

Interdomain routing

BGP-4

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BGP-4

- Border Gateway Protocol v4 (rfc4271) *VANICABEP*
 - inter-AS routing protocol (EGP)
 - inter-AS routing usually reflects **political and business relationships** between the ISP and organizations involved
 - intra-AS routing is optimized in accordance with the required technical demands
 - *How do I filter routing updates coming from a particular neighbor AS?*
 - *How do I make sure that I use this link or this provider rather than another one?*
- Exchange **network reachability information** between BGP speaking peers
- Supports **destination-based** forwarding only (CIDR)

SECURENT ROUTING:

SI PUÒ FAIRE TUTTO CON SOURCE Routing, SE NON È POSSIBILE EFFICIENTE

TRADITIONAL SOURCE
ROUTING

FORWARDING BASATO SUL
SU DESTINATION ADDRESS.

SECURENT ROUTING NASCE PERMETTE'

Vedere Source Routing IN modo

EFFICIENTE E PIÙ SEGRETE BASATO → dentro Pw' corrente del TRAFFICO

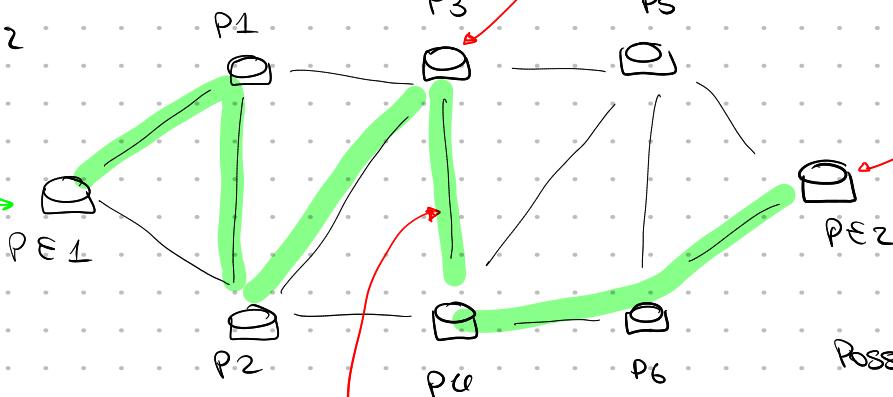
SUL MPLS.

Vedere CHE PASSI PER P3 E P4.

PATH: PE1 - PE2

DA: PE2

[PL]



Rosso ASSISTANTE

LABEL A CADA O

ROUTER

USARES MPLS SE IN MODO DIFFERITO.

LE LABEL POSSONO AVERE UN SCOPO LOCALE O GLOBAL,

DOBBIALE ASSICURARSI CHE IL LABEL 103 NON

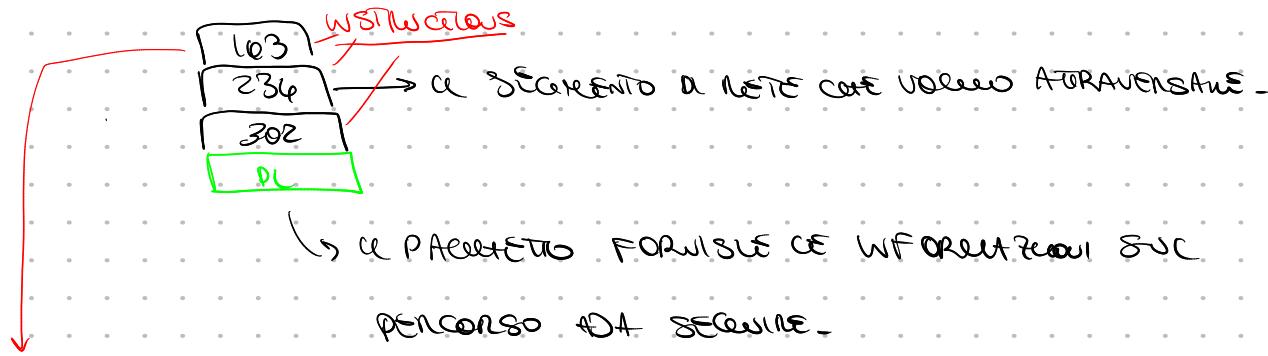
SIA ASSOCIAVATA AD UN ALTRIO ROUTER DELLA RETE.

LE LABEL POSSONO ANCHE AVERE SCOPO
LOCALE, COME NEL MPLS

→ TUTTI I ROUTER
SARanno CHE QUEL
ROUTER CHIAVE E'
UNIVOCAMENTE
IDENTIFICATO DALLA
LABEL

CON MPLS AVREMO CONCERNATO UN CSP CON RSVP - TE SETTANDO UN
ERO A TIPO COSE -

Con segmento multice non ho la spif di confronta LSP, ma faccio
direttamente a RSP del segmento del percorso.



Punto d'arrivo della strada, 103 identifica P3 e accoda come quale
è il successore next-hop per P3.

P3 riconosce che fa parte del secondo segmento, fa il pop e continua.

NON ABBIANO BANDWIDTH RESERVATION !!!

VALUTAZIONE: NON ABBIANO BISOGNO DI RSVP-TE



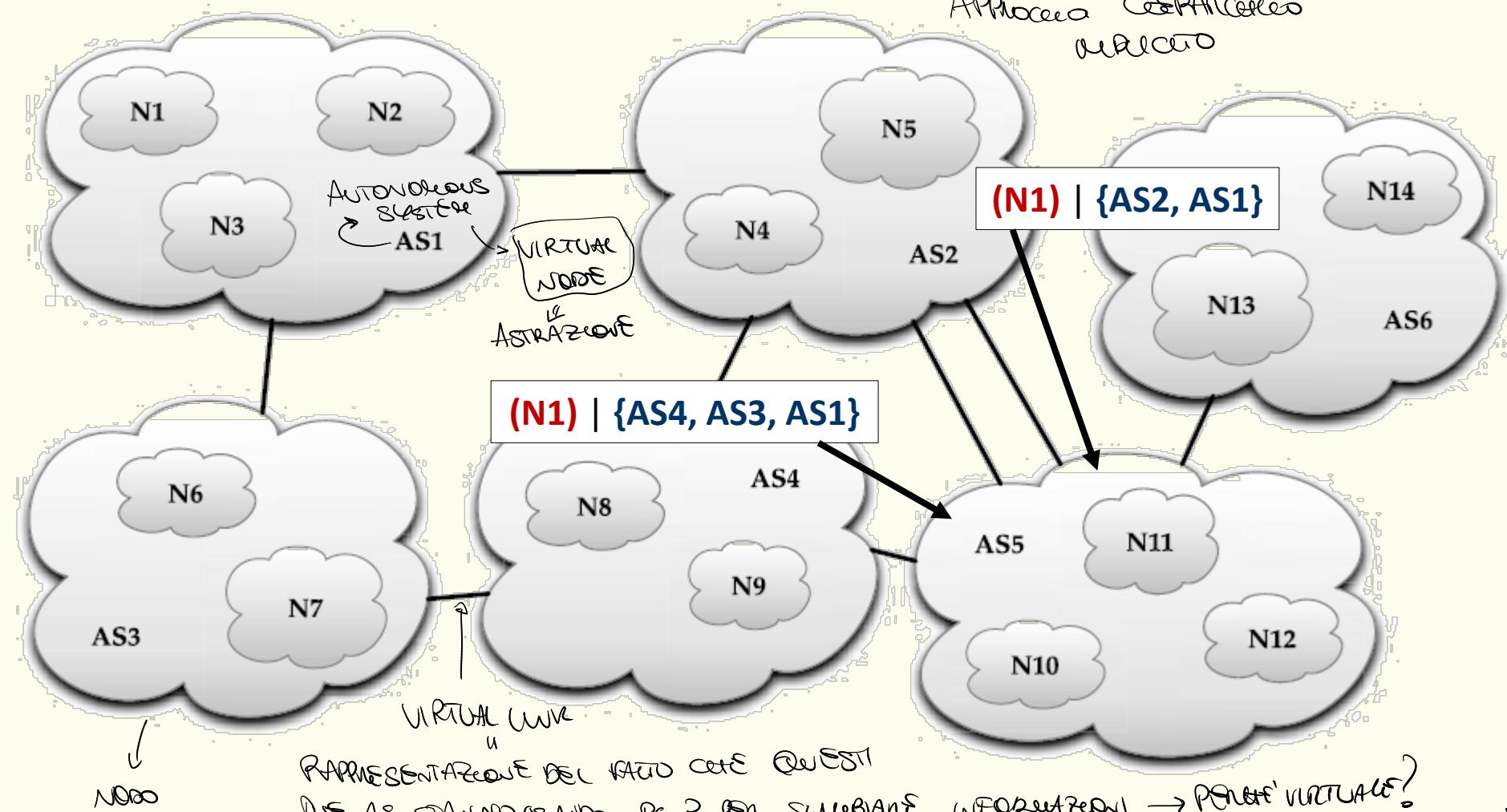
Path vector routing

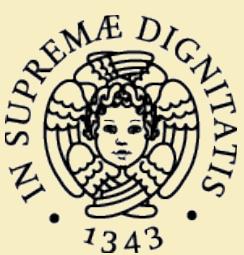
DISTANCE VECTOR → WORKING IN VARIOUS CASES OF NETS WE DON'T KNOW THE COSTS FOR
ROUTING

