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```

Basic

1.1 .vimrc

```
set ru nu cin cul sc so=3 ts=4 sw=4 bs=2 ls=2 mouse=a
ino {<CR> {<CR>}<C-o>0
ino jj <esc>
ino jk <esc>
map <F7> :w<CR>:!g++
    "%" -std=c++17 -Wall -Wextra -Wshadow -Wconversion
     -fsanitize=address,undefined -g && ./a.out<CR>
ca Hash w !cpp -dD -P -fpreprocessed
     \| tr -d "[:space:]" \| md5sum \| cut -c-6
```

1.2 PBDS

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
tree<int, null_type, less<int>, rb_tree_tag
     tree_order_statistics_node_update> bst;
// order_of_key(n): # of elements <= n</pre>
// find_by_order(n): 0-indexed
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
anu pbds::priority queue
    <int, greater<int>, thin_heap_tag> pq;
```

1.3 pragma

14

```
#pragma GCC optimize("Ofast,unroll-loops")
#pragma GCC target("avx,avx2,sse,sse2
    ,sse3,ssse3,sse4,popcnt,abm,mmx,fma,tune=native")
// chrono
    ::steady_clock::now().time_since_epoch().count()
```

Graph

2.1 2SAT/SCC

```
struct SAT { // O-base
   int low[N], dfn[N], bln[N], n, Time, nScc;
   bool instack[N], istrue[N];
   stack<int> st;
   vector<int> G[N], SCC[N];
   void init(int _n) {
     n = _n; // assert(n * 2 <= N);
     for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
   void add_edge(int a, int b) { G[a].emplace_back(b); }
   int rv(int a) {
     if (a >= n) return a - n;
     return a + n;
   void add_clause(int a, int b) {
     add_edge(rv(a), b), add_edge(rv(b), a);
   void dfs(int u) {
     dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
       do {
         tmp = st.top(), st.pop();
         instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != u);
       ++nScc;
     }
   bool solve() {
     Time = nScc = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;
     for (int i = 0; i < n + n; ++i)</pre>
       if (!dfn[i]) dfs(i);
     for (int i =
         0; i < n + n; ++i) SCC[bln[i]].emplace_back(i);
     for (int i = 0; i < n; ++i) {</pre>
       if (bln[i] == bln[i + n]) return false;
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true;
  }
};
```

2.2 BCC Vertex

```
int n, m, dfn[N], low[N], is_cut[N], nbcc = 0, t = 0;
vector<int> g[N], bcc[N], G[2 * N];
stack<int> st;
void tarjan(int p, int lp) {
  dfn[p] = low[p] = ++t;
  st.push(p);
  for (auto i : g[p]) {
    if (!dfn[i]) {
      tarjan(i, p);
      low[p] = min(low[p], low[i]);
      if (dfn[p] <= low[i]) {</pre>
        nbcc++;
```

```
is_cut[p] = 1;
         for (int x = 0; x != i; st.pop()) {
           x = st.top();
           bcc[nbcc].push_back(x);
         bcc[nbcc].push_back(p);
    } else low[p] = min(low[p], dfn[i]);
  }
}
void build() { // [n+1,n+nbcc] cycle, [1,n] vertex
  for (int i = 1; i <= nbcc; i++) {</pre>
    for (auto j : bcc[i]) {
       G[i + n].push_back(j);
       G[j].push_back(i + n);
  }
| }
```

2.3 BCC Edge

```
vector<int> tim, low;
     stack<int, vector<int>> st;
     int t, bcc_id;
     void dfs(int u, int p, const vector<vector</pre>
         <pair<int, int>>> &edge, vector<int> &pa) {
         tim[u] = low[u] = t++;
         st.push(u);
         for (const auto &[v, id] : edge[u]) {
             if (id == p)
                 continue;
             if (tim[v])
                 low[u] = min(low[u], tim[v]);
                 dfs(v, id, edge, pa);
                 if(low[v] > tim[u]) {
                     int x;
                     do {
                          pa[x = st.top()] = bcc_id;
                          st.pop();
                     } while (x != v);
                     bcc_id++;
                 else
                     low[u] = min(low[u], low[v]);
             }
         }
     vector<int> solve(const vector
         <vector<pair<int, int>>> &edge) { // (to, id)
         int n = edge.size();
         tim.resize(n);
         low.resize(n);
         t = bcc_id = 1;
         vector<int> pa(n);
         for (int i = 0; i < n; i++) {</pre>
             if (!tim[i]) {
                 dfs(i, -1, edge, pa);
                 while (!st.empty()) {
                     pa[st.top()] = bcc_id;
                     st.pop();
                 bcc_id++;
             }
         return pa;
    } // return bcc id(start from 1)
};
```

2.4 MinimumMeanCycle

```
| /* 0(V^3)
let dp[i][j] = min length from 1 to j exactly i edges
ans = min (dp[n + 1][u] - dp[i][u]) / (n + 1 - i) */
```

2.5 MaximumCliqueDyn

```
struct MaxClique { // fast when N <= 100</pre>
   bitset<N> G[N], cs[N];
   int ans, sol[N], q, cur[N], d[N], n;
   void init(int _n) {
    n = _n;
     for (int i = 0; i < n; ++i) G[i].reset();</pre>
   void add_edge(int u, int v) {
     G[v][v] = G[v][v] = 1;
   void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
     if (1 < 4) {
       for (int i : r) d[i] = (G[i] & mask).count();
       sort(all(r)
           , [&](int x, int y) { return d[x] > d[y]; });
     vector<int> c(r.size());
     int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
     cs[1].reset(), cs[2].reset();
     for (int p : r) {
       int k = 1;
       while ((cs[k] & G[p]).any()) ++k;
       if (k > rgt) cs[++rgt + 1].reset();
       cs[k][p] = 1;
       if (k < lft) r[tp++] = p;
     for (int k = lft; k <= rgt; ++k)</pre>
       for (int p = cs[k]._Find_first
           (); p < N; p = cs[k]._Find_next(p))
         r[tp] = p, c[tp] = k, ++tp;
     dfs(r, c, l + 1, mask);
  }
   void dfs(vector<</pre>
       int> &r, vector<int> &c, int l, bitset<N> mask) {
     while (!r.empty()) {
       int p = r.back();
       r.pop_back(), mask[p] = 0;
       if (q + c.back() <= ans) return;</pre>
       cur[q++] = p;
       vector<int> nr;
       for (int i : r) if (G[p][i]) nr.emplace_back(i);
       if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
       else if (q > ans) ans = q, copy_n(cur, q, sol);
       c.pop_back(), --q;
     }
  }
   int solve() {
     vector<int> r(n);
     ans = q = 0, iota(all(r), 0);
     pre_dfs(r, 0, bitset<N>(string(n, '1')));
     return ans;
  }
};
 2.6 DMST(slow)
```

```
struct zhu_liu { // O(VE)
  struct edge {
    int u, v;
    ll w;
  };
  vector<edge> E; // 0-base
  int pe[N], id[N], vis[N];
  ll in[N];
  void init() { E.clear(); }
  void add_edge(int u, int v, ll w) {
    if (u != v) E.emplace_back(edge{u, v, w});
  ll build(int root, int n) {
    ll ans = 0;
    for (;;) {
      fill_n(in, n, INF);
      for (int i = 0; i < E.size(); ++i)</pre>
        if (E[i].u != E[i].v && E[i].w < in[E[i].v])</pre>
          pe[E[i].v] = i, in[E[i].v] = E[i].w;
```

```
for (int u = 0; u < n; ++u) // no solution
        if (u != root && in[u] == INF) return -INF;
       int cntnode = 0;
      fill_n(id, n, -1), fill_n(vis, n, -1);
      for (int u = 0; u < n; ++u) {
        if (u != root) ans += in[u];
        int v = u:
        while (vis[v] != u && !~id[v] && v != root)
          vis[v] = u, v = E[pe[v]].u;
        if (v != root && !~id[v]) {
          for (int x = E[pe[v]].u; x != v;
                x = E[pe[x]].u)
             id[x] = cntnode;
          id[v] = cntnode++;
      }
      if (!cntnode) break; // no cycle
      for (int u = 0; u < n; ++u)
        if (!~id[u]) id[u] = cntnode++;
      for (int i = 0; i < E.size(); ++i) {</pre>
        int v = E[i].v;
        E[i].v = id[E[i].v], E[i].v = id[E[i].v];
        if (E[i].u != E[i].v) E[i].w -= in[v];
      n = cntnode, root = id[root];
    }
    return ans;
  }
|};
```

2.7 DMST

```
#define rep(i, a, b) for (int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef vector<int> vi;
struct RollbackUF {
 vi e;
 vector<pii> st;
 RollbackUF(int n) : e(n, -1) {}
 int size(int x) { return -e[find(x)]; }
 int find(int x) { return e[x] < 0 ? x : find(e[x]); }</pre>
 int time() { return sz(st); }
 void rollback(int t) {
   for (int i = time(); i-- > t;)
      e[st[i].first] = st[i].second;
   st.resize(t);
 }
 bool join(int a, int b) {
   a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]});
    st.push_back({b, e[b]});
    e[a] += e[b];
   e[b] = a;
    return true;
 }
};
struct Edge {
  int a, b;
 II w;
struct Node { /// lazy skew heap node
 Edge key;
  Node *l, *r;
  ll delta;
  void prop() {
   key.w += delta;
   if (l) l->delta += delta;
   if (r) r->delta += delta;
    delta = 0;
 }
 Edge top() {
   prop();
    return key;
```

```
|};
Node *merge(Node *a, Node *b) {
   if (!a || !b) return a ?: b;
   a->prop(), b->prop();
   if (a->key.w > b->key.w) swap(a, b);
   swap(a->l, (a->r = merge(b, a->r)));
   return a:
}
void pop(Node *&a) {
  a->prop();
  a = merge(a->l, a->r);
}
pair<ll, vi> dmst(int n, int r, vector<Edge> &g) {
  RollbackUF uf(n);
   vector<Node *> heap(n);
   for (Edge e : g)
    heap[e.b] = merge(heap[e.b], new Node{e});
   11 \text{ res} = 0;
   vi seen(n, -1), path(n), par(n);
   seen[r] = r;
   vector<Edge> Q(n), in(n, {-1, -1}), comp;
   deque<tuple<int, int, vector<Edge>>> cycs;
   rep(s, 0, n) {
     int u = s, qi = 0, w;
     while (seen[u] < 0) {
       if (!heap[u]) return {-1, {}};
       Edge e = heap[u]->top();
       heap[u]->delta -= e.w, pop(heap[u]);
       Q[qi] = e, path[qi++] = u, seen[u] = s;
       res += e.w, u = uf.find(e.a);
       if (seen[u] == s) { /// found cycle, contract
         Node *cyc = 0;
         int end = qi, time = uf.time();
         do cyc = merge(cyc, heap[w = path[--qi]]);
         while (uf.join(u, w));
         u = uf.find(u), heap[u] = cyc, seen[u] = -1;
         cycs.push_front({u, time, {&Q[qi], &Q[end]}});
    }
    rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
   for (auto &[u, t, comp] :
    cycs) { // restore sol (optional)
     uf.rollback(t);
     Edge inEdge = in[u];
     for (auto &e : comp) in[uf.find(e.b)] = e;
    in[uf.find(inEdge.b)] = inEdge;
  rep(i, 0, n) par[i] = in[i].a;
   return {res, par};
|}
2.8 VizingTheorem
```

```
namespace Vizing { // Edge coloring
                     // G: coloring adjM
int C[N][N], G[N][N];
void clear(int n) {
  for (int i = 0; i <= n; i++) {</pre>
    for (int j = 0; j <= n; j++) C[i][j] = G[i][j] = 0;</pre>
}
void solve(vector<pii> &E, int n, int m) {
  int X[n] = {}, a;
  auto update = [&](int u) {
    for (X[u] = 1; C[u][X[u]]; X[u]++);
  auto color = [&](int u, int v, int c) {
    int p = G[u][v];
    G[\upsilon][v] = G[v][\upsilon] = c;
    C[u][c] = v;
     C[v][c] = u;
    C[\upsilon][p] = C[\upsilon][p] = 0;
    if (p) X[u] = X[v] = p;
```

```
else update(u), update(v);
    return p;
 };
 auto flip = [&](int u, int c1, int c2) {
    int p = C[v][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
 };
  for (int i = 1; i <= n; i++) X[i] = 1;</pre>
 for (int t = 0; t < E.size(); t++) {</pre>
    int u = E[t].first, v0 = E[t].second, v = v0,
        c0 = X[u], c = c0, d;
    vector<pii> L;
    int vst[n] = {};
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
      if (!C[v][c])
        for (a = (int)L.size() - 1; a >= 0; a--)
          c = color(u, L[a].first, c);
      else if (!C[u][d])
        for (a = (int)L.size() - 1; a >= 0; a--)
          color(u, L[a].first, L[a].second);
      else if (vst[d]) break;
      else vst[d] = 1, v = C[v][d];
    if (!G[u][v0]) {
      for (; v; v = flip(v, c, d), swap(c, d));
      if (C[u][c0]) {
        for (a = (int)L.size() - 2;
             a >= 0 && L[a].second != c; a--)
        for (; a >= 0; a--)
          color(u, L[a].first, L[a].second);
      } else t--;
   }
 }
} // namespace Vizing
```

2.9 MinimumCliqueCover

```
struct Clique_Cover { // O-base, O(n2^n)
   int co[1 << N], n, E[N];</pre>
   int dp[1 << N];</pre>
   void init(int _n) {
     n = _n, fill_n(dp, 1 << n, 0);
     fill_n(E, n, 0), fill_n(co, 1 << n, 0);
   }
   void add_edge(int u, int v) {
     E[u] \mid = 1 << v, E[v] \mid = 1 << u;
   int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;
     dp[0] = (n \& 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {
       int t = i & -i;
       dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)</pre>
       co[i] = (co[i] \& i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic</pre>
       for (int i = 0; i < (1 << n); ++i)</pre>
          sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n;
  }
|};
```

2.10 CountMaximalClique

