

Linneuniversitetet

Kalmar Växjö

Report

Assignment 4



Author: LEONARDO PEDRO

Semester: Spring 2020 Email <u>lr222qp@student.lnu.se</u>

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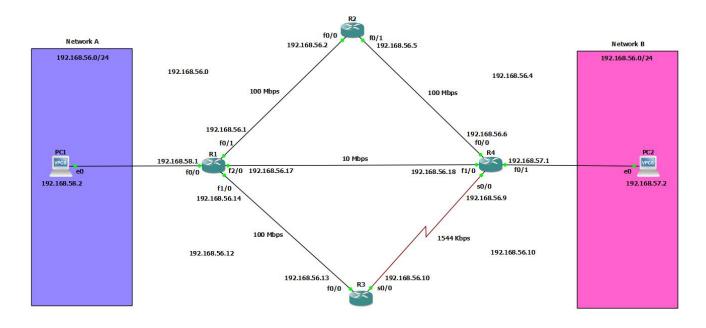


Figure 1. Topology picture from GNS3



Figure 2. Ping from R1 to VPC1

```
PC2
                                                                                    1 🕀
        PC1
                                 R2
                                              R3
                                                           R4
                                                                        R1 ×
                                                                                                              ×
*Mar 1 00:06:14.343: %SYS-5-CONFIG_I: Configured from console by console
R1#ping 192.168.58.2
Type escape sequence to abort.
Sending 5, 100-byte ICMP Echos to 192.168.58.2, timeout is 2 seconds:
Success rate is 0 percent (0/5)
R1#ping 192.168.58.2
Type escape sequence to abort.
Sending 5, 100-byte ICMP Echos to 192.168.58.2, timeout is 2 seconds:
Success rate is 0 percent (0/5)
R1#ping 192.168.56.18
Type escape sequence to abort.
Sending 5, 100-byte ICMP Echos to 192.168.56.18, timeout is 2 seconds:
Success rate is 100 percent (5/5), round-trip min/avg/max = 60/68/72 ms
R1#sh ip int br
                               IP-Address
                                                  OK? Method Status
FastEthernet0/0
Serial0/0
                               192.168.58.1
                                                 YES NVRAM up up
YES NVRAM administratively down down
                              unassigned
192.168.56.1
                                                 YES NVRAM up
YES NVRAM administratively down down
Serial0/1
                               unassigned
                                                 YES NVRAM administratively down down
YES NVRAM up up
                              unassigned
192.168.56.14
Serial0/2
                                                 YES NVRAM up
R1#ping 192.168.56.18
 ending 5, 100-byte ICMP Echos to 192.168.56.18, timeout is 2 seconds:
Success rate is 100 percent (5/5), round-trip min/avg/max = 40/65/76 ms
                      Solar-PuTTY free tool
                                                            © 2019 SolarWinds Worldwide, LLC. All rights reserved.
  solarwinds
```

Figure 3. Ping from R1 to R4

```
PC1> ping 192.168.57.2

*192.168.58.1 icmp_seq=1 ttl=255 time=9.797 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=2 ttl=255 time=3.978 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=3 ttl=255 time=0.988 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=4 ttl=255 time=2.157 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=5 ttl=255 time=1.463 ms (ICMP type:3, code:1, Destination host unreachable)

PC1> [
```

Figure 4. Ping from VPC1 to VPC2

NM-1FE-TX

NM-1FE-TX is a network module with one port fast Ethernet that offers a single auto sensing 10/100 TX connection with the use of an RJ-45 Connector and its interface is used on the Cisco 3600 and 3700 series routers. It also supports VLAN deployment in order to make groups work flow way more easy that enables to move and switch within the network.

We use NM-1FE-TX because NM-4T doesnt support async mode and NM-16ESW has 16 ports thats way more than what we need, NM-NAM doesnt have an external console port and NM-CIDS can only be configured with IDS (promiscuous mode).

WIC-1T

WIC-1T means 1-Port Serial WAN Interface Card that grants serial connections to remote destinations, such as websites or legacy serial network devices.

We use WIC-1T because WIC-2T supports 8Mbps max speed per port

Difference between /24 and /30 subnet?