→ PW SCALABLE
PW OVERHEAD
CENTRALIZED & CONVERGENCE

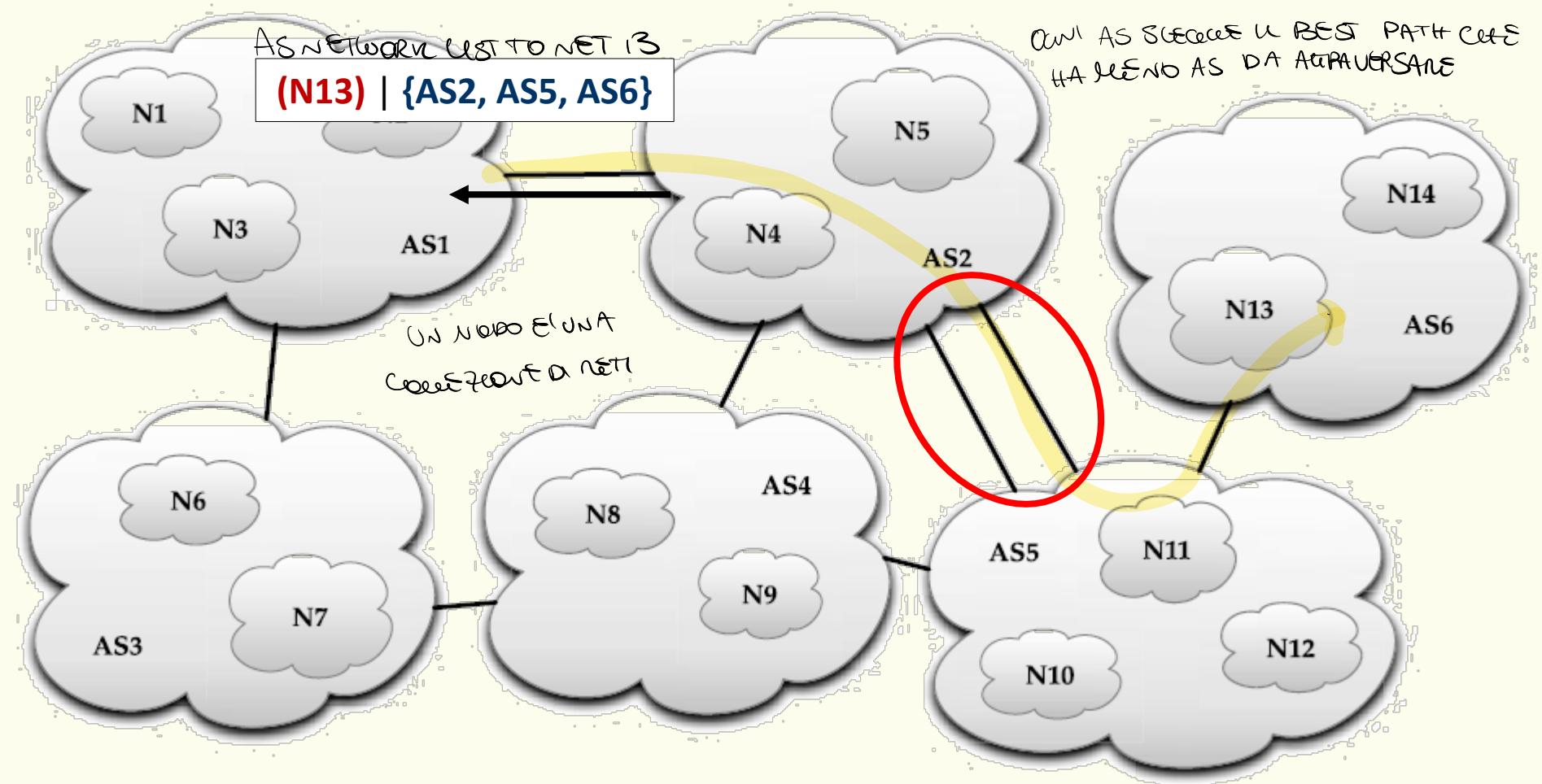
DOBBIA ASSURDARE CHE L'ALGORITMO CONVERGE

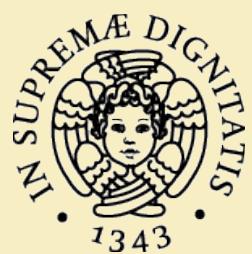
PROCESSO CONVERGENZA
REFLEXO



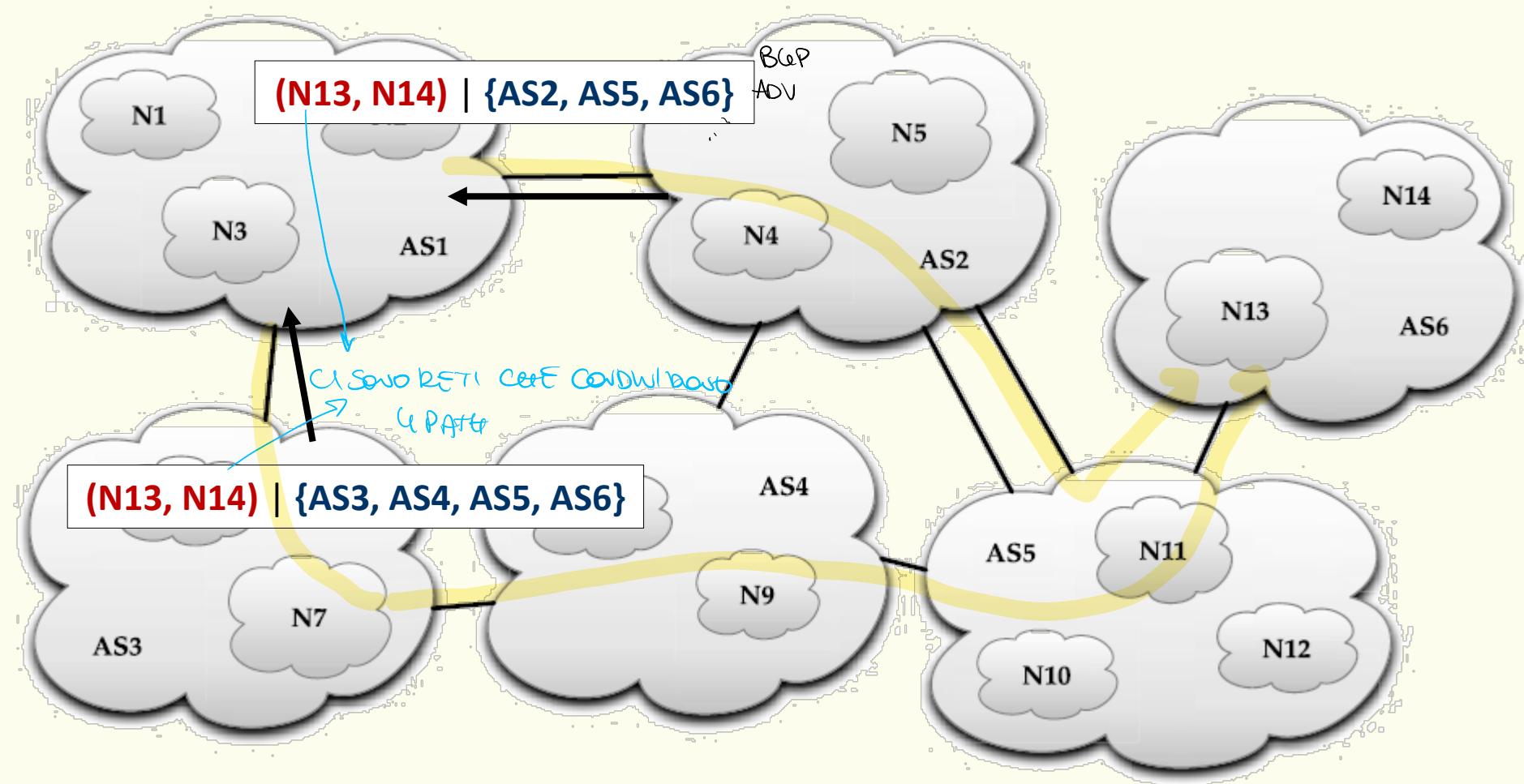


AS hop count metric



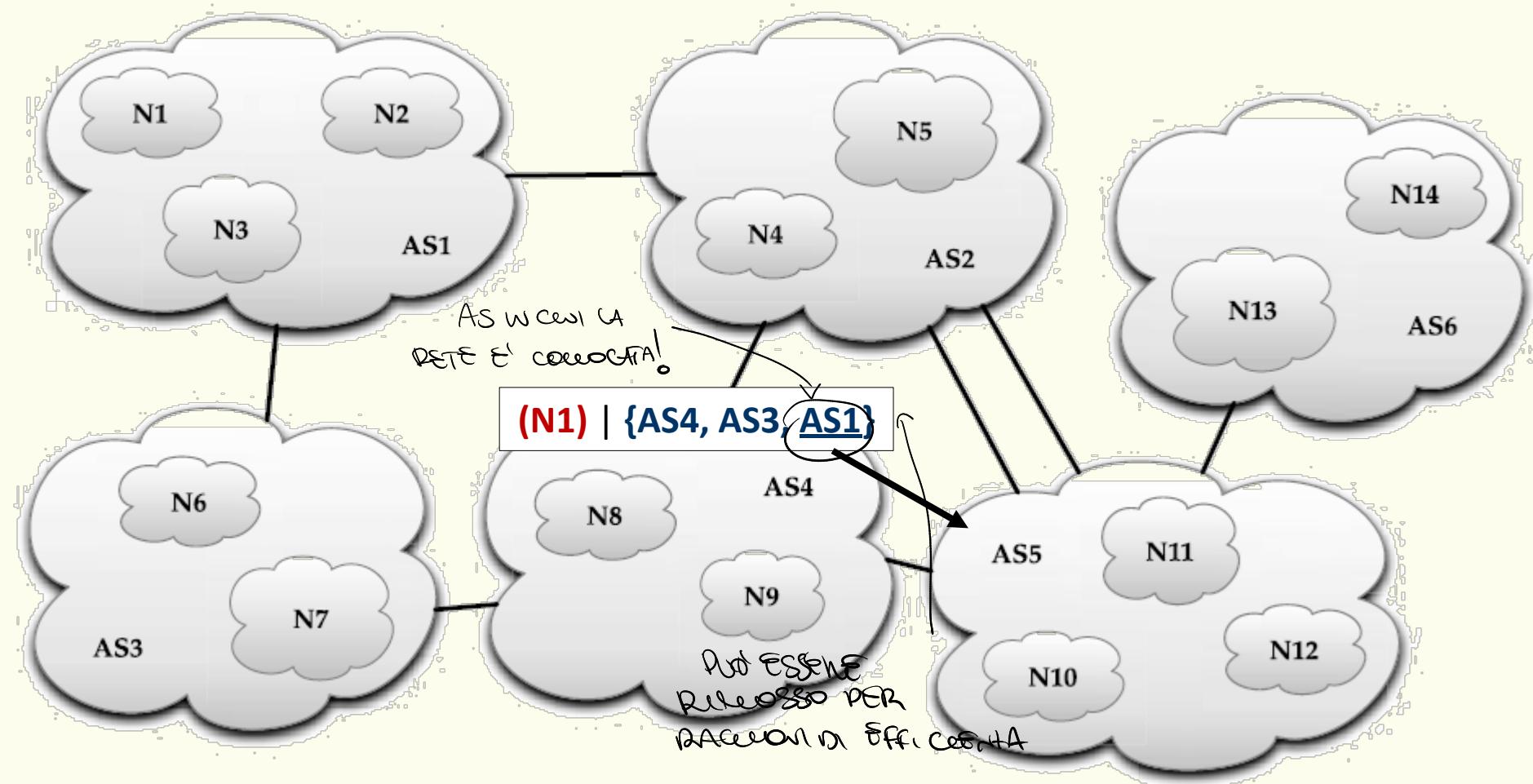


Set of networks advertisement



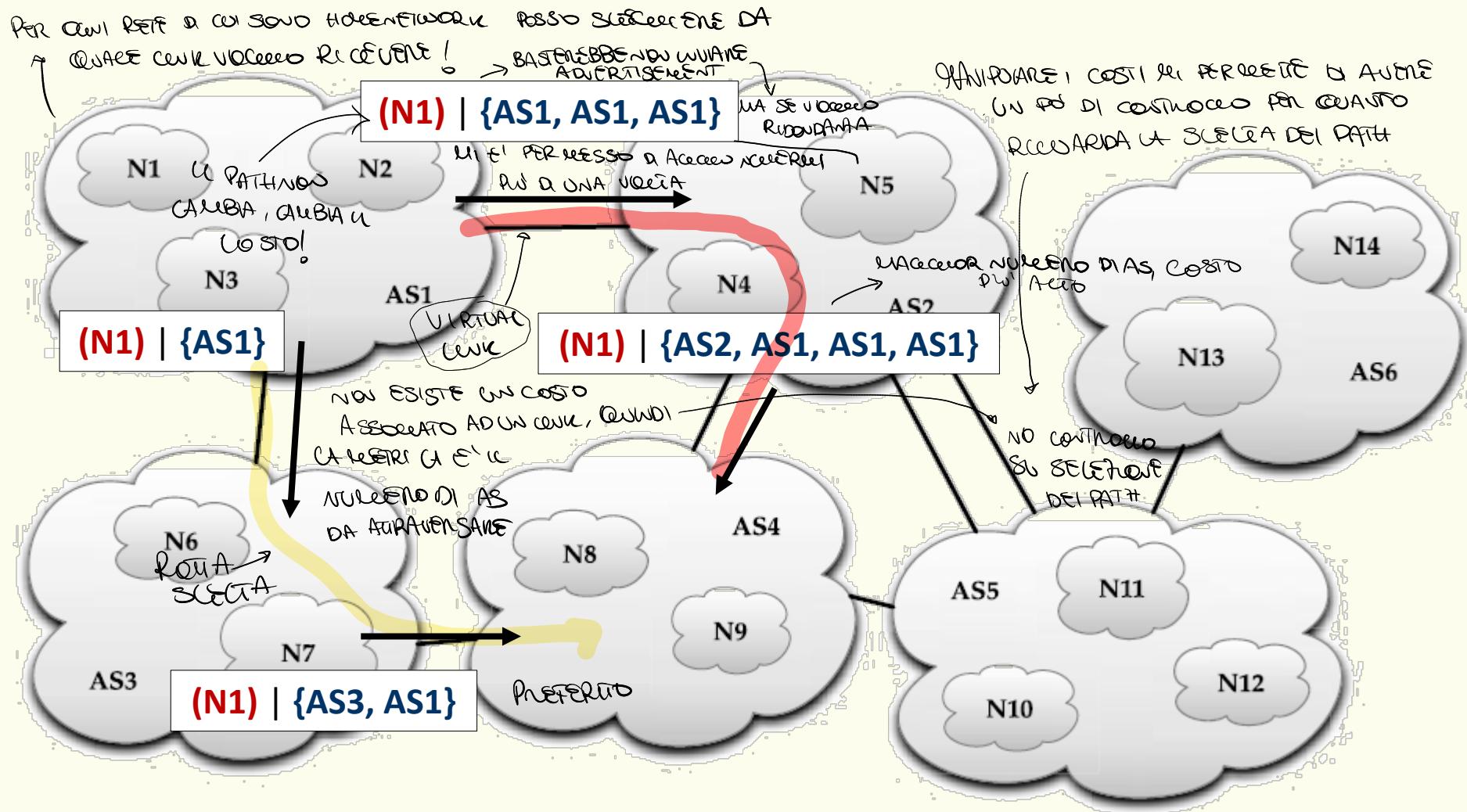


Home AS advertisement





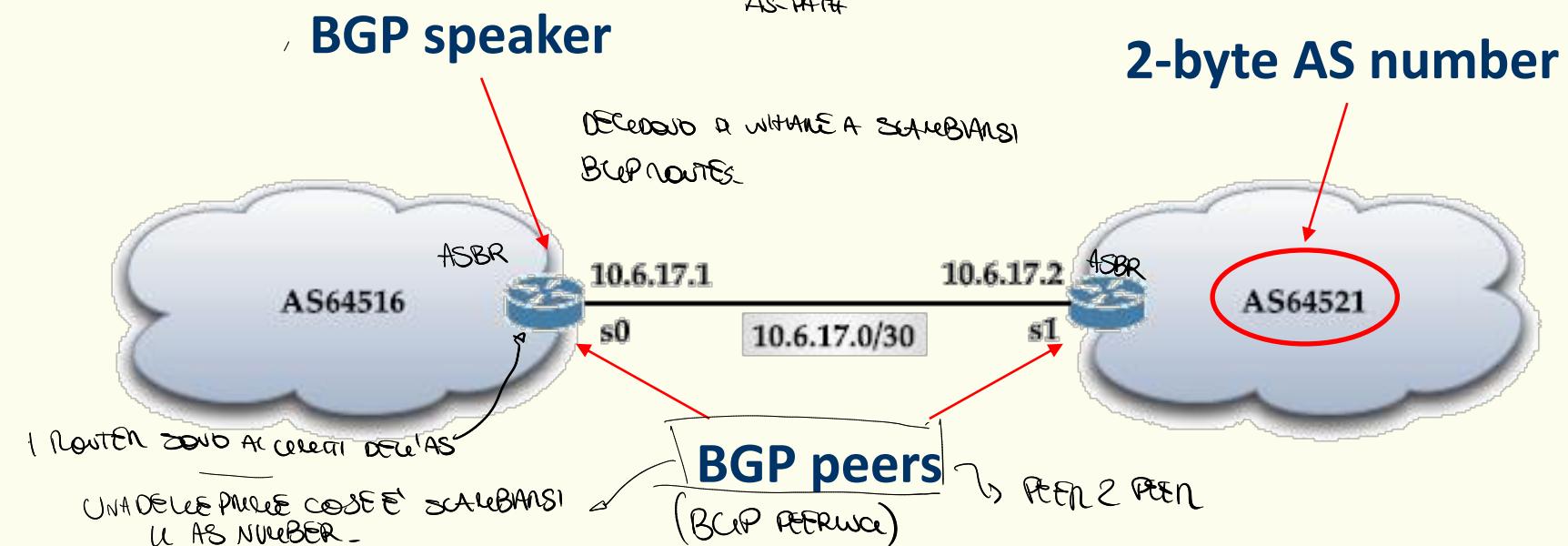
Path manipulation



Terminology

Network = IP prefix (A.B.C.D/n) → network layer reachability information (NLRI)

BGP route: IP prefix destination(s) → attributes of an AS-path by a receiving AS through an UPDATE message



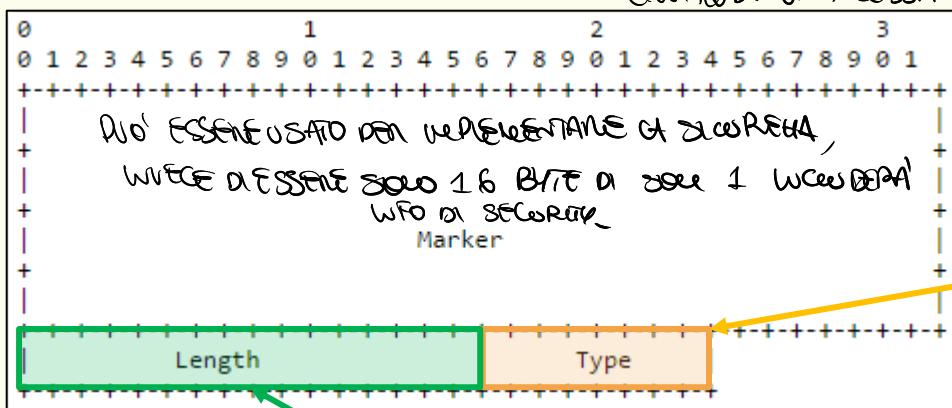
