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```

1 Basic

1.1 .vimrc

1.2 PBDS

```
#include <bits/extc++.h>
using namespace __gnu_pbds;

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
tree<int, null_type, less<int>, rb_tree_tag
    , tree_order_statistics_node_update> bst;
// order_of_key(n): # of elements <= n
// find_by_order(n): 0-indexed

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
__gnu_pbds::priority_queue
    <int, greater<int>, pairing_heap_tag> pq;
```

1.3 pargma

```
#pragma GCC optimize("Ofast,unroll-loops")
#pragma GCC target("avx,avx2,sse,sse2
    ,sse3,ssse3,sse4,popcnt,abm,mmx,fma,tune=native")
```

2 Graph

11

2.1 2SAT/SCC

```
struct SAT { // 0-base
  int low[N], dfn[N], bln[N], n, Time, nScc;
   bool instack[N], istrue[N];
  stack<int> st;
  vector < int > G[N], SCC[N];
  void init(int _n) {
     n = _n; // assert(n * 2 <= N);
     for (int i = 0; i < n + n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b) { G[a].emplace_back(b); }
   int rv(int a) {
     if (a >= n) return a - n;
     return a + n;
  void add_clause(int a, int b) {
     add_edge(rv(a), b), add_edge(rv(b), a);
  void dfs(int u) {
     dfn[u] = low[u] = ++Time;
     instack[u] = 1, st.push(u);
     for (int i : G[u])
       if (!dfn[i])
         dfs(i), low[u] = min(low[i], low[u]);
       else if (instack[i] && dfn[i] < dfn[u])</pre>
         low[u] = min(low[u], dfn[i]);
     if (low[u] == dfn[u]) {
       int tmp;
       do {
         tmp = st.top(), st.pop();
instack[tmp] = 0, bln[tmp] = nScc;
       } while (tmp != u);
  bool solve() {
     Time = nScc = 0;
     for (int i = 0; i < n + n; ++i)</pre>
     SCC[i].clear(), low[i] = dfn[i] = bln[i] = 0;

for (int i = 0; i < n + n; ++i)

if (!dfn[i]) dfs(i);
     for (int i =
          0; i < n + n; ++i) SCC[bln[i]].emplace_back(i);
     for (int i = 0; i < n; ++i) {
  if (bln[i] == bln[i + n]) return false;</pre>
       istrue[i] = bln[i] < bln[i + n];</pre>
       istrue[i + n] = !istrue[i];
     return true;
  }
};
```

2.2 BCC Vertex

```
int n, m, dfn[N], low[N], is_cut[N], nbcc = 0, t = 0;
vector < int > g[N], bcc[N], G[2 * N];
stack<int> st;
void tarjan(int p, int lp) {
  dfn[p] = low[p] = ++t;
  st.push(p);
  for (auto i : g[p]) {
     if (!dfn[i]) {
       tarjan(i, p);
        low[p] = min(low[p], low[i]);
       if (dfn[p] <= low[i]) {</pre>
          nbcc++
          is_cut[p] = 1;
          for (int x = 0; x != i; st.pop()) {
            x = st.top();
             bcc[nbcc].push_back(x);
          bcc[nbcc].push_back(p);
     } else low[p] = min(low[p], dfn[i]);
  }
void build() { // [n+1,n+nbcc] cycle, [1,n] vertex
for (int i = 1; i <= nbcc; i++) {
   for (auto j : bcc[i]) {</pre>
       G[i + n].push_back(j);
       G[j].push_back(i + n);
  }
```

2.3 MinimumMeanCycle

```
ll road[N][N]; // input here
struct MinimumMeanCycle {
  ll dp[N + 5][N], n;
  pll solve() {
    ll a = -1, b = -1, L = n + 1;
    for (int i = 2; i <= L; ++i)</pre>
       for (int k = 0; k < n; ++k)
         for (int j = 0; j < n; ++j)
           dp[i][j] =
             min(dp[i - 1][k] + road[k][j], dp[i][j]);
    for (int i = 0; i < n; ++i) {</pre>
       if (dp[L][i] >= INF) continue;
       ll ta = 0, tb = 1;
      for (int j = 1; j < n; ++j)
         if (dp[j][i] < INF &&
           ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
ta = dp[L][i] - dp[j][i], tb = L - j;
       if (ta == 0) continue;
      if (a == -1 || a * tb > ta * b) a = ta, b = tb;
    if (a != -1) {
      ll g = __gcd(a, b);
return pll(a / g, b / g);
    return pll(-1LL, -1LL);
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)</pre>
       for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;</pre>
};
```

2.4 MaximumCliqueDyn

```
struct MaxClique { // fast when N <= 100</pre>
 bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
   G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int l, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(all(r)
          , [&](int x, int y) { return d[x] > d[y]; });
    vector<int> c(r.size());
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first
          (); p < N; p = cs[k]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
 }
  void dfs(vector<</pre>
      int> &r, vector<int> &c, int l, bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.emplace_back(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
   }
  int solve() {
   vector<int> r(n);
    ans = q = 0, iota(all(r), 0);
```

```
pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
}
};
```

2.5 MinimumSteinerTree

```
struct SteinerTree { // 0-base
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcst[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {</pre>
       fill_n(dst[i], n, INF);
       dst[i][i] = vcst[i] = 0;
    }
  void chmin(int &x, int val) {
    x = min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)</pre>
       for (int i = 0; i < n; ++i)</pre>
         for (int j = 0; j < n; ++j)</pre>
           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
    shortest_path();
     int t = SZ(ter), full = (1 << t) - 1;</pre>
    for (int i = 0; i <= full; ++i)</pre>
       fill_n(dp[i], n, INF);
     copy_n(vcst, n, dp[0]);
     for (int msk = 1; msk <= full; ++msk) {</pre>
       if (!(msk & (msk - 1))) {
         int who = __lg(msk);
for (int i = 0; i < n; ++i)</pre>
           dp[msk
               ][i] = vcst[ter[who]] + dst[ter[who]][i];
       for (int i = 0; i < n; ++i)</pre>
         for (int sub = (
             msk - 1) & msk; sub; sub = (sub - 1) & msk)
           chmin(dp[msk][i],
               dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
       for (int i = 0; i < n; ++i) {</pre>
         tdst[i] = INF;
         for (int j = 0; j < n; ++j)</pre>
           chmin(tdst[i], dp[msk][j] + dst[j][i]);
       copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
}; // O(V 3^T + V^2 2^T)
```

2.6 DMST(slow)

```
struct zhu_liu { // O(VE)
 struct edge {
    int u, v;
    ll w;
 }:
 vector<edge> E; // 0-base
 int pe[N], id[N], vis[N];
 ll in[N];
 void init() { E.clear(); }
 void add_edge(int u, int v, ll w) {
    if (u != v) E.emplace_back(edge{u, v, w});
  ll build(int root, int n) {
    ll ans = 0;
    for (;;) {
      fill_n(in, n, INF);
      for (int i = 0; i < E.size(); ++i)</pre>
        if (E[i].u != E[i].v && E[i].w < in[E[i].v])</pre>
         pe[E[i].v] = i, in[E[i].v] = E[i].w;
      for (int u = 0; u < n; ++u) // no solution</pre>
        if (u != root && in[u] == INF) return -INF;
      int cntnode = 0;
      fill_n(id, n, -1), fill_n(vis, n, -1);
      for (int u = 0; u < n; ++u) {
        if (u != root) ans += in[u];
        int v = u;
```

```
while (vis[v] != u && !~id[v] && v != root)
         vis[v] = u, v = E[pe[v]].u;
if (v != root && !~id[v]) {
            for (int x = E[pe[v]].u; x != v;
                  x = E[pe[x]].u)
              id[x] = cntnode;
            id[v] = cntnode++;
         }
       if (!cntnode) break; // no cycle
       for (int u = 0; u < n; ++u)
  if (!~id[u]) id[u] = cntnode++;</pre>
       for (int i = 0; i < E.size(); ++i) {</pre>
         int v = E[i].v;
         E[i].u = id[E[i].u], E[i].v = id[E[i].v];
         if (E[i].u != E[i].v) E[i].w -= in[v];
       n = cntnode, root = id[root];
     return ans;
  }
};
```

2.7 DMST

```
#define rep(i, a, b) for (int i = a; i < (b); ++i)</pre>
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef vector<int> vi;
struct RollbackUF {
 vi e:
  vector<pii> st;
  RollbackUF(int n) : e(n, -1) {}
  int size(int x) { return -e[find(x)]; }
  int find(int x) { return e[x] < 0 ? x : find(e[x]); }</pre>
  int time() { return sz(st); }
  void rollback(int t) {
    for (int i = time(); i-- > t;)
      e[st[i].first] = st[i].second;
    st.resize(t);
 bool join(int a, int b) {
    a = find(a), b = find(b);
    if (a == b) return false;
    if (e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]});
    st.push_back({b, e[b]});
    e[a] += e[b];
   e[b] = a;
    return true:
 }
struct Edge {
 int a, b;
 ll w:
struct Node { /// lazy skew heap node
 Edge key;
  Node *l, *r;
  ll delta;
  void prop() {
    key.w += delta;
    if (l) l->delta += delta;
    if (r) r->delta += delta;
    delta = 0;
 Edge top() {
   prop();
    return key;
 }
Node *merge(Node *a, Node *b) {
 if (!a || !b) return a ?: b;
  a->prop(), b->prop();
 if (a->key.w > b->key.w) swap(a, b);
 swap(a->l, (a->r = merge(b, a->r)));
  return a;
void pop(Node *&a) {
 a->prop();
 a = merge(a->l, a->r);
pair<ll, vi> dmst(int n, int r, vector<Edge> &g) {
 RollbackUF uf(n);
 vector < Node *> heap(n);
```

```
for (Edge e : g)
    heap[e.b] = merge(heap[e.b], new Node{e});
  ll res = 0;
  vi seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector <Edge > Q(n), in(n, \{-1, -1\}), comp;
  deque<tuple<int, int, vector<Edge>>> cycs;
  rep(s, 0, n) {
    int u = s, qi = 0, w;
    while (seen[u] < 0) {</pre>
      if (!heap[u]) return {-1, {}};
      Edge e = heap[u]->top();
      heap[u]->delta -= e.w, pop(heap[u]);
Q[qi] = e, path[qi++] = u, seen[u] = s;
      res += e.w, u = uf.find(e.a);
      if (seen[u] == s) { /// found cycle, contract
        Node *cyc = 0;
        int end = qi, time = uf.time();
        do cyc = merge(cyc, heap[w = path[--qi]]);
        while (uf.join(u, w));
        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
        cycs.push_front({u, time, {&Q[qi], &Q[end]}});
      }
    rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
  for (auto &[u, t, comp] :
    cycs) { // restore sol (optional)
    uf.rollback(t);
    Edge inEdge = in[u];
for (auto &e : comp) in[uf.find(e.b)] = e;
    in[uf.find(inEdge.b)] = inEdge;
  rep(i, 0, n) par[i] = in[i].a;
  return {res, par};
2.8 VizingTheorem
```

```
namespace Vizing { // Edge coloring
                    // G: coloring adjM
int C[N][N], G[N][N];
void clear(int n) {
  for (int i = 0; i <= n; i++) {
    for (int j = 0; j <= n; j++) C[i][j] = G[i][j] = 0;</pre>
  }
void solve(vector<pii> &E, int n, int m) {
  int X[n] = {}, a;
  auto update = [&](int u) {
    for (X[u] = 1; C[u][X[u]]; X[u]++);
  auto color = [&](int u, int v, int c) {
    int p = G[u][v];
    G[u][v] = G[v][u] = c;
    C[u][c] = v;
    C[v][c] = u;
    C[u][p] = C[v][p] = 0;
    if (p) X[u] = X[v] = p;
    else update(u), update(v);
    return p;
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  for (int i = 1; i <= n; i++) X[i] = 1;</pre>
  for (int t = 0; t < E.size(); t++) {</pre>
    int u = E[t].first, v0 = E[t].second, v = v0,
        c0 = X[u], c = c0, d;
    vector<pii> L;
    int vst[n] = {};
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
      if (!C[v][c])
        for (a = (int)L.size() - 1; a >= 0; a--)
          c = color(u, L[a].first, c);
      else if (!C[u][d])
        for (a = (int)L.size() - 1; a >= 0; a--)
          color(u, L[a].first, L[a].second);
      else if (vst[d]) break;
```

