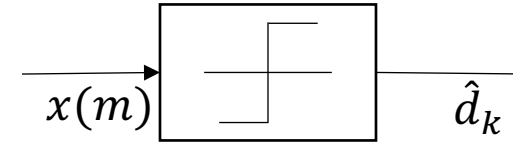
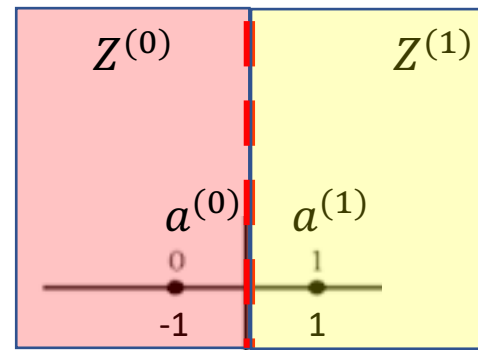


# Decision strategy

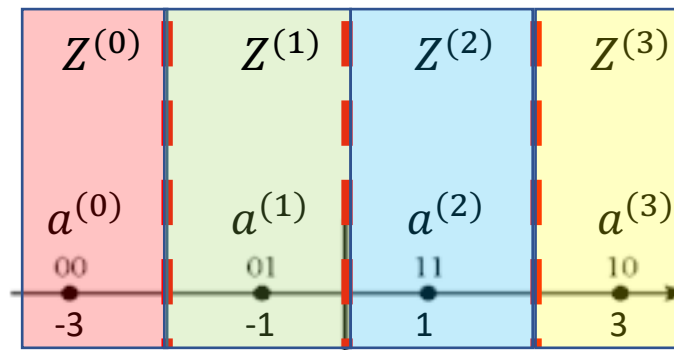


- Adopting the maximum likelihood criterion, we can partition the signal space in *zone of decisions*, where zone  $Z^{(i)}$  is the set of points that are closer to the symbol  $a^{(i)}$  than to any other symbol

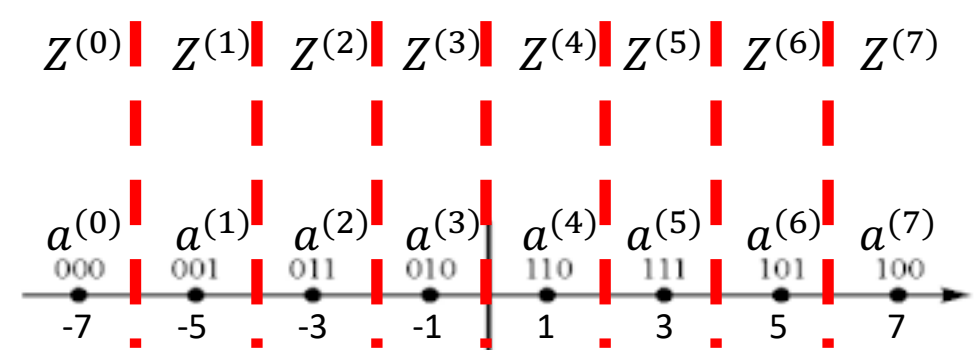
$$Z^{(i)} = \{x | d(x, a^{(i)}) < d(x, a^{(j)}), j \neq i, j = 1, \dots, M\}$$



$M = 2$   
 $m = 1$



$M = 4$   
 $m = 2$



$M = 8$   
 $m = 3$

The decision threshold are in the midpoints of the segment connecting any two adjacent symbols. For example, for  $M = 4$  the thresholds are in  $-2, 0$  and  $2$ .

# PAM error probability

- Even if the maximum likelihood decision strategy is optimal, the receiver still make errors due to the presence of noise.
- The error probability is averaged over all the symbol of the constellation

