

# *Orchestrating Decision Engines*

## University of Verona, 2011

Leonardo de Moura and Grant Passmore



# Satisfiability Modulo Theories (SMT)

A Satisfiability Checker  
with built-in support for useful theories

# Satisfiability Modulo Theories (SMT)

$b + 2 = c$  and  $f(\text{read}(\text{write}(a,b,3)), c-2) \neq f(c-b+1)$

# Satisfiability Modulo Theories (SMT)

$b + 2 = c$  and  $f(\text{read}(\text{write}(a,b,3), c-2)) \neq f(c-b+1)$

Arithmetic

# Satisfiability Modulo Theories (SMT)

$b + 2 = c$  and  $f(\text{read}(\text{write}(a,b,3), c-2)) \neq f(c-b+1)$

Array Theory

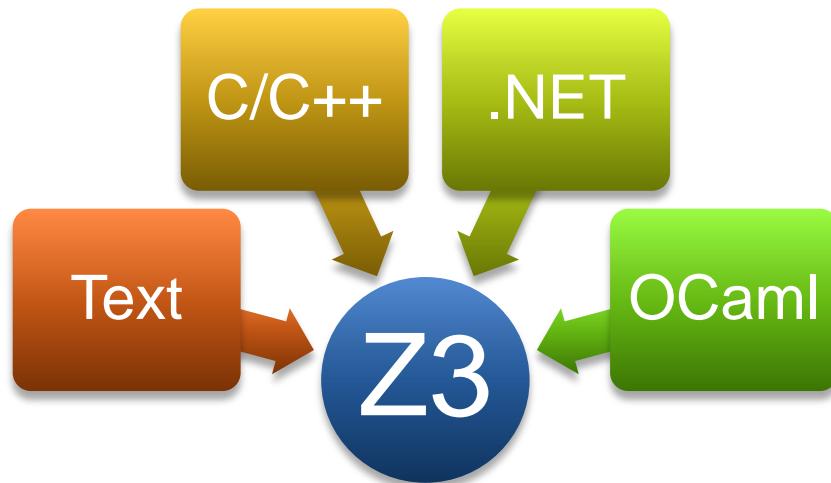
# Satisfiability Modulo Theories (SMT)

$b + 2 = c$  and  $f(\text{read}(\text{write}(a,b,3), c-2)) \neq f(c-b+1)$

Uninterpreted  
Functions

# SMT@Microsoft: Solver

- Z3 is a solver developed at Microsoft Research.
- Development/Research driven by internal customers.
- Free for non-commercial use.
- Interfaces:



