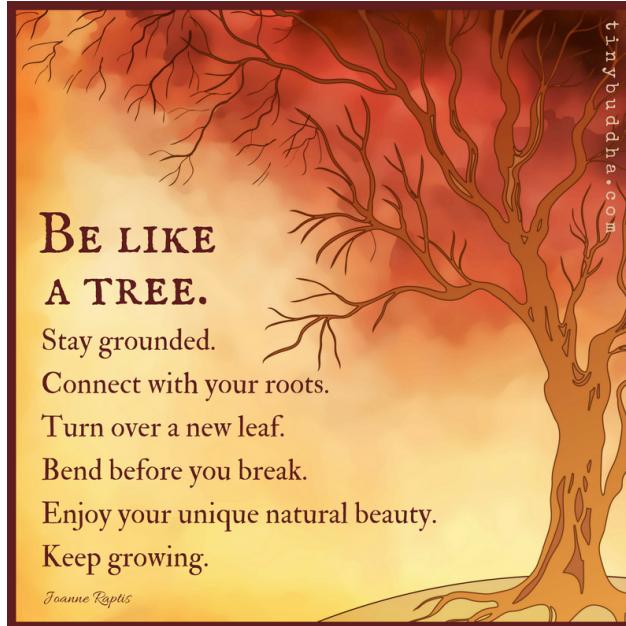


TREES 4 : LCA

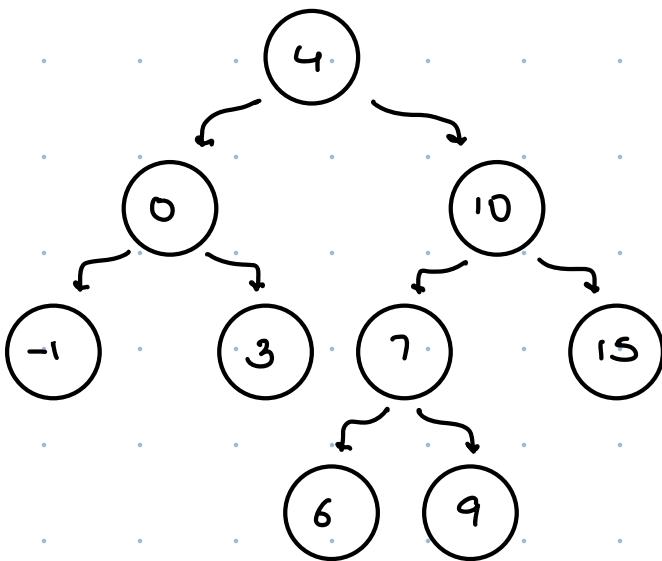


Good
Evening

Content

01. Kth Smallest Element in BST
02. Morris Inorder Traversal
03. LCA in Binary Tree
04. LCA in BST

Q1. Find Kth smallest element in BST



$$K = 3 \rightarrow 3$$

$$K = 5 \rightarrow 6$$

$$K = 8 \rightarrow 10$$

Idea → Do an inorder traversal on BST & store all elements in AL & return `al.get(k-1)`

```
main() {
    AL<I> ans = new AL<>();
    inorder (root, ans);
    return ans.get (k-1);
}
```

Tc : O(n)
Sc : O(n)

```
void inorder ( root, AL<I> ans) {
```

```
    if (root == null) return;
    inorder (root.left, ans);
    ans.add (root.val);
    inorder (root.right, ans);
```

Approach 2 → while doing inorder traversal, maintain a count

```
int count = 0  
int ans = -∞
```

```
void inorder (root , int k) {
```

```
    if (root == null) return;  
  
    inorder (root.left , k)  
    count = count + 1;  
  
    if (count == k)  
        ans = root.val;  
  
    return ;  
  
    if (ans == -∞) inorder (root.right , k)
```

Tc : O(n)
Sc : O(ht)

point left subtree
print root

Inorder Traversal point right subtree

Recursion

Sc : O(h)

