# Spell Checker

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## Spelling Correction

Spelling correction is an integral part of modern writing, ranging from **texting** and **emailing** to document creation and **web searches**. Despite their ubiquity, modern spell correctors aren't perfect, as evidenced by "autocorrect-gone-wrong" scenarios.



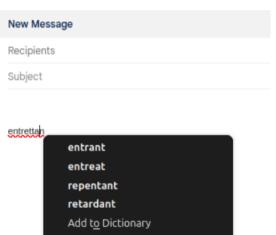


Figure 1: Spell checker.

## Applications of Spell Checking

- Text Writing
- Automated and Information Systems
  - Data Entry Systems
  - Search and Information Retrieval
  - Optical Character Recognition (OCR)
  - Chatbots
  - Translation Systems

### Automatic Spelling Correction Task

- 1 detection of an error;
- 2 generation of correction candidates;
- 3 ranking of candidate corrections;
- 4 perform automatic correction.

## Perspectives in Spelling Correction

- 1. Non-Word Spelling Correction
  - Detects and corrects errors where the **word does not exist** in the dictionary. Example:
    - Input: speling
    - Correction: spelling
- 2. Real-Word Spelling Correction
  - Detects and corrects errors where the **word exists** but is contextually wrong. Example:
    - Input: I no what to do.
    - Correction: I know what to do.

## Error Sources in Spelling

- Typographical Errors:
  - May change with input devices (physical or virtual keyboard, or OCR system) and environment conditions.
  - lacksquare Insertion: speeling o spelling
  - lacksquare Deletion: spelng ightarrow spelling
  - lacksquare Substitution: spolling o spelling
  - lacktriang Transposition: spelling o spelling
  - lacktriangle Diacritical marking: naive ightarrow naïve
- 2 Homophone Errors:
  - Homophones: their / there
  - Near-homophones: accept / except
- 3 Grammatical Errors:
  - among / between
- 4 Cross Word Boundary Errors:
  - maybe / may be

## Notable Algorithms and Tools

■ **Soundex** (1918): Phonetic algorithm that maps similar-sounding names.

```
Stephen \rightarrow S315, Perez \rightarrow P620, Juice \rightarrow J200, Robert \rightarrow R163
Steven \rightarrow S315, Powers \rightarrow P620, Juicy \rightarrow J200, Rupert \rightarrow R163
Stefan \rightarrow S315, Price \rightarrow P620, Juiced \rightarrow J230, Rubin \rightarrow R150
```

■ Shannon (1948): A Mathematical Theory of Communication.

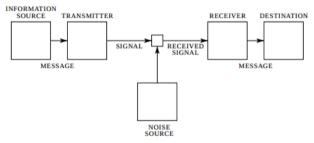


Fig. 1—Schematic diagram of a general communication system.

Figure 2: Noisy Channel.

■ Shannon (1950): Introduction of n-gram models in text analysis.



Figure 3: Word prediction is mean?

- Blair (1960): Early algorithm for spelling error correction.
  - Blair introduced the concept of similarity keys to group words based on their likelihood of being confused with one another.
  - r-letter abbreviation of an n-letter word
    - Information conveyed is inversely proportional to its a priori probability of occurrence.
    - $\blacksquare$  eliminate n-r letters in the order of their expected frequency
    - eliminate by their frequency of their occurrence as errors (best approach)
    - weight must also be given to the position of the letter in the word

#### EXAMPLE

A	$\mathbf{B}$	S (	ЭБ	В	$\mathbf{E}$	N '	$\Gamma$				$\mathbf{A}$	$\mathbf{B}$	S (	) F	В	A	N	$\mathbf{T}$
<b>5</b>	1	5	4	4	1	7	3	3.	Letter score	5	1	<b>5</b>	4	4	1	5	3	3
0	<b>2</b>	4	5	5	<b>5</b>	4	3	1	Position score	0	<b>2</b>	4	5	5	5	4	3	1
5	3	9	9	9	6	11	6	4	Sum of scores	5	3	9	9	9	6	9	6	4
		*	*	*		*	*		Delete			*	*	*		*	*	
					$\mathbf{A}$	$\mathbf{B}$	В	$\mathbf{T}$	Abbreviation	A	В	В	$\mathbf{T}$					

Figure 4: Example of Blair's algorithm.

