

1 Prim's

Algorithm 1: prims(G)

input : The original graph G
output: The minimum spanning tree of G (SPT)

```
1  $SPT \leftarrow smallestEdge(G)$ ;  
2 while  $|edges(SPT)|/2 < |vertices(G)| - 1$  do  
3    $cost \leftarrow \maxIntNumber$ ;  
4   for  $v$  in  $vertices(SPT)$  do  
5     for  $e$  in  $adjacents(G, v)$  do  
6       if  $getEdge(G, v, e) < cost$  and  $e \notin vertices(SPT)$  then  
7          $vertex \leftarrow v$ ;  
8          $minAdjacent \leftarrow e$ ;  
9          $cost \leftarrow getEdge(G, v, e)$ ;  
10      end  
11    end  
12  end  
13   $addVertex(SPT, minAdjacent)$ ;  
14   $addEdge(SPT, vertex, minAdjacent, cost)$ ;  
15 end  
16 return  $SPT$ ;
```

2 Kruskal's

Algorithm 2: `kruskals(G)`

```
input : The original graph  $G$ 
output: The minimum spanning tree of  $G$  ( $SPT$ )
1  $SPT \leftarrow smallestEdge(G)$ ;
2  $H = MinHeap()$ ;
3  $H \leftarrow H + smallestEdge(G)$ ;
4  $edgesCycles \leftarrow set()$ ;
5 while  $|edges(SPT)|/2 < |vertices(G)| - 1$  do
6    $(o, d, cost_{(o,d)}) \leftarrow H.pop()$ ;
7   if  $(o, d) \notin edges(SPT)$  then
8     if  $o \in vertices(SPT)$  then
9        $path \leftarrow path(SPT, o)$ ;
10      for  $p$  in  $path$  do
11        if  $p \in d$  then
12           $edgesCycles.add((o, d))$ 
13        end
14      end
15    end
16    if  $(o, d) \notin edgesCycles$  or  $(o, d) \notin vertices(SPT)$  then
17       $addVertices(o, d)$ ;
18       $addEdge(SPT, o, d, cost_{(o,d)})$ ;
19    end
20  end
21 end
22 return  $SPT$ ;
```