AS number space is divided into **public** (1-64511) and **private** (64512-65534)

4-byte AS number space is also defined → FUTURE, SEE IANA FACTORY

BGP operations



- Four BGP message types (+ 1) exchanged over TCP (port 179)
Now we have one stream of BGP



≤ 4096

- OPEN
 - UPDATE
 - NOTIFICATION
 - KEEPALIVE
 - ROUTE-REFRESH (optional)

OPTIONAL

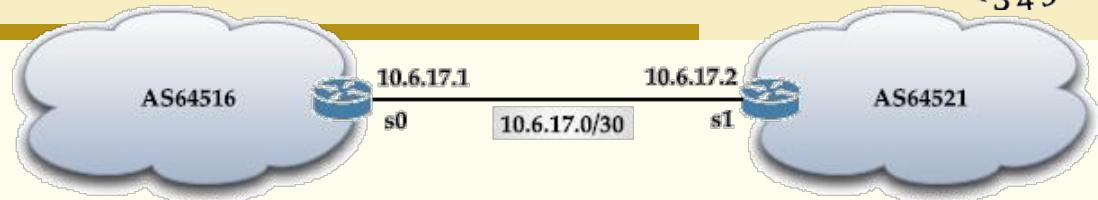
NON OBBLIGATORIO DA SUPPORTARE PER I
RESP SPEAKER.



BGP operations

- **OPEN message**

- First message sent to configured BGP peer (contains AS number)



- **KEEPALIVE**

- Exchanged periodically, approximately three times every *hold time* at most

- **NOTIFICATION**

- Closes a BGP session