■ Damerau-Levenshtein distance (1964, 1966): A string metric for measuring the edit distance between two sequences.

		s	u	n	d	a	у
	0	1	2	3	4	5	6
s	1	0	1 lelete	2	3	4	5
a	2	1	1	2	3	3	4
t	3	2	delete 2	2	3	4	4
u	4	3	2	3 repl	3 ace r	4 with r	5
r	5	4	3	3	4	4	5
d	6	5	4	4	3	4	5
а	7	6	5	5	4	3	4
у	8	7	6	6	5	4	3

Figure 5: How many operations does it take to turn Saturday into Sundat?

The Levenshtein distance between two strings a, b (of length |a| and |b| respectively) is given by

$$\mathsf{lev}(a,b) = \begin{cases} |a| & \text{if } |b| = 0, \\ |b| & \text{if } |a| = 0, \\ |\mathsf{lev}\left(\mathsf{tail}(a),\mathsf{tail}(b)\right) & \text{if } \mathsf{head}(a) = \mathsf{head}(b), \end{cases}$$

$$1 + \min \begin{cases} \mathsf{lev}\left(\mathsf{tail}(a),b\right) & \mathsf{deletion} \\ |\mathsf{lev}\left(\mathsf{tail}(b)\right) & \mathsf{insertion} \end{cases}$$

$$\mathsf{lev}\left(\mathsf{tail}(a),\mathsf{tail}(b)\right) & \mathsf{replacement} \end{cases}$$

Damerau-Levenshtein distance: also allows transposition of adjacent symbols.

Operations are expensive and language dependent: e.g. as of version 16.0, Unicode defines a total of 98682 Chinese characters.

- **BK-Trees** (1973): Efficient search for near matches using Levenshtein distance.
  - An arbitrary element *a* is selected as root node.
  - The k-th subtree is recursively built of all elements b such that d(a, b) = k.
  - Search idea: restrict the exploration of the tree to nodes that can only improve the best candidate found so far (use triangle inequality).

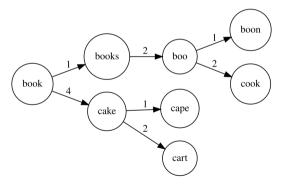
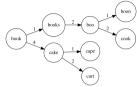


Figure 6: Burkhard-Keller Tree.

#### Search for w = 'cool'



- $d_u = d(w, u) = d('cool', 'book') = 2$ , set  $d_{best} = 2$ ;
- $v = \text{'books'}, |d_{uv} d_u| = |1 2| = 1 < d_{best}$ , then select v;
- $v = \text{'cake'}, |d_{\mu\nu} d_{\mu}| = |4 2| = 2 \not< d_{\text{best}}, \text{ do not select } v;$
- $d_u = d(w, u) = d('cool', 'books') = 3, d_u \not< d_{best};$
- $d_u = d(w, u) = d('cool', 'boo') = 2$ ,  $d_u \not< d_{best}$ ;
- $v = \text{'boon'}, |d_{uv} d_u| = |2 1| = 1 < d_{best}$ , then select v;
- $v = 'cook', |d_{uv} d_u| = |2 2| = 0 < d_{best}$ , then select v;
- $d_u = d(w, u) = d('cool', 'cook') = 1$ ,  $d_u < d_{best}$ , set  $d_{best} = 1$ ;
- $d_u = d(w, u) = d(\text{'cool'}, \text{'boon'}) = 2$ ,  $d_u \not< d_{\text{best}}$ ;
- 10 'cook' is returned as the answer with  $d_{\text{best}} = 1$ .