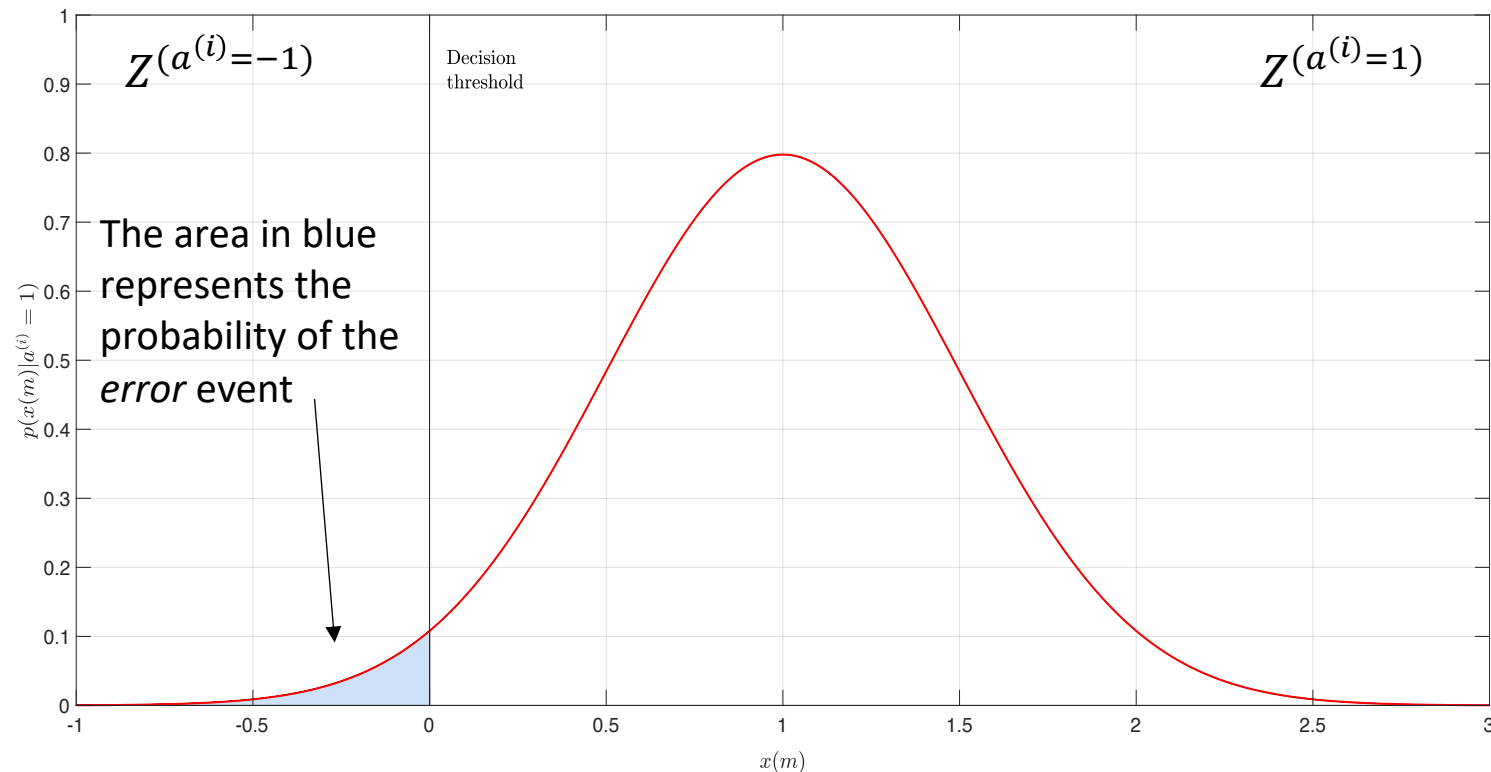
$$P_e = \lim_{N^{(s)} \rightarrow \infty} \frac{N_e^{(s)}}{N^{(s)}} = \frac{1}{M} \sum_{i=0}^{M-1} P(e|a^{(i)})$$

where  $N_e^{(s)}$  is the number of symbol errors and  $N^{(s)}$  is the number of transmitted symbols.

