

Hashing I

Today's Agenda:

- * HashMap & HashSet functionality
- * Question on Hashing.

Motel Business

R1

R2

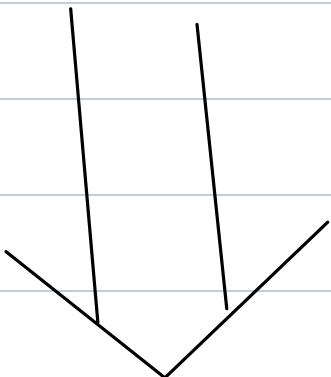
R3

R4

R5

Register:

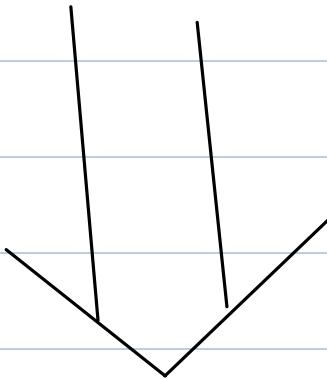
<u>Rooms</u>	<u>Occupancy</u>
R1	Occupied
R2	Occupied
R3	Vacant
R4	Occupied
R5	Vacant



1000 rooms

0	1	2	3		999	1000
	F	T	F		

Array of size 1001



Room no.s are still 1000 but numbering
to be done randomly between $\{1, 10^3\}$

Eg. 100035, 28, 34586

$$10^6 = 10^3 * 10^3 \quad \times$$

$$= 4 * 10^3 * 10^3$$

⇒ Create an array of size 10^9 .

Problem → Memory Wasteage
(only 1000 rooms are there) -

Solutions: Hashmap | Map | Dictionary -

→ Create a key value pair value for each room.

Key	Value
100035	Vacant
:	
28	Occupied
:	
34586	Occupied

Advantages:

- 1) space $\rightarrow 10^3$
- 2) Adding new room (a new Key value pair) \rightarrow constant time O(1)
- 3) Accessing availability of a room \rightarrow constant time O(1)

HashMap<Integer, String> hm;

Why? Constant time
will be having dedicated session
for it

Quiz 1: Population for every country.

String

Long

China → 1,300,000,000

India → 1,400,300,000

Nepal → 300,000

HashMap<String, Long> hm

Quiz 2: No. of state of every country.

String

Int

India → 29

Pakistan → 10

USA → 50

China → 15

HashMap<String, Integer> hm

Qviz 3: Name of all state of every country.

String

List<String>

India →

UP, MP, bca - - - -

USA →

California, Texas - - - -

Pakistan →

Karachi, Lahore

UK →

A, B, C

HashMap<String, List<String>> hm

QViz 4: Population of each state in every country.

String

India



HashMap<String, Integer>

U.P → 250,000

MP → 10,000,000

:

:

US



Texas → 250,000

California → 5,000,000

!

?

UK



A → 70,000

B → 30,000

!

!

0
0

;

0
0

;

Properties:

- 1) Keys are unique.
- 2) Values can be anything.
- 3) Searching will take constant time.

* If we are only interested in unique keys, value are irrelevant.

→ Use HashSet

HashMap

- **INSERT(Key,Value)**: new key-value pair is inserted. If the key already exists, it does no change.
- **SIZE**: returns the number of keys.
- **DELETE(Key)**: delete the key-value pair for given key.
- **UPDATE(Key,Value)**: previous value associated with the key is **overridden** by the new value.
- **SEARCH(Key)**: searches for the specified key.

HashSet

- **INSERT(Key)**: inserts a new key. If key already exists, it does no change.
- **SIZE**: returns number of keys.
- **DELETE(Key)**: deletes the given key.
- **SEARCH(Key)**: searches for the specified key.

Time Complexity of all the operations in both Hashmap and HashSet is **O(1)**.

Hashing Library Names in Different Languages

Java	C++	Python	Js	C#
HashMap	HashMap	unordered_map	dictionary	map
HashSet	HashSet	unordered_set	set	HashSet

Implementations :

HashMap

Java	HashMap<Integer, Integer> map = new HashMap<>();	HashSet set = new HashSet<>();
C++	unordered_map<int, int> map;	unordered_set set;
Python	dictionary = {}	set = set()
Js	let map = new Map();	let set = new Set();
C#	Dictionary<int, int> dict = new Dictionary<int, int>();	HashSet set = new HashSet();

HashSet

Q) Given N elements and Q queries. find the frequency of elements provided in the query.