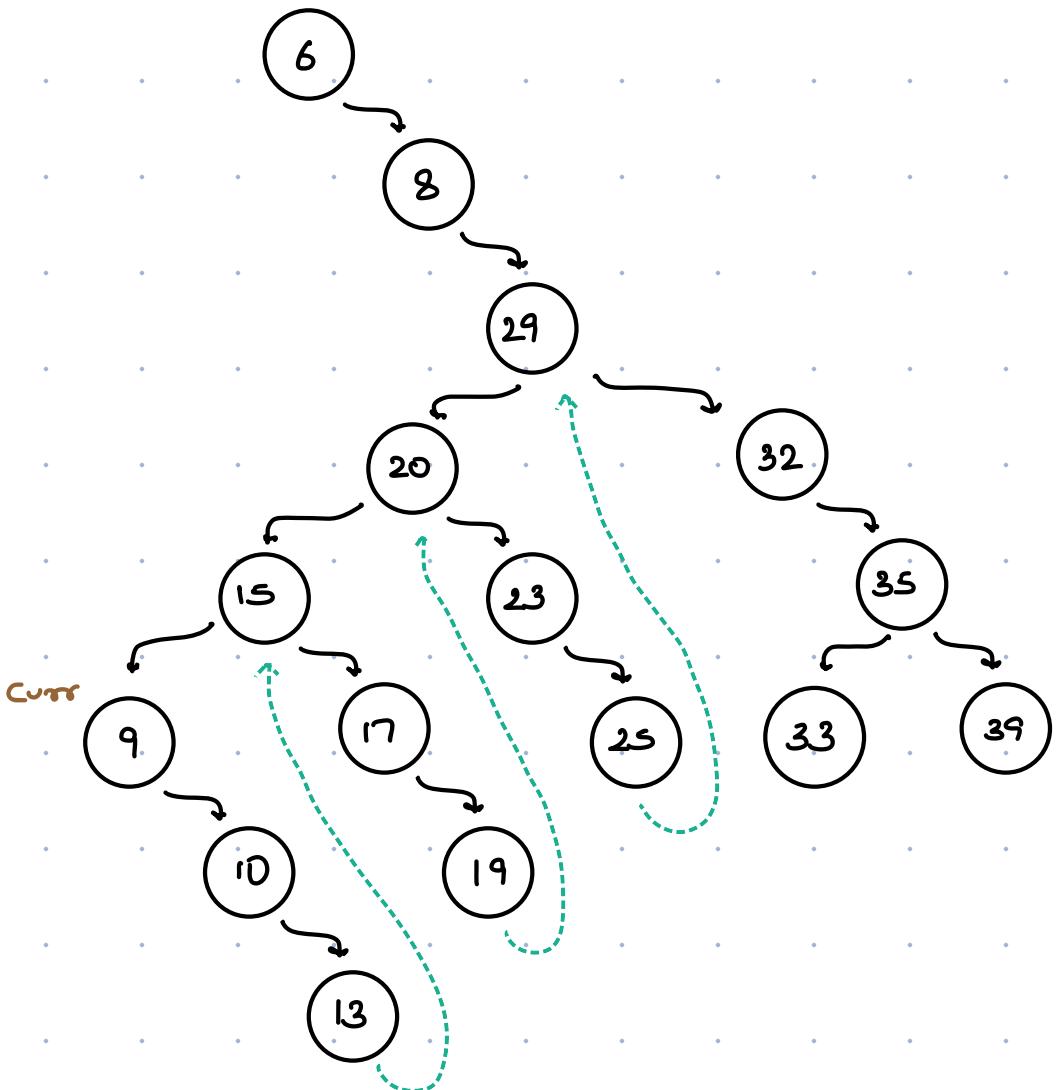
Iterative

Sc : O(n)

Morris Traversal

Sc : O(1)

