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```

# Basic

### 1.1 .vimrc

```
set ru nu cin cul sc so=3 ts=4 sw=4 bs=2 ls=2 mouse=a
inoremap {<CR> {<CR>}<C-o>0
map <F7> :w<CR>:!g++

"%" -Wall -Wextra -Wshadow -Wconversion -fsanitize
    =address,=undefined -D_GLIBCXX_DEBUG && ./a.out<CR>
```

### 1.2 PBDS

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
tree<int, null_type, less<int>, rb_tree_tag
     tree_order_statistics_node_update > bst;
// order_of_key(n): # of elements <= n</pre>
// find_by_order(n): 0-indexed
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
__gnu_pbds::priority_queue
    <int, greater<int>, pairing_heap_tag> pq;
```

# 1.3 pargma

```
#pragma GCC optimize("Ofast,unroll-loops")
#pragma GCC target("avx,avx2,sse,sse2
    ,sse3,ssse3,sse4,popcnt,abm,mmx,fma,tune=native")
```

# 2 Graph

# 2.1 2SAT/SCC

```
struct SAT { // 0-base
  int low[N], dfn[N], bln[N], n, Time, nScc;
   bool instack[N], istrue[N];
  stack<int> st;
  vector < int > G[N], SCC[N];
  void init(int _n) {
     n = _n; // assert(n * 2 <= N);
     for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
   void add_edge(int a, int b) { G[a].emplace_back(b); }
   int rv(int a) {
     if (a >= n) return a - n;
     return a + n;
  void add_clause(int a, int b) {
     add_edge(rv(a), b), add_edge(rv(b), a);
  void dfs(int u) {
     dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
       do {
         tmp = st.top(), st.pop();
instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != u);
  bool solve() {
     Time = nScc = 0;
     for (int i = 0; i < n + n; ++i)</pre>
     SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;

for (int i = 0; i < n + n; ++i)

if (!dfn[i]) dfs(i);
     for (int i =
          0; i < n + n; ++i) SCC[bln[i]].emplace_back(i);
     for (int i = 0; i < n; ++i) {
  if (bln[i] == bln[i + n]) return false;</pre>
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true:
  }
};
```

### 2.2 BCC Vertex

```
int n, m, dfn[N], low[N], is_cut[N], nbcc = 0, t = 0;
vector < int > g[N], bcc[N], G[2 * N];
stack<int> st;
void tarjan(int p, int lp) {
  dfn[p] = low[p] = ++t;
  st.push(p);
  for (auto i : g[p]) {
     if (!dfn[i]) {
       tarjan(i, p);
        low[p] = min(low[p], low[i]);
       if (dfn[p] <= low[i]) {</pre>
          nbcc++
          is_cut[p] = 1;
          for (int x = 0; x != i; st.pop()) {
            x = st.top();
             bcc[nbcc].push_back(x);
          bcc[nbcc].push_back(p);
     } else low[p] = min(low[p], dfn[i]);
  }
void build() { // [n+1,n+nbcc] cycle, [1,n] vertex
for (int i = 1; i <= nbcc; i++) {
   for (auto j : bcc[i]) {</pre>
       G[i + n].push_back(j);
       G[j].push_back(i + n);
  }
```

# 2.3 MinimumMeanCycle

```
ll road[N][N]; // input here
struct MinimumMeanCycle {
  ll dp[N + 5][N], n;
  pll solve() {
    ll a = -1, b = -1, L = n + 1;
    for (int i = 2; i <= L; ++i)</pre>
       for (int k = 0; k < n; ++k)
         for (int j = 0; j < n; ++j)
           dp[i][j] =
             min(dp[i - 1][k] + road[k][j], dp[i][j]);
    for (int i = 0; i < n; ++i) {</pre>
       if (dp[L][i] >= INF) continue;
       ll ta = 0, tb = 1;
      for (int j = 1; j < n; ++j)
         if (dp[j][i] < INF &&
           ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
ta = dp[L][i] - dp[j][i], tb = L - j;
       if (ta == 0) continue;
      if (a == -1 || a * tb > ta * b) a = ta, b = tb;
    if (a != -1) {
      ll g = __gcd(a, b);
return pll(a / g, b / g);
    return pll(-1LL, -1LL);
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)</pre>
       for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;</pre>
};
```

# 2.4 MaximumCliqueDyn

```
struct MaxClique { // fast when N <= 100</pre>
 bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
   G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(all(r)
          , [&](int x, int y) { return d[x] > d[y]; });
    vector<int> c(r.size());
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first
          (); p < N; p = cs[k]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
 }
  void dfs(vector<</pre>
      int> &r, vector<int> &c, int l, bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.emplace_back(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
   }
  int solve() {
   vector<int> r(n);
    ans = q = 0, iota(all(r), 0);
```

```
pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
}
};
```

### 2.5 MinimumSteinerTree

```
struct SteinerTree { // 0-base
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcst[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {</pre>
       fill_n(dst[i], n, INF);
       dst[i][i] = vcst[i] = 0;
    }
  void chmin(int &x, int val) {
    x = min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)</pre>
       for (int i = 0; i < n; ++i)</pre>
         for (int j = 0; j < n; ++j)</pre>
           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
    shortest_path();
     int t = SZ(ter), full = (1 << t) - 1;</pre>
    for (int i = 0; i <= full; ++i)</pre>
       fill_n(dp[i], n, INF);
     copy_n(vcst, n, dp[0]);
     for (int msk = 1; msk <= full; ++msk) {</pre>
       if (!(msk & (msk - 1))) {
         int who = __lg(msk);
for (int i = 0; i < n; ++i)</pre>
           dp[msk
               ][i] = vcst[ter[who]] + dst[ter[who]][i];
       for (int i = 0; i < n; ++i)</pre>
         for (int sub = (
             msk - 1) & msk; sub; sub = (sub - 1) & msk)
           chmin(dp[msk][i],
               dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
       for (int i = 0; i < n; ++i) {</pre>
         tdst[i] = INF;
         for (int j = 0; j < n; ++j)</pre>
           chmin(tdst[i], dp[msk][j] + dst[j][i]);
       copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
}; // O(V 3^T + V^2 2^T)
```

### 2.6 DMST(slow)

```
struct zhu_liu { // O(VE)
 struct edge {
    int u, v;
    ll w;
 }:
 vector<edge> E; // 0-base
 int pe[N], id[N], vis[N];
 ll in[N];
 void init() { E.clear(); }
 void add_edge(int u, int v, ll w) {
    if (u != v) E.emplace_back(edge{u, v, w});
  ll build(int root, int n) {
    ll ans = 0;
    for (;;) {
      fill_n(in, n, INF);
      for (int i = 0; i < E.size(); ++i)</pre>
        if (E[i].u != E[i].v && E[i].w < in[E[i].v])</pre>
         pe[E[i].v] = i, in[E[i].v] = E[i].w;
      for (int u = 0; u < n; ++u) // no solution</pre>
        if (u != root && in[u] == INF) return -INF;
      int cntnode = 0;
      fill_n(id, n, -1), fill_n(vis, n, -1);
      for (int u = 0; u < n; ++u) {
        if (u != root) ans += in[u];
        int v = u;
```

```
while (vis[v] != u && !~id[v] && v != root)
         vis[v] = u, v = E[pe[v]].u;
if (v != root && !~id[v]) {
            for (int x = E[pe[v]].u; x != v;
                  x = E[pe[x]].u)
              id[x] = cntnode;
            id[v] = cntnode++;
         }
       if (!cntnode) break; // no cycle
       for (int u = 0; u < n; ++u)
  if (!~id[u]) id[u] = cntnode++;</pre>
       for (int i = 0; i < E.size(); ++i) {</pre>
         int v = E[i].v;
         E[i].u = id[E[i].u], E[i].v = id[E[i].v];
         if (E[i].u != E[i].v) E[i].w -= in[v];
       n = cntnode, root = id[root];
     return ans;
  }
};
```

### 2.7 DMST

```
#define rep(i, a, b) for (int i = a; i < (b); ++i)</pre>
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef vector<int> vi;
struct RollbackUF {
 vi e:
  vector<pii> st;
  RollbackUF(int n) : e(n, -1) {}
  int size(int x) { return -e[find(x)]; }
  int find(int x) { return e[x] < 0 ? x : find(e[x]); }</pre>
  int time() { return sz(st); }
  void rollback(int t) {
    for (int i = time(); i-- > t;)
      e[st[i].first] = st[i].second;
    st.resize(t);
 bool join(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]});
    st.push_back({b, e[b]});
    e[a] += e[b];
   e[b] = a;
    return true:
 }
struct Edge {
 int a, b;
 ll w:
struct Node { /// lazy skew heap node
 Edge key;
  Node *l, *r;
  ll delta;
  void prop() {
    key.w += delta;
    if (l) l->delta += delta;
    if (r) r->delta += delta;
    delta = 0;
 Edge top() {
   prop();
    return key;
 }
Node *merge(Node *a, Node *b) {
 if (!a || !b) return a ?: b;
  a->prop(), b->prop();
 if (a->key.w > b->key.w) swap(a, b);
 swap(a->l, (a->r = merge(b, a->r)));
  return a;
void pop(Node *&a) {
 a->prop();
 a = merge(a->l, a->r);
pair<ll, vi> dmst(int n, int r, vector<Edge> &g) {
 RollbackUF uf(n);
 vector < Node *> heap(n);
```

```
for (Edge e : g)
    heap[e.b] = merge(heap[e.b], new Node{e});
  ll res = 0;
  vi seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector < Edge > Q(n), in(n, \{-1, -1\}), comp;
  deque<tuple<int, int, vector<Edge>>> cycs;
  rep(s, 0, n) {
    int u = s, qi = 0, w;
    while (seen[u] < 0) {</pre>
      if (!heap[u]) return {-1, {}};
      Edge e = heap[u]->top();
      heap[u]->delta -= e.w, pop(heap[u]);
Q[qi] = e, path[qi++] = u, seen[u] = s;
      res += e.w, u = uf.find(e.a);
      if (seen[u] == s) { /// found cycle, contract
        Node *cyc = 0;
         int end = qi, time = uf.time();
         do cyc = merge(cyc, heap[w = path[--qi]]);
        while (uf.join(u, w));
        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
        cycs.push\_front(\{u,\ time,\ \{\&Q[qi],\ \&Q[end]\}\});
      }
    rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
  for (auto &[u, t, comp] :
    cycs) { // restore sol (optional)
    uf.rollback(t);
    Edge inEdge = in[u];
for (auto &e : comp) in[uf.find(e.b)] = e;
    in[uf.find(inEdge.b)] = inEdge;
  rep(i, 0, n) par[i] = in[i].a;
  return {res, par};
2.8 VizingTheorem
```

