NTUEE CA 2021

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1. Briefly describe your CPU structure.

Basically, we design our CPU structure from concept of the lecture. We first, decode the instruction to realize what instruction it is, and set the correct control signals. Secondly, by the right control signals, we can compute the desired a ALU output from ALU operations. The next step is handling the memory writeing and register writing part from the control signals. Last, compute the correct PC.

2. Describe how you design the data path of instructions not referred in the lecture slides (jal, jalr, auipc, ...

- JAL set regWrite = 1, jump to PC + imm, and write rd as PC+4.
- JALR
 set regWrite = 1, jump to rs1 + imm, and write rd as PC+4.
- AUIPC set regWrite = 1, write rd as PC + imm.
- SLLIset regWrite = 1, and control alu to shift rs1 left imm bits.
- SRLI
 set regWrite = 1, and control alu to shift rs1 right imm bits.
- SLTI
 set regWrite = 1,ALU_src = 1, and set ALU as a SUB operation. If ALU_output < 0, set rd as 1, else set rd as 0.

3. Describe how you handle multi-cycle instructions (mul)

We instantiate mulDiv module in ALU module. Since mulDiv module requires clk and rst_n inputs, to let mulDiv module share the same clk and rst_n that were used by CHIP module, we added clk and rst_n inputs to ALU module, connect these inputs to clk and rst_n of CHIP module, and let clk and rst_n of mulDiv module connect to clk and rst_n of ALU module. Furthermore, mulDiv need 32 cycles to complete calculation, so we have to make sure the sequential part of CHIP module wait for the calculation to be done. That is, neither change PC nor write registers while ALU module is still computing. Hence, output done, which was assigned to the output ready of mulDiv module, was added to ALU module. done == 1 indicates that the calculation was finished. Finally, we can add conditions to the sequential part:

```
if (ALU_ctrl==3'b100) begin
                                      // If ALU module encounters a MUL instruction
 1
 2
                                     // Modify PC and registers when computation was done
         if (alu_done==1) begin
 3
             PC <= PC_nxt;
 4
             register[0] <= 32'd0;
 5
             for(i = 1; i < 32; i = i+1) begin
                 register[i] <= next_register[i];</pre>
 6
 7
             end
 8
        end
                       // Do not modify PC and registers before computation was done
 9
        else begin
             PC <= PC;
10
             register[0] <= register[0];</pre>
11
             for(i = 1; i < 32; i = i+1) begin
12
                 register[i] <= register[i];</pre>
13
14
             end
15
        end
16
    end
```

It can be seen easily that PC and registers change only when MUL was done.

- 4. Record total simulation time (CYCLE = 10 ns)
- leaf program

```
Simulation complete via $finish(1) at time 255 NS + 0
```

fact program

```
Simulation complete via $finish(1) at time 1575 NS + 0
```

■ hw1 program

```
Simulation complete via $finish(1) at time 485 NS + 0
```

5. Describe your observation

Basically, we follow the architecture of CPU based on what is taught through the lectures. It is quite neat and well-behaved. Furthermore, we find that the risc v instruction are well-organized that we can used the same block of instruction as rd the same block of instruction as op code, which is quite convenient for us to handle.

In bonus problem, we initially used MUL and DIVU instruction to implement multiplication and division. But when we used it, we found MUL and DIVU cost to many CPU cycles so as to easily reach the limit of total simulation time. In order to solve this problem, we changed the upper two instructions into SLLI and SRLI which only cost one CPU cycle and finally made our CPU work functionally.

6. Snapshot the "Register table" in Design Compiler

```
in routine CHIP line 343 in file
    '/home/raid7_2/userb06/b6502152/CA_Final_Project/Verilog/CHIP.v'.
       Register Name
                               | Flip-flop | 995
| Flip-flop | 29
           PC_reg
PC_reg
                                  | Flip-flop |
| Flip-flop |
Statistics for MUX OPs
Warning: /home/raid7_2/userb06/b6502152/CA_Final_Project/Verilog/CHIP.v:410: signed to unsigned conversion occurs. (VER-318)
Warning: /home/raid7_2/userb06/b6502152/CA_Final_Project/Verilog/CHIP.v:417: signed to unsigned conversion occurs. (VER-318)
Inferred memory devices in process
    in routine reg_file line 413 in file
                                 | Flip-flop | 995
| Flip-flop | 29
 statistics for case statements in always block at line 588 in file
'/home/raid7_2/userb06/b6502152/CA_Final_Project/Verilog/CHIP.v'
                 Line
                                        | auto/auto
Statistics for case statements in always block at line 606 in file
'/home/raid7_2/userb06/b6502152/CA_Final_Project/Verilog/CHIP.v'
       stics for case statements in always block at line 618 in file
'/home/raid7_2/userb06/b6502152/CA_Final_Project/Verilog/CHIP.v'
                                         | full/ parallel |
Statistics for case statements in always block at line 633 in file
'/home/raid7_2/userb06/b6502152/CA_Final_Project/Verilog/CHIP.v'
                                        | full/ parallel |
Inferred memory devices in process
in routine mulDiv line 661 in file
                                  | Flip-flop |
```

7. work distribution table

student id	Workload
B06502152 許書銓	Task 1 - leaf, Writing report
B07204024 李昱呈	Bonus - hw1, Writing report
B03208039 鐘友隸	Task 2 - fact, Writing report