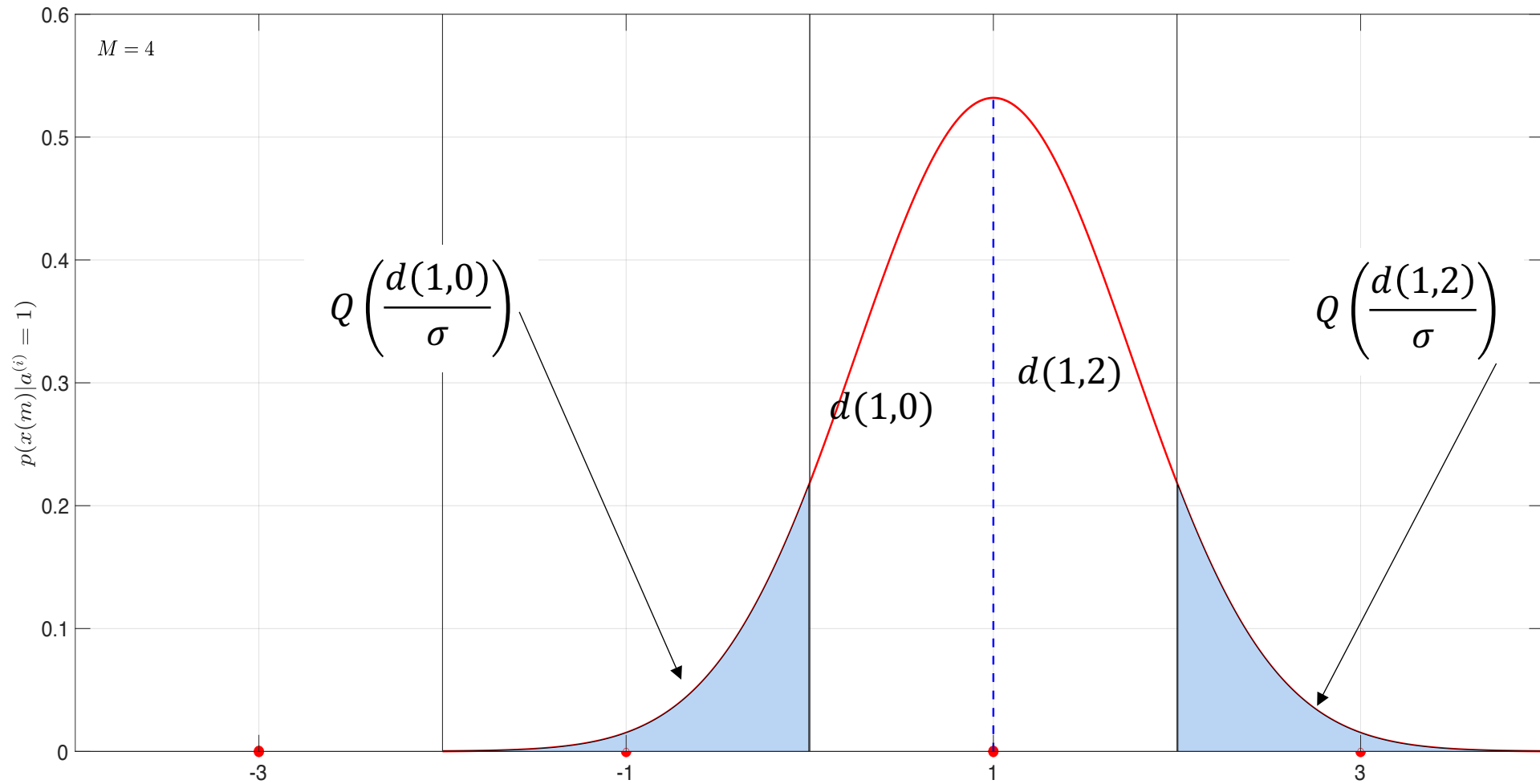
- The probability of error  $P(e|a^{(i)})$  is the probability that, having transmitted  $a^{(i)}$ , the decision variable  $x(m)$  does not fall in the decision region  $Z^{(i)}$ .

# PAM error probability

- To compute  $P(e|a^{(i)})$  we assume that the transmitted symbol is  $a_m = a^{(i)}$ , so that it is  $x(m) = a^{(i)} + n(m)$  and the probability of error is
$$P(e|a^{(i)}) = \Pr\{x(m) \notin Z^{(i)} | a_m = a^{(i)}\}$$



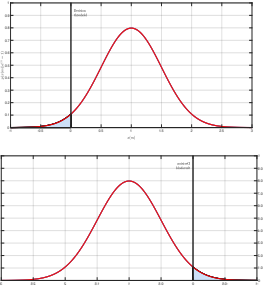
# PAM error probability



$$P(e|a^{(i)} = 1) = \int_{-\infty}^0 p(x|a^{(i)} = 1)dx + \int_2^{+\infty} p(x|a^{(i)} = 1)dx = Q\left(\frac{d(1,0)}{\sigma}\right) + Q\left(\frac{d(1,2)}{\sigma}\right) = 2Q\left(\frac{1}{\sigma}\right)$$

# PAM error probability: $Q$ -function

- The  $Q$ -function computes the integral of the *tail* of a Gaussian distribution.
- The probability that  $x \in \mathcal{N}(m, \sigma^2)$  is smaller than  $t_1$  or larger than  $t_2$  are the integral of Gaussian tails and they are computed as


$$\left. \begin{aligned} \int_{-\infty}^{t_1} pdf(x) dx &= Q\left(\frac{m - t_1}{\sigma}\right) \\ \int_{t_2}^{+\infty} pdf(x) dx &= Q\left(\frac{t_2 - m}{\sigma}\right) \end{aligned} \right\} = Q\left(\frac{d(t_i, m)}{\sigma}\right), i = 1, 2$$

- In our case,  $m$  is the symbol  $a^{(i)}$  and  $t_1$  or  $t_2$  are the detection thresholds.
- The main properties of the  $Q$ -function are
$$Q(-\infty) = 1, Q(\infty) = 0, Q(0) = 0.5, Q(-x) = 1 - Q(x).$$

# PAM error probability

- 2-PAM

$$P_e^{(2-PAM)} = \frac{1}{2} \left( Q \left( \frac{d(-1,0)}{\sigma} \right) + Q \left( \frac{d(1,0)}{\sigma} \right) \right) = Q \left( \frac{1}{\sigma} \right)$$