```
namespace Vizing { // Edge coloring
                    // G: coloring adjM
int C[N][N], G[N][N];
void clear(int n) {
  for (int i = 0; i <= n; i++) {
    for (int j = 0; j <= n; j++) C[i][j] = G[i][j] = 0;</pre>
  }
void solve(vector<pii> &E, int n, int m) {
  int X[n] = {}, a;
  auto update = [&](int u) {
    for (X[u] = 1; C[u][X[u]]; X[u]++);
  auto color = [&](int u, int v, int c) {
    int p = G[u][v];
    G[u][v] = G[v][u] = c;
    C[u][c] = v;
    C[v][c] = u;
    C[u][p] = C[v][p] = 0;
    if (p) X[u] = X[v] = p;
    else update(u), update(v);
    return p;
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  for (int i = 1; i <= n; i++) X[i] = 1;</pre>
  for (int t = 0; t < E.size(); t++) {</pre>
    int u = E[t].first, v0 = E[t].second, v = v0,
        c0 = X[u], c = c0, d;
    vector<pii> L;
    int vst[n] = {};
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
      if (!C[v][c])
        for (a = (int)L.size() - 1; a >= 0; a--)
          c = color(u, L[a].first, c);
      else if (!C[u][d])
        for (a = (int)L.size() - 1; a >= 0; a--)
          color(u, L[a].first, L[a].second);
      else if (vst[d]) break;
```

# 2.9 MinimumCliqueCover

```
struct Clique_Cover { // 0-base, 0(n2^n)
  int co[1 << N], n, E[N];</pre>
  int dp[1 << N];</pre>
  void init(int _n) {
    n = _n, fill_n(dp, 1 << n, 0);
</pre>
     fill_n(E, n, 0), fill_n(co, 1 << n, 0);
  void add_edge(int u, int v) {
    E[u] = 1 << v, E[v] = 1 << u;
  int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;
     dp[0] = (n \& 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {</pre>
       int t = i & -i;
       dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)</pre>
       co[i] = (co[i] & i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic</pre>
       for (int i = 0; i < (1 << n); ++i)</pre>
         sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n:
  }
};
```

# 2.10 CountMaximalClique

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
   g[u][v] = g[v][u] = 1;
 void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
      int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
      all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)
  if (g[v][some[d][j]])</pre>
           some[d + 1][tsn++] = some[d][j];
      for (int j = 0; j < nn; ++j)</pre>
        if (g[v][none[d][j]])
          none[d + 1][tnn++] = none[d][j];
      dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
   }
  int solve() {
    iota(some[0], some[0] + n, 1);
```

```
S = 0, dfs(0, 0, n, 0);
return S;
}
};
```

### 2.11 Theorems

|Maximum independent edge set| = |V| - |Minimum edge cover| |Maximum independent set| = |V| - |Minimum vertex cover|

# 3 Flow-Matching

#### 3.1 KM

```
struct KM { // 0-base
   ll w[N][N], hl[N], hr[N], slk[N];
   int \hat{f}[[\tilde{N}]], fr[\tilde{N}], pre[\tilde{N}], qu[\tilde{N}], ql, qr, n; bool vl[N], vr[N];
   void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i)</pre>
       fill_n(w[i], n, -INF);
   void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
   bool Check(int x) {
  if (vl[x] = 1, ~fl[x])
       return vr[qu[qr++] = fl[x]] = 1;
     while (~x) swap(x, fr[fl[x] = pre[x]]);
     return 0;
   void bfs(int s) {
     fill_n(slk
          , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
     ql = qr = 0, qu[qr++] = s, vr[s] = 1;
     for (ll d;;) {
       while (ql < qr)</pre>
          for (int x = 0, y = qu[ql++]; x < n; ++x)
            if (!vl[x] && slk
                [x] >= (d = hl[x] + hr[y] - w[x][y])) {
              if (pre[x] = y, d) slk[x] = d;
              else if (!Check(x)) return;
       d = INF:
       for (int x = 0; x < n; ++x)
         if (!vl[x] && d > slk[x]) d = slk[x];
       for (int x = 0; x < n; ++x) {
         if (vl[x]) hl[x] += d;
          else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
          if (!vl[x] && !slk[x] && !Check(x)) return;
    }
   ll solve() {
     fill_n(fl
          , n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
     for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(w[i], w[i] + n);
     for (int i = 0; i < n; ++i) bfs(i);</pre>
     ll res = 0;
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
     return res:
  }
};
```

#### 3.2 MCMF

```
struct MinCostMaxFlow { // 0-base
    struct Edge {
        ll from, to, cap, flow, cost, rev;
    } *past[N];
    vector < Edge > G[N];
    int inq[N], n, s, t;
    ll dis[N], up[N], pot[N];
    bool BellmanFord() {
        fill_n(dis, n, INF), fill_n(inq, n, 0);
        queue < int > q;
        auto relax = [&](int u, ll d, ll cap, Edge *e) {
        if (cap > 0 && dis[u] > d) {
            dis[u] = d, up[u] = cap, past[u] = e;
            if (!inq[u]) inq[u] = 1, q.push(u);
```

```
}
     }:
     relax(s, 0, INF, 0);
     while (!q.empty()) {
       int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u]) {
         ll d2 = dis[u] + e.cost + pot[u] - pot[e.to];
              (e.to, d2, min(up[u], e.cap - e.flow), &e);
       }
    }
     return dis[t] != INF;
  void solve(int
       , int _t, ll &flow, ll &cost, bool neg = true) {
= _s, t = _t, flow = 0, cost = 0;
     if (neg) BellmanFord(), copy_n(dis, n, pot);
     for (; BellmanFord(); copy_n(dis, n, pot)) {
       for (int
       \dot{i} = 0; i < n; ++i) dis[i] += pot[i] - pot[s]; flow += up[t], cost += up[t] * dis[t];
       for (int i = t; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         e.flow += up[t], G[e.to][e.rev].flow -= up[t];
       }
    }
  void init(int _n) {
    n = _n, fill_n(pot, n, 0);
     for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
  }
};
```

# 3.3 GeneralGraphMatching

```
struct Matching { // 0-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
    }
  void Blossom(int x, int y, int l) {
   for (; Find(x) != l; x = pre[y]) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
    }
  bool Bfs(int r) {
    iota(all(fa), 0); fill(all(s), -1);
    q = queue < int >(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : G[x])
        if (s[u] == -1) {
           if (pre[u] = x, s[u] = 1, match[u] == n) {
             for (int a = u, b = x, last;
                 b != n; a = last, b = pre[a])
               last =
                   match[b], match[b] = a, match[a] = b;
             return true;
        q.push(match[u]); s[match[u]] = 0;
} else if (!s[u] && Find(u) != Find(x)) {
           int l = LCA(u, x);
           Blossom(x, u, l); Blossom(u, x, l);
        }
    return false;
  Matching(int _n) : n(_n), fa(n + 1), s(n + 1), vis
  (n + 1), pre(n + 1, n), match(n + 1, n), G(n) {} void add_edge(int u, int v)
  { G[u].emplace_back(v), G[v].emplace_back(u); }
```

```
int solve() {
  int ans = 0;
  for (int x = 0; x < n; ++x)
   if (match[x] == n) ans += Bfs(x);
  return ans;
} // match[x] == n means not matched
```

# 3.4 MaxWeightMaching

```
#define rep(i, l, r) for (int i = (l); i <= (r); ++i)</pre>
struct WeightGraph { // 1-based
  struct edge {
    int u, v, w;
  };
  int n, nx;
  vector<int> lab;
  vector<vector<edge>> g;
  vector<int> slack, match, st, pa, S, vis;
  vector<vector<int>> flo, flo_from;
  queue < int > q;
  WeightGraph(int n_)
    : n(n_{-}), nx(n * 2), lab(nx + 1),
      g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1),
      flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
    match = st = pa = S = vis = slack;
    rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
  int ED(edge e) {
    return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
  void update_slack(int u, int x, int &s) {
    if (!s || ED(g[u][x]) < ED(g[s][x])) s = u;</pre>
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)
      if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
        update_slack(u, x, slack[x]);
  void q_push(int x) {
    if (x <= n) q.push(x);</pre>
      for (int y : flo[x]) q_push(y);
  void set_st(int x, int b) {
    st[x] = b;
    if(x > n)
      for (int y : flo[x]) set_st(y, b);
  vector<int> split_flo(auto &f, int xr) {
    auto it = find(ALL(f), xr);
    if (auto pr = it - f.begin(); pr % 2 == 1)
      reverse(1 + ALL(f)), it = f.end() - pr;
    auto res = vector(f.begin(), it);
    return f.erase(f.begin(), it), res;
  void set_match(int u, int v) {
    match[u] = g[u][v].v;
    if (u <= n) return;
int xr = flo_from[u][g[u][v].u];</pre>
    auto &f = flo[u], z = split_flo(f, xr);
    rep(i, 0, (int)z.size() - 1)
     set_match(z[i], z[i ^ 1]);
    set_match(xr, v);
    f.insert(f.end(), all(z));
  void augment(int u, int v) {
    for (;;) {
      int xnv = st[match[u]];
      set_match(u, v);
      if (!xnv) return;
      set_match(xnv, st[pa[xnv]]);
      u = st[pa[xnv]], v = xnv;
    }
  int lca(int u, int v) {
    static int t = 0:
    ++t;
    for (++t; u || v; swap(u, v))
      if (u) {
        if (vis[u] == t) return u;
        vis[u] = t;
        u = st[match[u]];
        if (u) u = st[pa[u]];
```

```
return 0:
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + all(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0;
  match[b] = match[o];
  vector<int> f = {o};
  for (int x = u, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  reverse(1 + all(f));
  for (int x = v, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  flo[b] = f;
  set_st(b, b);
for (int x = 1; x <= nx; ++x)</pre>
    g[b][x].w = g[x][b].w = 0;
  fill(all(flo_from[b]), 0);
  for (int xs : flo[b]) {
    for (int x = 1; x <= nx; ++x)</pre>
      if (g[b][x].w == 0 ||
         ED(g[xs][x]) < ED(g[b][x]))
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  }
  set_slack(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) {
      xs = x;
      continue:
    pa[xs] = g[x][xs].u;
    S[xs] = 1, S[x] = 0;
slack[xs] = 0;
    set_slack(x);
    q_push(x);
    xs = -1;
  for (int x : flo[b])
    if (x == xr) S[x] = 1, pa[x] = pa[b];
    else S[x] = -1, set_slack(x);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
    int nu = st[match[v]];
    pa[v] = e.u;
    S[v] = 1;
    slack[v] = slack[nu] = 0;
    S[nu] = 0;
    q_push(nu);
    else if (S[v] == 0) {
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
  }
  return false;
bool matching() {
  fill(all(S), -1), fill(all(slack), 0);
  q = queue < int >();
  for (int x = 1; x <= nx; ++x)</pre>
    if (st[x] == x && !match[x])
  pa[x] = 0, S[x] = 0, q_push(x);
if (q.empty()) return false;
  for (;;) {
    while (q.size()) {
      int u = q.front();
      q.pop();
       if (S[st[u]] == 1) continue;
       for (int v = 1; v <= n; ++v)
         if (g[u][v].w > 0 && st[u] != st[v]) {
           if (ED(g[u][v]) != 0)
           update_slack(u, st[v], slack[st[v]]);
else if (on_found_edge(g[u][v]))
             return true;
         }
    int d = INF;
    for (int b = n + 1; b <= nx; ++b)</pre>
      if (st[b] == b && S[b] == 1)
```

```
d = min(d, lab[b] / 2);
       for (int x = 1; x <= nx; ++x)</pre>
          if (int s = slack[x];
              st[x] == x && s && s[x] <= 0)
            d = min(d, ED(g[s][x]) / (S[x] + 2));
       for (int u = 1; u <= n; ++u)
  if (S[st[u]] == 1) lab[u] += d;</pre>
          else if (S[st[u]] == 0) {
  if (lab[u] <= d) return false;</pre>
            lab[u] -= d;
       rep(b, n + 1, nx) if (st[b] == b \&\& S[b] >= 0)
          lab[b] += d * (2 - 4 * S[b]);
        for (int x = 1; x <= nx; ++x)</pre>
          if (int s = slack[x]; st[x] == x && s &&
              st[s] != x \&\& ED(g[s][x]) == 0)
            if (on_found_edge(g[s][x])) return true;
       for (int b = n + 1; b <= nx; ++b)</pre>
          if (st[b] == b && S[b] == 1 && lab[b] == 0)
            expand_blossom(b);
     return false;
   pair<ll, int> solve() {
     fill(all(match), 0);
     rep(u, 0, n) st[u] = u, flo[u].clear();
     int w_max = 0;
     rep(u, 1, n) rep(v, 1, n) {
   flo_from[u][v] = (u == v ? u : 0);
       w_{max} = max(w_{max}, g[u][v].w);
     fill(all(lab), w_max);
     int n_matches = 0;
     ll tot_weight = 0;
     while (matching()) ++n_matches;
     rep(u, 1, n) if (match[u] && match[u] < u)
       tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
   void add_edge(int u, int v, int w) {
     g[u][v].w = g[v][u].w = w;
};
```

