

Recursion - I

Definition : Function calling itself.

Why? → A problem is broken down into smaller problem (**subproblem**) and the solution is generated using the subproblems.

Eg. Sum of first N natural numbers.

$$\text{sum}(5) = 1 + 2 + 3 + 4 + 5$$

$$= \text{sum}(4) + 5$$

Writing recursive code :

- 1) Assumptions : Decide what the function will do and its output.
- 2) Main Logic : Break the problem into subproblems to solve the assumption.
- 3) Base Case : Identify the inputs for which we will stop recursion.

Function call Tracing

It involves the sequence of function calls that are made when a program is executed.

825

```
int add(int x, int y)
{
    return x+y;
}
```

```
int mul(int x, int y)
{
    return x*y;
}
```

```
int sub(int x, int y)
{
    return x-y;
}
```

```
void print(int x)
{
    printf(x);
}
```

```
int main()
{
    int x=10;
    int y=20;
}
```

```
print(sub(mul(add(x,y),30),75));
return 0;
```

Main()

x=10, y=20

, print(sub(mul(add(x,y),30),75));

↓ ↑

print(sub(mul(add(x,y),30),75))

↓ 825

sub(mul(add(x,y),30),75)

↓ 900

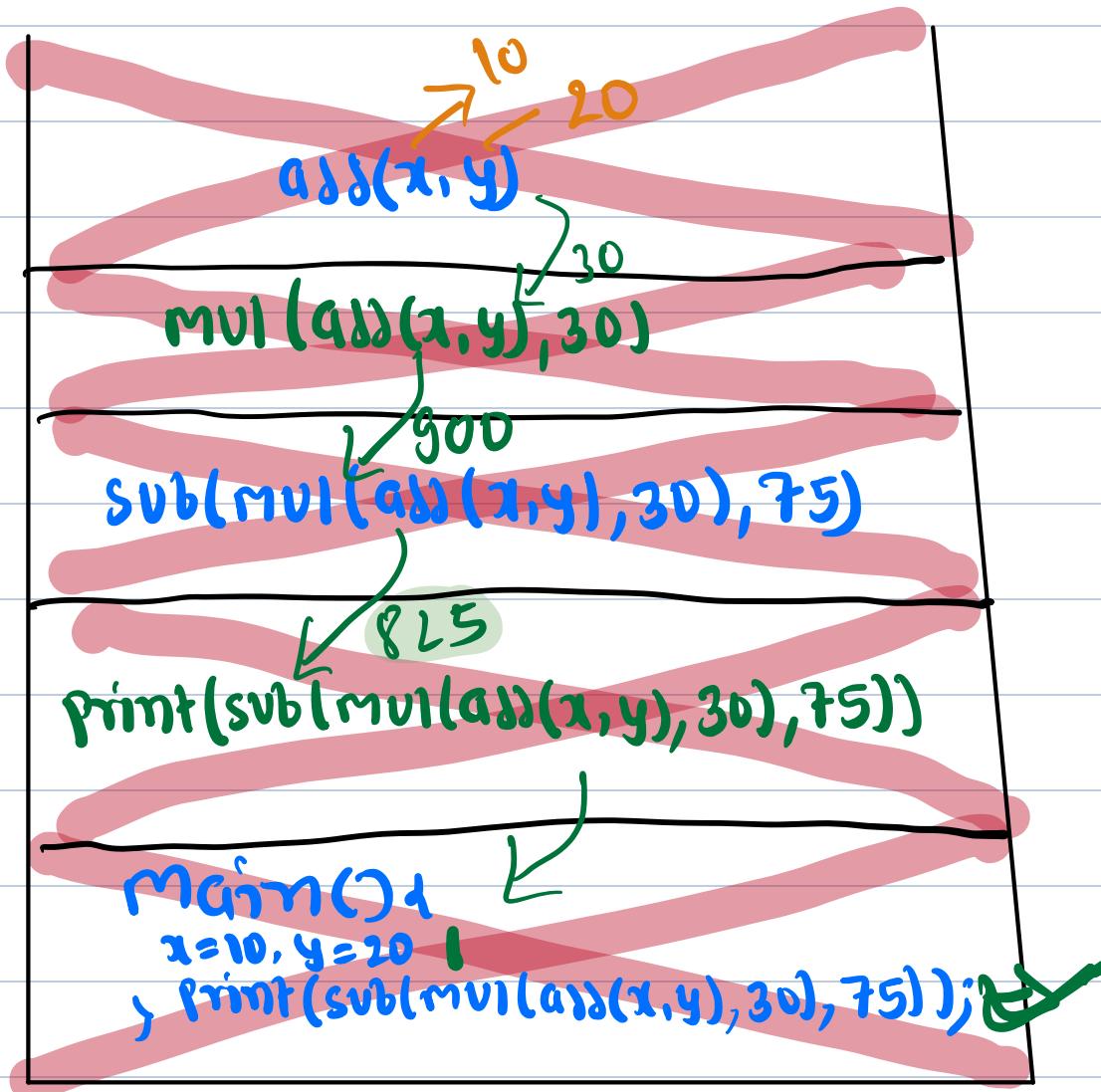
mul(900(x,y),30)