Inorder = 6 8

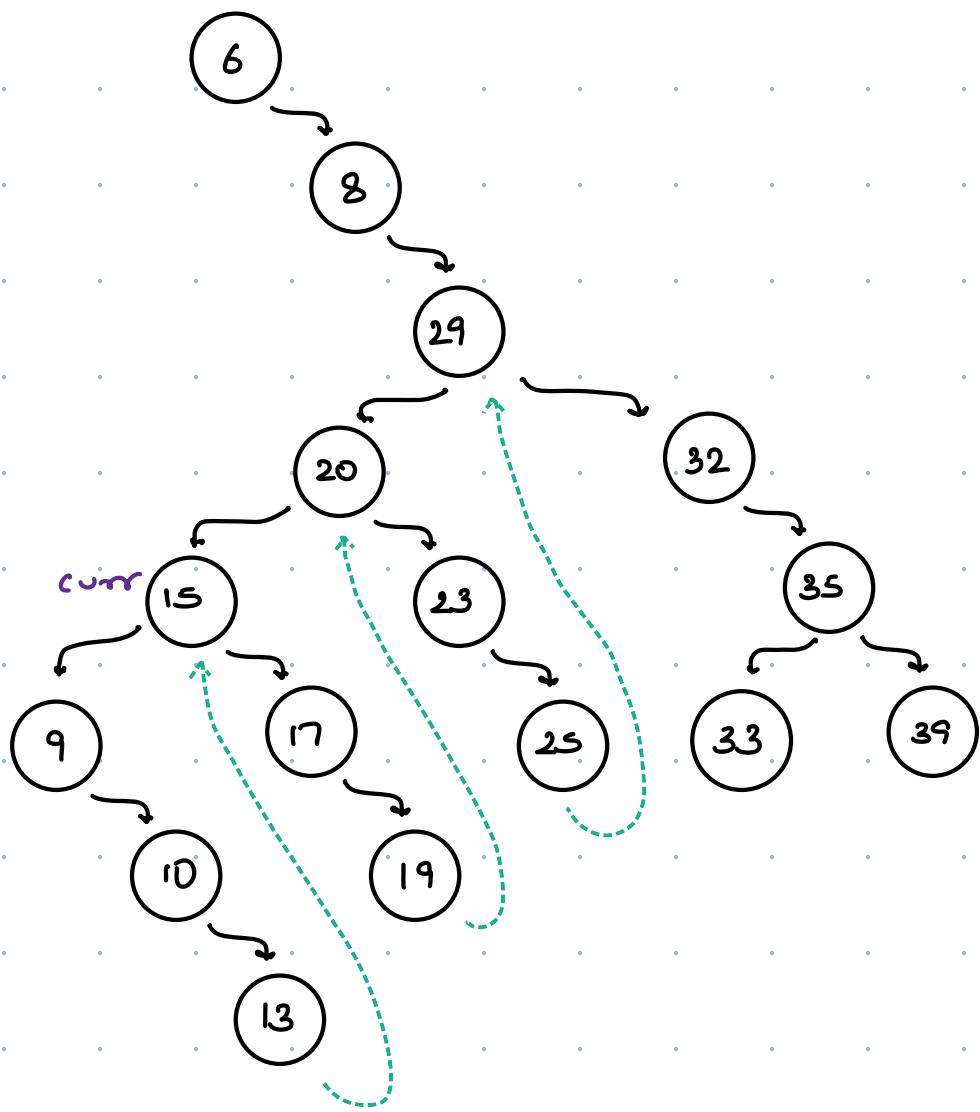


Obs = If (curr.left == null) → print curr;
 → Then move to RHS

Inorder predecessor

- Last node of left subtree to be printed OR
- right most child of left subtree OR
- Node in LST that doesn't right child

Inorder Traversal = 6 8 9 10 13



```

Node curr = root;

while (curr != null) {
    if (curr.left == null) {
        print(curr.data);
        curr = curr.right;
    }
    else {
        Node temp = curr.left;
        while (temp.right != null & temp.right != curr) {
            temp = temp.right;
        }
        if (temp.right == null) {
            temp.right = curr;
            curr = curr.left;
        }
        else if (temp.right == curr) {
            temp.right = null;
            print(curr.data);
            curr = curr.right;
        }
    }
}

```

TC: O(n)
SC: O(1)

create connection
break connection

Problem Statement

It is said that we all **humans** are related through some **common ancestor** at some point of time. Assume that a person can have **0, 1 or 2 children only**.

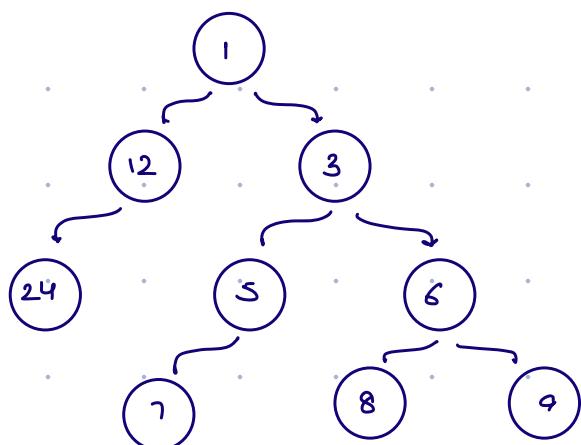
Given the Binary tree **A** representing the family tree, discover the earliest common family member who connects two given people **B** and **C** in a family tree.