```
struct BronKerbosch { // 1-base
   int n, a[N], g[N][N];
   int S, all[N][N], some[N][N], none[N][N];
   void init(int _n) {
     n = _n;
     for (int i = 1; i <= n; ++i)</pre>
       for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
   void add_edge(int u, int v) {
     g[v][v] = g[v][v] = 1;
   void dfs(int d, int an, int sn, int nn) {
     if (S > 1000) return; // pruning
     if (sn == 0 && nn == 0) ++S;
     int u = some[d][0];
     for (int i = 0; i < sn; ++i) {</pre>
       int v = some[d][i];
       if (g[v][v]) continue;
int tsn = 0, tnn = 0;
       copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)</pre>
         if (g[v][some[d][j]])
           some[d + 1][tsn++] = some[d][j];
       for (int j = 0; j < nn; ++j)</pre>
         if (g[v][none[d][j]])
           none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
       some[d][i] = 0, none[d][nn++] = v;
     }
   }
   int solve() {
     iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
     return S;
   }
};
```

2.11 Theorems

 $|\max \text{ independent edge set}| = |V| - |\min \text{ edge cover}| \\ |\max \text{ independent set}| = |V| - |\min \text{ vertex cover}| \\$

3 Flow-Matching

3.1 HopcroftKarp

```
struct hoperoftKarp { // O-based
  bool dfs(int a, int L, vector<vector<int>> &g,
    vector<int> &btoa, vector<int> &A,
    vector<int> &B) {
    if (A[a] != L) return 0;
    A[a] = -1;
    for (int b : g[a])
      if (B[b] == L + 1) {
        B[b] = 0;
        if (btoa[b] == -1 ||
          dfs(btoa[b], L + 1, g, btoa, A, B))
          return btoa[b] = a, 1;
    return 0;
  int solve(vector<vector<int>> &g, int m) {
    int res = 0;
    vector<int> btoa(-1, m), A(g.size()),
      B(btoa.size()), cur, next;
    for (;;) {
      fill(all(A), 0), fill(all(B), 0);
      cur.clear();
      for (int a : btoa)
        if (a != -1) A[a] = -1;
      for (int a = 0; a < g.size(); a++)</pre>
        if (A[a] == 0) cur.push_back(a);
      /// Find all layers using bfs.
      for (int lay = 1;; lay++) {
        bool islast = 0;
```

11 solve() {

fill_n(fl

ll res = 0;

return res;

for (int i = 0; i < n; ++i)</pre>

 $hl[i] = *max_element(w[i], w[i] + n);$

for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>

for (int i = 0; i < n; ++i) bfs(i);</pre>

, n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);

```
next.clear();
                                                              }
        for (int a : cur)
                                                            |};
          for (int b : g[a]) {
                                                             3.3 MCMF
             if (btoa[b] == -1) {
               B[b] = lay;
                                                            struct MinCostMaxFlow { // O-base
               islast = 1;
                                                               struct Edge {
             } else if (btoa[b] != a && !B[b]) {
                                                                 ll from, to, cap, flow, cost, rev;
               B[b] = lay;
                                                               } *past[N];
               next.push_back(btoa[b]);
                                                               vector<Edge> G[N];
                                                               int inq[N], n, s, t;
          }
                                                               ll dis[N], up[N], pot[N];
        if (islast) break;
                                                               bool BellmanFord() {
        if (next.empty()) return res;
                                                                 fill_n(dis, n, INF), fill_n(inq, n, 0);
        for (int a : next) A[a] = lay;
                                                                 queue<int> q;
        cur.swap(next);
                                                                 auto relax = [&](int u, ll d, ll cap, Edge *e) {
      }
                                                                   if (cap > 0 && dis[u] > d) {
      /// Use DFS to scan for augmenting paths.
                                                                     dis[v] = d, up[v] = cap, past[v] = e;
      for (int a = 0; a < g.size(); a++)</pre>
                                                                     if (!inq[u]) inq[u] = 1, q.push(u);
        res += dfs(a, 0, g, btoa, A, B);
                                                                  }
                                                                };
  }
                                                                 relax(s, 0, INF, 0);
|};
                                                                 while (!q.empty()) {
                                                                   int u = q.front();
3.2
      KM
                                                                   q.pop(), inq[v] = 0;
                                                                   for (auto &e : G[u]) {
struct KM { // O-base
                                                                     ll d2 = dis[v] + e.cost + pot[v] - pot[e.to];
  ll w[N][N], hl[N], hr[N], slk[N];
                                                                     relax
  int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
                                                                         (e.to, d2, min(up[u], e.cap - e.flow), &e);
  bool vl[N], vr[N];
                                                                  }
  void init(int _n) {
                                                                 }
    n = _n;
                                                                 return dis[t] != INF;
    for (int i = 0; i < n; ++i)</pre>
                                                              }
      fill_n(w[i], n, -INF);
                                                               void solve(int _s
  }
                                                                 , int _t, ll &flow, ll &cost, bool neg = true) { s = \_s, t = \_t, flow = 0, cost = 0;
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
                                                                 if (neg) BellmanFord(), copy_n(dis, n, pot);
  }
                                                                 for (; BellmanFord(); copy_n(dis, n, pot)) {
  bool Check(int x) {
                                                                   for (int
    if (vl[x] = 1, \sim fl[x])
                                                                       i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
      return vr[qu[qr++] = fl[x]] = 1;
                                                                   flow += up[t], cost += up[t] * dis[t];
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
                                                                   for (int i = t; past[i]; i = past[i]->from) {
    return 0;
                                                                     auto &e = *past[i];
                                                                     e.flow += up[t], G[e.to][e.rev].flow -= up[t];
  void bfs(int s) {
                                                                   }
    fill_n(slk
                                                                }
        , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
                                                               void init(int _n) {
    for (ll d;;) {
                                                                 n = _n, fill_n(pot, n, 0);
      while (ql < qr)</pre>
                                                                 for (int i = 0; i < n; ++i) G[i].clear();</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] && slk
                                                               void add_edge(ll a, ll b, ll cap, ll cost) {
               [x] >= (d = hl[x] + hr[y] - w[x][y])) {
                                                                 G[a].emplace_back
             if (pre[x] = y, d) slk[x] = d;
                                                                     (Edge{a, b, cap, 0, cost, (ll)G[b].size()});
             else if (!Check(x)) return;
                                                                 G[b].emplace_back
        }
                                                                     (Edge{b, a, 0, 0, -cost, (ll)G[a].size() - 1});
      d = INF;
      for (int x = 0; x < n; ++x)
                                                            |};
        if (!vl[x] && d > slk[x]) d = slk[x];
       for (int x = 0; x < n; ++x) {
                                                             3.4 GeneralGraphMatching
        if (vl[x]) hl[x] += d;
        else slk[x] -= d;
                                                            struct Matching { // O-base
        if (vr[x]) hr[x] -= d;
                                                               queue<int> q; int n;
                                                               vector<int> fa, s, vis, pre, match;
                                                               vector<vector<int>> G;
      for (int x = 0; x < n; ++x)
                                                               int Find(int u)
        if (!vl[x] && !slk[x] && !Check(x)) return;
```

```
struct Matching { // O-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
        if (vis[x] == tk) return x;
        vis[x] = tk;
        x = Find(pre[match[x]]);
    }
  void Blossom(int x, int y, int l) {
    for (; Find(x) != l; x = pre[y]) {
```

```
pre[x] = y, y = match[x];
       if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
  }
  bool Bfs(int r) {
    iota(all(fa), 0); fill(all(s), -1);
    q = queue<int>(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : G[x])
        if (s[v] == -1) {
          if (pre[u] = x, s[u] = 1, match[u] == n) {
            for (int a = u, b = x, last;
                b != n; a = last, b = pre[a])
               last =
                   match[b], match[b] = a, match[a] = b;
            return true;
          q.push(match[u]); s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int l = LCA(u, x);
          Blossom(x, u, l); Blossom(u, x, l);
        }
    return false;
  Matching(int_n): n(n), fa(n + 1), s(n + 1), vis
       (n + 1), pre(n + 1, n), match(n + 1, n), G(n) {}
  void add_edge(int u, int v)
  { G[u].emplace_back(v), G[v].emplace_back(u); }
  int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
      if (match[x] == n) ans += Bfs(x);
    return ans:
  } // match[x] == n means not matched
|};
```

3.5 MaxWeightMaching

```
#define rep(i, l, r) for (int i = (l); i <= (r); ++i)
struct WeightGraph { // 1-based
  struct edge {
    int u, v, w;
  int n, nx;
  vector<int> lab;
  vector<vector<edge>> g;
  vector<int> slack, match, st, pa, S, vis;
vector<vector<int>> flo, flo_from;
  queue<int> q;
  WeightGraph(int n_)
    : n(n_{-}), nx(n * 2), lab(nx + 1),
      g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1),
      flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
    match = st = pa = S = vis = slack;
    rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
  int ED(edge e) {
    return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
  void update_slack(int u, int x, int &s) {
    if (!s || ED(g[v][x]) < ED(g[s][x])) s = v;</pre>
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)
      if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
        update_slack(u, x, slack[x]);
  void q_push(int x) {
    if (x \le n) q.push(x);
      for (int y : flo[x]) q_push(y);
  void set_st(int x, int b) {
    st[x] = b;
```

```
if(x > n)
    for (int y : flo[x]) set_st(y, b);
vector<int> split_flo(auto &f, int xr) {
  auto it = find(ALL(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
    reverse(1 + ALL(f)), it = f.end() - pr;
  auto res = vector(f.begin(), it);
  return f.erase(f.begin(), it), res;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  int xr = flo_from[u][g[u][v].u];
  auto &f = flo[u], z = split_flo(f, xr);
  rep(i, 0, (int)z.size() - 1)
    set_match(z[i], z[i ^ 1]);
  set_match(xr, v);
  f.insert(f.end(), all(z));
void augment(int u, int v) {
  for (;;) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
}
int lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v))
    if (u) {
      if (vis[u] == t) return u;
      vis[u] = t;
      u = st[match[u]];
      if (u) u = st[pa[u]];
    ļ
  return 0;
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + all(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0;
  match[b] = match[o];
  vector<int> f = {o};
  for (int x = u, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  reverse(1 + all(f));
  for (int x = v, y; x != o; x = st[pa[y]])
    f.emplace_back(x),
      f.emplace_back(y = st[match[x]]), q_push(y);
  flo[b] = f;
  set_st(b, b);
  for (int x = 1; x <= nx; ++x)
    g[b][x].w = g[x][b].w = 0;
  fill(all(flo_from[b]), 0);
  for (int xs : flo[b]) {
    for (int x = 1; x <= nx; ++x)
      if (g[b][x].w == 0 ||
        ED(g[xs][x]) < ED(g[b][x])
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)</pre>
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  }
  set_slack(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) {
      xs = x;
      continue;
```

```
rep(u, 0, n) st[u] = u, flo[u].clear();
    pa[xs] = g[x][xs].u;
                                                              int w_max = 0;
    S[xs] = 1, S[x] = 0;
                                                              rep(u, 1, n) rep(v, 1, n) {
    slack[xs] = 0;
                                                                flo_from[u][v] = (u == v ? u : 0);
                                                                w_max = max(w_max, g[v][v].w);
    set_slack(x);
    q_push(x);
    xs = -1;
                                                              fill(all(lab), w_max);
                                                              int n_matches = 0;
                                                              ll tot_weight = 0;
  for (int x : flo[b])
                                                              while (matching()) ++n_matches;
    if (x == xr) S[x] = 1, pa[x] = pa[b];
                                                              rep(u, 1, n) if (match[u] \&\& match[u] < u)
    else S[x] = -1, set_slack(x);
                                                                tot_weight += g[u][match[u]].w;
  st[b] = 0;
                                                              return make_pair(tot_weight, n_matches);
}
bool on_found_edge(const edge &e) {
                                                            void add_edge(int u, int v, int w) {
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
                                                              g[v][v].w = g[v][v].w = w;
    int nu = st[match[v]];
    pa[v] = e.u;
                                                        |};
    S[v] = 1;
    slack[v] = slack[nu] = 0;
                                                          3.6 GlobalMinCut
    S[nu] = 0;
    q_push(nu);
  } else if (S[v] == 0) {
                                                         struct SW{ // global min cut, O(V^3)
    if (int o = lca(u, v)) add_blossom(u, o, v);
                                                            #define REP for (int i = 0; i < n; ++i)
    else return augment(u, v), augment(v, u), true;
                                                            static const int MXN = 514, INF = 2147483647;
                                                            int vst[MXN], edge[MXN][MXN], wei[MXN];
  return false;
                                                            void init(int n) {
}
                                                              REP fill_n(edge[i], n, 0);
bool matching() {
 fill(all(S), -1), fill(all(slack), 0);
                                                            void addEdge(int u, int v, int w){
  q = queue<int>();
                                                              edge[u][v] += w; edge[v][u] += w;
  for (int x = 1; x <= nx; ++x)</pre>
    if (st[x] == x \&\& !match[x])
                                                            int search(int &s, int &t, int n){
      pa[x] = 0, S[x] = 0, q_push(x);
                                                              fill_n(vst, n, 0), fill_n(wei, n, 0);
  if (q.empty()) return false;
                                                              s = t = -1;
                                                              int mx, cur;
  for (;;) {
    while (q.size()) {
                                                              for (int j = 0; j < n; ++j) {</pre>
                                                               mx = -1, cur = 0;
      int u = q.front();
                                                                REP if (wei[i] > mx) cur = i, mx = wei[i];
      q.pop();
                                                                vst[cur] = 1, wei[cur] = -1;
      if (S[st[u]] == 1) continue;
                                                                s = t; t = cur;
      for (int v = 1; v <= n; ++v)
                                                                REP if (!vst[i]) wei[i] += edge[cur][i];
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (ED(g[v][v]) != 0)
                                                              return mx;
            update_slack(u, st[v], slack[st[v]]);
          else if (on_found_edge(g[u][v]))
                                                            int solve(int n) {
            return true;
                                                              int res = INF;
                                                              for (int x, y; n > 1; n--){
                                                                res = min(res, search(x, y, n));
    int d = INF;
                                                                REP edge[i][x] = (edge[x][i] += edge[y][i]);
    for (int b = n + 1; b <= nx; ++b)</pre>
      if (st[b] == b && S[b] == 1)
                                                                  edge[y][i] = edge[n - 1][i];
        d = min(d, lab[b] / 2);
                                                                  edge[i][y] = edge[i][n - 1];
    for (int x = 1; x <= nx; ++x)</pre>
                                                                } // edge[y][y] = 0;
      if (int s = slack[x];
                                                              }
          st[x] == x && s && s[x] <= 0)
                                                              return res;
        d = min(d, ED(g[s][x]) / (S[x] + 2));
                                                           }
    for (int u = 1; u <= n; ++u)</pre>
                                                        |} sw;
      if (S[st[u]] == 1) lab[u] += d;
      else if (S[st[u]] == 0) {
                                                          3.7 BoundedFlow(Dinic)
        if (lab[u] <= d) return false;</pre>
                                                         struct BoundedFlow { // O-base
        lab[v] -= d;
                                                            struct edge {
                                                              int to, cap, flow, rev;
    rep(b, n + 1, nx) if (st[b] == b && S[b] >= 0)
      lab[b] += d * (2 - 4 * S[b]);
                                                            vector<edge> G[N];
    for (int x = 1; x <= nx; ++x)</pre>
                                                            int n, s, t, dis[N], cur[N], cnt[N];
      if (int s = slack[x]; st[x] == x && s &&
                                                            void init(int _n) {
          st[s] != x && ED(g[s][x]) == 0)
                                                              n = _n;
        if (on_found_edge(g[s][x])) return true;
                                                              for (int i = 0; i < n + 2; ++i)
    for (int b = n + 1; b <= nx; ++b)
                                                                G[i].clear(), cnt[i] = 0;
      if (st[b] == b && S[b] == 1 && lab[b] == 0)
        expand_blossom(b);
                                                            void add_edge(int u, int v, int lcap, int rcap) {
                                                              cnt[u] -= lcap, cnt[v] += lcap;
  return false;
                                                              G[u].emplace_back
}
                                                                  (edge{v, rcap, lcap, (int)G[v].size()});
pair<ll, int> solve() {
                                                              G[v].emplace_back
  fill(all(match), 0);
                                                                  (edge{u, 0, 0, (int)G[u].size() - 1});
```