- <http://research.microsoft.com/projects/z3>

# Some Applications @ Microsoft



The  
Spec#  
Programming System

HAVOC



Hyper-V

Microsoft®

Virtualization

Terminator T-2

VCC

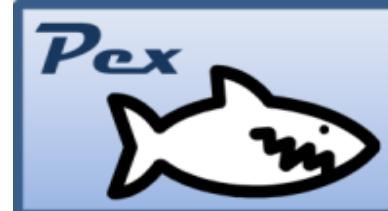
SLAM

NModel

Yogi

Vigilante

SpecExplorer

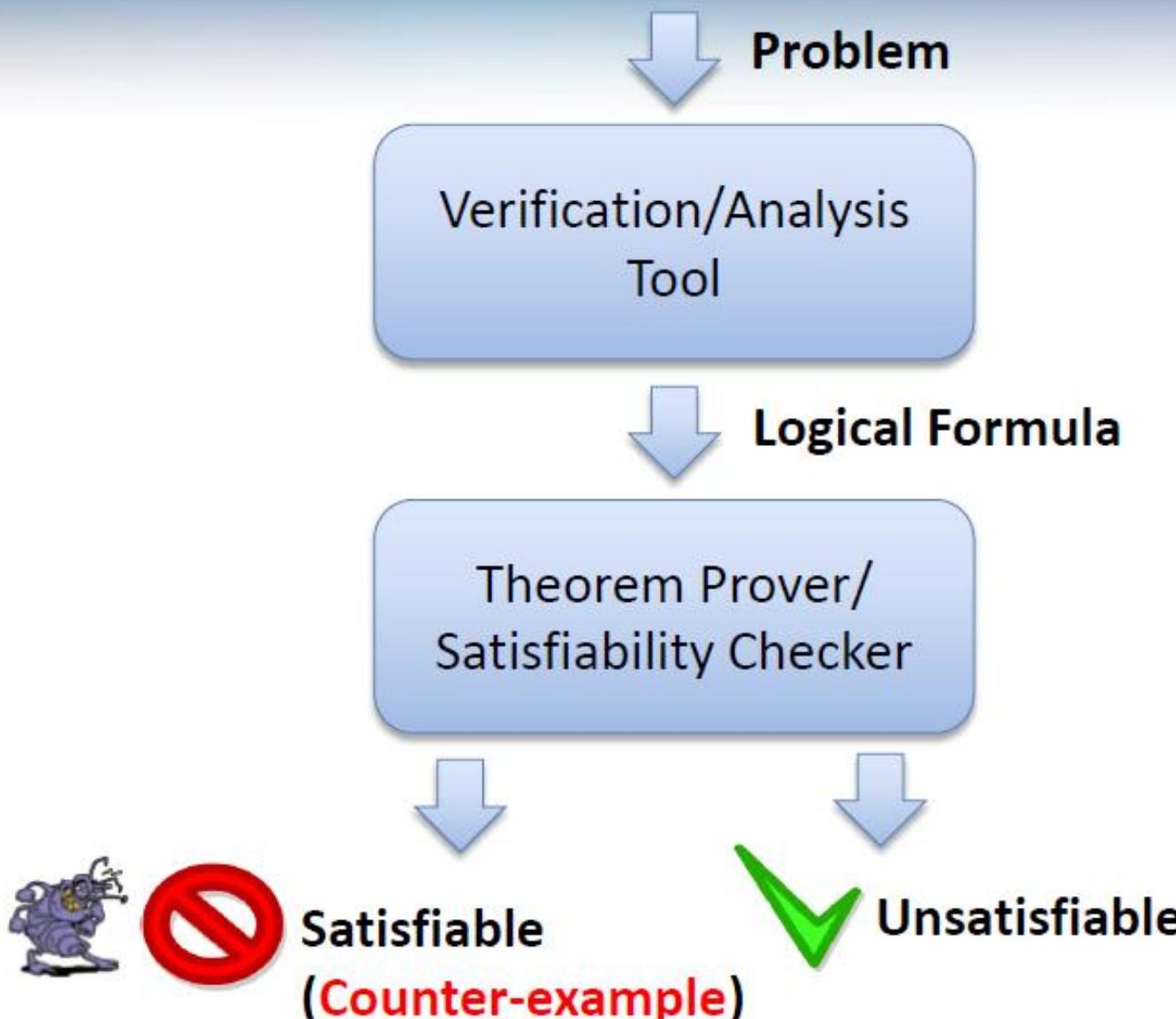


SAGE

F7

Microsoft®  
Research

# Verification/Analysis Tool: “Template”

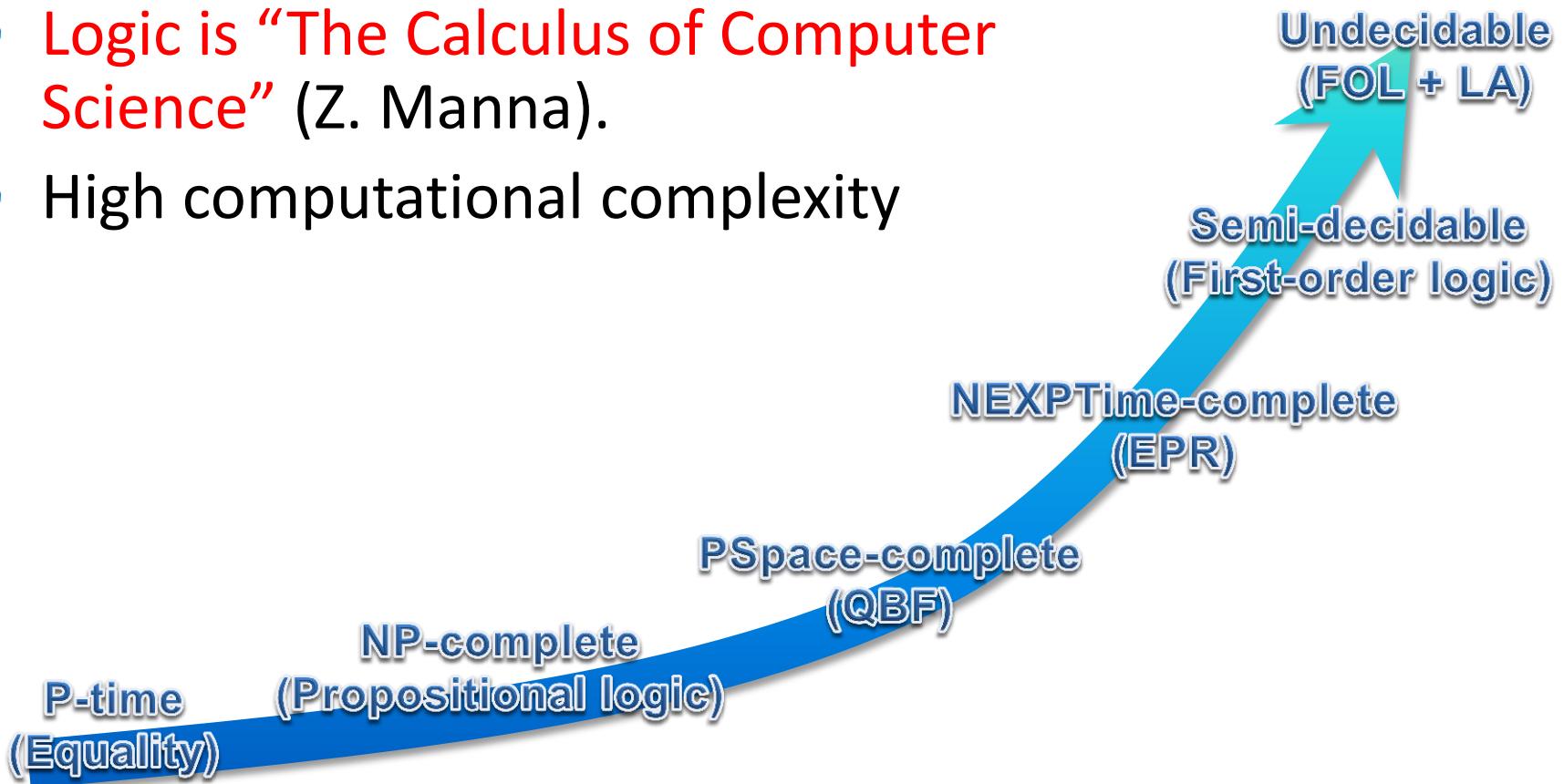


# Symbolic Reasoning

Verification/Analysis tools  
need some form of  
**Symbolic Reasoning**

# Symbolic Reasoning

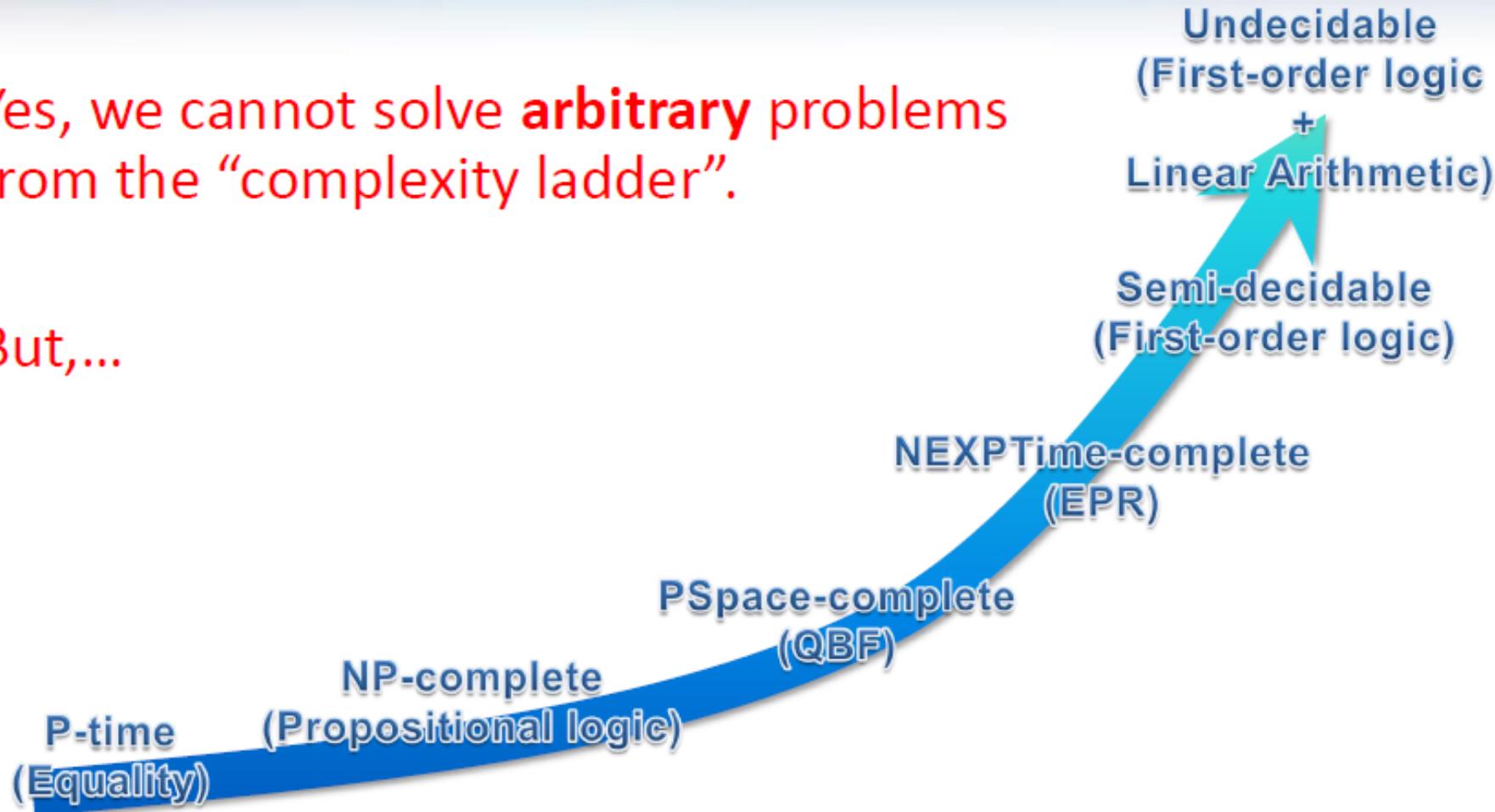
- Logic is “The Calculus of Computer Science” (Z. Manna).
- High computational complexity



# Symbolic Reasoning

Yes, we cannot solve **arbitrary** problems from the “complexity ladder”.

But,...



# Symbolic Reasoning

We can try to solve the  
problems we find in  
**real applications**

# Applications

Test case generation

Verifying Compilers

Predicate Abstraction

Invariant Generation

Type Checking

Model Based Testing

# Test-case generation

- Test (correctness + usability) is 95% of the deal:
  - Dev/Test is 1-1 in products.
  - Developers are responsible for unit tests.
- Tools:
  - Annotations and static analysis (SAL + ESP)
  - File Fuzzing
  - Unit test case generation

# Security is critical

- Security bugs can be very expensive:
  - Cost of each MS Security Bulletin: \$600k to \$Millions.
  - Cost due to worms: \$Billions.
  - The real victim is the customer.
- Most security exploits are initiated via files or packets.
  - Ex: Internet Explorer parses dozens of file formats.
- Security testing: hunting for million dollar bugs
  - Write A/V
  - Read A/V
  - Null pointer dereference
  - Division by zero

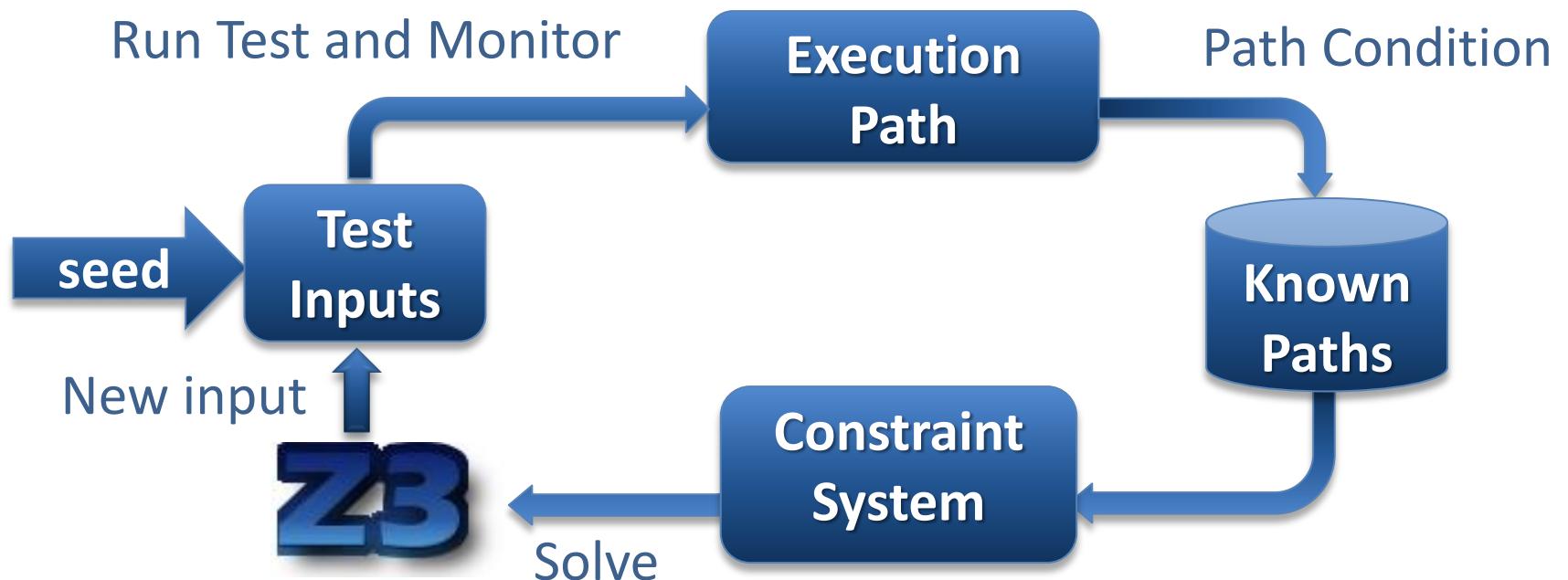


# Hunting for Security Bugs.

- Two main techniques used by “*black hats*”:
  - Code inspection (of binaries).
  - *Black box fuzz testing*.
- **Black box** fuzz testing:
  - A form of black box random testing.
  - Randomly *fuzz* (=modify) a well formed input.
  - Grammar-based fuzzing: rules to encode how to fuzz.
- **Heavily** used in security testing
  - At MS: several internal tools.
  - Conceptually simple yet effective in practice



# Directed Automated Random Testing ( DART)



# DARTish projects at Microsoft

PEX

Implements DART for .NET.

SAGE

Implements DART for x86 binaries.

YOGI

Implements DART to check the feasibility of program paths generated statically using a SLAM-like tool.

Vigilante

Partially implements DART to dynamically generate worm filters.

# What is *Pex*?

- Test input generator
  - Pex starts from parameterized unit tests
  - Generated tests are emitted as traditional unit tests

# ArrayList: The Spec

The screenshot shows two versions of the MSDN .NET Framework Developer Center page for the `ArrayList.Add` method.

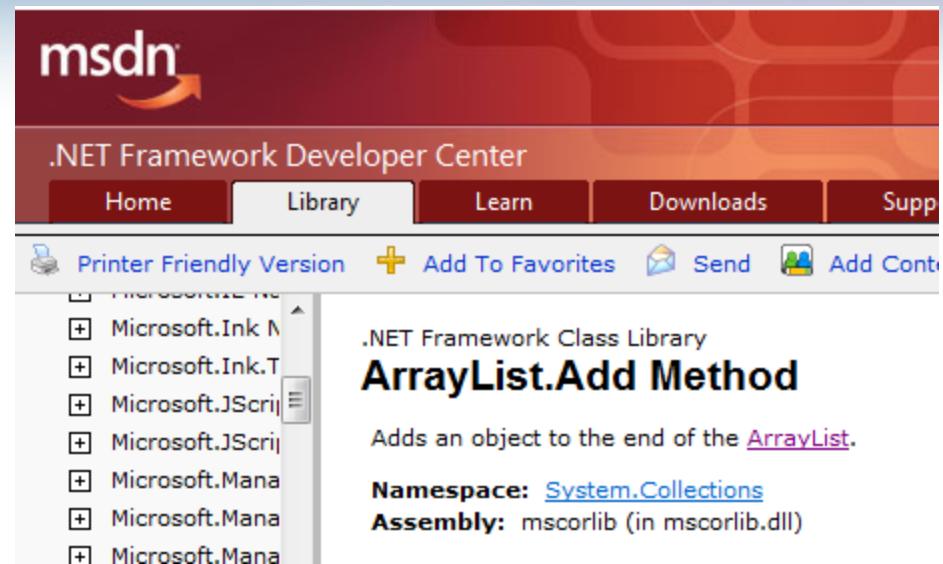
**Top Version:** This is a smaller screenshot of the page. It features the MSDN logo at the top left, followed by a navigation bar with links for Home, Library, Learn, Downloads, and Support. Below the navigation bar is a toolbar with links for Printer Friendly Version, Add To Favorites, Send, and Add Content. The main content area is titled ".NET Framework Class Library" and "ArrayList.Add Method". It describes the method as adding an object to the end of the `ArrayList`. The **Namespace:** is listed as `System.Collections` and the **Assembly:** as `mscorlib (in mscorlib.dll)`.

**Bottom Version:** This is a larger screenshot of the same page. It includes the same header and toolbar. The main content area is titled ".NET Framework Class Library" and "ArrayList.Add Method". It contains a sidebar with a tree view of related classes under the "Microsoft" namespace, including `Microsoft.Ink`, `Microsoft.Ink.T`, `Microsoft.JScri`, `Microsoft.JScri`, `Microsoft.Mana`, `Microsoft.Mana`, and `Microsoft.Mana`. The main content area also includes a "Remarks" section with a note about accepting null references and allowing duplicates, and a section about capacity and performance.

# ArrayList: AddItem Test

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```



The screenshot shows the MSDN .NET Framework Developer Center. The top navigation bar includes Home, Library (which is selected), Learn, Downloads, and Support. Below the navigation is a toolbar with links for Printer Friendly Version, Add To Favorites, Send, and Add Content. The main content area displays the documentation for the **ArrayList.Add Method**. It describes the method as adding an object to the end of the [ArrayList](#). The **Namespace:** is [System.Collections](#) and the **Assembly:** is [mscorlib \(in mscorlib.dll\)](#). A sidebar on the left lists various Microsoft namespaces.