Example

```

    Ram
   / \
Shyam Bob
 / \   \
Dholu Bholu Satish

-> Find LCA of Dholu and Bob
-> Answer = Ram
  
```



$$\text{LCA}(12, 5) \Rightarrow 1$$

$$\text{LCA}(8, 9) \Rightarrow 6$$

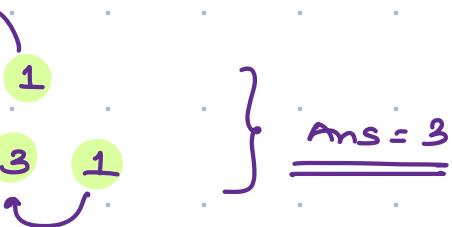
$$\text{LCA}(3, 8) \Rightarrow 3$$

Idea → Node to root path

$$\text{LCA}(5, 8)$$

$$\text{get NTR of } 5 = 5$$

$$\text{get NTR of } 8 = 8$$

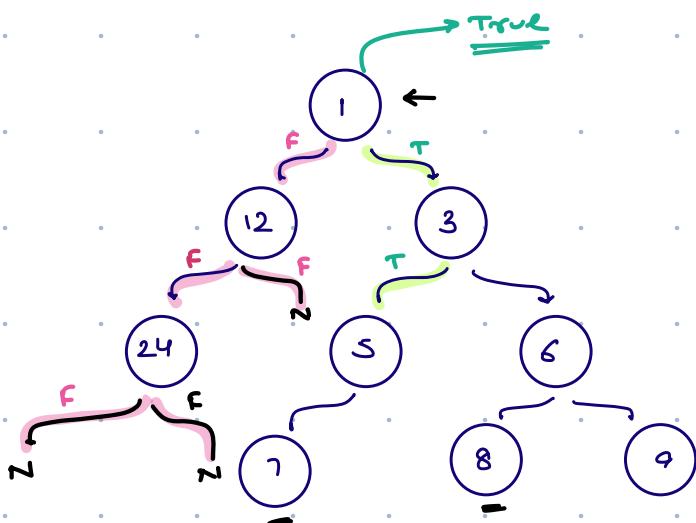


$$(5) \text{ RTN} = 1 \rightarrow \boxed{3} \rightarrow 5$$

$$(8) \text{ RTN} = 1 \rightarrow 3 \rightarrow 6 \quad 8$$

$$\} \quad \text{Ans} = 3$$

* Search for an element in Binary Tree



```
boolean search ( Node root, int k )
```

```
    if (root == null) return false;
```

```
    if (root.data == k) return true;
```

```
    if ( search (root.left, k) == true ) {
```

```
        return true;
```

```
    if ( search (root.right, k) == true ) {
```

```
        return true;
```

```
    return false;
```

```
}
```

* Node to Root Path

```
boolean search ( Node root, int k, ArrayList<Integer> ans )  
  
if (root == null) return false;  
  
if (root.data == k){  
    ans.add(root.data);  
    return true;  
}  
  
if (search (root.left, k) == true){  
    ans.add (root.data)  
    return true;  
}  
  
if (search (root.right, k) == true){  
    ans.add (root.data)  
    return true;  
}  
  
return false;
```