/30 stands for 255.255.255.252 /24 stands for 255.255.255.0

So /30 is way more efficent because whenever we need to connect to the network we wont be wasting any IP addresses.

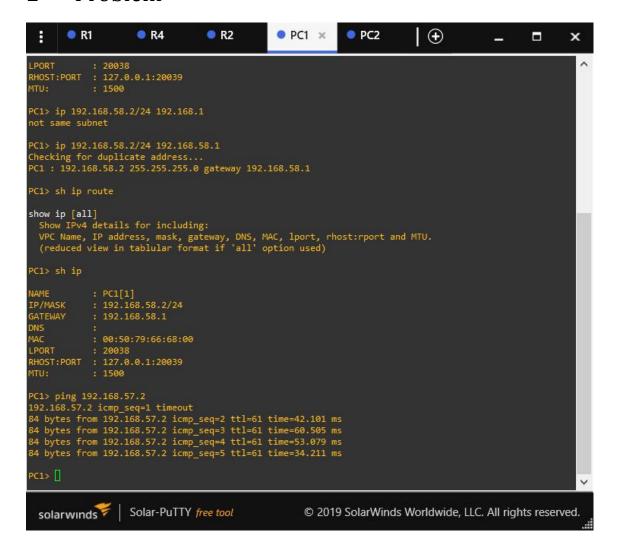


Figure 5. Ping from VPC1 to VPC2

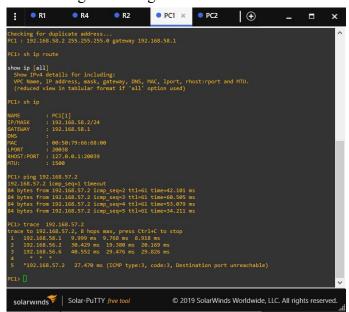


Figure 6. Tracing from VPC1 to VPC2

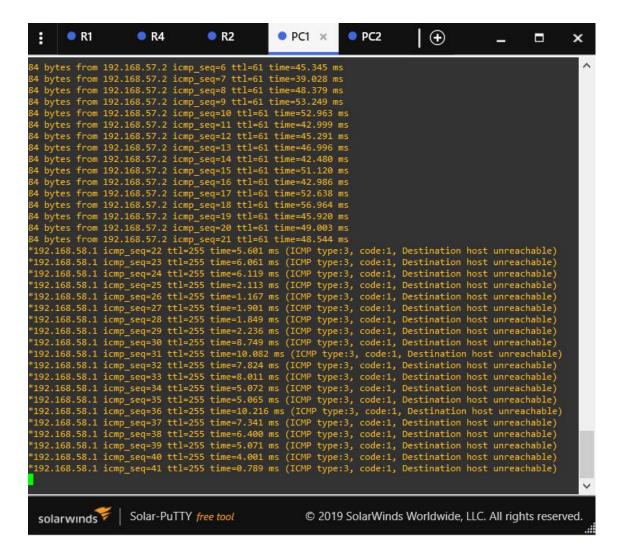


Figure 7. Shut down of the only active route

```
PCL) ping 192.168.57.2 icmp_seq=1 timeout
192.168.57.2 icmp_seq=2 timeout
192.168.57.2 icmp_seq=3 timeout
192.168.57.2 icmp_seq=3 timeout
192.168.57.2 icmp_seq=3 timeout
192.168.57.2 icmp_seq=3 timeout
192.168.57.2 icmp_seq=4 timeout
192.168.57.2 icmp_seq=5 timeout
192.168.57.2 icmp_seq=6 timeout
192.168.57.2 icmp_seq=6 timeout
192.168.57.2 icmp_seq=6 timeout
192.168.57.2 icmp_seq=6 timeout
192.168.57.2 icmp_seq=1 timeout
192.168.57.2 icmp_seq=2 timeout
192.168.57.2 icmp_seq=3 timeout
192.168.57.2 icmp_se
```