```
void add_edge(int u, int v, int cap) {
    G[u].emplace_back
         (edge{v, cap, 0, (int)G[v].size()});
    G[v].emplace_back
         (edge{u, 0, 0, (int)G[u].size() - 1});
  int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < G[u].size(); ++i) {</pre>
      edge &e = G[u][i];
      if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
         int df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df, G[e.to][e.rev].flow -= df;
           return df;
        }
      }
    dis[v] = -1;
    return 0;
  }
  bool bfs() {
    fill_n(dis, n + 3, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (edge &e : G[u])
         if (!~dis[e.to] && e.flow != e.cap)
           q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  }
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
    int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
    int sum = 0;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
         add_edge(n + 1, i, cnt[i]), sum += cnt[i];
       else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);</pre>
    if (sum != maxflow(n + 1, n + 2)) sum = -1;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
         G[n + 1].pop_back(), G[i].pop_back();
       else if (cnt[i] < 0)</pre>
         G[i].pop_back(), G[n + 2].pop_back();
    return sum != -1;
  }
  int solve(int _s, int _t) {
    add_edge(_t, _s, INF);
    if (!solve()) return -1; // invalid flow
    int x = G[_t].back().flow;
    return G[_t].pop_back(), G[_s].pop_back(), x;
  }
|};
```

3.8 GomoryHuTree

```
| MaxFlow Dinic;
| int g[N];
| void GomoryHu(int n) { // 0-base
| fill_n(g, n, 0);
| for (int i = 1; i < n; ++i) {
| Dinic.reset();
| add_edge(i, g[i], Dinic.maxflow(i, g[i]));
| for (int j = i + 1; j <= n; ++j)
| if (g[j] == g[i] && ~Dinic.dis[j])
```

```
g[j] = i;
```

} |}

3.9 MinCostCirculation

```
| struct MinCostCirculation { // O-base
  struct Edge {
     ll from, to, cap, fcap, flow, cost, rev;
  } *past[N];
  vector<Edge> G[N];
  ll dis[N], inq[N], n;
  void BellmanFord(int s) {
     fill_n(dis, n, INF), fill_n(inq, n, 0);
     queue<int> q;
     auto relax = [&](int u, ll d, Edge *e) {
       if (dis[u] > d) {
         dis[u] = d, past[u] = e;
         if (!inq[u]) inq[u] = 1, q.push(u);
    };
    relax(s, 0, 0);
    while (!q.empty()) {
       int u = q.front();
       q.pop(), inq[v] = 0;
       for (auto &e : G[v])
         if (e.cap > e.flow)
           relax(e.to, dis[u] + e.cost, &e);
    }
  }
  void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
     BellmanFord(cur.to);
     if (dis[cur.from] + cur.cost < 0) {</pre>
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
            i = cur.from; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         ++e.flow, --G[e.to][e.rev].flow;
    }
     ++cur.cap;
  }
  void solve(int mxlq) {
     for (int b = mxlg; b >= 0; --b) {
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           e.cap *= 2, e.flow *= 2;
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
             try_edge(e);
  }
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
     G[a].emplace_back(Edge{a, b,
          0, cap, 0, cost, (11)G[b].size() + (a == b));
     G[b].emplace_back(Edge
         {b, a, 0, 0, 0, -cost, (ll)G[a].size() - 1});
|} mcmf; // O(VE * ElogC)
```

3.10 FlowModelsBuilding

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - For each edge (x,y,l,u), connect x→y with capacity u-l.
 For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - of outgoing lower bounds. 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v). To maximize, connect $t\to s$ with capacity ∞ (skip

this in circulation problem), and let f be the

```
maximum flow from S to T . If f\!\neq\!\sum_{v\in V, in(v)>0}\!in(v) ,
          there's no solution. Otherwise, the maximum flow
          from s to t is the answer.
       – To minimize, let f be the maximum flow from {\cal S} to
         T. Connect t \to s with capacity \infty and let the flow from S to T be f'. If f+f' \neq \sum_{v \in V, in(v)>0} in(v),
          there's no solution. Otherwise, f^\prime is the answer.
  5. The solution of each edge e is \widetilde{l_e}+f_e, where f_e
      corresponds to the flow of edge e on the graph.
ullet Construct minimum vertex cover from maximum matching M
  on bipartite graph (X,Y)
  1. Redirect every edge: y \rightarrow x if (x,y) \in M, x \rightarrow y otherwise.
  2. DFS from unmatched vertices in X.
  3. x \in X is chosen iff x is unvisited.
  4. y \in Y is chosen iff y is visited.

    Minimum cost cyclic flow

  1. Consruct super source S and sink T
  2. For each edge (x,y,c), connect x \rightarrow y with (cost,cap) = (c,1)
      if c>0, otherwise connect y\to x with (cost, cap)=(-c,1)
  3. For each edge with c < 0, sum these cost as K, then
      increase d(y) by 1, decrease d(x) by 1
  4. For each vertex v with d(v) > 0, connect S \to v with
      (cost, cap) = (0, d(v))
      For each vertex v with d(v) < 0, connect v \to T with
      (cost, cap) = (0, -d(v))
  6. Flow from S to T, the answer is the cost of the flow
      C+K
  Maximum density induced subgraph
  1. Binary search on answer, suppose we're checking
      \hbox{answer } T
  2. Construct a max flow model, let K be the sum of all
      weights
  3. Connect source s \to v, v \in G with capacity K 4. For each edge (u,v,w) in G, connect u \to v and v \to u with
      capacity w
  5. For v\in G , connect it with sink v\to t with capacity K+2T-(\sum_{e\in E(v)}w(e))-2w(v)
  6. T is a valid answer if the maximum flow f < K|V|
  Minimum weight edge cover
  1. For each v \in V create a copy v', and connect u' \to v'
      with weight w(u,v).
  2. Connect v \! \to \! v' with weight 2\mu(v), where \mu(v) is the cost
      of the cheapest edge incident to v.
  3. Find the minimum weight perfect matching on G'.
 Project selection problem
  1. If p_v \! > \! 0, create edge (s,\!v) with capacity p_v; otherwise,
      create edge (v,t) with capacity -p_v.
  2. Create edge (u,v) with capacity w with w being the
      cost of choosing u without choosing v.
  3. The mincut is equivalent to the maximum profit of a
      subset of projects.

    Dual of minimum cost maximum flow

  1. Capacity c_{uv}, Flow f_{uv}, Cost w_{uv}, Required Flow
      difference for vertex b_u .
  2. If all w_{uv} are integers, then optimal solution can
      happen when all p_u are integers.
           \begin{aligned} \min & \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} & \geq -c_{uv} \Leftrightarrow \min & \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \end{aligned}
   \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_u
4.1 LCT
```

4 Data Struture

```
#define ls(x) Tree[x].son[0]
#define rs(x) Tree[x].son[1]
#define fa(x) Tree[x].fa
const int maxn = 600010;
struct node {
  int son[2], Min, id, fa, lazy;
} Tree[maxn];
int n, m, q, w[maxn], Min;
struct Node {
 int u, v, w;
} a[maxn];
inline bool IsRoot(int x) {
 return (ls(fa(x)) == x \mid\mid rs(fa(x)) == x) ? false
                                              : true;
inline void PushUp(int x) {
 Tree[x].Min = w[x], Tree[x].id = x;
```

```
if (ls(x) && Tree[ls(x)].Min < Tree[x].Min) {</pre>
    Tree[x].Min = Tree[ls(x)].Min;
    Tree[x].id = Tree[ls(x)].id;
  if (rs(x) && Tree[rs(x)].Min < Tree[x].Min) {</pre>
    Tree[x].Min = Tree[rs(x)].Min;
    Tree[x].id = Tree[rs(x)].id;
inline void Update(int x) {
  Tree[x].lazy ^= 1;
  swap(ls(x), rs(x));
inline void PushDown(int x) {
  if (!Tree[x].lazy) return;
  if (ls(x)) Update(ls(x));
  if (rs(x)) Update(rs(x));
  Tree[x].lazy = 0;
inline void Rotate(int x) {
  int y = fa(x), z = fa(y), k = rs(y) == x,
      w = Tree[x].son[!k];
  if (!IsRoot(y)) Tree[z].son[rs(z) == y] = x;
  fa(x) = z, fa(y) = x;
  if (w) fa(w) = y;
  Tree[x].son[!k] = y, Tree[y].son[k] = w;
  PushUp(y);
inline void Splay(int x) {
  stack<int> Stack;
  int y = x, z;
  Stack.push(y);
  while (!IsRoot(y)) Stack.push(y = fa(y));
  while (!Stack.empty())
    PushDown(Stack.top()), Stack.pop();
  while (!IsRoot(x)) {
    y = fa(x), z = fa(y);
    if (!IsRoot(y))
      Rotate((ls(y) == x) ^(ls(z) == y) ? x : y);
    Rotate(x);
  PushUp(x);
inline void Access(int root) {
  for (int x = 0; root; x = root, root = fa(root))
    Splay(root), rs(root) = x, PushUp(root);
inline void MakeRoot(int x) {
 Access(x), Splay(x), Update(x);
}
inline int FindRoot(int x) {
  Access(x), Splay(x);
  while (ls(x)) x = ls(x);
  return Splay(x), x;
inline void Link(int u, int v) {
  MakeRoot(u);
  if (FindRoot(v) != u) fa(u) = v;
}
inline void Cut(int u, int v) {
  MakeRoot(u);
  if (FindRoot(v) != u || fa(v) != u || ls(v)) return;
  fa(v) = rs(u) = 0;
}
inline void Split(int u, int v) {
  MakeRoot(u), Access(v), Splay(v);
}
inline bool Check(int u, int v) {
  return MakeRoot(u), FindRoot(v) == u;
inline int LCA(int root, int u, int v) {
  MakeRoot(root), Access(u), Access(v), Splay(u);
  if (!fa(u)) {
    Access(u), Splay(v);
    return fa(v);
```

```
}
 return fa(u);
}
每次進入節點和走邊都放入一次共 3n - 2
node(u) 表示進入節點 u 放入 treap 的位置
edge(u, v) 表示 u -> v 的邊放入 treap 的位置 (push v)
Makeroot u :
 L1 = [begin, node(u) - 1], L2 = [node(u), end]
  -> L2 + L1
Insert u, v:
 Tu \rightarrow L1 = [begin, node(u) - 1], L2 = [node(u), end]
  Tv \rightarrow L3 = [begin, node(v) - 1], L4 = [node(v), end]
  -> L2 + L1 + edge(u, v) + L4 + L3 + edge(v, u)
Delect u, v .
 maybe need swap u, v
 T -> L1 + edge(u, v) + L2 + edge(v, u) + L3
  -> L1 + L3, L2
```

4.2 Treap

```
int data, sz;
  node *l, *r;
 node(int k) : data(k), sz(1), l(0), r(0) {}
 void up() {
    sz = 1;
   if (l) sz += l->sz;
   if (r) sz += r->sz;
 }
 void down() {}
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
 if (!a || !b) return a ? a : b;
  if (rand() % (sz(a) + sz(b)) < sz(a))
    return a->down(), a->r = merge(a->r, b), a->up(),
 return b->down(), b->l = merge(a, b->l), b->up(), b;
void split(node *o, node *&a, node *&b, int k) {
 if (!o) return a = b = 0, void();
  o->down();
 if (o->data <= k)
   a = o, split(o->r, a->r, b, k), a->up();
 else b = o, split(o->l, a, b->l, k), b->up();
void split2(node *o, node *&a, node *&b, int k) {
 if (sz(0) <= k) return a = 0, b = 0, void();</pre>
 o->down();
 if (sz(o->l) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
 else b = o, split2(o->l, a, b->l, k);
 o->up();
node *kth(node *o, int k) {
 if (k <= sz(o->l)) return kth(o->l, k);
 if (k == sz(o->l) + 1) return o;
 return kth(o->r, k - sz(o->l) - 1);
int Rank(node *o, int key) {
 if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o->l) + 1 + Rank(o->r, key);
 else return Rank(o->l, key);
bool erase(node *&o, int k) {
 if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->l, o->r);
    delete t:
    return 1;
 node *&t = k < o->data ? o->l : o->r;
```