- 4-PAM

$$P_e^{(4-PAM)} = \frac{1}{4} \left( Q \left( \frac{d(-3,-2)}{\sigma} \right) + Q \left( \frac{d(-1,-2)}{\sigma} \right) + Q \left( \frac{d(-1,0)}{\sigma} \right) + Q \left( \frac{d(1,0)}{\sigma} \right) + Q \left( \frac{d(1,2)}{\sigma} \right) + Q \left( \frac{d(3,2)}{\sigma} \right) \right) = \frac{3}{2} Q \left( \frac{1}{\sigma} \right)$$

# PAM symbol error probability

- It is often useful to express the  $P_e$  in terms of  $E_s/N_0$ .

- 2-PAM:  $E_s = \frac{2^2-1}{6} = \frac{1}{2} \Rightarrow 2E_s = 1$  and  $\sigma^2 = N_0$ , and  $\frac{1}{\sigma} = \sqrt{\frac{1}{\sigma^2}} = \sqrt{\frac{2E_s}{N_0}}$ .

$$P_e^{(2-PAM)} = Q\left(\sqrt{\frac{2E_s}{N_0}}\right)$$

- 4-PAM:  $E_s = \frac{4^2-1}{6} = \frac{5}{2} \Rightarrow \frac{2}{5}E_s = 1$ , and  $\frac{1}{\sigma} = \sqrt{\frac{1}{\sigma^2}} = \sqrt{\frac{2E_s}{5N_0}}$ .

$$P_e^{(4-PAM)} = \frac{3}{2} Q\left(\sqrt{\frac{2E_s}{5N_0}}\right)$$

# PAM bit error probability

- To have a fair comparison, the modulation performance are expressed in terms of *bit error probability*  $P_e^{(b)}$  as function of  $E_b/N_0$ .
- The energy  $E_b$  per bit is computed as the energy per symbol divided by the number of bits per symbol

$$E_b = \frac{E_s}{\log_2 M}$$

- Although one symbol carries  $\log_2 M$  bits, it is reasonable to assume that in a well-designed system (*Gray mapping* and medium-high SNR) *a symbol error causes only one-bit errors*.
- If  $N^{(b)}$  and  $N_e^{(b)}$  are the number of transmitted bits and the number of bit errors, the bit error probability is computed as

$$P_e^{(b)} = \lim_{N^{(b)} \rightarrow \infty} \frac{N_e^{(b)}}{N^{(b)}} \approx \lim_{N^{(s)} \rightarrow \infty} \frac{N_e^{(s)}}{\log_2 M N^{(s)}} = \frac{1}{\log_2 M} \lim_{N^{(s)} \rightarrow \infty} \frac{N_e^{(s)}}{N^{(s)}} = \frac{1}{\log_2 M} P_e.$$



