

LinkedList 1



Good
Evening

* Agenda for Today

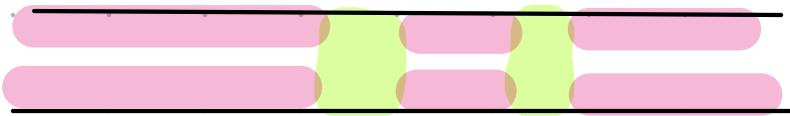
01. Basics of LinkedList
02. Accessing of element
03. Searching for a value
04. Insertion in LL
05. Deletion in LL
06. Reverse a LinkedList (v v I)
07. Check if palindrome

Linkedlist

→ Linear data structure

Issues with Array

- Fixed size
- Costly insert & delete operation
- Continuous memory location



- Linkedlist → Utilise all the free memory available whether its continuous or non continuous
- Insertion & Deletion in linkedlist is optimal as compared to Arrays.



```
class Node {
```

```
    int data;
```

```
    Node next;
```

```
    Node ( int d ) {
```

```
        data = d;
```

```
        next = null;
```

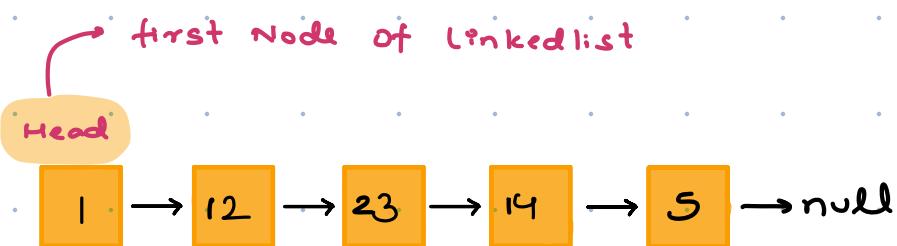
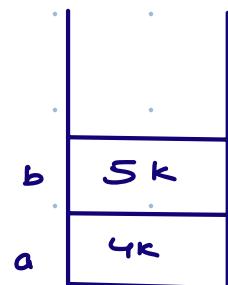
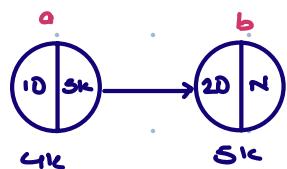
```
}
```

```

Node a = new Node(10);
Node b = new Node(20);

a.next = b

```



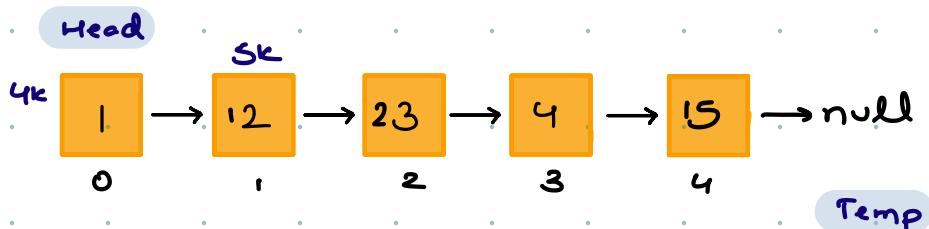
- Any linked list can be represented by the first node or head.

`int size (Node head){}` → Question signature

- * Never move head pointer
- * Null pointer Exception

* Operation on Linkedlist

01. Access the k^{th} element ($k=0 \rightarrow \text{represents first element}$)



(6)

$$k = 3 \longrightarrow \underline{\underline{\text{Ans} = 4}}$$

$$k = 2 \longrightarrow \text{Ans} = 23$$

$$k = 6 \longrightarrow \text{Ans} = -1$$

```
if (head == null) return -1;
```

```
Node temp = head;
```

```
for (i=0; i<k; i++) {
    if (temp == null) return -1;
    temp = temp.next
}
```

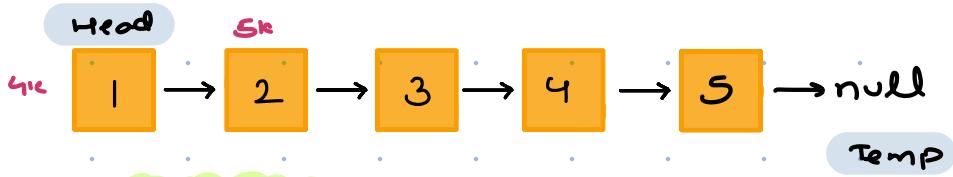
```
return temp.data;
```

$Tc: O(N)$
 $Sc: O(1)$

* Check for value X (searching)