2.9 MinimumCliqueCover

```
struct Clique_Cover { // 0-base, 0(n2^n)
  int co[1 << N], n, E[N];</pre>
  int dp[1 << N];</pre>
  void init(int _n) {
    n = _n, fill_n(dp, 1 << n, 0);
</pre>
     fill_n(E, n, 0), fill_n(co, 1 << n, 0);
  void add_edge(int u, int v) {
    E[u] = 1 << v, E[v] = 1 << u;
  int solve() {
     for (int i = 0; i < n; ++i)</pre>
       co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;
     dp[0] = (n \& 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {</pre>
       int t = i & -i;
       dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)</pre>
       co[i] = (co[i] & i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic</pre>
       for (int i = 0; i < (1 << n); ++i)</pre>
         sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n:
  }
};
```

2.10 CountMaximalClique

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
   g[u][v] = g[v][u] = 1;
 void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
      int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
      all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)
  if (g[v][some[d][j]])</pre>
           some[d + 1][tsn++] = some[d][j];
      for (int j = 0; j < nn; ++j)</pre>
        if (g[v][none[d][j]])
           none[d + 1][tnn++] = none[d][j];
      dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
   }
  int solve() {
    iota(some[0], some[0] + n, 1);
```

```
S = 0, dfs(0, 0, n, 0);
return S;
}
};
```

2.11 Theorems

|Maximum independent edge set| = |V| - |Minimum edge cover| |Maximum independent set| = |V| - |Minimum vertex cover|

3 Flow-Matching

3.1 KM

```
ll w[N][N], hl[N], hr[N], slk[N];
   int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
   bool vl[N], vr[N];
   void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i)</pre>
       fill_n(w[i], n, -INF);
   void add_edge(int a, int b, ll wei) {
     w[a][b] = wei;
   bool Check(int x) {
     if (vl[x] = 1, \sim fl[x])
       return vr[qu[qr++] = fl[x]] = 1;
     while (\sim x) swap(x, fr[fl[x] = pre[x]]);
     return 0;
   void bfs(int s) {
     fill_n(slk
         , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
     ql = qr = 0, qu[qr++] = s, vr[s] = 1;
     for (ll d;;) {
      while (ql < qr)
for (int x = 0, y = qu[ql++]; x < n; ++x)</pre>
           if (!vl[x] && slk
               [x] >= (d = hl[x] + hr[y] - w[x][y])) {
              if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
         }
       d = INF;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && d > slk[x]) d = slk[x];
       for (int x = 0; x < n; ++x) {
         if (vl[x]) hl[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !Check(x)) return;
    }
   il solve() {
     fill_n(fl
           n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
     for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(w[i], w[i] + n);
     for (int i = 0; i < n; ++i) bfs(i);</pre>
     ll res = 0;
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
     return res:
  }
};
```

3.2 MCMF

```
struct Edge {
    ll from, to, cap, flow, cost, rev;
} *past[N];
vector<Edge> G[N];
int inq[N], n, s, t;
ll dis[N], up[N], pot[N];
bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
        if (cap > 0 && dis[u] > d) {
            dis[u] = d, up[u] = cap, past[u] = e;
            if (!inq[u]) inq[u] = 1, q.push(u);
        }
};
```

```
relax(s, 0, INF, 0);
     while (!q.empty()) {
       int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u]) {
         ll d2 = dis[u] + e.cost + pot[u] - pot[e.to];
         relax
              (e.to, d2, min(up[u], e.cap - e.flow), &e);
     return dis[t] != INF;
   void solve(int
     , int _t, ll &flow, ll &cost, bool neg = true) { s = _s, t = _t, flow = 0. cost = \theta:
     if (neg) BellmanFord(), copy_n(dis, n, pot);
     for (; BellmanFord(); copy_n(dis, n, pot)) {
       for (int
            i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
       flow += up[t], cost += up[t] * dis[t];
       for (int i = t; past[i]; i = past[i]->from) {
         auto &e = *past[i];
         e.flow += up[t], G[e.to][e.rev].flow -= up[t];
    }
  }
   void init(int _n) {
     n = _n, fill_n(pot, n, 0);
     for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
   G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
     G[b].pb(Edge\{b, a, 0, 0, -cost, SZ(G[a]) - 1\});
};
```

3.3 GeneralGraphMatching

```
queue < int > q; int n;
vector<int> fa, s, vis, pre, match;
vector < vector < int >> G;
int Find(int u)
{ return u == fa[u] ? u : fa[u] = Find(fa[u]); }
int LCA(int x, int y) {
  static int tk = 0; tk++; x = Find(x); y = Find(y);
for (;; swap(x, y)) if (x != n) {
     if (vis[x] == tk) return x;
     vis[x] = tk;
    x = Find(pre[match[x]]);
  }
void Blossom(int x, int y, int l) {
   for (; Find(x) != l; x = pre[y]) {
    pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
     for (int z: \{x, y\}) if (fa[z] == z) fa[z] = l;
  }
bool Bfs(int r) {
  iota(ALL(fa), 0); fill(ALL(s), -1);
  q = queue < int >(); q.push(r); s[r] = 0;
  for (; !q.empty(); q.pop()) {
     for (int x = q.front(); int u : G[x])
       if (s[u] == -1) {
         if (pre[u] = x, s[u] = 1, match[u] == n) {
  for (int a = u, b = x, last;
                 b != n; a = last, b = pre[a])
               last =
                   match[b], match[b] = a, match[a] = b;
            return true;
          q.push(match[u]); s[match[u]] = 0;
       } else if (!s[u] && Find(u) != Find(x)) {
          int l = LCA(u, x);
         Blossom(x, u, l); Blossom(u, x, l);
  return false;
Matching( \begin{array}{c} \textbf{int} & \underline{\phantom{a}} \\ \textbf{n} \end{array}) \; : \; n(\underline{\phantom{a}} \\ \textbf{n}), \; fa(n + 1), \; s(n + 1), \; vis
     (n + 1), pre(n + 1, n), match(n + 1, n), G(n) {}
void add_edge(int u, int v)
{ G[u].pb(v), G[v].pb(u); }
int solve() {
  int ans = 0;
  for (int x = 0; x < n; ++x)
```

```
if (match[x] == n) ans += Bfs(x);
    return ans:
  } // match[x] == n means not matched
};
```

3.4 MaxWeightMaching

```
struct WeightGraph { // 1-based
  struct edge { int u, v, w; }; int n, nx;
  vector<int> lab; vector<vector<edge>> g;
  vector <int> slack, match, st, pa, S, vis;
vector <vector <int>> flo, flo_from; queue <int> q;
WeightGraph(int n_): n(n_), nx(n * 2), lab(nx + 1),
    g(nx + 1, vector < edge > (nx + 1)), slack(nx + 1)
    flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
match = st = pa = S = vis = slack;
    rep(u, 1, n) rep(v, 1, n) g[u][v] = {u, v, 0};
  int ED(edge e)
  { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
  void update_slack(int u, int x, int &s)
  { if (!s || ED(g[u][x]) < ED(g[s][x])) s = u; }
  void set_slack(int x) {
     slack[x] = 0;
     for (int u = 1; u <= n; ++u)</pre>
      if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
         update_slack(u, x, slack[x]);
  void q_push(int x) {
    if (x <= n) q.push(x);</pre>
    else for (int y : flo[x]) q_push(y);
  void set_st(int x, int b) {
    st[x] = b;
    if (x > n) for (int y : flo[x]) set_st(y, b);
  vector<int> split_flo(auto &f, int xr) {
     auto it = find(ALL(f), xr);
     if (auto pr = it - f.begin(); pr % 2 == 1)
      reverse(1 + ALL(f)), it = f.end() - pr;
     auto res = vector(f.begin(), it);
     return f.erase(f.begin(), it), res;
  void set_match(int u, int v) {
    match[u] = g[u][v].v;
    if (u <= n) return;</pre>
    int xr = flo_from[u][g[u][v].u];
     auto &f = flo[u], z = split_flo(f, xr);
    rep(i, 0, SZ(z)-1) set_match(z[i], z[i ^ 1]);
    set_match(xr, v); f.insert(f.end(), ALL(z));
  void augment(int u, int v) {
    for (;;) {
      int xnv = st[match[u]]; set_match(u, v);
      if (!xnv) return:
      set_match(xnv, st[pa[xnv]]);
      u = st[pa[xnv]], v = xnv;
  int lca(int u, int v) {
  static int t = 0; ++t;
     for (++t; u || v; swap(u, v)) if (u) {
      if (vis[u] == t) return u;
       vis[u] = t; u = st[match[u]];
      if (u) u = st[pa[u]];
    return 0;
  void add_blossom(int u, int o, int v) {
    int b = find(n + 1 + ALL(st), 0) - begin(st);
    lab[b] = 0, S[b] = 0; match[b] = match[o];
     vector<int> f = {o};
     for (int x = u, y; x != o; x = st[pa[y]])
      f.pb(x), f.pb(y = st[match[x]]), q_push(y);
     reverse(1 + ALL(f));
     for (int x = v, y; x != o; x = st[pa[y]])
      f.pb(x), f.pb(y = st[match[x]]), q_push(y);
     flo[b] = f; set_st(b, b);
     for (int x = 1; x <= nx; ++x)</pre>
      g[b][x].w = g[x][b].w = 0;
     fill(ALL(flo_from[b]), 0);
     for (int xs : flo[b]) {
       for (int x = 1; x <= nx; ++x)</pre>
         if (g[
             b][x].w == 0 || ED(g[xs][x]) < ED(g[b][x]))
           g[b][x] = g[xs][x], g[x][b] = g[x][xs];
```

```
for (int x = 1; x <= n; ++x)</pre>
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) { xs = x; continue; }
    pa[xs] = g[x][xs].u; S[xs] = 1, S[x] = 0;
    slack[xs] = 0; set_slack(x); q_push(x); xs = -1;
  for (int x : flo[b])
    if (x == xr) S[x] = 1, pa[x] = pa[b];
    else S[x] = -1, set_slack(x);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
    int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
    slack[v] = slack[nu] = 0; S[nu] = 0; q_push(nu);
  } else if (S[v] == 0) {
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
  return false;
bool matching() {
  fill(ALL(S), -1), fill(ALL(slack), 0);
  q = queue<int>();
  for (int x = 1; x <= nx; ++x)</pre>
    if (st[x] == x && !match[x])
      pa[x] = 0, S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
    while (SZ(q)) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)</pre>
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (ED(g[u][v]) != 0)
            update_slack(u, st[v], slack[st[v]]);
          else if
                (on_found_edge(g[u][v])) return true;
        }
    int d = INF;
    for (int b = n + 1; b <= nx; ++b)</pre>
      if (st[b] == b && S[b] == 1)
        d = min(d, lab[b] / 2);
    for (int x = 1; x <= nx; ++x)</pre>
      if (int
          s = slack[x]; st[x] == x && s && s[x] <= 0)
        d = min(d, ED(g[s][x]) / (S[x] + 2));
    for (int u = 1; u <= n; ++u)
      if (S[st[u]] == 1) lab[u] += d;
      else if (S[st[u]] == 0) {
        if (lab[u] <= d) return false;</pre>
        lab[u] -= d;
    rep(b, n + 1, nx) if (st[b] == b \&\& S[b] >= 0)
      lab[b] += d * (2 - 4 * S[b]);
    for (int x = 1; x <= nx; ++x)
  if (int s = slack[x]; st[x] == x &&</pre>
          s \&\& st[s] != x \&\& ED(g[s][x]) == 0)
        if (on_found_edge(g[s][x])) return true;
    for (int b = n + 1; b <= nx; ++b)</pre>
      if (st[b] == b && S[b] == 1 && lab[b] == 0)
        expand_blossom(b);
  return false;
pair<ll, int> solve() {
  fill(ALL(match), 0);
  rep(u, 0, n) st[u] = u, flo[u].clear();
  int w_max = 0;
  rep(u, 1, n) rep(v, 1, n) {
    flo_from[u][v] = (u == v ? u : 0);
    w_{max} = max(w_{max}, g[u][v].w);
  fill(ALL(lab), w_max);
  int n_matches = 0; ll tot_weight = 0;
  while (matching()) ++n_matches;
  rep(u, 1, n) if (match[u] && match[u] < u)</pre>
    tot_weight += g[u][match[u]].w;
```

