Spell Checker

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Spelling Correction

Spelling correction is an integral part of modern writing, ranging from **texting** and **emailing** to document creation and **web searches**. Despite their ubiquity, modern spell correctors aren't perfect, as evidenced by "autocorrect-gone-wrong" scenarios.



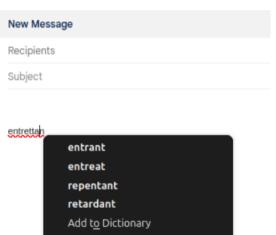


Figure 1: Spell checker.

Applications of Spell Checking

- Text Writing
- Automated and Information Systems
 - Data Entry Systems
 - Search and Information Retrieval
 - Optical Character Recognition (OCR)
 - Chatbots
 - Translation Systems

Automatic Spelling Correction Task

- 1 detection of an error;
- 2 generation of correction candidates;
- 3 ranking of candidate corrections;
- 4 perform automatic correction.

Perspectives in Spelling Correction

- 1. Non-Word Spelling Correction
 - Detects and corrects errors where the **word does not exist** in the dictionary. Example:
 - Input: speling
 - Correction: spelling
- 2. Real-Word Spelling Correction
 - Detects and corrects errors where the **word exists** but is contextually wrong. Example:
 - Input: I no what to do.
 - Correction: I know what to do.

Error Sources in Spelling

- Typographical Errors:
 - May change with input devices (physical or virtual keyboard, or OCR system) and environment conditions.
 - lacksquare Insertion: speeling o spelling
 - lacksquare Deletion: spelng o spelling
 - lacksquare Substitution: spolling o spelling
 - lacksquare Transposition: spelling ightarrow spelling
 - lacktriangle Diacritical marking: naive ightarrow naïve
- 2 Homophone Errors:
 - Homophones: their / there
 - Near-homophones: accept / except
- 3 Grammatical Errors:
 - among / between
- 4 Cross Word Boundary Errors:
 - maybe / may be

Notable Algorithms and Tools

■ **Soundex** (1918): Phonetic algorithm that maps similar-sounding names.

```
Stephen \rightarrow S315, Perez \rightarrow P620, Juice \rightarrow J200, Robert \rightarrow R163
Steven \rightarrow S315, Powers \rightarrow P620, Juicy \rightarrow J200, Rupert \rightarrow R163
Stefan \rightarrow S315, Price \rightarrow P620, Juiced \rightarrow J230, Rubin \rightarrow R150
```

■ Shannon (1948): A Mathematical Theory of Communication.

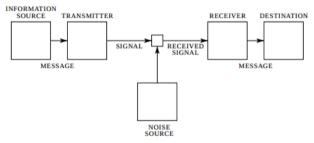


Fig. 1—Schematic diagram of a general communication system.

Figure 2: Noisy Channel.

■ Shannon (1950): Introduction of n-gram models in text analysis.



Figure 3: Word prediction is mean?

- Blair (1960): Early algorithm for spelling error correction.
- Damerau—Levenshtein distance (1964, 1966): A string metric for measuring the edit distance between two sequences.

		s	u	n	d	а	у
	0	1	2	3	4	5	6
s	1	0	1 lelete	2	3	4	5
а	2	1	1	2	3	3	4
t	3	2	delete 2	2	3	4	4
u	4	3	2	3 repl	3 ace r	4 with r	5
r	5	4	3	3	4	4	5
d	6	5	4	4	3	4	5
а	7	6	5	5	4	3	4
у	8	7	6	6	5	4	3

Figure 4: How many operations does it take to turn Saturday into Sundat?

