

KD-Trees

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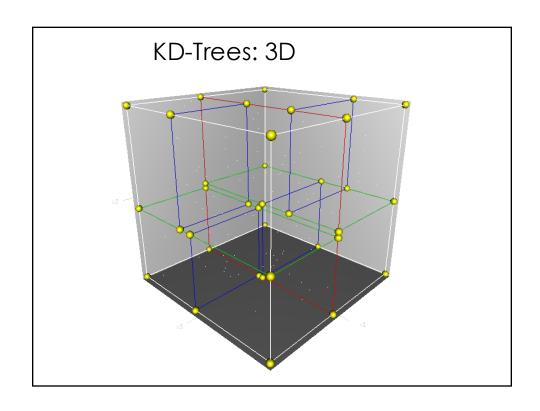
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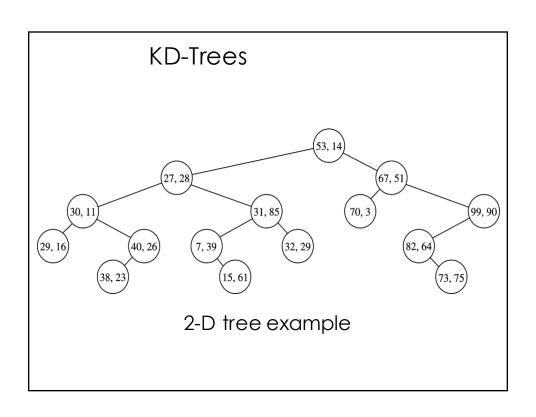
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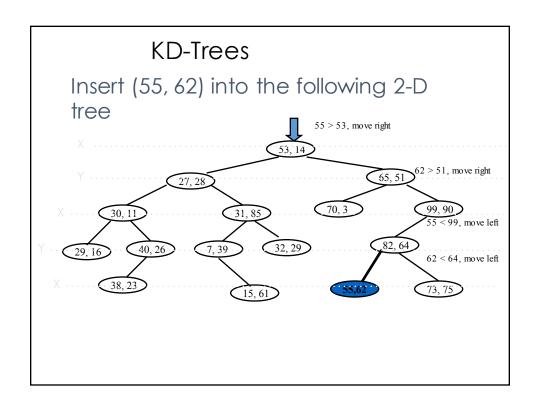
KD-Trees

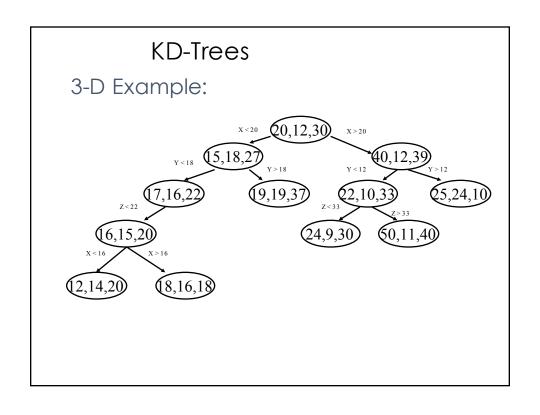
- No es necesario recorrer todos los datos para encontrar los vecinos cercanos
- Usando una estructura para almacenar los datos podemos realizar una búsqueda más eficiente.
- Un KD-Tree es una estructura de datos que particiona el espacio de tal manera de organizar los puntos k-dimensionales

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KD-Trees: Construction

Kd-Trees(object_list , depth = 0)
 if(object_list == empty) return empty set

K = size(object_list,2) % Número de variables donde viven los datos
 axis = mod(depth , k) % Dimensión en cuestión según la profundidad

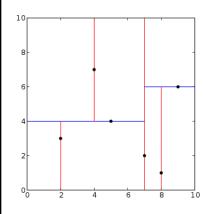
object_list.sort(axis)
largo = len(object_list)
median = object_list(largo// 2)

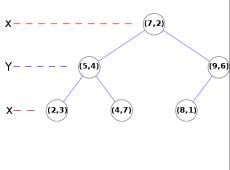
Node = []
Node.location = median
Node.left_child = Kd-Trees(object_list(0:median), depth+1)
Node.right_child = Kd-Trees(object_list(median+1:largo), depth+1)

return Node

KD-Trees: Construction Example

Object List = (2,3) (5,4) (9,6) (4,7) (8,1) (7,2)





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1- NN search Algorithm

- 1. Starting from the root move down the three recursively (like in the construction process)
- 2. Once a leaf node is reached:

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current best = leaf
current distance = dist(leaf,x)
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- 3. Unwind the recursion, move one level up:
 - if the current node is closer than the current best current best = current node current distance = dist(current node, x)
 - if the other child's size could be a closer point to x than the current best

 $\label{eq:KNN-KD-tree} \mbox{(KD_tree(root = current node), x)} \\ \mbox{else}$

Go to step 3