SETUP BGP session →

Alle SPEAKER Reconoscono SE
FANNO PARTE DELLO stesso AS ↗

Ora AS DIVERSI ↴

(THREEs) =) now C'è NIENTE COSA
S'AGGIUNTA.

LE NUCLEI SONO
DIVERSE

→ SE c'è uno SPAKER E' ANCORA ATTIVO O NO

↪ Quel rotta che PASSAVA DA QUEST'AS VIENE
INTONATA

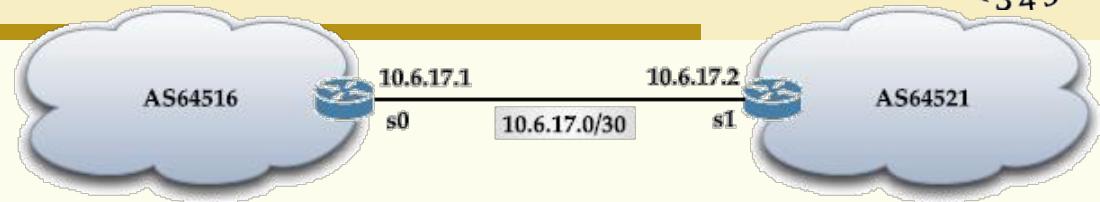
L'unico VIRTUALE rotto



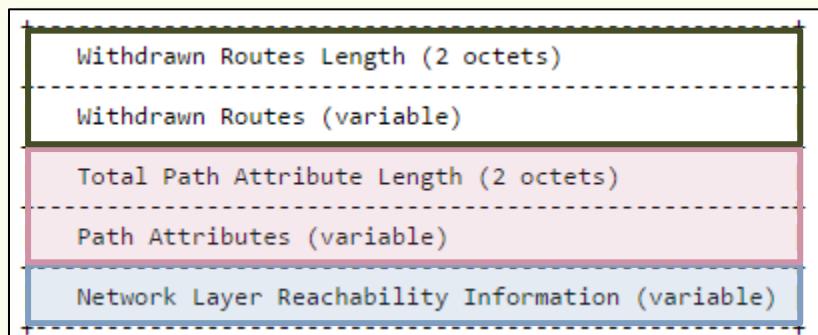
BGP operations

- **UPDATE message**

- Sent to exchange information about *networks*
- Works in PUSH mode
- Announced routes must be explicitly withdrawn



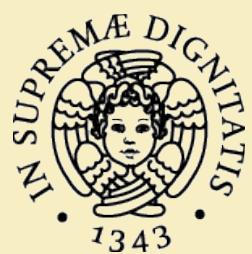
TRANSPORT BGP ROUTE → SET di PREF(SSI) IP CON ATTRIBUTI



A sequence of path attributes (TLV)

- ORIGIN
- AS_PATH
- NEXT_HOP
- LOCAL_PREF
- ...