## ■ **SPELL** (Unix, 1975)

- Error detection only.
- Prefix and suffix removal (reduces the list below 1/3);
  - buzzed  $\rightarrow$  buzz, mapping  $\rightarrow$  map, possibly  $\rightarrow$  possible, antisocial  $\rightarrow$  social, metaphysics  $\rightarrow$  physics.
- Hashing (discarding 60% of the remaining bits);

Examples of hashing functions:

- 1 Shift-and-Add:  $h = (h \ll 1) + \text{char}\%m$
- 2 Multiplicative Hashing:  $h = (a \cdot h + \text{char})\%m$  (with a typically 31 or 33)
- 3 XOR-based Hashing:  $h = h \oplus (char << k)$
- Words were represented by 16-bit machine words;
- Bloom filter:
- False Positives.

### ■ Jaro similarity (1989)

The Jaro similarity  $sim_j$  of two given strings  $s_1$  and  $s_2$  is

$$sim_j = \left\{ egin{array}{ll} 0 & ext{if } m=0 \ rac{1}{3} \left( rac{m}{|s_1|} + rac{m}{|s_2|} + rac{m-t}{m} 
ight) & ext{otherwise} \end{array} 
ight.$$

where:

- $\blacksquare$   $|s_i|$  is the length of the string  $s_i$ ;
- *m* is the number of "matching characters" (see below);
- lacktriangledown t is the number of "transpositions" (see below).

Jaro similarity score is 0 if the strings do not match at all, and 1 if they are an exact match. In the first step, each character of  $s_1$  is compared with all its matching characters in  $s_2$ . Two characters from  $s_1$  and  $s_2$  respectively, are considered **matching** only if they are the same and not farther than  $\left\lfloor \frac{\max(|s_1|,|s_2|)}{2} \right\rfloor - 1$  characters apart. **Transposition** is the number of matching characters that are not in the right order divided by two.

### ■ Jaro-Winkler similarity (1990)

- Introduces Winkler modification.
- Prefix Length  $\ell$ : if two strings share a common prefix, they are likely to be more similar.
- scale factor p: enhances the Jaro similarity score based on the length of the common prefix (usually set to 0.1 and should not exceed 0.25).

$$sim_w = sim_i + \ell p(1 - sim_i)$$

- **Metaphone** (1990), Double Metaphone (2000), Metaphone 3 (2009): Extracts phonetic information for better matching.
  - Set of rules to improves on the Soundex algorithm.
  - Smith → SMO, [SMO, XMT], Schmidt → SXMTT, [XMT, SMT],
  - Taylor → TLR, [TLR], Taylor → EFNS, [AFNS],
  - Roberts → RBRTS, [RPRTS]
  - spelling → SPLNK, [SPLNK], speling → SPLNK, [SPLNK], speeling → SPLNK, [SPLNK], sprlling → SPRLNK, [SPRLNK]

■ Noisy Channel Model (Kernighan et al., 1990 and Mays et al., 1991): Combined prior and likelihood models.

In the noisy channel model, we imagine that the surface form we see is actually a "distorted" form of an original word passed through a noisy channel. The decoder passes each hypothesis through a model of this channel and picks the word that best matches the surface noisy word. (Jurafsky and Martin, 2024)

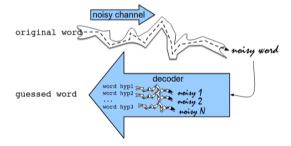


Figure 7: Noisy Channel Model

This noisy channel model is a kind of Bayesian inference.

Out of all possible words in the vocabulary V we want to find the word w such that P(w|x) is highest.

$$\hat{w} = \operatorname*{arg\,max}_{w \in V} P(w|x)$$

Using Bayes: P(x, w) = P(w|x)P(x) = P(x|w)P(w),

$$\hat{w} = \underset{w \in V}{\operatorname{arg \, max}} \frac{P(x|w)P(w)}{P(x)} = \underset{w \in V}{\operatorname{arg \, max}} \underbrace{P(x \mid w)P(w)}_{\text{channel model or likelihood}}$$

$$\hat{w} = \underset{w \in V}{\operatorname{arg\,max}} \left( \log P(x \mid w) + \log P(w) \right)$$

### **function** NOISY CHANNEL SPELLING(word x, dict D, lm, editprob) **returns** correction

```
if x \notin D

candidates, edits \leftarrow All strings at edit distance 1 from x that are \in D, and their edit

for each c, e in candidates, edits

channel \leftarrow editprob(e)

prior \leftarrow \operatorname{lm}(c)

score[c] = \log channel + \log prior

return \operatorname{argmax}_c score[c]
```

Figure 8: Noisy channel model for spelling correction for unknown words (Jurafsky and Martin, 2024).