# 3.5 GlobalMinCut

```
#undef INF
struct SW{ // global min cut, O(V^3)
  #define REP for (int i = 0; i < n; ++i)</pre>
  static const int MXN = 514, INF = 2147483647;
  int vst[MXN], edge[MXN][MXN], wei[MXN];
  void init(int n) {
    REP fill_n(edge[i], n, 0);
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  int search(int &s, int &t, int n){
    fill_n(vst, n, 0), fill_n(wei, n, 0);
    s = t = -1;
    int mx, cur;
    for (int j = 0; j < n; ++j) {</pre>
      mx = -1, cur = 0;
      REP if (wei[i] > mx) cur = i, mx = wei[i];
      vst[cur] = 1, wei[cur] = -1;
      s = t; t = cur;
REP if (!vst[i]) wei[i] += edge[cur][i];
    }
    return mx;
  int solve(int n) {
    int res = INF;
    for (int x, y; n > 1; n--){
       res = min(res, search(x, y, n));
       REP edge[i][x] = (edge[x][i] += edge[y][i]);
         edge[y][i] = edge[n - 1][i];
         edge[i][y] = edge[i][n - 1];
      } // edge[y][y] = 0;
    }
    return res;
  }
} sw;
```

# 3.6 BoundedFlow(Dinic)

```
struct BoundedFlow { // 0-base
  struct edge {
    int to, cap, flow, rev;
  vector<edge> G[N];
  int n, s, t, dis[N], cur[N], cnt[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n + 2; ++i)</pre>
      G[i].clear(), cnt[i] = 0;
  void add_edge(int u, int v, int lcap, int rcap) {
    cnt[u] -= lcap, cnt[v] += lcap;
    G[u].emplace_back
         (edge{v, rcap, lcap, (int)G[v].size()});
    G[v].emplace\_back
         (edge{u, 0, 0, (int)G[u].size() - 1});
  void add_edge(int u, int v, int cap) {
    G[u].emplace_back
         (edge{v, cap, 0, (int)G[v].size()});
    G[v].emplace_back
         (edge{u, 0, 0, (int)G[u].size() - 1});
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < G[u].size(); ++i) {</pre>
      edge &e = G[u][i];
      if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
         int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
           e.flow += df, G[e.to][e.rev].flow -= df;
           return df;
        }
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n + 3, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (edge &e : G[u])
        if (!~dis[e.to] && e.flow != e.cap)
          q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
    int sum = 0;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
      add_edge(n + 1, i, cnt[i]), sum += cnt[i];
else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);
    if (sum != maxflow(n + 1, n + 2)) sum = -1;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
        G[n + 1].pop_back(), G[i].pop_back();
      else if (cnt[i] < 0)</pre>
        G[i].pop_back(), G[n + 2].pop_back();
    return sum != -1;
  int solve(int _s, int _t) {
    add_edge(_t, _s, INF);
    if (!solve()) return -1; // invalid flow
    int x = G[_t].back().flow;
    return G[_t].pop_back(), G[_s].pop_back(), x;
  }
};
```

| MaxFlow Dinic;

```
3.7 GomoryHuTree
```

```
int g[N];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
for (int i = 1; i < n; ++i) {
     Dinic.reset();
     add_edge(i, g[i], Dinic.maxflow(i, g[i]));
for (int j = i + 1; j <= n; ++j)
        if (g[j] == g[i] && ~Dinic.dis[j])
           q[i] = i;
  }
}
```

#### 3.8 MinCostCirculation

```
struct MinCostCirculation { // 0-base
  struct Edge {
    ll from, to, cap, fcap, flow, cost, rev;
  } *past[N];
  vector < Edge > G[N];
  ll dis[N], inq[N], n;
  void BellmanFord(int s) {
     fill_n(dis, n, INF), fill_n(inq, n, 0);
     queue<int> q;
     auto relax = [&](int u, ll d, Edge *e) {
      if (dis[u] > d) {
        dis[u] = d, past[u] = e;
         if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, 0);
    while (!q.empty()) {
      int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u])
         if (e.cap > e.flow)
           relax(e.to, dis[u] + e.cost, &e);
    }
  }
  void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
     if (dis[cur.from] + cur.cost < 0) {</pre>
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
            i = cur.from; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         ++e.flow, --G[e.to][e.rev].flow;
      }
    }
    ++cur.cap;
  void solve(int mxlg) {
    for (int b = mxlg; b >= 0; --b) {
      for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
          e.cap *= 2, e.flow *= 2;
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].emplace_back(Edge{a, b,
          0, cap, 0, cost, (ll)G[b].size() + (a == b)});
    G[b].emplace back(Edge
         {b, a, 0, 0, 0, -cost, (ll)G[a].size() - 1});
} mcmf; // O(VE * ElogC)
```

#### 3.9 FlowModelsBuilding

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source  ${\cal S}$  and sink  ${\cal T}$  .
  - 2. For each edge (x,y,l,u), connect  $x \rightarrow y$  with capacity u-l.
  - 3. For each vertex  $\boldsymbol{v}$  , denote by  $in(\boldsymbol{v})$  the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S \to v$  with capacity in(v), otherwise, connect  $v \rightarrow T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from  $\boldsymbol{s}$  to  $\boldsymbol{t}$  is the answer.

- To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x,y) \in M$ ,  $x \rightarrow y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- Minimum cost cyclic flow
  - 1. Consruct super source  $\boldsymbol{S}$  and sink  $\boldsymbol{T}$
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap) = (c,1) if c>0, otherwise connect  $y \to x$  with (cost,cap) = (-c,1)
  - 3. For each edge with c<0 , sum these cost as K , then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0 , connect  $S\to v$  with  $(cost, cap)\,{=}\,(0, d(v))$
  - 5. For each vertex v with d(v)<0 , connect  $v\to T$  with  $(cost, cap)\,{=}\,(0, -d(v))$
  - 6. Flow from S to T , the answer is the cost of the flow  $C\!+\!K$
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer T
  - 2. Construct a max flow model, let  ${\cal K}$  be the sum of all weights
  - 3. Connect source  $s \! \to \! v$  ,  $v \! \in \! G$  with capacity K
  - 4. For each edge (u,v,w) in G , connect  $u\to v$  and  $v\to u$  with capacity w
  - 5. For  $v \in G$  , connect it with sink  $v \to t$  with capacity  $K + 2T (\sum_{e \in E(v)} w(e)) 2w(v)$
  - 6. T is a valid answer if the maximum flow  $f\!<\!K|V|$
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v' , and connect  $u' \to v'$  with weight w(u,v) .
  - 2. Connect  $v\to v'$  with weight  $2\mu(v)$  , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on  $G^{\prime}$  .
- Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ .
  - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
  - The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
  - 1. Capacity  $c_{uv}\,,$  Flow  $f_{uv}\,,$  Cost  $w_{uv}\,,$  Required Flow difference for vertex  $b_u\,.$
  - 2. If all  $\overset{.}{w}_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

$$\begin{split} \min & \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \ge -c_{uv} & \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{split}$$

# 4 Data Struture

#### 4.1 LCT

```
struct Splay { // xor-sum
  static Splay nil;
  Splay *ch[2], *f;
  int val, sum, rev, size;
Splay(int _val = 0)
    : val(_val), sum(_val), rev(0), size(1) {
    f = ch[0] = ch[1] = &nil;
  bool isr() {
    return f->ch[0] != this && f->ch[1] != this;
  int dir() { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void give_tag(int r) {
    if (r) swap(ch[0], ch[1]), rev ^= 1;
  void push() {
    if (ch[0] != &nil) ch[0]->give_tag(rev);
    if (ch[1] != &nil) ch[1]->give_tag(rev);
    rev = 0;
```

```
void pull() {
   // take care of the nil!
    size = ch[0] -> size + ch[1] -> size + 1;
    sum = ch[0] -> sum ^ ch[1] -> sum ^ val;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x -> f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d);
  p->pull(), x->pull();
void splay(Splay *x) {
  vector < Splay *> splayVec;
  for (Splay *q = x;; q = q->f) {
    splayVec.emplace_back(q);
    if (q->isr()) break;
  reverse(all(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir())
      rotate(x->f), rotate(x);
    else rotate(x), rotate(x);
 }
Splay *access(Splay *x) {
  Splay *q = nil;
  for (; x != nil; x = x->f)
    splay(x), x -> setCh(q, 1), q = x;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x) {
 root_path(x), x->give_tag(1);
  x->push(), x->pull();
void split(Splay *x, Splay *y) {
  chroot(x), root_path(y);
void link(Splay *x, Splay *y) {
 root_path(x), chroot(y);
  x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
  split(x, y);
  if (y->size != 5) return;
  y->push();
  y->ch[0] = y->ch[0]->f = nil;
Splay *get_root(Splay *x) {
  for (root_path(x); x->ch[0] != nil; x = x->ch[0])
   x->push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
  return get_root(x) == get_root(y);
Splay *lca(Splay *x, Splay *y) {
  access(x), root_path(y);
  if (y->f == nil) return y;
  return y->f;
void change(Splay *x, int val) {
  splay(x), x->val = val, x->pull();
int query(Splay *x, Splay *y) {
  split(x, y);
  return y->sum;
```

# 5 String

### 5.1 KMP

```
int KMP(string s, string t) {
  int n = t.size(), ans = 0;
  vector<int> f(t.size(), 0);
```

```
f[0] = -1:
  for (int i = 1, j = -1; i < t.size(); i++) {</pre>
    while (j >= 0)
      if (t[j + 1] == t[i]) break;
       else j = f[j];
    f[i] = ++j;
  for (int i = 0, j = 0; i < s.size(); i++) {</pre>
    while (j >= 0)
      if (t[j + 1] == s[i]) break;
       else j = f[j];
    if (++j + 1 == t.size()) ans++, j = f[j];
  return ans;
}
```

#### 5.2 Z

```
int Z[1000006];
void z(string s) {
  for (int i = 1, mx = 0; i < s.size(); i++) {
  if (i < Z[mx] + mx)</pre>
      Z[i] = min(Z[mx] - i + mx, Z[i - mx]);
    while (
      Z[i] + i < s.size() && s[i + Z[i]] == s[Z[i]])
      Z[i]++;
    if (Z[i] + i > Z[mx] + mx) mx = i;
 }
```