The Levenshtein distance between two strings a, b (of length |a| and |b| respectively) is given by

$$\mathsf{lev}(a,b) = \begin{cases} |a| & \text{if } |b| = 0, \\ |b| & \text{if } |a| = 0, \\ |\mathsf{lev}\left(\mathsf{tail}(a),\mathsf{tail}(b)\right) & \text{if } \mathsf{head}(a) = \mathsf{head}(b), \end{cases}$$

$$1 + \min \begin{cases} \mathsf{lev}\left(\mathsf{tail}(a),b\right) & \mathsf{deletion} \\ |\mathsf{lev}\left(\mathsf{tail}(b)\right) & \mathsf{insertion} \end{cases}$$

$$\mathsf{lev}\left(\mathsf{tail}(a),\mathsf{tail}(b)\right) & \mathsf{replacement} \end{cases}$$

Damerau-Levenshtein distance: also allows transposition of adjacent symbols.

Operations are expensive and language dependent: e.g. as of version 16.0, Unicode defines a total of 98682 Chinese characters.

- **BK-Trees** (1973): Efficient search for near matches using Levenshtein distance.
 - An arbitrary element *a* is selected as root node.
 - The k-th subtree is recursively built of all elements b such that d(a, b) = k.
 - Search idea: restrict the exploration of the tree to nodes that can only improve the best candidate found so far (use triangle inequality).

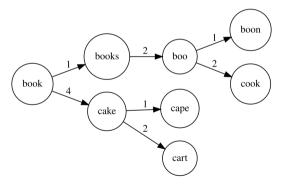
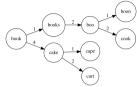


Figure 5: Burkhard-Keller Tree.

Search for w = 'cool'



- $d_u = d(w, u) = d('cool', 'book') = 2$, set $d_{best} = 2$;
- $v = \text{'books'}, |d_{uv} d_u| = |1 2| = 1 < d_{best}$, then select v;
- $v = \text{'cake'}, |d_{\mu\nu} d_{\mu}| = |4 2| = 2 \not< d_{\text{best}}, \text{ do not select } v;$
- $d_u = d(w, u) = d('cool', 'books') = 3, d_u \not< d_{best};$
- $d_u = d(w, u) = d('cool', 'boo') = 2$, $d_u \not< d_{best}$;
- $v = \text{'boon'}, |d_{uv} d_u| = |2 1| = 1 < d_{best}$, then select v;
- $v = 'cook', |d_{uv} d_u| = |2 2| = 0 < d_{best}$, then select v;
- $d_u = d(w, u) = d('cool', 'cook') = 1$, $d_u < d_{best}$, set $d_{best} = 1$;
- $d_u = d(w, u) = d(\text{'cool'}, \text{'boon'}) = 2, d_u \not< d_{\text{best}};$
- 10 'cook' is returned as the answer with $d_{\text{best}} = 1$.

Spell Checker

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- Hashing (discarding 60% of the remaining bits);

Examples of hashing functions:

- 1 Shift-and-Add: $h = (h \ll 1) + \text{char}\%m$
- 2 Multiplicative Hashing: $h = (a \cdot h + \text{char})\%m$ (with a typically 31 or 33)
- **3** XOR-based Hashing: $h = h \oplus (\text{char} << k)$
- Words were represented by 16-bit machine words;

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- Words were represented by 16-bit machine words;
- Bloom filter:
- False Positives.

- **Metaphone** (1990), Double Metaphone (2000), Metaphone 3 (2009): Extracts phonetic information for better matching.
 - Set of rules to improves on the Soundex algorithm.
 - Smith → SMO, [SMO, XMT], Schmidt → SXMTT, [XMT, SMT],
 - Taylor → TLR, [TLR], Taylor → EFNS, [AFNS],
 - Roberts → RBRTS, [RPRTS]
 - spelling → SPLNK, [SPLNK], speling → SPLNK, [SPLNK], speeling → SPLNK, [SPLNK], sprlling → SPRLNK, [SPRLNK]

■ Noisy Channel Model (Kernighan et al., 1990 and Mays et al., 1991): Combined prior and likelihood models.

