Case study: a TPC Benchmark for decision support

Leonardo Musini

Data Management for Big Data A.Y. 2022-2023

Contents

1	Introduction	2
2	Database definition 2.1 Database statistics	2 4
3	Query	4
4	Indexing on Tables	6
5	Materialization of Views	7
6	Indexing on Materialized Views	8
7	Horizontal Fragmentation 7.1 Hybrid and Vertical Fragmentations	9 9
8	Results	11
9	Conclusions	12

1 Introduction

The aim of this project is to implement and improve a business oriented ad-hoc query over the public TPC-H benchmark, in order to find the best optimization strategy.

Data have been generated using dbgen package of TPC-H benchmark, that can be downloaded from TPC official website https://www.tpc.org/tpch/, with a scale factor of 10. The optimization strategies that have been considered are the following:

- Indexing on tables;
- Materialization of views;
- Indexing on materialized views;
- Horizontal Fragmentation.

Each approach must keep the database size below 1.5 times the size of the original database. The DBMS chosen for the implementation is PostgreSQL, using pgAdmin 4 management tool. The specifications of the device used in this project are:

- Processor: Intel(R) Core(TM) i5-8250U CPU @ 1.60GHz 1.80 GHz;
- RAM: 8.00 GB;
- Storage: SSD;
- OS: Windows 11.

2 Database definition

The database is made up of 8 tables, whose structure definitions and constraints can be found in the following documentation https://www.tpc.org/tpc_documents_current_versions/pdf/tpc-h_v3.0.1.pdf. The SQL implementation is the following:

```
1 CREATE TABLE part (
      p_partkey INTEGER CHECK (p_partkey >= 0),
      p_name VARCHAR(55),
      p_mfgr CHAR(25),
      p_brand CHAR(10)
      p_type VARCHAR(25),
6
      p_size INTEGER CHECK (p_size >= 0),
      p_container CHAR(10),
      p_retailprice DECIMAL CHECK (p_retailprice >= 0),
9
      p_comment VARCHAR (23)
10
      PRIMARY KEY (p_partkey)
11
12 );
13
14 CREATE TABLE region (
      r_regionkey INTEGER CHECK (r_regionkey >= 0),
      r_name CHAR(25),
16
      r_comment VARCHAR (152),
17
      PRIMARY KEY(r_regionkey)
19 );
20
21 CREATE TABLE nation (
      n_nationkey INTEGER CHECK (n_nationkey >= 0),
22
      n_name CHAR(25),
23
```

```
n_regionkey INTEGER,
      n_comment VARCHAR (152),
25
      PRIMARY KEY(n_nationkey),
26
      FOREIGN KEY(n_regionkey) REFERENCES region(r_regionkey)
28);
29
30 CREATE TABLE supplier (
      s_suppkey INTEGER CHECK (s_suppkey >= 0),
31
32
      s_name CHAR(25),
      s_address VARCHAR(40),