0	1	2	3	4	5	6	7	8	9	10	
A[1] : →	2	6	3	8	2	8	2	3	8	10	6

Queries

$2 \rightarrow 3$

$8 \rightarrow 3$

$3 \rightarrow 2$

$5 \rightarrow 0$

BF: For each query iterate over an array and count its occurrence.

TC: $O(n * Q)$

SC: $O(1)$

Optimisation: We use HashMap.

	0	1	2	3	4	5	6	7	8	9	10
A[]:	2	6	3	8	2	8	2	3	8	10	6

2 → ~~1 2 3~~

6 → ~~1 2~~

3 → ~~1 2~~

8 → ~~1 2 3~~

10 → 1

HM

HashMap < Int, Int >

- ⇒ first create HashMap and store freq. of each element.
- ⇒ Iterate over query and check freq in HashMap in O(1)

Code:

HashMap<int, int> hm

```
for(i=0 ; i<n ; i++)
```

```
{  
    if(hm.search(A[i]) == true)
```

```
        hm.update(A[i], hm.get(A[i])+1);
```

y
else

```
{  
    hm.insert(A[i], 1);
```

TC: O(n+q)
SC: O(n)

```
for(i=0 ; i<q ; i++)
```

```
{  
    if(hm.search(qΣi) == true)
```

```
        print(hm.get(qΣi));
```

y
else
{
 print(0);

Q.2) Given N elements, find the first non repeating element.

Eg.

$$N=6 \quad \{1, 2, 3, 1, 2, 5\} \rightarrow 3$$

$$N=8 \quad \{4, 3, 3, 2, 5, 6, 4, 5\} \rightarrow 2$$

$$N=7 \quad \{2, 6, 8, 4, 7, 2, 9\} \rightarrow 6$$

Brute force: For each query iterate over an array and find its frequency. The first element found with no repeating occurrence is the answer.

TC: $O(N^2)$

SC: $O(1)$

$$N=8 \quad \{4, 3, 3, 2, 5, 6, 4, 5\} \rightarrow 2$$

4:2	2:1	6:1
3:2	5:2	

Optimisation :

- » Create a freq. Hashmap like previous question.
- 2) Iterate on the original array and check it's freq. The first element having freq. as 1 would be answer.

Code:

```
HashMap<int,int> hm;
for(i=0; i<n; i++)
{
    if(hm.search(A[i]) == true)
        hm.update(A[i], hm.get(A[i])+1);
    else
        hm.insert(A[i], 1);
}
for(i=0; i<n; i++)
{
    if(hm.get(A[i]) == 1)
        return A[i];
}
return -1;
```

TC: O(n)

SC: O(n)

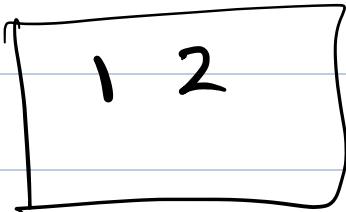
Break: 10-20

Q.3) Count of distinct elements-

$$N=5 \Rightarrow \{3, 5, 6, 5, 4\} \rightarrow 4$$

$$N=3 \Rightarrow \{3, 3, 3\} \rightarrow 1$$

$$N=5 \Rightarrow \{1, 1, 1, 2, 2\} \rightarrow 2$$



Approach: Create HashSet and store every element in HashSet, The size of HashSet will give the total no. of unique numbers.

```
Hashset<int> hs;  
for(i=0; i<n; i++)  
    2
```

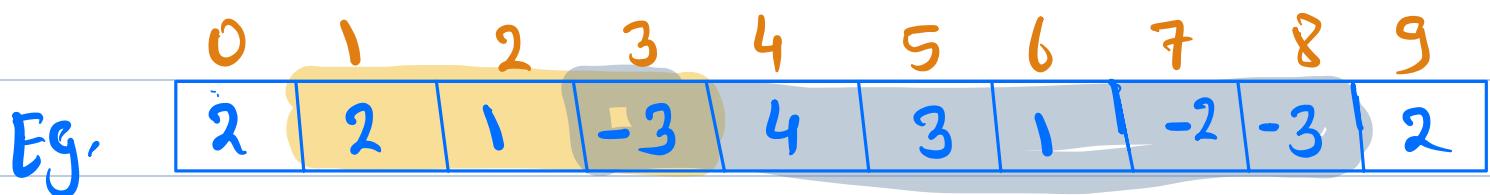
TC: O(n)

SC: O(n)

```
    hs.insert(A[i]);
```

```
return hs.size();
```

Q.4) Given an array of N elements, check if there exists a subarray with sum equal to 0.