In the noisy channel model, we imagine that the surface form we see is actually a "distorted" form of an original word passed through a noisy channel. The decoder passes each hypothesis through a model of this channel and picks the word that best matches the surface noisy word. (Jurafsky and Martin, 2024)

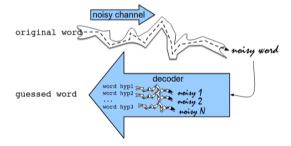


Figure 6: Noisy Channel Model

This noisy channel model is a kind of Bayesian inference.

Out of all possible words in the vocabulary V we want to find the word w such that P(w|x) is highest.

$$\hat{w} = \operatorname*{arg\,max}_{w \in V} P(w|x)$$

Using Bayes: P(x, w) = P(w|x)P(x) = P(x|w)P(w),

$$\hat{w} = \underset{w \in V}{\operatorname{arg\,max}} \frac{P(x|w)P(w)}{P(x)} = \underset{w \in V}{\operatorname{arg\,max}} \underbrace{P(x \mid w)P(w)}_{\text{channel model or likelihood}}$$

$$\hat{w} = \underset{w \in V}{\operatorname{arg\,max}} \left(\log P(x \mid w) + \log P(w) \right)$$

function NOISY CHANNEL SPELLING(word x, dict D, lm, editprob) **returns** correction

```
if x \notin D candidates, edits \leftarrow All strings at edit distance 1 from x that are \in D, and their edit for each c, e in candidates, edits channel \leftarrow editprob(e) prior \leftarrow \operatorname{lm}(c) score[c] = \log \operatorname{channel} + \log \operatorname{prior} return \operatorname{argmax}_c \operatorname{score}[c]
```

Figure 7: Noisy channel model for spelling correction for unknown words (Jurafsky and Martin, 2024).

Example

original word

actress cress caress access across acres



Figure 8: Example: misspelling acress.

		Transformation										
		Correct	Error	Position								
Error	Correction	Letter	Letter	(Letter #)	Type							
acress	actress	t	_	2	deletion							
acress	cress	_	a	0	insertion							
acress	caress	ca	ac	0	transposition							
acress	access	С	r	2	substitution							
acress	across	О	e	3	substitution							
acress	acres	_	S	5	insertion							
acress	acres	_	S	4	insertion							

Figure 9: Candidate corrections for the misspelling acress and the transformations that would have produced the error (after Kernighan et al. (1990)). "—" represents a null letter. (Jurafsky and Martin, 2024)

w	count(w)	p(w)
actress	9,321	.0000231
cress	220	.000000544
caress	686	.00000170
access	37,038	.0000916
across	120,844	.000299
acres	12,874	.0000318

Figure 10: Language model from the 404,253,213 words in the Corpus of Contemporary English (COCA) (Jurafsky and Martin, 2024).

Error model

- A perfect model would need all sorts of factors: who the typist was, whether the typist was left-handed or right-handed, and so on.
- We can get a pretty reasonable estimate of P(x|w) just by looking at **local context**: the identity of the correct letter itself, the misspelling, and the surrounding letters.
- Confusion Matrices:
 - del[x, y]: count(xy typed as x)
 - ins[x, y]: count(x typed as xy)
 - sub[x, y]: count(x typed as y)
 - trans[x, y]: count(xy typed as yx)