### Example

original word

actress cress caress access across acres



Figure 9: Example: misspelling acress.

		J			
		Correct	Error	Position	
Error	Correction	Letter	Letter	(Letter #)	Type
acress	actress	t	_	2	deletion
acress	cress	_	a	0	insertion
acress	caress	ca	ac	0	transposition
acress	access	С	r	2	substitution
acress	across	О	e	3	substitution
acress	acres	_	S	5	insertion
acress	acres	_	S	4	insertion

Figure 10: Candidate corrections for the misspelling acress and the transformations that would have produced the error (after Kernighan et al. (1990)). "—" represents a null letter. (Jurafsky and Martin, 2024)

w	count(w)	p(w)
actress	9,321	.0000231
cress	220	.000000544
caress	686	.00000170
access	37,038	.0000916
across	120,844	.000299
acres	12,874	.0000318

Figure 11: Language model from the 404,253,213 words in the Corpus of Contemporary English (COCA) (Jurafsky and Martin, 2024).

#### Error model

- A perfect model would need all sorts of factors: who the typist was, whether the typist was left-handed or right-handed, and so on.
- We can get a pretty reasonable estimate of P(x|w) just by looking at **local context**: the identity of the correct letter itself, the misspelling, and the surrounding letters.
- Confusion Matrices:
  - del[x, y]: count(xy typed as x)
  - ins[x, y]: count(x typed as xy)
  - sub[x, y]: count(x typed as y)
  - trans[x, y]: count(xy typed as yx)

					S	ub[]	X, Y	] =	Sub	stitı	ıtio					ect) i	for	Y (	corr	ect)						
X												Y	(co	rrect	)											
	a	<u>b</u>	С	d	c	f	g	h	i	j	k	1	m	n	0	p	q	r	S	t	u	v	w	X	уу	Z
a	0	0	7	1	342	0	0	2	118	0	1	0	0	3	76	0	0	1	35	9	9	0	1	0	5	0
b	0	0	9	9	2	2	3	1	0	0	0	5	11	5	0	10	0	0	2	1	0	0	8	0	0	0
c	6	5	0	16	0	9	5	0	0	0	1	0	7	9	1	10	2	5	39	40	1	3	7	1	1	0
d	1	10	13	0	12	0	5	5	0	0	2	3	7	3	0	1	0	43	30	22	0	0	4	0	2	0
c	388	0	3	11	0	2	2	0	89	0	0	3	0	5	93	0	0	14	12	6	15	0	1	0	18	0
f	0	15	0	3	1	0	5	2	0	0	0	3	4	1	0	0	0	6	4	12	0	0	2	0	0	0
g	4	1	11	11	9	2	0	0	0	1	1	3	0	0	2	1	3	5	13	21	0	0	1	0	3	0
h	1	8	0	3	0	0	0	0	0	0	2	0	12	14	2	3	0	3	1	11	0	0	2	0	0	0
í	103	0	0	0		0	1	0	0	0	0	6	0	0	49	0	0	0	2	1	47	0	2	1	15	0
j	0	1	1	9	0	0	1	0	0	0	0	2	1	0	0	0	0	0	5	0	0	0	0	0	0	0
k	1	2	8	4	1	1	2	5	0	0	0	0	5	0	2	0	0	0	6	0	0	0	. 4	0	0	3
1	2	10	1	4	0	4	5	6	13	0	1	0	0	14	2	5	0	11	10	2	0	0	0	0	0	0
m	1	3	7	8	0	2	0	6	0	0	4	4	0	180	0	6	0	0	9	15	13	3	2	2	3	0
n	2	7	6	5	3	0	1	19	1	0	4	35	78	0	0	.7	0	28	5	. 7	0	0	1	2	0	2
0	91	1	1	3	116	0	0	0	25	0	2	0	0	0	.0	14	0	2	4	14	39	0	0	0	18	0
P	0	11	1	2	0	6	5	0	2	9	0	2	7	6	15	0	0	1	3	6	0	4	1	0	0	0
q	0	0	1	0	0	0	27	0	0	0	Ű	0	0	0	0	0	0	0	.0	0	0	0	0	0	0	0
r	0	14	0	30	12	2	2	8	2	0	5	8	4	20	1	14	0	0	12	22	4	0	0	1	0	0
s	11	8	27	33	35	4	0	1	0	1	0	27	0	6	1	1	0	14	0	15	0	0	. 5	3	20	1
t	3	4	9	42	7	5	19	5	0	1	0	14	9	2	43	6	0	11	37	0	0	2	19	0	7	6
u	20	0	0	0	44	0	0	0	64	0	0	0	0	-	43	0	0	4	0	0	0	0	2	0	8	0
v	0	0		0	0	3	0	0	0	0	0	1	0	0	1	0	0	0	8	3	0	0	0	0	0	0
w	2	2	1	0	1	0	0	2	0	0	1	0	0	0	0	7	0	0	3	3	1	0	0	0	0	0
x	0	0	0	2	10	0	0	0	0	0	0	0	0	0	0	0	0	0	9	0	õ	0	0	0	0	0
У	0	0	2	0	15	0	0	0	15	0	0	0	2	0	6	1	0	2	36	8	3	0	0	1	0	0