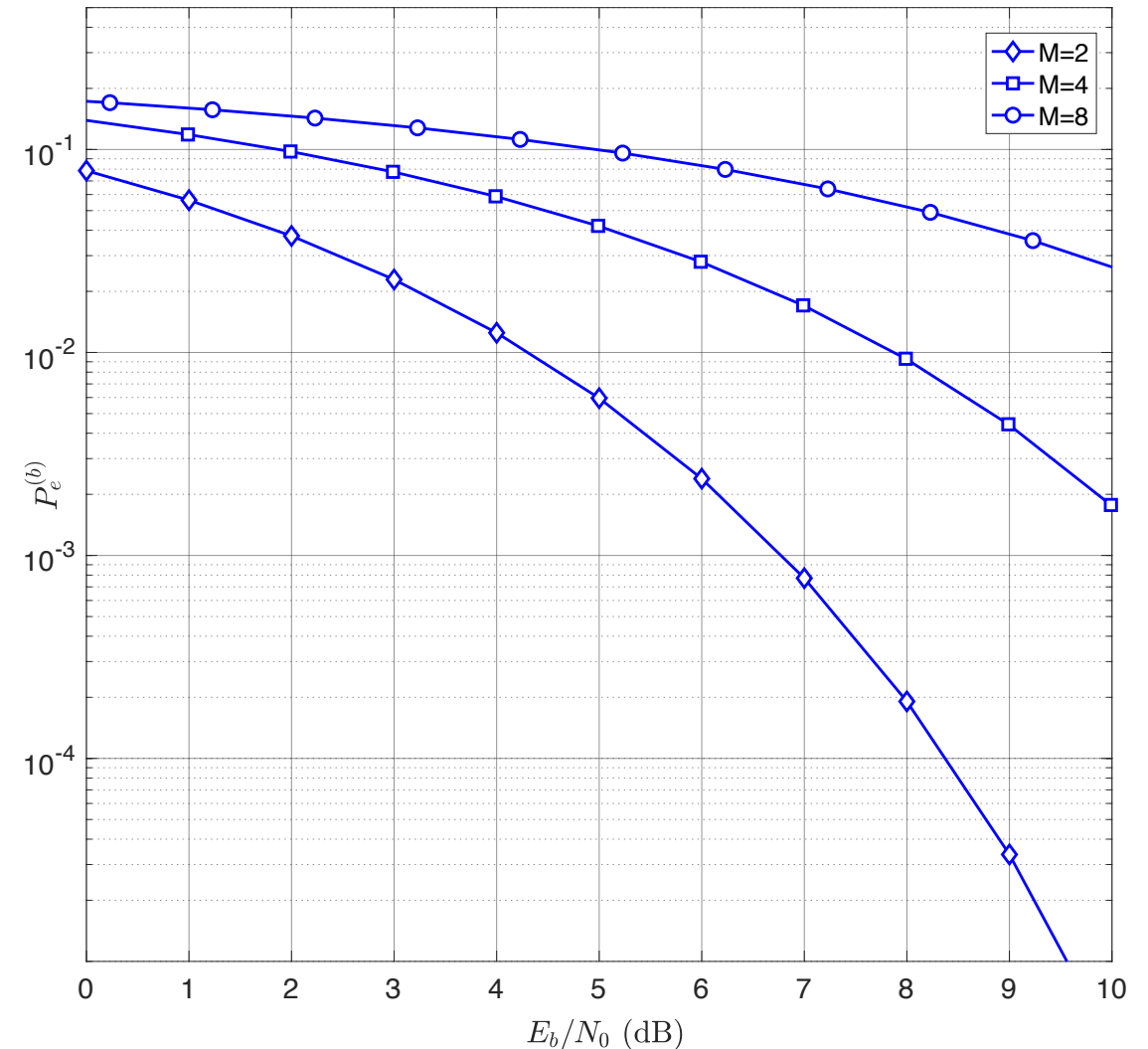
# PAM bit error probability

- 2-PAM:  $M = 2$ ,  $m = 1$  bit per symbol  $\Rightarrow P_e^{(b)} = P_e, E_b = E_s$

$$P_e^{(2-PAM),b} = Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

- 4-PAM:  $M = 4$ ,  $m = 2$  bit per symbol  $\Rightarrow P_e^{(b)} = \frac{1}{2}P_e, E_b = \frac{1}{2}E_s$

$$P_e^{(4-PAM),b} = \frac{3}{4}Q\left(\sqrt{\frac{4E_b}{5N_0}}\right)$$



# Digital communications

## Quadrature modulations (QAM)

# Quadrature modulations

- In analog modulations, QAM is obtained by transmitting two orthogonal DSB signals  $m_I(t)$ ,  $m_Q(t)$  and the complex envelope is