# 5.3 Manacher

```
int man[2000006];
int manacher(string s) {
  string t;
  for (int i = 0; i < s.size(); i++) {</pre>
    if (i) t.push_back('$');
    t.push_back(s[i]);
  int mx = 0, ans = 0;
  for (int i = 0; i < t.size(); i++) {</pre>
    man[i] = 1:
    man[i] = min(man[mx] + mx - i, man[2 * mx - i]);
    while (man[i] + i < t.size() && i - man[i] >= 0 &&
      t[i + man[i]] == t[i - man[i]])
      man[i]++;
    if (i + man[i] > mx + man[mx]) mx = i;
  for (int i = 0; i < t.size(); i++)</pre>
   ans = max(ans, man[i] - 1);
  return ans;
}
```

# 5.4 SuffixArray

```
vector<int> sa, cnt, rk, tmp, lcp;
void SA(string s) {
  int n = s.size();
  sa.resize(n), cnt.resize(n), rk.resize(n),
    tmp.resize(n):
  iota(all(sa), 0);
  sort(all(sa),
    [&](int i, int j) { return s[i] < s[j]; });
  rk[0] = 0;
  for (int i = 1; i < n; i++)</pre>
    rk[sa[i]] =
      rk[sa[i - 1]] + (s[sa[i - 1]] != s[sa[i]]);
  for (int k = 1; k <= n; k <<= 1) {</pre>
    fill(all(cnt), 0);
    for (int i = 0; i < n; i++)
  cnt[rk[(sa[i] - k + n) % n]]++;</pre>
    for (int i = 1; i < n; i++) cnt[i] += cnt[i - 1];</pre>
    for (int i = n - 1; i >= 0; i--)
       tmp[--cnt[rk[(sa[i] - k + n) % n]]] =
         (sa[i] - k + n) \% n;
    sa.swap(tmp);
    tmp[sa[0]] = 0;
    for (int i = 1; i < n; i++)</pre>
      tmp[sa[i]] = tmp[sa[i - 1]] +
         (rk[sa[i - 1]] != rk[sa[i]] ||
    rk[(sa[i - 1] + k) % n] !=
              rk[(sa[i] + k) % n]);
    rk.swap(tmp);
  }
}
```

```
void LCP(string s) {
  int n = s.size(), k = 0;
  lcp.resize(n);
  for (int i = 0; i < n; i++)</pre>
    if (rk[i] == 0) lcp[rk[i]] = 0;
    else
      if (k) k--;
      int j = sa[rk[i] - 1];
      while (
        i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
      lcp[rk[i]] = k;
```

### 5.5 SAIS

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N], _q[N * 2];
// zero based, string content MUST > 0
// SA[i]: SA[i]-th
     suffix is the i-th lexigraphically smallest suffix.
// H[i]: longest
     common prefix of suffix SA[i] and suffix SA[i - 1].
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, 0), copy_n(c, z, x); }
void induce
     (int *sa, int *c, int *s, bool *t, int n, int z) \{
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i)</pre>
    if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
  for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
       sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa
     , int *p, int *q, bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0,
      nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
    return:
  for (int i = n - 2; i >= 0; --i)
    t[i] = (
         s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i)
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
      bool neq = last < 0 || !equal
           (s + sa[i], s + p[q[sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns,
        nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i)
    sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
  induce(sa, c, s, t, n, z);
void mkhei(int n) {
  for (int i = 0, j = 0; i < n; ++i) {</pre>
    if (RA[i])
       for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
    H[RA[i]] = j, j = max(0, j - 1);
  }
}
void build(int *s, int n) {
 copy_n(s, n, _s), _s[n] = 0;
sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
copy_n(SA + 1, n, SA);
  for (int i = 0; i < n; ++i) RA[SA[i]] = i;</pre>
  mkhei(n);
```

#### 5.6 ACAutomaton

```
struct AhoCorasick {
  int ch[sumS][sigma] = {{}}, f[sumS] = {-1},
      tag[sumS], mv[sumS][sigma], jump[sumS],
      cnt[sumS];
  int idx = 0;
  int insert(string &s) {
    int j = 0;
    for (int i = 0; i < (int)s.size(); i++) {
  if (!ch[j][s[i] - base])</pre>
        ch[j][s[i] - base] = ++idx;
      j = ch[j][s[i] - base];
    tag[j] = 1;
    return j;
  int next(int u, int c) {
    return u < 0 ? 0 : mv[u][c];</pre>
  }
  void build() {
    queue < int > q;
    q.push(0);
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int v = 0; v < sigma; v++) {
        if (ch[u][v]) {
          f[ch[u][v]] = next(f[u], v);
           q.push(ch[u][v]);
        mv[u][v] =
           (ch[u][v] ? ch[u][v] : next(f[u], v));
      if (u) jump[u] = (tag[f[u]] ? f[u] : jump[f[u]]);
    }
  void match(string &s) {
    for (int i = 0; i <= idx; i++) cnt[i] = 0;</pre>
    for (int i = 0, j = 0; i < (int)s.size(); i++) {</pre>
      j = next(j, s[i] - base);
      cnt[j]++;
    vector<int> v;
    v.emplace_back(0);
    for (int i = 0; i < (int)v.size(); i++)</pre>
      for (int j = 0; j < sigma; j++)</pre>
        if (ch[v[i]][j]) v.emplace_back(ch[v[i]][j]);
    reverse(v.begin(), v.end());
    for (int i : v)
      if (f[i] > 0) cnt[f[i]] += cnt[i];
  }
} ac;
```

# 5.7 MinRotation

```
int mincyc(string s) {
  int n = s.size();
  s = s + s;
  int i = 0, ans = 0;
  while (i < n) {
    ans = i;
    int j = i + 1, k = i;
    while (j < s.size() && s[j] >= s[k]) {
        k = (s[j] > s[k] ? i : k + 1);
        ++j;
    }
    while (i <= k) i += j - k;
}
return ans;
}</pre>
```

### 5.8 ExtSAM

```
#define CNUM 26
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
    fill_n(next[tot], CNUM, 0);
    len[tot] = cnt[tot] = link[tot] = 0;
    return tot++;
}
```

```
void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
     int cur = next[last][c];
     len[cur] = len[last] + 1;
     int p = link[last];
    while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
     if (p == -1) return link[cur] = 0, cur;
     int q = next[p][c];
     if (len
         [p] + 1 == len[q]) return link[cur] = q, cur;
     int clone = newnode();
     for (int i = 0; i < CNUM; ++i)</pre>
      next[
           clone][i] = len[next[q][i]] ? next[q][i] : 0;
     len[clone] = len[p] + 1;
     while (p != -1 && next[p][c] == q)
       next[p][c] = clone, p = link[p];
     link[link[cur] = clone] = link[q];
     link[q] = clone;
     return cur;
  void insert(const string &s) {
     int cur = 0;
     for (auto ch : s) {
       int &nxt = next[cur][int(ch - 'a')];
       if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
  void build() {
     queue < int > q;
     q.push(0);
     while (!q.empty()) {
       int cur = q.front();
       q.pop();
       for (int i = 0; i < CNUM; ++i)</pre>
         if (next[cur][i])
           q.push(insertSAM(cur, i));
     vector<int> lc(tot);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
     partial_sum(all(lc), lc.begin());
     for (int i
         = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
   void solve() {
     for (int i = tot - 2; i >= 0; --i)
       cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
  }
};
```

# 5.9 PalindromeTree

```
struct palindromic tree {
  struct node {
    int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
    // pal. suf.
node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
    }
  vector < node > St;
  vector<char> s:
  int last, n;
  palindromic_tree() : St(2), last(1), n(0) {
    St[0].fail = 1, St[1].len = -1, s.emplace_back(-1);
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.emplace_back(0), St.emplace_back(-1);
    St[0].fail = 1, s.emplace_back(-1);
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
      x = St[x].fail;
    return x:
  inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = St.size();
      St.emplace_back(St[cur].len + 2);
      St[now].fail =
```

```
St[get_fail(St[cur].fail)].next[c];
St[cur].next[c] = now;
St[now].num = St[St[now].fail].num + 1;
}
last = St[cur].next[c], ++St[last].cnt;
}
inline void count() { // counting cnt
auto i = St.rbegin();
for (; i != St.rend(); ++i) {
   St[i->fail].cnt += i->cnt;
}
}
inline int size() { // The number of diff. pal.
   return (int)St.size() - 2;
}
};
```

# 6 Number Theory

# 6.1 Primes

#### 6.2 ExtGCD

```
// beware of negative numbers!
void extgcd(ll a, ll b, ll c, ll &x, ll &y) {
  if (b == 0) x = c / a, y = 0;
  else {
    extgcd(b, a % b, c, y, x);
    y -= x * (a / b);
  }
} // |x| <= b/2, |y| <= a/2</pre>
```

### 6.3 FloorCeil

```
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

#### 6.4 FloorSum

```
if (A == 0) return (N + 1) * (B / C);
if (A > C || B > C)
    return (N + 1) * (B / C) +
        N * (N + 1) / 2 * (A / C) +
        floorsum(A % C, B % C, C, N);
    ll M = (A * N + B) / C;
    return N * M - floorsum(C, C - B - 1, A, M - 1);
} // \sum^{n}_0 floor((ai + b) / m)
```

### 6.5 MillerRabin

```
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383 6 : primes <= 13
// n < 2^64 7 :
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool Miller_Rabin(ll a, ll n) {
   if ((a = a % n) == 0) return 1;
   if (n % 2 == 0) return n == 2;
   ll tmp = (n - 1) / ((n - 1) & (1 - n));
   ll t = __lg(((n - 1) & (1 - n))), x = 1;
   for (; tmp; tmp >>= 1, a = mul(a, a, n))
      if (tmp & 1) x = mul(x, a, n);
   if (x == 1 || x == n - 1) return 1;
   while (--t)
   if ((x = mul(x, x, n)) == n - 1) return 1;
   return 0;
}
```

#### 6.6 PollardRho

```
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2
       == 0) return PollardRho(n / 2), ++cnt[2], void();
  ll x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return:
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
    d = gcd(abs(x - y), n);
  }
}
```

#### 6.7 Fraction

```
ll n, d;
  fraction
       (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
     ll t = gcd(n, d);
    n /= t, d /= t;
    if (d < 0) n = -n, d = -d;
  fraction operator -() const
  { return fraction(-n, d); }
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator - (const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;
  }
};
```

### 6.8 ChineseRemainder

```
ll g = gcd(m1, m2);
if ((x2 - x1) % g) return -1; // no sol
m1 /= g; m2 /= g;
pll p = exgcd(m1, m2);
ll lcm = m1 * m2 * g;
ll res = p.first * (x2 - x1) * m1 + x1;
// be careful with overflow
return (res % lcm + lcm) % lcm;
}
```

# 6.9 Factorial $\mathsf{Mod} p^k$

```
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)
     if (i % p) prod[i] = prod[i - 1] * i % pk;
     else prod[i] = prod[i - 1];
  ll rt = 1;
  for (; n; n /= p) {
     rt = rt * mpow(prod[pk], n / pk, pk) % pk;
     rt = rt * prod[n % pk] % pk;
  }
  return rt;
} // (n! without factor p) % p^k</pre>
```