Figure 12: Confusion matrix for spelling errors (Kernighan et al., 1990).

## Estimating the channel model

$$P(x|w) = \begin{cases} \frac{\text{del}[x_{i-1}, w_i]}{\text{count}[x_{i-1}w_i]}, & \text{if deletion} \\ \frac{\text{ins}[x_{i-1}, w_i]}{\text{count}[w_{i-1}]}, & \text{if insertion} \\ \frac{\text{sub}[x_i, w_i]}{\text{count}[w_i]}, & \text{if substitution} \\ \frac{\text{trans}[w_i, w_{i+1}]}{\text{count}[w_i, w_{i+1}]}, & \text{if transposition} \end{cases}$$

Candidate	Correct	Error		
Correction	Letter	Letter	$\mathbf{x} \mathbf{w}$	P(x w)
actress	t	-	c ct	.000117
cress	-	a	a #	.00000144
caress	ca	ac	ac ca	.00000164
access	С	r	r c	.000000209
across	0	e	e o	.0000093
acres	-	S	es e	.0000321
acres	-	S	ss s	.0000342

Figure 13: Channel model for acress; the probabilities are taken from the del[], ins[], sub[], and trans[] confusion matrices as shown in Kernighan et al. (1990).

Final probabilities for each of the potential corrections

Correct	Error				
Letter	Letter	$\mathbf{x} \mathbf{w}$	P(x w)	P(w)	$10^9 *P(\mathbf{x} \mathbf{w})P(\mathbf{w})$
t	-	c ct	.000117	.0000231	2.7
-	a	a #	.00000144	.000000544	0.00078
ca	ac	ac ca	.00000164	.00000170	0.0028
С	r	r c	.000000209	.0000916	0.019
0	e	elo	.0000093	.000299	2.8
-	S	es e	.0000321	.0000318	1.0
-	S	ss s	.0000342	.0000318	1.0
	t - ca c o -	t - a ca ac c r o e - s	Letter         Letter         x w           t         -         c ct           -         a a #           ca         ac ca           c         r r c           o         e e o           -         s es e	Letter         Letter         x w         P(x w)           t         -         c ct         .000117           -         a #         .00000144           ca         ac ca         .00000164           c         r c         .000000209           o         e o         .00000321	Letter         Letter         x w         P(x w)         P(w)           t         -         c ct         .000117         .0000231           -         a  #         .00000144         .000000544           ca         ac ca         .00000164         .00000170           c         r c         .000000209         .0000916           o         e o         .0000093         .000299           -         s e e o         .0000321         .0000318

Figure 14: Computation of the ranking for each candidate correction, using the language model shown earlier and the error model. The final score is multiplied by  $10^9$  for readability (Jurafsky and Martin, 2024).