Figure 8. Shut down of every single route 1 by 1

```
PC1> trace 192.168.57.2 trace to 192.168.57.2, 8 hops max, press Ctrl+C to stop

1 192.168.58.1 9.374 ms 10.578 ms 9.921 ms

2 *192.168.56.2 25.569 ms 15.650 ms

3 *192.168.56.6 42.328 ms 52.295 ms

4 * * *

5 *192.168.57.2 49.167 ms (ICMP type:3, code:3, Destination port unreachable)

PC1> trace 192.168.57.2 trace to 192.168.57.2 trace to 192.168.57.2 shops max, press Ctrl+C to stop

1 192.168.58.1 4.357 ms 9.909 ms 9.098 ms

2 *192.168.56.18 25.067 ms 32.434 ms

3 *192.168.57.2 39.567 ms (ICMP type:3, code:3, Destination port unreachable)

PC1> trace 192.168.57.2 trace to 192.168.57.2 trace to 192.168.57.2, 8 hops max, press Ctrl+C to stop

1 192.168.58.1 9.581 ms 9.872 ms 9.315 ms

2 *192.168.56.13 12.295 ms 19.729 ms

3 192.168.56.9 18.562 ms 31.158 ms 29.296 ms

4 * * *

5 *192.168.57.2 36.867 ms (ICMP type:3, code:3, Destination port unreachable)
```

Figure 9. Tracing VPC2 from VPC1

```
ip route [ip] [mask] [router_interface] [metric]

Ip - Ip address of the router you want to connect
Mask - Subnet mask of the correspondent Ip
Router_interface - Interface where the router is connected to (ex. f0/0, f1/1..)
Metric - Number that decides priority
```

```
PC1> trace 192.168.57.2
trace to 192.168.57.2, 8 hops max, press Ctrl+C to stop
1 192.168.58.1 9.524 ms 9.194 ms 9.691 ms
2 192.168.56.18 19.360 ms 19.563 ms 29.787 ms
3 * * *
4 *192.168.57.2 37.619 ms (ICMP type:3, code:3, Destination port unreachable)
```

Figure 10. Tracing VPC1 to VPC2 (before)

```
## A bytes from 192.168.57.2 icmp_seq=25 ttl=62 time=39.506 ms

44 bytes from 192.168.57.2 icmp_seq=26 ttl=62 time=33.375 ms

45 bytes from 192.168.57.2 icmp_seq=27 ttl=62 time=33.375 ms

46 bytes from 192.168.57.2 icmp_seq=28 ttl=62 time=31.049 ms

46 bytes from 192.168.57.2 icmp_seq=28 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=28 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=31 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=32 ttl=62 time=34.299 ms

48 bytes from 192.168.57.2 icmp_seq=33 ttl=62 time=35.986 ms

48 bytes from 192.168.57.2 icmp_seq=33 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=33 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=35 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=35 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=31.049 ms

48 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=31.049 ms

49 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=31.049 ms

49 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=31.049 ms

40 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=38.933 ms

40 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=38.933 ms

40 bytes from 192.168.57.2 icmp_seq=36 ttl=62 time=38.933 ms

40 bytes from 192.168.57.2 icmp_seq=37 ttl=62 time=38.933 ms

40 bytes from 192.168.57.2 icmp_seq=41 ttl=25 time=1.049 ms

40 bytes from 192.168.57.2 icmp_seq=41 ttl=25 time=1.049 ms

40 bytes from 192.168.57.2 icmp_seq=45 ttl=62 time=4.049 ms

40 bytes from 192.168.57.2 icmp_seq=45 ttl=62 time=4.049 ms

40 bytes from 192.168.57.2 icmp_seq=55 timeout

40 bytes from 192.168.57.2 icmp_seq=55 timeout

40 bytes from 192.168.57.2 icmp_seq=55 timeout

40 bytes from 192.168.57.2 icmp_seq=56 ttl=61 time=41.168 ms

40 bytes from 192.168.57.2 icmp_seq=56 ttl=61 time=42.838 ms

40 bytes from 192.168.57.2 icmp_seq=65 ttl=61 time=42.838 ms

40 bytes from 192.168.57.2 icmp_seq=65 ttl=61 time=42.838 ms

40 bytes from 192.168.57.2 icmp_seq=65 ttl=
```

Figure 11. Shut down of continuous ping between VPC1 and VPC2

```
PC1> trace 192.168.57.2

trace to 192.168.57.2, 8 hops max, press Ctrl+C to stop

1 192.168.58.1 10.064 ms 8.955 ms 9.056 ms

2 192.168.56.13 30.411 ms 30.773 ms 30.103 ms

3 192.168.56.9 29.292 ms 29.241 ms 19.091 ms

4 * * *

5 *192.168.57.2 28.215 ms (ICMP type:3, code:3, Destination port unreachable)
```

Figure 12. Tracing VPC1 to VPC2 (after)

RIPv2 is a Distance Vector Routing Protocol and it uses hops as a mesure to know the router path, any number above 15 hops its considered unreachable.