# ArrayList: Starting Pex...

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
...  
}
```

## Inputs

# ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

Inputs

(0, null)

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
...  
}
```

# ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Inputs	Observed Constraints
(0,null)	$!(c < 0)$

c < 0 → false

# ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item);  
    }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length) 0 == c → true  
        ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

Inputs	Observed Constraints
(0,null)	!(c<0) && 0==c

# ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

Inputs	Observed Constraints
(0,null)	!(c<0) && 0==c

item == item → true

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

# ArrayList: Picking the next branch to cover

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && <b>0!=c</b>		

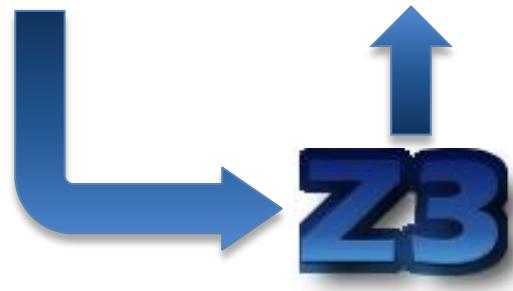


# ArrayList: Solve constraints using SMT solver

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	



# ArrayList: Run 2, (1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length) 0 == c → false  
        ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c

# ArrayList: Pick new branch

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0		

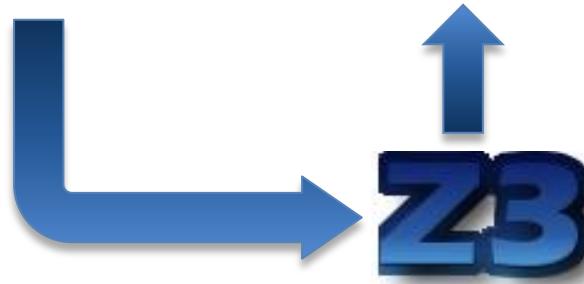


# ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0	(-1,null)	



# ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0	(-1,null)	c<0

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

c < 0 → true

# ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0	(-1,null)	c<0

# ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0	(-1,null)	c<0



Login Join Twitter!

Once again, Pex blows my mind. It's utterly amazing the bugs that it can find :).

*about 7 hours ago from TweetDeck*



jasonbock  
Jason Bock

# White box testing in practice

## How to test this code?

(Real code from .NET base class libraries.)

```
[SecurityPermissionAttribute(SecurityAction.LinkDemand, Flags=SecurityPermissionFlag.SerializationFormatter)]
public ResourceReader(Stream stream)
{
    if (stream==null)
        throw new ArgumentNullException("stream");
    if (!stream.CanRead)
        throw new ArgumentException(Environment.GetResourceString("Argument_StreamNotReadable"));

    _resCache = new Dictionary<String, ResourceLocator>(FastResourceComparer.Default);
    _store = new BinaryReader(stream, Encoding.UTF8);
    // We have a faster code path for reading resource files from an assembly.
    _ums = stream as UnmanagedMemoryStream;

    BCLDebug.Log("RESMGRFILEFORMAT", "ResourceReader .ctor(Stream). UnmanagedMemoryStream: "+(_ums!=null));
    ReadResources();
}
```

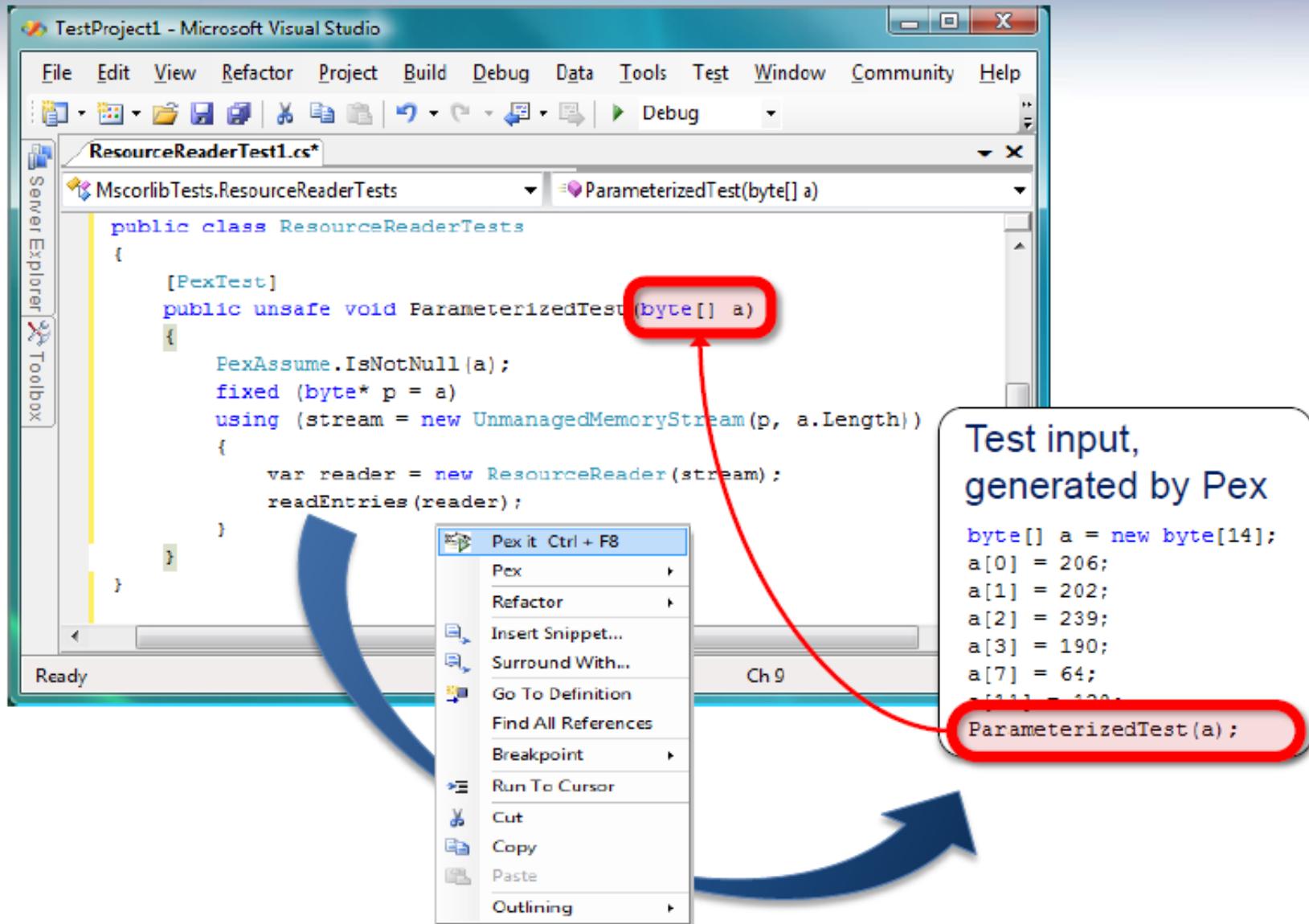
# White box testing in practice

```
// Reads in the header information for a .resources file. Verifies some
// of the assumptions about this resource set, and builds the class table
// for the default resource file format.
private void ReadResources() {
    BCLDebug.Assert(_store != null, "ResourceReader is closed!");
    BinaryFormatter bf = new BinaryFormatter(null, new StreamingContext(StreamingContextStates.File |
#if !FEATURE_PAL
    _typeLimitingBinder = new TypeLimitingDeserializationBinder();
    bf.Binder = _typeLimitingBinder;
#endif
    _objFormatter = bf;
    try {
        // Read ResourceManager header
        // Check for magic number
        int magicNum = _store.ReadInt32();
        if (magicNum != ResourceManager.MagicNumber)
            throw new ArgumentException(Environment.GetResourceString("Resources_StreamNotValid"));
        // Assuming this is ResourceManager header v1 or greater, hopefully,
        // after the version number there is a number of bytes to skip
        // to bypass the rest of the ResMgr header.
        int resMgrHeaderVersion = _store.ReadInt32();
        if (resMgrHeaderVersion > 1) {
            int numBytesToSkip = _store.ReadInt32();
            BCLDebug.Log("RESMGRFILEFORMAT", LogLevel.Status, "ReadResources: Unexpected ResMgr header");
            BCLDebug.Assert(numBytesToSkip >= 0, "numBytesToSkip in ResMgr header should be positive!");
            _store.BaseStream.Seek(numBytesToSkip, SeekOrigin.Current);
        } else {
            BCLDebug.Log("RESMGRFILEFORMAT", "ReadResources: Parsing ResMgr header v1.");
            SkipInt32();      // We don't care about numBytesToSkip.
            // Read in type name for a suitable ResourceReader
            //
```

# White box testing in practice

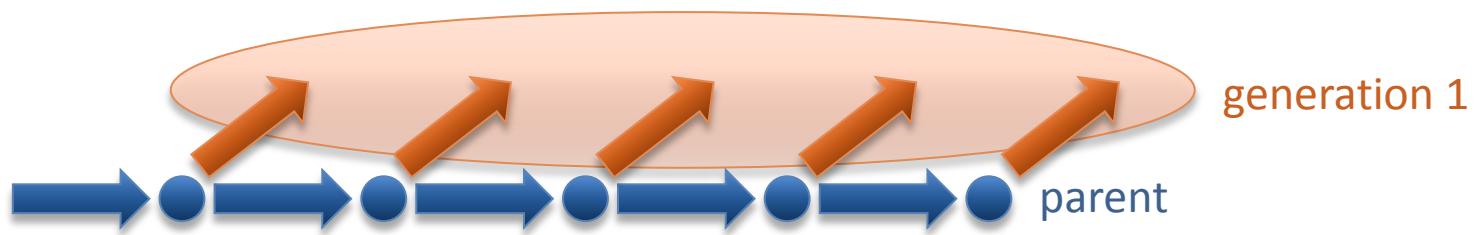
```
// Reads in the header information for a .resources file. Verifies some
// of the assumptions about this resource set, and builds the class table
// for the default resource file format.
private void ReadResources() {
    BCLDebug.Assert(_store != null, "ResourceReader is closed!");
    BinaryFormatter bf = new BinaryFormatter(null, new StreamingContext(StreamingContextStates.File |
#endif !FEATURE_PAL
    _typeLimitingBinder = new TypeLimitingDeserializationBinder();
    bf.Binder = _typeLimitingBinder;
#endif
    _objFormatter = bf;
    try {
        // Read ResourceManager header
        // Check for magic number
        int magicNum = _store.ReadInt32();
        if public virtual int ReadInt32() {
            if (m_isMemoryStream) {
                // read directly from MemoryStream buffer
                MemoryStream mStream = m_stream as MemoryStream;
                BCLDebug.Assert(mStream != null, "m_stream as MemoryStream != null");
                int
                if
                    return mStream.InternalReadInt32();
                }
                else
                {
                    FillBuffer(4);
                }
            }
            return (int)(m_buffer[0] | m_buffer[1] << 8 | m_buffer[2] << 16 | m_buffer[3] << 24);
        }
    }
    // Read in type name for a suitable ResourceReader
    //
```

# Pex – Test Input Generation



# SAGE

- Apply DART to large applications (not units).
- Start with well-formed input (not random).
- Combine with generational search (not DFS).
  - Negate 1-by-1 each constraint in a path constraint.
  - Generate many children for each parent run.



# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; .....
```

Generation 0 – seed file

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 00 00 00 00 00 00 00 00 00 00 00 00 ; RIFF.....  
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000060h: 00 00 00 00 ; ....
```

Generation 1

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 00 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF....***....  
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....  
00000060h: 00 00 00 00
```

Generation 2

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFFE...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; .....
```

Generation 3

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 00 00 00 ; .....strh.....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; .....
```

Generation 4

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 76 69 64 73 ; ....strh... vids
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 5

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 76 69 64 73 ; .....strh.....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 00 00 00 00 ; .....strf.....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; .....
```

Generation 6

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 00 28 00 00 00 ; ....strf....(...
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00
```

Generation 7

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 00 28 00 00 00 ; ....strf....(...
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 C9 9D E4 4E ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 8

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 00 28 00 00 00 ; ....strf....(...
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 9

# Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 B2 75 76 3A 28 00 00 00 ; ....strh^uv:(...
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 10 – CRASH

# SAGE (cont.)

- SAGE is very effective at finding bugs.
- Works on large applications.
- Fully automated
- Easy to deploy (x86 analysis – any language)
- Used in various groups inside Microsoft
- Powered by Z3.

# SAGE $\leftrightarrow$ Z3

- Formulas are usually big conjunctions.
- SAGE uses only the bitvector and array theories.
- Pre-processing step has a huge performance impact.
  - Eliminate variables.
  - Simplify formulas.
- Early unsat detection.

# Application Scenarios

“Big” and hard formulas

Thousands of “small” and easy formulas

Short timeout (< 5secs)

# Application Scenarios

“Big” and hard formulas



**HAVOC**

**VCC**

Thousands of “small” and easy formulas

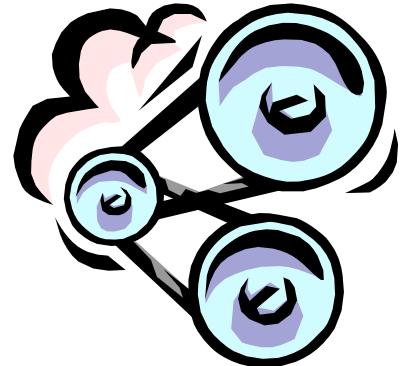
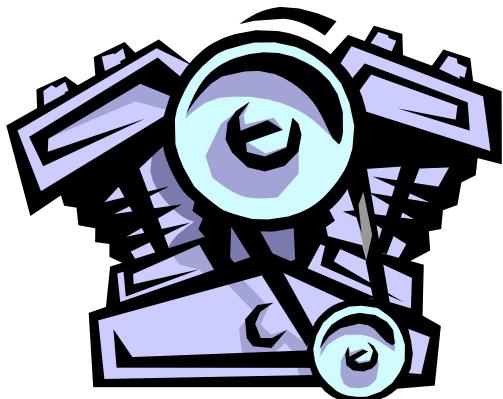


Short timeout (< 5secs)

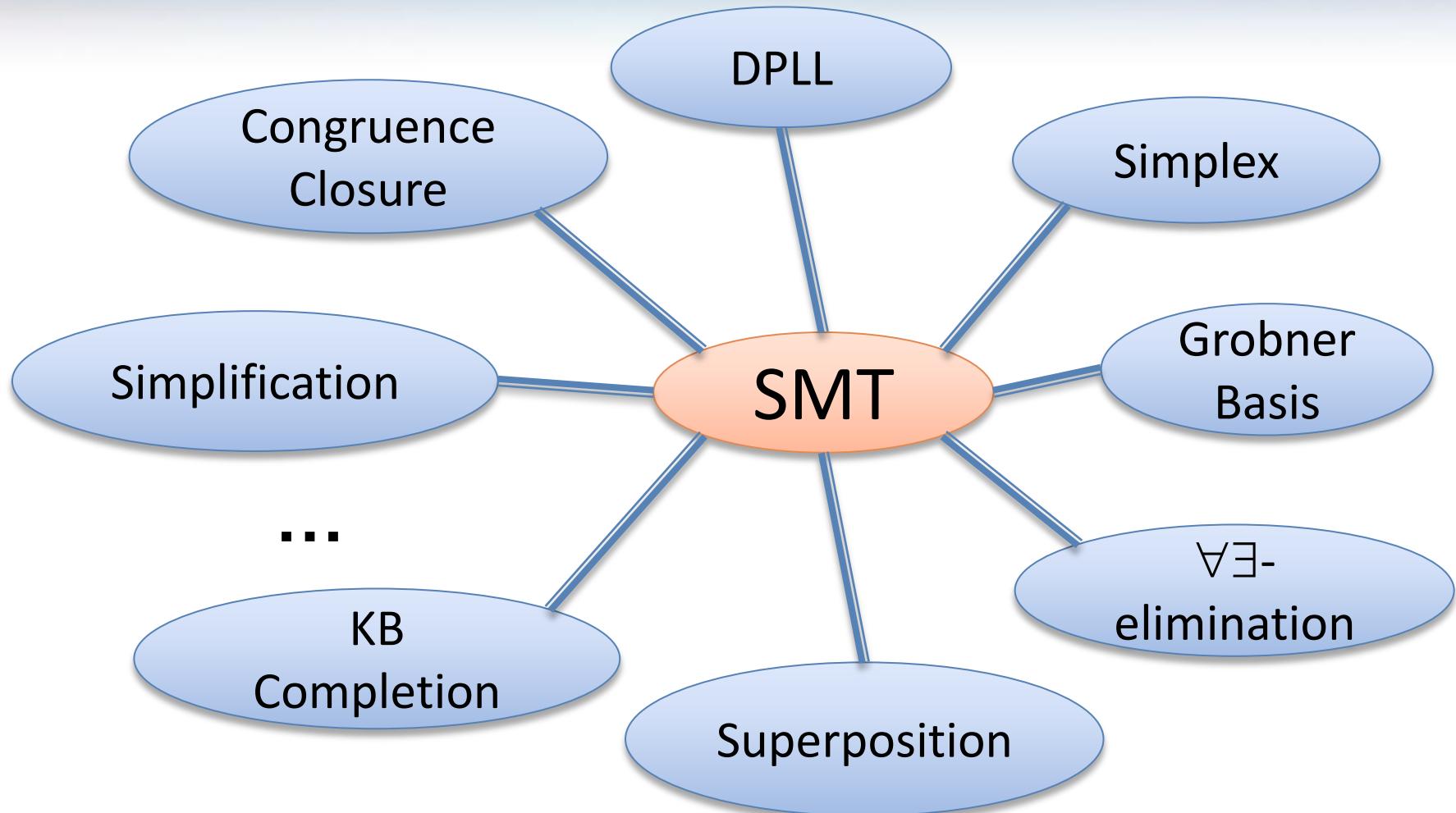
**SAGE**

# Combining Engines

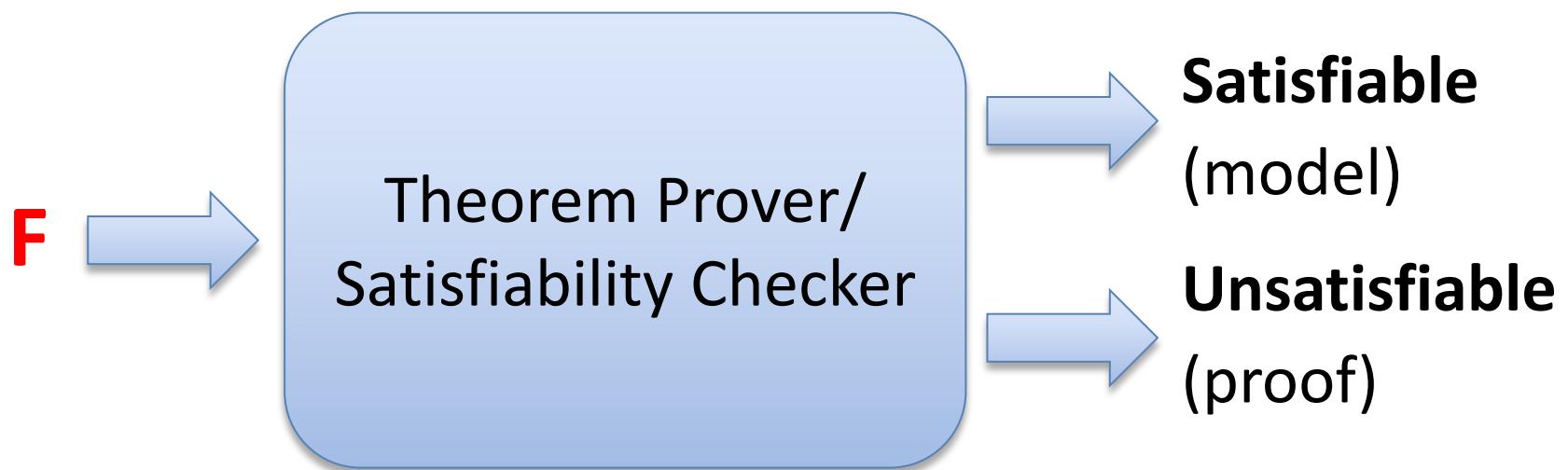
Current SMT solvers provide  
a combination  
of different engines



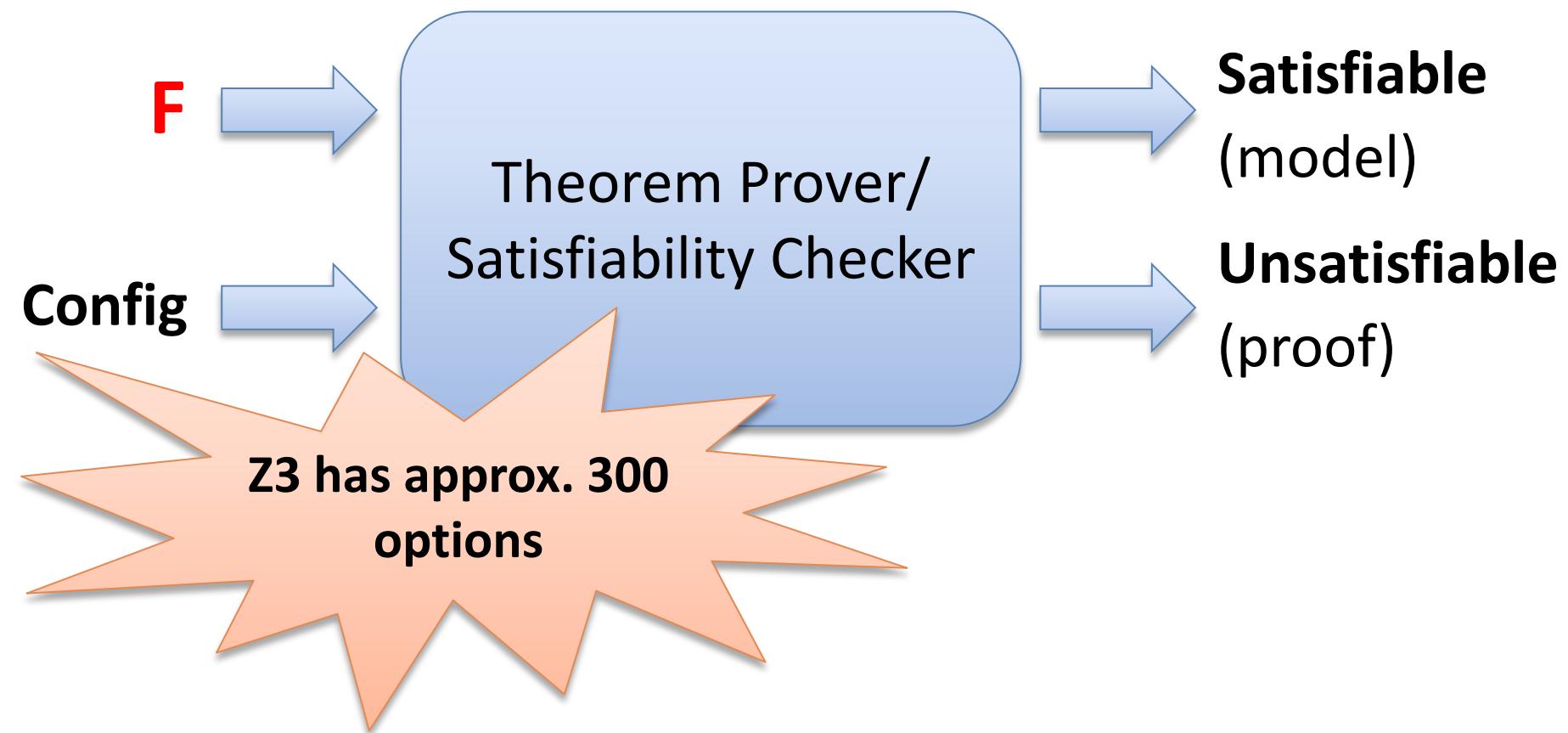
# Combining Engines



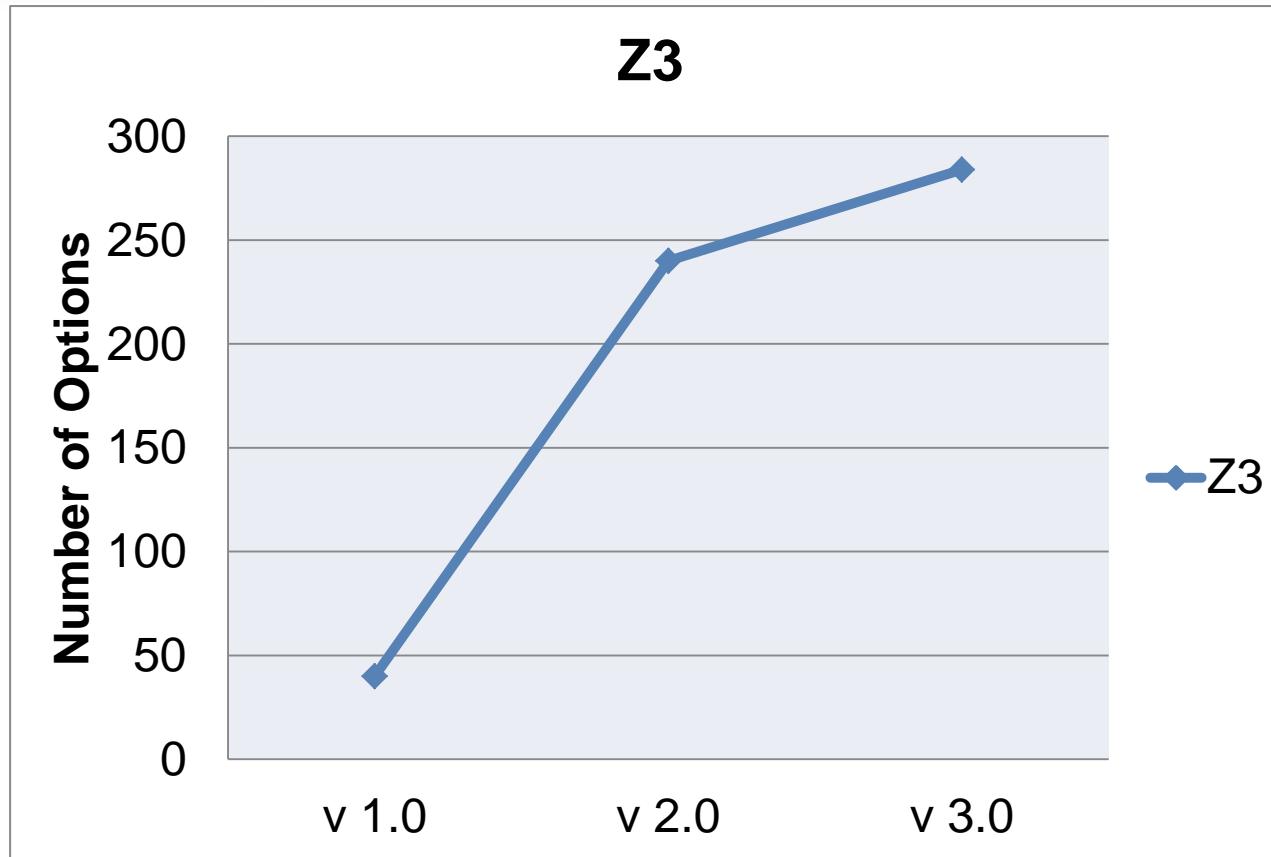
# Theorem Provers & Satisfiability Checkers



# Theorem Provers & Satisfiability Checkers



# Z3 number of options evolution



# Opening the “Black Box”

Actual feedback provided by Z3 users:

“Could you send me your CNF converter?”