$$\tilde{s}_{QAM}(t) = m_I(t) + jm_Q(t)$$

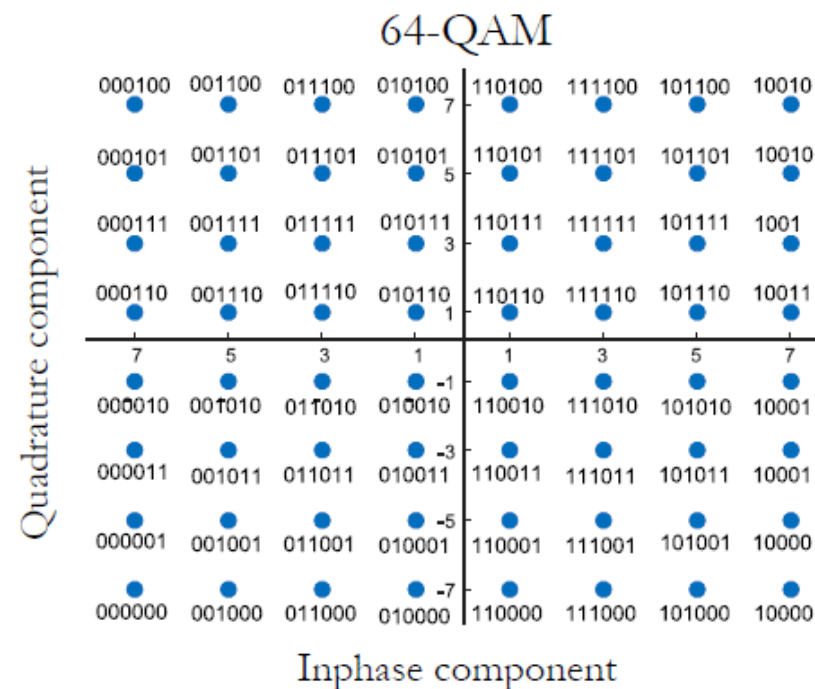
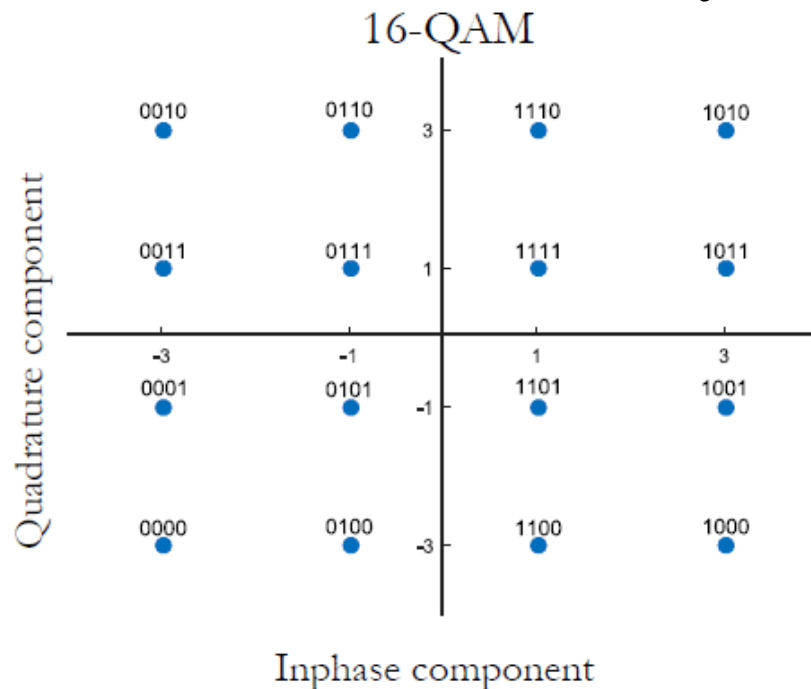
- Quadrature PAM is obtained exactly in the same manner by transmitting two PAM signals in quadrature  $m_I(t) = \sum_i a_i g_T(t - iT)$  and  $m_Q(t) = \sum_i b_i g_T(t - iT)$ , with  $a_i, b_i$  PAM symbols.
- The QAM signal is

$$s_{QAM}(t) = \sum_i (a_i + jb_i) g_T(t - iT) = \sum_i c_i g_T(t - iT)$$

and the QAM complex symbols take the form  $c_i = a_i + jb_i$ .

# QAM symbols

- Because QAM is the combination of two orthogonal PAM, the values of  $M_{QAM} = M_{PAM}^2$  are squared powers of 2, i.e.  $m$  is always even.
  - If the two PAMs have  $M_{PAM} = 4$  symbols than the QAM has  $M_{QAM} = 16$  symbols, if  $M_{PAM} = 8$  than  $M_{QAM} = 64$ .



# Energy of a QAM symbol

- In the computation of power and energy the only difference between PAM and QAM is in the mean square value of the symbols.
- Keeping in mind that the in-phase and quadrature symbols are independent and zero-mean, it is

$$A = E\{c_i c_i^*\} = E\{a_i^2\} + E\{b_i^2\} = 2 \frac{M_{PAM}^2 - 1}{3} = 2 \frac{M_{QAM} - 1}{3}$$

- The energy per symbol is

$$E_s = \frac{A}{2} = \frac{M_{QAM} - 1}{3}$$

- QAM constellation is much more compact and requires less energy per symbol compared to PAM.

- $A^{(4-PAM)} = \frac{16-1}{3} = 5;$        $A^{(4-QAM)} = 2 \frac{4-1}{3} = 2$
  - $A^{(16-PAM)} = \frac{256-1}{3} = 85;$        $A^{(16-QAM)} = 2 \frac{16-1}{3} = 10.$

# QAM error probability

- The *complex* decision variable is

$$\begin{aligned}x(m) &= c_m + n(m) = (a_m + jb_m) + (n_I(m) + jn_Q(m)) \\ &= a_m + n_I(m) + j(b_m + n_Q(m))\end{aligned}$$

- The in-phase and quadrature noise components  $n_I(m)$  and  $n_Q(m)$  are independent.
- Error events depends on noise. If the noise is independent also the error events on the two components are independent.
- The error probability can be *approximated* as the sum of the probability of making an error on the in-phase symbol  $a_m$  and probability of making an error on the quadrature symbol  $b_m$ .

