# 4. Number Theory and Cryptography

# 4.1: Divisibility and Modular Arithmetic

#### **Divisibility and Its Properties**

Divisibility is fundamental in number theory. If a and b are integers with a≠0, then a divides b (denoted a|b) if there exists an integer c such that:

b=ac

For example, 5|20 because  $20=5\times4$ , but  $5\nmid22$  since 22 is not a multiple of 5.

#### **Properties of Divisibility:**

- If a|b and a|c, then a|(b+c).
- If a|b, then a|bc for any integer c.
- If a|b and b|c, then a|c.

### The Division Algorithm

For any integer a and positive integer d, there exist unique integers q (quotient) and r (remainder) such that:

a=dq+r, 0≤r<d

#### **Example:**

Dividing 101 by 11:

101=11×9+2

Here, q=9 and r=2.

### **Modular Arithmetic**

In modular arithmetic, two integers are **congruent modulo m** if they have the same remainder when divided by m:

a≡b(mod m) if and only if mI(a-b)

#### **Example:**

29≡5(mod 12)

because 29–5=24, which is divisible by 12.

#### **Properties of Modular Arithmetic:**

- 1.  $(a+c)\equiv (b+d) \pmod{m}$
- 2.  $(a \cdot c) \equiv (b \cdot d) \pmod{m}$

# 4.2: Integer Representations and Algorithms

## **Base b Expansion**

Any integer n can be expressed in base b as:

$$n = a_k b^k + a_{k-1} b^{k-1} + .... + a_1 b + a_0$$

where 0≤a<sub>i</sub><b.

#### **Example:**

Converting 241 to binary:

 $(241)_{10} = (11110001)_2$ 

### **Algorithms for Base Conversion**

To convert from decimal to base b, repeatedly divide by b, recording the remainders.

### **Binary Arithmetic**

Binary addition follows:

- 0+0=0
- 0+1=1
- 1+1=10(carry 1)

Multiplication is performed using shift-and-add methods.

#### **Modular Exponentiation**

Used in cryptography for fast exponentiation:

b<sup>n</sup> mod m

is computed efficiently by **exponentiation by squaring**.

### 4.3: Primes and Greatest Common Divisors

### **Prime Numbers and Factorization**

A prime number is a number greater than 1 with exactly two distinct divisors: 1 and itself.

#### **Prime Factorization Example:**

 $120=2^3\times3\times5$ 

### **Greatest Common Divisor (GCD)**

The GCD of a and b is the largest number dividing both.

Using the **Euclidean Algorithm**:

gcd(a,b)=gcd(b,a mod b)

### **Example:**

gcd(252,198)=18

### **Bézout's Identity**

For any integers a and b, there exist integers x and y such that:

gcd(a,b)=ax+by

# 4.4: Solving Congruences

# **Linear Congruences**

Equations of the form:

ax≡b(mod m) have a solution if gcd(a,m)|b.

## The Chinese Remainder Theorem (CRT)

If we have:

 $x\equiv a_1 \pmod{m_1}, x\equiv a_2 \pmod{m_2}$ 

where  $m_1$  and  $m_2$  are coprime, then there is a unique solution modulo  $m_1m_2$ .

# **4.5: Applications of Congruences**

### **Hashing Functions**

Used to store data efficiently:

 $h(k)=k \mod m$ 

#### **Pseudorandom Number Generation**

Using the linear congruential generator (LCG):

 $x_{n+1}=(ax_n+c) \mod m$ 

### **Check Digit Schemes**

Used in ISBN and credit card validation.

## **Example (ISSN Checksum):**

 $d_8 \equiv 3d_1 + 4d_2 + 5d_3 + \dots + 9d_7 \pmod{11}$ 

# 4.6: Cryptography

## **Classical Ciphers**

The Caesar Cipher shifts letters by kkk:

 $f(p)=(p+k)\mod 26$ 

# **Public-Key Cryptography (RSA)**

- 1. Choose primes p and q, compute n=pq.
- 2. Compute  $\varphi(n)=(p-1)(q-1)$ .
- 3. Choose e with  $gcd(e, \varphi(n))=1$ .
- 4. Compute  $d\equiv e^{-1} \pmod{\varphi(n)}$ .
- 5. Encrypt: c=me mod n.
- 6. Decrypt: m=c<sup>d</sup> mod n.

## **Diffie-Hellman Key Exchange**

Allows two parties to share a secret key securely.

## **Homomorphic Encryption**

Allows computations on encrypted data, useful for secure cloud computing.