```
return erase(t, k) ? o->up(), 1 : 0;
| }
void insert(node *&o, int k) {
  node *a, *b;
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
void interval(node *&o, int l, int r) {
  node *a, *b, *c;
  split2(o, a, b, l - 1), split2(b, b, c, r);
  // operate
  o = merge(a, merge(b, c));
```

4.3 CentroidDecomposition

```
vector<pll> G[N];
pll info[N]; // store info. of itself
pll upinfo[N]; // store info. of climbing up
int n, pa[N], layer[N], sz[N], done[N];
ll dis[__lg(N) + 1][N];
void init(int _n) {
  n = _n, layer[0] = -1;
  fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
  for (int i = 1; i <= n; ++i) G[i].clear();</pre>
void add_edge(int a, int b, int w) {
  G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
void get_cent(
  int u, int f, int &mx, int &c, int num) {
  int mxsz = 0;
  sz[u] = 1;
  for (pll e : G[u])
    if (!done[e.X] && e.X != f) {
      get_cent(e.X, u, mx, c, num);
      sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
  if (mx > max(mxsz, num - sz[u]))
    mx = max(mxsz, num - sz[u]), c = u;
void dfs(int u, int f, ll d, int org) {
  // if required, add self info or climbing info
  dis[layer[org]][u] = d;
  for (pll e : G[u])
    if (!done[e.X] && e.X != f)
      dfs(e.X, u, d + e.Y, org);
int cut(int u, int f, int num) {
  int mx = 1e9, c = 0, lc;
  get_cent(u, f, mx, c, num);
  done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
  for (pll e : G[c])
    if (!done[e.X]) {
      if (sz[e.X] > sz[c])
        lc = cut(e.X, c, num - sz[c]);
      else lc = cut(e.X, c, sz[e.X]);
      upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    }
  return done[c] = 0, c;
void build() { cut(1, 0, n); }
void modify(int u) {
  for (int a = u, ly = layer[a]; a;
       a = pa[a], --ly) {
    info[a].X += dis[ly][u], ++info[a].Y;
      upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
 }
ll query(int u) {
  ll rt = 0;
  for (int a = u, ly = layer[a]; a;
       a = pa[a], --ly) {
    rt += info[a].X + info[a].Y * dis[ly][u];
    if (pa[a])
```

4.4 HeavylightDecomposition

```
int n, ulink[N], deep[N], mxson[N], w[N], pa[N];
  int t, pl[N], data[N], dt[N], bln[N], edge[N], et;
  vector<pii> G[N];
  void init(int _n) {
  n = _n, t = 0, et = 1;
    for (int i = 1; i <= n; ++i)</pre>
       G[i].clear(), mxson[i] = 0;
  void add_edge(int a, int b, int w) {
     G[a].pb(pii(b, et));
     G[b].pb(pii(a, et));
     edge[et++] = w;
  void dfs(int u, int f, int d) {
    w[u] = 1, pa[u] = f, deep[u] = d++;
    for (auto &i : G[u])
       if (i.X != f) {
         dfs(i.X, u, d), w[u] += w[i.X];
         if (w[mxson[u]] < w[i.X]) mxson[u] = i.X;</pre>
       } else bln[i.Y] = u, dt[u] = edge[i.Y];
  }
  void cut(int u, int link) {
     data[pl[v] = t++] = dt[v], vlink[v] = link;
     if (!mxson[u]) return;
    cut(mxson[u], link);
     for (auto i : G[u])
       if (i.X != pa[u] && i.X != mxson[u])
         cut(i.X, i.X);
  void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
  int query(int a, int b) {
     int ta = ulink[a], tb = ulink[b], re = 0;
     while (ta != tb)
       if (deep[ta] < deep[tb])</pre>
         /*query*/, tb = ulink[b = pa[tb]];
       else /*query*/, ta = ulink[a = pa[ta]];
     if (a == b) return re;
     if (pl[a] > pl[b]) swap(a, b);
     /*query*/
     return re:
|};
```

5 String

5.1 KMP

```
int KMP(string s, string t) {
  t = " "s + t; // consistency with ACa
  int n = t.size(), ans = 0;
  vector<int> f(t.size(), 0);
  f[0] = -1;
  for (int i = 1, j = -1; i < t.size(); i++) {</pre>
    while (j \ge 0 \&\& t[j + 1] != t[i])
       j = f[j];
    f[i] = ++j;
  for (int i = 0, j = 0; i < s.size(); i++) {</pre>
    while (j \ge 0 \&\& t[j + 1] != s[i])
       j = f[j];
    if (++j + 1 == t.size()) ans++, j = f[j];
  }
  return ans;
|}
```

5.2 Z

```
int Z[1000006];
void z(string s) {
  for (int i = 1, mx = 0; i < s.size(); i++) {
    if (i < Z[mx] + mx)
        Z[i] = min(Z[mx] - i + mx, Z[i - mx]);
    while (
        Z[i] + i < s.size() && s[i + Z[i]] == s[Z[i]])
        Z[i]++;
    if (Z[i] + i > Z[mx] + mx) mx = i;
  }
}
```

5.3 Manacher

```
int man[2000006];
int manacher(string s) {
   string t;
   for (int i = 0; i < s.size(); i++) {</pre>
     if (i) t.push_back('$');
     t.push_back(s[i]);
   int mx = 0, ans = 0;
  for (int i = 0; i < t.size(); i++) {</pre>
     man[i] = 1;
     man[i] = min(man[mx] + mx - i, man[2 * mx - i]);
     while (man[i] + i < t.size() && i - man[i] >= 0 &&
       t[i + man[i]] == t[i - man[i]])
       man[i]++;
     if (i + man[i] > mx + man[mx]) mx = i;
   for (int i = 0; i < t.size(); i++)</pre>
     ans = max(ans, man[i] - 1);
   return ans;
| }
```

5.4 SuffixArray

```
vector<int> sa, cnt, rk, tmp, lcp;
void SA(string s) {
   int n = s.size();
   sa.resize(n), cnt.resize(n), rk.resize(n),
     tmp.resize(n);
   iota(all(sa), 0);
   sort(all(sa),
     [&](int i, int j) { return s[i] < s[j]; });
   rk[0] = 0;
   for (int i = 1; i < n; i++)</pre>
     rk[sa[i]] =
       rk[sa[i - 1]] + (s[sa[i - 1]] != s[sa[i]]);
   for (int k = 1; k <= n; k <<= 1) {</pre>
     fill(all(cnt), 0);
     for (int i = 0; i < n; i++)</pre>
       cnt[rk[(sa[i] - k + n) % n]]++;
     for (int i = 1; i < n; i++) cnt[i] += cnt[i - 1];</pre>
     for (int i = n - 1; i >= 0; i--)
       tmp[--cnt[rk[(sa[i] - k + n) % n]]] =
         (sa[i] - k + n) % n;
     sa.swap(tmp);
     tmp[sa[0]] = 0;
     for (int i = 1; i < n; i++)
       tmp[sa[i]] = tmp[sa[i - 1]] +
         (rk[sa[i - 1]] != rk[sa[i]] ||
           rk[(sa[i - 1] + k) % n] !=
             rk[(sa[i] + k) % n]);
     rk.swap(tmp);
  }
}
void LCP(string s) {
  int n = s.size(), k = 0;
  lcp.resize(n);
   for (int i = 0; i < n; i++)</pre>
     if (rk[i] == 0) lcp[rk[i]] = 0;
     else {
       if (k) k--;
```

```
int j = sa[rk[i] - 1];
while (
    i + k < n && j + k < n && s[i + k] == s[j + k])
    k++;
    lcp[rk[i]] = k;
}
5.5 SAIS</pre>
```

```
namespace sfx {
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
// zero based, string content MUST > 0
// SA[i]: SA[i]-th
    suffix is the i-th lexigraphically smallest suffix.
// H[i]: longest
    common prefix of suffix SA[i] and suffix SA[i - 1].
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, 0), copy_n(c, z, x); }
void induce
    (int *sa, int *c, int *s, bool *t, int n, int z) {
  copy_n(c, z - 1, x + 1);
 for (int i = 0; i < n; ++i)</pre>
   if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
 for (int i = n - 1; i >= 0; --i)
   if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa
    , int *p, int *q, bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0,
      nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  fill_n(c, z, 0);
 for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
 partial_sum(c, c + z, c);
  if (uniq) {
   for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
    return;
 }
 for (int i = n - 2; i >= 0; --i)
    t[i] = (
        s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
 pre(sa, c, n, z);
  for (int i = 1; i <= n - 1; ++i)</pre>
    if (t[i] && !t[i - 1])
      sa[--x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
 for (int i = 0; i < n; ++i)</pre>
   if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
      bool neq = last < 0 || !equal
          (s + sa[i], s + p[q[sa[i]] + 1], s + last);
      ns[q[last = sa[i]]] = nmxz += neq;
   }
  sais(ns,
       nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
 pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i)
   sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
 induce(sa, c, s, t, n, z);
void mkhei(int n) {
 for (int i = 0, j = 0; i < n; ++i) {</pre>
   if (RA[i])
      for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
   H[RA[i]] = j, j = max(0, j - 1);
 }
void build(int *s, int n) {
 copy_n(s, n, _s), _s[n] = 0;
  sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
 copy_n(SA + 1, n, SA);
```

```
for (int i = 0; i < n; ++i) RA[SA[i]] = i;</pre>
   mkhei(n);
| } }
 5.6 ACAutomaton
 #define sumS 500005
 #define sigma 26
 #define base 'a
 struct AhoCorasick {
   int ch[sumS][sigma] = {{}}, f[sumS] = {-1},
       tag[sumS], mv[sumS][sigma], jump[sumS],
       cnt[sumS];
   int idx = 0, t = -1;
   vector<int> E[sumS], q;
   pii o[sumS];
   int insert(string &s, int t) {
     int j = 0;
     for (int i = 0; i < (int)s.size(); i++) {</pre>
       if (!ch[j][s[i] - base])
         ch[j][s[i] - base] = ++idx;
       j = ch[j][s[i] - base];
     tag[j] = 1;
     return j;
   int next(int u, int c) {
     return u < 0 ? 0 : mv[u][c];</pre>
   void dfs(int u) {
     o[v].F = ++t;
     for (auto v : E[u]) dfs(v);
     o[u].S = t;
   void build() {
     int k = -1;
     q.emplace_back(0);
     while (++k < q.size()) {</pre>
       int u = q[k];
       for (int v = 0; v < sigma; v++) {</pre>
         if (ch[u][v]) {
           f[ch[u][v]] = next(f[u], v);
           q.emplace_back(ch[u][v]);
         mv[v][v] =
            (ch[u][v] ? ch[u][v] : next(f[u], v));
       }
       if (u) jump[u] = (tag[f[u]] ? f[u] : jump[f[u]]);
     reverse(q.begin(), q.end());
     for (int i = 1; i <= idx; i++)</pre>
       E[f[i]].emplace_back(i);
     dfs(0);
   void match(string &s) {
     fill(cnt, cnt + idx + 1, 0);
     for (int i = 0, j = 0; i < (int)s.size(); i++)</pre>
       cnt[j = next(j, s[i] - base)]++;
     for (int i : q)
       if (f[i] > 0) cnt[f[i]] += cnt[i];
   }
|} ac;
 5.7 MinRotation
int mincyc(string s) {
   int n = s.size();
   s = s + s;
   int i = 0, ans = 0;
   while (i < n) {
     ans = i:
     int j = i + 1, k = i;
     while (j < s.size() && s[j] >= s[k]) {
       k = (s[j] > s[k] ? i : k + 1);
       ++j;
```

while (i <= k) i += j - k;</pre>

```
return ans:
| }
5.8 ExtSAM
#define CNUM 26
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
  int cnt[N * 2]; // occurence
  int newnode() {
     fill_n(next[tot], CNUM, 0);
     len[tot] = cnt[tot] = link[tot] = 0;
     return tot++;
  }
  void init() { tot = 0, newnode(), link[0] = -1; }
  int insertSAM(int last, int c) {
     int cur = next[last][c];
    len[cur] = len[last] + 1;
     int p = link[last];
     while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
     if (p == -1) return link[cur] = 0, cur;
     int q = next[p][c];
     if (len
         [p] + 1 == len[q]) return link[cur] = q, cur;
     int clone = newnode();
     for (int i = 0; i < CNUM; ++i)</pre>
      next[
           clone][i] = len[next[q][i]] ? next[q][i] : 0;
    len[clone] = len[p] + 1;
     while (p != -1 && next[p][c] == q)
       next[p][c] = clone, p = link[p];
     link[link[cur] = clone] = link[q];
    link[q] = clone;
     return cur;
  }
  void insert(const string &s) {
     int cur = 0;
     for (auto ch : s) {
       int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
  }
  void build() {
     queue<int> q;
     q.push(0);
     while (!q.empty()) {
      int cur = q.front();
      q.pop();
       for (int i = 0; i < CNUM; ++i)</pre>
         if (next[cur][i])
           q.push(insertSAM(cur, i));
     vector<int> lc(tot);
    for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
     partial_sum(all(lc), lc.begin());
     for (int i
         = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i; | \} // |x| <= b/2, |y| <= a/2
  void solve() {
     for (int i = tot - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
|};
5.9
      PalindromeTree
```

```
struct palindromic_tree {
  struct node {
   int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
                  // pal. suf.
   node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
```

```
for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
    }
   };
   vector<node> St;
   vector<char> s;
   int last, n;
   palindromic_tree() : St(2), last(1), n(0) {
     St[0].fail = 1, St[1].len = -1, s.emplace_back(-1);
   inline void clear() {
     St.clear(), s.clear(), last = 1, n = 0;
     St.emplace_back(0), St.emplace_back(-1);
     St[0].fail = 1, s.emplace_back(-1);
   inline int get_fail(int x) {
     while (s[n - St[x].len - 1] != s[n])
       x = St[x].fail;
     return x;
   inline void add(int c) {
     s.push_back(c -= 'a'), ++n;
     int cur = get_fail(last);
     if (!St[cur].next[c]) {
       int now = St.size();
       St.emplace_back(St[cur].len + 2);
       St[now].fail =
         St[get_fail(St[cur].fail)].next[c];
       St[cur].next[c] = now;
       St[now].num = St[St[now].fail].num + 1;
     last = St[cur].next[c], ++St[last].cnt;
   inline void count() { // counting cnt
     auto i = St.rbegin();
     for (; i != St.rend(); ++i) {
       St[i->fail].cnt += i->cnt;
  }
  inline int size() { // The number of diff. pal.
     return (int)St.size() - 2;
};
```

Number Theory

6.1 Primes

12721 13331 14341 75577 123457 222557 556679 999983 1097774749 1076767633 100102021 999997771 1001010013 98789101 1000512343 987654361 999991231 999888733 987777733 999991921 1010101333 1010102101 1000000000039 1000000000000037 2305843009213693951 4611686018427387847 9223372036854775783 18446744073709551557

6.2 ExtGCD

