]) q_push(y);
void set_st(int x, int b) {
  st[x] = b;
  if(x > n)
    for (int y : flo[x]) set_st(y, b);
vector<int> split_flo(auto &f, int xr) {
  auto it = find(ALL(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
    reverse(1 + ALL(f)), it = f.end() - pr;
  auto res = vector(f.begin(), it);
  return f.erase(f.begin(), it), res;
}
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  int xr = flo_from[u][g[u][v].u];
  auto &f = flo[u], z = split_flo(f, xr);
  rep(i, 0, (int)z.size() - 1)
    set_match(z[i], z[i ^ 1]);
  set_match(xr, v);
  f.insert(f.end(), all(z));
}
void augment(int u, int v) {
  for (;;) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
}
int lca(int u, int v) {
  static int t = 0;
  ++t;
  for (++t; u || v; swap(u, v))
    if (u) {
      if (vis[u] == t) return u;
      vis[u] = t;
      u = st[match[u]];
      if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + all(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0;
  match[b] = match[o];
  vector<int> f = {o};
  for (int x = u, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
```

```
reverse(1 + all(f));
  for (int x = v, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  flo[b] = f;
  set_st(b, b);
  for (int x = 1; x <= nx; ++x)</pre>
    g[b][x].w = g[x][b].w = 0;
  fill(all(flo_from[b]), 0);
  for (int xs : flo[b]) {
    for (int x = 1; x <= nx; ++x)
      if (g[b][x].w == 0 ||
        ED(g[xs][x]) < ED(g[b][x])
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)</pre>
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
}
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) {
      xs = x:
      continue;
    }
    pa[xs] = g[x][xs].u;
    S[xs] = 1, S[x] = 0;
    slack[xs] = 0;
    set_slack(x);
    q_push(x);
    xs = -1;
  }
  for (int x : flo[b])
    if (x == xr) S[x] = 1, pa[x] = pa[b];
    else S[x] = -1, set_slack(x);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
    int nu = st[match[v]];
    pa[v] = e.u;
    S[v] = 1;
    slack[v] = slack[nu] = 0;
    S[nu] = 0;
    q_push(nu);
  } else if (S[v] == 0) {
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
  return false;
bool matching() {
  fill(all(S), -1), fill(all(slack), 0);
  q = queue<int>();
  for (int x = 1; x <= nx; ++x)
    if (st[x] == x \&\& !match[x])
      pa[x] = 0, S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
    while (q.size()) {
      int u = q.front();
      q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (ED(g[v][v]) != 0)
            update_slack(u, st[v], slack[st[v]]);
          else if (on_found_edge(g[u][v]))
            return true;
        }
    int d = INF;
    for (int b = n + 1; b <= nx; ++b)
```

```
if (st[b] == b && S[b] == 1)
           d = min(d, lab[b] / 2);
       for (int x = 1; x <= nx; ++x)</pre>
         if (int s = slack[x];
             st[x] == x && s && s[x] <= 0)
           d = min(d, ED(g[s][x]) / (S[x] + 2));
       for (int u = 1; u <= n; ++u)</pre>
         if (S[st[v]] == 1) lab[v] += d;
         else if (S[st[u]] == 0) {
           if (lab[u] <= d) return false;</pre>
           lab[u] -= d;
       rep(b, n + 1, nx) if (st[b] == b \&\& S[b] >= 0)
         lab[b] += d * (2 - 4 * S[b]);
       for (int x = 1; x <= nx; ++x)
         if (int s = slack[x]; st[x] == x && s &&
             st[s] != x \&\& ED(g[s][x]) == 0)
           if (on_found_edge(g[s][x])) return true;
       for (int b = n + 1; b <= nx; ++b)
         if (st[b] == b && S[b] == 1 && lab[b] == 0)
           expand_blossom(b);
     return false;
  }
  pair<ll, int> solve() {
    fill(all(match), 0);
     rep(u, 0, n) st[u] = u, flo[u].clear();
     int w_max = 0;
     rep(u, 1, n) rep(v, 1, n) {
       flo_from[u][v] = (u == v ? u : 0);
       w_{max} = max(w_{max}, g[v][v].w);
    fill(all(lab), w_max);
     int n_matches = 0;
     tot_weight = 0;
     while (matching()) ++n_matches;
     rep(u, 1, n) if (match[u] \&\& match[u] < u)
       tot_weight += g[u][match[u]].w;
    return make_pair(tot_weight, n_matches);
  }
  void add_edge(int u, int v, int w) {
     g[v][v].w = g[v][v].w = w;
  }
|};
```

GlobalMinCut 3.6

```
#undef INF
struct SW{ // global min cut, O(V^3)
 #define REP for (int i = 0; i < n; ++i)
  static const int MXN = 514, INF = 2147483647;
 int vst[MXN], edge[MXN][MXN], wei[MXN];
 void init(int n) {
   REP fill_n(edge[i], n, 0);
 void addEdge(int u, int v, int w){
    edge[v][v] += w; edge[v][v] += w;
 int search(int &s, int &t, int n){
   fill_n(vst, n, 0), fill_n(wei, n, 0);
    s = t = -1;
    int mx, cur;
    for (int j = 0; j < n; ++j) {</pre>
     mx = -1, cur = 0;
     REP if (wei[i] > mx) cur = i, mx = wei[i];
     vst[cur] = 1, wei[cur] = -1;
     s = t; t = cur;
     REP if (!vst[i]) wei[i] += edge[cur][i];
   return mx;
 }
 int solve(int n) {
    int res = INF;
    for (int x, y; n > 1; n--){
     res = min(res, search(x, y, n));
     REP edge[i][x] = (edge[x][i] += edge[y][i]);
     REP {
```

```
edge[y][i] = edge[n - 1][i];
         edge[i][y] = edge[i][n - 1];
      } // edge[y][y] = 0;
    return res;
|} sw;
```

```
3.7 BoundedFlow(Dinic)
struct BoundedFlow { // O-base
  struct edge {
    int to, cap, flow, rev;
  vector<edge> G[N];
  int n, s, t, dis[N], cur[N], cnt[N];
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n + 2; ++i)</pre>
      G[i].clear(), cnt[i] = 0;
  void add_edge(int u, int v, int lcap, int rcap) {
    cnt[u] -= lcap, cnt[v] += lcap;
    G[u].emplace_back
        (edge{v, rcap, lcap, (int)G[v].size()});
    G[v].emplace_back
        (edge{u, 0, 0, (int)G[u].size() - 1});
  void add_edge(int u, int v, int cap) {
    G[u].emplace_back
        (edge{v, cap, 0, (int)G[v].size()});
    G[v].emplace_back
        (edge{u, 0, 0, (int)G[u].size() - 1});
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < G[u].size(); ++i) {</pre>
      edge &e = G[v][i];
      if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
        int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df, G[e.to][e.rev].flow -= df;
          return df;
     }
    dis[u] = -1;
    return 0;
  }
  bool bfs() {
    fill_n(dis, n + 3, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (edge &e : G[u])
        if (!~dis[e.to] && e.flow != e.cap)
          q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
    int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
    int sum = 0;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
        add_edge(n + 1, i, cnt[i]), sum += cnt[i];
      else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);
```

```
if (sum != maxflow(n + 1, n + 2)) sum = -1;
for (int i = 0; i < n; ++i)
    if (cnt[i] > 0)
        G[n + 1].pop_back(), G[i].pop_back();
    else if (cnt[i] < 0)
        G[i].pop_back(), G[n + 2].pop_back();
    return sum != -1;
}
int solve(int _s, int _t) {
    add_edge(_t, _s, INF);
    if (!solve()) return -1; // invalid flow
    int x = G[_t].back().flow;
    return G[_t].pop_back(), G[_s].pop_back(), x;
}
};</pre>
```

3.8 GomoryHuTree

```
| MaxFlow Dinic;
| int g[N];
| void GomoryHu(int n) { // O-base
| fill_n(g, n, 0);
| for (int i = 1; i < n; ++i) {
| Dinic.reset();
| add_edge(i, g[i], Dinic.maxflow(i, g[i]));
| for (int j = i + 1; j <= n; ++j)
| if (g[j] == g[i] && ~Dinic.dis[j])
| g[j] = i;
| }
| }
```

3.9 MinCostCirculation

```
struct MinCostCirculation { // O-base
  struct Edge {
    tl from, to, cap, fcap, flow, cost, rev;
  } *past[N];
  vector<Edge> G[N];
 ll dis[N], inq[N], n;
  void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, Edge *e) {
      if (dis[v] > d) {
        dis[v] = d, past[v] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
      }
   };
    relax(s, 0, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[v] = 0;
      for (auto &e : G[u])
        if (e.cap > e.flow)
          relax(e.to, dis[u] + e.cost, &e);
  }
 void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
    if (dis[cur.from] + cur.cost < 0) {</pre>
      ++cur.flow, --G[cur.to][cur.rev].flow;
      for (int
           i = cur.from; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        ++e.flow, --G[e.to][e.rev].flow;
      }
    ++cur.cap;
 }
  void solve(int mxlg) {
    for (int b = mxlg; b >= 0; --b) {
      for (int i = 0; i < n; ++i)</pre>
        for (auto &e : G[i])
          e.cap *= 2, e.flow *= 2;
      for (int i = 0; i < n; ++i)</pre>
        for (auto &e : G[i])
```

3.10 FlowModelsBuilding

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \rightarrow y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - of outgoing lower bounds. 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - there's no solution. Otherwise, f' is the answer. 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
- 1. Redirect every edge: $y \rightarrow x$ if $(x,y) \in M$, $x \rightarrow y$ otherwise.
- 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x,y,c), connect $x \rightarrow y$ with (cost,cap) = (c,1) if c>0, otherwise connect $y \rightarrow x$ with (cost,cap) = (-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost,cap)=(0,d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) = (0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
 - 3. Connect source $s \rightarrow v$, $v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K + 2T (\sum_{e \in E(v)} w(e)) 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v .
 - 3. Find the minimum weight perfect matching on G^\prime .
- Project selection problem
 - 1. If $p_v\!>\!0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