# 6.10 QuadraticResidue

```
ll trial(ll y, ll z, ll m) {
    ll a0 = 1, a1 = 0, b0 = z, b1 = 1, p = (m - 1) / 2;
   while (p) {
      if (p & 1)
        tie(a0, a1) =
           make_pair((a1 * b1 % m * y + a0 * b0) % m,
             (a0 * b1 + a1 * b0) % m);
      tie(b0, b1) =
        make_pair((b1 * b1 % m * y + b0 * b0) % m,
  (2 * b0 * b1) % m);
      p >>= 1;
   if (a1) return inv(a1, m);
   return -1;
 mt19937 rd(49);
 ll psqrt(ll y, ll p) {
  if (fpow(y, (p - 1) / 2, p) != 1) return -1;
  for (int i = 0; i < 30; i++) {</pre>
      ll z = rd() \% p;
      if (z * z % p == y) return z;
      ll x = trial(y, z, p);
      if (x == -1) continue;
      return x;
   return -1;
}
```

#### 6.11 MeisselLehmer

```
if (n <= 1) return 0;</pre>
int v = sqrt(n), s = (v + 1) / 2, pc = 0;
vector<int> smalls(v + 1), skip(v + 1), roughs(s);
vector<ll> larges(s);
for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;
for (int i = 0; i < s; ++i) {</pre>
  roughs[i] = 2 * i + 1;
  larges[i] = (n / (2 * i + 1) + 1) / 2;
for (int p = 3; p <= v; ++p) {
  if (smalls[p] > smalls[p - 1]) {
    int q = p * p;
    ++pc;
    if (1LL * q * q > n) break;
    skip[p] = 1;
    for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
    int ns = 0;
    for (int k = 0; k < s; ++k) {</pre>
      int i = roughs[k];
      if (skip[i]) continue;
ll d = 1LL * i * p;
       larges[ns] = larges[k] - (d <= v ? larges</pre>
           [smalls[d] - pc] : smalls[n / d]) + pc;
       roughs[ns++] = i;
    }
    s = ns;
    for (int j = v / p; j >= p; --j) {
      int c =
             smalls[j] - pc, e = min(j * p + p, v + 1);
      for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
    }
  }
for (int k = 1; k < s; ++k) {</pre>
  const ll m = n / roughs[k];
  ll t = larges[k] - (pc + \hat{k} - 1);
  for (int l = 1; l < k; ++l) {</pre>
    int p = roughs[l];
    if (1LL * p * p > m) break;
    t -= smalls[m / p] - (pc + l - 1);
  larges[0] -= t;
return larges[0];
```

# 6.12 DiscreteLog

```
constexpr int kStep = 32000;
  unordered_map<int, int> p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
    y = 1LL * y * x % m;
    b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;</pre>
    if (p.find(s) != p.end()) return i + kStep - p[s];
  }
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
s = 1LL * s * x % m;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
```

### 6.13 Theorems

• Cramer's rule

$$ax+by = e \qquad x = \frac{ed-bf}{ad-bc}$$

$$cx+dy = f \Rightarrow y = \frac{af-ec}{ad-bc}$$

• Vandermonde's Identity

$$C(n\!+\!m,\!k)\!=\!\sum_{i=0}^{k}\!C(n,\!i)C(m,\!k\!-\!i)$$

• Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}\!=\!d(i)$  ,  $L_{ij}\!=\!-c$  where c is the number of edge  $(i,\!j)$  in G .

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$  .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a n imes n matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$  , otherwise  $d_{ij} = -d_{ji}$  .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1, d_2, ..., d_n$  for each  $\emph{labeled}$  vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
  - Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\dots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$  .
- Erdős-Gallai theorem

A sequence of nonnegative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of a finite simple graph on  $\boldsymbol{n}$  vertices if

and only if 
$$d_1+\cdots+d_n$$
 is even and  $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$ 

holds for every  $1 \le k \le n$ .

· Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq \cdots \geq a_n$  and  $b_1, \ldots, b_n$ is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$ 

holds for every  $1 \le k \le n$ .

Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,\,b_1),\,...\,,(a_n,\,b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum a_i = \sum b_i$  and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

• Pick's theorem

For simple polygon, when points are all integer, we have  $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1.$ 

• Möbius inversion formula

- 
$$f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$$
  
-  $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$ 

- Spherical cap
  - A portion of a sphere cut off by a plane.
  - $r\colon$  sphere radius,  $a\colon$  radius of the base of the cap,  $h\colon$  height of the cap,  $\theta$ : arcsin(a/r).
  - Volume  $=\pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\theta)$  $\cos\theta)^2/3$ .
  - Area  $= 2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$ .
- Lagrange multiplier
  - Optimize  $f(x_1,...,x_n)$  when k constraints  $g_i(x_1,...,x_n)\!=\!0$  .
  - Lagrangian function  $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=$  $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n).$
  - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines
  - Line 1: $m{v}_1 = m{p}_1 + t_1 m{d}_1$
  - Line 2: ${m v}_2\!=\!{m p}_2\!+\!t_2{m d}_2$ -  $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
  - $\boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}$
  - $n_2 = d_2 \times n$
  - $c_1 = p_1 + \frac{(p_2 p_1) \cdot n_2}{d_1 \cdot n_2} d_1$
  - $c_2 = p_2 + \frac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2$

#### 6.14 Estimation

- Estimation
  - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
  - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30for  $n=0\sim 9$ , 627 for n=20,  $\sim 2e5$  for n=50,  $\sim 2e8$  for n=100.
  - Total number of partitions of  $\boldsymbol{n}$  distinct elements: B(n) = 1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437,190899322,....

#### 6.15 **EuclideanAlgorithms**

- $m = |\frac{an+b}{a}|$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \text{ mod } c,b \text{ mod } c,c,n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c,c-b-1,a,m-1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod{c}, b \bmod{c}, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c - b - 1, a, m - 1) \\ -h(c, c - b - 1, a, m - 1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c,b \bmod c,c,n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c,b \bmod c,c,n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c,b \bmod c,c,n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c,c-b-1,a,m-1) \\ - 2f(c,c-b-1,a,m-1) - f(a,b,c,n), & \text{otherwise} \end{cases} \end{split}$$

### 6.16 Numbers

Bernoulli numbers

ullet Stirling numbers of the second kind Partitions of n distinct elements into exactly  $\boldsymbol{k}$  groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), \\ S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n \\ x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem

ntagonal number theorem 
$$\prod_{n=1}^{\infty}(1-x^n)=1+\sum_{k=1}^{\infty}(-1)^k\Big(x^{k(3k+1)/2}+x^{k(3k-1)/2}\Big)$$
 talan numbers

• Catalan numbers 1

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1j:s s.t.  $\pi(j) \ge j$ , k j:s s.t.  $\pi(j) > j$ . E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)

$$E(n,0) = E(n,n-1) = 1$$
  
 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$ 

# 6.17 GeneratingFunctions

- Ordinary Generating Function  $A(x)\!=\!\sum_{i>0}\!a_ix^i$ 
  - $A(rx) \Rightarrow r^n a_n$
  - $A(x) + B(x) \Rightarrow a_n + b_n$
  - $A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
  - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $xA(x)' \Rightarrow na_n$
  - $\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function  $A(x) = \sum_{i \geq 0} \frac{a_i}{i!} x_i$ 
  - $A(x)+B(x) \Rightarrow a_n+b_n$
  - $A^{(k)}(x) \Rightarrow a_{n+k}$
  - $A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
  - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $xA(x) \Rightarrow na_n$
- Special Generating Function
  - $(1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i$  $-\frac{1}{(1-x)^n} = \sum_{i>0} {i \choose n-1} x^i$

- 
$$S_k = \sum_{x=1}^n x^k$$
:  $S = \sum_{p=0}^\infty x^p = \frac{e^x - e^{x(n+1)}}{1 - e^x}$ 

# Linear Algebra

# 7.1 GuassianElimination

```
int n, m;
   fraction M[MAXN][MAXN + 1], sol[MAXN];
  int solve() { //-1: inconsistent, >= 0: rank
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
       for (int j = 0; j < n; ++j) {</pre>
         if (i == j) continue;
         fraction tmp = -M[j][piv] / M[i][piv];
for (int k = 0; k <=</pre>
               m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
     int rank = 0;
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m && M[i][m].n) return -1;
       else if (piv
             < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
     return rank;
};
```

# 7.2 BerlekampMassey

```
vector<T> BerlekampMassey(const vector<T> &output) {
   vector<T> d(SZ(output) + 1), me, he;
   for (int f = 0, i = 1; i <= SZ(output); ++i) {</pre>
      for (int j = 0; j < SZ(me); ++j)
  d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;</pre>
      if (me.empty()) {
         me.resize(f = i);
          continue;
      vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.pb(-k);
       for (T x : he) o.pb(x * k);
      o.resize(max(SZ(o), SZ(me)));

for (int j = 0; j < SZ(me); ++j) o[j] += me[j];

if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
      me = o;
   return me;
}
```

#### 7.3 Simplex

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^{n} A_{ij} \bar{x}_j = b_j$  holds.

- 1. In case of minimization, let  $c_i'\!=\!-c_i$
- 2.  $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} A_{ji} x_i \leq -b_j$
- 3.  $\sum_{1 \le i \le n} A_{ji} x_i = b_j$   $\bullet \sum_{1 \le i \le n} A_{ji} x_i \le b_j$   $\bullet \sum_{1 \le i \le n} A_{ji} x_i \le b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

```
// infeasible < 0, unbounded = inf, Ax <= b, max
struct simplex {
  const double inf = 1 / .0, eps = 1e-9;
  int n, m = 0;
  double A[205][205], B[205];
  void init(int _n) { n = _n; }
void equation(vector<double> a, double b) {
    ++m;
    for (int i = 0; i < n; i++) A[m][i] = a[i];</pre>
    A[m][n + m] = 1, B[m] = b;
  double solve(vector<double> c) {
    for (int i = 0; i < n; i++) A[0][i] = -c[i];</pre>
    A[0][n] = 1;
    int flag = 1;
    while (flag--)
for (int i = 0; i <= n + m; i++)</pre>
```

```
if (A[0][i] < -eps) {</pre>
             double bx = inf;
             int piv = -1;
             for (int j = 1; j <= m; j++)</pre>
               if (0 <= A[j][i] && B[j] / A[j][i] <= bx)</pre>
             piv = j, bx = B[j] / A[j][i];
if (piv == -1) continue;
             if (bx == inf) return inf;
             flag = 1;
             for (int j = 0; j <= m; j++)</pre>
               if (j != piv) {
                 for (int k = 0; k <= n + m; k++)</pre>
                    if (k != i)
                      A[j][k] -=
                         A[piv][k] * A[j][i] / A[piv][i];
                 B[j] -= B[piv] * A[j][i] / A[piv][i];
                 A[j][i] = 0;
     for (int i = 0; i <= m; i++)
  if (B[i] < -eps) return -inf;</pre>
     return B[0] / A[0][n];
} lp;
```