# 4-QAM error probability

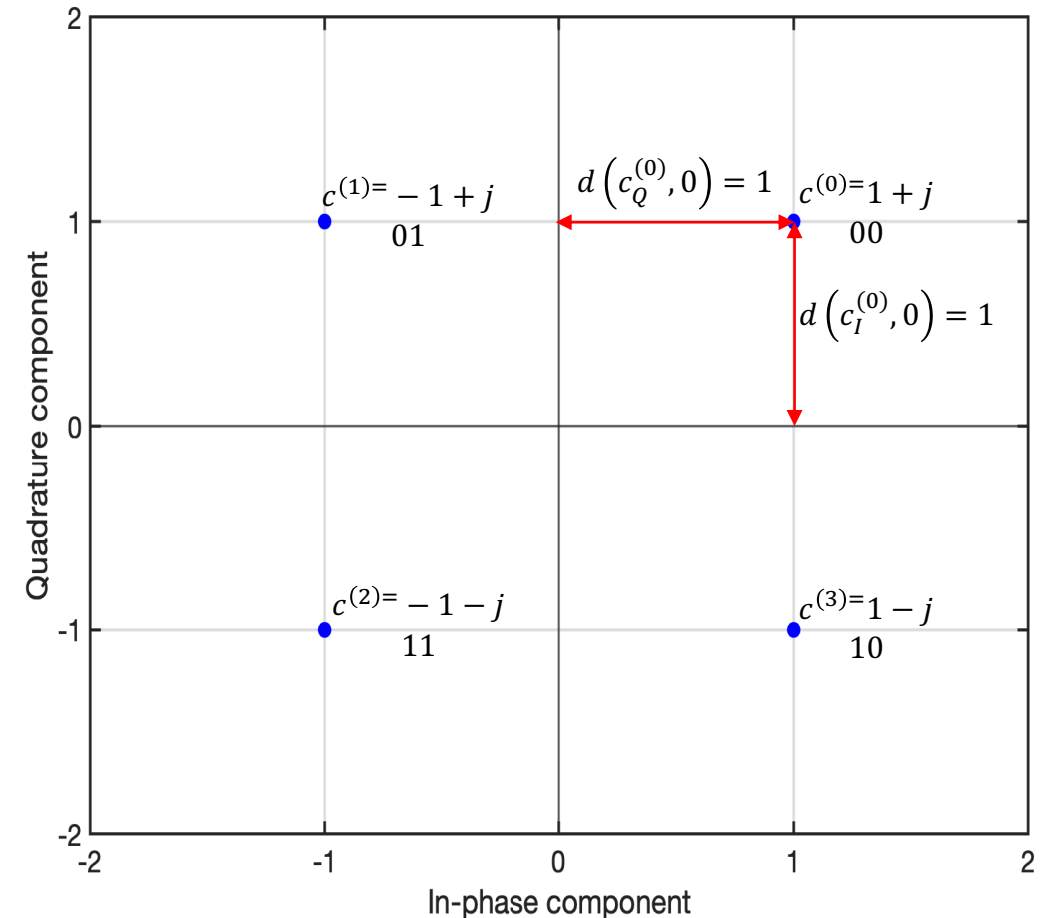
- 4-QAM is obtained as the composition of two 2-PAM in quadrature.
- The symbol error probability is

$$P_e^{(4-QAM)} = \frac{1}{4} \sum_{i=0}^3 P(e|c^{(i)}) = P(e|c^{(0)})$$

$$P(e|c^{(0)}) \approx Q\left(\frac{d(c_I^{(0)}, 0)}{\sigma_{n_I}}\right) + Q\left(\frac{d(c_Q^{(0)}, 0)}{\sigma_{n_Q}}\right)$$