“I want to implement my own search strategy.”

“I want to include these rewriting rules in Z3.”

“I want to apply a substitution to term  $t$ .”

“I want to compute the set of implied equalities.”

# The Strategy Challenge

To build theoretical and practical tools  
allowing users to exert strategic control  
over core heuristic aspects of high  
performance SMT solvers.

# What is a strategy?

Theorem proving as an exercise of combinatorial search

Strategies are adaptations of general search mechanisms which reduce the search space by tailoring its exploration to a particular class of formulas.

# The Need for “Strategies”

Different Strategies for Different Domains.

# The Need for “Strategies”

Different Strategies for Different Domains.

From timeout to 0.05 secs...

# Example in Quantified Bit-Vector Logic (QBVF)

Joint work with C. Wintersteiger and Y. Hamadi  
FMCAD 2010

QBVF = Quantifiers + Bit-vectors + uninterpreted functions

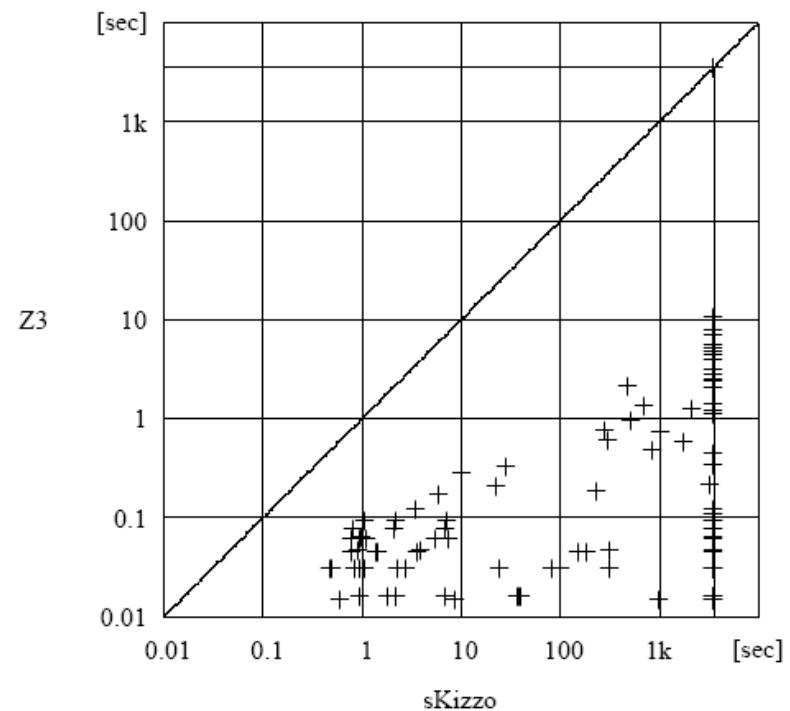
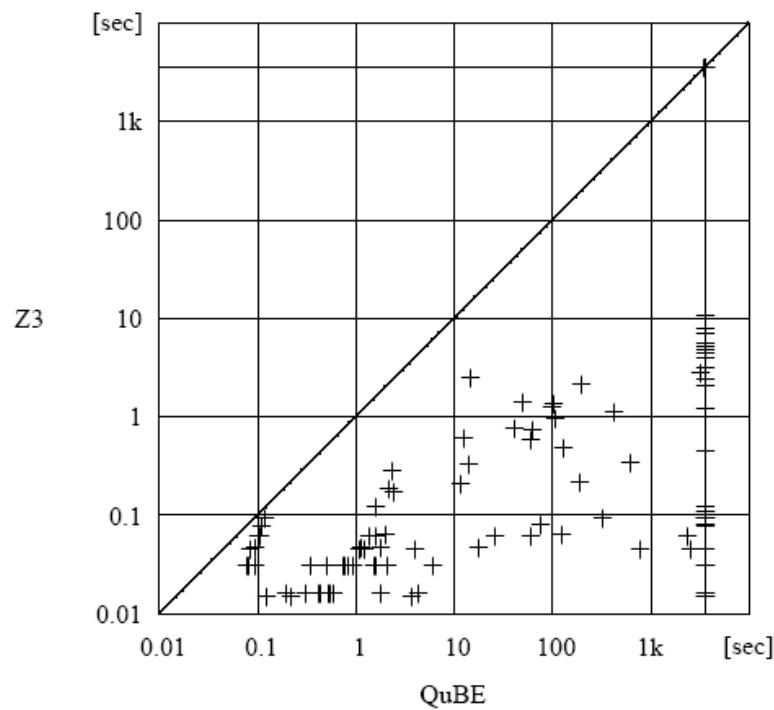
Hardware Fixpoint Checks.

Given:  $I[x]$  and  $T[x, x']$

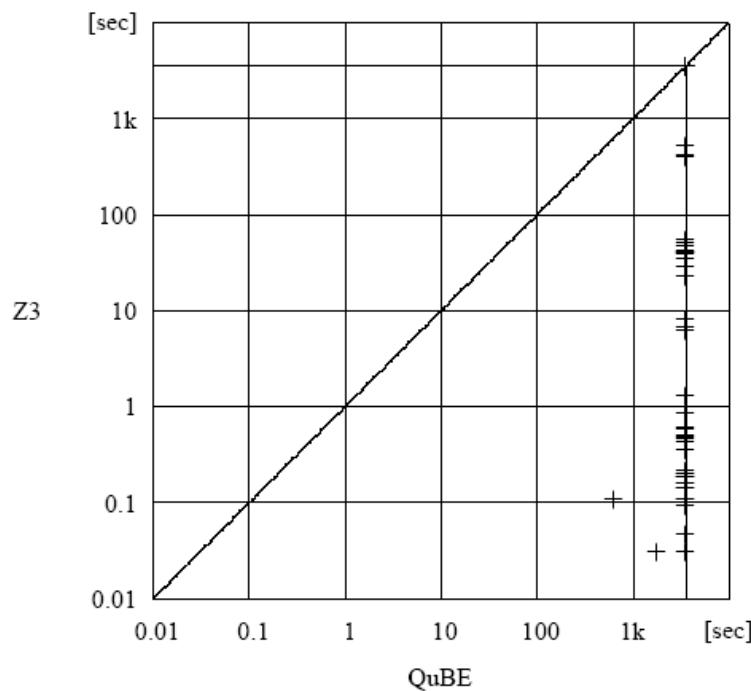
$$\forall x, x' . I[x] \wedge T^k[x, x'] \rightarrow \exists y, y' . I[y] \wedge T^{k-1}[y, y']$$

Ranking function synthesis.

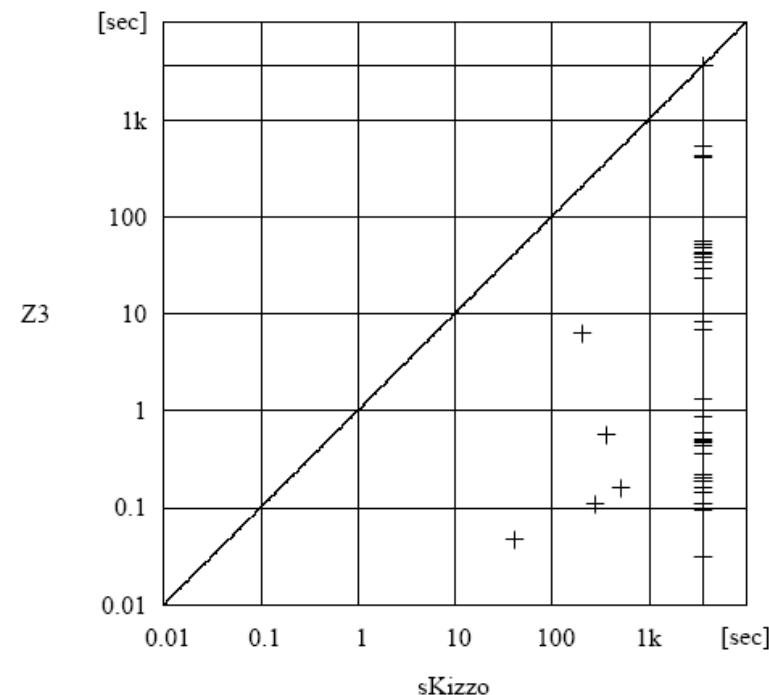
# Hardware Fixpoint Checks



# Ranking Function Synthesis



QuBE



sKizzo

# Why is Z3 so fast in these benchmarks?

Z3 is using different engines:  
rewriting, simplification, model checking, SAT, ...

Z3 is using a customized **strategy**.

We could do it because  
we have access to the source code.

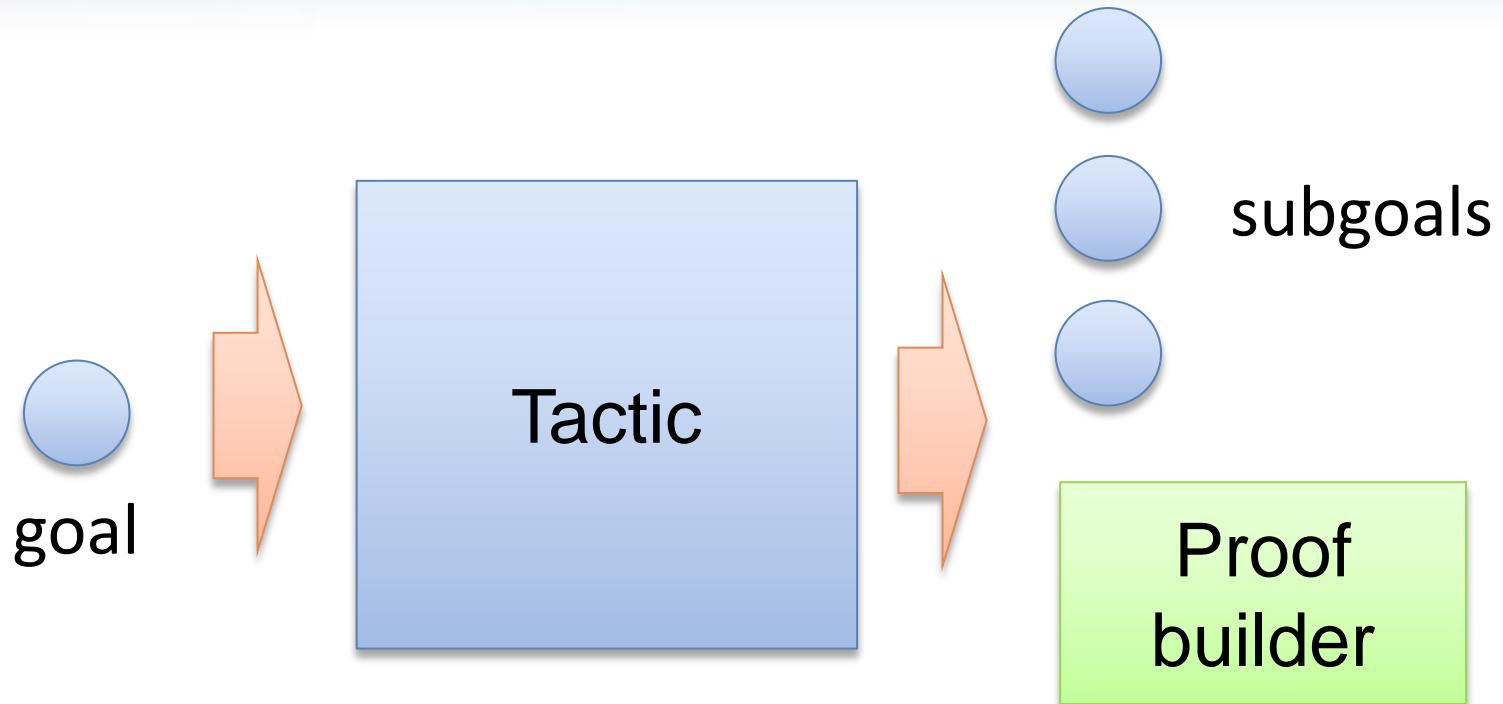
# The "Message"

SMT solvers are collections of little engines.

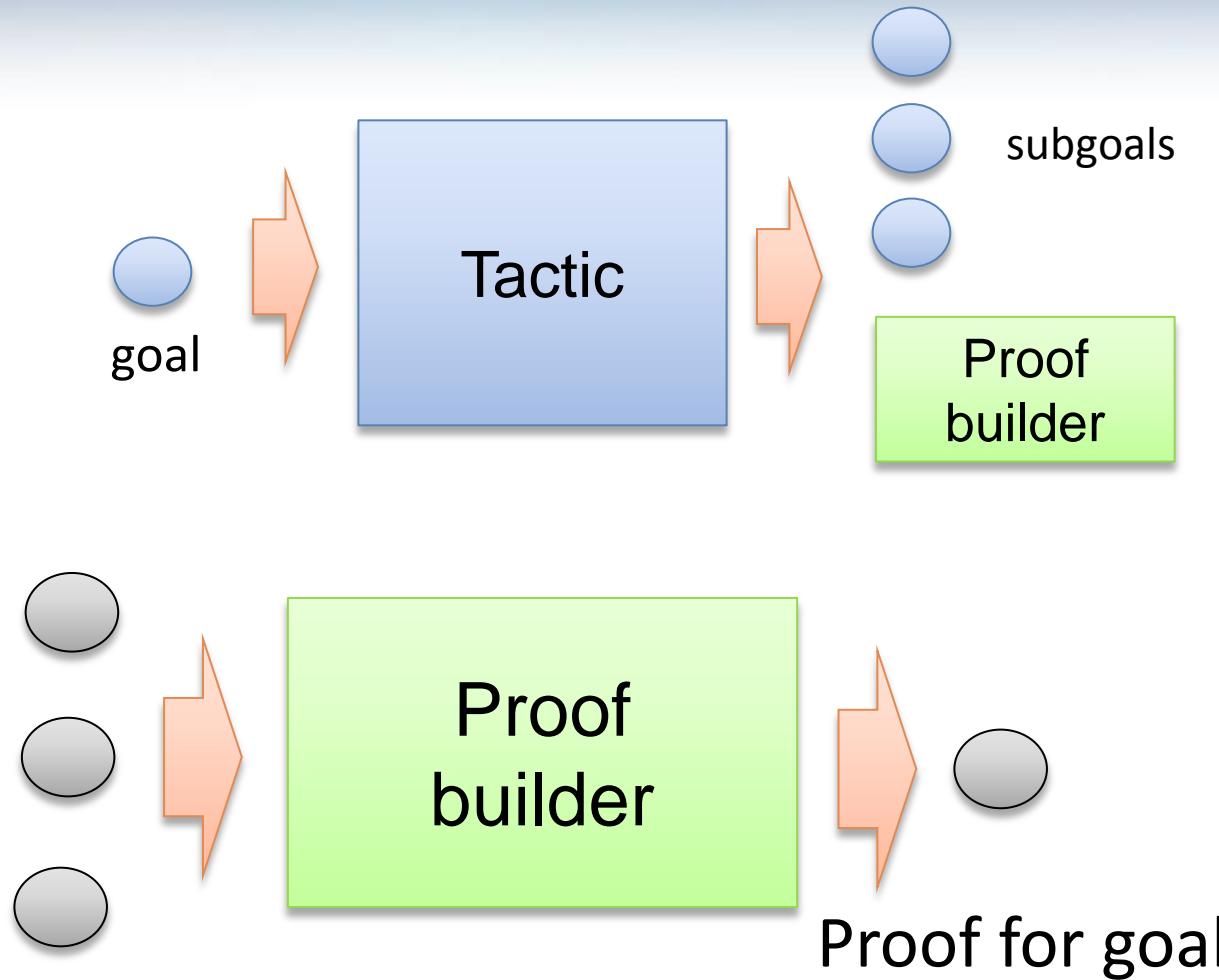
They should provide access to these engines.

Users should be able to define their own strategies.

# Main inspiration: LCF-approach

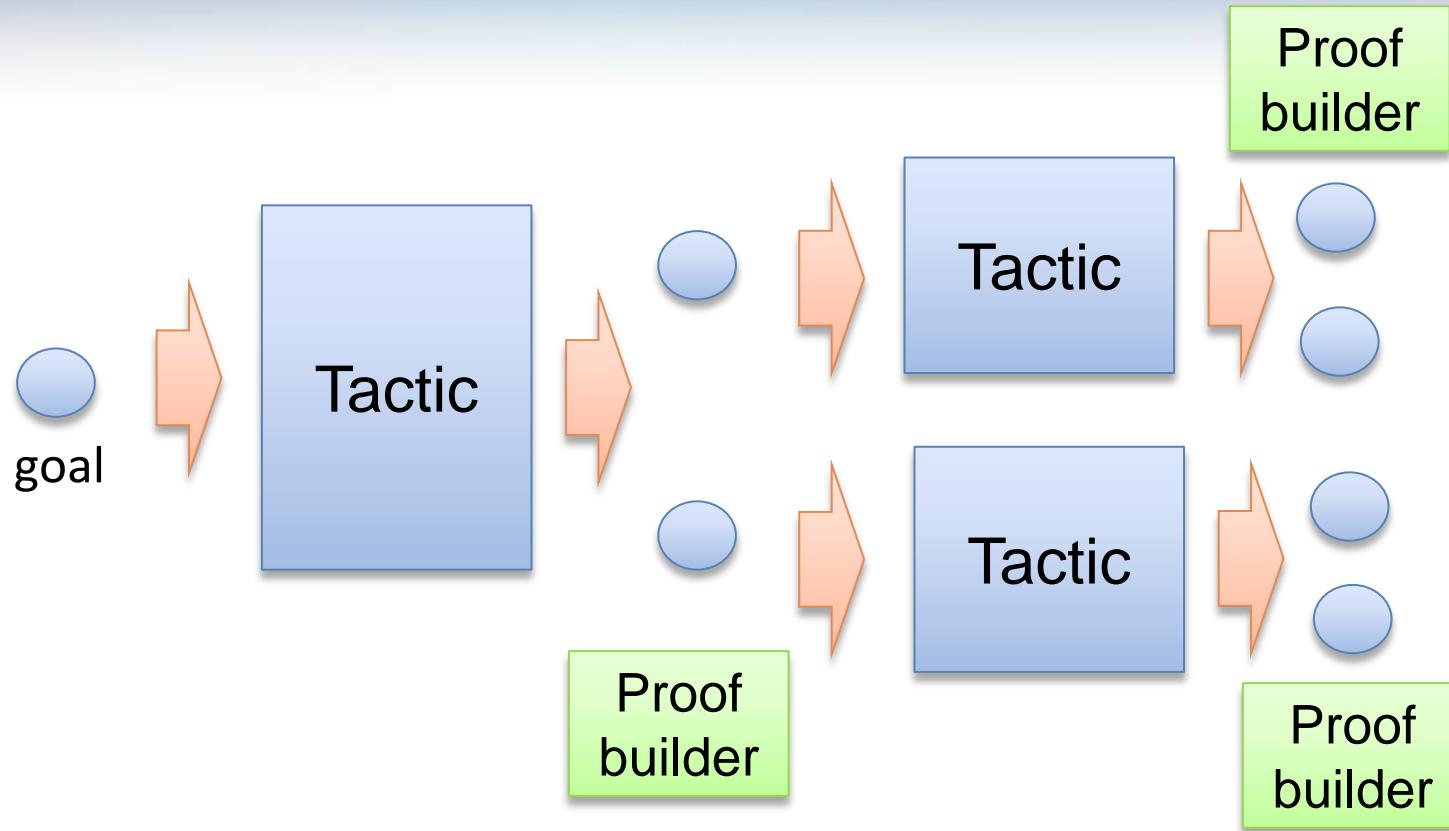


# Main inspiration: LCF-approach

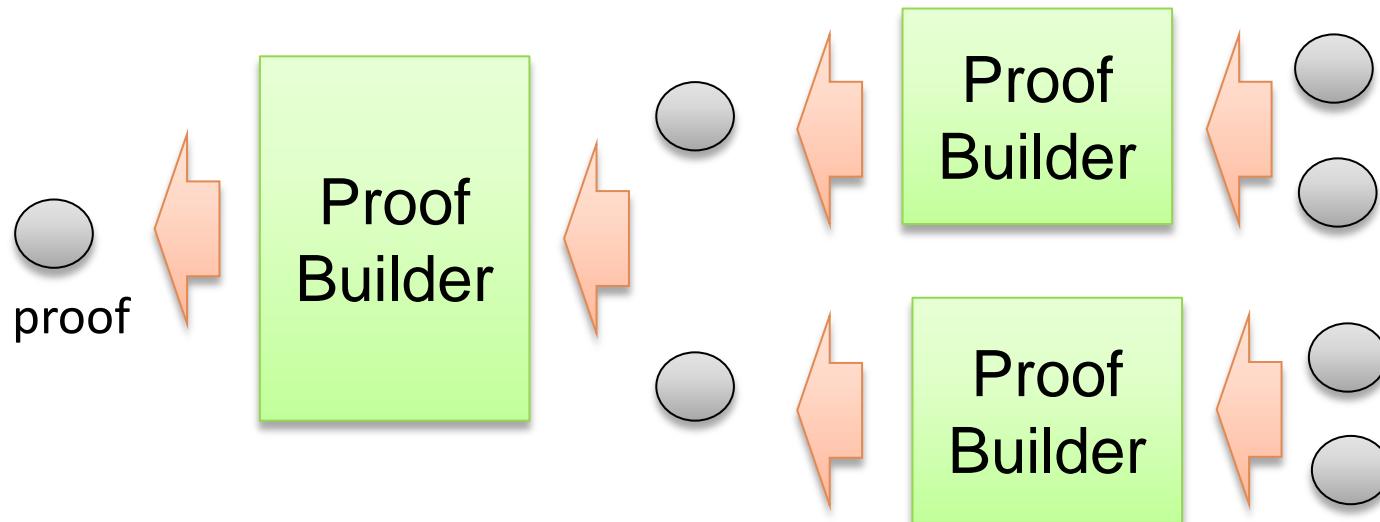


Proofs for subgoals

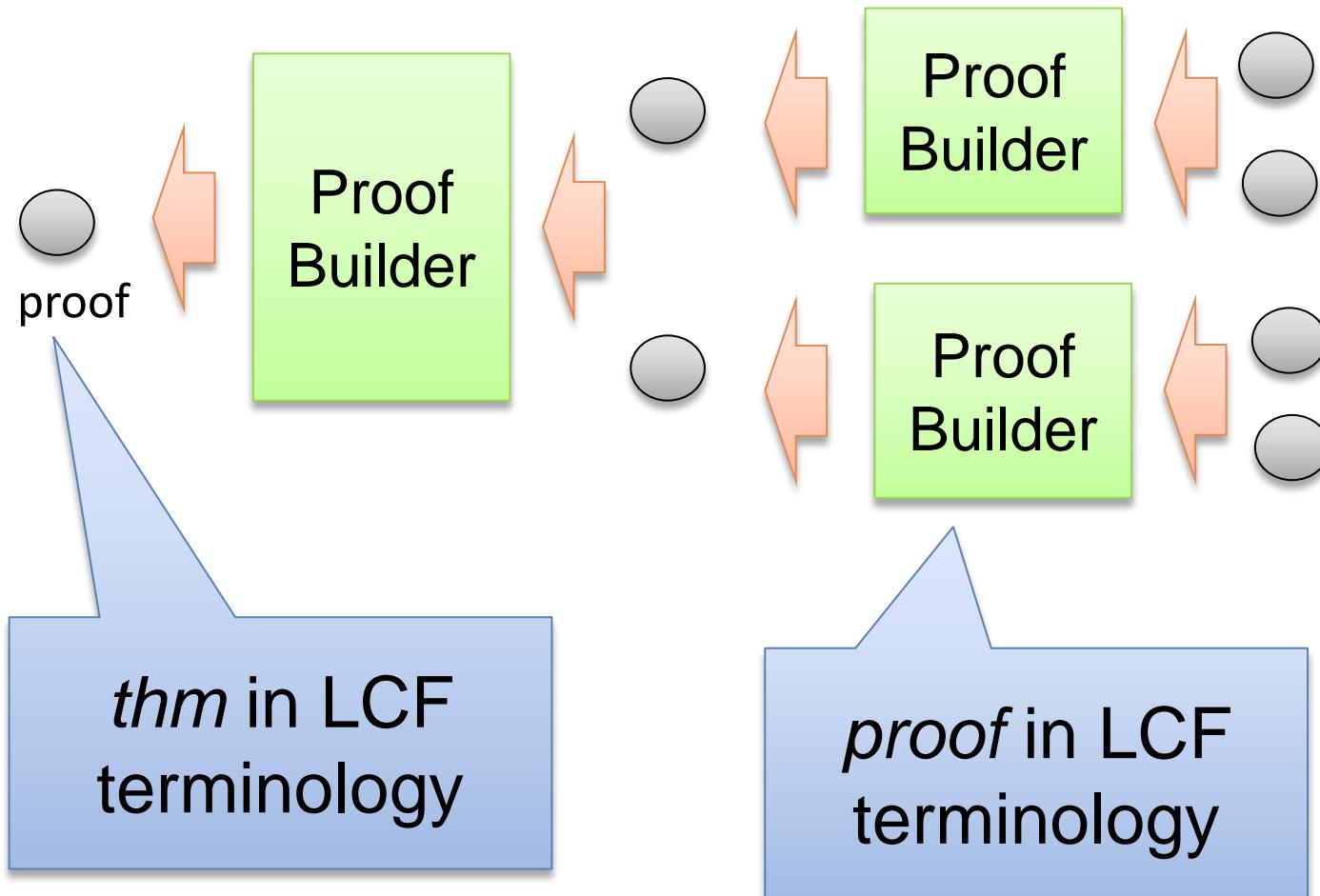
# Main inspiration: LCF-approach



# Main inspiration: LCF-approach



# Main inspiration: LCF-approach



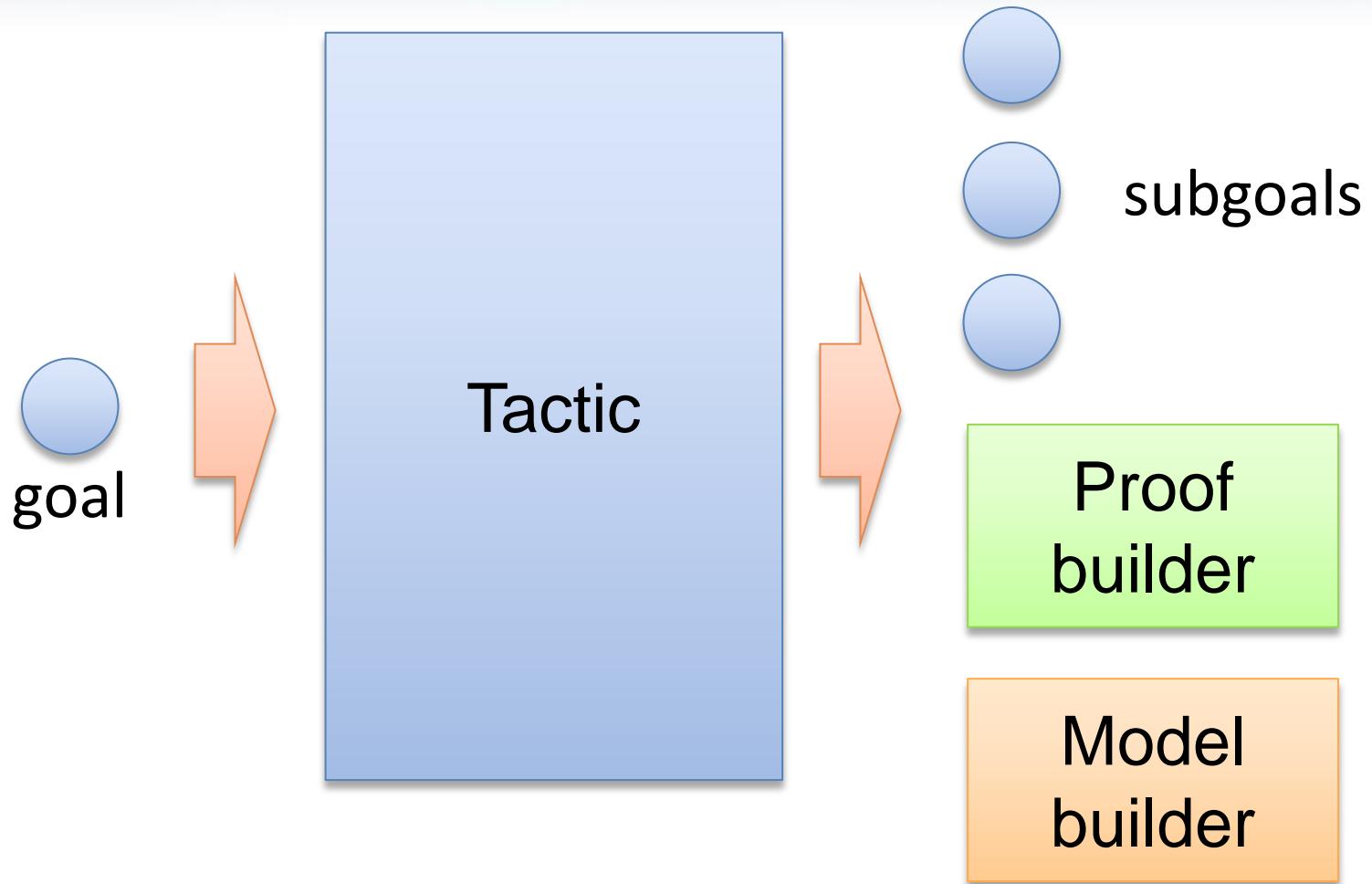
# Tacticals aka Combinators

then( Tactic , Tactic ) = Tactic

orelse( Tactic , Tactic ) = Tactic

repeat( Tactic ) = Tactic

# SMT Tactic



# SMT Tactic

*goal*        = *formula sequence*  $\times$  *attribute sequence*

*proofconv* = *proof sequence*  $\rightarrow$  *proof*

*modelconv* = *model*  $\times$  *nat*  $\rightarrow$  *model*

*trt*        = **sat** *model*

    |    **unsat** *proof*

    |    **unknown** *goal sequence*  $\times$  *modelconv*  $\times$  *proofconv*

    |    **fail**

*tactic*        = *goal*  $\rightarrow$  *trt*

# SMT Tactic

*goal*       $= \text{formula sequence} \times \text{attribute sequence}$

*proofconv*     $= \text{proof sequence} \rightarrow \text{proof}$

*modelconv*     $= \text{model} \times \text{nat} \rightarrow \text{model}$

*trt*             $= \text{sat model}$

  |    *unsat proof*

  |    *unknown goal sequence*  $\times$  *modelconv*  $\times$  *proofconv*

  |    *fail*