↓ ↑ 30

add(x,y)

↓ ↓ 10 20

DS involved in function calls



Q1) Given a positive integer N , find the factorial of N .

$$N = 5$$

$$5! = 5 * 4 * 3 * 2 * 1 = 120$$

$$\text{fact}(5) = \text{fact}(4) * 5$$

Solution:

Assumption: Create a function which takes N as an input and return $\text{fact}(N)$ as an output.

Main logic: $\text{fact}(n) = \text{fact}(n-1) * n$

Base case:

If ($n == 1$)

return ;

```
int fact( n ) {  
    if( n == 1 ) return 1;  
    return n * fact( n - 1 );
```

6

Call Tracing: print(fact(3));

```
int fact( n ) {  
    if( n == 1 ) return 1;  
    return n * fact( n - 1 );
```

3
2
1

```
int fact( n ) {  
    if( n == 1 ) return 1;  
    return n * fact( n - 1 );
```

Q) Print number in increasing order from 1 to N.

Eg. $N=5$

1 2 3 4 5

Qviz:

Solution:

Assumption: Create a function which takes N as an input and print all no. from 1 to N.

Main logic :

$$\text{inc}(n) = \text{inc}(n-1) \downarrow \text{print}(n);$$

Base case:

$$\text{if } (n == 1) \left\{ \begin{array}{l} \text{print } 1; \\ \text{return} \end{array} \right.$$

void inc(n) {
 if (n == 1) {
 print(n);
 return;
 }
 inc(n-1);
 print(n);
}

void inc(n) {
 if ($n == 1$) { print(n);
 return;
 }
}

inc($n - 1$);
 print(n);
}

void inc(n) {
 if ($n == 1$) { print(n);
 return;
 }
}

inc($n - 1$);
 print(n);
}

void inc(n) {
 if ($n == 1$) { print(n);
 return;
 }
}

inc($n - 1$);
 print(n);
}

Output: 1 2 3

9>

Problem

Whirlpool wants to design a timer for their **washing machines**. This feature is a simple countdown timer. When a user sets a time, for example, 10 minutes, the washing machine needs to show each minute passing, counting down until it reaches 0.

Your task is to write a program that takes an integer **A** (the time in minutes set by the user) and then prints out each minute as it counts down to 0. The requirement is that after a user set a timer for the washing machine for some time say **A**, the washing machine should display each minute after that decremented one by one till the time becomes 0.

Simplified Problem statement

Given **N**, print all numbers from **1 to N** in decreasing order.