Steps:

- Ping VPC1 VPC2
- Shutdown R1 interface f0/1
- Shutdown R4 interface f0/0

This will make a significant packet lost due to the timeout to change paths. Afte the timeout it will find a new path to reach VPC2.

```
PCI> trace 192.168.57.2 8 hops max, press Ctrl+C to stop

1 192.168.56.1 2.707 ms 22.556 ms 8.008 ms

2 192.168.56.2 30.070 ms 19.607 ms 29.452 ms

3 192.168.56.6 30.793 ms 30.324 ms 30.306 ms

4 *192.168.57.2 41.653 ms (ICMP type:3, code:3, Destination port unreachable)

PCI> ping 192.168.57.2 icmp_seq=1 ttl=61 time=38.911 ms

84 bytes from 192.168.57.2 icmp_seq=2 ttl=61 time=42.533 ms

84 bytes from 192.168.57.2 icmp_seq=3 ttl=61 time=40.952 ms

84 bytes from 192.168.57.2 icmp_seq=3 ttl=61 time=44.449 ms

*192.168.58.1 icmp_seq=5 ttl=255 time=24.529 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=5 ttl=255 time=6.319 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=5 ttl=255 time=3.014 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=8 ttl=255 time=3.014 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=8 ttl=255 time=3.014 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=8 ttl=255 time=4.476 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=8 ttl=255 time=4.476 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=8 ttl=255 time=3.014 ms (ICMP type:3, code:1, Destination host unreachable)

*192.168.58.1 icmp_seq=8 ttl=255 time=3.014 ms (ICMP type:3, code:1, Destination host unreachable)

*4 bytes from 192.168.57.2 icmp_seq=11 ttl=62 time=33.591 ms

84 bytes from 192.168.57.2 icmp_seq=12 ttl=62 time=37.591 ms

84 bytes from 192.168.57.2 icmp_seq=13 ttl=62 time=32.908 ms

*C

PCI> trace 192.168.57.2 8 hops max, press Ctrl+C to stop

1 192.168.58.1 9.580 ms 8.976 ms 9.830 ms

2 192.168.56.18 29.811 ms 19.160 ms 19.129 ms

3 *192.168.57.2 29.280 ms (ICMP type:3, code:3, Destination port unreachable)

PCI>
```

Figure 13.Shutdown R1 and R4 interface

OSPF is a link-state routing protocol that sends a packet from the router to discover and memorize all the designed neighbour routers, that first packet contains the list with all the neighbours as well as the link state data. Also for instace, R1 acts as an ABR (Area Border Router) that means that this router is gonna have the task to forward all the routing information through all the next areas. OSPF uses the bandwidth as a mesurement to calculate the cost of a certain path what it makes easier to know what path is the fastest.

5 Problem 5

Conclusion

Static

Its based on given configuration what makes it easier to setup while being a bit time consuming.

In case of failure, the package lost is very high (\sim =70%)

It is more efficient if used on a small network

RIPv2

Its based on distance vector that uses the Distance Vectoring Protocol what makes it take the shortest routing path but not the fastest, despite of that it is easier and faster to setup. The ideal scenario for this method is when all the networks have the same speed but different paths where it can mange really well all the path priorities.

In case of failure, the package lost is average (\sim =40%)

It is more efficient if used on a small network which are not very dynamic and have less than 15 hops, after 15 hops will show that the network was unreachable.

OSPF

Its based on link-state what makes it a bit harder but can be manged farelly quick. This method calculates the cost of reaching from one destination to another making it choosing the fastest path to a specific destination what makes it so definitly the most complex out of the three methods used.

In case of failure, the package lost is really low what makes it really reliable (\sim =10%)

It is the most efficient to use on both small and large networks.