```
\begin{split} \min & \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \ge -c_{uv} &\Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{split}
```

4 Data Struture 4.1 LCT

```
#define ls(x) Tree[x].son[0]
#define rs(x) Tree[x].son[1]
#define fa(x) Tree[x].fa
const int maxn = 600010;
struct node {
  int son[2], Min, id, fa, lazy;
} Tree[maxn];
int n, m, q, w[maxn], Min;
struct Node {
 int U, V, W;
} a[maxn];
inline bool IsRoot(int x) {
 return (ls(fa(x)) == x \mid\mid rs(fa(x)) == x) ? false
inline void PushUp(int x) {
  Tree[x].Min = w[x], Tree[x].id = x;
  if (ls(x) && Tree[ls(x)].Min < Tree[x].Min) {</pre>
    Tree[x].Min = Tree[ls(x)].Min;
    Tree[x].id = Tree[ls(x)].id;
 }
 if (rs(x) && Tree[rs(x)].Min < Tree[x].Min) {</pre>
    Tree[x].Min = Tree[rs(x)].Min;
    Tree[x].id = Tree[rs(x)].id;
 }
inline void Update(int x) {
 Tree[x].lazy ^= 1;
  swap(ls(x), rs(x));
inline void PushDown(int x) {
 if (!Tree[x].lazy) return;
  if (ls(x)) Update(ls(x));
  if (rs(x)) Update(rs(x));
  Tree[x].lazy = 0;
inline void Rotate(int x) {
 int y = fa(x), z = fa(y), k = rs(y) == x,
      w = Tree[x].son[!k];
  if (!IsRoot(y)) Tree[z].son[rs(z) == y] = x;
 fa(x) = z, fa(y) = x;
  if (w) fa(w) = y;
 Tree[x].son[!k] = y, Tree[y].son[k] = w;
 PushUp(y);
inline void Splay(int x) {
  stack<int> Stack;
  int y = x, z;
 Stack.push(y);
 while (!IsRoot(y)) Stack.push(y = fa(y));
 while (!Stack.empty())
    PushDown(Stack.top()), Stack.pop();
 while (!IsRoot(x)) {
    y = fa(x), z = fa(y);
    if (!IsRoot(y))
      Rotate((ls(y) == x) ^(ls(z) == y) ? x : y);
    Rotate(x):
 }
 PushUp(x);
}
inline void Access(int root) {
  for (int x = 0; root; x = root, root = fa(root))
    Splay(root), rs(root) = x, PushUp(root);
inline void MakeRoot(int x) {
 Access(x), Splay(x), Update(x);
```

```
| inline int FindRoot(int x) {
  Access(x), Splay(x);
  while (ls(x)) x = ls(x);
  return Splay(x), x;
|}
inline void Link(int u, int v) {
  MakeRoot(u);
  if (FindRoot(v) != u) fa(u) = v;
inline void Cut(int u, int v) {
  MakeRoot(u);
  if (FindRoot(v) != u || fa(v) != u || ls(v)) return;
  fa(v) = rs(u) = 0;
}
inline void Split(int u, int v) {
  MakeRoot(u), Access(v), Splay(v);
}
inline bool Check(int u, int v) {
  return MakeRoot(u), FindRoot(v) == u;
inline int LCA(int root, int u, int v) {
  MakeRoot(root), Access(u), Access(v), Splay(u);
  if (!fa(u)) {
    Access(u), Splay(v);
    return fa(v);
  }
  return fa(u);
}
/* ETT
每次進入節點和走邊都放入一次共 3n - 2
node(u) 表示進入節點 u 放入 treap 的位置
edge(u, v) 表示 u -> v 的邊放入 treap 的位置 (push v)
Makeroot u
  L1 = [begin, node(u) - 1], L2 = [node(u), end]
  -> L2 + L1
Insert u, v:
  Tu \rightarrow L1 = [begin, node(u) - 1], L2 = [node(u), end]
  Tv -> L3 = [begin, node(v) - 1], L4 = [node(v), end]
  -> L2 + L1 + edge(u, v) + L4 + L3 + edge(v, u)
Delect u, v .
  maybe need swap u, v
  T -> L1 + edge(u, v) + L2 + edge(v, u) + L3
  -> L1 + L3, L2
     String
5
5.1 KMP
int KMP(string s, string t) {
  t = " "s + t; // consistency with ACa
  int n = t.size(), ans = \theta;
  vector<int> f(t.size(), 0);
  f[0] = -1;
  for (int i = 1, j = -1; i < t.size(); i++) {</pre>
    while (j >= 0 && t[j + 1] != t[i])
      j = f[j];
    f[i] = ++j;
  for (int i = 0, j = 0; i < s.size(); i++) {</pre>
    while (j \ge 0 \&\& t[j + 1] != s[i])
      j = f[j];
    if (++j + 1 == t.size()) ans++, j = f[j];
  return ans;
}
```

5.2 Z

int Z[1000006];

while (

void z(string s) {

if (i < Z[mx] + mx)

for (int i = 1, mx = 0; i < s.size(); i++) {</pre>

Z[i] = min(Z[mx] - i + mx, Z[i - mx]);

man[i] = min(man[mx] + mx - i, man[2 * mx - i]);

while (man[i] + i < t.size() && i - man[i] >= 0 &&

for (int i = 0; i < t.size(); i++) {</pre>

t[i + man[i]] == t[i - man[i]])

for (int i = 0; i < t.size(); i++)</pre>

ans = max(ans, man[i] - 1);

if (i + man[i] > mx + man[mx]) mx = i;

5.4 SuffixArray

man[i] = 1;

man[i]++;

return ans;

|}

```
vector<int> sa, cnt, rk, tmp, lcp;
void SA(string s) {
  int n = s.size();
  sa.resize(n), cnt.resize(n), rk.resize(n),
     tmp.resize(n);
  iota(all(sa), 0);
  sort(all(sa),
     [&](int i, int j) { return s[i] < s[j]; });
  rk[0] = 0;
  for (int i = 1; i < n; i++)</pre>
     rk[sa[i]] :
       rk[sa[i - 1]] + (s[sa[i - 1]] != s[sa[i]]);
  for (int k = 1; k <= n; k <<= 1) {</pre>
     fill(all(cnt), 0);
     for (int i = 0; i < n; i++)</pre>
       cnt[rk[(sa[i] - k + n) % n]]++;
     for (int i = 1; i < n; i++) cnt[i] += cnt[i - 1];</pre>
     for (int i = n - 1; i >= 0; i--)
       tmp[--cnt[rk[(sa[i] - k + n) % n]]] =
         (sa[i] - k + n) % n;
     sa.swap(tmp);
     tmp[sa[0]] = 0;
    for (int i = 1; i < n; i++)</pre>
       tmp[sa[i]] = tmp[sa[i - 1]] +
         (rk[sa[i - 1]] != rk[sa[i]] ||
           rk[(sa[i - 1] + k) % n] !=
             rk[(sa[i] + k) % n]);
     rk.swap(tmp);
  }
void LCP(string s) {
  int n = s.size(), k = 0;
  lcp.resize(n);
  for (int i = 0; i < n; i++)</pre>
     if (rk[i] == 0) lcp[rk[i]] = 0;
     else {
       if (k) k--;
       int j = sa[rk[i] - 1];
       while (
         i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
       lcp[rk[i]] = k;
| }
```

5.5 SAIS

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
// zero based, string content MUST > 0
// SA[i]: SA[i]-th
     suffix is the i-th lexigraphically smallest suffix.
// H[i]: longest
    common prefix of suffix SA[i] and suffix SA[i - 1].
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, 0), copy_n(c, z, x); }
void induce
     (int *sa, int *c, int *s, bool *t, int n, int z) {
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i)</pre>
    if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
  for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
       sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
}
void sais(int *s, int *sa
     , int *p, int *q, bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0,
      nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
  for (int i = n - 2; i >= 0; --i)
    t[i] = (
         s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
  for (int i = 0; i < n; ++i)</pre>
    if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
      bool neg = last < 0 || !equal
           (s + sa[i], s + p[q[sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
    }
  sais(ns,
       nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
|}
void mkhei(int n) {
  for (int i = 0, j = 0; i < n; ++i) {
    if (RA[i])
      for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
    H[RA[i]] = j, j = max(0, j - 1);
  }
}
void build(int *s, int n) {
  copy_n(s, n, _s), _s[n] = 0;
  sais(_s, SA, _p, _q, _t, _c, _n + 1, 256); copy_n(SA + 1, _n, SA);
  for (int i = 0; i < n; ++i) RA[SA[i]] = i;</pre>
  mkhei(n);
5.6 ACAutomaton
```

```
#define sumS 500005
#define sigma 26
#define base 'a'
```

```
struct AhoCorasick {
  int ch[sumS][sigma] = {{}}, f[sumS] = {-1},
      tag[sumS], mv[sumS][sigma], jump[sumS],
      cnt[sumS];
  int idx = 0, t = -1;
  vector<int> E[sumS], q;
  pii o[sumS];
  int insert(string &s, int t) {
    int j = 0;
    for (int i = 0; i < (int)s.size(); i++) {</pre>
      if (!ch[j][s[i] - base])
        ch[j][s[i] - base] = ++idx;
      j = ch[j][s[i] - base];
    tag[j] = 1;
    return j;
  int next(int u, int c) {
    return u < 0 ? 0 : mv[u][c];</pre>
  }
  void dfs(int u) {
    o[u].F = ++t;
    for (auto v : E[u]) dfs(v);
    o[u].S = t;
  void build() {
    int k = -1;
    q.emplace_back(0);
    while (++k < q.size()) {</pre>
      int u = q[k];
      for (int v = 0; v < sigma; v++) {</pre>
        if (ch[u][v]) {
          f[ch[u][v]] = next(f[u], v);
           q.emplace_back(ch[u][v]);
        }
        mv[v][v] =
           (ch[u][v] ? ch[u][v] : next(f[u], v));
      }
      if (u) jump[u] = (tag[f[u]] ? f[u] : jump[f[u]]);
    reverse(q.begin(), q.end());
    for (int i = 1; i <= idx; i++)</pre>
      E[f[i]].emplace_back(i);
    dfs(0);
  void match(string &s) {
    fill(cnt, cnt + idx + 1, 0);
    for (int i = 0, j = 0; i < (int)s.size(); i++)</pre>
      cnt[j = next(j, s[i] - base)]++;
    for (int i : q)
      if (f[i] > 0) cnt[f[i]] += cnt[i];
  }
} ac;
5.7 MinRotation
int mincyc(string s) {
  int n = s.size();
                                                            };
  int i = 0, ans = 0;
  while (i < n) {
    ans = i;
int j = i + 1, k = i;
    while (j < s.size() && s[j] >= s[k]) {
      k = (s[j] > s[k] ? i : k + 1);
      ++j;
```