```
// beware of negative numbers!
void extgcd(ll a, ll b, ll c, ll &x, ll &y) {
  if (b == 0) x = c / a, y = 0;
  else {
    extgcd(b, a % b, c, y, x);
    y -= x * (a / b);
```

6.3 FloorCeil

```
int floor(int a, int b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
|{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

6.4 FloorSum

```
| II floorsum(II A, II B, II C, II N) {
  if (A == 0) return (N + 1) * (B / C);
  if (A > C || B > C)
     return (N + 1) * (B / C) +
      N * (N + 1) / 2 * (A / C) +
```

```
floorsum(A % C, B % C, C, N);
  11 M = (A * N + B) / C;
  return N * M - floorsum(C, C - B - 1, A, M - 1);
|} // \sum^{n}_0 floor((ai + b) / m)
6.5 MillerRabin
// n < 4,759,123,141 3 : 2, 7, 61
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 4,759,123,141
// n < 3,474,749,660,383 6 : primes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool Miller_Rabin(ll a, ll n) {
  if ((a = a % n) == 0) return 1;
  if (n % 2 == 0) return n == 2;
  ll tmp = (n - 1) / ((n - 1) & (1 - n));
  for (; tmp; tmp >>= 1, a = mul(a, a, n))
```

if ((x = mul(x, x, n)) == n - 1) return 1;

if (tmp & 1) x = mul(x, a, n);

if (x == 1 || x == n - 1) return 1;

6.6 PollardRho

while (--t)

return 0;

|}

```
|map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
       == 0) return PollardRho(n / 2), ++cnt[2], void();
  11 x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return;
    }
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
    d = gcd(abs(x - y), n);
  }
|}
```

6.7 Fraction

```
struct fraction {
   ll n, d;
   fraction
       (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
    ll t = gcd(n, d);
    n /= t, d /= t;
    if (d < 0) n = -n, d = -d;
  }
  fraction operator-() const
  { return fraction(-n, d); }
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator-(const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;
  }
|};
```

6.8 ChineseRemainder

```
|ll solve(ll x1, ll m1, ll x2, ll m2) {
| ll g = gcd(m1, m2);
| if ((x2 - x1) % g) return -1; // no sol
```

```
m1 /= g; m2 /= g;
ll x, y;
extgcd(m1, m2, __gcd(m1, m2), x, y);
ll lcm = m1 * m2 * g;
ll res = x * (x2 - x1) * m1 + x1;
// be careful with overflow
return (res % lcm + lcm) % lcm;
}
```

6.9 Factorial $\mathsf{Mod} p^k$

```
// 0(p^k + log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
   prod[0] = 1;
   for (int i = 1; i <= pk; ++i)
        if (i % p) prod[i] = prod[i - 1] * i % pk;
        else prod[i] = prod[i - 1];
        ll rt = 1;
   for (; n; n /= p) {
        rt = rt * mpow(prod[pk], n / pk, pk) % pk;
        rt = rt * prod[n % pk] % pk;
   }
   return rt;
} // (n! without factor p) % p^k</pre>
```

6.10 QuadraticResidue

```
|// Berlekamp-Rabin, log^2(p)
ll trial(ll y, ll z, ll m) {
  11 a0 = 1, a1 = 0, b0 = z, b1 = 1, p = (m - 1) / 2;
  while (p) {
     if (p & 1)
      tie(a0, a1) =
         make_pair((a1 * b1 % m * y + a0 * b0) % m,
           (a0 * b1 + a1 * b0) % m);
     tie(b0, b1) =
      make_pair((b1 * b1 % m * y + b0 * b0) % m,
         (2 * b0 * b1) % m);
    p >>= 1;
  if (a1) return inv(a1, m);
  return -1;
}
mt19937 rd(49);
ll psqrt(ll y, ll p) {
  if (fpow(y, (p - 1) / 2, p) != 1) return -1;
  for (int i = 0; i < 30; i++) {
    ll z = rd() \% p;
    if (z * z % p == y) return z;
    ll x = trial(y, z, p);
    if (x == -1) continue;
    return x;
  return -1;
}
```

6.11 MeisselLehmer

```
| II PrimeCount(II n) \{ // n \sim 10^13 => < 2s \}
   if (n <= 1) return 0;</pre>
   int v = sqrt(n), s = (v + 1) / 2, pc = 0;
   vector<int> smalls(v + 1), skip(v + 1), roughs(s);
   vector<ll> larges(s);
   for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
   for (int i = 0; i < s; ++i) {</pre>
     roughs[i] = 2 * i + 1;
     larges[i] = (n / (2 * i + 1) + 1) / 2;
   for (int p = 3; p <= v; ++p) {</pre>
     if (smalls[p] > smalls[p - 1]) {
       int q = p * p;
       ++pc;
       if (1LL * q * q > n) break;
       skip[p] = 1;
       for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
       int ns = 0;
       for (int k = 0; k < s; ++k) {</pre>
```

```
int i = roughs[k];
        if (skip[i]) continue;
        11 d = 1LL * i * p;
        larges[ns] = larges[k] - (d <= v ? larges
             [smalls[d] - pc] : smalls[n / d]) + pc;
        roughs[ns++] = i;
      }
      s = ns:
      for (int j = v / p; j >= p; --j) {
              smalls[j] - pc, e = min(j * p + p, v + 1);
        for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
      }
    }
  }
  for (int k = 1; k < s; ++k) {
    const ll m = n / roughs[k];
    ll t = larges[k] - (pc + k - 1);
    for (int l = 1; l < k; ++l) {</pre>
      int p = roughs[l];
      if (1LL * p * p > m) break;
      t -= smalls[m / p] - (pc + l - 1);
    larges[0] -= t;
  }
  return larges[0];
}
```

6.12 DiscreteLog

```
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered_map<int, int> p;
   int b = 1:
   for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
     y = 1LL * y * x % m;
     b = 1LL * b * x % m;
  }
   for (int i = 0; i < m + 10; i += kStep) {</pre>
     s = 1LL * s * b % m;
     if (p.find(s) != p.end()) return i + kStep - p[s];
  }
   return -1:
 int DiscreteLog(int x, int y, int m) {
   if (m == 1) return 0;
   int s = 1;
   for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
     s = 1LL * s * x % m;
   if (s == y) return 100;
   int p = 100 + DiscreteLog(s, x, y, m);
   if (fpow(x, p, m) != y) return -1;
   return p;
j }
```

6.13 Theorems

· Cramer's rule

$$\begin{array}{c} ax+by=e \\ cx+dy=f \\ \end{array} \Rightarrow \begin{array}{c} x=\frac{ed-bf}{ad-bc} \\ y=\frac{af-ec}{ad-bc} \end{array}$$

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}\!=\!d(i)$, $L_{ij}\!=\!-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(L_{rr})|$.
- Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij}\!=\!-d_{ji}$. $rac{rank(D)}{2}$ is the maximum matching on G .

- Cayley's Formula
 - Given a degree sequence $d_1,d_2,...,d_n$ for each labeled vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős-Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+\cdots+d_n$ is even and

$$\sum_{i=1}^k \! d_i \! \leq \! k(k-1) + \sum_{i=k+1}^n \! \min(d_i,\!k) \text{ holds for every } 1 \! \leq \! k \! \leq \! n \text{.}$$

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and $b_1,...,b_n$ is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$$
 holds for every $1 \leq k \leq n$.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),...,(a_n,b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$

and
$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k)$$
 holds for every

• Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\mbox{lattice points in the interior}\}+$ #{lattice points on the boundary} -1.

- Möbius inversion formula

 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$ $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- · Spherical cap
 - A portion of a sphere cut off by a plane.
 - r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : $\arcsin(a/r)$.
 - Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos\theta) (1 \cos\theta)^2/3$.
 - Area $=2\pi r \hat{h} = \pi (a^2 + \hat{h}^2) = 2\pi r^2 (1 \cos\theta)$.
- Lagrange multiplier
 - Optimize $f(x_1,...,x_n)$ when k constraints $g_i(x_1,...,x_n)=0$.
 - Lagrangian function $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=f(x_1,\ldots,x_n)$
 - $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n)$. The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines
 - Line 1: $v_1 = p_1 + t_1 d_1$
 - Line 2: ${m v}_2\!=\!{m p}_2\!+\!t_2{m d}_2$
 - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
 - $\boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}$ - $\boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$

 - $\begin{array}{lll} & c_1 \!=\! p_1 \!+\! \frac{(p_2 \!-\! p_1) \cdot n_2}{d_1 \cdot n_2} d_1 \\ & c_2 \!=\! p_2 \!+\! \frac{(p_1 \!-\! p_2) \cdot n_1}{d_2 \cdot n_1} d_2 \end{array}$

6.14 Estimation

- Estimation
 - The number of divisors of n is at most around 100 for $n\!<\!$ 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200000 for n < 1e19.
 - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30 for $n=0\sim 9$, 627 for n=20, $\sim 2e5$ for n = 50, $\sim 2e8$ for n = 100.
 - Total number of partitions of n distinct elements: B(n) = 1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437,190899322,....

6.15 EuclideanAlgorithms

- $m = \lfloor \frac{an+b}{a} \rfloor$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \mod c, b \mod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \text{ mod } c,b \text{ mod } c,c,n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ &= \begin{cases} \frac{1}{2} \cdot (n(n+1)m - f(c,c-b-1,a,m-1)) \\ -h(c,c-b-1,a,m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ - 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.16 Numbers

Bernoulli numbers

$$\begin{split} &B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ &\sum_{j=0}^m \binom{m+1}{j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!} \text{.} \\ &S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k^+ n^{m+1-k} \end{split}$$

ullet Stirling numbers of the second kind Partitions of ndistinct elements into exactly k groups. S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
• Pentagonal number theorem
$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
• Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)^n+1} {kn \choose n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. $k \ j$:s s.t. $\pi(j) > \pi(j+1), \ k+1 \ j \colon \text{s.t.} \ \pi(j) \geq j, \ k \ j \colon \text{s.t.} \ \pi(j) > j.$ E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

6.17 GeneratingFunctions

- Ordinary Generating Function $A(x) = \sum_{i \geq 0} a_i x^i$
 - $A(rx) \Rightarrow r^n a_n$
 - $A(x) + B(x) \Rightarrow a_n + b_n$
 - $A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\cdots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x)' \Rightarrow na_n$
 - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i > 0} \frac{a_i}{i!} x_i$
 - $A(x)+B(x) \Rightarrow a_n+b_n$

 - $A^{(k)}(x) \Rightarrow a_{n+k}$ $A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} {n \choose i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x) \Rightarrow na_n$
- Special Generating Function
 - $(1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i$
 - $-\frac{1}{(1-x)^n} = \sum_{i>0} \binom{i}{n-1} x^i$

```
- S_k = \sum_{x=1}^n x^k: S = \sum_{p=0}^\infty x^p = \frac{e^x - e^{x(n+1)}}{1 - e^x}
```

7 Linear Algebra

7.1 GuassianElimination

```
#undef M
 struct matrix { //m variables, n equations
   int n, m;
   fraction M[N][N + 1], sol[N];
   int solve() { //-1: inconsistent, >= 0: rank
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
       for (int j = 0; j < n; ++j) {</pre>
         if (i == j) continue;
         fraction tmp = -M[j][piv] / M[i][piv];
         for (int k = 0; k <=</pre>
               m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
       }
     int rank = 0;
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m && M[i][m].n) return -1;
       else if (piv
             < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
     return rank;
|};
```

BerlekampMassey

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(output.size() + 1), me, he;
  for (int f = 0, i = 1; i <= output.size(); ++i) {</pre>
     for (int j = 0; j < me.size(); ++j)</pre>
       d[i] += output[i - j - 2] * me[j];
     if ((d[i] -= output[i - 1]) == 0) continue;
     if (me.empty()) {
       me.resize(f = i);
       continue;
     vector<T> o(i - f - 1);
     T k = -d[i] / d[f];
     o.emplace_back(-k);
     for (T x : he) o.emplace_back(x * k);
     o.resize(max(o.size(), me.size()));
     for (int j = 0; j < me.size(); ++j) o[j] += me[j];</pre>
     if (i - f + (int
         )he.size()) >= (int)me.size()) he = me, f = i;
  return me;
}
```

```
7.3 Simplex
    Standard form: maximize \mathbf{c}^T\mathbf{x} subject to A\mathbf{x} \leq \mathbf{b} and \mathbf{x} \geq 0.
Dual LP: minimize \mathbf{b}^T\mathbf{y} subject to \tilde{A^T}\mathbf{y} \geq \mathbf{c} and \mathbf{y} \geq \mathbf{0}. \bar{\mathbf{x}} and \bar{\mathbf{y}} are optimal if and only if for all i \in [1,n], either
ar{x}_i\!=\!0 or \sum_{j=1}^m A_{ji}ar{y}_j\!=\!c_i holds and for all i\!\in\![1,m] either ar{y}_i\!=\!0
or \sum_{j=1}^{n} A_{ij} \bar{x}_j = b_j holds.
1. In case of minimization, let c_i' = -c_i
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j
      • \sum_{1 \leq i \leq n}^{-} A_{ji} x_i \leq b_j
• \sum_{1\leq i\leq n}A_{ji}x_i\geq b_j
4. If x_i has no lower bound, replace x_i with x_i-x_i'
// using N + 2M variables
const int mxM = 25;
const int mxN = 25 + 2 * mxM;
struct simplex {
    const double inf = 1 / .0, eps = 1e-9;
   int n, m, k, var[mxN], inv[mxN], art[mxN];
```

#define base ll // complex<double>

#define N 524288