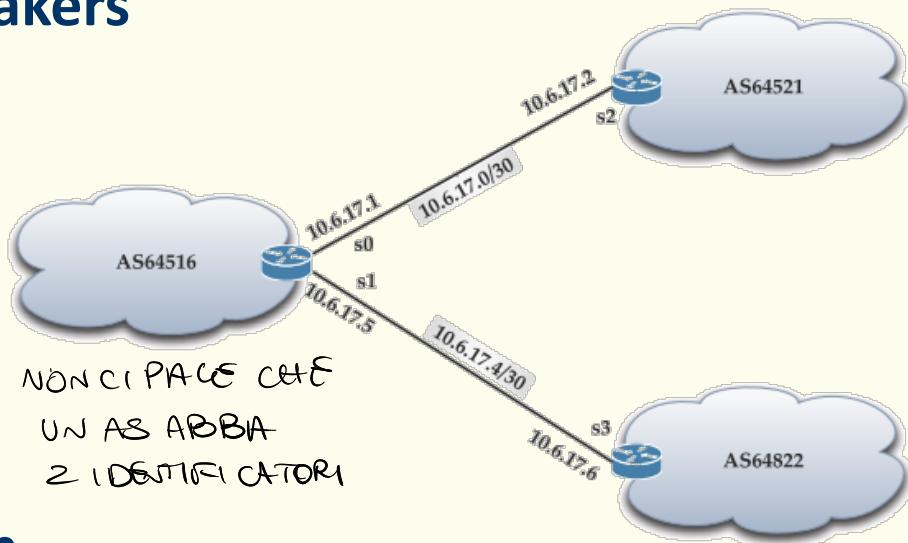
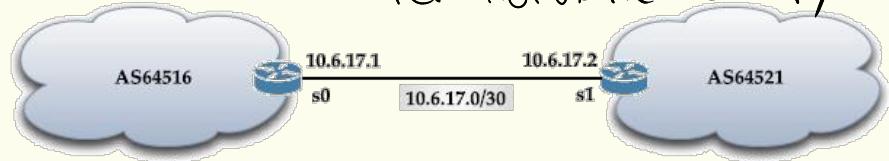
List of networks {address/prefix length}