```
return make_pair(tot_weight, n_matches);
}
void add_edge(int u, int v, int w)
{ g[u][v].w = g[v][u].w = w; }
};
```

3.5 GlobalMinCut

```
#define REP for (int i = 0; i < n; ++i)
static const int MXN = 514, INF = 2147483647;</pre>
   int vst[MXN], edge[MXN][MXN], wei[MXN];
   void init(int n) {
    REP fill_n(edge[i], n, 0);
   void addEdge(int u, int v, int w){
     edge[u][v] += w; edge[v][u] += w;
   int search(int &s, int &t, int n){
     fill_n(vst, n, 0), fill_n(wei, n, 0);
     s = t = -1;
     int mx, cur;
     for (int j = 0; j < n; ++j) {</pre>
       mx = -1, cur = 0;
       REP if (wei[i] > mx) cur = i, mx = wei[i];
       vst[cur] = 1, wei[cur] = -1;
       s = t; t = cur;
REP if (!vst[i]) wei[i] += edge[cur][i];
     return mx;
   int solve(int n) {
     int res = INF;
     for (int x, y; n > 1; n--){
       res = min(res, search(x, y, n));
       REP edge[i][x] = (edge[x][i] += edge[y][i]);
         edge[y][i] = edge[n - 1][i];
         edge[i][y] = edge[i][n - 1];
       return res;
  }
} sw;
```

3.6 BoundedFlow(Dinic)

```
struct edge {
  int to, cap, flow, rev;
vector<edge> G[N];
int n, s, t, dis[N], cur[N], cnt[N];
void init(int _n) {
  n = _n;
for (int i = 0; i < n + 2; ++i)</pre>
    G[i].clear(), cnt[i] = 0;
void add_edge(int u, int v, int lcap, int rcap) {
  cnt[u] -= lcap, cnt[v] += lcap;
G[u].pb(edge{v, rcap, lcap, SZ(G[v])});
G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
void add_edge(int u, int v, int cap) {
  G[u].pb(edge{v, cap, 0, SZ(G[v])});
G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
int dfs(int u, int cap) {
  if (u == t || !cap) return cap;
  for (int &i = cur[u]; i < SZ(G[u]); ++i) {</pre>
    edge &e = G[u][i];
    if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
       int df = dfs(e.to, min(e.cap - e.flow, cap));
       if (df) {
         e.flow += df, G[e.to][e.rev].flow -= df;
         return df;
      }
    }
  dis[u] = -1;
  return 0;
bool bfs() {
  fill_n(dis, n + 3, -1);
  queue < int > q;
  q.push(s), dis[s] = 0;
  while (!q.empty()) {
    int u = q.front();
```

```
q.pop();
       for (edge &e : G[u])
         if (!~dis[e.to] && e.flow != e.cap)
           q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
       while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
    int sum = 0;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
       add_edge(n + 1, i, cnt[i]), sum += cnt[i];
else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int i = 0; i < n; ++i)</pre>
       if (cnt[i] > 0)
       G[n + 1].pop_back(), G[i].pop_back();
else if (cnt[i] < 0)</pre>
         G[i].pop_back(), G[n + 2].pop_back();
     return sum != -1;
  int solve(int _s, int
                            _t) {
    add_edge(_t, _s, INF);
     if (!solve()) return -1; // invalid flow
     int x = G[_t].back().flow;
    return G[_t].pop_back(), G[_s].pop_back(), x;
  }
};
```

3.7 GomoryHuTree

```
int g[MAXN];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
  for (int i = 1; i < n; ++i) {
    Dinic.reset();
    add_edge(i, g[i], Dinic.maxflow(i, g[i]));
    for (int j = i + 1; j <= n; ++j)
        if (g[j] == g[i] && ~Dinic.dis[j])
        g[j] = i;
  }
}</pre>
```

3.8 MinCostCirculation

```
struct Edge {
  ll from, to, cap, fcap, flow, cost, rev;
} *past[N];
vector < Edge > G[N];
ll dis[N], inq[N], n;
void BellmanFord(int s) {
  fill_n(dis, n, INF), fill_n(inq, n, 0);
  queue<int> q;
  auto relax = [&](int u, ll d, Edge *e) {
    if (dis[u] > d) {
      dis[u] = d, past[u] = e;
if (!inq[u]) inq[u] = 1, q.push(u);
    }
  };
  relax(s, 0, 0);
  while (!q.empty()) {
  int u = q.front();
    q.pop(), inq[u] = 0;
    for (auto &e : G[u])
      if (e.cap > e.flow)
         relax(e.to, dis[u] + e.cost, &e);
void try_edge(Edge &cur) {
  if (cur.cap > cur.flow) return ++cur.cap, void();
  BellmanFord(cur.to);
  if (dis[cur.from] + cur.cost < 0) {</pre>
    ++cur.flow, --G[cur.to][cur.rev].flow;
    for (int
          i = cur.from; past[i]; i = past[i]->from) {
      auto &e = *past[i];
++e.flow, --G[e.to][e.rev].flow;
```

```
}
      ++cur.cap;
   }
   void solve(int mxlg) {
     for (int b = mxlg; b >= 0; --b) {
  for (int i = 0; i < n; ++i)</pre>
           for (auto &e : G[i])
             e.cap *= 2, e.flow *= 2;
        for (int i = 0; i < n; ++i)</pre>
           for (auto &e : G[i])
             if (e.fcap >> b & 1)
                try_edge(e);
   void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
   void add_edge(ll a, ll b, ll cap, ll cost) {
     G[a].pb(Edge
     {a, b, 0, cap, 0, cost, SZ(G[b]) + (a == b)});
G[b].pb(Edge{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1});
} mcmf; // O(VE * ElogC)
```

3.9 FlowModelsBuilding

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink \dot{T} .
 - 2. For each edge (x,y,l,u), connect $x \rightarrow y$ with capacity u-l.
 - 3. For each vertex v , denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
 - 1. Redirect every edge: $y \rightarrow x$ if $(x,y) \in M$, $x \rightarrow y$ otherwise.
 - 2. DFS from unmatched vertices in \widetilde{X} .
 - 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
 - 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0 , connect $S\to v$ with $(cost, cap)\,{=}\,(0, d(v))$
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow $C\!+\!K$
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \! \to \! v$, $v \! \in \! G$ with capacity K
 - 4. For each edge (u,v,w) in G , connect $u\to v$ and $v\to u$ with capacity w
 - 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K + 2T (\sum_{e \in E(v)} w(e)) 2w(v)$
 - 6. T is a valid answer if the maximum flow $f\!<\!K|V|$
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v' , and connect $u' \to v'$ with weight w(u,v) .
 - 2. Connect $v\to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v .
 - 3. Find the minimum weight perfect matching on G^{\prime} .
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .

```
2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers. \min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \geq -c_{uv} \Leftrightarrow \min \sum_{u} b_u p_u + \sum_{uv} c_{uv} \max(0, p_v - p_u - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_u \qquad \qquad p_u \geq 0
```

4 Data Struture

4.1 LCT

```
static Splay nil;
  Splay *ch[2], *f;
  int val, sum, rev, size;
  Splay (int
      _val = 0) : val(_val), sum(_val), rev(0), size(1)
  {f = ch[0] = ch[1] = &nil; }
  bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
  int dir()
  { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void give_tag(int r) {
    if (r) swap(ch[0], ch[1]), rev ^= 1;
  void push() {
    if (ch[0] != &nil) ch[0]->give_tag(rev);
    if (ch[1] != &nil) ch[1]->give_tag(rev);
  void pull() {
    // take care of the nil!
    size = ch[0] -> size + ch[1] -> size + 1;
    sum = ch[0] -> sum ^ ch[1] -> sum ^ val;
    if (ch[0] != &nil) ch[0]->f = this;
if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x - > f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x - > f = p - > f;
 p->setCh(x->ch[!d], d);
 x->setCh(p, !d);
  p->pull(), x->pull();
void splay(Splay *x) {
  vector < Splay *> splay Vec;
  for (Splay *q = x;; q = q->f) {
    splayVec.pb(q);
    if (q->isr()) break;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir())
    rotate(x->f), rotate(x);
else rotate(x), rotate(x);
 }
Splay* access(Splay *x) {
 Splay *q = nil;
for (; x != nil; x = x->f)
   splay(x), x -> setCh(q, 1), q = x;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x){
  root_path(x), x->give_tag(1);
 x->push(), x->pull();
void split(Splay *x, Splay *y) {
  chroot(x), root_path(y);
void link(Splay *x, Splay *y) {
 root_path(x), chroot(y);
 x->setCh(y, 1);
```

```
void cut(Splay *x, Splay *y) {
  split(x, y);
  if (y->size != 5) return;
  y->push();
 y - > ch[0] = y - > ch[0] - > f = nil;
Splay* get_root(Splay *x) {
  for (root_path(x); x->ch[0] != nil; x = x->ch[0])
    x->push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
  return get_root(x) == get_root(y);
Splay* lca(Splay *x, Splay *y) {
 access(x), root_path(y);
  if (y->f == nil) return y;
  return v->f;
void change(Splay *x, int val) {
 splay(x), x->val = val, x->pull();
int query(Splay *x, Splay *y) {
 split(x, y);
  return y->sum;
```

5 String

5.1 KMP

```
int n = t.size(), ans = 0;
vector <int > f(t.size(), 0);
f[0] = -1;
for (int i = 1, j = -1; i < t.size(); i++) {
    while (j >= 0)
        if (t[j + 1] == t[i]) break;
        else j = f[j];
    f[i] = ++j;
}
for (int i = 0, j = 0; i < s.size(); i++) {
    while (j >= 0)
        if (t[j + 1] == s[i]) break;
        else j = f[j];
    if (++j + 1 == t.size()) ans++, j = f[j];
}
return ans;
}
```

5.2 Z

```
void z(string s) {
  for (int i = 1, mx = 0; i < s.size(); i++) {
    if (i < Z[mx] + mx)
        Z[i] = min(Z[mx] - i + mx, Z[i - mx]);
    while (
        Z[i] + i < s.size() && s[i + Z[i]] == s[Z[i]])
        Z[i]++;
    if (Z[i] + i > Z[mx] + mx) mx = i;
  }
}
```

5.3 Manacher

```
int manacher(string s) {
  string t;
   for (int i = 0; i < s.size(); i++) {</pre>
     if (i) t.push_back('$');
     t.push_back(s[i]);
   int mx = 0, ans = 0;
  for (int i = 0; i < t.size(); i++) {</pre>
    man[i] = 1;
     man[i] = min(man[mx] + mx - i, man[2 * mx - i]);
     while (man[i] + i < t.size() && i - man[i] >= 0 &&
       t[i + man[i]] == t[i - man[i]])
       man[i]++;
    if (i + man[i] > mx + man[mx]) mx = i;
  for (int i = 0; i < t.size(); i++)</pre>
    ans = max(ans, man[i] - 1);
   return ans;
}
```

5.4 SuffixArray