	<pre>sub[X, Y] = Substitution of X (incorrect) for Y (correct)</pre>																									
X												Y	(co	rrect)											
	a	b	С	d	c	f	g	h	i	j	k	1	m	n	0	p	q	r	S	t	u	V	w	х	У	Z
a	0	0	7	1	342	0	0	2	118	0	1	0	0	3	76	0	0	1	35	9	9	0	1	0	5	0
b	0	0	9	9	2	2	3	1	0	0	0	5	11	5	0	10	0	0	2	1	0	0	8	0	0	0
c	6	5	0	16	0	9	5	0	0	0	1	0	7	9	1	10	2	5	39	40	1	3	7	1	1	0
d	1	10	13	0	12	0	5	5	0	0	2	3	7	3	0	1	0	43	30	22	0	0	4	0	2	0
С	388	0	3	11	0	2	2	0	89	0	0	3	0	5	93	0	0	14	12	6	15	0	1	0	18	0
f	0	15	0	3	1	0	5	2	0	0	0	3	4	1	0	0	0	6	4	12	0	0	2	0	0	0
g	4	1	11	11	9	2	0	0	0	1	1	3	0	0	2	1	3	5	13	21	0	0	1	0	3	0
h	1	8	0	3	0	0	0	0	0	0	2	0	12	14	2	3	0	3	1	11	0	0	2	0	0	0
í	103	0	0	0	146	0	1	0	0	0	0	6	0	0	49	0	0	0	2	1	47	0	2	1	15	0
j	0	1	1	9	0	0	1	0	0	0	0	2	1	0	0	0	0	0	5	0	0	0	0	0	0	0
k	1	2	8	4	1	1	2	5	0	0	0	0	5	0	2	0	0	0	6	0	0	0	. 4	0	0	3
1	2	10	1	4	0	4	5	6	13	0	1	0	0	14	2	5	0	11	10	2	0	0	0	0	0	0
m	1	3	7	8	0	2	0	6	0	0	4	4	0	180	0	6	0	0	9	15	13	3	2	2	3	0
n	2	7	6	5	3	0	1	19	1	0	4	35	78	0	0	7	0	28	5	7	0	0	1	2	0	2
0	91	1	1	3	116	0	0	0	25	0	2	0	0	0	0	14	0	2	4	14	39	0	0	0	18	0
p	0	11	1	2	0	6	5	0	2	9	0	2	7	6	15	0	0	1	3	6	0	4	1	0	0	0
q	0	0	1	0	0	0	27	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
r	0	14	0	30	12	2	2	8	2	0	5	8	4	20	1	14	0	0	12	22	4	0	0	1	0	0
S	11	8	27	33	35	4	0	1	0	1	0	27	0	6	1	7	0	14	0	15	0	0	5	3	20	1
t	3	4	9	42	7	5	19	5	0	1	0	14	9	5	5	6	0	11	37	0	0	2	19	0	7	6
u	20	0	0	0	44	0	0	0	64	0	0	0	0	2	43	0	0	4	0	0	0	0	2	0	8	0
v	0	0	7	0	0	3	0	0	0	0	0	1	0	0	1	0	0	0	8	3	0	0	0	0	0	0
w	2	2	1	0	1	0	0	2	0	0	1	0	0	0	0	7	0	6	3	3	1	0	0	0	0	0
x	0	0	0	2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	9	0	0	0	0	0	0	0
У	0	0	2	0	15	0	1	7	15	0	0	0	2	0	6	1	0	7	36	8	5	0	0	1	0	0
-	()	Λ	0	7	Λ	Λ	0	Λ	Λ	Λ	Ω	"1	- 5	0	0	n	0	2	21	2	Ω	Ω	0	Ω	2	0

Figure 11: Confusion matrix for spelling errors (Kernighan et al., 1990).

Estimating the channel model

$$P(x|w) = \begin{cases} \frac{\text{del}[x_{i-1}, w_i]}{\text{count}[x_{i-1}w_i]}, & \text{if deletion} \\ \frac{\text{ins}[x_{i-1}, w_i]}{\text{count}[w_{i-1}]}, & \text{if insertion} \\ \frac{\text{sub}[x_i, w_i]}{\text{count}[w_i]}, & \text{if substitution} \\ \frac{\text{trans}[w_i, w_{i+1}]}{\text{count}[w_i, w_{i+1}]}, & \text{if transposition} \end{cases}$$

Candidate	Correct	Error		
Correction	Letter	Letter	$\mathbf{x} \mathbf{w}$	P(x w)
actress	t	-	c ct	.000117
cress	-	a	a #	.00000144
caress	ca	ac	ac ca	.00000164
access	С	r	r c	.000000209
across	О	e	e o	.0000093
acres	-	s	es e	.0000321
acres	-	S	ss s	.0000342

Figure 12: Channel model for acress; the probabilities are taken from the del[], ins[], sub[], and trans[] confusion matrices as shown in Kernighan et al. (1990).