5.8 ExtSAM

return ans:

while (i <= k) i += j - k;</pre>

```
| #define CNUM 26
| struct exSAM {
```

```
int len[N * 2], link[N * 2]; // maxlength, suflink
int next[N * 2][CNUM], tot; // [0, tot), root = 0
int lenSorted[N * 2]; // topo. order
int cnt[N * 2]; // occurence
int newnode() {
  fill_n(next[tot], CNUM, 0);
  len[tot] = cnt[tot] = link[tot] = 0;
  return tot++;
void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
  int cur = next[last][c];
  len[cur] = len[last] + 1;
  int p = link[last];
  while (p != -1 && !next[p][c])
    next[p][c] = cur, p = link[p];
  if (p == -1) return link[cur] = 0, cur;
  int q = next[p][c];
  if (len
      [p] + 1 == len[q]) return link[cur] = q, cur;
  int clone = newnode();
  for (int i = 0; i < CNUM; ++i)</pre>
        clone][i] = len[next[q][i]] ? next[q][i] : 0;
  len[clone] = len[p] + 1;
  while (p != -1 && next[p][c] == q)
    next[p][c] = clone, p = link[p];
  link[link[cur] = clone] = link[q];
  link[q] = clone;
  return cur;
void insert(const string &s) {
  int cur = 0;
  for (auto ch : s) {
    int &nxt = next[cur][int(ch - 'a')];
    if (!nxt) nxt = newnode();
    cnt[cur = nxt] += 1;
void build() {
  queue<int> q;
  q.push(0);
  while (!q.empty()) {
    int cur = q.front();
    q.pop();
    for (int i = 0; i < CNUM; ++i)</pre>
      if (next[cur][i])
        q.push(insertSAM(cur, i));
  }
  vector<int> lc(tot);
  for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
  partial_sum(all(lc), lc.begin());
  for (int i
      = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
void solve() {
  for (int i = tot - 2; i >= 0; --i)
    cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
}
```

5.9 PalindromeTree

```
St[0].fail = 1, St[1].len = -1, s.emplace_back(-1);
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.emplace_back(0), St.emplace_back(-1);
    St[0].fail = 1, s.emplace_back(-1);
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
      x = St[x].fail;
    return x;
  }
  inline void add(int c) {
    s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = St.size();
      St.emplace_back(St[cur].len + 2);
      St[now].fail =
         St[get_fail(St[cur].fail)].next[c];
       St[cur].next[c] = now;
      St[now].num = St[St[now].fail].num + 1;
    last = St[cur].next[c], ++St[last].cnt;
  }
  inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
      St[i->fail].cnt += i->cnt;
  }
  inline int size() { // The number of diff. pal.
    return (int)St.size() - 2;
|};
```

6 Number Theory

6.1 Primes

6.2 ExtGCD

```
|// beware of negative numbers!
|void extgcd(ll a, ll b, ll c, ll &x, ll &y) {
| if (b == 0) x = c / a, y = 0;
| else {
| extgcd(b, a % b, c, y, x);
| y -= x * (a / b);
| }
|} // |x| <= b/2, |y| <= a/2</pre>
```

6.3 FloorCeil

```
| int floor(int a, int b)
|{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
| int ceil(int a, int b)
|{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

6.4 FloorSum

```
| Il floorsum(Il A, Il B, Il C, Il N) {
| if (A == 0) return (N + 1) * (B / C);
| if (A > C || B > C)
| return (N + 1) * (B / C) +
| N * (N + 1) / 2 * (A / C) +
| floorsum(A % C, B % C, C, N);
| Il M = (A * N + B) / C;
| return N * M - floorsum(C, C - B - 1, A, M - 1);
| } // \sum^{n}_0 floor((ai + b) / m)
```

6.5 MillerRabin

```
| / / n < 4,759,123,141
                           3 : 2, 7, 61
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383 6 : primes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool Miller_Rabin(ll a, ll n) {
  if ((a = a % n) == 0) return 1;
   if (n % 2 == 0) return n == 2;
  ll tmp = (n - 1) / ((n - 1) & (1 - n));
  ll t = _{-}lg(((n - 1) & (1 - n))), x = 1;
   for (; tmp; tmp >>= 1, a = mul(a, a, n))
    if (tmp \& 1) x = mul(x, a, n);
   if (x == 1 || x == n - 1) return 1;
   while (--t)
     if ((x = mul(x, x, n)) == n - 1) return 1;
   return 0;
|}
```

6.6 PollardRho

```
|map<ll, int> cnt;
void PollardRho(ll n) {
   if (n == 1) return;
   if (prime(n)) return ++cnt[n], void();
   if (n % 2
       == 0) return PollardRho(n / 2), ++cnt[2], void();
   11 x = 2, y = 2, d = 1, p = 1;
   #define f(x, n, p) ((mul(x, x, n) + p) % n)
   while (true) {
     if (d != n && d != 1) {
       PollardRho(n / d);
       PollardRho(d);
       return;
     }
     if (d == n) ++p;
     x = f(x, n, p), y = f(f(y, n, p), n, p);
     d = gcd(abs(x - y), n);
|}
```

6.7 Fraction

```
struct fraction {
   ll n, d;
   fraction
       (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
     ll t = gcd(n, d);
    n /= t, d /= t;
     if (d < 0) n = -n, d = -d;
  fraction operator-() const
   { return fraction(-n, d); }
   fraction operator+(const fraction &b) const
   { return fraction(n * b.d + b.n * d, d * b.d); }
   fraction operator-(const fraction &b) const
   { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
   { return fraction(n * b.n, d * b.d); }
   fraction operator/(const fraction &b) const
   { return fraction(n * b.d, d * b.n); }
   void print() {
     cout << n;
     if (d != 1) cout << "/" << d;
|};
```

6.8 ChineseRemainder

```
|ll solve(ll x1, ll m1, ll x2, ll m2) {
    ll g = gcd(m1, m2);
    if ((x2 - x1) % g) return -1; // no sol
    m1 /= g; m2 /= g;
    ll x, y;
    extgcd(m1, m2, __gcd(m1, m2), x, y);
    ll lcm = m1 * m2 * g;
    ll res = x * (x2 - x1) * m1 + x1;
```

```
// be careful with overflow
  return (res % lcm + lcm) % lcm;
| }
6.9 Factorial\mathsf{Mod}p^k
// O(p^k + \log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)</pre>
     if (i % p) prod[i] = prod[i - 1] * i % pk;
  else prod[i] = prod[i - 1];
ll rt = 1;
  for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
  }
  return rt;
|} // (n! without factor p) % p^k
6.10 QuadraticResidue
```

```
// Berlekamp-Rabin, log^2(p)
ll trial(ll y, ll z, ll m) {
  ll a0 = 1, a1 = 0, b0 = z, b1 = 1, p = (m - 1) / 2;
  while (p) {
    if (p & 1)
      tie(a0, a1) =
        make_pair((a1 * b1 % m * y + a0 * b0) % m,
           (a0 * b1 + a1 * b0) % m);
    tie(b0, b1) =
      make_pair((b1 * b1 % m * y + b0 * b0) % m,
        (2 * b0 * b1) % m);
    p >>= 1:
  if (a1) return inv(a1, m);
  return -1;
mt19937 rd(49);
ll psqrt(ll y, ll p) {
  if (fpow(y, (p - 1) / 2, p) != 1) return -1;
  for (int i = 0; i < 30; i++) {</pre>
    11 z = rd() \% p;
    if (z * z % p == y) return z;
    ll x = trial(y, z, p);
    if (x == -1) continue;
    return x:
  }
  return -1;
| }
```

6.11 MeisselLehmer

