

# Fully Automated Formal Verification: How far can we go?

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```
mapping (address => uint) balance;
address immutable owner;
constructor(uint amt) {
   return balance[ user]:
   uint prevBal = balance[msg.sender] + balance[_to];
    balance[ to] += amt:
   uint postBal = balance[msg.sender] + balance[ to]:
    assert(prevBal == postBal);
```

```
contract AMM is ERC20 {
   ERC20 token0:
   ERC20 token1:
   constructor(address token0, address token1) {
       token0 =
       token1 =
                ERC20( token1):
   function swap(address src. address dst. uint amt) external {
       require(src != dst, "no self swap");
       uint out
       uint kpost = token0.balance0f(address(this)) * token1.balance0f(address(this));
```

```
IERC777Sender(from).tokensToSend(operator, from, to, amount);
function callTokensReceived(address operator, address from, address to, uint256 amount) private {
```

- ♦ Mythril symbolic execution
- ♦ hevm symbolic execution, invariant fuzzing
- ♦ Echidna invariant fuzzing
- ♦ SMTChecker model checking
- ♦ solc-verify model checking
- ♦ VeriSmart model checking

Tools that claim to try to prove/break properties automatically and are publicly available.

Which tool can either prove correctness or find bugs in all the examples?

None

Automated formal verification is undecidable

# Target Solidity

- ♦ A lot to encode: high level features, various data types, pointers, inheritance.
- ♦ Results rely on compiler correctness.

#### Pros

- ♦ Gives more structure information, for example, loops, external calls, functions.
- ♦ Can try harder problems, involving loop/contract invariants.

Model checking: SMTChecker, solc-verify, VeriSmart

## Target EVM bytecode Cons

- ♦ Not a lot of structure.
- ♦ Hard to track storage, external calls, functions.
- ♦ Needs to verify ABI encoding/decoding.

#### Pros

- ♦ Easier to encode.
- Results are closer to the deployed object.

Symbolic execution and fuzzing: Mythril, hevm, Echidna

### Experiment:

Use each tool on each example, first automatically then tweaking parameters and writing specs taylored to the tool.

#### Prove functional correctness of transfer function

```
contract Token {
   mapping (address => uint) balance;
   address immutable owner;
    function balanceOf(address _user) public view returns (uint) {
       uint prevBal = balance[msg.sender] + balance[_to];
       uint postBal = balance[msg.sender] + balance[ to]:
```

### Prove functional correctness of transfer function

- ♦ Mythril OK
- ♦ hevm OK
- ♦ Echidna OK (no bugs found)
- ♦ SMTChecker OK
- ♦ solc-verify OK
- ♦ VeriSmart OK

### Find bug in buggy transfer function

```
contract Token {
   mapping (address => uint) balance:
   address immutable owner;
       return balance[ user]:
       uint prevBal = balance[msg.sender] + balance[ to]:
       uint postBal = balance[msg.sender] + balance[_to];
       assert(prevBal == postBal):
```

### Find bug in buggy transfer function

- ♦ Mythril OK, with counterexample
- ♦ hevm OK, with counterexample
- ♦ Echidna OK, with counterexample
- ♦ SMTChecker OK, with counterexample
- ♦ solc-verify OK, no counterexample
- ♦ VeriSmart No

### AMM swap functional correctness (?)

```
ERC20 token0:
ERC20 token1:
    token0 =
    uint out
    uint kpost = token0.balanceOf(address(this)) * token1.balanceOf(address(this)):
```

AMM swap functional correctness (?)

Symbolic execution and model checking tools could not prove/disprove the assertion

# AMM swap functional correctness (?) Fuzzing?

```
Tests found: 1

Tests found: 1
```

## AMM swap functional correctness Fuzzing with hevm

```
new User(token0, token1, pair);
```

# AMM swap functional correctness Fuzzing with Echidna

### AMM swap functional correctness

```
function swap(address src, address dst, uint amt) external {
   require(src != dst, "no self swap");
                  address(token0) || src
                                             address(token1), "src not in pair");
                  address(token0) || dst =
   KPrev
            token0.balanceOf(address(this)) * token1.balanceOf(address(this));
   ERC20(src).transferFrom(msg.sender, address(this), amt);
   uint out
       ERC20(dst).balanceOf(address(this))
        (KPrev / ERC20(src).balanceOf(address(this)) + 1);
   ERC20(dst).transfer(msg.sender, out);
   KPost
           token0.balanceOf(address(this)) * token1.balanceOf(address(this));
```

# Model checkers still cannot prove correctness, but fuzzers could not find any other problems.