# 8 Polynomials

# 8.1 NTT (FFT)

```
#define base ll // complex < double >
#define N 524288
// const double PI = acosl(-1);
const ll mod = 998244353, g = 3;
base omega[4 * N], omega_[4 * N];
int rev[4 * N];
ll fpow(ll b, ll p);
ll inverse(ll a) { return fpow(a, mod - 2); }
void calcW(int n) {
  ll r = fpow(g, (mod - 1) / n), invr = inverse(r);
  omega[0] = omega_[0] = 1;
  for (int i = 1; i < n; i++) {
  omega[i] = omega[i - 1] * r % mod;</pre>
    omega_[i] = omega_[i - 1] * invr % mod;
  // double arg = 2.0 * PI / n;
  // for (int i = 0; i < n; i++)
// {
  //
          omega[i] = base(cos(i * arg), sin(i * arg));
          omega_[i] = base(cos(-i * arg), sin(-i *
  //
          arg));
  // }
}
void calcrev(int n) {
 int k = __lg(n);
for (int i = 0; i < n; i++) rev[i] = 0;
for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < k; j++)</pre>
      if (i & (1 << j)) rev[i] ^= 1 << (k - j - 1);</pre>
vector<base> NTT(vector<base> poly, bool inv) {
  base *w = (inv ? omega_ : omega);
  int n = poly.size();
  for (int i = 0; i < n; i++)</pre>
    if (rev[i] > i) swap(poly[i], poly[rev[i]]);
  for (int len = 1; len < n; len <<= 1) {</pre>
    int arg = n / len / 2;
    for (int i = 0; i < n; i += 2 * len)</pre>
      for (int j = 0; j < len; j++) {</pre>
         base odd =
           w[j * arg] * poly[i + j + len] % mod;
         poly[i + j + len] :
           (poly[i + j] - odd + mod) % mod;
         poly[i + j] = (poly[i + j] + odd) \% mod;
  if (inv)
    for (auto &a : poly) a = a * inverse(n) % mod;
  return poly;
```

```
vector<base> mul(vector<base> f, vector<base> g) {
   int sz = 1 << (__lg(f.size() + g.size() - 1) + 1);
   f.resize(sz), g.resize(sz);
   calcrev(sz);
   calcW(sz);
   f = NTT(f, 0), g = NTT(g, 0);
   for (int i = 0; i < sz; i++)
      f[i] = f[i] * g[i] % mod;
   return NTT(f, 1);
}</pre>
```

#### 8.2 FHWT

```
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
  for (int i = 0; i < n; i += L)</pre>
       for (int j = i; j < i + (L >> 1); ++j)
          a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[
    N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
void
     subset_convolution(int *a, int *b, int *c, int L) {
   // c_k = \sum_{i=0}^{n} \{i \mid j = k, i \& j = 0\} a_i * b_j
  int n = 1 << L;
  for (int i = 1; i < n; ++i)</pre>
     ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
     f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)</pre>
     fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)</pre>
     for (int j = 0; j <= i; ++j)</pre>
       for (int x = 0; x < n; ++x)
h[i][x] += f[j][x] * g[i - j][x];</pre>
  for (int i = 0; i <= L; ++i)</pre>
  fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)
     c[i] = h[ct[i]][i];
```

# 8.3 PolynomialOperations

```
poly inv(poly A) {
   A.resize(1 << (__lg(A.size() - 1) + 1));</pre>
  poly B = {inverse(A[0])};
  for (int n = 1; n < A.size(); n += n) {</pre>
    poly pA(A.begin(), A.begin() + 2 * n);
    calcrev(4 * n);
    calcW(4 * n);
    pA.resize(4 * n);
    B.resize(4 * n);
    pA = NTT(pA, 0);
    B = NTT(B, 0);
    for (int i = 0; i < 4 * n; i++)</pre>
      B[i] =
        ((B[i] * 2 - pA[i] * B[i] % mod * B[i]) % mod +
          mod) %
        mod:
    B = NTT(B, 1);
   B.resize(2 * n);
  return B;
pair<poly, poly> div(poly A, poly B) {
  if (A.size() < B.size()) return make_pair(poly(), A);</pre>
  int n = A.size(), m = B.size();
  poly revA = A, invrevB = B;
  reverse(revA.begin(), revA.end());
  reverse(invrevB.begin(), invrevB.end());
  revA.resize(n - m + 1);
  invrevB.resize(n - m + 1);
  invrevB = inv(invrevB);
  poly Q = mul(revA, invrevB);
  Q.resize(n - m + 1);
  reverse(Q.begin(), Q.end());
  poly R = mul(Q, B);
  R.resize(m - 1);
```

```
for (int i = 0; i < m - 1; i++)
  R[i] = (A[i] - R[i] + mod) % mod;</pre>
  return make_pair(Q, R);
ll fast_kitamasa(ll k, poly A, poly C) {
  int n = A.size();
  C.emplace_back(mod - 1);
  poly Q, R = \{0, 1\}, F = \{1\};
  R = div(R, C);
  while (k) {
  if (k & 1) F = div(mul(F, R), C);
    R = div(mul(R, R), C);
  ll ans = 0;
  for (int i = 0; i < F.size(); i++)</pre>
    ans = (ans + A[i] * F[i]) % mod;
  return ans;
}
vector<ll> fpow(vector<ll> f, ll p, ll m) {
  int b = 0;
  while (b < f.size() && f[b] == 0) b++;</pre>
  f = vector<ll>(f.begin() + b, f.end());
  int n = f.size();
  f.emplace_back(0);
  vector<ll> q(min(m, b * p), 0);
  q.emplace_back(fpow(f[0], p));
  for (int k = 0; q.size() < m; k++) {</pre>
    ll res = 0;
    for (int i = 0; i < min(n, k + 1); i++)</pre>
      res = (res +
               p * (i + 1) % mod * f[i + 1] % mod *
                 q[k - i + b * p]) %
        mod;
    for (int i = 1; i < min(n, k + 1); i++)</pre>
      res = (res -
               f[i] * (k - i + 1) % mod *
                 q[k - i + 1 + b * p]) %
        mod;
    res = (res < 0 ? res + mod : res) *
      inv(f[0] * (k + 1) % mod) % mod;
    q.emplace_back(res);
  }
  return q;
```

### 8.4 NewtonMethod+MiscGF

Given  $F(\boldsymbol{x})$  where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0$  (mod  $x^{2^k}$ ), then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$

- $\bullet \ A^{-1} \colon \ B_{k+1} \!=\! B_k (2 \!-\! A B_k) \ \ \mathrm{mod} x^{2^{k+1}}$
- $\ln A$ :  $(\ln A)' = \frac{A'}{A}$
- $\bullet \ \exp\!A \colon \ B_{k+1} \!=\! B_k (1 \!+\! A \!-\! \ln\! B_k) \ \bmod x^{2^{k+1}}$
- $\sqrt{A}$ :  $B_{k+1} = \frac{1}{2}(B_k + AB_k^{-1}) \mod x^{2^{k+1}}$

# 9 Geometry

### 9.1 Basic

```
typedef pair < pdd, pdd > Line;
struct Cir{ pdd 0; double R; };
const double eps = 1e - 8;
pdd operator + (pdd a, pdd b)
{ return pdd(a.X + b.X, a.Y + b.Y); }
pdd operator - (pdd a, pdd b)
{ return pdd(a.X - b.X, a.Y - b.Y); }
pdd operator * (pdd a, double b)
{ return pdd(a.X * b, a.Y * b); }
pdd operator / (pdd a, double b)
{ return pdd(a.X / b, a.Y * b); }
double dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
double cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
```

```
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
{ return sign(cross(b - a, c - a)); }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
  if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= 0;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
    return btw(p1, p2, p3) || btw(p1, p2, p4) || btw(p3, p4, p1) || btw(p3, p4, p2);
  return a123 * a124 <= 0 && a341 * a342 <= 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
double a124 = cross(p2 - p1, p4 - p1);
  return (p4
      * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
    ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
    (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) \{
  pdd dp = p1 - p0
, dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  \} // from line p0--p1 to q0--q1, apply to r
```

## 9.2 ConvexHull

### 9.3 SortByAngle

```
#define is_neg(k) (
    sign(k.Y) < 0 || (sign(k.Y) == 0 && sign(k.X) < 0))
int A = is_neg(a), B = is_neg(b);
if (A != B)
    return A < B;
if (sign(cross(a, b)) == 0)
    return same ? abs2(a) < abs2(b) : -1;
return sign(cross(a, b)) > 0;
}
```

# 9.4 DisPointSegment

### 9.5 PointInCircle

```
bool in_cc(const array<pll, 3> &p, pll q) {
   __int128 det = 0;
   for (int i = 0; i < 3; ++i)
    det += __int128(abs2(p[i]) - abs2(q)) *
        cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);</pre>
```

```
return det > 0; // in: >0, on: =0, out: <0
```

# 9.6 PointInConvex

# 9.7 PointTangentConvex

```
return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
      { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

# 9.8 CircTangentCirc

```
// sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = abs2(c1.0 - c2.0);
  if (sign(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d;
  if (c * c > 1) return ret;
  double h = sqrt(max(0.0, 1.0 - c * c));
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
       v.Y * c + sign2 * h * v.X);
     pdd p1 = c1.0 + n * c1.R;
     pdd p2 = c2.0 + n * (c2.R * sign1);
     if (sign(p1.X - p2.X) == 0 and
        sign(p1.Y - p2.Y) == 0)
       p2 = p1 + perp(c2.0 - c1.0);
     ret.pb(Line(p1, p2));
  return ret;
}
```

#### 9.9 LineConvexIntersect

```
return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 || cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r; ) {
      int m = ((l + r + (l < r ? 0 : n)) / 2) % n;</pre>
      (cmpL(m) == cmpL(t) ? l : r) = m;
    return (l + !cmpL(r)) % n;
 pii res = pii(gao(B, A), gao(A, B)); // (i, j)
if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
  if (!
      cmpL(res.X) && !cmpL(res.Y)) // along side i, i+1
    switch ((res.X - res.Y + n + 1) % n) {
      case 0: return pii(res.X, res.X);
      case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned
       in the same order as the line hits the convex */
```

```
return res;
} // convex cut: (r, l]
```

### 9.10 CircIntersectCirc

```
pdd o1 = a.0, o2 = b.0;
double r1 =
        a.R, r2 = b.R, d2 = abs2(o1 - o2), d = sqrt(d2);
if(d < max
        (r1, r2) - min(r1, r2) || d > r1 + r2) return 0;
pdd u = (o1 + o2) * 0.5
        + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
double A = sqrt((r1 + r2 + d) *
        (r1 - r2 + d) * (r1 + r2 - d) * (-r1 + r2 + d));
pdd v
        = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A / (2 * d2);
p1 = u + v, p2 = u - v;
return 1;
}
```

# 9.11 PolyIntersectCirc

```
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r && B
         < PI/2) S -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S;
double area_poly_circle(const
     vector<pdd> poly,const pdd &0,const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=_area(poly[i]-0,poly[(i+1)%SZ(poly
         )]-0,r)*ori(0,poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
```

#### 9.12 MinMaxEnclosingRect

```
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
 hull(dots);
  double Max = 0, Min = INF, deg;
  int n = SZ(dots);
 dots.pb(dots[0]);
 for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {</pre>
   pll nw = vec(i + 1);
   while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
     u = (u + 1) \% n;
   while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
     r = (r + 1) \% n;
   if (!i) l = (r + 1) % n;
   while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))</pre>
     l = (l + 1) \% n;
   Min = min(Min, (double)(dot(nw, vec(r)) - dot
       (nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));
   deg = acos(dot(diff(r
       , l), vec(u)) / abs(diff(r, l)) / abs(vec(u)));
   deg = (qi - deg) / 2;
   return pdd(Min, Max);
```

### 9.13 MinEnclosingCircle

```
pdd cent;
random_shuffle(ALL(dots));
cent = dots[0], r = 0;
```

#### 9.14 CircleCover