```
void SA(string s) {
  int n = s.size();
  sa.resize(n), cnt.resize(n), rk.resize(n),
    tmp.resize(n);
  iota(sa.begin(), sa.end(), 0);
  sort(sa.begin(), sa.end(),
     [&](int i, int j) { return s[i] < s[j]; });</pre>
  rk[0] = 0;
  for (int i = 1; i < n; i++)</pre>
    rk[sa[i]] =
      rk[sa[i - 1]] + (s[sa[i - 1]] != s[sa[i]]);
  for (int k = 1; k <= n; k <<= 1) {</pre>
    fill(cnt.begin(), cnt.end(), 0);
    for (int i = 0; i < n; i++)
  cnt[rk[(sa[i] - k + n) % n]]++;</pre>
    for (int i = 1; i < n; i++) cnt[i] += cnt[i - 1];
for (int i = n - 1; i >= 0; i--)
       tmp[--cnt[rk[(sa[i] - k + n) % n]]] =
         (sa[i] - k + n) \% n;
    sa.swap(tmp);
     tmp[sa[0]] = 0;
    for (int i = 1; i < n; i++)</pre>
      tmp[sa[i]] = tmp[sa[i - 1]] +
         (rk[sa[i - 1]] != rk[sa[i]] ||
           rk[(sa[i - 1] + k) % n] !=
             rk[(sa[i] + k) % n]);
    rk.swap(tmp);
  }
}
void LCP(string s) {
  int n = s.size(), k = 0;
  lcp.resize(n);
  for (int i = 0; i < n; i++)</pre>
    if (rk[i] == 0) lcp[rk[i]] = 0;
     else {
       if (k) k--;
      int j = sa[rk[i] - 1];
       while (
         i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
       lcp[rk[i]] = k;
    }
}
```

5.5 SAIS

```
bool _t[N * 2];
int SA[N * 2], H[N], RA[N];
int _s[N * 2], _c[N * 2], x[N], _p[N],
// zero based, string content MUST > 0
                                   _p[N], _q[N * 2];
// SA[i]: SA[i]-th
    suffix is the i-th lexigraphically smallest suffix.
// H[i]: longest
    common prefix of suffix SA[i] and suffix SA[i - 1].
void pre(int *sa, int *c, int n, int z)
{ fill_n(sa, n, 0), copy_n(c, z, x); }
void induce
    (int *sa, int *c, int *s, bool *t, int n, int z) {
  copy_n(c, z - 1, x + 1);
  for (int i = 0; i < n; ++i)</pre>
    if (sa[i] && !t[sa[i] - 1])
      sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  copy_n(c, z, x);
  for (int i = n - 1; i >= 0; --i)
    if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa
     int *p, int *q, bool *t, int *c, int n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0,
      nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  fill_n(c, z, 0);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;</pre>
  partial_sum(c, c + z, c);
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;</pre>
    return;
      (int i = n - 2; i >= 0; --i)
    t[i] = (
        s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
```

```
for (int i = 1; i <= n - 1; ++i)
  if (t[i] && !t[i - 1])</pre>
        sa[--x[s[i]]] = p[q[i] = nn++] = i;
   induce(sa, c, s, t, n, z);
   for (int i = 0; i < n; ++i)</pre>
     if (sa[i] && t[sa[i]] && !t[sa[i] - 1]) {
  bool neq = last < 0 || !equal</pre>
             (s + sa[i], s + p[q[sa[i]] + 1], s + last);
        ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns,
         nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
   pre(sa, c, n, z);

for (int i = nn - 1; i >= 0; --i)
     sa[--x[s[p[nsa[i]]]]] = p[nsa[i]];
   induce(sa, c, s, t, n, z);
void mkhei(int n) {
   for (int i = 0, j = 0; i < n; ++i) {</pre>
     if (RA[i])
     for (; _s[i + j] == _s[SA[RA[i] - 1] + j]; ++j);
H[RA[i]] = j, j = max(0, j - 1);
   }
void build(int *s, int n) {
   copy_n(s, n, _s), _s[n] = 0;
   sais(_s, SA, _p, _q, _t, _c, n + 1, 256);
   copy_n(SA + 1, n, SA);
   for (int i = 0; i < n; ++i) RA[SA[i]] = i;</pre>
   mkhei(n);
}}
```

5.6 ACAutomaton

```
#define sigma 26
#define base 'a
struct AhoCorasick {
  int ch[sumS][sigma] = {{}}, f[sumS] = {-1},
       tag[sumS], mv[sumS][sigma], jump[sumS],
       cnt[sumS];
  int idx = 0;
  int insert(string &s) {
    int j = 0;
    for (int i = 0; i < (int)s.size(); i++) {</pre>
      if (!ch[j][s[i] - base])
   ch[j][s[i] - base] = ++idx;
      j = ch[j][s[i] - base];
    tag[j] = 1;
    return i:
  int next(int u, int c) {
    return u < 0 ? 0 : mv[u][c];</pre>
  void build() {
    queue<int> q;
    q.push(0);
    while (!q.empty()) {
      int u = q.front();
      a.pop():
       for (int v = 0; v < sigma; v++) {</pre>
         if (ch[u][v]) {
           f[ch[u][v]] = next(f[u], v);
           q.push(ch[u][v]);
         }
         mv[u][v] =
           (ch[u][v] ? ch[u][v] : next(f[u], v));
       if (u) jump[u] = (tag[f[u]] ? f[u] : <math>jump[f[u]]);
    }
  void match(string &s) {
    for (int i = 0; i <= idx; i++) cnt[i] = 0;</pre>
    for (int i = 0, j = 0; i < (int)s.size(); i++) {</pre>
      j = next(j, s[i] - base);
      cnt[j]++;
    vector<int> v:
    v.emplace_back(0);
    for (int i = 0; i < (int)v.size(); i++)
  for (int j = 0; j < sigma; j++)</pre>
         if (ch[v[i]][j]) v.emplace_back(ch[v[i]][j]);
    reverse(v.begin(), v.end());
    for (int i : v)
      if (f[i] > 0) cnt[f[i]] += cnt[i];
```

```
National Taiwan University RngBased (forked from 8BQube)
|} ac:
 5.7 MinRotation
  int n = s.size();
   s = s + s;
   int i = 0, ans = 0;
   while (i < n) {</pre>
     ans = i;
     int j = i + 1, k = i;
     while (j < s.size() && s[j] >= s[k]) {
       k = (s[j] > s[k] ? i : k + 1);
     while (i <= k) i += j - k;</pre>
   }
   return ans;
 5.8 ExtSAM
   int len[N * 2], link[N * 2]; // maxlength, suflink
int next[N * 2][CNUM], tot; // [0, tot), root = 0
   int lenSorted[N * 2]; // topo. order
   int cnt[N * 2]; // occurence
   int newnode() {
     fill n(next[tot], CNUM, 0);
     len[tot] = cnt[tot] = link[tot] = 0;
     return tot++;
   void init() { tot = 0, newnode(), link[0] = -1; }
   int insertSAM(int last, int c) {
     int cur = next[last][c];
     len[cur] = len[last] + 1;
     int p = link[last];
     while (p != -1 && !next[p][c])
     next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;
     int q = next[p][c];
     if (len
         [p] + 1 == len[q]) return link[cur] = q, cur;
     int clone = newnode();
     for (int i = 0; i < CNUM; ++i)</pre>
       next[
            clone][i] = len[next[q][i]] ? next[q][i] : 0; |};
     len[clone] = len[p] + 1;
     while (p != -1 && next[p][c] == q)
       next[p][c] = clone, p = link[p];
     link[link[cur] = clone] = link[q];
     link[q] = clone;
     return cur;
   void insert(const string &s) {
     int cur = 0;
     for (auto ch : s) {
       int &nxt = next[cur][int(ch - 'a')];
       if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
   void build() {
     queue < int > q;
     q.push(0);
     while (!q.empty()) {
       int cur = q.front();
       q.pop();
       for (int i = 0; i < CNUM; ++i)</pre>
         if (next[cur][i])
           q.push(insertSAM(cur, i));
     vector<int> lc(tot);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
     partial_sum(ALL(lc), lc.begin());
     for (int i
         = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
   void solve() {
     for (int i = tot - 2; i >= 0; --i)
       cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
  }
}:
```

5.9 PalindromeTree

```
struct node {
  int next[26], fail, len;
```

```
int cnt, num; // cnt: appear times, num: number of
                 // pal. suf.
  node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
   for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
};
vector < node > St;
vector<char> s;
int last, n;
palindromic_tree() : St(2), last(1), n(0) {
  St[0].fail = 1, St[1].len = -1, s.pb(-1);
inline void clear() {
  St.clear(), s.clear(), last = 1, n = 0;
  St.pb(0), St.pb(-1);
  St[0].fail = 1, s.pb(-1);
inline int get_fail(int x) {
  while (s[n - St[x].len - 1] != s[n])
   x = St[x].fail;
  return x;
inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
  int cur = get_fail(last);
  if (!St[cur].next[c]) {
    int now = SZ(St);
    St.pb(St[cur].len + 2);
    St[now].fail
      St[get_fail(St[cur].fail)].next[c];
    St[cur].next[c] = now;
    St[now].num = St[St[now].fail].num + 1;
  last = St[cur].next[c], ++St[last].cnt;
inline void count() { // counting cnt
  auto i = St.rbegin();
  for (; i != St.rend(); ++i) {
    St[i->fail].cnt += i->cnt;
inline int size() { // The number of diff. pal.
  return SZ(St) - 2;
```

Number Theory

6.1 Primes

6.2 ExtGCD

```
// beware of negative numbers!
void extgcd(ll a, ll b, ll c, ll &x, ll &y) {
  if (b == 0) x = c / a, y = 0;
    extgcd(b, a % b, c, y, x);
    y -= x * (a / b);
  }
| \} // |x| \ll b/2, |y| \ll a/2
```

6.3 FloorCeil

```
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

6.4 FloorSum

```
if (A == 0) return (N + 1) * (B / C);
  if (A > C || B > C)
    return (N + 1) * (B / C) +
      N * (N + 1) / 2 * (A / C) +
      floorsum(A % C, B % C, C, N);
  ll M = (A * N + B) / C;
  return N * M - floorsum(C, C - B - 1, A, M - 1);
} // \sum^{n}_0 floor((ai + b) / m)
```

6.5 MillerRabin

```
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383 6 : primes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool Miller_Rabin(ll a, ll n) {
```

```
if ((a = a % n) == 0) return 1;
if (n % 2 == 0) return n == 2;
ll tmp = (n - 1) / ((n - 1) & (1 - n));
ll t = __lg(((n - 1) & (1 - n))), x = 1;
for (; tmp; tmp >>= 1, a = mul(a, a, n))
if (tmp & 1) x = mul(x, a, n);
if (x == 1 || x == n - 1) return 1;
while (--t)
  if ((x = mul(x, x, n)) == n - 1) return 1;
```

6.6 PollardRho

```
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2
       == 0) return PollardRho(n / 2), ++cnt[2], void();
  ll x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return:
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
d = gcd(abs(x - y), n);
  }
}
```

6.7 Fraction

```
ll n, d;
  fraction
      (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
    ll t = gcd(n, d);
    n /= t, d /= t;
if (d < 0) n = -n, d = -d;
  fraction operator -() const
  { return fraction(-n, d); }
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator - (const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); }
  fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;
 }
};
```

```
6.8 Simplex
    Standard form: maximize \mathbf{c}^T \mathbf{x} subject to A \mathbf{x} \leq \mathbf{b} and \mathbf{x} \geq 0.
Dual LP: minimize {f b}^T{f y} subject to A^T{f y}\!\ge\!{f c} and {f y}\!\ge\!{f 0}. ar{{f x}} and ar{{f y}} are optimal if and only if for all i\in[1,n], either \bar{x}_i=0
or \sum_{j=1}^m A_{ji} ar{y}_j = c_i holds and for all i \in [1,m] either ar{y}_i = 0 or
\sum_{j=1}^{n} A_{ij} \bar{x}_j = b_j holds.
1. In case of minimization, let c_i'\!=\!-c_i
2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
3. \sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j
     • \sum_{1 \le i \le n} A_{ji} x_i \le b_j
     • \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j
4. If x_i has no lower bound, replace x_i with x_i - x_i'
// infeasible < 0, unbounded = inf, Ax <= b, max
struct simplex {
   const double inf = 1 / .0, eps = 1e-9;
    int n, m = 0;
   double A[205][205], B[205];
void init(int _n) { n = _n; }
void equation(vector<double> a, double b) {
```

for (int i = 0; i < n; i++) A[m][i] = a[i];</pre>

for (int i = 0; i < n; i++) A[0][i] = -c[i];</pre>

A[m][n + m] = 1, B[m] = b;double solve(vector<double> c) {

A[0][n] = 1;

```
int flag = 1;
     while (flag--)
       for (int i = 0; i <= n + m; i++)</pre>
          if (A[0][i] < -eps) {</pre>
            double bx = inf;
            int piv = -1;
            for (int j = 1; j <= m; j++)
  if (0 <= A[j][i] && B[j] / A[j][i] <= bx)</pre>
                 piv = j, bx = B[j] / A[j][i];
            if (piv == -1) continue;
            if (bx == inf) return inf;
            flag = 1;
             for (int j = 0; j <= m; j++)</pre>
               if (j != piv) {
                 for (int k = 0; k <= n + m; k++)</pre>
                   if (k != i)
                     A[j][k]
                        A[piv][k] * A[j][i] / A[piv][i];
                 B[j] -= B[piv] * A[j][i] / A[piv][i];
                 A[j][i] = 0;
     for (int i = 0; i <= m; i++)</pre>
       if (B[i] < -eps) return -inf;</pre>
     return B[0] / A[0][n];
} lp;
```