Final probabilities for each of the potential corrections

Candidate	Correct	Error				
Correction	Letter	Letter	$\mathbf{x} \mathbf{w}$	P(x w)	P(w)	$10^9 *P(x w)P(w)$
actress	t	-	c ct	.000117	.0000231	2.7
cress	-	a	a #	.00000144	.000000544	0.00078
caress	ca	ac	ac ca	.00000164	.00000170	0.0028
access	С	r	r c	.000000209	.0000916	0.019
across	0	e	e o	.0000093	.000299	2.8
acres	-	s	es e	.0000321	.0000318	1.0
acres	-	s	ss s	.0000342	.0000318	1.0
ucres			00 0	10000012	.0000070	1.0

Figure 13: Computation of the ranking for each candidate correction, using the language model shown earlier and the error model. The final score is multiplied by 10^9 for readability (Jurafsky and Martin, 2024).



Unfortunately, the algorithm was wrong here; the writer's intention becomes clear from the context: ... was called a "stellar and versatile **acress** whose combination of sass and glamour has defined her ...". The surrounding words make it clear that actress and not across was the intended word. (Jurafsky and Martin, 2024)

Using the *Corpus of Contemporary American English* to compute bigram probabilities for the words *actress* and *across* in their context using add-one smoothing, we get the following probabilities:

$$P(\text{actress}|\text{versatile}) = .000021$$

 $P(\text{across}|\text{versatile}) = .000021$
 $P(\text{whose}|\text{actress}) = .0010$
 $P(\text{whose}|\text{across}) = .000006$

Multiplying these out gives us the language model estimate for the two candidates in context:

$$P(\text{versatile actress whose}) = .000021 \times .0010 = 210 \times 10^{-10}$$

 $P(\text{versatile across whose}) = .000021 \times .000006 = 1 \times 10^{-10}$

Spell Checker

Jurafsky, D., & Martin, J. H. (2024). Speech and Language Processing.

Kernighan, M. D. et al. (1990). A spelling correction program based on a noisy channel model.

Mays, E. et al. (1991). Context based spelling correction.

- Noisy Channel Model
 - Correct (Unix, 1990): Takes inputs from SPELL rejected words and provides candidates. Operations: Insertion, Deletion, Substitution, Reversal. Uses error probabilities.

- **Aspell** (2000): Combines spelling and phonetic correction.
 - Hashing for Spell Checking: Efficient candidate lookup using hash tables.
 - Metaphone Algorithm: Handles phonetic corrections by matching words that sound similar.
 - Ispell's Near Miss Strategy:
 - Focuses on edit distance 1 to reduce the search space.
 - Early Dictionary Filtering: Prunes invalid candidates during generation.

Example - Handling Homophones in Aspell

- Misspelled word ther
- Candidates: there, their, they're
- 1 Metaphone
 - The Metaphone algorithm transforms words into phonetic codes based on pronunciation.
 - Phonetic codes for the candidate words:
 - \blacksquare there \rightarrow OR
 - lacksquare their ightarrow OR
 - lacktriangle they're ightarrow OR
 - Homophones share the same code (OR).

2 Workflow:

- Input: Misspelled word ther.
- Step 1: Generate candidates using **edit distance 2 or less**:
 - Candidates: there, their, thee, thor, her, the, they're.
- Step 2: Compute Metaphone codes for all candidates:
 - Candidates phonetically similar to ther (OR) rank higher: there, their, thor, they're.
- Step 3: Rank and suggest based on:
 - Word frequency, Edit distance, Phonetic Similarity, Error Likelihood.

3 Limitations:

- Metaphone matches words by sound but lacks contextual understanding.
- Example:
 - Input: "I went to ther house."
 - Suggestions: thee, their, there, therm, the, her, Thar, Thea, Thor, Thur.
 - Aspell cannot infer the correct word (their) without considering the sentence's context.