33
      s_nationkey INTEGER,
34
      s_phone CHAR(15),
      s_acctbal DECIMAL,
36
      s_comment VARCHAR (101),
37
      PRIMARY KEY(s_suppkey),
38
      FOREIGN KEY (s_nationkey) REFERENCES nation(n_nationkey)
39
40 );
41
_{\rm 42} CREATE TABLE partsupp (
      ps_partkey INTEGER CHECK (ps_partkey >= 0),
      ps_suppkey INTEGER,
44
      ps_availqty INTEGER CHECK (ps_availqty >= 0),
45
      ps_supplycost DECIMAL CHECK (ps_supplycost >= 0),
      ps_comment VARCHAR(199),
47
      PRIMARY KEY (ps_partkey, ps_suppkey),
48
      FOREIGN KEY (ps_partkey) REFERENCES part(p_partkey),
49
      FOREIGN KEY (ps_suppkey) REFERENCES supplier(s_suppkey)
50
51 );
52
53 CREATE TABLE customer (
      c_custkey INTEGER CHECK (c_custkey >= 0),
      c_name VARCHAR(25),
55
56
      c_address VARCHAR(40),
57
      c_nationkey INTEGER,
      c_phone CHAR(15),
58
      c_acctbal DECIMAL,
      c_mktsegment CHAR (10),
60
      c comment VARCHAR (117).
61
      PRIMARY KEY(c_custkey),
62
      FOREIGN KEY (c_nationkey) REFERENCES nation(n_nationkey)
63
64 );
66 CREATE TABLE orders (
67
      o_orderkey INTEGER,
      o_custkey INTEGER,
68
      o_orderstatus CHAR(1),
69
      o_totalprice DECIMAL CHECK (o_totalprice >= 0),
70
      o_orderdate DATE,
71
      o_orderpriority CHAR(15),
72
73
      o_clerk CHAR(15),
      o\_shippriority\ INTEGER,
74
75
      o_comment VARCHAR(79),
      PRIMARY KEY(o_orderkey),
76
      FOREIGN KEY (o_custkey) REFERENCES customer(c_custkey)
77
78 );
79
80 CREATE TABLE lineitem (
      l_orderkey INTEGER,
81
      l_partkey \ INTEGER,
82
83
      l_suppkey INTEGER,
      l_linenumber INTEGER,
84
      1_quantity DECIMAL CHECK (1_quantity >= 0),
```

```
1_extendedprice DECIMAL CHECK (1_extendedprice >= 0),
       1_discount DECIMAL CHECK (1_discount BETWEEN 0.00 AND 1.00),
87
       1_tax DECIMAL CHECK (1_tax >= 0),
88
       l_returnflag CHAR(1),
       l_linestatus CHAR(1)
90
91
       l_shipdate DATE,
       l_commitdate DATE,
92
       1_receiptdate DATE CHECK (1_shipdate <= 1_receiptdate),</pre>
93
       l_shipinstruct CHAR(25),
94
       l_shipmode CHAR(10),
95
       1_comment VARCHAR (44),
96
97
       PRIMARY KEY (1_orderkey, 1_linenumber),
       FOREIGN KEY (1_orderkey) REFERENCES orders(o_orderkey),
98
       FOREIGN KEY (1_partkey) REFERENCES part(p_partkey),
99
       FOREIGN KEY (1_suppkey) REFERENCES supplier(s_suppkey),
100
       FOREIGN KEY (1_partkey, 1_suppkey) REFERENCES partsupp(ps_partkey, ps_suppkey)
101
102 );
```