6.9 GuassianElimination

```
fraction M[MAXN][MAXN + 1], sol[MAXN];
  int solve() { //-1: inconsistent, >= 0: rank
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
      while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
      for (int j = 0; j < n; ++j) {</pre>
         if (i == j) continue;
         fraction tmp = -M[j][piv] / M[i][piv];
         for (int k = 0; k <=</pre>
              m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
      }
    int rank = 0;
    for (int i = 0; i < n; ++i) {</pre>
      int piv = 0;
      while (piv < m && !M[i][piv].n) ++piv;</pre>
      if (piv == m && M[i][m].n) return -1;
      else if (piv
            < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
    return rank;
  }
};
```

6.10 ChineseRemainder

```
ll g = gcd(m1, m2);
if ((x2 - x1) % g) return -1; // no sol
m1 /= g; m2 /= g;
pll p = exgcd(m1, m2);
ll lcm = m1 * m2 * g;
ll res = p.first * (x2 - x1) * m1 + x1;
// be careful with overflow
return (res % lcm + lcm) % lcm;
```

6.11 Factorial $\mathsf{Mod} p^k$

```
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)</pre>
     if (i % p) prod[i] = prod[i - 1] * i % pk;
     else prod[i] = prod[i - 1];
  ll rt = 1;
  for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
     rt = rt * prod[n % pk] % pk;
  return rt;
} // (n! without factor p) % p^k
```

6.12 QuadraticResidue

```
ll trial(ll y, ll z, ll m) {
  ll a0 = 1, a1 = 0, b0 = z, b1 = 1, p = (m - 1) / 2;
  while (p) {
    if (p & 1)
      tie(a0, a1) =
        make_pair((a1 * b1 % m * y + a0 * b0) % m,
          (a0 * b1 + a1 * b0) % m);
    tie(b0, b1) =
      make_pair((b1 * b1 % m * y + b0 * b0) % m,
        (2 * b0 * b1) % m);
  if (a1) return inv(a1, m);
  return -1;
mt19937 rd(49);
ll psqrt(ll y, ll p) {
  if (fpow(y, (p - 1) / 2, p) != 1) return -1;
for (int i = 0; i < 30; i++) {</pre>
    ll z = rd() \% p;
    if (z * z % p == y) return z;
    \hat{x} = trial(y, z, p);
    if (x == -1) continue;
    return x;
  return -1:
```

6.13 MeisselLehmer

```
if (n <= 1) return 0;</pre>
int v = sqrt(n), s = (v + 1) / 2, pc = 0;
vector<int> smalls(v + 1), skip(v + 1), roughs(s);
vector<ll> larges(s);
for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
for (int i = 0; i < s; ++i) {</pre>
  roughs[i] = 2 * i + 1;
  larges[i] = (n / (2 * i + 1) + 1) / 2;
for (int p = 3; p <= v; ++p) {</pre>
  if (smalls[p] > smalls[p - 1]) {
    int q = p * p;
    ++ DC:
    if (1LL * q * q > n) break;
    skip[p] = 1;
    for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
    int ns = 0;
    for (int k = 0; k < s; ++k) {</pre>
      int i = roughs[k];
      if (skip[i]) continue;
      ll d = 1LL * i * p;
      larges[ns] = larges[k] - (d <= v ? larges</pre>
           [smalls[d] - pc] : smalls[n / d]) + pc;
      roughs[ns++] = i;
    s = ns;
    for (int j = v / p; j >= p; --j) {
      int c =
            smalls[j] - pc, e = min(j * p + p, v + 1);
      for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
 }
for (int k = 1; k < s; ++k) {
  const ll m = n / roughs[k];</pre>
  ll t = larges[k] - (pc + k - 1);
  for (int l = 1; l < k; ++l) {</pre>
    int p = roughs[l];
    if (1LL * p * p > m) break;
    t -= smalls[m / p] - (pc + l - 1);
  larges[0] -= t;
}
return larges[0];
```

6.14 DiscreteLog

```
constexpr int kStep = 32000;
unordered_map < int, int > p;
int b = 1;
```

```
for (int i = 0; i < kStep; ++i) {
    p[y] = i;
    y = 1LL * y * x % m;
    b = 1LL * b * x % m;
}
for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;
    if (p.find(s) != p.end()) return i + kStep - p[s];
}
return -1;
}
int DiscreteLog(int x, int y, int m) {
    if (m == 1) return 0;
    int s = 1;
    for (int i = 0; i < 100; ++i) {
        if (s == y) return i;
        s = 1LL * s * x % m;
    }
    if (s == y) return 100;
    int p = 100 + DiscreteLog(s, x, y, m);
    if (fpow(x, p, m) != y) return -1;
    return p;
}</pre>
```

6.15 BerlekampMassey

```
vector<T> BerlekampMassey(const vector<T> &output) {
  vector<T> d(SZ(output) + 1), me, he;
  for (int f = 0, i = 1; i <= SZ(output); ++i) {
    for (int j = 0; j < SZ(me); ++j)
        d[i] += output[i - j - 2] * me[j];
    if ((d[i] -= output[i - 1]) == 0) continue;
    if (me.empty()) {
        me.resize(f = i);
        continue;
    }
    vector<T> o(i - f - 1);
    T k = -d[i] / d[f]; o.pb(-k);
    for (T x : he) o.pb(x * k);
        o.resize(max(SZ(o), SZ(me)));
    for (int j = 0; j < SZ(me); ++j) o[j] += me[j];
    if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
    me = o;
    }
    return me;
}
```

6.16 Theorems

• Cramer's rule

• Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

• Kirchhoff's Theorem

Denote L be a $n\times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.

 Let $T_{n,k}$ be the number of labeled forests on n vertices with
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős—Gallai theorem

A sequence of nonnegative integers $d_1 \ge \cdots \ge d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if

and only if $d_1+\dots+d_n$ is even and $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$

holds for every $1 \le k \le n$.

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,\,b_1),\,\ldots\,,\,(a_n,\,b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1.$

- Möbius inversion formula
 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
 - $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Spherical cap
 - A portion of a sphere cut off by a plane.
 - $r\colon$ sphere radius, $a\colon$ radius of the base of the cap, $h\colon$ height of the cap, θ : $\operatorname{arcsin}(a/r)$.
 - Volume $=\pi h^2(3r-h)/3 = \pi h(3a^2+h^2)/6 = \pi r^3(2+\cos\theta)(1-\cos\theta)$ $\cos\theta)^2/3$.
 - Area $= 2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$.
- Lagrange multiplier
 - Optimize $f(x_1,...,x_n)$ when k constraints $g_i(x_1,...,x_n) = 0$.
 - Lagrangian function $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)=$ $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n).$
 - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines
 - Line 1: ${m v}_1\!=\!{m p}_1\!+\!t_1{m d}_1$
 - Line 2: $m{v}_2\!=\!m{p}_2\!+\!t_2m{d}_2$
 - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
 - $\boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}$
 - $\boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$
 - $c_1 = p_1 + \frac{(p_2 p_1) \cdot n_2}{d_1 \cdot n_2} d_1$
 - $c_2 = p_2 + \frac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2$

Estimation 6.17

- Estimation
 - The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 2000000 for n < 1e19.
 - The number of ways of writing n as a sum of positive integers, disregarding the order of the summands. 1,1,2,3,5,7,11,15,22,30
 - for $n=0\sim 9$, 627 for n=20, $\sim 2e5$ for n=50, $\sim 2e8$ for n=100. Total number of partitions of n distinct elements: B(n) = 1,1,2,5,15,52,203,877,4140,21147,115975,678570,4213597,27644437,190899322,....

6.18 EuclideanAlgorithms

- $m = |\frac{an+b}{a}|$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c - b - 1, a, m - 1) \\ -h(c, c - b - 1, a, m - 1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{b} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ - 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.19 Numbers

• Bernoulli numbers

$$\begin{split} &B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ &\sum_{j=0}^m \binom{m+1}{j} B_j = 0 \text{, EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!} \\ &S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k^+ n^{m+1-k} \end{split}$$

ullet Stirling numbers of the second kind Partitions of n distinct elements into exactly \boldsymbol{k} groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), \\ S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n = \sum_{i=0}^n S(n,i)(x)_i$$
 • Pentagonal number theorem

• Pentagonal number theorem
$$\prod_{n=1}^{\infty}(1-x^n)=1+\sum_{k=1}^{\infty}(-1)^k\left(x^{k(3k+1)/2}+x^{k(3k-1)/2}\right)$$
 • Catalan numbers
$$C_n^{(k)}=\frac{1}{(k-1)n+1}\binom{kn}{n}$$

$$C_n^{(k)}(x)=1+\sum_{k=1}^{\infty}(C_n^{(k)}(x))^k$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1j:s s.t. $\pi(j) \ge j$, k j:s s.t. $\pi(j) > j$. E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k-1) + ($$

6.20 GenerationFunctions

- Ordinary Generating Function $A(x)\!=\!\sum_{i\geq 0}\!a_ix^i$
 - $A(rx) \Rightarrow r^n a_n$
 - $A(x) + B(x) \Rightarrow a_n + b_n$
 - $A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\cdots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x)' \Rightarrow na_n$
 - $\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i > 0} \frac{a_i}{i!} x_i$
 - $A(x)+B(x) \Rightarrow a_n+b_n$

 - $A^{(k)}(x) \Rightarrow a_{n+k}$ $A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
 - $A(x)^k \Rightarrow \sum_{i_1+i_2+\cdots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $xA(x) \Rightarrow na_n$
- Special Generating Function $(1+x)^n = \sum_{i \geq 0} \binom{n}{i} x^i$

 - $\frac{1}{(1-x)^n} = \sum_{i \ge 0} {i \choose n-1} x^i$
 - $S_k = \sum_{x=1}^n x^k$: $S = \sum_{p=0}^\infty x^p = \frac{e^x e^{x(n+1)}}{1 e^x}$