```
ll PrimeCount(ll n) { // n \sim 10^13 \Rightarrow < 2s
  if (n <= 1) return 0;
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector<int> smalls(v + 1), skip(v + 1), roughs(s);
  vector<ll> larges(s);
  for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
  for (int i = 0; i < s; ++i) {
    roughs[i] = 2 * i + 1;
    larges[i] = (n / (2 * i + 1) + 1) / 2;
  }
  for (int p = 3; p <= v; ++p) {</pre>
    if (smalls[p] > smalls[p - 1]) {
      int q = p * p;
      ++pc;
      if (1LL * q * q > n) break;
      skip[p] = 1;
      for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
      int ns = 0;
      for (int k = 0; k < s; ++k) {</pre>
        int i = roughs[k];
        if (skip[i]) continue;
        ll d = 1LL * i * p;
        larges[ns] = larges[k] - (d <= v ? larges
             [smalls[d] - pc] : smalls[n / d]) + pc;
```

```
roughs[ns++] = i;
       s = ns:
       for (int j = v / p; j >= p; --j) {
        int c =
              smalls[j] - pc, e = min(j * p + p, v + 1);
         for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
      }
    }
  }
  for (int k = 1; k < s; ++k) {
    const ll m = n / roughs[k];
    ll t = larges[k] - (pc + k - 1);
     for (int l = 1; l < k; ++l) {</pre>
       int p = roughs[l];
       if (1LL * p * p > m) break;
       t -= smalls[m / p] - (pc + l - 1);
    larges[0] -= t;
  return larges[0];
}
```

6.12 DiscreteLog

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
   unordered_map<int, int> p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {</pre>
     p[y] = i;
     y = 1LL * y * x % m;
     b = 1LL * b * x % m;
   }
   for (int i = 0; i < m + 10; i += kStep) {</pre>
     s = 1LL * s * b % m;
     if (p.find(s) != p.end()) return i + kStep - p[s];
   return -1;
}
int DiscreteLog(int x, int y, int m) {
   if (m == 1) return 0;
   int s = 1;
   for (int i = 0; i < 100; ++i) {
     if (s == y) return i;
     s = 1LL * s * x % m;
  if (s == y) return 100;
   int p = 100 + DiscreteLog(s, x, y, m);
   if (fpow(x, p, m) != y) return -1;
   return p;
}
```

6.13 Theorems

• Cramer's rule

$$ax+by=e cx+dy=f y = af - ec d-bc$$

Vandermonde's Identity

$$C(n\!+\!m,\!k)\!=\!\sum_{i=0}^{k}\!C(n,\!i)C(m,\!k\!-\!i)$$

• Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}\!=\!d(i)$, $L_{ij}\!=\!-c$ where c is the number of edge $(i,\!j)$ in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence $d_1,d_2,...,d_n$ for each labeled vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.

- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős-Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+\cdots+d_n$ is even and

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k) \text{ holds for every } 1 \leq k \leq n \text{.}$$

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \! \geq \! \cdots \! \geq \! a_n$ and $b_1, ..., b_n$ is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$$
 holds for every $1 \leq k \leq n$.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),...,(a_n,b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$

and
$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n \text{.}$$
 Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}$ + #{lattice points on the boundary} -1.

- Möbius inversion formula

 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$ $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- · Spherical cap
 - A portion of a sphere cut off by a plane.
 - r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : $\arcsin(a/r)$.
 - Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos\theta) (1 \cos\theta)^2/3$.
 - Area $=2\pi rh=\pi(a^2+h^2)=2\pi r^2(1-\cos\theta)$.
- Lagrange multiplier
 - Optimize $f(x_1,...,x_n)$ when k constraints $g_i(x_1,...,x_n)=0$.
 - Lagrangian function $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=f(x_1,\ldots,x_n)$
 - $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n)$. The solution corresponding to the original control of the solution corresponding to the original control of the solution o strained optimization is always a saddle point of the Lagrangian function.
- · Nearest points of two skew lines
 - Line 1: ${m v}_1\!=\!{m p}_1\!+\!t_1{m d}_1$
 - Line 2: $v_2 = p_2 + t_2 d_2$
 - $n = d_1 \times d_2$
 - $\boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}$
 - $\boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$
 - $c_1 = p_1 + \frac{(p_2 p_1) \cdot n_2}{d_1 \cdot n_2} d_1$
 - $c_2 = p_2 + \frac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2$

6.14 Estimation

- Estimation
 - The number of divisors of n is at most around 100 for $n\!<\!$ 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
 - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for $n=0\sim 9$, 627 for n=20, $\sim 2e5$ for $n\!=\!50$, $\sim\!2e8$ for $n\!=\!100$.
 - Total number of partitions of n distinct elements: B(n) = 1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437,190899322,....

6.15 EuclideanAlgorithms

- m = | ^{an+b}/_{an+b} |
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \mod c,b \mod c,c,n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ &= \begin{pmatrix} \frac{1}{2} \cdot (n(n+1)m - f(c,c-b-1,a,m-1) \\ -h(c,c-b-1,a,m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) = & \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2 \\ = & \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \text{ mod } c,b \text{ mod } c,c,n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \text{ mod } c,b \text{ mod } c,c,n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \text{ mod } c,b \text{ mod } c,c,n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c,c-b-1,a,m-1) \\ -2f(c,c-b-1,a,m-1) - f(a,b,c,n), & \text{otherwise} \end{cases} \end{split}$$

6.16 Numbers

Bernoulli numbers

$$\begin{split} &B_0-1, B_1^{\pm}=\pm \tfrac{1}{2}, B_2=\tfrac{1}{6}, B_3=0\\ &\sum_{j=0}^m \binom{m+1}{j} B_j=0\text{, EGF is } B(x)=\tfrac{x}{e^x-1}=\sum_{n=0}^\infty B_n \frac{x^n}{n!}\text{.}\\ &S_m(n)=\sum_{k=1}^n k^m=\frac{1}{m+1}\sum_{k=0}^m \binom{m+1}{k} B_k^+ n^{m+1-k} \end{split}$$

ullet Stirling numbers of the second kind Partitions of ndistinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem
$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
 • Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C_n^{(k)}(x) = \frac{1}{(k-1)n+1} (x^{k(k)}(x))^{k}$$

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. $k \ j \colon \mathbf{s} \ \mathbf{s} \cdot \mathbf{t}$ $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \ge j$, k j:s s.t. $\pi(j) > j$. E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

6.17 GeneratingFunctions

- Ordinary Generating Function $A(x) = \sum_{i \geq 0} a_i x^i$
 - $A(rx) \Rightarrow r^n a_n$

 - $A(x)+B(x) \Rightarrow a_n+b_n$ $A(x)B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\cdots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x)' \Rightarrow na_n$
 - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$
 - $A(x)+B(x) \Rightarrow a_n+b_n$

 - $A^{(k)}(x) \Rightarrow a_{n+k}$ $A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} {n \choose i_1, i_2, \dots, i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x) \Rightarrow na_n$
- Special Generating Function
 - $(1+x)^n = \sum_{i \ge 0} \binom{n}{i} x^i$
 - $-\frac{1}{(1-x)^n} = \sum_{i>0} {i \choose n-1} x^i$

```
- S_k = \sum_{x=1}^n x^k: S = \sum_{p=0}^\infty x^p = \frac{e^x - e^{x(n+1)}}{1 - e^x}
```

7 Linear Algebra

7.1 GuassianElimination

```
#undef M
 struct matrix { //m variables, n equations
  int n, m;
   fraction M[N][N + 1], sol[N];
   int solve() { //-1: inconsistent, >= 0: rank
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
       for (int j = 0; j < n; ++j) {</pre>
         if (i == j) continue;
         fraction tmp = -M[j][piv] / M[i][piv];
         for (int k = 0; k <=</pre>
               m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
       }
     }
     int rank = 0;
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m && M[i][m].n) return -1;
       else if (piv
             < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
     return rank;
  }
|};
```

7.2 BerlekampMassey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(output.size() + 1), me, he;
  for (int f = 0, i = 1; i <= output.size(); ++i) {</pre>
    for (int j = 0; j < me.size(); ++j)</pre>
      d[i] += output[i - j - 2] * me[j];
    if ((d[i] -= output[i - 1]) == 0) continue;
    if (me.empty()) {
      me.resize(f = i);
       continue;
    }
    vector<T> o(i - f - 1);
    T k = -d[i] / d[f];
    o.emplace_back(-k);
    for (T x : he) o.emplace_back(x * k);
    o.resize(max(o.size(), me.size()));
    for (int j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
    if (i - f + (int
         )he.size()) >= (int)me.size()) he = me, f = i;
  return me;
| }
```

7.3 Simplex

```
Standard form: maximize \mathbf{c}^T\mathbf{x} subject to A\mathbf{x} \leq \mathbf{b} and \mathbf{x} \geq 0. Dual LP: minimize \mathbf{b}^T\mathbf{y} subject to A^T\mathbf{y} \geq \mathbf{c} and \mathbf{y} \geq \mathbf{0}. \bar{\mathbf{x}} and \bar{\mathbf{y}} are optimal if and only if for all i \in [1,n], either \bar{x}_i = 0 or \sum_{j=1}^m A_{ji}\bar{x}_j = c_i holds and for all i \in [1,m] either \bar{y}_i = 0 or \sum_{j=1}^n A_{ij}\bar{x}_j = b_j holds.