2.1 Database statistics

In the following section there is a list of some statistics about the tables and some of their main attributes used in the query, except for primary keys or attributes not strictly necessary:

- part, with 9 columns, 2000000 rows and a size of 363 MB:
 - o p_type, with 150 distinct values;
- region, with 3 columns, 5 rows and a size of 24 kB;
- nation, with 4 columns, 25 rows and a size of 24 kB;
- supplier, with 7 columns, 100000 rows and a size of 22 MB;
- partsupp, with 5 columns, 8000000 rows and a size of 2908 MB;
- customer, with 8 columns, 1500000 rows and a size of 313 MB;
- orders, with 9 columns, 15000000 rows and a size of 4404 MB:
 - o o_orderdate, with 2406 distinct values, that range from 1992-01-01 to 1998-08-02;
- lineitem, with 16 columns, 59986052 rows and a size of 10 GB:
 - o l_extendedprice, with 1351462 distinct values, that range from 900.91 to 104949.50,
 - o l_discount, with 11 distinct values, that range from 0.00 to 0.10.

The total size of the database is 18 GB, then the maximum extension of the database space on the disk must not exceed 27 GB.

3 Query

The query schema asks to implement the aggregation of the export/import of revenue of lineitems between two different nations (\mathbf{E},\mathbf{I}) where \mathbf{E} is the nation of the lineitem supplier and \mathbf{I} the nation of the lineitem customer (export means that the supplier is in the nation \mathbf{E} and import means the customer is in the nation \mathbf{I}). The **revenue** is obtained by:

```
revenue = l_extendedprice * (1 - l_discount)
```

The aggregations should be performed with the following roll-up:

- Month \rightarrow Quarter \rightarrow Year;
- Type;
- Nation \rightarrow Region.

The slicing is over **Type** and **Exporting nation**.

The SQL implementation is the following:

```
1 WITH supplier_nation_region AS (
    SELECT
      s_suppkey,
3
      s_name,
      s_nationkey,
      n_name AS s_nationname,
      r_name AS s_regionname
    FROM supplier
      JOIN nation ON (s_nationkey = n_nationkey)
      JOIN region ON (n_regionkey = r_regionkey)
_{\rm 11} ), customer_nation_region AS (
12
    SELECT
13
      c_custkey,
      c_name,
14
      c_nationkey,
     n_name AS c_nationname,
16
      r_name AS c_regionname
17
    FROM customer
     JOIN nation ON (c_nationkey = n_nationkey)
19
      JOIN region ON (n_regionkey = r_regionkey)
21 ), lineitem_orders AS (
   SELECT
22
      l_partkey,
      l_suppkey,
24
25
      o_orderdate,
      o_custkey,
      l_extendedprice * (1 - l_discount) AS revenue
27
    FROM lineitem JOIN orders ON (l_orderkey = o_orderkey)
28
29 )
30
31 SELECT
   EXTRACT (YEAR FROM o_orderdate) AS o_year,
32
    EXTRACT (QUARTER FROM o_orderdate) AS o_quarter,
33
    EXTRACT (MONTH FROM o_orderdate) AS o_month,
35
    p_type,
36
    s_regionname,
    s_nationname,
    s_name,
38
    c_regionname,
40
    c_nationname,
41
    c_name,
    SUM (revenue) AS total_revenue
43 FROM lineitem_orders
    JOIN part ON l_partkey = p_partkey
    JOIN supplier_nation_region ON (s_suppkey = 1_suppkey)
    JOIN customer_nation_region ON (c_custkey = o_custkey)
46
47 WHERE
    s_nationkey <> c_nationkey
    AND s_nationname = 'UNITED KINGDOM'
    AND p_type = 'LARGE POLISHED STEEL'
50
51 GROUP BY
```

```
o_year,
    o_quarter,
53
54
    o_month,
    p_type,
    c_regionname,
56
57
     c_nationname,
58
    c_name,
59
    s_regionname,
     s_nationname,
60
61
     s_name
62 ;
```

In the previous query, since we need the nation of suppliers and the nation of customers, we create two Common Table Expressions (CTEs), supplier_nation_region and customer_nation_region, combining data respectively between supplier or customer and nation and region tables in the JOIN clause.

A third CTE, lineitem_orders, has been created as a join between lineitem and order tables in which we select the date of the orders, represented by o_orderdate attribute, some foreign keys needed for joining and we calculate the revenue as defined before.

In the main part of the query, we extract year, quarter and month from o_orderdate, select the attributes needed, calculate the total revenue as the sum of single group revenues, join lineitem_orders with part, supplier_nation_region and customer_nation_region, and in conjunction with the GROUP BY clause, we achieve the aggregation with the roll-up previously described.

The slicing is performed in the WHERE clause over attributes p_type of table part and s_nationname of supplier_nation_region.