```
struct CircleCover {
  int C;
  Cir c[N];
  bool g[N][N], overlap[N][N];
  // Area[i] : area covered by at least i circles
  double Area[ N ];
  void init(int _C){ C = _C;}
  struct Teve {
    pdd p; double ang; int add;
Teve() {}
    Teve(pdd _a
    , double _b, int _c):p(_a), ang(_b), add(_c){}
bool operator<(const Teve &a)const</pre>
    {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
  bool disjuct(Cir &a, Cir &b, int x)
{return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
  bool contain(int i, int j) {
    /* c[j] is non-strictly in c[i]. */
    return (sign
         (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
         R) == 0 && i < j)) && contain(c[i], c[j], -1);
  void solve(){
    fill_n(Area, C + 2, 0);
    for(int i = 0; i < C; ++i)</pre>
       for(int j = 0; j < C; ++j)</pre>
         overlap[i][j] = contain(i, j);
    for(int i = 0; i < C; ++i)
  for(int j = 0; j < C; ++j)</pre>
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
              disjuct(c[i], c[j],
                                     -1));
    for(int i = 0; i < C; ++i){</pre>
       int E = 0, cnt = 1;
       for(int j = 0; j < C; ++j)</pre>
         if(j != i && overlap[j][i])
            ++cnt;
       for(int j = 0; j < C; ++j)</pre>
         if(i != j && g[i][j]) {
  pdd aa, bb;
           CCinter(c[i], c[j], aa, bb);
           double A =
                 atan2(aa.Y - c[i].O.Y, aa.X - c[i].O.X);
           double B =
                 atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
           eve[E++] = Teve
                (bb, B, 1), eve[E++] = Teve(aa, A, -1);
           if(B > A) ++cnt;
       if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
       else{
         sort(eve, eve + E);
         eve[E] = eve[0];
         for(int j = 0; j < E; ++j){</pre>
           cnt += eve[j].add;
           Area[cnt
                ] += cross(eve[j].p, eve[j + 1].p) * .5;
           double theta = eve[j + 1].ang - eve[j].ang;
           if (theta < 0) theta += 2. * pi;
Area[cnt] += (theta</pre>
                  - sin(theta)) * c[i].R * c[i].R * .5;
      }
    }
  }
```

# 9.15 Trapezoidalization

};

```
struct SweepLine {
  struct cmp {
    cmp(const SweepLine &_swp): swp(_swp) {}
    bool operator()(int a, int b) const -
      if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
        return swp.slope_cmp(a, b);
      return swp.get_y(a) + swp.eps < swp.get_y(b);</pre>
    }
    const SweepLine &swp;
  } _cmp;
  T curTime, eps, curQ;
  vector < Line > base;
  multiset<int, cmp> sweep;
  multiset<pair<T, int>> event;
  vector<typename multiset<int, cmp>::iterator> its;
  vector
      <typename multiset<pair<T, int>>::iterator> eits;
  bool slope_cmp(int a, int b) const {
    assert(a != -1);
    if (b == -1) return 0;
    return sign(cross(base
        [a].Y - base[a].X, base[b].Y - base[b].X)) < 0;</pre>
    get_y(int idx) const {
    if (idx == -1) return curQ;
    Line l = base[idx];
    if (l.X.X == l.Y.X) return l.Y.Y;
    return ((curTime - l.X.X) * l.Y.Y
+ (l.Y.X - curTime) * l.X.Y) / (l.Y.X - l.X.X);
  void insert(int idx) {
    its[idx] = sweep.insert(idx);
    if (its[idx] != sweep.begin())
      update_event(*prev(its[idx]));
    update_event(idx);
    event.emplace(base[idx].Y.X, idx + 2 * SZ(base));
  void erase(int idx) {
    assert(eits[idx] == event.end());
    auto p = sweep.erase(its[idx]);
    its[idx] = sweep.end();
    if (p != sweep.begin())
      update_event(*prev(p));
  void update_event(int idx) {
    if (eits[idx] != event.end())
      event.erase(eits[idx]);
    eits[idx] = event.end();
    auto nxt = next(its[idx]);
    if (nxt ==
         sweep.end() || !slope_cmp(idx, *nxt)) return;
    auto t = intersect(base[idx].
        X, base[idx].Y, base[*nxt].X, base[*nxt].Y).X;
    if (t + eps < curTime || t</pre>
          >= min(base[idx].Y.X, base[*nxt].Y.X)) return;
    eits[idx] = event.emplace(t, idx + SZ(base));
  void swp(int idx) {
    assert(eits[idx] != event.end());
    eits[idx] = event.end();
    int nxt = *next(its[idx]);
    swap((int&)*its[idx], (int&)*its[nxt]);
    swap(its[idx], its[nxt]);
    if (its[nxt] != sweep.begin())
      update_event(*prev(its[nxt]));
    update_event(idx);
  // only expected to call the functions below
  SweepLine(T t, T e, vector
      <Line> vec): _cmp(*this), curTime(t), eps(e)
      , curQ(), base(vec), sweep(_cmp), event(), its(SZ
(vec), sweep.end()), eits(SZ(vec), event.end()) {
    for (int i = 0; i < SZ(base); ++i) {</pre>
      auto &[p, q] = base[i];
      if (p > q) swap(p, q);
      if (p.X <= curTime && curTime <= q.X)</pre>
        insert(i);
      else if (curTime < p.X)</pre>
        event.emplace(p.X, i);
    }
  void setTime(T t, bool ers = false) {
```

```
assert(t >= curTime):
     while (!event.empty() && event.begin()->X <= t) {
  auto [et, idx] = *event.begin();</pre>
       int s = idx / SZ(base);
       idx %= SZ(base);
       if (abs(et - t) <= eps && s == 2 && !ers) break;</pre>
       curTime = et:
       event.erase(event.begin());
       if (s == 2) erase(idx);
       else if (s == 1) swp(idx);
       else insert(idx);
    }
     curTime = t;
   T nextEvent() {
     if (event.empty()) return INF;
     return event.begin()->X;
   int lower_bound(T y) {
    curQ = y;
     auto p = sweep.lower_bound(-1);
     if (p == sweep.end()) return -1;
     return *p;
};
```

# 9.16 TriangleHearts

```
p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
  double m = 2. * (x1 * y2 - y1 * x2);
 center.X = (x1 * x1)
       * y2 - x2 * x2 * y1 + y1 * y2 * (y1 - y2)) / m;
 center.Y = (x1 * x2)
      * (x2 - x1) - y1 * y1 * x2 + x1 * y2 * y2) / m;
  return center + p0;
pdd incenter
    (pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
  double a =
      abs(p2 - p3), b = abs(p1 - p3), c = abs(p1 - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter
    (p1, p2, p3) * 3 - circenter(p1, p2, p3) * 2; }
```

### 9.17 HalfPlaneIntersect

```
{ return pll(cross(a.Y
      bool isin(Line l0, Line l1, Line l2)
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(l0, l2);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;</pre>
  return (__int128)
a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(ALL(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
    return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  }):
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(
         dq.back().Y - dq.back().X, p.Y - p.X, 0) == -1)
       continue:
    while (SZ(dq
         ) >= 2 && !isin(p, dq[SZ(dq) - 2], dq.back()))
       dq.pop_back();
    while (SZ(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
      dq.pop_front();
    dq.pb(p);
  while (SZ(dq)
       >= 3 && !isin(dq[0], dq[SZ(dq) - 2], dq.back()))
    dq.pop_back();
  while (SZ(dq) >= 3 && !isin(dq.back(), dq[0], dq[1]))
```

```
dq.pop_front();
return vector<Line>(ALL(dq));
}
```

# 9.18 RotatingSweepLine

```
int n = SZ(ps), m = 0;
   vector<int> id(n), pos(n);
   vector<pii> line(n * (n - 1));
   for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++j)</pre>
       if (i != j) line[m++] = pii(i, j);
   sort(ALL(line), [&](pii a, pii b) {
     return cmp(ps[a.Y] - ps[a.X], ps[b.Y] - ps[b.X]);
   }); // cmp(): polar angle compare
   iota(ALL(id), 0);
   sort(ALL(id), [&](int a, int b) {
  if (ps[a].Y != ps[b].Y) return ps[a].Y < ps[b].Y;</pre>
     return ps[a] < ps[b];</pre>
   \}); // initial order, since (1, 0) is the smallest
   for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
   for (int i = 0; i < m; ++i) {</pre>
     auto l = line[i];
     // do something
     tie(pos[l.X], pos[l.Y], id[pos[l.X]], id[pos[l.Y
          ]]) = make_tuple(pos[l.Y], pos[l.X], l.Y, l.X);
   }
}
```

# 9.19 DelaunayTriangulation

```
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)%3], u.p[(i+2)%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const
     ll inf = MAXC * MAXC * 100; // lower_bound unknown
struct Tri;
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
  pll p[3];
  Edge edge[3];
Tri* chd[3];
  Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)
      if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
       return 0;
    return 1;
pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
    the_root
         = // Tri should at least contain all points
      new(tris++) Tri(pll(-inf, -inf),
           pll(inf + inf, -inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
  void add point(const
       pll &p) { add_point(find(the_root, p), p); }
  Tri* the_root;
```

```
static Tri* find(Tri* root, const pll &p) {
    while (1) {
       if (!root->has_chd())
         return root;
       for (int i = 0; i < 3 && root->chd[i]; ++i)
         if (root->chd[i]->contains(p)) {
           root = root->chd[i];
           break;
    assert(0); // "point not found"
  void add_point(Tri* root, pll const& p) {
    Tri* t[3];
      /* split it into three triangles */
    for (int i = 0; i < 3; ++i)</pre>
       t[i] = new(tris)
           ++) Tri(root->p[i], root->p[(i + 1) % 3], p);
     for (int i = 0; i < 3; ++i)</pre>
       edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)</pre>
       edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
     for (int i = 0; i < 3; ++i)</pre>
       root->chd[i] = t[i];
     for (int i = 0; i < 3; ++i)
       flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p
         [0], tri->p[1], tri->p[2], trj->p[pj])) return;
     /* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri
         (tri->p[(pi + 1) % 3], trj->p[pj], tri->p[pi]);
     Tri* trl = new(tris++) Tri
         (trj->p[(pj + 1) % 3], tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
     edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd
         [0] = trk; tri->chd[1] = trl; tri->chd[2] = 0;
    trj->chd
         [0] = trk; trj->chd[1] = trl; trj->chd[2] = 0;
    flip(trk, 1); flip(trk, 2);
flip(trl, 1); flip(trl, 2);
  }
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
  if (vst.find(now) != vst.end())
    return:
  vst.insert(now);
  if (!now->has_chd())
  return triang.pb(now);
for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(ps, ps + n);
  Trig tri; // the triangulation structure for (int i = 0; i < n; ++i)
    tri.add_point(ps[i]);
  go(tri.the_root);
}
```

### 9.20 VonoroiDiagram

```
vector<vector<Line>> vec;
void build_voronoi_line(int n, pll *arr) {
  tool.init(n, arr); // Delaunay
  vec.clear(), vec.resize(n);
  for (int i = 0; i < n; ++i)</pre>
    for (auto e : tool.head[i]) {
      int u = tool.oidx[i], v = tool.oidx[e.id];
      pll m = (arr[v]
           ] + arr[u]) / 2LL, d = perp(arr[v] - arr[u]);
      vec[u].pb(Line(m, m + d));
}
```

# 9.21 3DConeDistance