1. In case of minimization, let c_i' = -c_i
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j
• \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j
• \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
4. If x_i has no lower bound, replace x_i with x_i - x_i'
// using N + 2M variables const int mxM = 25; const int mxN = 25 + 2 * mxM; struct simplex {
```

const double inf = 1 / .0, eps = 1e-9;

int n, m, k, var[mxN], inv[mxN], art[mxN];

```
double A[mxM][mxN], B[mxM], x[mxN];
   void init(int _n) { n = _n, m = 0; }
   void equation(vector<double> a, double b) {
     for (int i = 0; i < n; i++) A[m][i] = a[i];</pre>
     B[m] = b, var[m] = n + m, ++m;
   void pivot(int r, int c, double bx) {
     for (int i = 0; i <= m + 1; i++)
       if (i != r && abs(A[i][c]) > eps) {
         x[var[i]] -= bx * A[i][c] / A[i][var[i]];
         double f = A[i][c] / A[r][c];
         for (int j = 0; j <= n + m + k; j++)</pre>
           A[i][j] -= A[r][j] * f;
         B[i] -= B[r] * f;
  }
   double phase(int p) {
     while (true) {
       int in = min_element(
                  A[m + p], A[m + p] + n + m + k + 1) -
         A[m + p];
       if (A[m + p][in] >= -eps) break;
       double bx = inf;
       int piv = -1;
       for (int i = 0; i < m; i++)</pre>
         if (A[i][in] > eps && B[i] / A[i][in] <= bx)</pre>
           piv = i, bx = B[i] / A[i][in];
       if (piv == -1) return inf;
       int out = var[piv];
       pivot(piv, in, bx);
       x[out] = 0, x[in] = bx, var[piv] = in;
     }
     return x[n + m];
   double solve(vector<double> c) {
     auto invert = [&](int r) {
       for (int j = 0; j <= n + m; j++) A[r][j] *= -1;</pre>
       B[r] *= -1;
     };
     k = 1:
     for (int i = 0; i < n; i++) A[m][i] = -c[i];</pre>
     fill(A[m + 1], A[m + 1] + mxN, 0.0);
     for (int i = 0; i <= m + 1; i++)</pre>
       fill(A[i] + n, A[i] + n + m + 2, 0.0),
         var[i] = n + i, A[i][n + i] = 1;
     for (int i = 0; i < m; i++) {
       if (B[i] < 0) {</pre>
         ++k;
         for (int j = 0; j <= n + m; j++)</pre>
           A[m + 1][j] += A[i][j];
         var[i] = n + m + k, A[i][var[i]] = 1,
         art[var[i]] = n + i;
       x[var[i]] = B[i];
     phase(1):
     if (*max_element(
           x + (n + m + 2), x + (n + m + k + 1)) > eps)
       return .0 / .0;
     for (int i = 0; i <= m; i++)</pre>
       if (var[i] > n + m)
         var[i] = art[var[i]], invert(i);
     k = 0;
     return phase(0);
|} lp;
```

8 Polynomials 8.1 NTT (FFT)

```
|// 9223372036737335297, 3
|#define base ll // complex<double>
|#define N 524288
```

| }

```
// const double PI = acosl(-1);
                                                             const int N = 21;
const ll mod = 998244353, g = 3;
                                                             int f[
base omega[4 * N], omega_[4 * N];
                                                                 N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
int rev[4 * N];
                                                             void
                                                                 subset_convolution(int *a, int *b, int *c, int L) {
ll fpow(ll b, ll p);
                                                               // c_k = \sum_{i = 0} a_i * b_j
                                                               int n = 1 << L;
ll inverse(ll a) { return fpow(a, mod - 2); }
                                                               for (int i = 1; i < n; ++i)</pre>
                                                                 ct[i] = ct[i \& (i - 1)] + 1;
void calcW(int n) {
                                                               for (int i = 0; i < n; ++i)</pre>
  ll r = fpow(g, (mod - 1) / n), invr = inverse(r);
                                                                 f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  omega[0] = omega_[0] = 1;
                                                               for (int i = 0; i <= L; ++i)</pre>
  for (int i = 1; i < n; i++) {</pre>
                                                                 fwt(f[i], n, 1), fwt(g[i], n, 1);
    omega[i] = omega[i - 1] * r % mod;
                                                               for (int i = 0; i <= L; ++i)
    omega_[i] = omega_[i - 1] * invr % mod;
                                                                 for (int j = 0; j <= i; ++j)</pre>
                                                                   for (int x = 0; x < n; ++x)
  // double arg = 2.0 * PI / n;
                                                                     h[i][x] += f[j][x] * g[i - j][x];
  // for (int i = 0; i < n; i++)
                                                               for (int i = 0; i <= L; ++i)</pre>
  // {
                                                                 fwt(h[i], n, -1);
  //
      omega[i] = base(cos(i * arg), sin(i * arg));
                                                               for (int i = 0; i < n; ++i)</pre>
  //
       omega_[i] = base(cos(-i * arg), sin(-i * arg));
                                                                 c[i] = h[ct[i]][i];
  // }
                                                            |}
}
                                                             8.3 PolynomialOperations
void calcrev(int n) {
                                                             #define poly vector<ll>
  int k = __lg(n);
                                                             poly inv(poly A) {
  for (int i = 0; i < n; i++) rev[i] = 0;</pre>
                                                               A.resize(1 << (__lg(A.size() - 1) + 1));
  for (int i = 0; i < n; i++)</pre>
                                                               poly B = {inverse(A[0])};
    for (int j = 0; j < k; j++)
                                                               for (int n = 1; n < A.size(); n += n) {</pre>
      if (i & (1 << j)) rev[i] ^= 1 << (k - j - 1);</pre>
                                                                 poly pA(A.begin(), A.begin() + 2 * n);
                                                                 calcrev(4 * n);
                                                                 calcW(4 * n);
vector<base> NTT(vector<base> poly, bool inv) {
                                                                 pA.resize(4 * n);
  base *w = (inv ? omega_ : omega);
                                                                 B.resize(4 * n);
  int n = poly.size();
                                                                 pA = NTT(pA, 0);
  for (int i = 0; i < n; i++)</pre>
                                                                 B = NTT(B, 0);
    if (rev[i] > i) swap(poly[i], poly[rev[i]]);
                                                                 for (int i = 0; i < 4 * n; i++)
                                                                   B[i] =
  for (int len = 1; len < n; len <<= 1) {</pre>
                                                                     ((B[i] * 2 - pA[i] * B[i] % mod * B[i]) % mod +
    int arg = n / len / 2;
    for (int i = 0; i < n; i += 2 * len)</pre>
                                                                     mod:
      for (int j = 0; j < len; j++) {</pre>
                                                                 B = NTT(B, 1);
         base odd =
           w[j * arg] * poly[i + j + len] % mod;
                                                                 B.resize(2 * n);
         poly[i + j + len] =
           (poly[i + j] - odd + mod) % mod;
                                                               return B;
         poly[i + j] = (poly[i + j] + odd) \% mod;
                                                             pair<poly, poly> div(poly A, poly B) {
  }
                                                               if (A.size() < B.size()) return make_pair(poly(), A);</pre>
  if (inv)
                                                               int n = A.size(), m = B.size();
    for (auto &a : poly) a = a * inverse(n) % mod;
                                                               poly revA = A, invrevB = B;
  return poly;
                                                               reverse(revA.begin(), revA.end());
                                                               reverse(invrevB.begin(), invrevB.end());
                                                               revA.resize(n - m + 1);
vector<base> mul(vector<base> f, vector<base> g) {
                                                               invrevB.resize(n - m + 1);
  int sz = 1 << (__lg(f.size() + g.size() - 1) + 1);</pre>
                                                               invrevB = inv(invrevB);
  f.resize(sz), g.resize(sz);
  calcrev(sz);
                                                               poly Q = mul(revA, invrevB);
  calcW(sz):
                                                               Q.resize(n - m + 1);
  f = NTT(f, 0), q = NTT(q, 0);
                                                               reverse(Q.begin(), Q.end());
  for (int i = 0; i < sz; i++)</pre>
                                                               poly R = mul(Q, B);
    f[i] = f[i] * g[i] % mod;
                                                               R.resize(m - 1);
  return NTT(f, 1);
                                                               for (int i = 0; i < m - 1; i++)
| }
                                                                 R[i] = (A[i] - R[i] + mod) \% mod;
8.2 FHWT
                                                               return make_pair(Q, R);
                                                             }
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
                                                             ll fast_kitamasa(ll k, poly A, poly C) {
xor: (x, y = (x + y) * op, (x - y) * op)
                                                               int n = A.size();
invop: or, and, xor = -1, -1, 1/2 */
                                                               C.emplace_back(mod - 1);
void fwt(int *a, int n, int op) { //or
                                                               poly Q, R = \{0, 1\}, F = \{1\};
  for (int L = 2; L <= n; L <<= 1)
                                                               R = div(R, C);
    for (int i = 0; i < n; i += L)</pre>
                                                               while (k) {
      for (int j = i; j < i + (L >> 1); ++j)
                                                                 if (k & 1) F = div(mul(F, R), C);
         a[j + (L >> 1)] += a[j] * op;
                                                                 R = div(mul(R, R), C);
```

k >>= 1;