4 Indexing on Tables

The first optimization was done by using indexes on tables. The EXPLAIN ANALYZE clause can help to understand which are the more costly operations: JOIN, GROUP BY and WHERE clauses are the ones that impact mostly on the query efficiency. Therefore, we have created indexes for attributes in these clauses, except for primary keys since PostgreSQL already creates indexes for them

Here we have the implementation:

```
1 CREATE INDEX IF NOT EXISTS l_partkey_index
2    ON lineitem (l_partkey);
3
4 CREATE INDEX IF NOT EXISTS l_suppkey_index
5    ON lineitem (l_suppkey);
6
7 CREATE INDEX IF NOT EXISTS o_orderdate_index
8    ON orders (o_orderdate ASC NULLS LAST);
9
10 CREATE INDEX IF NOT EXISTS o_custkey_index
11    ON orders (o_custkey);
12
13 CREATE INDEX IF NOT EXISTS p_type_index
14    ON part (p_type ASC NULLS LAST);
15
16 CREATE INDEX IF NOT EXISTS s_name_index
17    ON supplier (s_name ASC NULLS LAST);
```

```
18
19 CREATE INDEX IF NOT EXISTS s_nationkey_index
20 ON supplier (s_nationkey);
21
22 CREATE INDEX IF NOT EXISTS c_name_index
23 ON customer (c_name ASC NULLS LAST);
24
25 CREATE INDEX IF NOT EXISTS c_nationkey_index
26 ON customer (c_nationkey);
27
28 CREATE INDEX IF NOT EXISTS n_regionkey_index
29 ON nation (n_regionkey);
```

The total database size is of 21 GB, respecting the size request.

5 Materialization of Views

Another way to optimize the query is to create materialized views in order to store important information into the disk and then to have a easier and faster access.

For this query we have materialized the three CTEs: in this way we can have a direct access to the tables generated from the joins in the WITH clause.

In the following code we show how to create the materialized views:

```
1 CREATE MATERIALIZED VIEW mv_supplier_nation_region AS
    SELECT
      s_suppkey,
      s_name,
      s_nationkey,
      n_name AS s_nationname,
      r_name AS s_regionname
    FROM supplier
      JOIN nation ON (s_nationkey = n_nationkey)
10
      JOIN region ON (n_regionkey = r_regionkey);
12 CREATE MATERIALIZED VIEW mv_customer_nation_region AS
    SELECT
      c_custkey,
14
1.5
      c_name,
      c_nationkey,
      n_name AS c_nationname,
17
      r_name AS c_regionname
18
    FROM customer
19
      JOIN nation ON (c_nationkey = n_nationkey)
20
21
      JOIN region ON (n_regionkey = r_regionkey);
22
23 CREATE MATERIALIZED VIEW mv_lineitem_orders AS
    SELECT
      l_partkey,
25
26
      l_suppkey,
27
      o_orderdate,
28
      o custkev.
      l_extendedprice * (1 - l_discount) AS revenue
    FROM lineitem JOIN orders ON (l_orderkey = o_orderkey);
```

We can observe some statistics of the materialized views:

- mv_supplier_nation_region, with 5 columns, 100000 rows and a size of 12 MB;
- mv_customer_nation_region, with 5 columns, 1500000 rows and a size of 157 MB;

• mv_lineitem_orders, with 5 columns, 59986052 rows and a size of 3390 MB.

The database has a total size of approximately 21.5 GB, so the request of 1.5 the size of the original database is respected.

Consequently, we modify the query removing the WITH clauses as follows:

```
1 SELECT
    EXTRACT (YEAR FROM o_orderdate) AS o_year,
    EXTRACT (QUARTER FROM o_orderdate) AS o_quarter,
    EXTRACT (MONTH FROM o_orderdate) AS o_month,
    p_type,
    s_regionname,
    s_nationname,
    s_name,
    c_regionname,
   c_nationname,
    c name.
11
    SUM(revenue) AS total_revenue
13 FROM mv_lineitem_orders
    JOIN part ON l_partkey = p_partkey
    JOIN mv_supplier_nation_region ON (s_suppkey = l_suppkey)
    JOIN mv_customer_nation_region ON (c_custkey = o_custkey)
17 WHERE
    s_nationkey <> c_nationkey
    AND p_type = 'LARGE POLISHED STEEL'
19
    AND s_nationname = 'UNITED KINGDOM'
21 GROUP BY
   o_year,
22
   o_quarter,
    o_month,
24
25
    p_type,
   c_regionname,
27
    c nationname.
28
    c_name,
   s_regionname,
29
    s_nationname,
30
31
    s_name
32 ;
```