BF Approach: Consider all the subarrays.

```
for (s=0; s<n; s++)  
{
```

```
    for (e=s; e<n; e++)
```

```
    {
```

sum = 0

```
        for (k=s; k<=e; k++)
```

```
            sum = sum + A[k];
```

```
            if (sum == 0) return true;
```

```
    }
```

TC: $O(n^3)$
SC: $O(1)$

return false;

* Using PF sum or CF $\rightarrow O(n^2)$

Optimisation

0	1	2	3	4	5	6	7	8	9
2	2	1	-3	4	3	1	-2	-3	2

PF	2	4	5	2	6	9	10	8	5	7
----	---	---	---	---	---	---	----	---	---	---

Eg. $\{ -2, -1, 3, 5 \}$

PF $\{-2, -3, 0, 5\}$

Code:

|| Create Prefix Array.

Hashset<int> hs;

hs.insert(0);

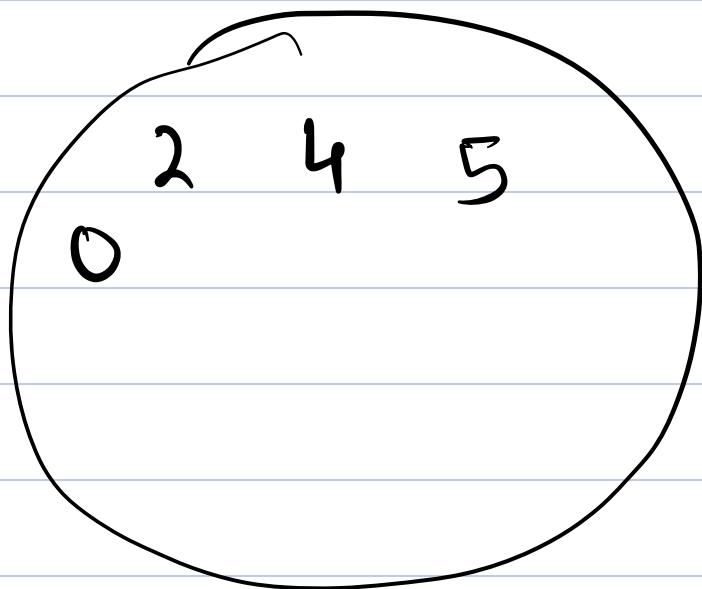
for (i=0; i < PF.length; i++)

if (hs.search(PF[i]) == true)

return true;

else
} } ns.insert(PF[i]);
return False;

2	4	5	2	6	9	10	8	5	7
---	---	---	---	---	---	----	---	---	---



OUTPUT: true.

Tc: O(n)

Sc: O(n)

NW: Try to solve without using PF.

Q. 5)

SCALER has a series of contests to help learners improve their coding skills. Each learner tries these contests to practice, but not everyone participates the same amount. **SCALER** wants to know which learner is participating the least so they can help them more or figure out how to get them more involved.

Problem

You have a list of id where each id represents a learner who tried a contest. A learner's number shows up more times if they try more contests. Your job is to write a program that looks at this list and finds out who has tried the **least number of contests**.

$$A[] = \{101, 102, 103, 101, 102, 101, 104, 105, 106, 105, 105\}$$

Learner Id	Contest Tried
101	3
102	2
103	1
104	1
105	3
106	1

Output: 3

▷ frequency map where key → learner id
value and value → Contest attempted.

- 2) find the minimum count.
- 3) no. of entry/key having that minimum count.

Code:

HashMap<Integer, Integer> hm = new HashMap<();

for (i=0 ; i < N; i++)
{

 if (hm.containskey(A[i]))
 {

 hm.put(A[i], hm.get(A[i]) + 1);

 else
 {

 hm.put(A[i], 1);

 }

```
int minCount = Integer.MAX_VALUE;  
int ans = 0
```

```
for (int tries: hm.values())
```

```
{  
    if (minCount > tries)
```

```
        minCount = tries;  
        ans = 1
```

```
    } else if (tries == minCount)  
    {  
        ans++;
```

```
return ans;
```