Unfortunately, the algorithm was wrong here; the writer's intention becomes clear from the context: ... was called a "stellar and versatile **acress** whose combination of sass and glamour has defined her ...". The surrounding words make it clear that actress and not across was the intended word. (Jurafsky and Martin, 2024)

Using the *Corpus of Contemporary American English* to compute bigram probabilities for the words *actress* and *across* in their context using add-one smoothing, we get the following probabilities:

$$P(\text{actress}|\text{versatile}) = .000021$$
  
 $P(\text{across}|\text{versatile}) = .000021$   
 $P(\text{whose}|\text{actress}) = .0010$   
 $P(\text{whose}|\text{across}) = .000006$ 

Multiplying these out gives us the language model estimate for the two candidates in context:

$$P(\text{versatile actress whose}) = .000021 \times .0010 = 210 \times 10^{-10}$$
  
 $P(\text{versatile across whose}) = .000021 \times .000006 = 1 \times 10^{-10}$ 

Spell Checker

Jurafsky, D., & Martin, J. H. (2024). Speech and Language Processing.

Kernighan, M. D. et al. (1990). A spelling correction program based on a noisy channel model.

Mays, E. et al. (1991). Context based spelling correction.

- Noisy Channel Model
  - Correct (Unix, 1990): Takes inputs from SPELL rejected words and provides candidates. Operations: Insertion, Deletion, Substitution, Reversal. Uses error probabilities.

- **Brill-Moore channel model** (2000): String to string edits.
  - Let  $\Sigma$  be an alphabet, the model allows all edit operations of the form  $\alpha \to \beta$ , where  $\alpha, \beta \in \Sigma^*$ .
  - $P(\alpha \to \beta)$  is the probability that when users intends to type  $\alpha$  and they typed  $\beta$  instead.
  - $\blacksquare$   $P(\alpha \to \beta | PNS)$  probability conditioned by the position on the string
    - $\blacksquare$  P(e | a) does not vary greatly with position.
    - P(ent | ant) is highly dependent upon position.
    - People rarely mistype antler as entler, but often mistype reluctant as reluctent.

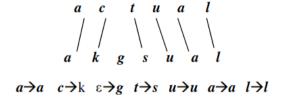


Figure 15: String alignment (Brill and Moore, 2000).

Brill, E. and Moore, R. C. (2000). *An Improved Error Model for Noisy Channel Spelling Correction*.

- **Aspell** (2000): Combines spelling and phonetic correction.
  - Hashing for Spell Checking: Efficient candidate lookup using hash tables.
  - Metaphone Algorithm: Handles phonetic corrections by matching words that sound similar.
  - Ispell's Near Miss Strategy:
    - Focuses on edit distance 1 to reduce the search space.
    - Early Dictionary Filtering: Prunes invalid candidates during generation.

# Example - Handling Homophones in Aspell

- Misspelled word ther
- Candidates: there, their, they're
- 1 Metaphone
  - The Metaphone algorithm transforms words into phonetic codes based on pronunciation.
  - Phonetic codes for the candidate words:
    - $\blacksquare$  there  $\rightarrow$  OR
    - lacksquare their ightarrow OR
    - lacktriangle they're ightarrow OR
  - Homophones share the same code (OR).

# 2 Workflow:

- Input: Misspelled word ther.
- Step 1: Generate candidates using **edit distance 2 or less**:
  - Candidates: there, their, thee, thor, her, the, they're.
- Step 2: Compute Metaphone codes for all candidates:
  - Candidates phonetically similar to ther (OR) rank higher: there, their, thor, they're.
- Step 3: Rank and suggest based on:
  - Word frequency, Edit distance, Phonetic Similarity, Error Likelihood.