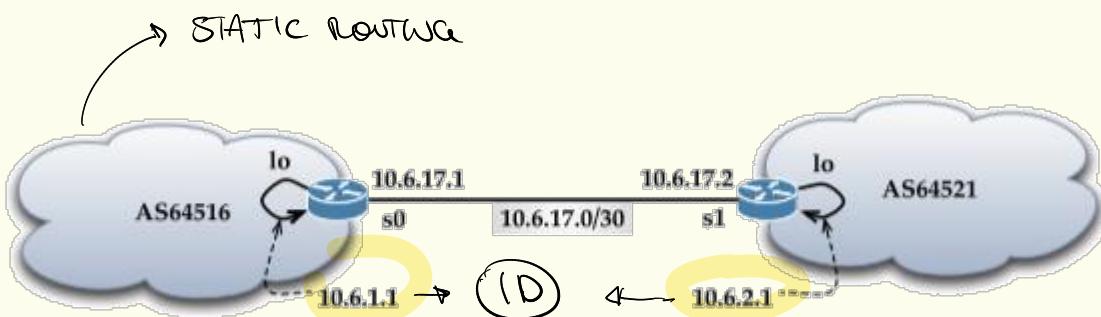
BGP initialization

p2p connections between BGP speakers

PER CONFI CEN PIAZIATE NOU
PER AUTOMATIC DISCOVERY

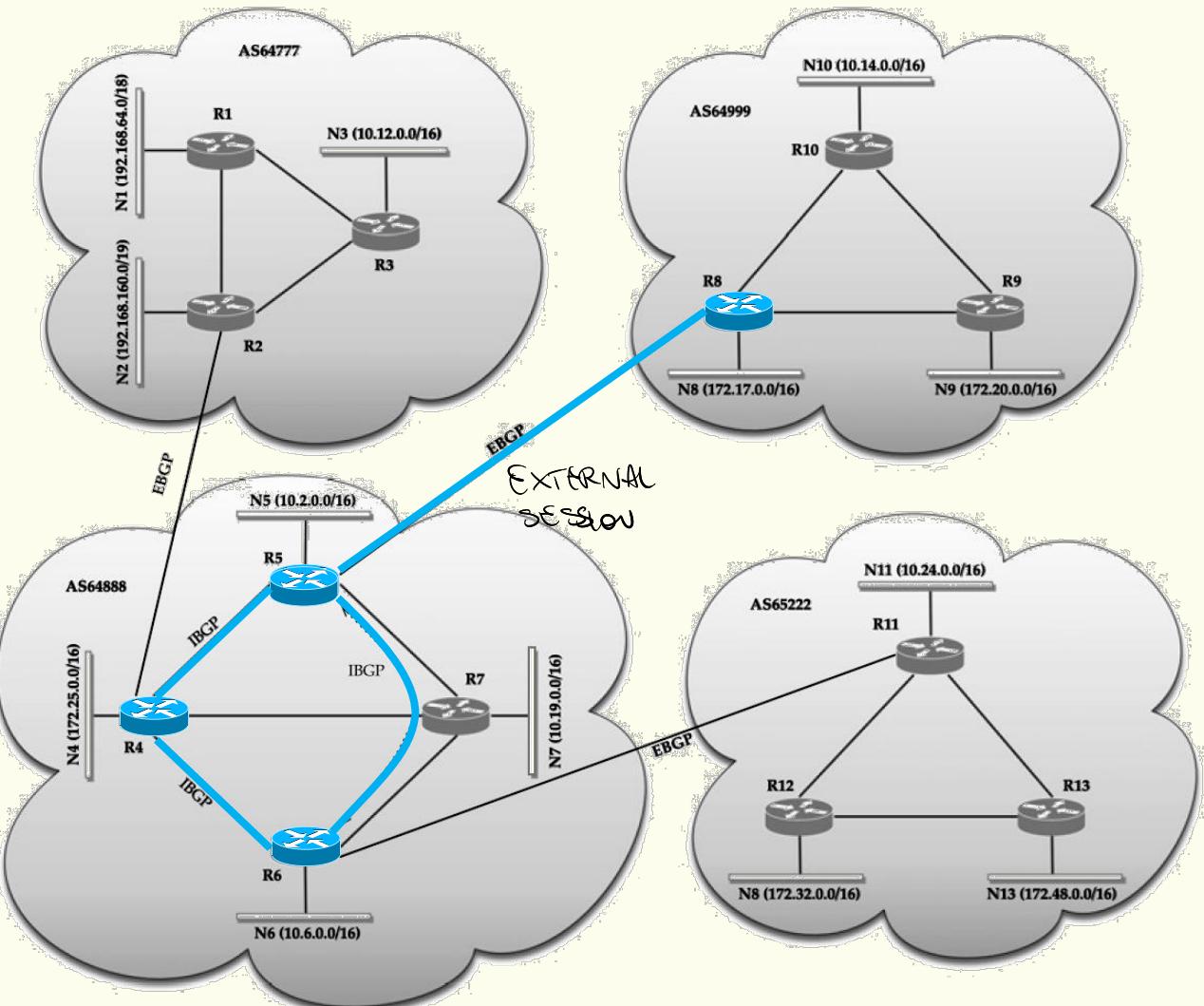


Loopback interface-based approach

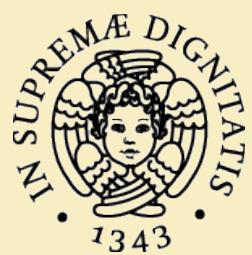


(external) eBGP vs. (internal) iBGP

eBGP: BGP speakers are in different ASes, e.g. R5 and R8



iBGP: BGP speakers are in the same AS, e.g. R5 and R6

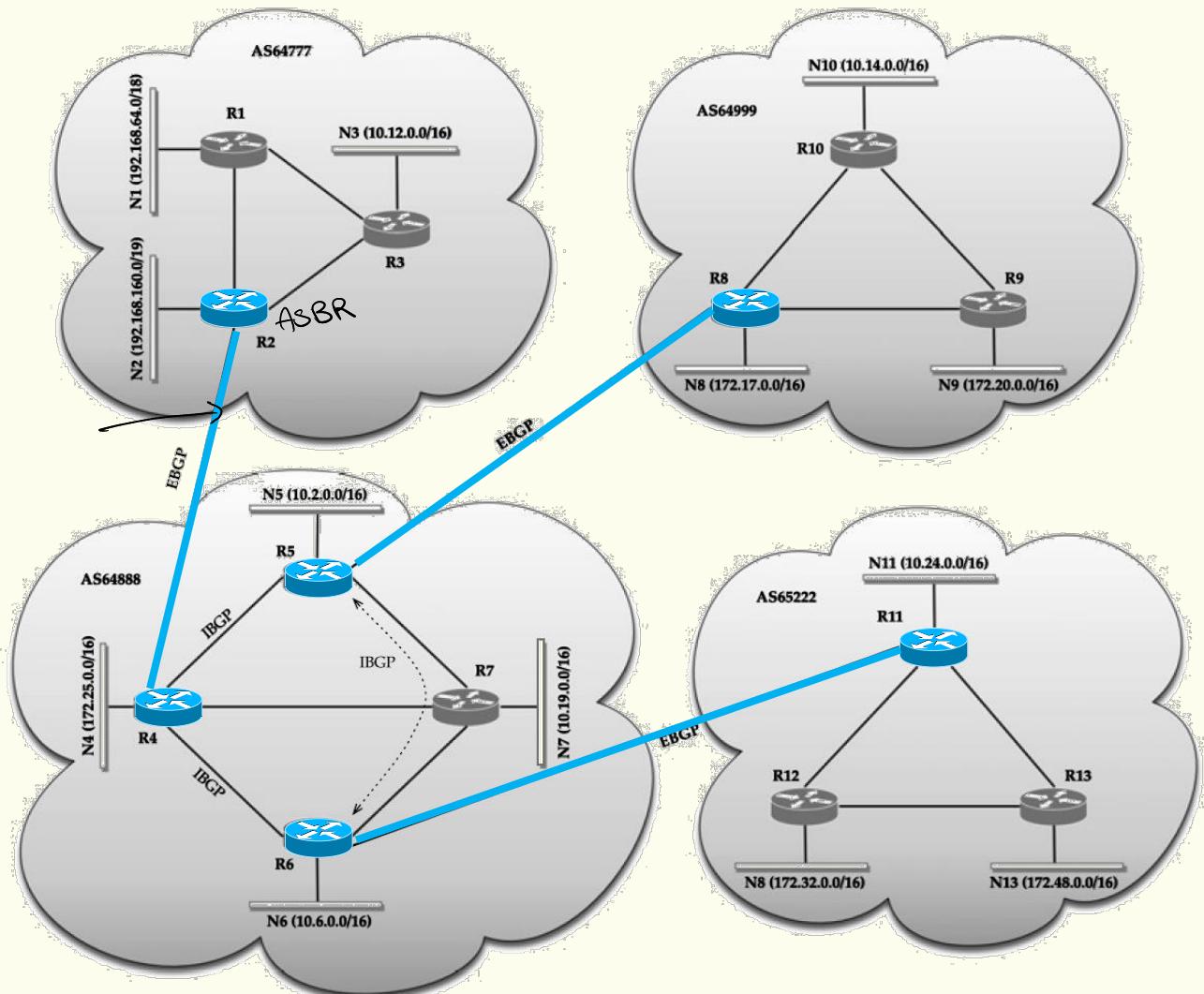


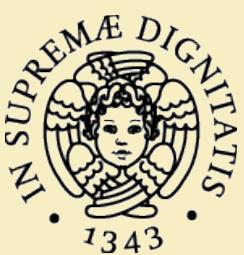
Why iBGP?

Stub AS (e.g. AS64777):
no need for eBGP speakers
to redistribute BGP routes
within the AS

PER SAPERE CERI
ACERI PREFISSI
RACEA IUNGIBILE

Transit AS (e.g. AS64888):
BGP routes learned by one
eBGP speaker need be
redistributed to the other
eBGP speakers





Why iBGP?

PERCHÉ DOVREMO AVERE BGP SPEAKERS
ACI INTERNO DELLO STESSO AS?

CONNESSO AD
UN SOLO ARCO
AS

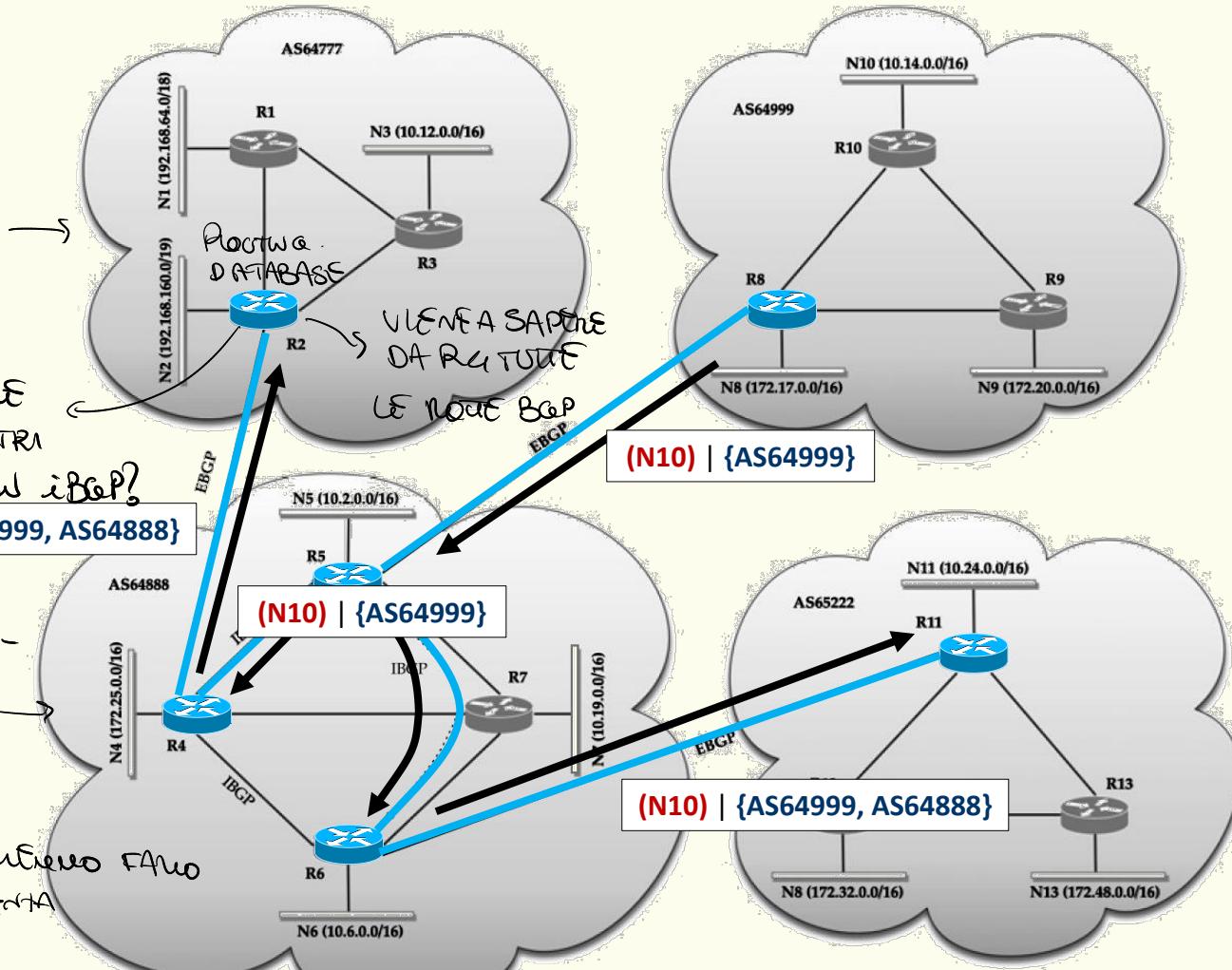
DEVE PROPAGARE
O NO AGLI ALTRI
ROUTER CON iBGP?
(N10) | {AS64999, AS64888}

IN GENERALE NO,
DEFAULT ROUTE

TRANSIT
AS

TRA GLI ALTRI PORTANO FAI
CON OSPF, SENZA
USARE iBGP.

iBGP E' SEGRE PER I TRANSIT AS NON PER GLI STUB AS.