3

* LCA in Binary Tree for k_1 & k_2

01. Get NTR for k_1 → Reverse it

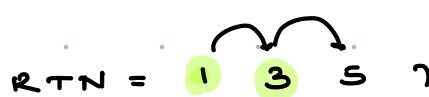
TC: $O(n)$

SC: $O(h)$

$O(\text{ht of tree})$

02. Get NTR for k_2 → Reverse it

$$k_1 = 7$$

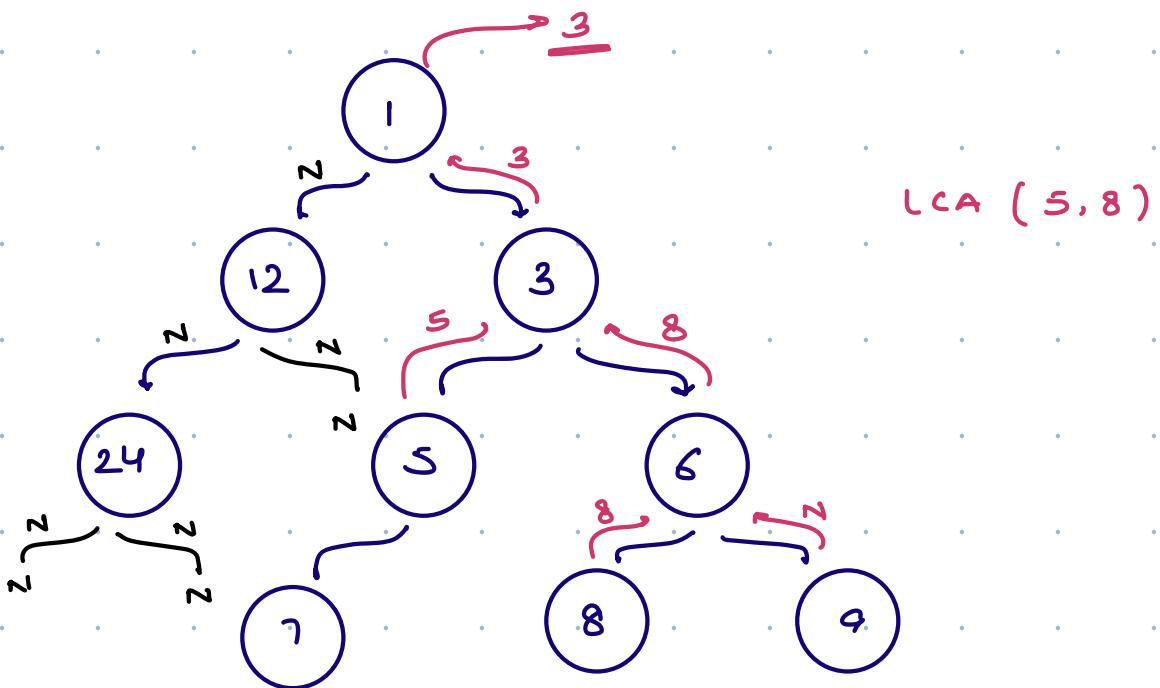


$$k_2 = 8$$



$$\text{Ans} = 3$$

03. Traverse from start & return last common node.



Node LCA (root, p, q)

```

if (root == null) return null;

if (root == p || root == q) return root;

l = LCA (root.left, p, q);
r = LCA (root.right, p, q);

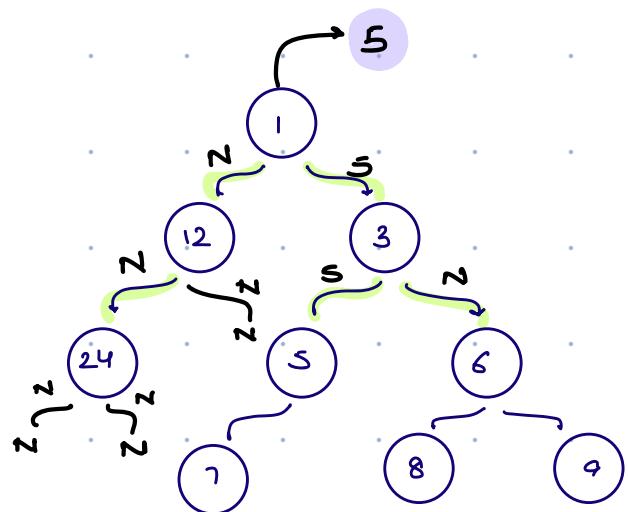
if (l != null & r != null) return root;

return l == null ? r : l;
    
```