6 Indexing on Materialized Views

An improvement of the previous optimization is to create indexes on materialized views. Attributes in JOIN, GROUP BY and WHERE clauses were considered in this case too:

```
1 CREATE INDEX IF NOT EXISTS c_nationkey_index
2    ON mv_customer_nation_region (c_nationkey);
3
4 CREATE INDEX IF NOT EXISTS c_regionname_index
5    ON mv_customer_nation_region (c_regionname ASC NULLS LAST);
6
7 CREATE INDEX IF NOT EXISTS c_nationname_index
8    ON mv_customer_nation_region (c_nationname ASC NULLS LAST);
9
10 CREATE INDEX IF NOT EXISTS c_name_index
11    ON mv_customer_nation_region (c_name ASC NULLS LAST);
12
13 CREATE INDEX IF NOT EXISTS s_nationkey_index
14    ON mv_supplier_nation_region (s_nationkey);
15
```

The total size of the database is 22 GB.

7 Horizontal Fragmentation

In the WHERE clause of the query, we are slicing over a specific p_type and a specific s_nationname, which means we are doing a selection over those tables. Moreover, we also want that the supplier nation and customer nations are different. For this reasons, in order to reduce the execution times, we can perform an Horizontal Fragmentation (HF) over tables part, supplier and customer:

```
1 CREATE TABLE part_large_polished_steel AS (
      SELECT *
2
      FROM part
      WHERE p_type = 'LARGE POLISHED STEEL');
6 CREATE TABLE supplier_uk AS (
      SELECT supplier.*
      FROM supplier
      JOIN nation ON s_nationkey = n_nationkey
9
      WHERE n_name = 'UNITED KINGDOM');
10
11
12 CREATE TABLE customer_not_uk AS (
      SELECT customer.*
13
14
      FROM customer
      JOIN nation ON c_nationkey = n_nationkey
      WHERE n_name <> 'UNITED KINGDOM');
```

As we will see in the **Results** section, the HF shows only a slight improvement of the execution times. Then it could be useful to mix it with Vertical Fragmentation (VF), in order to reduce the workload of the query.