```
double A[mxM][mxN], B[mxM], x[mxN];
                                                             // const double PI = acosl(-1);
                                                             const 11 mod = 998244353, g = 3;
  void init(int _n) { n = _n, m = 0; }
                                                             base omega[4 * N], omega_[4 * N];
  void equation(vector<double> a, double b) {
    for (int i = 0; i < n; i++) A[m][i] = a[i];</pre>
                                                             int rev[4 * N];
    B[m] = b, var[m] = n + m, ++m;
                                                             ll fpow(ll b, ll p);
  void pivot(int r, int c, double bx) {
                                                             ll inverse(ll a) { return fpow(a, mod - 2); }
    for (int i = 0; i <= m + 1; i++)</pre>
       if (i != r && abs(A[i][c]) > eps) {
                                                             void calcW(int n) {
         x[var[i]] = bx * A[i][c] / A[i][var[i]];
                                                               ll r = fpow(g, (mod - 1) / n), invr = inverse(r);
         double f = A[i][c] / A[r][c];
                                                               omega[0] = omega_[0] = 1;
         for (int j = 0; j <= n + m + k; j++)</pre>
                                                               for (int i = 1; i < n; i++) {</pre>
           A[i][j] -= A[r][j] * f;
                                                                 omega[i] = omega[i - 1] * r % mod;
         B[i] -= B[r] * f;
                                                                 omega_[i] = omega_[i - 1] * invr % mod;
  }
                                                               // double arg = 2.0 * PI / n;
  double phase(int p) {
                                                               // for (int i = 0; i < n; i++)
    while (true) {
                                                               // {
      int in = min_element(
                                                               //
                                                                   omega[i] = base(cos(i * arg), sin(i * arg));
                  A[m + p], A[m + p] + n + m + k + 1) -
                                                               //
                                                                    omega_[i] = base(cos(-i * arg), sin(-i * arg));
         A[m + p];
                                                               // }
      if (A[m + p][in] >= -eps) break;
                                                             }
      double bx = inf;
      int piv = -1;
                                                             void calcrev(int n) {
      for (int i = 0; i < m; i++)
                                                               int k = __lg(n);
         if (A[i][in] > eps && B[i] / A[i][in] <= bx)</pre>
                                                               for (int i = 0; i < n; i++) rev[i] = 0;</pre>
           piv = i, bx = B[i] / A[i][in];
                                                               for (int i = 0; i < n; i++)</pre>
      if (piv == -1) return inf;
                                                                 for (int j = 0; j < k; j++)
      int out = var[piv];
                                                                   if (i & (1 << j)) rev[i] ^= 1 << (k - j - 1);</pre>
      pivot(piv, in, bx);
      x[out] = 0, x[in] = bx, var[piv] = in;
                                                             vector<base> NTT(vector<base> poly, bool inv) {
    return x[n + m];
                                                               base *w = (inv ? omega_ : omega);
  }
                                                               int n = poly.size();
  double solve(vector<double> c) {
                                                               for (int i = 0; i < n; i++)</pre>
    auto invert = [&](int r) {
                                                                 if (rev[i] > i) swap(poly[i], poly[rev[i]]);
      for (int j = 0; j <= n + m; j++) A[r][j] *= -1;</pre>
      B[r] *= -1;
                                                               for (int len = 1; len < n; len <<= 1) {</pre>
    };
                                                                 int arg = n / len / 2;
    k = 1:
                                                                 for (int i = 0; i < n; i += 2 * len)</pre>
    for (int i = 0; i < n; i++) A[m][i] = -c[i];</pre>
                                                                   for (int j = 0; j < len; j++) {</pre>
    fill(A[m + 1], A[m + 1] + mxN, 0.0);
                                                                     base odd =
    for (int i = 0; i <= m + 1; i++)</pre>
                                                                       w[j * arg] * poly[i + j + len] % mod;
      fill(A[i] + n, A[i] + n + m + 2, 0.0),
                                                                     poly[i + j + len] =
         var[i] = n + i, A[i][n + i] = 1;
                                                                        (poly[i + j] - odd + mod) % mod;
                                                                     poly[i + j] = (poly[i + j] + odd) \% mod;
    for (int i = 0; i < m; i++) {
      if (B[i] < 0) {
                                                               }
         ++k;
                                                               if (inv)
         for (int j = 0; j <= n + m; j++)</pre>
                                                                 for (auto &a : poly) a = a * inverse(n) % mod;
           A[m + 1][j] += A[i][j];
                                                               return poly;
         invert(i);
         var[i] = n + m + k, A[i][var[i]] = 1,
         art[var[i]] = n + i;
                                                             vector<base> mul(vector<base> f, vector<base> g) {
      }
                                                               int sz = 1 << (__lg(f.size() + g.size() - 1) + 1);</pre>
      x[var[i]] = B[i];
                                                               f.resize(sz), g.resize(sz);
                                                               calcrev(sz);
                                                               calcW(sz);
    phase(1):
                                                               f = NTT(f, 0), q = NTT(q, 0);
    if (*max_element(
                                                               for (int i = 0; i < sz; i++)</pre>
           x + (n + m + 2), x + (n + m + k + 1)) > eps)
                                                                 f[i] = f[i] * g[i] % mod;
      return .0 / .0;
                                                               return NTT(f, 1);
    for (int i = 0; i <= m; i++)</pre>
                                                            |}
       if (var[i] > n + m)
        var[i] = art[var[i]], invert(i);
                                                             8.2 FHWT
    k = 0;
    return phase(0);
                                                            | /* x: a[j], y: a[j + (L >> 1)]
                                                             or: (y += x * op), and: (x += y * op)
|} lp;
                                                             xor: (x, y = (x + y) * op, (x - y) * op)
                                                             invop: or, and, xor = -1, -1, 1/2 */
     Polynomials
                                                             void fwt(int *a, int n, int op) { //or
                                                               for (int L = 2; L <= n; L <<= 1)
8.1 NTT (FFT)
                                                                 for (int i = 0; i < n; i += L)</pre>
                                                                   for (int j = i; j < i + (L >> 1); ++j)
// 9223372036737335297, 3
```

|}

a[j + (L >> 1)] += a[j] * op;

```
const int N = 21;
int f[
    N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
void
    subset_convolution(int *a, int *b, int *c, int L) {
  // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
  int n = 1 << L;
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i \& (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)</pre>
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)</pre>
    for (int j = 0; j <= i; ++j)</pre>
       for (int x = 0; x < n; ++x)</pre>
         h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i)</pre>
    fwt(h[i], n, -1);
  for (int i = 0; i < n; ++i)</pre>
    c[i] = h[ct[i]][i];
}
```

8.3 PolynomialOperations

#define poly vector<ll>

```
poly inv(poly A) {
 A.resize(1 << (__lg(A.size() - 1) + 1));
 poly B = {inverse(A[0])};
 for (int n = 1; n < A.size(); n += n) {</pre>
   poly pA(A.begin(), A.begin() + 2 * n);
    calcrev(4 * n);
   calcW(4 * n);
    pA.resize(4 * n);
   B.resize(4 * n);
   pA = NTT(pA, 0);
    B = NTT(B, 0);
    for (int i = 0; i < 4 * n; i++)
      Blil =
        ((B[i] * 2 - pA[i] * B[i] % mod * B[i]) % mod +
        mod:
   B = NTT(B, 1);
   B.resize(2 * n);
 return B;
pair<poly, poly> div(poly A, poly B) {
 if (A.size() < B.size()) return make_pair(poly(), A);</pre>
 int n = A.size(), m = B.size();
 poly revA = A, invrevB = B;
 reverse(revA.begin(), revA.end());
 reverse(invrevB.begin(), invrevB.end());
 revA.resize(n - m + 1);
 invrevB.resize(n - m + 1);
 invrevB = inv(invrevB);
 poly Q = mul(revA, invrevB);
 Q.resize(n - m + 1);
 reverse(Q.begin(), Q.end());
 poly R = mul(Q, B);
 R.resize(m - 1);
 for (int i = 0; i < m - 1; i++)</pre>
    R[i] = (A[i] - R[i] + mod) \% mod;
 return make_pair(Q, R);
ll fast_kitamasa(ll k, poly A, poly C) {
 int n = A.size();
 C.emplace_back(mod - 1);
 poly Q, R = \{0, 1\}, F = \{1\};
 R = div(R, C);
 while (k) {
   if (k & 1) F = div(mul(F, R), C);
   R = div(mul(R, R), C);
   k >>= 1;
```

```
ll ans = 0;
   for (int i = 0; i < F.size(); i++)</pre>
     ans = (ans + A[i] * F[i]) % mod;
   return ans;
vector<ll> fpow(vector<ll> f, ll p, ll m) {
   int b = 0;
   while (b < f.size() && f[b] == 0) b++;</pre>
   f = vector<ll>(f.begin() + b, f.end());
   int n = f.size();
   f.emplace_back(0);
   vector<ll> q(min(m, b * p), 0);
   q.emplace_back(fpow(f[0], p));
   for (int k = 0; q.size() < m; k++) {</pre>
     ll res = 0;
     for (int i = 0; i < min(n, k + 1); i++)</pre>
       res = (res +
               p * (i + 1) % mod * f[i + 1] % mod *
                  q[k - i + b * p]) %
         mod;
     for (int i = 1; i < min(n, k + 1); i++)</pre>
       res = (res ·
                f[i] * (k - i + 1) % mod *
                  q[k - i + 1 + b * p]) %
     res = (res < 0 ? res + mod : res) *
       inv(f[0] * (k + 1) % mod) % mod;
     q.emplace_back(res);
  }
   return q;
}
```

8.4 NewtonMethod+MiscGF

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that F(P) = 0 can be found iteratively. Denote by Q_k the polynomial such that $F(Q_k) = 0 \pmod{x^{2^k}}$, then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$

- A^{-1} : $B_{k+1} = B_k(2 AB_k) \mod x^{2^{k+1}}$
- $\ln A$: $(\ln A)' = \frac{A'}{A}$
- $\exp A$: $B_{k+1} = B_k(1 + A \ln B_k) \mod x^{2^{k+1}}$
- \sqrt{A} : $B_{k+1} = \frac{1}{2}(B_k + AB_k^{-1}) \mod x^{2^{k+1}}$

Geometry

9.1 Basic

```
typedef pair<pdd, pdd> Line;
struct Cir{ pdd 0; double R; };
const double eps = 1e-8;
pdd operator+(pdd a, pdd b)
{ return pdd(a.F + b.S, a.S + b.S); }
pdd operator-(pdd a, pdd b)
{ return pdd(a.F - b.S, a.S - b.S); }
pdd operator*(pdd a, double b)
{ return pdd(a.F * b, a.S * b); }
pdd operator/(pdd a, double b)
{ return pdd(a.F / b, a.S / b); }
double dot(pdd a, pdd b)
{ return a.F * b.F + a.S * b.S; }
double cross(pdd a, pdd b)
{ return a.F * b.S - a.S * b.F; }
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
```

```
{ return sign(cross(b - a, c - a)); }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
  if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= 0;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
  if (a123 == 0 && a124 == 0)
    return btw(p1, p2, p3) || btw(p1, p2, p4) ||
btw(p3, p4, p1) || btw(p3, p4, p2);
  return a123 * a124 <= 0 && a341 * a342 <= 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
  double a124 = cross(p2 - p1, p4 - p1);
  return (p4
       * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2
pdd perp(pdd p1)
{ return pdd(-p1.S, p1.F); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
    ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
     (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) {
  pdd dp = p1 - p0
       , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  return q0 + pdd(
       cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
} // from line p0--p1 to q0--q1, apply to r
```

9.2 ConvexHull

9.3 SortByAngle

```
int cmp(pll a, pll b, bool same = true) {
#define is_neg(k) (
    sign(k.S) < 0 || (sign(k.S) == 0 && sign(k.F) < 0))
    int A = is_neg(a), B = is_neg(b);
    if (A != B)
        return A < B;
    if (sign(cross(a, b)) == 0)
        return same ? abs2(a) < abs2(b) : -1;
    return sign(cross(a, b)) > 0;
}
```

9.4 DisPointSegment

```
| double PointSegDist(pdd q0, pdd q1, pdd p) {
| if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
| if (sign(dot(q1 - q0,
| p - q0)) >= 0 && sign(dot(q0 - q1, p - q1)) >= 0)
| return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
| return min(abs(p - q0), abs(p - q1));
|}
```

9.5 PointInCircle

```
// return q'
    s relation with circumcircle of tri(p[0],p[1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
    __int128 det = 0;
    for (int i = 0; i < 3; ++i)
        det += __int128(abs2(p[i]) - abs2(q)) *
            cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

9.6 PointInConvex

9.7 PointTangentConvex

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
/* bool pred(int a, int b);
f(0) \sim f(n-1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int l = 0, r = n; bool rv = pred(1, 0);
  while (r - l > 1) {
    int m = (l + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
    else l = m;
  return pred(l, r % n) ? l : r % n;
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch((int)C.size(), [&](int x, int y)
    { return ori(p, C[x], C[y]) == s; });
  return pii(gao(1), gao(-1));
|} // return (a, b), ori(p, C[a], C[b]) >= 0
```

9.8 CircTangentCirc

```
vector<Line
      > go( const Cir& c1 , const Cir& c2 , int sign1 ){
   // sign1 = 1 for outer tang, -1 for inter tang
   vector<Line> ret;
   double d_{sq} = abs2(c1.0 - c2.0);
   if (sign(d_sq) == 0) return ret;
   double d = sqrt(d_sq);
   pdd v = (c2.0 - c1.0) / d;
   double c = (c1.R - sign1 * c2.R) / d;
   if (c * c > 1) return ret;
   double h = sqrt(max(0.0, 1.0 - c * c));
   for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
      pdd n = pdd(v.F * c - sign2 * h * v.S,
       v.S * c + sign2 * h * v.F);
      pdd p1 = c1.0 + n * c1.R;
      pdd p2 = c2.0 + n * (c2.R * sign1);
     if (sign(p1.F - p2.F) == 0 and
    sign(p1.S - p2.S) == 0)
        p2 = p1 + perp(c2.0 - c1.0);
```