```
ll ans = 0;
  for (int i = 0; i < F.size(); i++)</pre>
    ans = (ans + A[i] * F[i]) % mod;
  return ans;
vector<ll> fpow(vector<ll> f, ll p, ll m) {
  int b = 0;
  while (b < f.size() && f[b] == 0) b++;</pre>
  f = vector<ll>(f.begin() + b, f.end());
  int n = f.size();
  f.emplace_back(0);
  vector<ll> q(min(m, b * p), 0);
  q.emplace_back(fpow(f[0], p));
  for (int k = 0; q.size() < m; k++) {</pre>
    ll res = 0;
    for (int i = 0; i < min(n, k + 1); i++)</pre>
      res = (res +
               p * (i + 1) % mod * f[i + 1] % mod *
                 q[k - i + b * p]) %
         mod;
    for (int i = 1; i < min(n, k + 1); i++)</pre>
      res = (res -
               f[i] * (k - i + 1) % mod *
                 q[k - i + 1 + b * p]) %
    res = (res < 0 ? res + mod : res) *
      inv(f[0] * (k + 1) % mod) % mod;
    q.emplace_back(res);
  }
  return q;
|}
```

8.4 NewtonMethod+MiscGF

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by Q_k the polynomial such that $F(Q_k)\!=\!0\pmod{x^{2^k}}$, then

$$\begin{split} Q_{k+1} \!=\! Q_k \!-\! \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}} \\ \bullet \ A^{-1} \!: \ B_{k+1} \!=\! B_k(2 \!-\! A B_k) \mod x^{2^{k+1}} \\ \bullet \ \ln\! A \!: \ (\ln\! A)' \!=\! \frac{A'}{A} \\ \bullet \ \exp\! A \!: \ B_{k+1} \!=\! \frac{A}{2}(B_k \!+\! A B_k^{-1}) \mod x^{2^{k+1}} \\ \bullet \ \sqrt{A} \!: \ B_{k+1} \!=\! \frac{1}{2}(B_k \!+\! A B_k^{-1}) \mod x^{2^{k+1}} \end{split}$$

9 Geometry

9.1 Basic

```
typedef pair<pdd, pdd> Line;
struct Cir{ pdd 0; double R; };
const double eps = 1e-8;
pdd operator+(pdd a, pdd b)
{ return pdd(a.F + b.S, a.S + b.S); }
pdd operator-(pdd a, pdd b)
{ return pdd(a.F - b.S, a.S - b.S); }
pdd operator*(pdd a, double b)
{ return pdd(a.F * b, a.S * b); }
pdd operator/(pdd a, double b)
{ return pdd(a.F / b, a.S / b); }
double dot(pdd a, pdd b)
{ return a.F * b.F + a.S * b.S; }
double cross(pdd a, pdd b)
{ return a.F * b.S - a.S * b.F; }
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
```

```
{ return sign(cross(b - a, c - a)); }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
  if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= 0;</pre>
}
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
    return btw(p1, p2, p3) || btw(p1, p2, p4) ||
      btw(p3, p4, p1) || btw(p3, p4, p2);
  return a123 * a124 <= 0 && a341 * a342 <= 0;
}
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
  double a124 = cross(p2 - p1, p4 - p1);
  return (p4
      * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
pdd perp(pdd p1)
{ return pdd(-p1.S, p1.F); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
    ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
    (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) {
  pdd dp = p1 - p0
      , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  return q0 + pdd(
      cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
|} // from line p0--p1 to q0--q1, apply to r
```

9.2 ConvexHull

```
void hull(vector<pll> &dots) { // n=1 => ans = {}
  sort(dots.begin(), dots.end());
  vector<pll> ans(1, dots[0]);
  for (int ct = 0; ct < 2; ++ct, reverse(all(dots)))
    for (int i = 1, t = ans.size()
        ; i < dots.size(); ans.emplace_back(dots[i++]))
    while (ans.size() > t &&
        ori(ans.end()[-2], ans.back(), dots[i]) <= 0)
        ans.pop_back();
    ans.pop_back(), ans.swap(dots);
}</pre>
```

9.3 SortByAngle

```
int cmp(pll a, pll b, bool same = true) {
    #define is_neg(k) (
        sign(k.S) < 0 || (sign(k.S) == 0 && sign(k.F) < 0))
    int A = is_neg(a), B = is_neg(b);
    if (A != B)
        return A < B;
    if (sign(cross(a, b)) == 0)
        return same ? abs2(a) < abs2(b) : -1;
    return sign(cross(a, b)) > 0;
}
```

9.4 DisPointSegment

```
|double PointSegDist(pdd q0, pdd q1, pdd p) {
| if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
| if (sign(dot(q1 - q0,
| p - q0)) >= 0 && sign(dot(q0 - q1, p - q1)) >= 0)
| return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
| return min(abs(p - q0), abs(p - q1));
|}
```

9.5 PointInCircle

9.6 PointInConvex

9.7 PointTangentConvex

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
/* bool pred(int a, int b);
f(0) \sim f(n-1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int l = 0, r = n; bool rv = pred(1, 0);
  while (r - l > 1) {
    int m = (l + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
    else l = m;
  }
  return pred(l, r % n) ? l : r % n;
}
pii get_tangent(vector<pll>> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch((int)C.size(), [&](int x, int y)
    { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
|} // return (a, b), ori(p, C[a], C[b]) >= 0
```

9.8 CircTangentCirc

```
vector<Line
    > go( const Cir& c1 , const Cir& c2 , int sign1 ){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = abs2(c1.0 - c2.0);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
    pdd n = pdd(v.F * c - sign2 * h * v.S,
      v.S * c + sign2 * h * v.F);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sign(p1.F - p2.F) == 0 and
        sign(p1.S - p2.S) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
```

```
ret.emplace_back(Line(p1, p2));
}
return ret;
.
```

9.9 LineCircleIntersect

9.10 LineConvexIntersect

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch((int)C.size(), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
}
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = (int)C.size();
  if (cmpL(A) < 0 \mid \mid cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r; ) {
      int m = ((l + r + (l < r? 0 : n)) / 2) % n;
       (cmpL(m) == cmpL(t) ? l : r) = m;
    }
    return (l + !cmpL(r)) % n;
  };
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.F == res.S) // touching the corner i
    return pii(res.F, -1);
  if (!
       cmpL(res.F) && !cmpL(res.S)) // along side i, i+1
    switch ((res.F - res.S + n + 1) % n) {
       case 0: return pii(res.F, res.F);
      case 2: return pii(res.S, res.S);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned
       in the same order as the line hits the convex */
  return res;
|} // convex cut: (r, l]
```

9.11 CircIntersectCirc

9.12 PolyIntersectCirc

```
|// Divides into multiple triangle, and sum up
|const double PI = acos(-1);
|double _area(pdd pa, pdd pb, double r) {
```

```
if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
  if (abs(pb) < eps) return 0;</pre>
 double S, h, theta;
 double a = abs(pb), b = abs(pa), c = abs(pb - pa);
  double cosB = dot(pb, pb - pa) / a / c,
         B = acos(cosB);
  double cosC = dot(pa, pb) / a / b, C = acos(cosC);
 if (a > r) {
   S = (C / 2) * r * r;
   h = a * b * sin(C) / c;
    if (h < r && B < PI / 2)
      S = (acos(h / r) * r * r -
        h * sqrt(r * r - h * h));
 } else if (b > r) {
    theta = PI - B - asin(sin(B) / r * a);
    S = .5 * a * r * sin(theta) +
      (C - theta) / 2 * r * r;
  } else S = .5 * sin(C) * a * b;
double area_poly_circle(const vector<pdd> poly,
  const pdd &0, const double r) {
  double S = 0;
 for (int i = 0; i < (int)poly.size(); ++i)</pre>
    S += _area(poly[i] - 0,
           poly[(i + 1) % (int)poly.size()] - 0, r) *
      ori(
        0, poly[i], poly[(i + 1) % (int)poly.size()]);
 return fabs(S);
```

9.13 MinkowskiSum

```
vector<pll> Minkowski
     (vector<pll> A, vector<pll> B) { // |A|, |B| > 3
   hull(A), hull(B);
   vector<pll> C(1, A[0] + B[0]), s1, s2;
  for (int i = 0; i < A.size(); ++i)</pre>
     s1.emplace_back(A[(i + 1) % A.size()] - A[i]);
   for (int i = 0; i < B.size(); i++)</pre>
     s2.emplace_back(B[(i + 1) % B.size()] - B[i]);
  for (int i = 0, j = 0; i < A.size() || j < B.size();)</pre>
     if (j >= B.size()
          || (i < A.size() \&\& cross(s1[i], s2[j]) >= 0))
       C.emplace_back(B[j % B.size()] + A[i++]);
     else
       C.emplace_back(A[i % A.size()] + B[j++]);
   return hull(C), C;
|}
```

9.14 MinMaxEnclosingRect

```
const double INF = 1e18, qi = acos(-1) / 2 * 3;
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
 hull(dots);
  double Max = 0, Min = INF, deg;
 int n = (int)dots.size();
 dots.emplace_back(dots[0]);
 for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
    pll nw = vec(i + 1);
    while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
     u = (u + 1) \% n;
    while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
     r = (r + 1) \% n;
    if (!i) l = (r + 1) % n;
    while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))
     l = (l + 1) \% n;
   Min = min(Min, (double)(dot(nw, vec(r)) - dot
        (nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));
    deg = acos(dot(diff(r
        , l), vec(u)) / abs(diff(r, l)) / abs(vec(u)));
    deg = (qi - deg) / 2;
    Max = max(Max, abs(diff))
        (r, l)) * abs(vec(u)) * sin(deg) * sin(deg));
```