7.1 Hybrid and Vertical Fragmentations

Selecting only the attributes that we need in the query may help to speed up the process. We had implemented a VF on the previous tables (performing an Hybrid Fragmentation) and on the tables lineitem and orders:

```
1 CREATE TABLE part_large_polished_steel AS (
2     SELECT p_partkey, p_type
3     FROM part
4     WHERE p_type = 'LARGE POLISHED STEEL');
5     CREATE TABLE supplier_uk AS (
7     SELECT s_suppkey, s_name, s_nationkey
```

```
FROM supplier
       JOIN nation ON s_nationkey = n_nationkey
9
       WHERE n_name = 'UNITED KINGDOM');
10
12 CREATE TABLE customer_not_uk AS (
       SELECT c_custkey, c_name, c_nationkey
13
       FROM customer
14
       JOIN nation ON c_nationkey = n_nationkey
15
       WHERE n_name <> 'UNITED KINGDOM');
16
17
18 CREATE TABLE lineitem_reducted AS (
       SELECT 1_orderkey, 1_partkey, 1_suppkey, 1_extendedprice, 1_discount
       FROM lineitem);
20
21
22 CREATE TABLE orders_reducted AS (
       {\tt SELECT \ o\_orderkey} \;, \; \; {\tt o\_custkey} \;, \; \; {\tt o\_orderdate}
23
       FROM orders):
     The query had to be modified in the following way:
1 WITH supplier_nation_region AS (
     SELECT
       s_suppkey,
       s_name,
       s_nationkey,
      n_name AS s_nationname,
6
       r_name AS s_regionname
     FROM supplier_uk
9
       JOIN nation ON (s_nationkey = n_nationkey)
       JOIN region ON (n_regionkey = r_regionkey)
_{\rm 11} ), customer_nation_region AS (
12
   SELECT
13
      c_custkey,
      c_name,
14
      c_nationkey,
      n_name AS c_nationname,
16
       r_name AS c_regionname
17
    FROM customer_not_uk
      JOIN nation ON (c_nationkey = n_nationkey)
JOIN region ON (n_regionkey = r_regionkey)
19
20
21 ), lineitem_orders AS (
    SELECT
22
23
       l_partkey,
      l_suppkey,
24
      o_orderdate,
25
26
       o_custkey,
       l_extendedprice * (1 - l_discount) AS revenue
27
28
    FROM lineitem_reducted JOIN orders_reducted ON (1_orderkey = o_orderkey)
29 )
30
31 SELECT
     EXTRACT (YEAR FROM o_orderdate) AS o_year,
     EXTRACT (QUARTER FROM o_orderdate) AS o_quarter,
33
    EXTRACT (MONTH FROM o_orderdate) AS o_month,
    p_type,
35
36
     s_regionname,
     s_nationname,
    s_name,
38
39
     c_regionname,
40
    c_nationname,
41
     c_name,
    SUM(revenue) AS total_revenue
43 FROM lineitem_orders
```

```
JOIN part_large_polished_steel ON l_partkey = p_partkey
    JOIN supplier_nation_region ON (s_suppkey = l_suppkey)
45
    JOIN customer_nation_region ON (c_custkey = o_custkey)
46
47 GROUP BY
    o_year,
48
49
    o_quarter,
    o_month,
50
51
    p_type,
52
    c_regionname,
    c_nationname,
53
54
    c_name,
    s_regionname,
    s_nationname,
56
57
    s_name
58 ;
```

8 Results

In the following table the execution times of 10 runs are shown, comparing the standard query with the optimizations described previously:

Execution Times						
Standard	Indexed Tables	Materialized Views	Indexed MV	HF & VF		
37.511	34.305	14.331	15.325	18.665		
37.538	34.257	13.905	15.320	17.892		
38.919	34.429	13.754	15.286	17.492		
37.754	33.626	14.051	15.364	17.543		
38.490	34.099	14.372	15.296	18.108		
38.747	33.136	13.919	15.367	17.640		
37.958	34.280	13.800	15.899	18.526		
37.479	33.720	14.396	15.372	17.955		
37.620	33.620	13.714	15.404	17.384		
38.214	34.382	13.803	15.938	17.628		

Table 1: Execution times are in seconds (s)

In the optimization using horizontal fragmentation on tables, we have said that just a slight improvement on the performances of the query can be observed. In the following table we show the results of 5 runs with HF:

Horizontal Fragmentation						
36.845	37.112	36.796	37.129	36.994		

Table 2: Execution times are in seconds (s)

9 Conclusions

The goal of this study is to create a database, implement the query proposed and then find the best optimization in order to achieve the best running times, without exceeding the original database size by 1.5.

We can compare the performance of the standard query with that of its optimizations in the last table, which reports means and standard deviations of execution times and sizes of the database:

	Standard	Indexed Tables	Materialized View	Indexed MV	HF & VF
μ	38.027	34.085	14.004	15.457	17.883
σ	0.511	0.410	0.254	0.233	0.415
size	18	21	21.5	22	18

Table 3: Execution times are in seconds (s), sizes are in giga-bytes (GB)

As we can observe in the table above, the best improvement in performance is reached by the query that uses materialized views, with a significant reduction in execution times. Even though we expected a further improvement using indexes on materialized views, the results shows a slight increase in executions. The fragmentation also shows very good performance improvements with the positive effect of keeping the database size constant.

In conclusion, the best optimization for this query is the one that uses materialized views that speeds the execution times approximately of 2.7 times and only increases the database size of about 1.2 times.