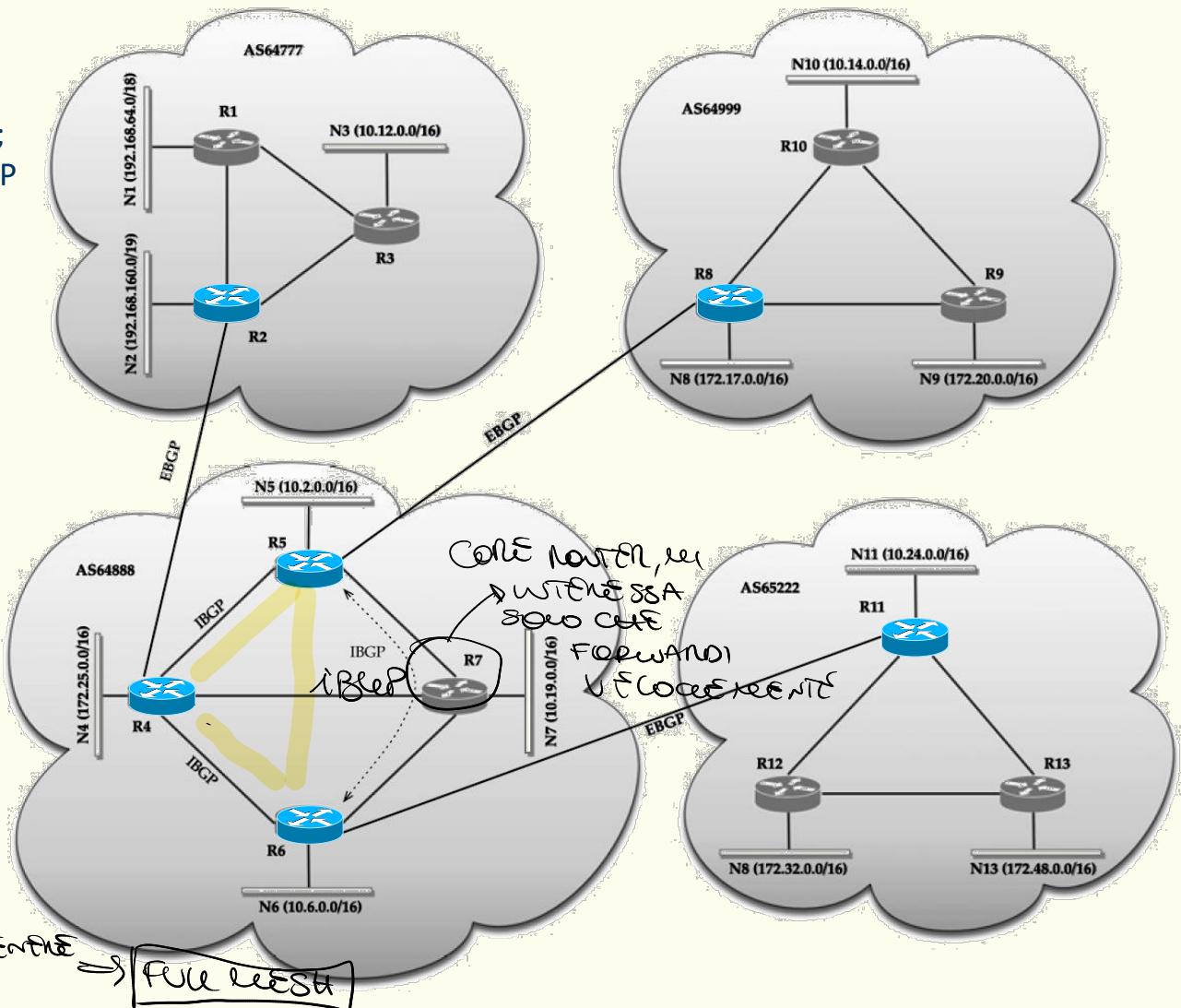
Rules

Rule 1 A BGP speaker can advertise IP prefixes it has learned from an eBGP speaker to a neighboring iBGP speaker; similarly, a BGP speaker can advertise IP prefixes it has learned from an iBGP speaker to an eBGP speaker

Rule 2 An iBGP speaker cannot advertise IP prefixes it has learned from an iBGP speaker to another peer iBGP speaker

Two reasons:

1. Avoid looping of BGP route updates within the AS
2. No need to advertise internal routes





Rules

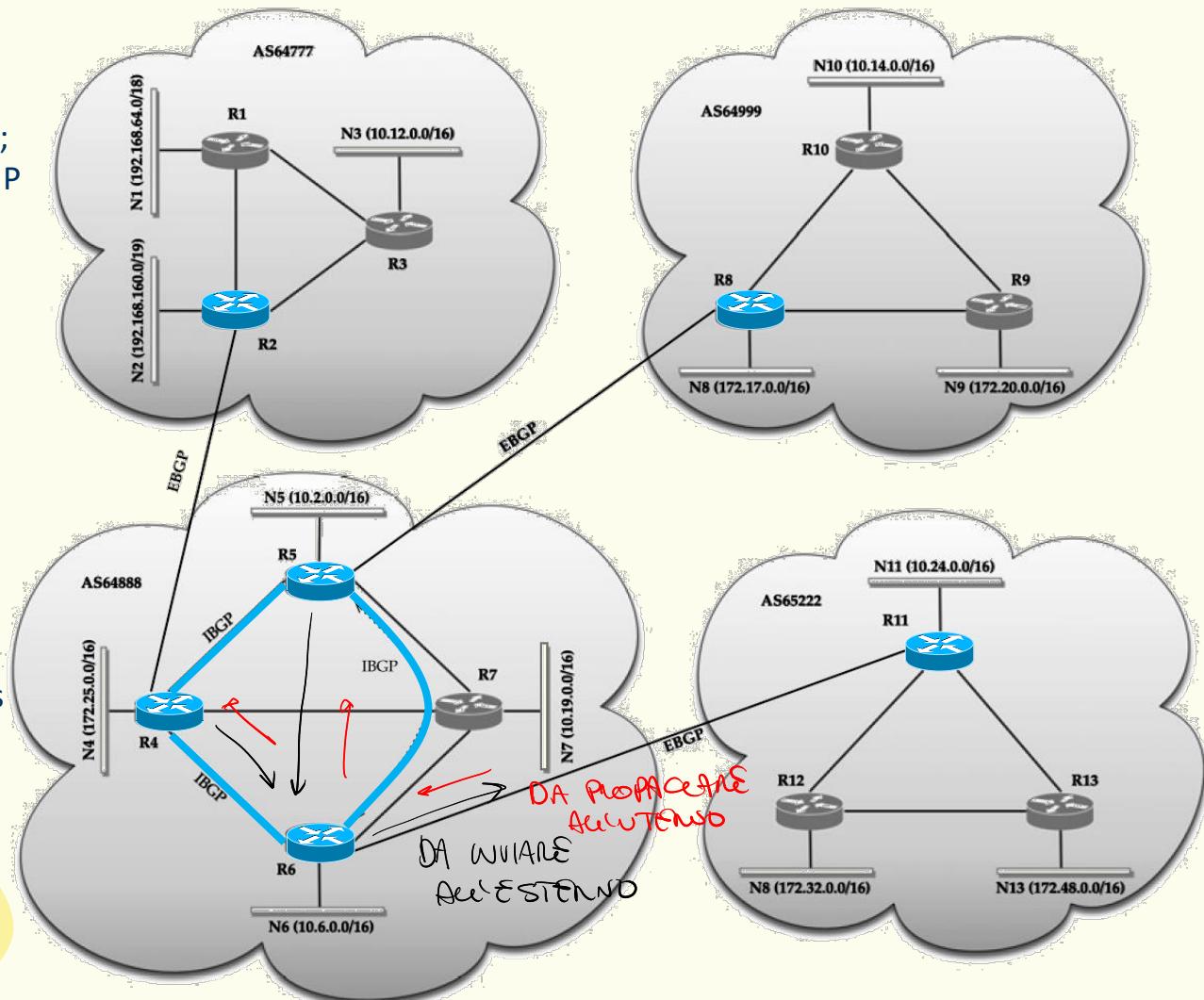
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Two reasons:

1. Avoid looping of BGP route updates within the AS
2. No need to advertise internal routes

A full mesh iBGP connectivity is needed

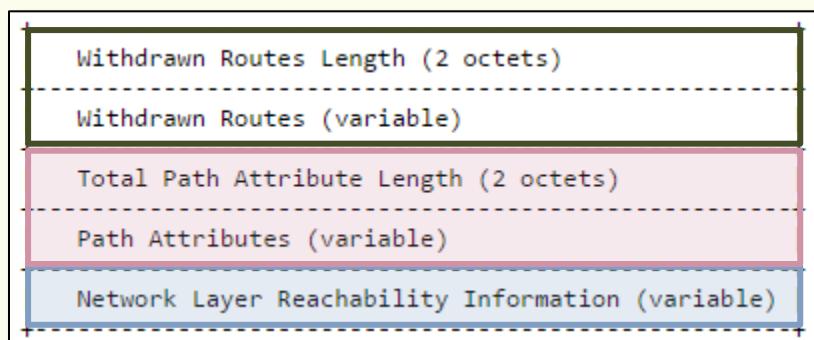
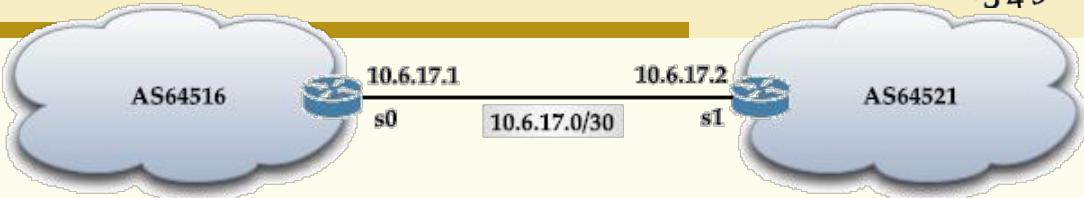




PATH attributes

- **UPDATE message**

- Path attributes advertised for a set of routes
- Keep track of route-specific information such as path information, and degree of preference of a route
- Used in the BGP filtering and route decision process



A sequence of path attributes (TLV)

- ORIGIN
- AS_PATH
- NEXT_HOP
- LOCAL_PREF
- ...