```
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    ret.emplace_back(Line(p1, p2));
                                                             if (abs(pa) < abs(pb)) swap(pa, pb);</pre>
                                                             if (abs(pb) < eps) return 0;</pre>
  return ret;
                                                             double S, h, theta;
|}
                                                             double a = abs(pb), b = abs(pa), c = abs(pb - pa);
                                                             double cosB = dot(pb, pb - pa) / a / c,
      LineCircleIntersect
                                                                    B = acos(cosB);
                                                             double cosC = dot(pa, pb) / a / b, C = acos(cosC);
vector<pdd> circleLine(pdd c, double r, pdd a, pdd b) {
                                                             if (a > r) {
  pdd p
                                                               S = (C / 2) * r * r;
        = a + (b - a) * dot(c - a, b - a) / abs2(b - a);
                                                               h = a * b * sin(C) / c;
  double s = cross
                                                               if (h < r && B < PI / 2)
       (b - a, c - a), h2 = r * r - s * s / abs2(b - a);
                                                                 S = (acos(h / r) * r * r -
  if (h2 < 0) return {};
                                                                   h * sqrt(r * r - h * h));
  if (h2 == 0) return {p};
                                                             } else if (b > r) {
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
                                                               theta = PI - B - asin(sin(B) / r * a);
  return {p - h, p + h};
                                                               S = .5 * a * r * sin(theta) +
| }
                                                                 (C - theta) / 2 * r * r;
9.10
        LineConvexIntersect
                                                             } else S = .5 * sin(C) * a * b;
int TangentDir(vector<pll> &C, pll dir) {
   return cyc_tsearch((int)C.size(), [&](int a, int b) {
                                                           double area_poly_circle(const vector<pdd> poly,
    return cross(dir, C[a]) > cross(dir, C[b]);
                                                             const pdd &0, const double r) {
                                                             double S = 0;
                                                             for (int i = 0; i < (int)poly.size(); ++i)</pre>
                                                               S += _area(poly[i] - 0,
#define cmpL(i) sign(cross(C[i] - a, b - a))
                                                                      poly[(i + 1) % (int)poly.size()] - 0, r) *
pii lineHull(pll a, pll b, vector<pll> &C) {
                                                                 ori(
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
                                                                   0, poly[i], poly[(i + 1) % (int)poly.size()]);
                                                             return fabs(S);
  int n = (int)C.size();
  if (cmpL(A) < 0 \mid \mid cmpL(B) > 0)
    return pii(-1, -1); // no collision
                                                                   MinkowskiSum
                                                           9.13
  auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r; ) {
                                                           vector<pll> Minkowski
       int m = ((l + r + (l < r? 0 : n)) / 2) % n;
                                                                (vector<pll> A, vector<pll> B) { // |A|, |B|>=3
       (cmpL(m) == cmpL(t) ? l : r) = m;
                                                             hull(A), hull(B);
    }
                                                             vector<pll> C(1, A[0] + B[0]), s1, s2;
    return (l + !cmpL(r)) % n;
                                                             for (int i = 0; i < A.size(); ++i)</pre>
  };
                                                               s1.emplace_back(A[(i + 1) % A.size()] - A[i]);
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
                                                             for (int i = 0; i < B.size(); i++)</pre>
  if (res.F == res.S) // touching the corner i
                                                               s2.emplace_back(B[(i + 1) % B.size()] - B[i]);
    return pii(res.F, -1);
                                                             for (int i = 0, j = 0; i < A.size() || j < B.size();)</pre>
  if (!
                                                               if (j >= B.size()
       cmpL(res.F) && !cmpL(res.S)) // along side i, i+1
                                                                     || (i < A.size() \&\& cross(s1[i], s2[j]) >= 0))
    switch ((res.F - res.S + n + 1) % n) {
                                                                 C.emplace_back(B[j % B.size()] + A[i++]);
       case 0: return pii(res.F, res.F);
                                                               else
      case 2: return pii(res.S, res.S);
                                                                 C.emplace_back(A[i % A.size()] + B[j++]);
                                                             return hull(C), C;
  /* crossing sides (i, i+1) and (j, j+1)
                                                          }
  crossing corner i is treated as side (i, i+1)
  returned
                                                           9.14 MinMaxEnclosingRect
       in the same order as the line hits the convex */
  return res;
                                                           const double INF = 1e18, qi = acos(-1) / 2 * 3;
|} // convex cut: (r, l]
                                                           pdd solve(vector<pll> &dots) {
                                                           #define diff(u, v) (dots[u] - dots[v])
9.11 CircIntersectCirc
                                                           #define vec(v) (dots[v] - dots[i])
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
                                                             hull(dots);
                                                             double Max = 0, Min = INF, deg;
  pdd o1 = a.0, o2 = b.0;
   double r1 =
                                                             int n = (int)dots.size();
        a.R, r2 = b.R, d2 = abs2(o1 - o2), d = sqrt(d2);
                                                             dots.emplace_back(dots[0]);
  if(d < max
                                                             for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
       (r1, r2) - min(r1, r2) \mid | d > r1 + r2) return 0;
                                                               pll nw = vec(i + 1);
  pdd \ U = (o1 + o2) * 0.5
                                                               while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
       + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
                                                                 u = (u + 1) \% n;
  double A = sqrt((r1 + r2 + d) *
                                                               while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
```

r = (r + 1) % n;if (!i) l = (r + 1) % n;

l = (l + 1) % n;

deg = acos(dot(diff(r

deg = (qi - deg) / 2;

Max = max(Max, abs(diff))

while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))

Min = min(Min, (double)(dot(nw, vec(r)) - dot

(nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));

, l), vec(u) / abs(diff(r, l)) / abs(vec(u)));

(r, l)) * abs(vec(u)) * sin(deg) * sin(deg));

9.12 PolyIntersectCirc

p1 = v + v, p2 = v - v;

return 1;

|}

```
// Divides into multiple triangle, and sum up
const double PI = acos(-1);
double _area(pdd pa, pdd pb, double r) {
```

(r1 - r2 + d) * (r1 + r2 - d) * (-r1 + r2 + d));

= pdd(o1.S - o2.S, -o1.F + o2.F) * A / (2 * d2);

double B =

atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);

```
}
                                                                       eve[E++] = Teve
  return pdd(Min, Max);
                                                                            (bb, B, 1), eve[E++] = Teve(aa, A, -1);
| }
                                                                       if(B > A) ++cnt;
9.15 MinEnclosingCircle
                                                                   if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
                                                                   else{
pdd Minimum_Enclosing_Circle
                                                                     sort(eve, eve + E);
     (vector<pdd> dots, double &r) {
   pdd cent;
                                                                     eve[E] = eve[0];
   random_shuffle(all(dots));
                                                                     for(int j = 0; j < E; ++j){</pre>
                                                                       cnt += eve[j].add;
  cent = dots[0], r = 0;
  for (int i = 1; i < (int)dots.size(); ++i)</pre>
     if (abs(dots[i] - cent) > r) {
                                                                            ] += cross(eve[j].p, eve[j + 1].p) * .5;
                                                                       double theta = eve[j + 1].ang - eve[j].ang;
       cent = dots[i], r = 0;
       for (int j = 0; j < i; ++j)
                                                                       if (theta < 0) theta += 2. * pi;</pre>
         if (abs(dots[j] - cent) > r) {
                                                                       Area[cnt] += (theta
                                                                             - sin(theta)) * c[i].R * c[i].R * .5;
           cent = (dots[i] + dots[j]) / 2;
           r = abs(dots[i] - cent);
           for(int k = 0; k < j; ++k)</pre>
                                                                   }
                                                                 }
             if(abs(dots[k] - cent) > r)
               cent = excenter
                                                              }
                   (dots[i], dots[j], dots[k], r);
                                                            |};
         }
                                                             9.17
                                                                     LineCmp
  return cent;
| }
                                                            using Line = pair<pll, pll>;
                                                             struct lineCmp {
9.16 CircleCover
                                                               bool operator()(Line l1, Line l2) const {
                                                                 int X =
const int N = 1021:
                                                                   (\max(l1.F.F, l2.F.F) + \min(l1.S.F, l2.S.F)) / 2;
struct CircleCover {
                                                                 ll p1 =
  int C;
                                                                      (X - l1.F.F) * l1.S.S + (l1.S.F - X) * l1.F.S,
  Cir c[N];
                                                                    p2 =
  bool g[N][N], overlap[N][N];
                                                                      (X - 12.F.F) * 12.S.S + (12.S.F - X) * 12.F.S,
  // Area[i] : area covered by at least i circles
                                                                    q1 = (l1.S.F - l1.F.F), q2 = (l2.S.F - l2.F.F);
  double Area[ N ];
                                                                 if (q1 == 0) p1 = l1.F.S + l1.S.S, q1 = 2;
   void init(int _C){ C = _C;}
                                                                 if (q2 == 0) p2 = l2.F.S + l2.S.S, q2 = 2;
   struct Teve {
                                                                 if (l1.F == l2.F || l2.F == l2.S) l1 = l2;
     pdd p; double ang; int add;
                                                                 return make_tuple((__int128)(p1 * q2), l1) <</pre>
     Teve() {}
                                                                   make_tuple((__int128)(p2 * q1), l2);
     Teve(pdd _a
         , double _b, int _c):p(_a), ang(_b), add(_c){}
                                                            };
    bool operator<(const Teve &a)const
     {return ang < a.ang;}
                                                             9.18
                                                                    Trapezoidalization
  eve[N * 2];
   // strict: x = 0, otherwise x = -1
                                                            | struct SweepLine {
  bool disjuct(Cir &a, Cir &b, int x)
                                                               struct cmp {
   {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
                                                                 cmp(const SweepLine &_swp): swp(_swp) {}
  bool contain(Cir &a, Cir &b, int x)
                                                                 bool operator()(int a, int b) const {
   {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
                                                                   if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
  bool contain(int i, int j) {
                                                                     return swp.slope_cmp(a, b);
     /* c[j] is non-strictly in c[i]. */
                                                                   return swp.get_y(a) + swp.eps < swp.get_y(b);</pre>
    return (sign
         (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
                                                                 const SweepLine &swp;
         R) == 0 \& i < j) && contain(c[i], c[j], -1);
                                                                 _cmp;
                                                               T curTime, eps, curQ;
                                                               vector<Line> base;
  void solve(){
                                                               multiset<int, cmp> sweep;
     fill_n(Area, C + 2, 0);
                                                               multiset<pair<T, int>> event;
     for(int i = 0; i < C; ++i)</pre>
                                                               vector<typename multiset<int, cmp>::iterator> its;
       for(int j = 0; j < C; ++j)</pre>
                                                               vector
         overlap[i][j] = contain(i, j);
                                                                   <typename multiset<pair<T, int>>::iterator> eits;
     for(int i = 0; i < C; ++i)</pre>
                                                               bool slope_cmp(int a, int b) const {
       for(int j = 0; j < C; ++j)</pre>
                                                                 assert(a != -1);
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                                                                 if (b == -1) return 0;
             disjuct(c[i], c[j], -1));
                                                                 return sign(cross(base
     for(int i = 0; i < C; ++i){</pre>
                                                                     [a].Y - base[a].X, base[b].Y - base[b].X)) < 0;
       int E = 0, cnt = 1;
       for(int j = 0; j < C; ++j)</pre>
                                                               T get_y(int idx) const {
         if(j != i && overlap[j][i])
                                                                 if (idx == -1) return curQ;
           ++cnt;
                                                                 Line l = base[idx];
       for(int j = 0; j < C; ++j)</pre>
                                                                 if (l.X.X == l.Y.X) return l.Y.Y;
         if(i != j && g[i][j]) {
                                                                 return ((curTime - l.X.X) * l.Y.Y
           pdd aa, bb;
                                                                     + (l.Y.X - curTime) * l.X.Y) / (l.Y.X - l.X.X);
           CCinter(c[i], c[j], aa, bb);
           double A =
                                                               void insert(int idx) {
                atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
```

its[idx] = sweep.insert(idx);

if (its[idx] != sweep.begin())