#### 3 Limitations:

- Metaphone matches words by sound but lacks contextual understanding.
- Example:
  - Input: "I went to ther house."
  - Suggestions: thee, their, there, therm, the, her, Thar, Thea, Thor, Thur.
  - Aspell cannot infer the correct word (their) without considering the sentence's context.

## **GNU** Aspell

■ Hunspell (2002): Morphological analyzer with affix rules and phonetic matching.

#### Key Features:

- Morphological Analysis:
  - Supports complex languages with rich morphology (e.g., Hungarian, Turkish, Finnish).
  - Handles word roots, prefixes, and suffixes using affix rules.
- Dictionary System:
  - Two components:
    - 1 Dictionary File: Contains root forms of words.
    - 2 Affix File: Defines rules for combining roots with prefixes/suffixes.
- Levenshtein Distance:
  - Uses *edit distance* to generate and rank candidate corrections.
- Phonetic Matching:
  - Uses a table-driven phonetic transcription algorithm borrowed from Aspell. It is useful for languages with not pronunciation based orthography.
- n-gram similarity:
  - Improve suggestions.

- Multilingual Support:
  - Available for 98 languages with extensive dictionaries.

## Applications:

- Integrated into tools like LibreOffice, Firefox, and Chrome for multilingual spell checking.
- Supports custom dictionaries for specialized fields (e.g., medical, legal).

Hunspell at GitHub

■ Norvig's Algorithm (2007): Uses Damerau-Levenshtein distance to generate candidates.

## Key Features:

- Eit Distance:
  - Generates all possible words within a given edit distance (e.g., 1 or 2) from the misspelled word.
  - Handles insertion, deletion, substitution, and transposition.
- Dictionary Lookup:
  - Filters candidates by validating them against a word dictionary.
- Ranking:
  - Ranks valid candidates based on:
    - Word Frequency: More frequent words are prioritized.
    - Likelihood of Errors: Based on the Noisy Channel Model (optional).

How to Write a Spelling Corrector

■ QWERTY Weighted Levenshtein Distance: takes keyboard distance into account.

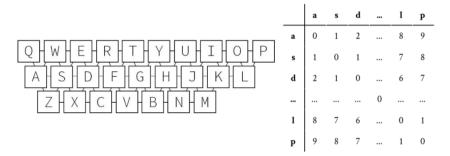


Figure 16: QWERTY keyboard and keyboard distance matrix.

■ Distance between keys are in [0,9]. They are multiplied by 2/9.

- Deletion: weighted by the average of the distances to the adjacent characters in the string.
- Insertion: unchanged, weight 1.
- Substitution: weighted according to the distance between the character that is removed and the character that is inserted.
- Transposition: unchanged, weight 1.

Samuelsson, 2017

- Neural-Based Models: Leverage deep learning for advanced error detection and correction.
  - Utilize deep learning techniques to improve spellchecking:
    - Recurrent Neural Networks (RNNs)
    - Word Embeddings
    - Transformers
  - Contextual Awareness
  - Learning from Data
  - Handling Typos

# Examples:

- Google's Smart Compose
- Grammarly
- Microsoft Editor
- LanguageTool

# **Lexical Similarity Metrics**

**1 Levenshtein Distance** (Edit Distance):

Minimal number of insertions, deletions, and replacements to transform one word into another.

2 Jaro Similarity:

Measures similarity based on matching characters and transpositions.

**3 Keyboard Distance:** 

Considers physical proximity of keys.

4 Phonetic Matching:

Algorithms like Soundex and Metaphone to identify similar-sounding words.

# Domain-Specific Spell Checkers

- 1 Medical
  - MedSpell: a medical spelling and autocorrect application
  - OpenMedSpel (open-source)
- 2 Programming
  - CodeSpell: designed primarily for checking misspelled words in source code
- 3 Learning
  - Kidspell: A child-oriented, rule-based, phonetic spellchecker
- 4 Accessibility
  - Real Check: A Spellchecker for Dyslexia
- 5 Custom Dictionaries
  - Hunspell and Aspell: Add specialized vocabularies

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