LearnerId

101

102

103

104

105

106

Context Tri

3

2

1

1

3

1



minCount = ~~MAXVAL~~ 221

ans = ~~X~~ 3

```
counts = {}
for id in A:
    if id in counts:
        counts[id] += 1
    else:
        counts[id] = 1

minCount = float('inf')
ans = 0
for tries in counts.values():
    if tries < minCount:
        minCount = tries
        ans = 1
    elif tries == minCount:
        ans += 1
```

Doubt Session

```
class Student {  
    int age;  
    String name;  
  
    void display(){  
        System.out.println("My name is " + this.name + ". I am " + this.age + " years old");  
    }  
  
    void sayHello(String name){  
        System.out.println(this.name + " says hello to " + name);  
    }  
}  
  
public class Client {  
    public static void main(String[] args) {  
        Student s1 = new Student();  
        s1.age = 10;  
        s1.name = "A";  
  
        Student s2 = new Student();  
        s2.age = 20;  
        s2.name = "B";  
  
        swap(s1, s2);  
  
        s1.display();  
    }  
}  
  
private static void swap(Student s1, Student s2) {  
    int tage = s1.age; tage = 10  
    s1.age = s2.age; s1.age = 20  
    s2.age = tage; s2.age = 10  
  
    String tname = s1.name; tname = A  
    s1.name = s2.name; s1.name = B  
    s2.name = tname; s2.name = A  
}
```

5K **6K**

The diagram illustrates the state of two objects, *s1* and *s2*, before and after they are swapped. It shows two boxes labeled **5K** and **6K**. Box **5K** contains initial values: *s1* has age=10 and name=A. Box **6K** contains initial values: *s2* has age=20 and name=B. Arrows point from the variable names to their respective boxes. After the swap operation, the values are updated: *s1* now has age=20 and name=B, while *s2* has age=10 and name=A. The final values are highlighted with red boxes.

```

public class Student {
    int age;
    String name;

    void display(){
        System.out.println("My name is " + this.name + ". I am " + this.age + " years old");
    }

    void sayHello(String name){
        System.out.println(this.name + " says hello to " + name);
    }
}

public class Client {
    public static void main(String[] args) {
        Student s1 = new Student();
        s1.age = 10;
        s1.name = "A";

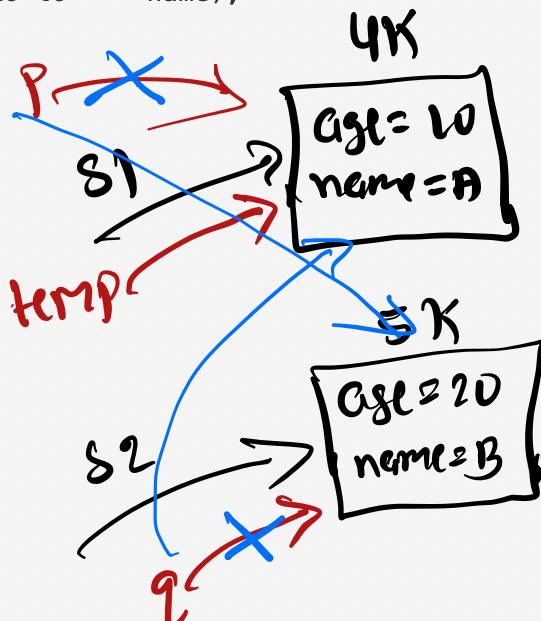
        Student s2 = new Student();
        s2.age = 20;
        s2.name = "B";

        swap(s1, s2);

        s1.display();
    }

    private static void swap(Student p, Student q) {
        Student temp = p; // temp = 4K
        p = q; // p -> q
        q = temp; // q -> p
    }
}

```



Shallow Copy

```
public class P {
    protected int d1 = 10;
    protected int d = 100;

    protected void fun1(){
        System.out.println("P's fun1");
    }

    protected void fun(){
        System.out.println("P's fun");
    }

    static protected void sfun(){
        System.out.println("P's sfun");
    }
}

public class C extends P{
    protected int d2 = 20;
    protected int d = 200;

    protected void fun2(){
        System.out.println("C's fun2");
    }

    protected void fun(){
        System.out.println("C's fun");
    }

    static protected void sfun(){
        System.out.println("C's sfun");
    }
}

public class Client {
    public static void main(String[] args){
        P obj = new C();
        System.out.println(obj.d1);
        System.out.println(obj.d);
        obj.fun1();
        obj.fun();
        obj.sfun();
    }
}
```