Polynomials

7.1 NTT (FFT)

```
#define base ll // complex < double >
#define N 524288
// const double PI = acosl(-1);
const ll mod = 998244353, g = 3;
base omega[4 * N], omega_[4 * N];
int rev[4 * N];
ll fpow(ll b, ll p);
ll inverse(ll a) { return fpow(a, mod - 2); }
void calcW(int n) {
  ll r = fpow(g, (mod - 1) / n), invr = inverse(r);
  omega[0] = omega_[0] = 1;
for (int i = 1; i < n; i++) {
  omega[i] = omega[i - 1] * r % mod;</pre>
     omega_[i] = omega_[i - 1] * invr % mod;
```

```
// double arg = 2.0 * PI / n;
  // for (int i = 0; i < n; i++)
// {
  //
           omega[i] = base(cos(i * arg), sin(i * arg));
  //
           omega_[i] = base(cos(-i * arg), sin(-i *
  11
           ara));
  // }
}
void calcrev(int n) {
  int k = __lg(n);
for (int i = 0; i < n; i++) rev[i] = 0;</pre>
   for (int i = 0; i < n; i++)</pre>
     for (int j = 0; j < k; j++)
       if (i & (1 << j)) rev[i] ^= 1 << (k - j - 1);</pre>
vector<base> NTT(vector<base> poly, bool inv) {
  base *w = (inv ? omega_ : omega);
   int n = poly.size();
   for (int i = 0; i < n; i++)</pre>
     if (rev[i] > i) swap(poly[i], poly[rev[i]]);
   for (int len = 1; len < n; len <<= 1) {</pre>
     int arg = n / len / 2;
     for (int i = 0; i < n; i += 2 * len)
        for (int j = 0; j < len; j++) {</pre>
          base odd =
           w[j * arg] * poly[i + j + len] % mod;
          poly[i + j + len] =
   (poly[i + j] - odd + mod) % mod;
          poly[i + j] = (poly[i + j] + odd) % mod;
   if (inv)
     for (auto &a : poly) a = a * inverse(n) % mod;
   return poly;
vector<base> mul(vector<base> f, vector<base> g) {
  int sz = 1 << (__lg(f.size() + g.size() - 1) + 1);</pre>
   f.resize(sz), g.resize(sz);
  calcrev(sz);
  calcW(sz);
  f = NTT(f, 0), g = NTT(g, 0);
  for (int i = 0; i < sz; i++)
f[i] = f[i] * g[i] % mod;</pre>
   return NTT(f, 1);
}
```

7.2 FHWT

```
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
     for (int i = 0; i < n; i += L)</pre>
       for (int j = i; j < i + (L >> 1); ++j)
         a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[
    N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
     subset_convolution(int *a, int *b, int *c, int L) {
   // c_k = \sum_{i=0}^{n} \{i \mid j = k, i \& j = 0\} a_i * b_j
  int n = 1 << L;</pre>
  for (int i = 1; i < n; ++i)</pre>
     ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
     f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)
  fwt(f[i], n, 1), fwt(g[i], n, 1);
for (int i = 0; i <= L; ++i)</pre>
     for (int j = 0; j <= i; ++j)</pre>
       for (int x = 0; x < n; ++x)
h[i][x] += f[j][x] * g[i - j][x];</pre>
  for (int i = 0; i <= L; ++i)
  fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)</pre>
     c[i] = h[ct[i]][i];
```

7.3 PolynomialOperations

```
poly inv(poly A) {
   A.resize(1 << (__lg(A.size() - 1) + 1));</pre>
  poly B = {inverse(A[0])};
  for (int n = 1; n < A.size(); n += n) {</pre>
    poly pA(A.begin(), A.begin() + 2 * n);
    calcrev(4 * n);
    calcW(4 * n);
    pA.resize(4 * n);
    B.resize(4 * n);
    pA = NTT(pA, 0);
    B = NTT(B, 0);
    for (int i = 0; i < 4 * n; i++)
      B[i] =
        ((B[i] * 2 - pA[i] * B[i] % mod * B[i]) % mod +
          mod) %
    B = NTT(B, 1);
B.resize(2 * n);
  }
  return B;
pair<poly, poly> div(poly A, poly B) {
  if (A.size() < B.size()) return make_pair(poly(), A);</pre>
  int n = A.size(), m = B.size();
  poly revA = A, invrevB = B;
  reverse(revA.begin(), revA.end());
  reverse(invrevB.begin(), invrevB.end());
  revA.resize(n - m + 1);
  invrevB.resize(n - m + 1);
  invrevB = inv(invrevB);
  poly Q = mul(revA, invrevB);
  Q.resize(n - m + 1);
  reverse(Q.begin(), Q.end());
  poly R = mul(Q, B);
  R.resize(m - 1);
  for (int i = 0; i < m - 1; i++)
    R[i] = (A[i] - R[i] + mod) \% mod;
  return make_pair(Q, R);
ll fast_kitamasa(ll k, poly A, poly C) {
  int n = A.size();
  C.emplace_back(mod - 1);
  poly Q, R = \{0, 1\}, F = \{1\};
  R = div(R, C);
  while (k) {
    if (k & 1) F = div(mul(F, R), C);
    R = div(mul(R, R), C);
    k >>= 1:
  ll ans = 0;
  for (int i = 0; i < F.size(); i++)</pre>
    ans = (ans + A[i] * F[i]) % mod;
  return ans;
vector<ll> fpow(vector<ll> f, ll p, ll m) {
  int b = 0;
  while (b < f.size() && f[b] == 0) b++;</pre>
  f = vector<ll>(f.begin() + b, f.end());
  int n = f.size():
  f.emplace_back(0);
  vector<ll> q(min(m, b * p), 0);
  q.emplace_back(fpow(f[0], p));
  for (int k = 0; q.size() < m; k++) {</pre>
    ll res = 0;
    for (int i = 0; i < min(n, k + 1); i++)</pre>
      res = (res +
               p * (i + 1) % mod * f[i + 1] % mod *
                 q[k - i + b * p]) %
        mod:
    for (int i = 1; i < min(n, k + 1); i++)</pre>
      res = (res
               f[i] * (k - i + 1) % mod *
                 q[k - i + 1 + b * p]) %
        mod;
    res = (res < 0 ? res + mod : res) *
      inv(f[0] * (k + 1) % mod) % mod;
    q.emplace back(res);
  }
  return q;
```

7.4 NewtonMethod+MiscGF

Given F(x) where

| }

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that F(P) = 0 can be found iteratively. Denote by Q_k the polynomial such that $F(Q_k) = 0$ $\pmod{x^{2^k}}$, then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$

- $\bullet \ A^{-1} \colon \ B_{k+1} \! = \! B_{k}(2 \! \! AB_{k}) \ \ \mathrm{mod}x^{2^{k+1}}$
- $\ln A$: $(\ln A)' = \frac{A'}{A}$
- $\exp A$: $B_{k+1} = B_k(1 + A \ln B_k) \mod x^{2^{k+1}}$
- \sqrt{A} : $B_{k+1} = \frac{1}{2}(B_k + AB_k^{-1}) \mod x^{2^{k+1}}$

Geometry

8.1 Basic

```
typedef pair<pdd, pdd> Line;
struct Cir{ pdd 0; double R; };
const double eps = 1e-8;
pdd operator+(pdd a, pdd b)
{ return pdd(a.X + b.X, a.Y + b.Y); }
pdd operator - (pdd a, pdd b) { return pdd(a.X - b.X, a.Y - b.Y); }
pdd operator*(pdd a, double b)
{ return pdd(a.X * b, a.Y * b); }
pdd operator/(pdd a, double b)
{ return pdd(a.X / b, a.Y / b); }
double dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
double cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
double abs2(pdd a)
{ return dot(a, a); }
double abs(pdd a)
{ return sqrt(dot(a, a)); }
int sign(double a)
{ return fabs(a) < eps ? 0 : a > 0 ? 1 : -1; }
int ori(pdd a, pdd b, pdd c)
{ return sign(cross(b - a, c - a)); }
bool collinearity(pdd p1, pdd p2, pdd p3)
{ return sign(cross(p1 - p3, p2 - p3)) == 0; }
bool btw(pdd p1, pdd p2, pdd p3) {
  if (!collinearity(p1, p2, p3)) return 0;
  return sign(dot(p1 - p3, p2 - p3)) <= 0;</pre>
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  int a123 = ori(p1, p2, p3);
  int a124 = ori(p1, p2, p4);
  int a341 = ori(p3, p4, p1);
  int a342 = ori(p3, p4, p2);
if (a123 == 0 && a124 == 0)
    return btw(p1, p2, p3) || btw(p1, p2, p4) ||
      btw(p3, p4, p1) || btw(p3, p4, p2);
  return a123 * a124 <= 0 && a341 * a342 <= 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4) {
  double a123 = cross(p2 - p1, p3 - p1);
double a124 = cross(p2 - p1, p4 - p1);
  return (p4
       * a123 - p3 * a124) / (a123 - a124); // C^3 / C^2 | } // return (a, b), ori(p, C[a], C[b]) >= 0
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
    p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
   ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
    (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) \{
  pdd dp = p1 - p0
      , dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
  return q0 + pdd(
      cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
\} // from line p0--p1 to q0--q1, apply to r
```

8.2 ConvexHull

```
sort(dots.begin(), dots.end());
vector<pll> ans(1, dots[0]);
for (int ct = 0; ct < 2; ++ct, reverse(ALL(dots)))</pre>
 for (int i = 1,
       t = SZ(ans); i < SZ(dots); ans.pb(dots[i++]))</pre>
    while (SZ(ans) > t && ori
       (ans[SZ(ans) - 2], ans.back(), dots[i]) <= 0)
      ans.pop_back();
ans.pop_back(), ans.swap(dots);
```

8.3 SortByAngle

```
#define is_neg(k) (
    sign(k.Y) < 0 \mid \mid (sign(k.Y) == 0 \&\& sign(k.X) < 0))
  int A = is_neg(a), B = is_neg(b);
  if (A != B)
    return A < B;
  if (sign(cross(a, b)) == 0)
    return same ? abs2(a) < abs2(b) : -1;</pre>
  return sign(cross(a, b)) > 0;
```

8.4 DisPointSegment

```
if (sign(abs(q0 - q1)) == 0) return abs(q0 - p);
  return fabs(cross(q1 - q0, p - q0) / abs(q0 - q1));
return min(abs(p - q0), abs(p - q1));
}
```

8.5 PointInCircle

```
bool in_cc(const array<pll, 3> &p, pll q) {
    _int128 det = 0;
  for (int i = 0; i < 3; ++i)
    det += __int128(abs2(p[i]) - abs2(q)) *
         cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
  return det > 0; // in: >0, on: =0, out: <0
}
```

8.6 PointInConvex

```
int a = 1, b = SZ(C) - 1, r = !strict;
   if (SZ(C) == 0) return false;
   if (SZ(C) < 3) return r && btw(C[0], C.back(), p);</pre>
   if (ori(C[0], C[a], C[b]) > 0) swap(a, b);
       (C[0], C[a], p) >= r \mid\mid ori(C[0], C[b], p) <= -r)
     return false;
   while (abs(a - b) > 1) {
  int c = (a + b) / 2;
     (ori(C[0], C[c], p) > 0 ? b : a) = c;
   return ori(C[a], C[b], p) < r;</pre>
}
```

8.7 PointTangentConvex

```
return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
 auto gao = [&](int s) {
   return cyc_tsearch(SZ(C), [&](int x, int y)
    { return ori(p, C[x], C[y]) == s; });
  return pii(gao(1), gao(-1));
```

8.8 CircTangentCirc

```
// sign1 = 1 for outer tang, -1 for inter tang
vector<Line> ret;
double d_sq = abs2(c1.0 - c2.0);
if (sign(d_sq) == 0) return ret;
double d = sqrt(d_sq);
pdd v = (c2.0 - c1.0) / d;
double c = (c1.R - sign1 * c2.R) / d;
if (c * c > 1) return ret;
double h = sqrt(max(0.0, 1.0 - c * c));
for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
    v.Y * c + sign2 * h * v.X);
  pdd p1 = c1.0 + n * c1.R;
  pdd p2 = c2.0 + n * (c2.R * sign1);
```

8.9 LineConvexIntersect

```
return cyc_tsearch(SZ(C), [&](int a, int b) {
     return cross(dir, C[a]) > cross(dir, C[b]);
  });
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 || cmpL(B) > 0)
  return pii(-1, -1); // no collision
auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r; ) {
      int m = ((l + r + (l < r? 0 : n)) / 2) % n;
       (cmpL(m) == cmpL(t) ? l : r) = m;
    return (l + !cmpL(r)) % n;
  };
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
  if (!
       cmpL(res.X) && !cmpL(res.Y)) // along side i, i+1
    switch ((res.X - res.Y + n + 1) % n) {
       case 0: return pii(res.X, res.X);
       case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned
        in the same order as the line hits the convex */
  return res;
} // convex cut: (r, l]
```

8.10 CircIntersectCirc

8.11 PolyIntersectCirc

```
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r && B
        < PI/2) S -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S:
double area_poly_circle(const
     vector<pdd> poly,const pdd &0,const double r){
```