```
return pdd(Min, Max);
}
 9.15 MinEnclosingCircle
| pdd Minimum_Enclosing_Circle
     (vector<pdd> dots, double &r) {
   pdd cent;
   random_shuffle(all(dots));
   cent = dots[0], r = 0;
   for (int i = 1; i < (int)dots.size(); ++i)</pre>
     if (abs(dots[i] - cent) > r) {
       cent = dots[i], r = 0;
       for (int j = 0; j < i; ++j)</pre>
         if (abs(dots[j] - cent) > r) {
           cent = (dots[i] + dots[j]) / 2;
           r = abs(dots[i] - cent);
           for(int k = 0; k < j; ++k)</pre>
             if(abs(dots[k] - cent) > r)
               cent = excenter
                    (dots[i], dots[j], dots[k], r);
         }
  return cent;
}
 9.16 CircleCover
const int N = 1021:
struct CircleCover {
   int C;
   Cir c[N];
   bool g[N][N], overlap[N][N];
   // Area[i] : area covered by at least i circles
   double Area[ N ];
   void init(int _C){ C = _C;}
   struct Teve {
     pdd p; double ang; int add;
     Teve() {}
     Teve(pdd _a
          , double _b, int _c):p(_a), ang(_b), add(_c){}
     bool operator<(const Teve &a)const</pre>
     {return ang < a.ang;}
   eve[N * 2];
   // strict: x = 0, otherwise x = -1
   bool disjuct(Cir &a, Cir &b, int x)
   {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
   bool contain(Cir &a, Cir &b, int x)
   {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
   bool contain(int i, int j) {
     /* c[j] is non-strictly in c[i]. */
     return (sign
         (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
         R) == 0 \& i < j) && contain(c[i], c[j], -1);
   void solve(){
     fill_n(Area, C + 2, 0);
     for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         overlap[i][j] = contain(i, j);
     for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
             disjuct(c[i], c[j], -1));
     for(int i = 0; i < C; ++i){</pre>
       int E = 0, cnt = 1;
       for(int j = 0; j < C; ++j)</pre>
         if(j != i && overlap[j][i])
           ++cnt;
       for(int j = 0; j < C; ++j)</pre>
         if(i != j && g[i][j]) {
           pdd aa, bb;
           CCinter(c[i], c[j], aa, bb);
           double A =
```

atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);

atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);

double B =

```
eve[E++] = Teve
             (bb, B, 1), eve[E++] = Teve(aa, A, -1);
        if(B > A) ++cnt;
    if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
    else{
      sort(eve, eve + E);
      eve[E] = eve[0];
      for(int j = 0; j < E; ++j){</pre>
        cnt += eve[j].add;
             ] += cross(eve[j].p, eve[j + 1].p) * .5;
        double theta = eve[j + 1].ang - eve[j].ang;
        if (theta < 0) theta += 2. * pi;</pre>
        Area[cnt] += (theta
              - sin(theta)) * c[i].R * c[i].R * .5;
    }
  }
}
```

9.17 LineCmp

```
using Line = pair<pll, pll>;
struct lineCmp {
  bool operator()(Line l1, Line l2) const {
    int X =
       (\max(l1.F.F, l2.F.F) + \min(l1.S.F, l2.S.F)) / 2;
    ll p1 =
          (X - l1.F.F) * l1.S.S + (l1.S.F - X) * l1.F.S,
        p2 =
         (X - 12.F.F) * 12.S.S + (12.S.F - X) * 12.F.S,
        q1 = (l1.S.F - l1.F.F), q2 = (l2.S.F - l2.F.F);
    if (q1 == 0) p1 = l1.F.S + l1.S.S, q1 = 2;
    if (q2 == 0) p2 = l2.F.S + l2.S.S, q2 = 2;
    if (l1.F == l2.F || l2.F == l2.S) l1 = l2;
    return make_tuple((__int128)(p1 * q2), l1) <</pre>
       make_tuple((__int128)(p2 * q1), l2);
  }
};
```

9.18 Trapezoidalization

```
struct SweepLine {
  struct cmp {
    cmp(const SweepLine &_swp): swp(_swp) {}
    bool operator()(int a, int b) const {
      if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
        return swp.slope_cmp(a, b);
      return swp.get_y(a) + swp.eps < swp.get_y(b);</pre>
   const SweepLine &swp;
    _cmp;
  T curTime, eps, curQ;
  vector<Line> base;
  multiset<int, cmp> sweep;
  multiset<pair<T, int>> event;
  vector<typename multiset<int, cmp>::iterator> its;
  vector
      <typename multiset<pair<T, int>>::iterator> eits;
 bool slope_cmp(int a, int b) const {
    assert(a != -1);
    if (b == -1) return 0;
    return sign(cross(base
        [a].Y - base[a].X, base[b].Y - base[b].X)) < 0;
 }
 T get_y(int idx) const {
    if (idx == -1) return curQ;
    Line l = base[idx];
    if (l.X.X == l.Y.X) return l.Y.Y;
    return ((curTime - l.X.X) * l.Y.Y
        + (l.Y.X - curTime) * l.X.Y) / (l.Y.X - l.X.X);
  void insert(int idx) {
    its[idx] = sweep.insert(idx);
    if (its[idx] != sweep.begin())
```

```
update_event(*prev(its[idx]));
  update_event(idx);
  event.emplace
      (base[idx].Y.X, idx + 2 * (int)base.size());
void erase(int idx) {
  assert(eits[idx] == event.end());
  auto p = sweep.erase(its[idx]);
  its[idx] = sweep.end();
  if (p != sweep.begin())
    update_event(*prev(p));
void update_event(int idx) {
  if (eits[idx] != event.end())
    event.erase(eits[idx]);
  eits[idx] = event.end();
  auto nxt = next(its[idx]);
  if (nxt ==
       sweep.end() || !slope_cmp(idx, *nxt)) return;
  auto t = intersect(base[idx].
      X, base[idx].Y, base[*nxt].X, base[*nxt].Y).X;
  if (t + eps < curTime || t</pre>
       >= min(base[idx].Y.X, base[*nxt].Y.X)) return;
  eits[idx
      ] = event.emplace(t, idx + (int)base.size());
void swp(int idx) {
  assert(eits[idx] != event.end());
  eits[idx] = event.end();
  int nxt = *next(its[idx]);
  swap((int&)*its[idx], (int&)*its[nxt]);
  swap(its[idx], its[nxt]);
  if (its[nxt] != sweep.begin())
    update_event(*prev(its[nxt]));
  update_event(idx);
// only expected to call the functions below
SweepLine(T t, T e, vector<Line> vec): _cmp
    (*this), curTime(t), eps(e), curQ(), base(vec),
     sweep(_cmp), event(), its((int)vec.size(), sweep
    .end()), eits((int)vec.size(), event.end()) {
  for (int i = 0; i < (int)base.size(); ++i) {</pre>
    auto &[p, q] = base[i];
    if (p > q) swap(p, q);
    if (p.X <= curTime && curTime <= q.X)</pre>
      insert(i);
    else if (curTime < p.X)</pre>
      event.emplace(p.X, i);
}
void setTime(T t, bool ers = false) {
  assert(t >= curTime);
  while (!event.empty() && event.begin()->X <= t) {</pre>
    auto [et, idx] = *event.begin();
    int s = idx / (int)base.size();
    idx %= (int)base.size();
    if (abs(et - t) <= eps && s == 2 && !ers) break;</pre>
    curTime = et;
    event.erase(event.begin());
    if (s == 2) erase(idx);
    else if (s == 1) swp(idx);
    else insert(idx);
  curTime = t;
T nextEvent() {
  if (event.empty()) return INF;
  return event.begin()->X;
int lower_bound(T y) {
  curQ = y;
  auto p = sweep.lower_bound(-1);
  if (p == sweep.end()) return -1;
  return *p;
```

9.19 TriangleHearts

|};

```
| pdd circenter(
  pdd p0, pdd p1, pdd p2) { // radius = abs(center)
  p1 = p1 - p0, p2 = p2 - p0;
double x1 = p1.F, y1 = p1.S, x2 = p2.F, y2 = p2.S;
  double m = 2. * (x1 * y2 - y1 * x2);
  pdd center = pdd((x1 * x1 * y2 - x2 * x2 * y1 +
                       y1 * y2 * (y1 - y2)) / m,
     (x1 * x2 * (x2 - x1) - y1 * y1 * x2 +
       x1 * y2 * y2) / m);
  return center + p0;
}
pdd incenter(
  pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
   double a = abs(p2 - p3), b = abs(p1 - p3),
          c = abs(p1 - p2);
   double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3) {
  return (p1 + p2 + p3) / 3;
pdd orthcenter(pdd p1, pdd p2, pdd p3) {
   return masscenter(p1, p2, p3) * 3 -
     circenter(p1, p2, p3) * 2;
```

9.20 HalfPlaneIntersect

```
pll area_pair(Line a, Line b)
{ return pll(cross(a.S
     - a.F, b.F - a.F), cross(a.S - a.F, b.S - a.F));    }
bool isin(Line l0, Line l1, Line l2) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(l0, l2);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
  return (__int128)
       a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(all(arr), [&](Line a, Line b) -> int {
    if (cmp(a.S - a.F, b.S - b.F, 0) != -1)
      return cmp(a.S - a.F, b.S - b.F, 0);
    return ori(a.F, a.S, b.S) < 0;</pre>
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(
        dq.back().S - dq.back().F, p.S - p.F, 0) == -1)
      continue;
    while ((int)dq.size() >= 2
        && !isin(p, dq[(int)dq.size() - 2], dq.back()))
      dq.pop_back();
    while
        ((int)dq.size() \ge 2 \&\& !isin(p, dq[0], dq[1]))
      dq.pop_front();
    dq.emplace_back(p);
  while ((int)dq.size() >= 3 &&
       !isin(dq[0], dq[(int)dq.size() - 2], dq.back()))
    dq.pop_back();
  while ((int)
      dq.size() >= 3 \& !isin(dq.back(), dq[0], dq[1]))
    dq.pop_front();
  return vector<Line>(all(dq));
```

9.21 RotatingSweepLine