*tactic*         $= \text{goal} \rightarrow \text{trt}$

end-game tactics:  
never return unknown(sb, mc, pc)

# SMT Tactic

*goal*       $= \text{formula sequence} \times \text{attribute sequence}$

*proofconv*     $= \text{proof sequence} \rightarrow \text{proof}$

*modelconv*     $= \text{model} \times \text{nat} \rightarrow \text{model}$

*trt*             $= \text{sat model}$

|    *unsat proof*

|    *unknown goal sequence*  $\times$  *modelconv*  $\times$  *proofconv*

|    *fail*

*tactic*         $= \text{goal} \rightarrow \text{trt}$

non-branching tactics:  
sb is a singleton in  
*unknown*(sb, mc, pc)

# Trivial goals

**Empty goal** [ ] is trivially satisfiable

**False goal** [ ..., false, ...] is trivially unsatisfiable

basic : tactic

# SMT Tactic example

$$[ a = b + 1, (a < 0 \vee a > 0), b > 3 ]$$


Tactic:  
elim-vars

Proof  
builder

$$[ (b + 1 < 0 \vee b + 1 > 0), b > 3 ]$$


Model  
builder

# SMT Tactic example

$$[ a = b + 1, (a < 0 \vee a > 0), b > 3 ]$$

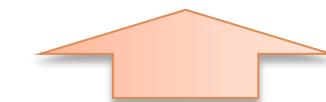
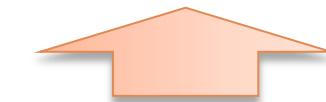

Tactic:  
elim-vars

$$M, M(a) = M(b) + 1$$

Proof  
builder

$$[ (b + 1 < 0 \vee b + 1 > 0), b > 3 ]$$


Model  
builder



M

# SMT Tactic example

$$[ a = b + 1, (a < 0 \vee a > 0), b > 3 ]$$


Tactic:  
split-or



Proof  
builder

$$\begin{aligned} & [ a = b + 1, a < 0, b > 3 ] \\ & [ a = b + 1, a > 0, b > 3 ] \end{aligned}$$

Model  
builder

# SMT Tactics

simplify

nnf

cnf

tseitin

lift-if

bitblast

gb

vts

propagate-bounds

propagate-values

split-ineqs

split-eqs

rewrite

p-cad

sat

solve-eqs

# SMT Tacticals

`then` :  $(tactic \times tactic) \rightarrow tactic$

`then`( $t_1, t_2$ ) applies  $t_1$  to the given goal and  $t_2$  to every subgoal produced by  $t_1$ .

`then*` :  $(tactic \times tactic\ sequence) \rightarrow tactic$

`then*`( $t_1, [t_{2_1}, \dots, t_{2_n}]$ ) applies  $t_1$  to the given goal, producing subgoals  $g_1, \dots, g_m$ .

If  $n \neq m$ , the tactic fails. Otherwise, it applies  $t_{2_i}$  to every goal  $g_i$ .

`orelse` :  $(tactic \times tactic) \rightarrow tactic$

`orelse`( $t_1, t_2$ ) first applies  $t_1$  to the given goal, if it fails then returns the result of  $t_2$  applied to the given goal.

`par` :  $(tactic \times tactic) \rightarrow tactic$

`par`( $t_1, t_2$ ) executes  $t_1$  and  $t_2$  in parallel.

# SMT Tacticals

$$\text{then}(\text{skip}, t) = \text{then}(t, \text{skip}) = t$$

$$\text{orelse}(\text{fail}, t) = \text{orelse}(t, \text{fail}) = t$$

# SMT Tacticals

`repeat : tactic → tactic`

Keep applying the given tactic until no subgoal is modified by it.

`repeatupto : tactic × nat → tactic`

Keep applying the given tactic until no subgoal is modified by it, or the maximum number of iterations is reached.

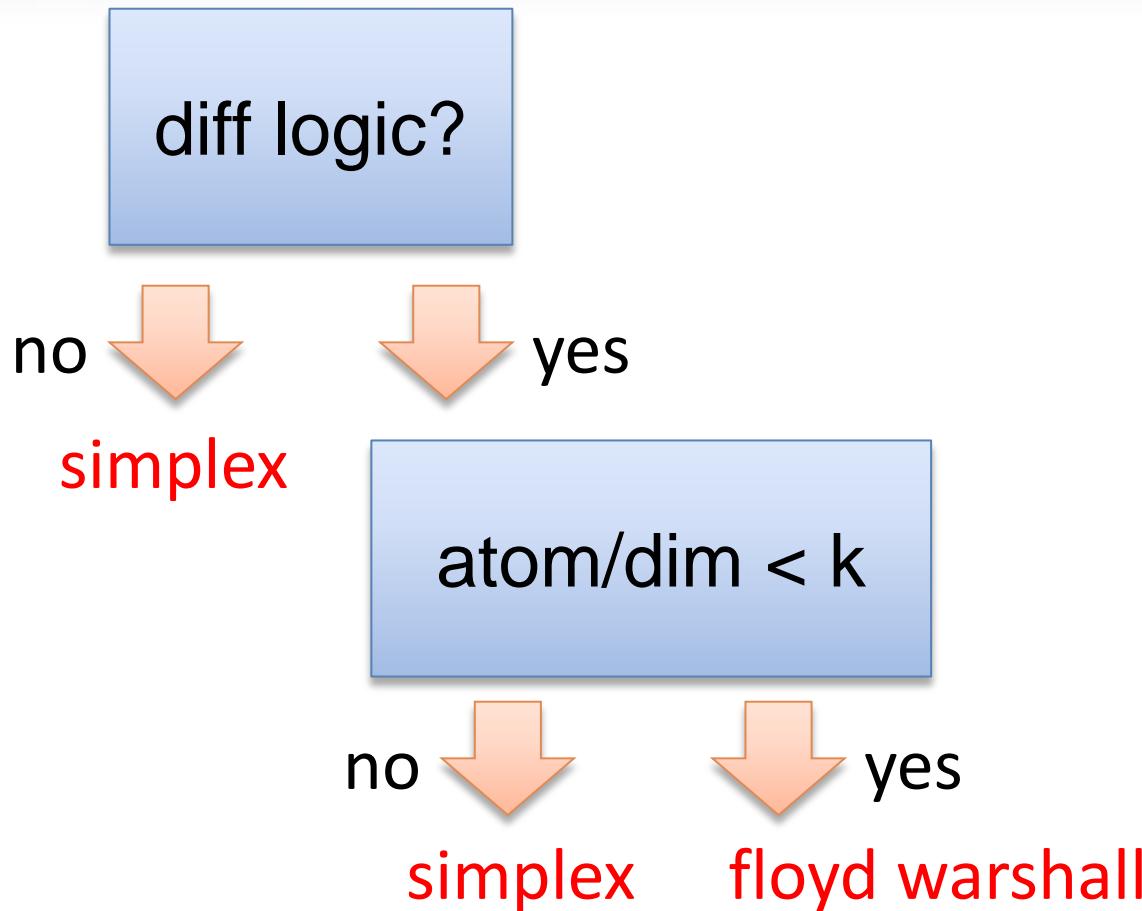
`tryfor : tactic × seconds → tactic`

`tryfor( $t, k$ )` returns the value computed by tactic  $t$  applied to the given goal if this value is computed within  $k$  seconds, otherwise it fails.

# Feature / Measures

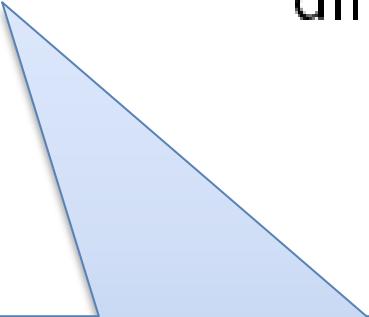
Probing structural features of formulas.

# Feature / Measures: Yices Strategy



# Feature / Measures: Yices Strategy