\downarrow

Value



$$x = 3 \longrightarrow \underline{\underline{\text{Ans} = \text{true}}}$$

```

if (head == null) return false;
Node temp = head;

while (temp != null) {
    if (temp.data == x) return true;
    temp = temp.next;
}

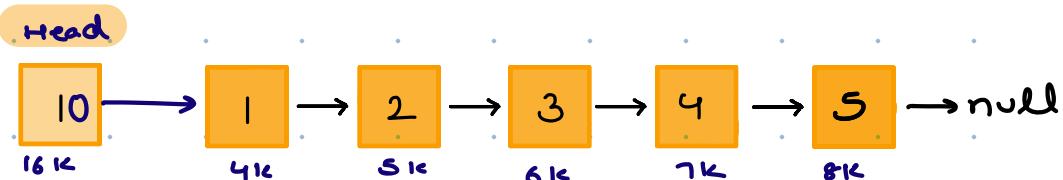
return false;

```

TC: O(n)
SC: O(1)

* Insert a new Node with Data at kth position

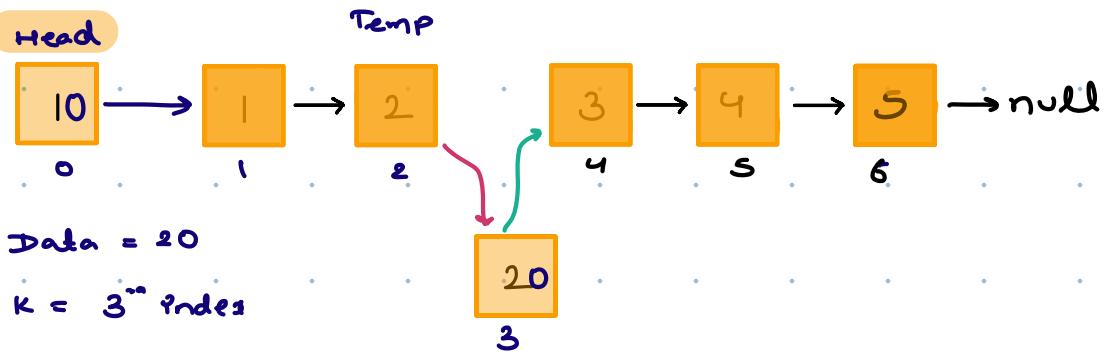
Data = 10
K = 0 } Add it in front



```

Node nn = new Node(data);
nn.next = head;
head = nn;

```



```
Node nn = new Node(20);
```

Edge cases

```
Node temp = head
```

```
i=0
```

```
while ( i < k-1 ) {
```

```
    temp = temp.next
```

```
    i++;
```

```
}
```

```
nn.next = temp.next
```

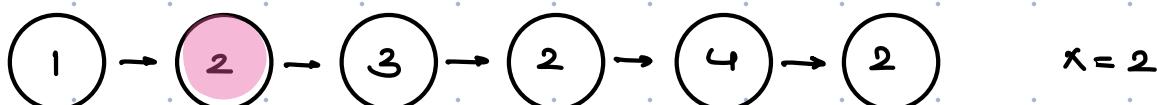
```
temp.next = nn;
```

01. Null case while traversing

02 Head is null

* Deletion in Linkedlist

→ Delete the first occurrence of value x given in linkedlist. If x is not present, leave LL as it is.