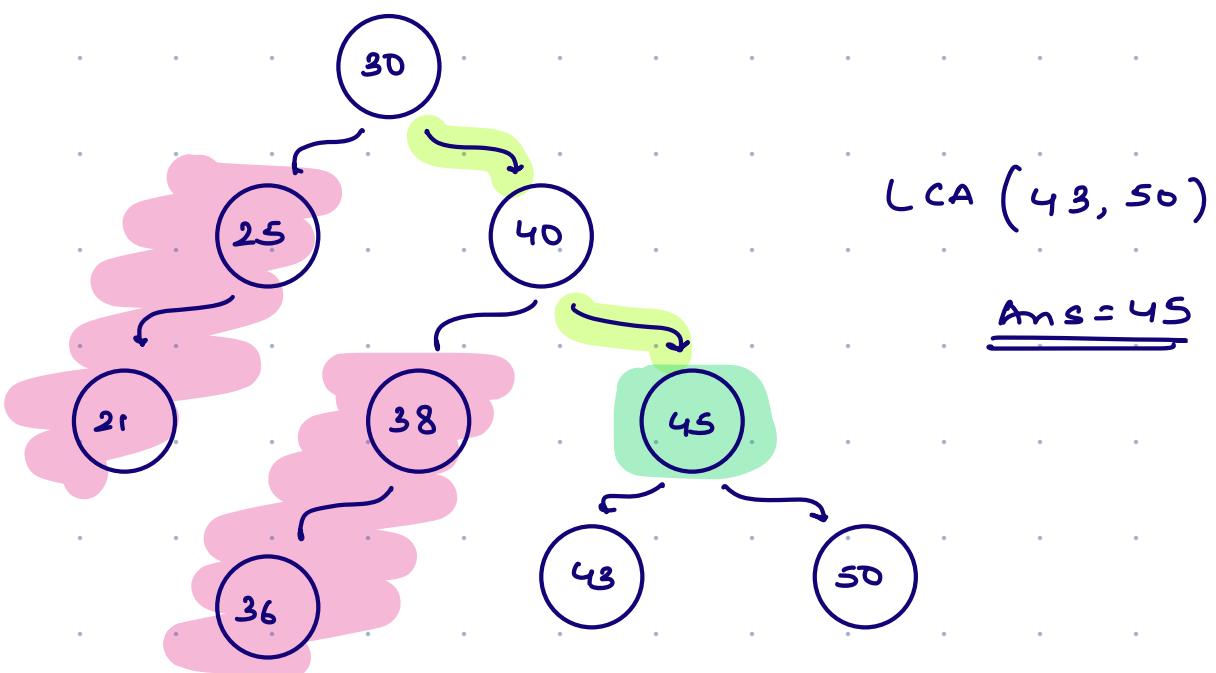
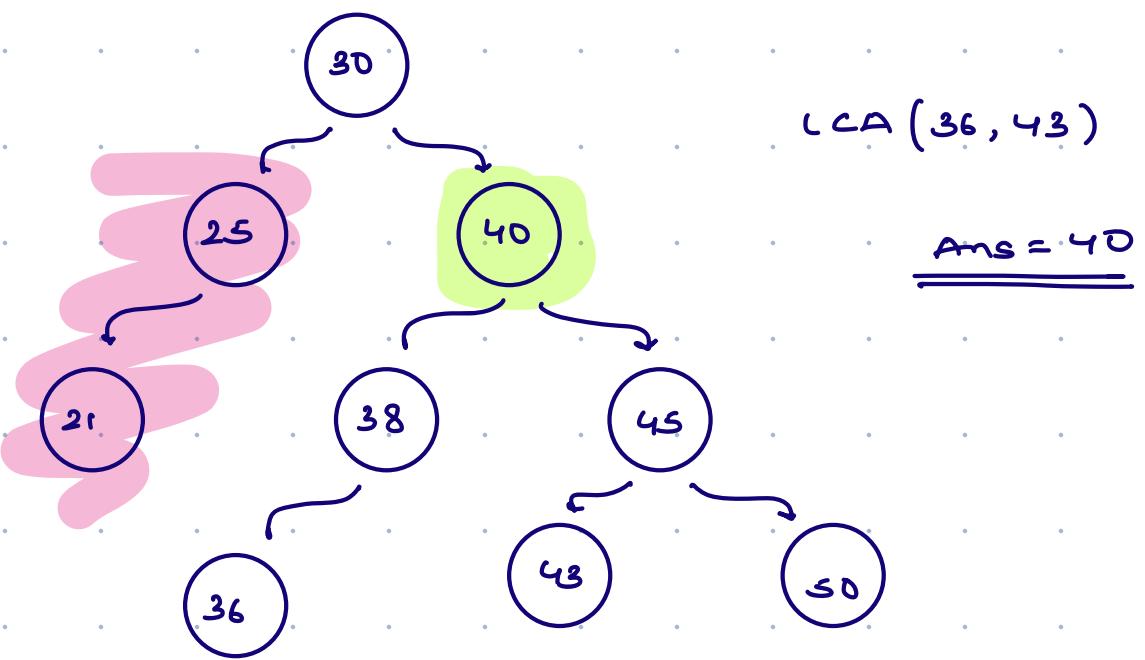
TC: O(n)
SC: O(ht)