8.12 MinMaxEnclosingRect

```
pdd solve(vector<pll>> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
  hull(dots);
  double Max = 0, Min = INF, deg;
  int n = SZ(dots):
  dots.pb(dots[0]);
  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {</pre>
    pll nw = vec(i + 1);
    while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
      u = (u + 1) \% n;
    while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
      r = (r + 1) \% n;
    if (!i) l = (r + 1) % n;
    while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))
    l = (l + 1) % n;</pre>
    Min = min(Min, (double)(dot(nw, vec(r)) - dot
        (nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));
    deg = acos(dot(diff(r
    , l), vec(u)) / abs(diff(r, l)) / abs(vec(u))); deg = (qi - deg) / 2;
    return pdd(Min, Max);
```

8.13 MinEnclosingCircle

8.14 CircleCover

```
struct CircleCover {
  int C:
  Cir c[N];
  bool g[N][N], overlap[N][N];
  // Area[i] : area covered by at least i circles
  double Area[ N ];
  void init(int _C){ C = _C;}
  struct Teve {
    pdd p; double ang; int add;
    Teve() {}
    Teve(pdd _a
    , double _b, int _c):p(_a), ang(_b), add(_c){}
bool operator<(const Teve &a)const</pre>
    {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
bool disjuct(Cir &a, Cir &b, int x)
  {return sign(abs(a.0 - b.0) - a.R - b.R) > x;}
bool contain(Cir &a, Cir &b, int x)
  {return sign(a.R - b.R - abs(a.0 - b.0)) > x;}
  bool contain(int i, int j) {
    /* c[j] is non-strictly in c[i]. */
    return (sign
         (c[i].R - c[j].R) > 0 \mid | (sign(c[i].R - c[j].
         R) == 0 && i < j)) && contain(c[i], c[j], -1);
  void solve(){
```

```
fill_n(Area, C + 2, 0);
    for(int i = 0; i < C; ++i)</pre>
      for(int j = 0; j < C; ++j)</pre>
        overlap[i][j] = contain(i, j);
    for(int i = 0; i < C; ++i)</pre>
      for(int j = 0; j < C; ++j)</pre>
        g[i][j] = !(overlap[i][j] || overlap[j][i] ||
             disjuct(c[i], c[j], -1));
    for(int i = 0; i < C; ++i){</pre>
      int E = 0, cnt = 1;
      for(int j = 0; j < C; ++j)</pre>
        if(j != i && overlap[j][i])
           ++cnt;
      for(int j = 0; j < C; ++j)</pre>
        if(i != j && g[i][j]) {
           pdd aa, bb;
           CCinter(c[i], c[j], aa, bb);
           double A =
                atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);
           double B =
                atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
           eve[E++] = Teve
               (bb, B, 1), eve[E++] = Teve(aa, A, -1);
           if(B > A) ++cnt;
      if(E == 0) Area[cnt] += pi * c[i].R * c[i].R;
      else{
        sort(eve, eve + E);
         eve[E] = eve[0];
        for(int j = 0; j < E; ++j){
           cnt += eve[j].add;
           Area[cnt
               ] += cross(eve[j].p, eve[j + 1].p) * .5;
           double theta = eve[j + 1].ang - eve[j].ang;
           if (theta < 0) theta += 2. * pi;</pre>
           Area[cnt] += (theta
                - sin(theta)) * c[i].R * c[i].R * .5;
      }
    }
 }
};
```

8.15 Trapezoidalization

```
struct SweepLine {
  struct cmp {
    cmp(const SweepLine &_swp): swp(_swp) {}
    bool operator()(int a, int b) const {
  if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
        return swp.slope_cmp(a, b);
      return swp.get_y(a) + swp.eps < swp.get_y(b);</pre>
    const SweepLine &swp;
  } _cmp;
  T curTime, eps, curQ;
  vector<Line> base;
  multiset < int , cmp > sweep;
 multiset<pair<T, int>> event;
  vector<typename multiset<int, cmp>::iterator> its;
      <typename multiset<pair<T, int>>::iterator> eits;
  bool slope_cmp(int a, int b) const {
    assert(a != -1);
    if (b == -1) return 0;
    return sign(cross(base
        [a].Y - base[a].X, base[b].Y - base[b].X)) < 0;</pre>
  T get_y(int idx) const {
    if (idx == -1) return curQ;
    Line l = base[idx];
    if (l.X.X == l.Y.X) return l.Y.Y;
    return ((curTime - l.X.X) * l.Y.Y
        + (l.Y.X - curTime) * l.X.Y) / (l.Y.X - l.X.X);
  void insert(int idx) {
    its[idx] = sweep.insert(idx);
    if (its[idx] != sweep.begin())
      update_event(*prev(its[idx]));
    update_event(idx);
    event.emplace(base[idx].Y.X, idx + 2 * SZ(base));
  void erase(int idx) {
    assert(eits[idx] == event.end());
    auto p = sweep.erase(its[idx]);
    its[idx] = sweep.end();
```

```
if (p != sweep.begin())
      update_event(*prev(p));
  void update_event(int idx) {
    if (eits[idx] != event.end())
      event.erase(eits[idx]);
    eits[idx] = event.end();
    auto nxt = next(its[idx]);
    if (nxt ==
          sweep.end() || !slope_cmp(idx, *nxt)) return;
    auto t = intersect(base[idx].
        X, base[idx].Y, base[*nxt].X, base[*nxt].Y).X;
    if (t + eps < curTime || t</pre>
          >= min(base[idx].Y.X, base[*nxt].Y.X)) return;
    eits[idx] = event.emplace(t, idx + SZ(base));
  void swp(int idx) {
    assert(eits[idx] != event.end());
    eits[idx] = event.end();
    int nxt = *next(its[idx]);
    swap((int&)*its[idx], (int&)*its[nxt]);
    swap(its[idx], its[nxt]);
    if (its[nxt] != sweep.begin())
      update_event(*prev(its[nxt]));
    update_event(idx);
  }
  // only expected to call the functions below
  SweepLine(T t, T e, vector
       <Line> vec): _cmp(*this), curTime(t), eps(e)
        curQ(), base(vec), sweep(_cmp), event(), its(SZ
    (vec), sweep.end()), eits(SZ(vec), event.end()) {
for (int i = 0; i < SZ(base); ++i) {</pre>
      auto &[p, q] = base[i];
       if (p > q) swap(p, q);
      if (p.X <= curTime && curTime <= q.X)</pre>
        insert(i);
       else if (curTime < p.X)</pre>
         event.emplace(p.X, i);
  void setTime(T t, bool ers = false) {
    assert(t >= curTime);
    while (!event.empty() && event.begin()->X <= t) {</pre>
      auto [et, idx] = *event.begin();
      int s = idx / SZ(base);
      idx %= SZ(base);
      if (abs(et - t) <= eps && s == 2 && !ers) break;</pre>
      curTime = et;
      event.erase(event.begin());
      if (s == 2) erase(idx);
      else if (s == 1) swp(idx);
      else insert(idx);
    curTime = t;
  T nextEvent() {
    if (event.empty()) return INF;
    return event.begin()->X;
  int lower_bound(T y) {
    curQ = y;
    auto p = sweep.lower_bound(-1);
    if (p == sweep.end()) return -1;
    return *p;
  }
};
```

8.16 TriangleHearts

```
p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
double m = 2. * (x1 * y2 - y1 * x2);
  center.X = (x1 * x1)
        * y2 - x2 * x2 * y1 + y1 * y2 * (y1 - y2)) / m;
  center.Y = (x1 * x2)
       * (x2 - x1) - y1 * y1 * x2 + x1 * y2 * y2) / m;
  return center + p0;
pdd incenter
    (pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
  double a =
      abs(p2 - p3), b = abs(p1 - p3), c = abs(p1 - p2);
  double s = a + b + c:
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
```

```
{ return (p1 + p2 + p3) / 3; } pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter
     (p1, p2, p3) * 3 - circenter(p1, p2, p3) * 2; }
```

8.17 HalfPlaneIntersect

```
{ return pll(cross(a.Y
- a.X, b.X - a.X), cross(a.Y - a.X, b.Y - a.X)); }
bool isin(Line l0, Line l1, Line l2) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(l0, l2);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
        n (__int128)
a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
  return (
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(ALL(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
       return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(
         dq.back().Y - dq.back().X, p.Y - p.X, 0) == -1)
       continue:
     while (SZ(dq
         ) >= 2 && !isin(p, dq[SZ(dq) - 2], dq.back()))
    dq.pop_back();
while (SZ(dq) >= 2 && !isin(p, dq[0], dq[1]))
      dq.pop_front();
    dq.pb(p);
  while (SZ(dq)
        >= 3 && !isin(dq[0], dq[SZ(dq) - 2], dq.back()))
     dq.pop_back();
  while (SZ(dq) >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
    dq.pop_front();
  return vector < Line > (ALL(dq));
```

8.18 RotatingSweepLine

```
int n = SZ(ps), m = 0;
  vector<int> id(n), pos(n);
vector<pii> line(n * (n - 1));
   for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++j)
       if (i != j) line[m++] = pii(i, j);
   sort(ALL(line), [&](pii a, pii b) {
     return cmp(ps[a.Y] - ps[a.X], ps[b.Y] - ps[b.X]);
  }); // cmp(): polar angle compare
   iota(ALL(id), 0);
  sort(ALL(id), [&](int a, int b) {
  if (ps[a].Y != ps[b].Y) return ps[a].Y < ps[b].Y;</pre>
     return ps[a] < ps[b];</pre>
  \}); // initial order, since (1, 0) is the smallest
  for (int i = 0; i < n; ++i) pos[id[i]] = i;
for (int i = 0; i < m; ++i) {</pre>
     auto l = line[i];
     // do something
     tie(pos[l.X], pos[l.Y], id[pos[l.X]], id[pos[l.Y
          ]]) = make_tuple(pos[l.Y], pos[l.X], l.Y, l.X);
  }
}
```

8.19 DelaunayTriangulation

```
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri,
each points are u.p[(i+1)\%3], u.p[(i+2)\%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const
     ll inf = MAXC * MAXC * 100; // lower_bound unknown
```

```
struct Tri:
struct Edge {
  Tri* tri; int side;
  Edge(): tri(0), side(0){}
  Edge(Tri* _tri, int _side): tri(_tri), side(_side){}
struct Tri {
  pll p[3];
  Edge edge[3];
  Tri* chd[3];
  Tri() {}
  Tri(const pll& p0, const pll& p1, const pll& p2) {
    p[0] = p0; p[1] = p1; p[2] = p2;
    chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return !!chd[0] + !!chd[1] + !!chd[2];
  bool contains(pll const& q) const {
    for (int i = 0; i < 3; ++i)</pre>
      if (ori(p[i], p[(i + 1) % 3], q) < 0)</pre>
        return 0;
    return 1;
pool[N * 10], *tris;
void edge(Edge a, Edge b) {
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig() {
    the_root
         = // Tri should at least contain all points
      new(tris++) Tri(pll(-inf, -inf),
            pll(inf + inf, -inf), pll(-inf, inf + inf));
  Tri* find(pll p) { return find(the_root, p); }
  void add_point(const
       pll &p) { add_point(find(the_root, p), p); }
  Tri* the_root;
  static Tri* find(Tri* root, const pll &p) {
    while (1) {
      if (!root->has_chd())
        return root;
      for (int i = 0; i < 3 && root->chd[i]; ++i)
        if (root->chd[i]->contains(p)) {
          root = root->chd[i];
          break;
        }
    assert(0); // "point not found"
  void add_point(Tri* root, pll const& p) {
    Tri* t[3];
     /* split it into three triangles */
    for (int i = 0; i < 3; ++i)</pre>
      t[i] = new(tris
           ++) Tri(root->p[i], root->p[(i + 1) % 3], p);
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 0), Edge(t[(i + 1) \% 3], 1));
    for (int i = 0; i < 3; ++i)</pre>
      edge(Edge(t[i], 2), root->edge[(i + 2) % 3]);
    for (int i = 0; i < 3; ++i)</pre>
      root->chd[i] = t[i];
    for (int i = 0; i < 3; ++i)</pre>
      flip(t[i], 2);
  void flip(Tri* tri, int pi) {
    Tri* trj = tri->edge[pi].tri;
    int pj = tri->edge[pi].side;
    if (!trj) return;
    if (!in_cc(tri->p
    [0], tri->p[1], tri->p[2], trj->p[pj])) return;
/* flip edge between tri,trj */
    Tri* trk = new(tris++) Tri
        (tri->p[(pi + 1) % 3], trj->p[pj], tri->p[pi]);
    Tri* trl = new(tris++) Tri
    (trj->p[(pj + 1) % 3], tri->p[pi], trj->p[pj]);
edge(Edge(trk, 0), Edge(trl, 0));
    edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
    edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
    edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
    edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
    tri->chd
        [0] = trk; tri->chd[1] = trl; tri->chd[2] = 0;
```

```
tri->chd
         [0] = trk; trj->chd[1] = trl; trj->chd[2] = 0;
    flip(trk, 1); flip(trk, 2);
flip(trl, 1); flip(trl, 2);
};
vector<Tri*> triang; // vector of all triangle
set<Tri*> vst;
void go(Tri* now) { // store all tri into triang
 if (vst.find(now) != vst.end())
   return:
  vst.insert(now);
 if (!now->has_chd())
    return triang.pb(now);
  for (int i = 0; i < now->num_chd(); ++i)
    go(now->chd[i]);
void build(int n, pll* ps) { // build triangulation
  tris = pool; triang.clear(); vst.clear();
  random_shuffle(ps, ps + n);
 Trig tri; // the triangulation structure
for (int i = 0; i < n; ++i)</pre>
    tri.add_point(ps[i]);
  go(tri.the_root);
```