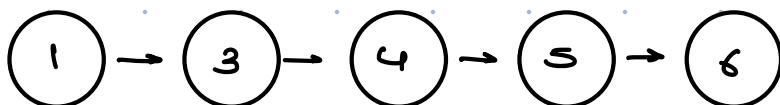
Ans ⇒

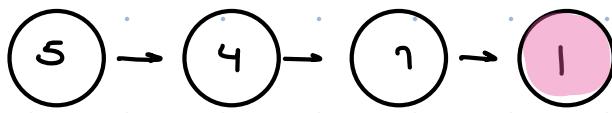


$x=2$

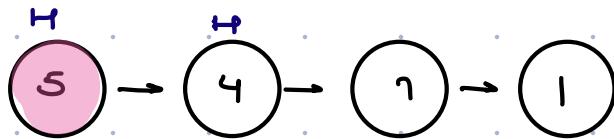


Ans =

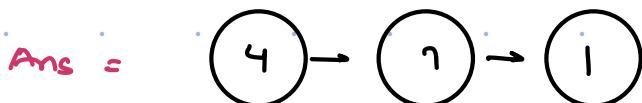




$x = 1$



$x = 5$



Node deletefirst (Node head) {

If (head == null) return head;

if (head.data == x) {

head = head.next;

return head;

temp = head;

while (temp != null) {

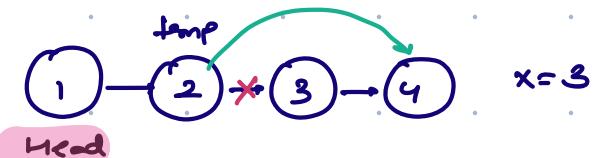
if (temp.next != null & temp.next.data == x)

temp.next = temp.next.next;

return head;

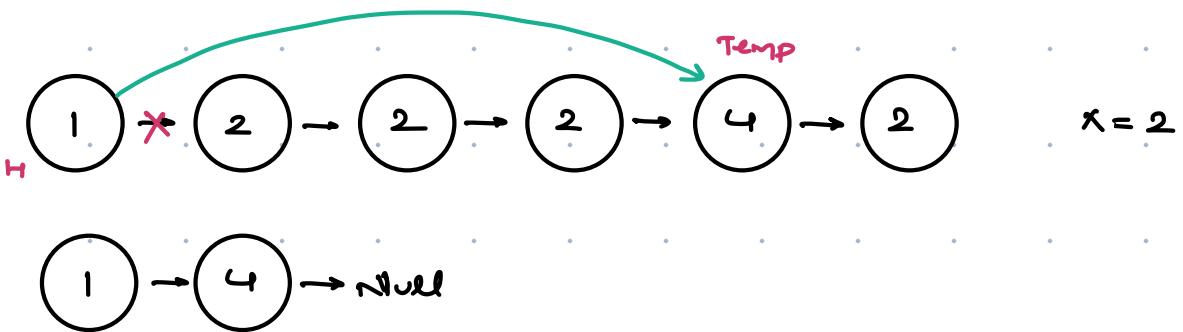
temp = temp.next;

return head;



$x = 3$

* Delete all Occurrences of x



Idea → check for all value of x & move your next ptr accordingly

while (`head != null && head.data == x`) {

`head = head.next;`

 Node temp = head;