Eg. $N=5$

5 4 3 2 1

```
void des(N) {  
    if (N == 1) {  
        print(n);  
        return;  
    }  
    print(n);  
    des(N-1);  
}
```

TC: O(n)
SC: O(n)

Time and Space Complexity of Recursion

Time complexity = $O(\text{no. of function call} * \text{time per function call})$

Space complexity = $O(\text{maximum depth of recursion tree} * \text{stack space + stack per fun call})$

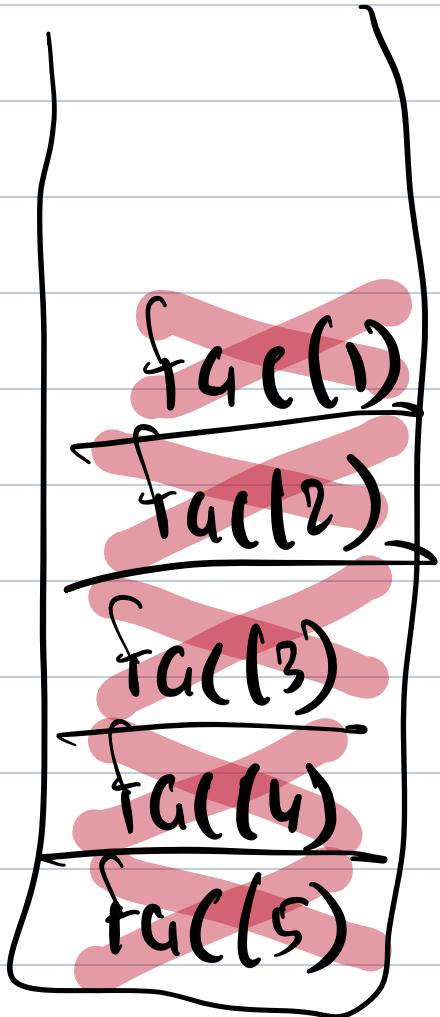
Factorial of number

Code

```
int fac(int n){  
    if (N == 0) {  
        return 1;  
    }  
    return n * fac(n-1);
```

$N = 5$

$\text{fac}(5)$
↓
 $5 * \text{fac}(4)$
↓
 $4 * \text{fac}(3)$
↓
 $3 * \text{fac}(2)$
↓
 $2 * \text{fac}(1)$
↓
 $\text{fac}(1)$



TC: $O(n)$
SC: $O(n)$

Q) Printing number in Decreasing Increasing order.

Eg. $N=1 \rightarrow 1$
 $N=3 \rightarrow 3\ 2\ 1\ 1\ 2\ 3$

$N=5 \rightarrow 5\ 4\ 3\ 2\ 1\ 1\ 2\ 3\ 4\ 5$

$\text{DesInc}(5) = \text{print}(5) + \text{desInc}(4) + \text{print}(5)$

↓
subproblem

Output: 4 3 2 1 1 2 3 4

Code:

function desinc(n)

d

If ($n == 0$) return;

print(n);
desinc($n - 1$);
print(n);

DI(4)

↓ ↗

DI(3)

↓ ↗

DI(2)

↓ ↗

DI(1)

↓ ↗

DI(0)

Break:
10-30

8) Printing N^{th} number in fibonacci series

Series: 0 1 1 2 3 5 8 13
0 1 2 3 4 5 6 7

$N=6$

Qviz $N=7$. $\text{fib}(7) = \text{fib}(5) + \text{fib}(6)$
 $= 5 + 8$
 $= 13$

int fib(n) {
 if($n==0$ || $n==1$) {return n}
 return fib(n-1) + fib(n-2);
}

Solution:

Assumption: Create a fun which takes a integer n as input and return N^{th} term of fibonacci series.