8.20 VonoroiDiagram

8.21 3DConeDistance

```
const double eps = 1e-12, pi = acos(-1);
struct point {
 ld x, y, z;
ld r, h;
bool on_edge(point p) {
  return abs(p.z) < eps &&</pre>
    abs(p.x * p.x + p.y * p.y - r * r) < eps;
bool on bottom(point p) { return abs(p.z) < eps; }</pre>
ld straight(point a, point b) {
  return sqrtl((a.x - b.x) * (a.x - b.x) +
    (a.y - b.y) * (a.y - b.y) +
    (a.z - b.z) * (a.z - b.z));
ld arc(point a, point b) {
 if (abs(a.x) < eps && abs(a.y) < eps &&</pre>
    abs(a.z - h) < eps)
    return straight(a, b);
  swap(a, b);
  if (abs(a.x) < eps && abs(a.y) < eps &&</pre>
   abs(a.z - h) < eps)
    return straight(a, b);
  swap(a, b);
  ld ans = 1e18;
  for (int t = 0; t <= 1; t++) {</pre>
    double alpha = (a.x < 0 ? pi + atan2l(-a.y, -a.x)
                              : atan2l(a.y, a.x));
    if (alpha < 0) alpha += 2 * pi;</pre>
    alpha *= r / sqrtl(h * h + r * r);
    double lA = sqrtl(
      a.x * a.x + a.y * a.y + (a.z - h) * (a.z - h));
    double beta = (b.x < 0 ? pi + atan2l(-b.y, -b.x)
                            : atan2l(b.y, b.x));
    if (beta < 0) beta += 2 * pi;</pre>
    beta *= r / sqrtl(h * h + r * r);
```

```
double lB = sqrtl(
      b.x * b.x + b.y * b.y + (b.z - h) * (b.z - h));
    double Ax = lA * cosl(alpha),
            Ay = lA * sinl(alpha), Bx = lB * cosl(beta),
           By = lB * sinl(beta);
    ans = min(ans,
      sqrtl((Ax - Bx) * (Ax - Bx) +
        (Ay - By) * (Ay - By)));
    a.x = -a.x;
    b.x = -b.x;
  }
  return ans;
}
ld distance(point a, point b) {
  if (on_bottom(a) && !on_edge(a) && !on_bottom(b))
    return 1e18;
  if (on_bottom(b) && !on_edge(b) && !on_bottom(a))
    return 1e18;
  if (on_bottom(a) && on_bottom(b))
    return straight(a, b);
  return arc(a, b);
const int s = 90;
ld ans = 1e18, theta = pi / 180.0;
point a, b;
vector<point> p;
vector<ld> ar;
vector<ld> dis, from, center = {0.0, pi};
vector<int> vis;
void iterate() {
  p = {a, b};
  ar = {-1e9, -1e9};
  dis.clear();
  from.clear();
  vis.clear();
  for (auto d : center)
    for (int i = -s; i <= s; i++) {</pre>
      p.emplace_back(point{r * cosl(d + i * theta),
    r * sinl(d + i * theta), 0});
      ar.emplace_back(d + i * theta);
  int n = p.size();
  dis.resize(n);
  from.resize(n);
  vis.resize(n);
  for (int i = 1; i < n; i++) dis[i] = 1e18;</pre>
  for (int t = 0; t < n; t++) {</pre>
    int u = 0;
    ld d = 1e18;
    for (int i = 0; i < n; i++)</pre>
      if (!vis[i] && dis[i] < d) u = i, d = dis[i];</pre>
    vis[u] = 1;
    for (int v = 0; v < n; v++)</pre>
      if (d + distance(p[u], p[v]) < dis[v]) {
        from[v] = u;
        dis[v] = distance(p[u], p[v]) + d;
      }
  ans = min(ans, dis[1]);
  center.clear();
  if (from[1] > 1) center.emplace_back(ar[from[1]]);
  if (from[from[1]] > 1)
    center.emplace_back(ar[from[from[1]]]);
  theta /= 20.0;
double cone_distance(
  double r, double h, point a, point b) {
  if (on_bottom(a) && on_bottom(b))
    return straight(a, b);
  else {
    for (int t = 1; t <= 10; t++) iterate();</pre>
    return ans;
```

9 Misc

9.1 MoAlgoWithModify

```
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
  int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
    L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
  bool operator<(const Query &q) const {</pre>
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    if (RBid != q.RBid) return RBid < q.RBid;
    return T < b.T;
void solve(vector<Query> query) {
  sort(ALL(query));
  int L=0, R=0, T=-1;
  for (auto q : query) {
    while (T < q.T) addTime(L, R, ++T); // TODO
    while (T > q.T) subtime(L, R, T--); // TODO while (R < q.R) add(arr[++R]); // TODO
    while (L > q.L) add(arr[--L]); // TODO
    while (R > q.R) sub(arr[R--]); // TODO
    while (L < q.L) sub(arr[L++]); // TODO
    // answer query
```

9.2 MoAlgoOnTree

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord[in[u]] = ord[out[u]] = u
4) bitset < MAXN > inset
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
    int c = LCA(u, v);
    if (c == u // c == v)
    q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
    else
    q.lca = c, q.L = out[v], q.R = in[u];
q.Lid = q.L / blk;
  bool operator < (const Query &q) const {
    if (LBid != q.LBid) return LBid < q.LBid;</pre>
    return R < q.R;
7:
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
    else add(arr[x]); // TODO
    inset[x] = ~inset[x];
void solve(vector<Query> query) {
  sort(ALL(query));
  int L = 0, R = 0;
  for (auto q : query) {
    while (R < q.R) flip(ord[++R]);
    while (L > q.L) flip(ord[--L]);
    while (R > q.R) flip(ord[R--]);
    while (L < q.L) flip(ord[L++]);
    if (~q.lca) add(arr[q.lca]);
    // answer query
    if (~q.lca) sub(arr[q.lca]);
```

9.3 MoAlgoAdvanced

- Mo's Algorithm With Addition Only
 - Sort querys same as the normal Mo's algorithm.
 - For each query $\left[l,r\right] :$
 - If $l/blk\!=\!r/blk$, brute-force.
 - If $l/blk \neq curL/blk$, initialize $curL := (l/blk+1) \cdot blk, curR := curL-1$
 - If $r\!>\!cur R$, increase cur R
 - decrease $\operatorname{\it cur} L$ to fit l , and then undo after answering
- Mo's Algorithm With Offline Second Time

- Require: Changing answer \equiv adding $f([l,\!r],\!r\!+\!1)$.
- Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1).
- Part1: Answer all f([1,r],r+1) first.
- Part2: Store $curR \to R$ for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
- Note: You must do the above symmetrically for the left boundaries.

9.4 HilbertCurve

```
ll res = 0;
for (int s = n / 2; s; s >>= 1) {
   int rx = (x & s) > 0;
   int ry = (y & s) > 0;
   res += s * 1ll * s * ((3 * rx) ^ ry);
   if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
   }
}
return res;
} // n = 2^k
```

9.5 SternBrocotTree

- Construction: Root $\frac{1}{1}$, left/right neighbor $\frac{0}{1},\frac{1}{0}$, each node is sum of last left/right neighbor: $\frac{a}{b},\frac{c}{d}\to\frac{a+c}{b+d}$
- Search known $\frac{p}{q}$: keep L-R alternative. Each step can calcaulated in O(1) \Rightarrow total $O(\log C)$.
- Search unknown $\frac{p}{q}$: keep L-R alternative. Each step can calcaulated in $O(\log C)$ checks \Rightarrow total $O(\log^2 C)$ checks.

9.6 CyclicLCS

```
#define LU 1
#define U 2
const int mov[3][2] = \{0, -1, -1, -1, -1, 0\};
int al, bl;
char a[MAXL * 2], b[MAXL * 2]; // 0-indexed
int dp[MAXL * 2][MAXL];
char pred[MAXL * 2][MAXL];
inline int lcs_length(int r) {
  int i = r + al, j = bl, l = 0;
  while (i > r) {
    char dir = pred[i][j];
if (dir == LU) l++;
    i += mov[dir][0];
    j += mov[dir][1];
  return l:
inline void reroot(int r) { // r = new base row
  int i = r, j = 1;
  while (j <= bl && pred[i][j] != LU) j++;</pre>
  if (j > bl) return;
  pred[i][j] = L;
  while (i < 2 * al && j <= bl) {
    if (pred[i + 1][j] == U) {
      pred[i][j] = L;
    } else if (j < bl && pred[i + 1][j + 1] == LU) {</pre>
      i++;
      j++;
      pred[i][j] = L;
    } else {
      j++;
    }
 }
int cyclic_lcs() {
   // a, b, al, bl should be properly filled
  // note: a WILL be altered in process
              -- concatenated after itself
  char tmp[MAXL];
  if (al > bl) {
    swap(al, bl);
    strcpy(tmp, a);
    strcpy(a, b);
    strcpy(b, tmp);
  strcpy(tmp, a);
  strcat(a, tmp);
  // basic lcs
  for (int i = 0; i <= 2 * al; i++) {</pre>
    dp[i][0] = 0;
    pred[i][0] = U;
```

```
for (int j = 0; j <= bl; j++) {</pre>
    dp[0][j] = 0;
    pred[0][j] = L;
  if (dp[i][j - 1] == dp[i][j]) pred[i][j] = L;
else if (a[i - 1] == b[j - 1]) pred[i][j] = LU;
       else pred[i][j] = U;
  // do cyclic lcs
  int clcs = 0;
  for (int i = 0; i < al; i++) {
   clcs = max(clcs, lcs_length(i));</pre>
    reroot(i + 1);
  // recover a
  a[al] = '\0';
  return clcs;
9.7 AllLCS
  vector<int> h(SZ(t));
  iota(ALL(h), 0);
  for (int a = 0; a < SZ(s); ++a) {</pre>
    int v = -1;
    for (int c = 0; c < SZ(t); ++c)</pre>
      if (s[a] == t[c] || h[c] < v)</pre>
        swap(h[c], v);
    // LCS(s[0, a], t[b, c]) =
// c - b + 1 - sum([h[i] >= b] | i <= c)
    // h[i] might become -1 !!
```

9.8 SimulatedAnnealing

```
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
    if (exp(-(nw - ans
        ) / factor) >= (double)(rnd() % base) / base)
        ans = nw;
    factor *= 0.999995;
}
```

9.9 Python

}

|math.isqrt(2) # integer sqrt