```
| void rotatingSweepLine(vector<pii> &ps) {
| int n = (int)ps.size(), m = 0;
```

```
vector<int> id(n), pos(n);
   vector<pii> line(n * (n - 1));
   for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++j)</pre>
       if (i != j) line[m++] = pii(i, j);
   sort(all(line), [&](pii a, pii b) {
     return cmp(ps[a.S] - ps[a.F], ps[b.S] - ps[b.F]);
   }); // cmp(): polar angle compare
   iota(all(id), 0);
   sort(all(id), [&](int a, int b) {
     if (ps[a].S != ps[b].S) return ps[a].S < ps[b].S;</pre>
     return ps[a] < ps[b];</pre>
   }); // initial order, since (1, 0) is the smallest
   for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
   for (int i = 0; i < m; ++i) {</pre>
     auto l = line[i];
     // do something
     tie(pos[l.F], pos[l.S], id[pos[l.F]], id[pos[l.S
         ]]) = make_tuple(pos[l.S], pos[l.F], l.S, l.F);
}
```

9.22 DelaunayTriangulation

```
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const
     ll inf = MAXC * MAXC * 100; // Lower_bound unknown
struct Tri;
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
};
struct Tri {
  pll p[3];
  Edge edge[3];
  Tri* chd[3];
  Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)</pre>
      if (ori(p[i], p[(i + 1) % 3], q) < 0)
        return 0;
    return 1;
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
    the_root
         = // Tri should at least contain all points
      new(tris++) Tri(pll(-inf, -inf),
           pll(inf + inf, -inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
```

```
void add_point(const
        Tri* the_root;
   static Tri* find(Tri* root, const pll &p) {
    while (1) {
      if (!root->has_chd())
        return root;
       for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break;
        }
    }
    assert(0); // "point not found"
  }
  void add_point(Tri* root, pll const& p) {
    /* split it into three triangles */
    for (int i = 0; i < 3; ++i)
      t[i] = new(tris
           ++) Tri(root->p[i], root->p[(i + 1) % 3], p);
    for (int i = 0; i < 3; ++i)</pre>
       edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)</pre>
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)</pre>
       flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p
         [0], tri->p[1], tri->p[2], trj->p[pj])) return;
     /* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri
         (tri->p[(pi + 1) % 3], trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri
         (trj->p[(pj + 1) % 3], tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
    edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd
         [0] = trk; tri->chd[1] = trl; tri->chd[2] = 0;
    trj->chd
         [0] = trk; trj->chd[1] = trl; trj->chd[2] = 0;
    flip(trk, 1); flip(trk, 2);
    flip(trl, 1); flip(trl, 2);
  }
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
  if (vst.find(now) != vst.end())
    return;
  vst.insert(now);
  if (!now->has_chd())
    return triang.emplace_back(now);
  for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(ps, ps + n);
   Trig tri; // the triangulation structure
  for (int i = 0; i < n; ++i)</pre>
    tri.add_point(ps[i]);
   go(tri.the_root);
| }
```

9.23 VonoroiDiagram

10 Misc

10.1 MoAlgoWithModify

```
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
   int L, R, LBid, RBid, T;
   Query(int l, int r, int t):
     L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
   bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     if (RBid != q.RBid) return RBid < q.RBid;</pre>
     return T < b.T;</pre>
  }
};
void solve(vector<Query> query) {
   sort(ALL(query));
   int L=0, R=0, T=-1;
   for (auto q : query) {
     while (T < q.T) addTime(L, R, ++T); // TODO
     while (T > q.T) subTime(L, R, T--); // TODO
     while (R < q.R) add(arr[++R]); // TODO</pre>
     while (L > q.L) add(arr[--L]); //
     while (R > q.R) sub(arr[R--]); // TODO
     while (L < q.L) sub(arr[L++]); // TODO</pre>
     // answer query
   }
}
```

10.2 MoAlgoOnTree

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord[in[u]] = ord[out[u]] = u
4) bitset<MAXN> inset
*/
struct Query -
  int L, R, LBid, lca;
  Query(int u, int v) {
    int c = LCA(u, v);
    if (c == u || c == v)
      q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
    else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
    else
      q.lca = c, q.L = out[v], q.R = in[u];
    q.Lid = q.L / blk;
  bool operator<(const Query &q) const {</pre>
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    return R < q.R;</pre>
 }
};
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
    else add(arr[x]); // TODO
    inset[x] = ~inset[x];
```

```
void solve(vector<Query> query) {
   sort(ALL(query));
   int L = 0, R = 0;
  for (auto q : query) {
     while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
     while (R > q.R) flip(ord[R--]);
     while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
|}
```

10.3 MoAlgoAdvanced

- Mo's Algorithm With Addition Only
 - Sort querys same as the normal Mo's algorithm.

 - For each query [l,r]:
 If l/blk = r/blk, brute-force.
 - If $l/blk \neq curL/blk$, initialize $curL := (l/blk+1) \cdot blk$, curR :=
 - If r > curR, increase curR
- decrease curL to fit l, and then undo after answering
- · Mo's Algorithm With Offline Second Time
 - Require: Changing answer \equiv adding f([l,r],r+1).
 - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1).
 - Part1: Answer all f([1,r],r+1) first. Part2: Store $curR \to R$ for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
 - Note: You must do the above symmetrically for the left boundaries.

10.4 HilbertCurve

```
ll hilbert(int n, int x, int y) {
  ll res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y \& s) > 0;
    res += s * 111 * s * ((3 * rx) ^ ry);
    if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
  }
  return res;
|} // n = 2^k
```

10.5 SternBrocotTree

- Construction: Root $\frac{1}{1}$, left/right neighbor $\frac{0}{1}, \frac{1}{0}$, each node is sum of last left/right neighbor: $\frac{a}{b}, \frac{c}{d} \rightarrow \frac{a+c}{b+d}$
- Property: Adjacent (mid-order DFS) $\frac{a}{b},\frac{c}{d}\!\Rightarrow\!bc\!-\!ad\!=\!1$.
- Search known $\frac{p}{a}$: keep L-R alternative. Each step can calcaulated in $O(1) \Rightarrow$ total $O(\log C)$.
- Search unknown $\frac{p}{a}$: keep L-R alternative. Each step can calcaulated in $O(\log C)$ checks \Rightarrow total $O(\log^2 C)$ checks.

10.6 AllLCS

```
void all_lcs(string s, string t) { // O-base
  vector<int> h((int)t.size());
   iota(all(h), 0);
   for (int a = 0; a < (int)s.size(); ++a) {</pre>
    int v = -1;
     for (int c = 0; c < (int)t.size(); ++c)</pre>
       if (s[a] == t[c] || h[c] < v)
         swap(h[c], v);
     // LCS(s[0, a], t[b, c]) =
     // c - b + 1 - sum([h[i] >= b] | i <= c)
     // h[i] might become -1 !!
  }
į }
```

10.7 SimulatedAnnealing

```
double factor = 100000:
 const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {</pre>
         answer, nw: current value, rnd(): mt19937 rnd()
     if (exp(-(nw - ans
         ) / factor) >= (double)(rnd() % base) / base)
         ans = nw;
     factor *= 0.99995;
}
```

10.8 SMAWK

```
int opt[N];
ll A(int x, int y); // target func
void smawk(vector<int> &r, vector<int> &c);
void interpolate(vector<int> &r, vector<int> &c) {
  int n = (int)r.size();
  vector<int> er;
  for (int i = 1; i < n; i += 2) er.emplace_back(r[i]);</pre>
  smawk(er, c);
  for (int i = 0, j = 0; j < c.size(); j++) {</pre>
     if (A(r[i], c[j]) < A(r[i], opt[r[i]]))</pre>
      opt[r[i]] = c[i];
     if (i + 2 < n \&\& c[j] == opt[r[i + 1]])
       j--, i += 2;
void reduce(vector<int> &r, vector<int> &c) {
  int n = (int)r.size();
  vector<int> nc;
  for (int i : c) {
    int j = (int)nc.size();
     while (
       j \& A(r[j-1], nc[j-1]) > A(r[j-1], i))
       nc.pop_back(), j--;
     if (nc.size() < n) nc.emplace_back(i);</pre>
  smawk(r, nc);
}
void smawk(vector<int> &r, vector<int> &c) {
  if (r.size() == 1 && c.size() == 1) opt[r[0]] = c[0];
  else if (r.size() >= c.size()) interpolate(r, c);
  else reduce(r, c);
|}
```

10.9 Python

|math.isqrt(2) # integer sqrt

10.10 LineContainer

```
struct Line {
  mutable ll k, m, p;
  bool operator<(const Line &o) const {</pre>
    return k < o.k;
 bool operator<(ll x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>> {
  // (for doubles, use inf = 1/.0, div(a,b) = a/b)
static const ll inf = LLONG_MAX;
  ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b);
  bool isect(iterator x, iterator y) {
    if (y == end()) return x->p = inf, 0;
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
    else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
  void add(ll k, ll m) {
    auto z = insert(\{k, m, 0\}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
```

```
if (x != begin() && isect(--x, y))
    isect(x, y = erase(y));
while ((y = x) != begin() && (--x)->p >= y->p)
    isect(x, erase(y));
}
ll query(ll x) {
    assert(!empty());
    auto l = *lower_bound(x);
    return l.k * x + l.m;
};
```