```
orelse(then(failif(diff ∧  $\frac{\text{atom}}{\text{dim}} > k$ ), simplex), floydwarshall)
```



Fail if condition is not satisfied.  
Otherwise, do nothing.

# Feature / Measures: Examples

`bw`: Sum total bit-width of all rational coefficients of polynomials in case.

`diff`: True if the formula is in the difference logic fragment.

`linear`: True if all polynomials are linear.

`dim`: Number of arithmetic constants.

`atoms`: Number of atoms.

`degree`: Maximal total multivariate degree of polynomials.

`size`: Total formula size.

# Tacticals: syntax sugar

`if( $c$ ,  $t_1$ ,  $t_2$ ) = orelse(then(failif( $\neg c$ ),  $t_1$ ),  $t_2$ )`

`when( $c$ ,  $t$ ) = if( $c$ ,  $t$ , skip)`

# Under/Over-Approximations

Under-approximation

unsat answers cannot be trusted

Over-approximation

sat answers cannot be trusted

# Under/Over-Approximations

Under-approximation  
model finders

Over-approximation  
proof finders

# Under/Over-Approximations

Under-approximation

$$S \rightarrow S \cup S'$$

Over-approximation

$$S \rightarrow S \setminus S'$$

# Under/Over-Approximations

Under-approximation

Example: QF\_NIA model finders  
add bounds to unbounded variables (and blast)

Over-approximation

Example: Boolean abstraction

# Under/Over-Approximations

Combining under and over is bad!  
sat and unsat answers cannot be trusted.

# Tracking: under/over-approximations

In principle, proof and model converters can check if the resultant models and proofs are valid.

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**Problem: if it fails what do we do?**

# Tracking: under/over-approximations

In principle, proof and model converters can check if the resultant models and proofs are valid.

Problem: if it fails what do we do?

We want to write tactics that can check whether a goal is the result of an abstraction or not.

# Tracking: under/over-approximations

## Solution

Associate an **precision attribute** to each goal.

# Goal Attributes

Store extra logical information

Examples:

precision markers

goal depth

polynomial factorizations

# Decision Engines as Tacticals

AP-CAD ( tactic ) = tactic

# Decision Engines as Tactics

