

Topological Data Analysis

2022–2023

Solutions of Exercises

24 November 2022

① Suppose that $f: (V, \pi) \rightarrow (V', \pi')$ is an isomorphism. Then there is a morphism $g: (V', \pi') \rightarrow (V, \pi)$ such that $g \circ f = \text{id}_V$ and $f \circ g = \text{id}_{V'}$. Therefore, for all $t \in \mathbb{R}$,

$$\begin{cases} g_t \circ f_t = (g \circ f)_t = (\text{id}_V)_t = \text{id}_{V_t} \\ f_t \circ g_t = (f \circ g)_t = (\text{id}_{V'})_t = \text{id}_{V'_t} \end{cases}$$

This implies that $f_t: V_t \rightarrow V'_t$ and $g_t: V'_t \rightarrow V_t$ are mutually inverse isomorphisms.

Conversely, suppose that $f_t: V_t \rightarrow V'_t$ is an isomorphism for all t .

Pick an inverse $g_t: V'_t \rightarrow V_t$ for each t . We need to prove that

$\{g_t\}_{t \in \mathbb{R}}$ is a morphism of persistence modules. Indeed, since f is a morphism, for all $s \leq t$ we have $f_t \circ \pi_{s,t} = \pi'_{s,t} \circ f_s$ and hence

$$g_t \circ f_t \circ \pi_{s,t} = g_t \circ \pi'_{s,t} \circ f_s; \quad \pi_{s,t} = g_t \circ \pi'_{s,t} \circ f_s;$$

$$\pi_{s,t} \circ g_s = g_t \circ \pi'_{s,t} \circ f_s \circ g_s; \quad \pi_{s,t} \circ g_s = g_t \circ \pi'_{s,t} \quad \text{so } g \text{ is a morphism.}$$

Moreover, $g \circ f = \text{id}_V$ since $g_t \circ f_t = \text{id}_{V_t} \forall t$, and likewise $f \circ g = \text{id}_{V'}$.

② Let (V, π) and (V', π') be persistence modules of finite type and suppose that $f: V \rightarrow V'$ is an isomorphism.

Let A and A' be the respective spectra.

Let $a \in A$. We need to prove that $a \in A'$. Suppose the contrary.

If $a \notin A'$, then there is a $\delta > 0$ such that $\pi'_{s,t}$ is an isomorphism for $a - \delta < s \leq t < a + \delta$. Since $a \in A$, there is an $\varepsilon > 0$ such that $\pi_{s,a}$ is not an isomorphism for $a - \varepsilon < s < a$.

Now we have, for $\lambda = \min\{\delta, \varepsilon\}$ and $a - \lambda < s < a$, that

$$\begin{array}{ccc} V_s & \xrightarrow{\pi_{s,a}} & V_a \\ f_s \downarrow \cong & & \cong \downarrow f_a \\ V'_s & \xrightarrow[\pi'_{s,a}]{\cong} & V'_a \end{array}$$

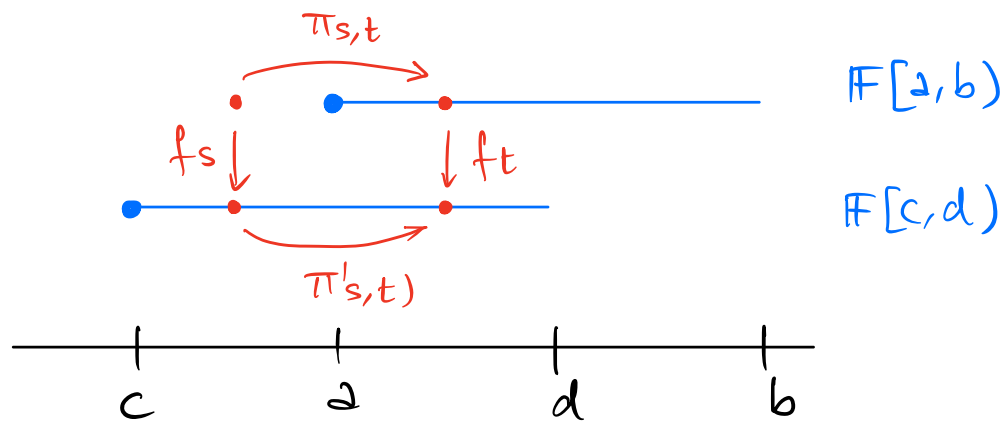
$$f_a \circ \pi_{s,a} = \pi'_{s,a} \circ f_s$$

where f_a, f_s and $\pi'_{s,a}$ are isomorphisms while $\pi_{s,a}$ is not.