$$= 2Q\left(\frac{1}{\sigma}\right)$$

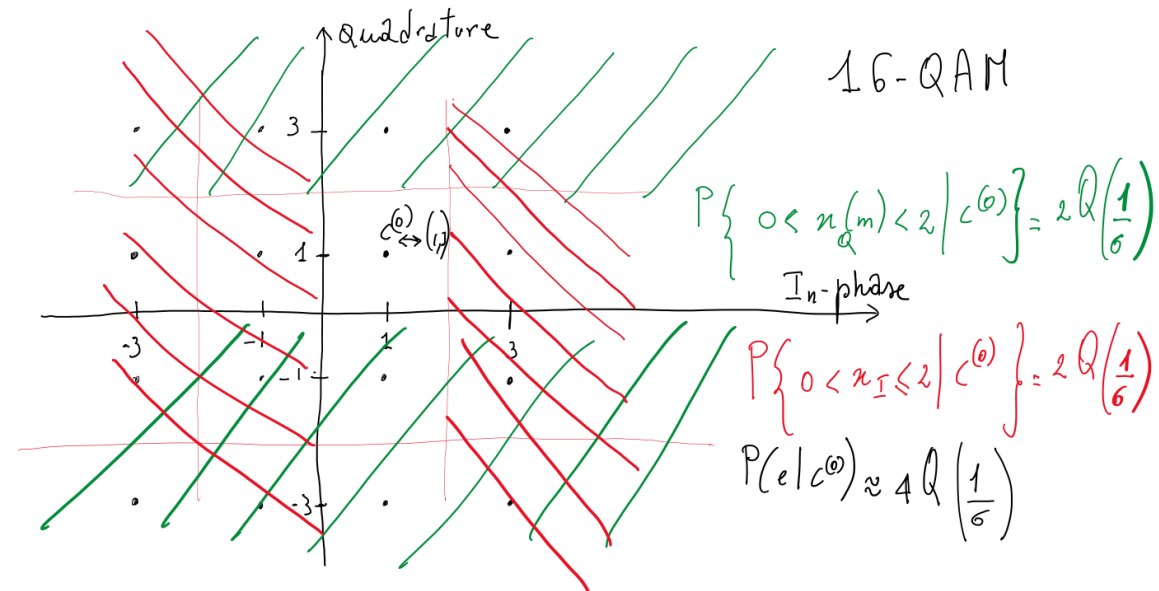
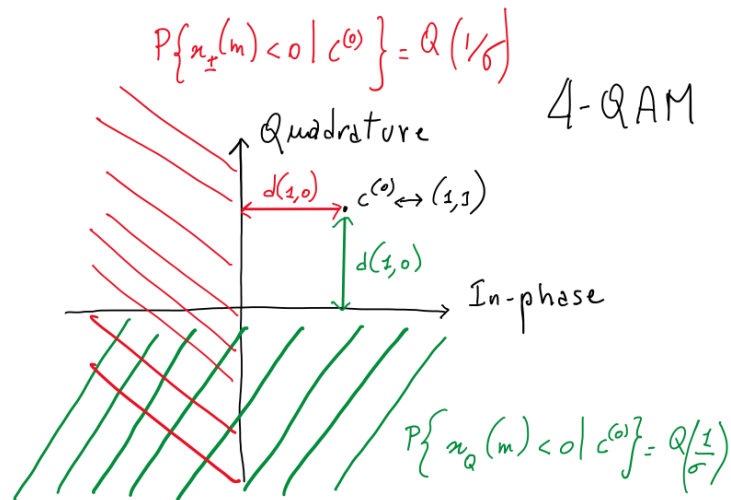
$$P_e^{(4-QAM)} \approx 2P_e^{(2-PAM)}$$



# M-QAM error probability

- M-QAM is obtained as the composition of two PAM in quadrature, each with  $\sqrt{M}$  symbols.
- The symbol error probability can always be approximated as

$$P_e^{(M\text{-QAM})} \approx 2P_e^{(\sqrt{M}\text{-PAM})}$$





# QAM symbol error probability

- 4-QAM:  $E_s = \frac{4-1}{3} = 1 \Rightarrow E_s = 1, M = 4, m = 2$  and  $\frac{1}{\sigma} = \sqrt{\frac{1}{\sigma^2}} = \sqrt{\frac{E_s}{N_0}}$ .

$$P_e^{(4-QAM)} \approx 2Q\left(\sqrt{\frac{E_s}{N_0}}\right);$$

16-QAM:  $E_s = \frac{16-1}{3} = 5 \Rightarrow \frac{1}{5}E_s = 1, M = 16, m = 4$  and  $\frac{1}{\sigma} = \sqrt{\frac{1}{\sigma^2}} = \sqrt{\frac{E_s}{5N_0}}$ .

$$P_e^{(16-QAM)} \approx 2\frac{3}{2}\left(\sqrt{\frac{1}{\sigma}}\right) = 3Q\left(\sqrt{\frac{E_s}{5N_0}}\right); P_e^{(16-QAM),b} \approx \frac{1}{4}3Q\left(\sqrt{\frac{4E_b}{5N_0}}\right) = \frac{3}{4}Q\left(\sqrt{\frac{4E_b}{5N_0}}\right)$$

# M-QAM error probability

- The total number of M-QAM transmitted bits is the sum of the number of bits transmitted on the in-phase and quadrature channels.
- Because each channel is independent, the bit error probability per channel is independent.
- Accordingly,  $P_e^{(M-QAM),b}$  can be *exactly* computed as the sum of the bit error probability on the in-phase channel and the quadrature channel, divided by two.

$$P_e^{(M-QAM),b} = \frac{1}{2} 2P_e^{(\sqrt{M}-PAM),b} = P_e^{(\sqrt{M}-PAM),b}$$

# QAM bit error probability

- 4-QAM:

$$P_e^{(4-QAM),b} = P_e^{(2-PAM),b} \approx Q\left(\sqrt{\frac{2E_b}{N_0}}\right)$$

- 16-QAM

$$P_e^{(16-QAM),b} = P_e^{(4-PAM),b} \approx \frac{3}{4} Q\left(\sqrt{\frac{4E_b}{5N_0}}\right)$$

