## A Finite Presentation of CNOT-dihedral Operators

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#### Outline

Quantum circuits

Clifford+T and universal gate sets

CNOT-dihedral gates

Rules for CNOT-dihedral gates

Normal forms

Phase polynomials and uniqueness

Open questions

# Phase polynomials

Action of a diagonal gate:

$$D|x\rangle = \omega^{p_D(x)}|x\rangle, \quad p_D: \mathbb{Z}_2^n \to \mathbb{Z}_8,$$

$$p_D(x) = \sum_{i=1}^k a_i g_i(x).$$

where  $a_i \in \mathbb{Z}_8$  and  $g_i : \mathbb{Z}_2^n \to \mathbb{Z}_2$  are terms on at most n variables.

$$\omega^{k} |x\rangle = \omega^{k} |x\rangle$$

$$T^{k} |x_{1}\rangle = \omega^{kx_{1}} |x_{1}\rangle$$

$$U^{k} |x_{1}x_{2}\rangle = \omega^{k(x_{1} \oplus x_{2})} |x_{1}x_{2}\rangle$$

$$V^{k} |x_{1}x_{2}x_{3}\rangle = \omega^{k(x_{1} \oplus x_{2} \oplus x_{3})} |x_{1}x_{2}x_{3}\rangle$$

Therefore for D in normal form we have

$$p_D(x) = a_0 + \sum_i a_i x_i + \sum_{i < j} b_{ij}(x_i \oplus x_j) + \sum_{i < j < k} c_{ijk}(x_i \oplus x_j \oplus x_k)$$

Where  $a_i \in \mathbb{Z}_8$ ,  $b_{ij} \in \mathbb{Z}_4$ ,  $c_{ijk} \in \mathbb{Z}_2$  (from the bounds on  $deg_X(D)$  for X = T, U, V obtained in normal form).

### Uniqueness

Suppose D and D' are distinct diagonal normal forms, then by construction  $p_D(x) \neq p_{D'}(x)$  as polynomials. However this does not mean that  $\exists y \in \mathbb{Z}_2^n$  s.t.  $p_D(y) \neq p_{D'}(y)$ .

Counterexample:  $4x_1 + 4x_2 + 4(x_1 \oplus x_2) = 0 \mod 8$  for any  $x_1, x_2 \in \mathbb{Z}_2$ 

To show there exists such a y need a technical lemma: construct a multilinear polynomial  $q:\mathbb{Z}_8^n\to\mathbb{Z}_8$  such that  $p_D(y)-p_{D'}(y)=q(y)\ \forall y\in\mathbb{Z}_2^n.$ 

Do this by translating from mod 2 to mod 8:

$$x_i \oplus x_j = x_i + x_j - 2x_i x_j$$

$$x_i \oplus x_j \oplus x_k = x_i + x_j + x_k - 2x_i x_j - 2x_i x_k - 2x_j x_k + 4x_i x_j x_k$$

Then  $p_D(x) - p_{D'}(x) \neq 0 \implies q(x) \neq 0$ , because of the bounds on  $a_i$ ,  $b_{ij}$  and  $c_{ijk}$ . Pick a non-zero term  $dx_{i_1} \dots x_{i_k}$  in q(x) then letting y be the vector with 1's in  $i_j$ th positions and zero everywhere else we obtain  $q(y) = d \neq 0$  and so  $D|y\rangle \neq D'|y\rangle$ .

### Open questions

- ▶ Interaction between ZX and CNOT-dihedral rules (rule 13?)
- ► The paper doesn't contain an algorithm for normalizing CNOT-dihedral circuits, it uses the properties of the ambient symmetric monoidal structure non-constructively. What could be a rewrite system? Algebraic vs combinatorial description.
- Complexity of CNOT-dihedral circuits.