```
update_event(*prev(its[idx]));
                                                        |};
  update_event(idx);
                                                         9.19 TriangleHearts
  event.emplace
      (base[idx].Y.X, idx + 2 * (int)base.size());
                                                        | pdd circenter(
}
                                                           pdd p0, pdd p1, pdd p2) { // radius = abs(center)
void erase(int idx) {
                                                           p1 = p1 - p0, p2 = p2 - p0;
double x1 = p1.F, y1 = p1.S, x2 = p2.F, y2 = p2.S;
  assert(eits[idx] == event.end());
  auto p = sweep.erase(its[idx]);
                                                           double m = 2. * (x1 * y2 - y1 * x2);
  its[idx] = sweep.end();
                                                           pdd center = pdd((x1 * x1 * y2 - x2 * x2 * y1 +
  if (p != sweep.begin())
                                                                               y1 * y2 * (y1 - y2)) / m,
    update_event(*prev(p));
                                                              (x1 * x2 * (x2 - x1) - y1 * y1 * x2 +
}
                                                                x1 * y2 * y2) / m);
void update_event(int idx) {
                                                           return center + p0;
  if (eits[idx] != event.end())
                                                         }
    event.erase(eits[idx]);
                                                         pdd incenter(
  eits[idx] = event.end();
                                                           pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
  auto nxt = next(its[idx]);
                                                           double a = abs(p2 - p3), b = abs(p1 - p3),
  if (nxt ==
                                                                  c = abs(p1 - p2);
       sweep.end() || !slope_cmp(idx, *nxt)) return;
                                                           double s = a + b + c;
  auto t = intersect(base[idx].
                                                           return (a * p1 + b * p2 + c * p3) / s;
      X, base[idx].Y, base[*nxt].X, base[*nxt].Y).X;
  if (t + eps < curTime || t</pre>
                                                         pdd masscenter(pdd p1, pdd p2, pdd p3) {
       >= min(base[idx].Y.X, base[*nxt].Y.X)) return;
                                                           return (p1 + p2 + p3) / 3;
  eits[idx
      ] = event.emplace(t, idx + (int)base.size());
                                                         pdd orthcenter(pdd p1, pdd p2, pdd p3) {
}
                                                           return masscenter(p1, p2, p3) * 3 -
void swp(int idx) {
                                                             circenter(p1, p2, p3) * 2;
  assert(eits[idx] != event.end());
                                                         |}
  eits[idx] = event.end();
                                                          9.20 HalfPlaneIntersect
  int nxt = *next(its[idx]);
  swap((int&)*its[idx], (int&)*its[nxt]);
                                                         pll area_pair(Line a, Line b)
  swap(its[idx], its[nxt]);
                                                         { return pll(cross(a.S
  if (its[nxt] != sweep.begin())
                                                               - a.F, b.F - a.F), cross(a.S - a.F, b.S - a.F)); }
    update_event(*prev(its[nxt]));
                                                         bool isin(Line l0, Line l1, Line l2) {
  update_event(idx);
                                                           // Check inter(l1, l2) strictly in l0
                                                           auto [a02X, a02Y] = area_pair(l0, l2);
// only expected to call the functions below
                                                           auto [a12X, a12Y] = area_pair(l1, l2);
SweepLine(T t, T e, vector<Line> vec): _cmp
                                                           if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
    (*this), curTime(t), eps(e), curQ(), base(vec),
                                                           return (__int128)
     sweep(_cmp), event(), its((int)vec.size(), sweep
                                                                 a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
    .end()), eits((int)vec.size(), event.end()) {
  for (int i = 0; i < (int)base.size(); ++i) {</pre>
                                                         /* Having solution, check size > 2 */
    auto &[p, q] = base[i];
                                                         /* --^-- Line.X --^-- Line.Y --^-- */
    if (p > q) swap(p, q);
                                                         vector<Line> halfPlaneInter(vector<Line> arr) {
    if (p.X <= curTime && curTime <= q.X)</pre>
                                                           sort(all(arr), [&](Line a, Line b) -> int {
      insert(i);
                                                              if (cmp(a.S - a.F, b.S - b.F, 0) != -1)
    else if (curTime < p.X)</pre>
                                                                return cmp(a.S - a.F, b.S - b.F, 0);
      event.emplace(p.X, i);
                                                             return ori(a.F, a.S, b.S) < 0;</pre>
}
                                                           deque<Line> dq(1, arr[0]);
void setTime(T t, bool ers = false) {
                                                           for (auto p : arr) {
  assert(t >= curTime);
                                                              if (cmp(
  while (!event.empty() && event.begin()->X <= t) {</pre>
                                                                  dq.back().S - dq.back().F, p.S - p.F, 0) == -1)
    auto [et, idx] = *event.begin();
                                                                continue;
    int s = idx / (int)base.size();
                                                              while ((int)dq.size() >= 2
    idx %= (int)base.size();
                                                                  && !isin(p, dq[(int)dq.size() - 2], dq.back()))
    if (abs(et - t) <= eps && s == 2 && !ers) break;</pre>
                                                                dq.pop_back();
    curTime = et;
                                                              while
    event.erase(event.begin());
                                                                  ((int)dq.size() >= 2 \&\& !isin(p, dq[0], dq[1]))
    if (s == 2) erase(idx);
                                                                dq.pop_front();
    else if (s == 1) swp(idx);
                                                             dq.emplace_back(p);
    else insert(idx);
  }
                                                           while ((int)dq.size() >= 3 &&
  curTime = t;
                                                                 !isin(dq[0], dq[(int)dq.size() - 2], dq.back()))
}
                                                              dq.pop_back();
T nextEvent() {
                                                           while ((int)
  if (event.empty()) return INF;
                                                                dq.size() >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
  return event.begin()->X;
                                                              dq.pop_front();
                                                           return vector<Line>(all(dq));
int lower_bound(T y) {
  curQ = y;
  auto p = sweep.lower_bound(-1);
                                                          9.21 RotatingSweepLine
  if (p == sweep.end()) return -1;
  return *p;
                                                         void rotatingSweepLine(vector<pii> &ps) {
```

int n = (int)ps.size(), m = 0;

```
vector<int> id(n), pos(n);
   vector<pii> line(n * (n - 1));
  for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++j)</pre>
       if (i != j) line[m++] = pii(i, j);
   sort(all(line), [&](pii a, pii b) {
     return cmp(ps[a.S] - ps[a.F], ps[b.S] - ps[b.F]);
  }); // cmp(): polar angle compare
   iota(all(id), 0);
   sort(all(id), [&](int a, int b) {
     if (ps[a].S != ps[b].S) return ps[a].S < ps[b].S;</pre>
     return ps[a] < ps[b];</pre>
  }); // initial order, since (1, 0) is the smallest
  for (int i = 0; i < n; ++i) pos[id[i]] = i;</pre>
  for (int i = 0; i < m; ++i) {</pre>
     auto l = line[i];
     // do something
     tie(pos[l.F], pos[l.S], id[pos[l.F]], id[pos[l.S
         ]]) = make_tuple(pos[l.S], pos[l.F], l.S, l.F);
  }
| }
```

9.22 DelaunayTriangulation

```
'* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const
     Il inf = MAXC * MAXC * 100; // Lower_bound unknown
struct Tri;
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
 Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
 pll p[3];
  Edge edge[3];
 Tri* chd[3];
 Tri() {}
 Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
 bool has_chd() const { return chd[0] != 0; }
 int num_chd() const {
   return !!chd[0] + !!chd[1] + !!chd[2];
 bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)</pre>
      if (ori(p[i], p[(i + 1) % 3], q) < 0)
        return 0;
   return 1;
} pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
 Triq() {
    the_root
         = // Tri should at least contain all points
      new(tris++) Tri(pll(-inf, -inf),
           pll(inf + inf, -inf), pll(-inf, inf + inf));
 Tri* find(pll p) { return find(the_root, p); }
```

```
void add_point(const
       pll &p) { add_point(find(the_root, p), p); }
  Tri* the_root;
  static Tri* find(Tri* root, const pll &p) {
    while (1) {
      if (!root->has_chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break;
        }
    }
    assert(0); // "point not found"
  void add_point(Tri* root, pll const& p) {
    /* split it into three triangles */
    for (int i = 0; i < 3; ++i)
      t[i] = new(tris
          ++) Tri(root->p[i], root->p[(i + 1) % 3], p);
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 0), Edge(t[(i + 1) % 3], 1));
    for (int i = 0; i < 3; ++i)
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)</pre>
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edqe[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p
        [0], tri->p[1], tri->p[2], trj->p[pj])) return;
    /* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri
        (tri->p[(pi + 1) % 3], trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri
        (trj->p[(pj + 1) % 3], tri->p[pi], trj->p[pj]);
    edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
    edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd
        [0] = trk; tri->chd[1] = trl; tri->chd[2] = 0;
    trj->chd
        [0] = trk; trj->chd[1] = trl; trj->chd[2] = 0;
    flip(trk, 1); flip(trk, 2);
    flip(trl, 1); flip(trl, 2);
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
  if (vst.find(now) != vst.end())
    return;
  vst.insert(now);
  if (!now->has_chd())
    return triang.emplace_back(now);
  for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(ps, ps + n);
  Trig tri; // the triangulation structure
  for (int i = 0; i < n; ++i)</pre>
    tri.add_point(ps[i]);
  go(tri.the_root);
9.23 VonoroiDiagram
```

```
// all coord. is even
      you may want to call halfPlaneInter after then
vector<vector<Line>> vec;
void build_voronoi_line(int n, pll *arr) {
  tool.init(n, arr); // Delaunay
  vec.clear(), vec.resize(n);
  for (int i = 0; i < n; ++i)</pre>
    for (auto e : tool.head[i]) {
      int u = tool.oidx[i], v = tool.oidx[e.id];
      pll m = (arr[v
           ] + arr[u]) / 2LL, d = perp(arr[v] - arr[u]);
       vec[u].emplace_back(Line(m, m + d));
|}
```

10 Misc

10.1 MoAlgoWithModify

```
.
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
  int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
     L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
  bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     if (RBid != q.RBid) return RBid < q.RBid;</pre>
     return T < b.T;</pre>
  }
};
void solve(vector<Query> query) {
  sort(ALL(query));
  int L=0, R=0, T=-1;
  for (auto q : query) {
     while (T < q.T) addTime(L, R, ++T); // TODO
     while (T > q.T) subTime(L, R, T--); // TODO
     while (R < q.R) add(arr[++R]); // TODO</pre>
     while (L > q.L) add(arr[--L]); // TODO
     while (R > q.R) sub(arr[R--]); // TODO
     while (L < q.L) sub(arr[L++]); // TODO</pre>
     // answer query
| }
```

10.2 MoAlgoOnTree

```
Mo's Algorithm On Tree
Preprocess:

    LCA

2) dfs with in[u] = dft++, out[u] = dft++
3) ord[in[u]] = ord[out[u]] = u
4) bitset<MAXN> inset
*/
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
    int c = LCA(u, v);
    if (c == u || c == v)
      q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
    else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
    else
      q.lca = c, q.L = out[v], q.R = in[v];
    q.Lid = q.L / blk;
  }
  bool operator<(const Query &q) const {</pre>
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    return R < q.R;</pre>
};
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
    else add(arr[x]); // TODO
    inset[x] = ~inset[x];
```

```
| }
void solve(vector<Query> query) {
   sort(ALL(query));
   int L = 0, R = 0;
   for (auto q : query) {
     while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
     while (R > q.R) flip(ord[R--]);
     while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
}
```

10.3 MoAlgoAdvanced

- Mo's Algorithm With Addition Only
 - Sort querys same as the normal Mo's algorithm.

 - For each query [l,r]:
 If l/blk = r/blk, brute-force.
 - If $l/blk \neq curL/blk$, initialize $curL := (l/blk+1) \cdot blk$, curR :=curL-1
 - If r > curR, increase curR
 - decrease curL to fit l, and then undo after answering
- · Mo's Algorithm With Offline Second Time
 - Require: Changing answer \equiv adding f([l,r],r+1).
 - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1) .
 - Part1: Answer all f([1,r],r+1) first.
 - Part2: Store $curR \rightarrow R$ for curL (reduce the space to $\mathcal{O}(N)$), and then answer them by the second offline algorithm.
 - Note: You must do the above symmetrically for the left boundaries.

10.4 HilbertCurve

```
ll hilbert(int n, int x, int y) {
  ll res = 0;
   for (int s = n / 2; s; s >>= 1) {
     int rx = (x \& s) > 0;
     int ry = (y \& s) > 0;
     res += s * 111 * s * ((3 * rx) ^ ry);
     if (ry == 0) {
       if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
  }
  return res;
| \} // n = 2^k
```

10.5 SternBrocotTree

- Construction: Root $\frac{1}{1}$, left/right neighbor $\frac{0}{1},\frac{1}{0}$, each node is sum of last left/right neighbor: $\frac{a}{b},\frac{c}{d} \rightarrow \frac{a+c}{b+d}$
- Property: Adjacent (mid-order DFS) $\frac{a}{b},\frac{c}{d}\!\Rightarrow\!bc\!-\!ad\!=\!1$.
- Search known $\frac{p}{a}$: keep L-R alternative. Each step can calcaulated in $O(1) \Rightarrow$ total $O(\log C)$.
- Search unknown $\frac{p}{a}$: keep L-R alternative. Each step can calcaulated in $O(\log C)$ checks \Rightarrow total $O(\log^2 C)$ checks.

10.6 AULLCS

```
void all_lcs(string s, string t) { // O-base
   vector<int> h((int)t.size());
   iota(all(h), 0);
   for (int a = 0; a < (int)s.size(); ++a) {</pre>
     int v = -1;
     for (int c = 0; c < (int)t.size(); ++c)</pre>
       if (s[a] == t[c] || h[c] < v)
         swap(h[c], v);
     // LCS(s[0, a], t[b, c]) =
     // c - b + 1 - sum([h[i] >= b] | i <= c)
     // h[i] might become -1 !!
  }
}
```

10.7 SimulatedAnnealing

```
| double factor = 100000;
| const int base = 1e9; // remember to run ~ 10 times
| for (int it = 1; it <= 1000000; ++it) {
| // ans:
| answer, nw: current value, rnd(): mt19937 rnd()
| if (exp(-(nw - ans
| ) / factor) >= (double)(rnd() % base) / base)
| ans = nw;
| factor *= 0.99995;
| }
```

if (x != begin() && isect(--x, y)) isect(x, y = erase(y)); while ((y = x) != begin() && (--x)->p >= y->p) isect(x, erase(y)); } ll query(ll x) { assert(!empty()); auto l = *lower_bound(x); return l.k * x + l.m; };

10.8 SMAWK

```
int opt[N];
ll A(int x, int y); // target func
 void smawk(vector<int> &r, vector<int> &c);
 void interpolate(vector<int> &r, vector<int> &c) {
   int n = (int)r.size();
   vector<int> er;
   for (int i = 1; i < n; i += 2) er.emplace_back(r[i]);</pre>
   smawk(er, c):
   for (int i = 0, j = 0; j < c.size(); j++) {</pre>
      \textbf{if} \ (\texttt{A}(\texttt{r[i]}, \ \texttt{c[j]}) < \texttt{A}(\texttt{r[i]}, \ \texttt{opt[r[i]]})) 
       opt[r[i]] = c[i];
     if (i + 2 < n \&\& c[j] == opt[r[i + 1]])
       j--, i += 2;
  }
 }
 void reduce(vector<int> &r, vector<int> &c) {
   int n = (int)r.size();
   vector<int> nc;
   for (int i : c) {
     int j = (int)nc.size();
     while (
       j \& A(r[j-1], nc[j-1]) > A(r[j-1], i))
       nc.pop_back(), j--;
     if (nc.size() < n) nc.emplace_back(i);</pre>
   }
   smawk(r, nc);
 void smawk(vector<int> &r, vector<int> &c) {
   if (r.size() == 1 && c.size() == 1) opt[r[0]] = c[0];
   else if (r.size() >= c.size()) interpolate(r, c);
   else reduce(r, c);
| }
```

10.9 Python

|math.isqrt(2) # integer sqrt

10.10 LineContainer

```
struct Line {
  mutable ll k, m, p;
 bool operator<(const Line &o) const {</pre>
    return k < o.k;</pre>
 }
 bool operator<(ll x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>>> {
 // (for doubles, use inf = 1/.0, div(a,b) = a/b)
static const ll inf = LLONG_MAX;
 ll div(ll a, ll b) { // floored division
    return a / b - ((a ^ b) < 0 && a % b);
 bool isect(iterator x, iterator y) {
    if (y == end()) return x->p = inf, 0;
    if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
    else x->p = div(y->m - x->m, x->k - y->k);
    return x->p >= y->p;
 }
 void add(ll k, ll m) {
    auto z = insert({k, m, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
```