```
Fuzzing invariant
Running 1 tests for src/Amm.t.sol:TestAMM
[PASS] invariant_k() (runs: 100, depth: 20)
```

### Prove function correctness of deposit

Prove function correctness of deposit:

No tool could prove that the assertion is not reachable automatically.

#### Modified version for hevm

### hevm results

```
Running Deposit contract
checking postcondition...
Q.E.D.
Explored: 295 branches without assertion violations
```

#### Modified version for SMTChecker

```
// Add deposit data root to Merkle tree (update a single 'branch' node)
deposit_count ** 1;
uint size = deposit_count;
for (uint height = 0; height < DEPOSIT_CONTRACT_TREE_DEPTH; height **) {
    if (fisize & 1) == 1) {
        if (fisize & 2) == 1) {
            branch[height] = node;
            return;
    }
    node = sha256(abi.encodePacked(branch[height], node));
    size /= 2;
}
// As the loop should always end prematurely with the 'return' statement,
// this code should be unreachable. We assert 'false' just to be safe.
assert(false);</pre>
```

+ use non default Horn solver Eldarica via solc-js' SMT callback + use Eldarica's abstract:off option

# SMTChecker's inductive invariant for the loop before the assertion

### Prove ERC777 property in the presence of external calls

```
callTokensToSend(from, from, recipient, amount):
IERC777Recipient(to).tokensReceived(operator, from, to, amount):
```

### SMTChecker says that the property does not hold

```
arning: CHC: Assertion violation happens here.
Counterexample:
totalSupply = 0
amount = 1
prevTotal = 115792089237316195423570985008687907853269984665640564039457584007913129639897
Transaction trace:
ERC777.constructor(115792089237316195423570985008687907853269984665640564039457584007913129639897)
State: totalSupply = 115792089237316195423570985008687907853269984665640564039457584007913129639897
FRC777, burn(1)
   Context, msgSender() -- internal call
   ERC777, burn(7720, 1) -- internal call
       Context, msgSender() -- internal call
       ERC777, callTokensToSend(7720, 7720, 0, 1) -- internal call
            IERC777Sender(from).tokensToSend(operator, from, to, amount) -- untrusted external call, synthesized as:
               ERC777.burn(115792089237316195423570985008687907853269984665640564039457584007913129639897){ value: 14 } -- reentrant call
                   Context, msgSender(){ value: 14 } -- internal call
                   ERC777, burn(3565, 11579208923731619542357098500868790785326998466564039457584007913129639897){ value: 14 } -- internal call
                       Context. msgSender(){ value: 14 } -- internal call
                       FRC777, callTokensToSend(3565, 3565, 0. 11579208923731619542357098500868790785326998466564039457584007913129639897){ value: 14 } -- internal call
                           IERC777Sender(from).tokensToSend(operator, from, to, amount) -- untrusted external call
```

- Remove strings (hard for counterexamples)

#### Let's test that

```
IERC777 immutable erc:
function tokensReceived(address, address, uint256 amount) public override
function tokensToSend(address, address, address, uint256 amount) public override {}
Rec rec;
function tokensToSend(address, address, address, uint256 amount) public override {}
```

### Property is indeed broken

```
Running 1 tests for src/FRC777.t.sol:FRC777Test
    - create FRC77700xCe71065D4017F316FC606Fe4422e11eB2c47c246 (src/ERC777.t.sol:29)
     -call ERC777::transfer(address.uint256)(@0x185a4dc360CE69bDCceE33b3784B0282f7961aea, 100) (src/ERC777.t.sol:34)
      -call ERC777Test::tokensToSend(address.address.address.uint256)(ERC777Test@0xb4c79daB8f259C7Aee6E5b2Aa729821864227e84, ERC777Test@0xb4c79daB8f259C7A
ee6E5b2Aa729821864227e84. @0x185a4dc360CE69bDCceE33b3784B0282f7961aea. 100) (src/ERC777.sol:67)
        Sent(100) (src/ERC777.sol:323)
       -Transfer(100) (src/ERC777.sol:324)
       -call 0x185a4dc360CE69bDCceE33b3784B0282f7961aea::tokensReceived(ERC777Test@0xb4c79daB8f259C7Aee6E5b2Aa729821864227e84. ERC777Test@0xb4c79daB8f259C7
Aee6E5b2Aa729821864227e84. @0x185a4dc360CE69bDCceE33b3784B0282f7961aea. 100) (src/ERC777.sol:84)
         -call ERC777::burn(uint256)(100) (src/ERC777.t.sol:16)
           Lcall 0x185a4dc360CE69bDCceE33b3784B0282f7961aea::tokensToSend(@0x185a4dc360CE69bDCceE33b3784B0282f7961aea, @0x185a4dc360CE69bDCceE33b3784B028
-Transfer(100) (src/FRC777.sol:307)
```

### We can forbid reentrancy

# Are we safe now? Reentrancy guard blocks the tx before the assertion failure