Main logic:

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$$

Base case:

If($n == 0 \text{ || } n == 1$)
 { return n ;
 }



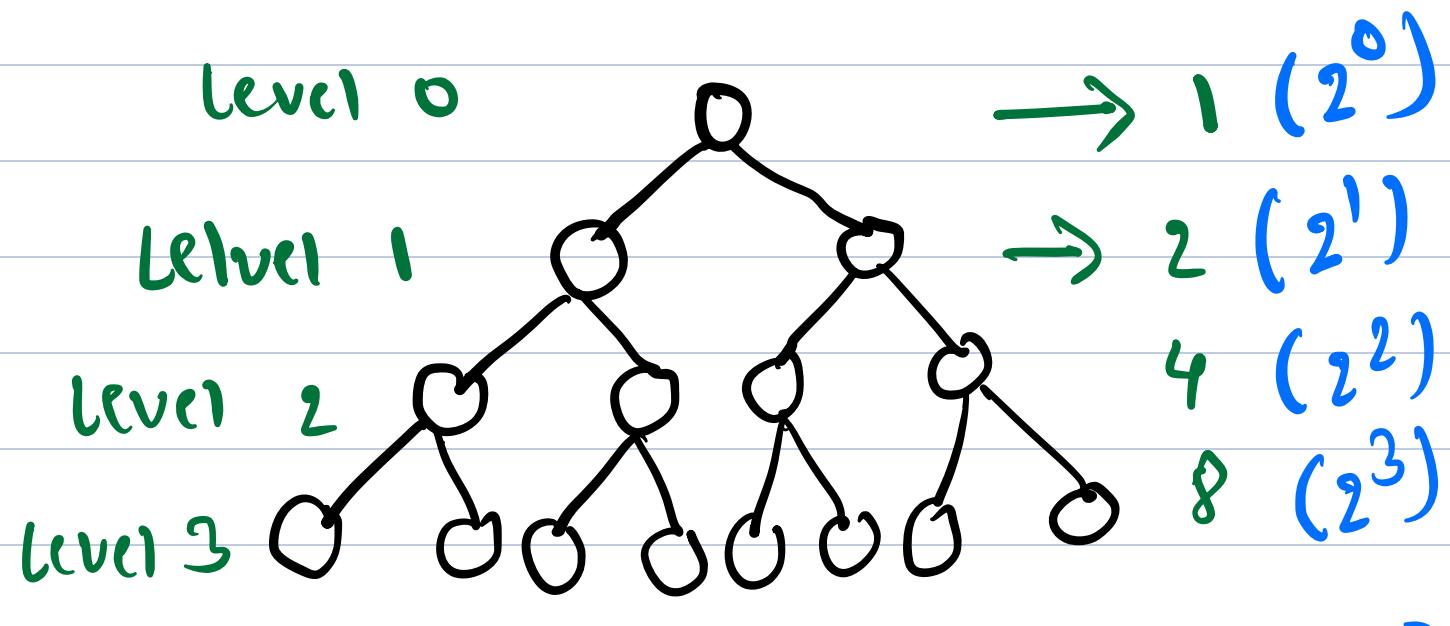
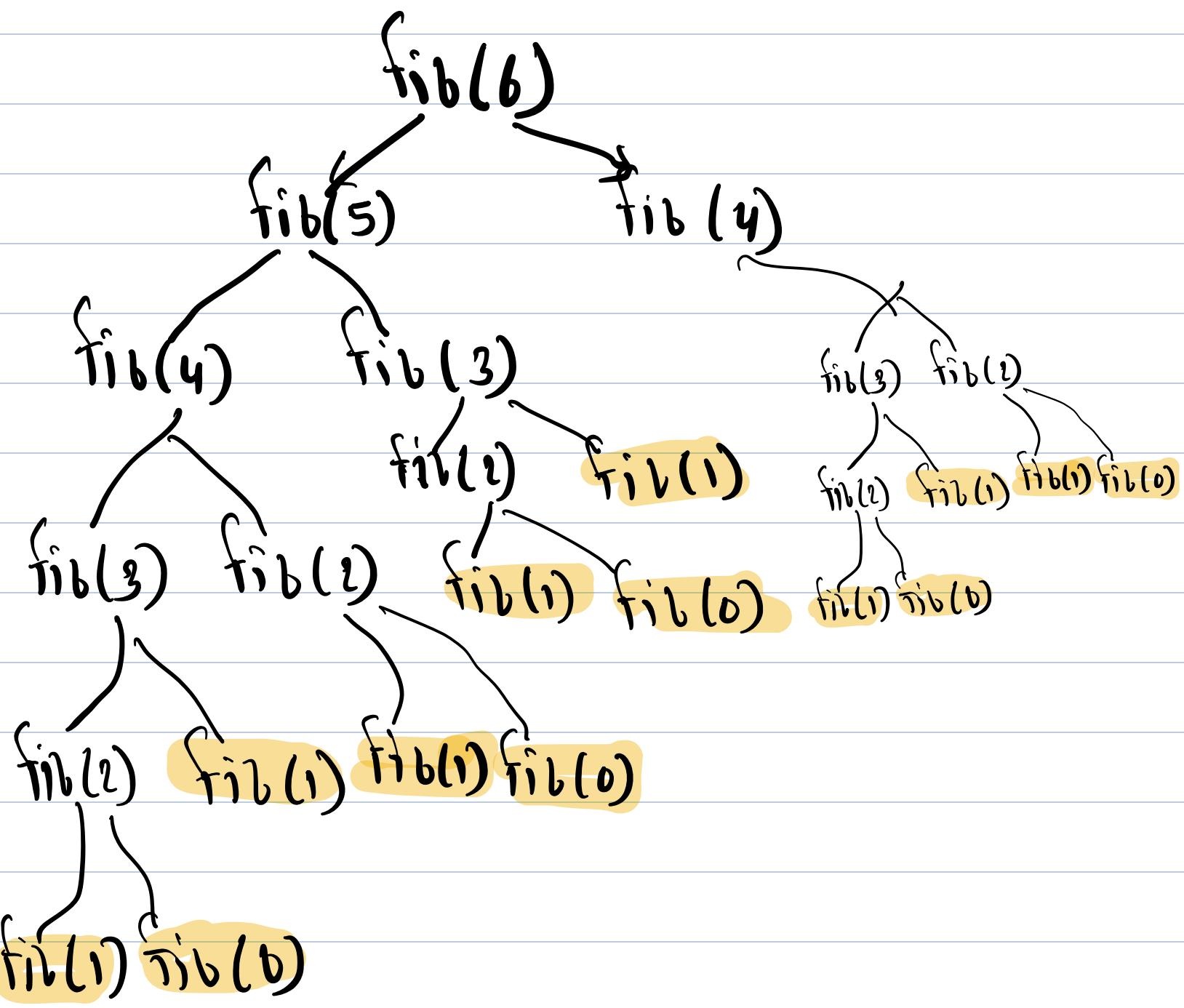
```
int fib(n) {  
    if (n==0 || n==1) return n;  
    return fib(n-1)+fib(n-2);  
}
```

Diagram illustrating the execution flow of the fib function for $n=1$. The flow starts at the top node, which calls fib(1). This leads to the base case check if $n=0$ or $n=1$. Since $n=1$, it returns 1. A green arrow points from this return value to the bottom node, which then returns 1.

```
int fib(n) {  
    if (n==0 || n==1) return n;  
    return fib(n-1)+fib(n-2);  
}
```

```
int fib(n) {  
    if (n==0 || n==1) return n;  
    return fib(n-1)+fib(n-2);  
}
```

```
int fib(n) {  
    if (n==0 || n==1) return n;  
    return fib(n-1)+fib(n-2);  
}
```



level 4

16 (2^4)

level 5

32 (2^5)

$$2^0 + 2^1 + 2^1 + 2^2 \dots - 2^x$$

$$\frac{a(r^n - 1)}{(r - 1)}$$

$$= \frac{1 * (2^n - 1)}{(2 - 1)}$$

$$= 2^n - 1$$

TC $\approx O(2^n)$

SC = $O(N)$

MAX STACK SPACE = ~~0X2345~~

CURRENT STACK SPACE = ~~0X2345~~

~~X5K345~~
~~3234321~~