```
const double eps = 1e-12, pi = acos(-1);
struct point {
 ld x, y, z;
};
ld r, h;
bool on_edge(point p) {
  return abs(p.z) < eps &&
    abs(p.x * p.x + p.y * p.y - r * r) < eps;
}
bool on_bottom(point p) { return abs(p.z) < eps; }</pre>
ld straight(point a, point b) {
  return sqrtl((a.x - b.x) * (a.x - b.x) +
    (a.y - b.y) * (a.y - b.y) +
    (a.z - b.z) * (a.z - b.z));
ld arc(point a, point b) {
  if (abs(a.x) < eps && abs(a.y) < eps &&</pre>
    abs(a.z - h) < eps)
    return straight(a, b);
  swap(a, b);
  if (abs(a.x) < eps && abs(a.y) < eps &&</pre>
    abs(a.z - h) < eps)
    return straight(a, b);
  swap(a, b);
  ld\ ans = 1e18;
  for (int t = 0; t <= 1; t++) {</pre>
    double alpha = (a.x < 0 ? pi + atan2l(-a.y, -a.x)
                              : atan2l(a.y, a.x));
    if (alpha < 0) alpha += 2 * pi;</pre>
    alpha *= r / sqrtl(h * h + r * r);
    double lA = sqrtl(
      a.x * a.x + a.y * a.y + (a.z - h) * (a.z - h));
    double beta = (b.x < 0 ? pi + atan2l(-b.y, -b.x)
                            : atan2l(b.y, b.x));
    if (beta < 0) beta += 2 * pi;</pre>
    beta *= r / sqrtl(h * h + r * r);
    double lB = sqrtl(
      b.x * b.x + b.y * b.y + (b.z - h) * (b.z - h));
    double Ax = lA * cosl(alpha),
           Ay = lA * sinl(alpha), Bx = lB * cosl(beta),
           By = lB * sinl(beta);
    ans = min(ans,
      sqrtl((Ax - Bx) * (Ax - Bx) +
        (Ay - By) * (Ay - By)));
    a.x = -a.x;
    b.x = -b.x;
  }
  return ans;
ld distance(point a, point b) {
  if (on_bottom(a) && !on_edge(a) && !on_bottom(b))
    return 1e18;
  if (on_bottom(b) && !on_edge(b) && !on_bottom(a))
    return 1e18:
  if (on_bottom(a) && on_bottom(b))
    return straight(a, b);
  return arc(a, b);
}
const int s = 90;
ld ans = 1e18, theta = pi / 180.0;
point a, b;
vector<point> p;
vector<ld> ar;
vector<ld> dis, from, center = {0.0, pi};
vector<int> vis:
void iterate() {
  p = {a, b};
  ar = {-1e9, -1e9};
  dis.clear();
  from.clear();
  vis.clear();
```

```
for (auto d : center)
     for (int i = -s; i <= s; i++) {
   p.emplace_back(point{r * cosl(d + i * theta),}</pre>
         r * sinl(d + i * theta), 0});
       ar.emplace_back(d + i * theta);
  int n = p.size();
  dis.resize(n);
  from.resize(n);
  vis.resize(n);
  for (int i = 1; i < n; i++) dis[i] = 1e18;</pre>
  for (int t = 0; t < n; t++) {</pre>
     int u = 0;
     ld d = 1e18;
     for (int i = 0; i < n; i++)</pre>
       if (!vis[i] && dis[i] < d) u = i, d = dis[i];</pre>
     vis[u] = 1;
     for (int v = 0; v < n; v++)
       if (d + distance(p[u], p[v]) < dis[v]) 
         from[v] = u;
         dis[v] = distance(p[u], p[v]) + d;
  }
  ans = min(ans, dis[1]);
  center.clear();
   if (from[1] > 1) center.emplace_back(ar[from[1]]);
  if (from[from[1]] > 1)
     center.emplace_back(ar[from[from[1]]]);
  theta /= 20.0;
double cone_distance(
  double r, double h, point a, point b) {
if (on_bottom(a) && on_bottom(b))
     return straight(a, b);
  else {
     for (int t = 1; t <= 10; t++) iterate();</pre>
     return ans;
}
```

### 10 Misc

# 10.1 MoAlgoWithModify

```
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
  int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
   L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
  bool operator<(const Query &q) const {</pre>
    if (LBid != q.LBid) return LBid < q.LBid;
if (RBid != q.RBid) return RBid < q.RBid;</pre>
    return T < b.T;
void solve(vector<Query> query) {
  sort(ALL(query));
  int L=0, R=0, T=-1;
  for (auto q : query) {
    while (T < q.T) addTime(L, R, ++T); // TODO
    while (T > q.T) subTime(L, R, T--); // TODO
    while (R < q.R) add(arr[++R]); // TODO
    while (L > q.L) add(arr[--L]); // TODO
    while (R > q.R) sub(arr[R--]); // TODO
    while (L < q.L) sub(arr[L++]); // TODO
    // answer query
```

#### 10.2 MoAlgoOnTree

```
| Mo's Algorithm On Tree
| Preprocess:
| 1) LCA
| 2) dfs with in[u] = dft++, out[u] = dft++
| 3) ord[in[u]] = ord[out[u]] = u
| 4) bitset<MAXN> inset
| */
| struct Query {
| int L, R, LBid, lca;
```

```
Query(int u, int v) {
    int c = LCA(u, v);
     if (c == u || c == v)
      q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
    else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
    else
      q.lca = c, q.L = out[v], q.R = in[u];
    q.Lid = q.L / blk;
  bool operator < (const Query &q) const {
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    return R < q.R;
}:
void flip(int x) {
   if (inset[x]) sub(arr[x]); // TODO
    else add(arr[x]); // TODO
    inset[x] = ~inset[x];
void solve(vector<Query> query) {
  sort(ALL(query));
  int L = 0, R = 0;
  for (auto q : query) {
    while (R < q.R) flip(ord[++R]);
    while (L > q.L) flip(ord[--L]);
    while (R > q.R) flip(ord[R--]);
    while (L < q.L) flip(ord[L++]);
    if (~q.lca) add(arr[q.lca]);
    // answer query
    if (~q.lca) sub(arr[q.lca]);
}
```

# 10.3 MoAlgoAdvanced

- Mo's Algorithm With Addition Only
  - Sort querys same as the normal Mo's algorithm.
  - For each query [l,r]:
  - If l/blk = r/blk, brute-force.
  - If  $l/blk \neq curL/blk$ , initialize  $curL := (l/blk+1) \cdot blk, curR := curL-1$
  - If  $r\!>\!cur R$ , increase cur R
  - decrease  $\operatorname{\it cur} L$  to fit l , and then undo after answering
- Mo's Algorithm With Offline Second Time
  - Require: Changing answer  $\equiv$  adding f([l,r],r+1).
  - Require:  $f([l,r],r+1) = f([1,r],r+1) \hat{f}([1,l),r+1)$ . Part1: Answer all f([1,r],r+1) first.
  - Part2: Store  $curR \to R$  for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
  - Note: You must do the above symmetrically for the left boundaries.

### 10.4 HilbertCurve

```
ll res = 0;
for (int s = n / 2; s; s >>= 1) {
   int rx = (x & s) > 0;
   int ry = (y & s) > 0;
   res += s * 1ll * s * ((3 * rx) ^ ry);
   if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
   }
}
return res;
} // n = 2^k
```

#### 10.5 SternBrocotTree

- Construction: Root  $\frac{1}{1}$ , left/right neighbor  $\frac{0}{1},\frac{1}{0}$ , each node is sum of last left/right neighbor:  $\frac{a}{b},\frac{c}{d}\to\frac{a+c}{b+d}$
- Property: Adjacent (mid-order DFS)  $\frac{a}{b}, \frac{c}{d} \Rightarrow bc ad = 1$ .
- Search known  $\frac{p}{q}$ : keep L-R alternative. Each step can calcaulated in O(1)  $\Rightarrow$  total  $O(\log C)$ .
- Search unknown  $\frac{p}{q}$ : keep L-R alternative. Each step can calcaulated in  $O(\log C)$  checks  $\Rightarrow$  total  $O(\log^2 C)$  checks.

# 10.6 CyclicLCS

```
#define LU 1
#define U 2
const int mov[3][2] = {0, -1, -1, -1, -1, 0};
int al, bl;
char a[MAXL * 2], b[MAXL * 2]; // 0-indexed
int dp[MAXL * 2][MAXL];
char pred[MAXL * 2][MAXL];
inline int lcs_length(int r) {
  int i = r + al, j = bl, l = 0;
  while (i > r) {
    char dir = pred[i][j];
    if (dir == LU) l++;
    i += mov[dir][0];
    j += mov[dir][1];
  }
  return l;
inline void reroot(int r) { // r = new base row
  int i = r, j = 1;
  while (j <= bl && pred[i][j] != LU) j++;</pre>
  if (j > bl) return;
  pred[i][j] = L;
while (i < 2 * al && j <= bl) {</pre>
    if (pred[i + 1][j] == U) {
      i++:
      pred[i][j] = L;
    } else if (j < bl && pred[i + 1][j + 1] == LU) {</pre>
      i++:
       j++;
      pred[i][j] = L;
    } else {
      j++;
    }
  }
int cyclic_lcs() {
  // a, b, al, bl should be properly filled
  // note: a WILL be altered in process
              -- concatenated after itself
  char tmp[MAXL];
  if (al > bl) {
    swap(al, bl);
    strcpy(tmp, a);
    strcpy(a, b);
strcpy(b, tmp);
  strcpy(tmp, a);
  strcat(a, tmp);
  // basic lcs
  for (int i = 0; i <= 2 * al; i++) {</pre>
    dp[i][0] = 0;
    pred[i][0] = U;
  for (int j = 0; j <= bl; j++) {</pre>
    dp[0][j] = 0;
    pred[0][j] = L;
  for (int i = 1; i <= 2 * al; i++) {</pre>
    for (int j = 1; j <= bl; j++) {
  if (a[i - 1] == b[j - 1])</pre>
       dp[i][j] = dp[i - 1][j - 1] + 1;
else dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
      if (dp[i][j - 1] == dp[i][j]) pred[i][j] = L;
else if (a[i - 1] == b[j - 1]) pred[i][j] = LU;
       else pred[i][j] = U;
    }
  }
  // do cyclic lcs
  int clcs = 0;
  for (int i = 0; i < al; i++) {</pre>
    clcs = max(clcs, lcs_length(i));
    reroot(i + 1);
  // recover a
  a[al] = '\0';
  return clcs;
```

# 10.7 AllLCS

```
vector < int > h(SZ(t));
iota(ALL(h), 0);
for (int a = 0; a < SZ(s); ++a) {
  int v = -1;</pre>
```

```
for (int c = 0; c < SZ(t); ++c)
    if (s[a] == t[c] || h[c] < v)
        swap(h[c], v);
    // LCS(s[0, a], t[b, c]) =
    // c - b + 1 - sum([h[i] >= b] | i <= c)
    // h[i] might become -1 !!
}
</pre>
```

# 10.8 SimulatedAnnealing

```
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
    if (exp(-(nw - ans
        ) / factor) >= (double)(rnd() % base) / base)
        ans = nw;
    factor *= 0.99995;
}
```

# 10.9 Python

```
math.isqrt(2) # integer sqrt
```