This is a contradiction.

This argument shows that $A \subseteq A'$, and $A' \subseteq A$ by symmetry.

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- If $c \leq a < d \leq b$, we define $f: F[a, b) \rightarrow F[c, d)$ as follows:

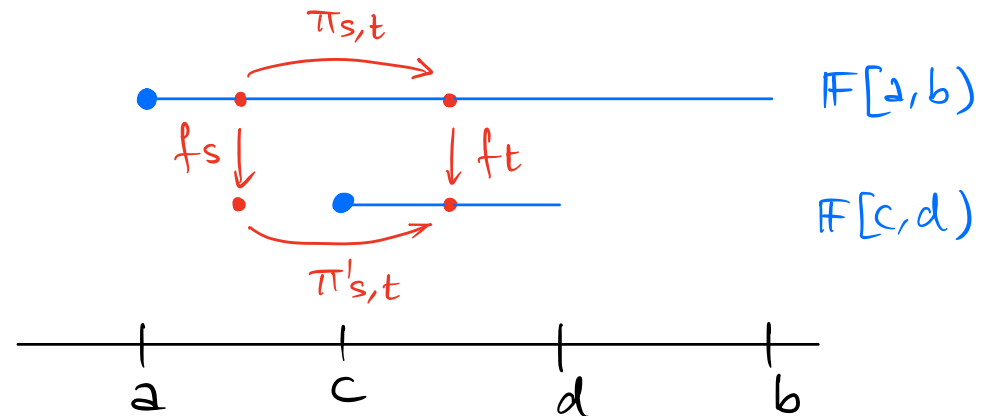
$$f_t = \begin{cases} \text{id}_F & \text{if } a \leq t < d \\ 0 & \text{otherwise.} \end{cases}$$

Then $f_t \circ \pi_{s,t} = \pi'_{s,t} \circ f_s \quad \forall s \leq t$.

- If $a < c$, then for $a \leq s < c$ and $c \leq t < d$ we have either

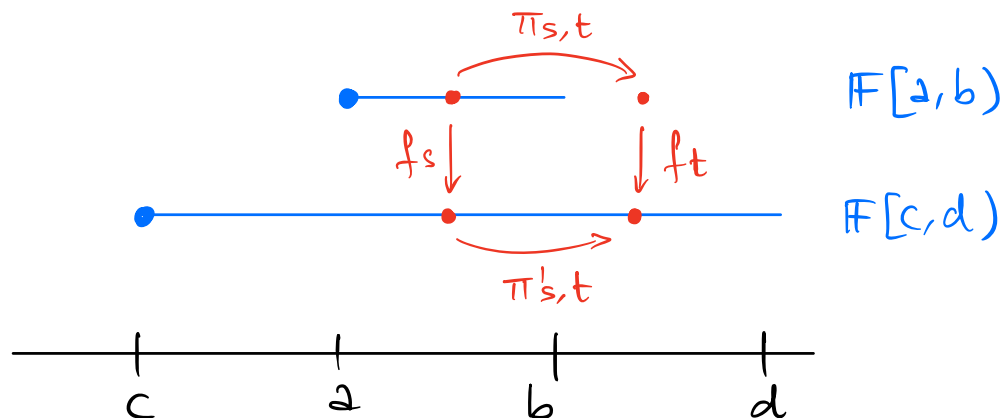
$$\begin{array}{ccc} F & \xrightarrow{\sim} & F \quad \text{or} \quad F \longrightarrow 0 \\ f_s \downarrow & & \downarrow f_t \\ 0 & \longrightarrow & F \\ \text{if } t < b & & \text{if } t \geq b, \end{array}$$

which forces $f_t = 0$ for $c \leq t < d$ and hence $f = 0$.



- If $c \leq a$ and $b < d$, then for $a \leq s < b$ and $b \leq t < d$ we have

$$\begin{array}{ccc}
 \mathbb{F} & \longrightarrow & 0 \\
 f_s \downarrow & & \downarrow f_t \\
 \mathbb{F} & \xrightarrow{\cong} & \mathbb{F}
 \end{array}$$



which forces $f_s = 0$ for $a \leq s < b$ and therefore $f = 0$.

- If $d \leq a$, then necessarily $f_t = 0$ for all t . Hence $f = 0$ is the only possibility.