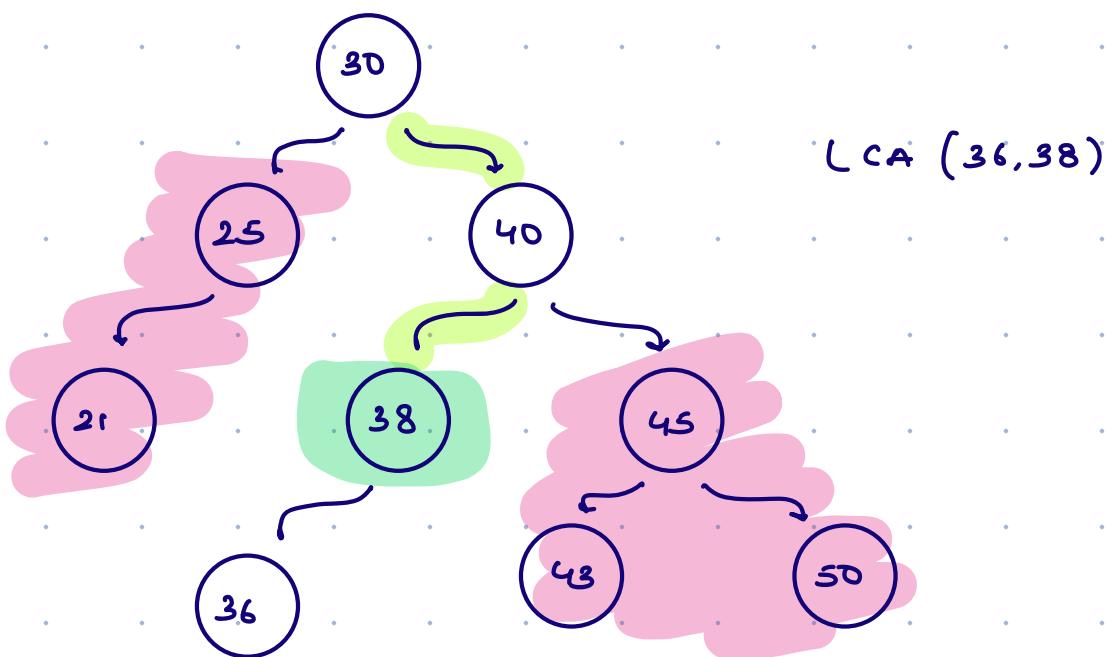
3

$p = 5$
 $q = 7$



* Lowest common Ancestor in Binary Search tree





* Code in whatsapp group (very easy)

longest substring of valid parenthesis



cnt : 1

Valid parentheses count

Problem Description

Alice is a mathematician and she loves to solve puzzles. One day, her friend Bob gave her a string **A** of size **N** consisting of opening and closing parentheses and asked her to find the length of the **longest** valid parentheses substring.

Note: A valid parentheses string means that each opening symbol has a corresponding closing symbol and the pairs of parentheses are properly nested.

```
public int solve(final String A) {  
    int n = A.length();  
    int ans = 0;  
    Stack < Integer > st = new Stack < > ();  
    st.push(-1);  
  
    for (int i = 0; i < n; i++) {  
        if (A.charAt(i) == '(') {  
            st.push(i);  
        }  
        else {  
            if (!st.empty()) {  
                st.pop();  
            }  
            if (!st.empty()) {  
                ans = Math.max(ans, i - st.peek());  
            }  
            if (st.empty()) {  
                st.push(i);  
            }  
        }  
    }  
    return ans;  
}
```



(() ()
o 1 2 3 4

(()) (()) (())

(()) (())