```
Running 1 tests for src/ERC777.t.sol:ERC777Test
      create ERC777Mutex00xCe71065D4017F316EC606Fe4422e11eB2c47c246 (src/ERC777.t.sol:40)
       call ERC0777Mutex::transfer(address,uint256)(@0x185a4dc360CE69bDCceE33b378480282f7961aea, 100) (src/ERC777.t.sol:45)
        -call ERC777Test::tokensToSend(address,address,address,uint256)(ERC777Test@0xb4c79daB8f259C7Aee6E5b2Aa729821864227e84, ERC777Test@0xb4c79daB8f259C7Aee6E5b2Aa72982
1864227e84. @0x185a4dc360CE69bDCceE33b3784B0282f7961aea. 100) (src/ERC777Mutex.sol:74)
           Sent(100) (src/ERC777Mutex.sol:327)
          Transfer(100) (src/FRC777Mutex.sol:328)
          -call 0x185a4dc368CE69bDCceE33b3784B0282f796laea::tokensReceived(ERC777Test00xb4c79daB8f259C7Aee6E5b2Aa729821864227e84. ERC777Test00xb4c79daB8f259C7Aee6E5b2Aa7298
 21864227e84 @0x185a4dc360CF69bDCceF33b3784R0282f7961aea 100) (src/FRC777Mutex.sol:88)
            - call ERC777Mutex::burn(uint256)(100) (src/ERC777.t.sol:21)
                Revert 0x (src/ERC777Mutex.sol:88)
             Revert Av (src/FRC777.t.sel:45)
```

### SMTChecker proves that now the assertions are safe

```
Contract invariants and external call properties for ERC777Mutex.sol:ERC777Mutex:
(((cerrorCode> >= ) & (llock || ((_totalSupply + ((- 1) * _totalSupply*)) <= 0)) && (lock || ((_totalSupply + ((- 1) * _totalSupply + ((- 1) * _totalSupply*)) <= 0)) & ((llock || ((_totalSupply + ((- 1) * _totalSupply + ((- 1) * _totalS
```

$$\mathsf{lock} \implies (\mathsf{\_totalSupply} = \mathsf{\_totalSupply'})$$

- hevm internals Symbolic execution
- ♦ Execute EVM while building an AST for symbolic terms
- ♦ Upon JUMPI, explore both paths and add new constraints
- ♦ Avoid infinite loops by `--max-iterations`
- ♦ For all end states violating the property, check if constraints are satisfiable

## hevm internals - Symbolic ds-tests & rpc

#### data StorageModel

- = ConcreteS -- ^ Uses `Concrete` Storage. Reading / Writing from abstract
  - -- locations causes a runtime failure.
  - -- Can be nicely combined with RPC.
  - SymbolicS -- ^ Uses `Symbolic` Storage. Reading / Writing never
    - -- reaches RPC, but always done using an SMT array
    - -- with no default value.
  - InitialS -- ^ Uses `Symbolic` Storage. Reading / Writing never
    - -- reaches RPC, but always done using an SMT array
    - -- with 0 as the default value.

# hevm internals - Equivalence checking

- ♦ Make two copies of an abstract VM, differing only in the code of the current contract
- ♦ Execute both symbolically and take all pairs of endstates
- ♦ For every pair, check if both constraint sets are satisfiable and if results differ
- ♦ If no such pair exists, the contracts are equivalent

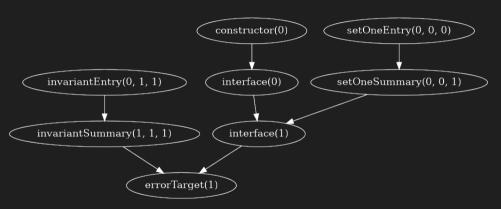
#### SMTChecker internals