 while (`temp != null && temp.next != null`) {

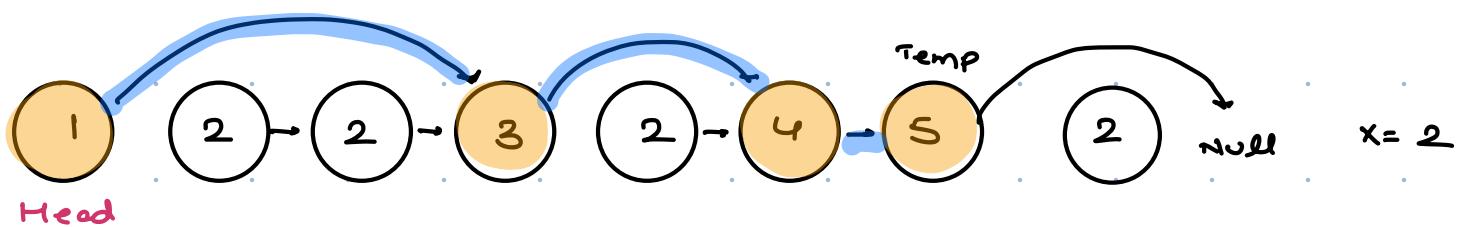
 if (`temp.next.data == x`) {

`temp.next = temp.next.next;`

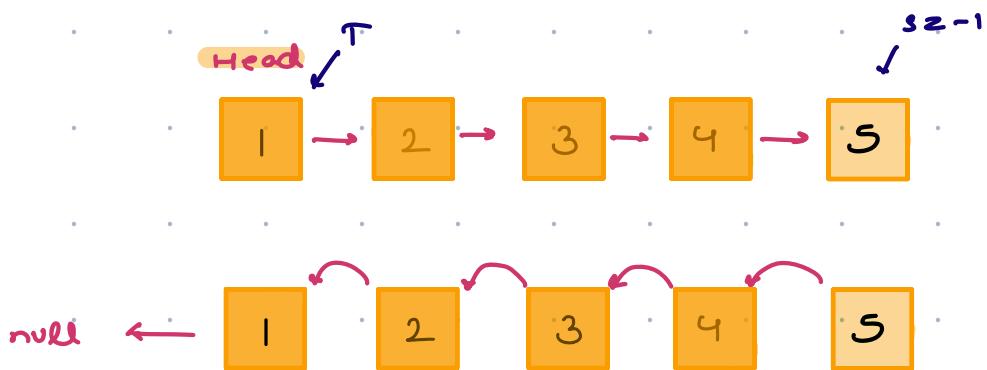
 } else {

`temp = temp.next;`

 return head;



* Reverse a Linkedlist



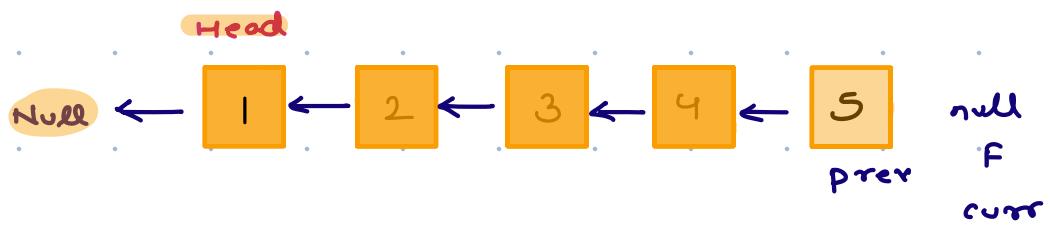
* Brute force

01. size of LL

02. temp = 0
last = $sz - 1$ } swapping

03. temp = 1
last = $sz - 2$ } swapping

* Use 3 pointers



```
Node prev = null;  
Node curr = head
```

```
while (curr != null) {
```

```
    Node form = curr.next;
```

```
    curr.next = prev;
```

```
    prev = curr;
```

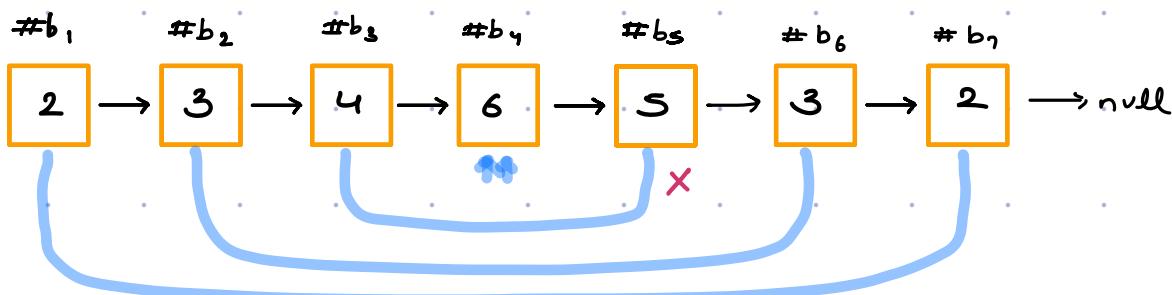
```
    curr = form;
```

3

```
return prev;
```

TC : O(N)
SC : O(1)

* Given a Linkedlist, check if it is a **palindrome**.



Not a palindrome

{ A C A
 N a m a n
 M a a m }

1

= Valid palindrome

Solution 1

- Reverse the entire linkedlist
- Compare the old linkedlist & reversed LL.

Solution 2

- Get the mid of Linkedlist → $\text{size}/2$
- Reverse the second half
- Compare the first half with reversed second half.

Try this Approach