GNU Aspell

■ Hunspell (2002): Morphological analyzer with affix rules and phonetic matching.

Key Features:

- Morphological Analysis:
 - Supports complex languages with rich morphology (e.g., Hungarian, Turkish, Finnish).
 - Handles word roots, prefixes, and suffixes using affix rules.
- Dictionary System:
 - Two components:
 - 1 Dictionary File: Contains root forms of words.
 - 2 Affix File: Defines rules for combining roots with prefixes/suffixes.
- Levenshtein Distance:
 - Uses *edit distance* to generate and rank candidate corrections.
- Phonetic Matching:
 - Uses a table-driven phonetic transcription algorithm borrowed from Aspell. It is useful for languages with not pronunciation based orthography.
- n-gram similarity:
 - Improve suggestions.

- Multilingual Support:
 - Available for 98 languages with extensive dictionaries.

Applications:

- Integrated into tools like LibreOffice, Firefox, and Chrome for multilingual spell checking.
- Supports custom dictionaries for specialized fields (e.g., medical, legal).

Hunspell at GitHub

■ Norvig's Algorithm (2007): Uses Damerau-Levenshtein distance to generate candidates.

Key Features:

- Eit Distance:
 - Generates all possible words within a given edit distance (e.g., 1 or 2) from the misspelled word.
 - Handles insertion, deletion, substitution, and transposition.
- Dictionary Lookup:
 - Filters candidates by validating them against a word dictionary.
- Ranking:
 - Ranks valid candidates based on:
 - Word Frequency: More frequent words are prioritized.
 - Likelihood of Errors: Based on the Noisy Channel Model (optional).

How to Write a Spelling Corrector

Spell Checker

■ **Neural-Based Models**: Leverage deep learning for advanced error detection and correction.

Key Techniques in Spelling Correction

The Noisy Channel Model

■ Bayesian Inference:

Finds the word (w) that maximizes (P(w|x)):

$$\hat{w} = \arg\max_{w \in V} P(x|w)P(w)$$

- Components:
 - **Prior Probability**: (P(w)) (from n-grams like unigrams, bigrams, etc.)
 - Error Model: (P(x|w)) (using confusion matrices).

Confusion Matrix Examples:

- **Deletion**: (del[x, y])
- Insertion: (ins[x, y])
- Substitution: (sub[x, y])
- **Transposition**: (trans[x, y])

Lexical Similarity Metrics

1 Levenshtein Distance (Edit Distance):

Minimal number of insertions, deletions, and replacements to transform one word into another.

2 Jaro Similarity:

Measures similarity based on matching characters and transpositions.

3 Keyboard Distance:

Considers physical proximity of keys.

4 Phonetic Matching:

Algorithms like Soundex and Metaphone to identify similar-sounding words.

Spell Checker

Probabilistic Models

■ Example:

Correction for real-word spelling errors.

Input: Only two of thew apples

Candidates: two of the, two of threw, two of thaw, etc.

Use trigram probabilities (e.g., (P(the|two of) = 0.476012)).

Neural Networks and Deep Learning

■ Models like RNNs, Transformers, and LSTMs learn patterns from large datasets.

Real-World Spelling Correction

- Input: Sentence with an error.
- Output: Highest-probability correction based on:
 - 1 Language Model ((P(W)))
 - 2 Channel Model ((P(x|W))).

Domain-Specific Spell Checkers

- 1 Medical
 - MedSpell: a medical spelling and autocorrect application
 - OpenMedSpel (open-source)
- 2 Programming
 - CodeSpell: designed primarily for checking misspelled words in source code
- 3 Learning
 - Kidspell: A child-oriented, rule-based, phonetic spellchecker
- 4 Accessibility
 - Real Check: A Spellchecker for Dyslexia
- 5 Custom Dictionaries
 - Hunspell and Aspell: Add specialized vocabularies

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