```
then(preprocess, smt(finalcheck))
```

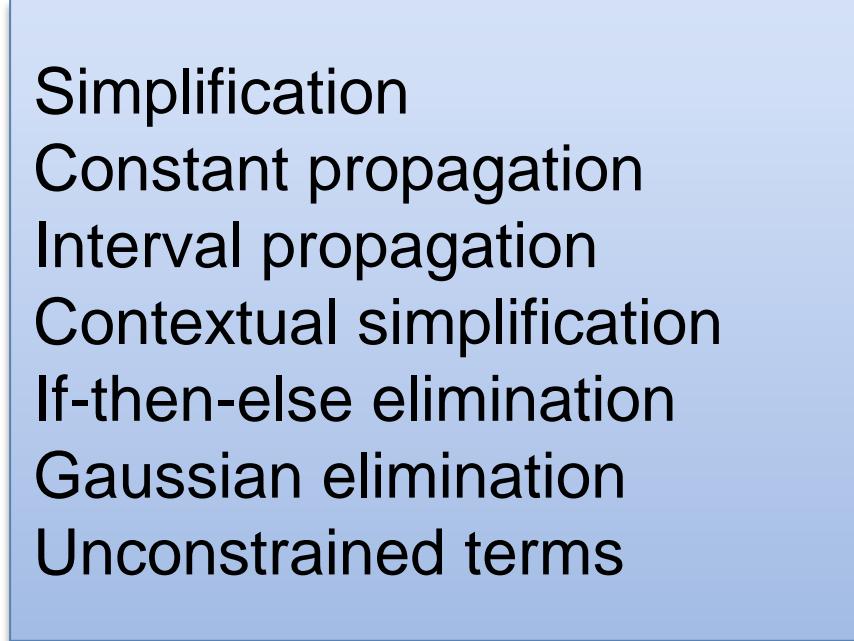
# Strategy: Example

```
then(then(simplify, gaussian), orelse(modelfinder, smt(apcad(icp))))
```

# RAHD Calculemus Strategy

# Z3 QF\_LIA Strategy

then(preamble, orelse(mf, pb, bounded, smt))



- Simplification
- Constant propagation
- Interval propagation
- Contextual simplification
- If-then-else elimination
- Gaussian elimination
- Unconstrained terms

# Challenge: small step configuration

proof procedure as a transition system

Abstract DPLL, DPLL(T), Abstract GB, cutsat, ...

UnitPropagate :

$$M \parallel F, C \vee l \implies M l \parallel F, C \vee l \quad \text{if } \begin{cases} M \models \neg C \\ l \text{ is undefined in } M \end{cases}$$

PureLiteral :

$$M \parallel F \implies M l \parallel F \quad \text{if } \begin{cases} l \text{ occurs in some clause of } F \\ \neg l \text{ occurs in no clause of } F \\ l \text{ is undefined in } M \end{cases}$$

Decide :

$$M \parallel F \implies M l^d \parallel F \quad \text{if } \begin{cases} l \text{ or } \neg l \text{ occurs in a clause of } F \\ l \text{ is undefined in } M \end{cases}$$

Fail :

$$M \parallel F, C \implies FailState \quad \text{if } \begin{cases} M \models \neg C \\ M \text{ contains no decision literals} \end{cases}$$

Backtrack :

$$M l^d N \parallel F, C \implies M \neg l \parallel F, C \quad \text{if } \begin{cases} M l^d N \models \neg C \\ N \text{ contains no decision literals} \end{cases}$$

# Challenge: small step configuration

proof procedure as a transition system

Abstract DPLL, DPLL(T), Abstract GB, cutsat, ...

## Challenge:

Efficient strategic control

$$M \parallel F \quad \Rightarrow \quad M l^d \parallel F \quad \text{if } \begin{cases} l \text{ or } \neg l \text{ occurs in a clause of } F \\ l \text{ is undefined in } M \end{cases}$$

Fail :

$$M \parallel F, C \quad \Rightarrow \quad FailState \quad \text{if } \begin{cases} M \models \neg C \\ M \text{ contains no decision literals} \end{cases}$$

Backtrack :

$$M l^d N \parallel F, C \quad \Rightarrow \quad M \neg l \parallel F, C \quad \text{if } \begin{cases} M l^d N \models \neg C \\ N \text{ contains no decision literals} \end{cases}$$

# Conclusion

Different domains need different strategies.

We must expose the little engines in SMT solvers.

Interaction between different engines is a must.

Tactic and Tactics: **big step approach**.

More transparency.