```
contract BinaryMachine {
    function setOne() public {
    function setZero() public {
    function invariant() public view {
        assert(x < 1);
```

```
error = 0 \land x = 0 \implies constructor(error. x)
                                     error = 0 \land x = x' \implies setOneEntry(error, x, x')
                              setOneEntry(error, x, x') \implies setOneSummary(error, x, 1)
                                     error = 0 \land x = x' \implies setZeroEntry(error, x, x')
                             setZeroEntry(error, x, x') \implies setZeroSummary(error, x, 0)
                                     error = 0 \land x = x' \implies invariantEntry(error, x, x')
invariantEntry(error, x, x') \land x > 2 \land errorCode = 1 \implies invariantSummary(errorCode, x, x')
                   invariantEntry(error, x, x') \land x < 2 \implies invariantSummary(error, x, 1)
```

```
error = 0 \land constructor(error, x) \implies interface(x)
 interface(x) \land setOneSummary(error, x, x') \land error = 0 \implies interface(x')
  interface(x) \land setOneSummary(error, x, x') \land error > 0 \implies errorTarget(error)
 interface(x) \land setZeroSummary(error, x, x') \land error = 0 \implies interface(x')
 interface(x) \land setZeroSummary(error, x, x') \land error > 0 \implies errorTarget(error)
interface(x) \land invariantSummary(error, x, x') \land error = 0 \implies interface(x')
interface(x) \land invariantSummary(error, x, x') \land error > 0 \implies errorTarget(error)
```

# Counterexample graph



#### Correct code

```
contract BinaryMachine {
         assert(x <= 2);</pre>
```

#### Inductive invariants

$$interface(e,x)=x<2$$
 
$$invariantSummary(e,x,x')=x'=0 \ \forall \ x'=1$$
 
$$constructorSummary(e,x)=e=0 \ \land \ x'=0$$
 
$$setOneSummary(e,x,x')=e=0 \ \land \ x'=1$$
 
$$setZeroSummary(e,x,x')=e=0 \ \land \ x'=0$$

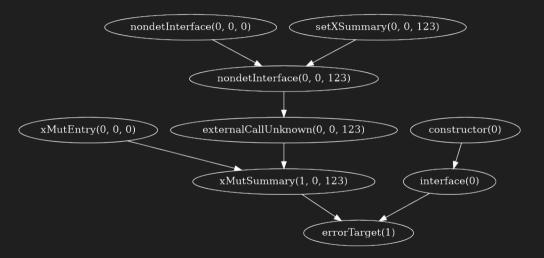
#### SMTChecker internals - External calls

```
interface Unknown {
   function callMe() external;
contract ExtCall {
   uint x;
   function setX(uint _x) public { x = _x; }
    function xMut(Unknown _u) public {
       uint xPrev = x;
       _u.callMe();
        assert(xPrev == x);
```

New rules help with synthesis of externally called unknown functions

$$error = 0 \implies nondetInterface(error, x, x')$$
 
$$error = 0 \land nondetInterface(error, x, x') \land setX(error', x', x'') \implies nondetInterface(error', x, x'')$$
 
$$error = 0 \land nondetInterface(error, x, x') \land xMut(error', x', x'') \implies nondetInterface(error', x, x'')$$

### Counterexample graph synthesizing external calls



#### SMTChecker internals - External calls

```
interface Unknown {
    function callMe() external;
contract ExtCall {
    uint x;
    function xMut(Unknown _u) public {
        uint xPrev = x:
       _u.callMe();
        assert(xPrev == x);
```

Inductive invariants and external call properties

externalCallUnknown(e, x, x') = 
$$(x = 0 \implies e = 0) \land (x = 0 \implies x' = 0)$$
  
interface(x) =  $x = 0$ 

#### Conclusions

- ♦ Automated FV tools can be quite powerful but...
- lacktriangle No automated tool will do the job everytime
- ♦ FV is still an expert domain
- ♦ Specific tool knowledge is required to extract full potential
- ♦ Playing with the tool is essential to get good results

# Thank you!