List of networks {address/prefix length}

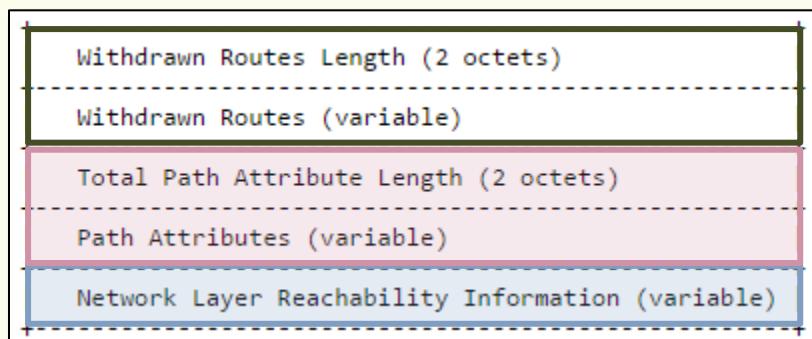


PATH attributes

• UPDATE message

UN BUP PEER DEV' ESSERE W CERTO IN A PROCESSO

- **Well-known mandatory:** Must be present in all UPDATE messages
- **Well-known discretionary:** Could be present in UPDATE messages
- **Optional transitive:** If not recognized, are propagated to other neighbors
- **Optional non-transitive:** Discarded if not recognized



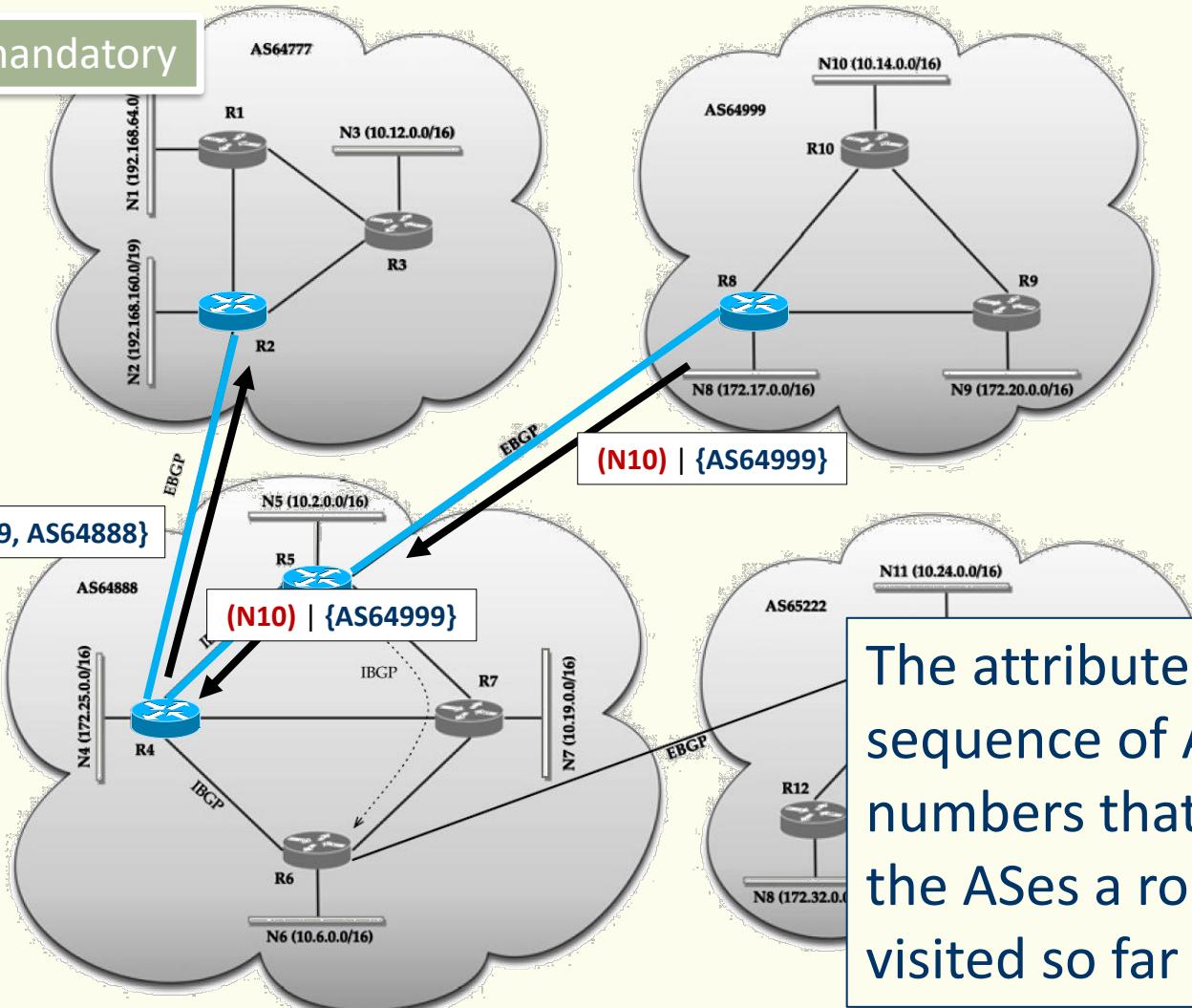
A sequence of path attributes (TLV)

- ORIGIN |
IGP EGP UNKNOWN
- AS_PATH
- NEXT_HOP
- LOCAL_PREF
- ...

List of networks {address/prefix length}

AS_PATH

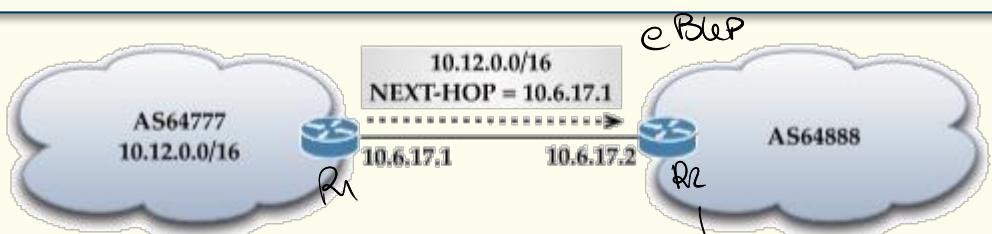
Well-known mandatory



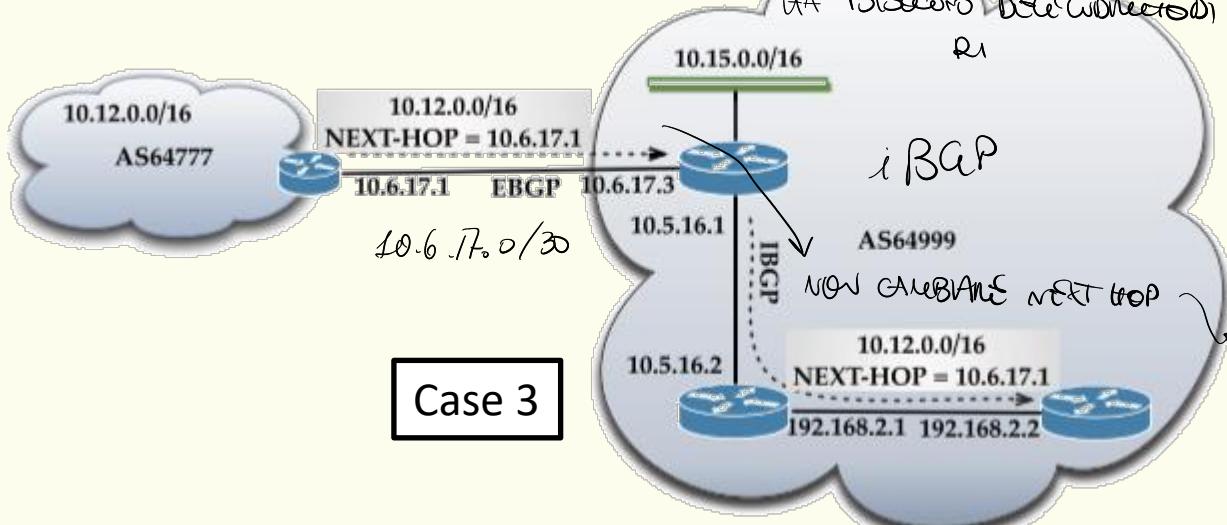
The attribute stores a sequence of AS numbers that identifies the ASes a route has visited so far

NEXT_HOP

The attribute defines the IP address of the router that **SHOULD** be used as the next hop (not necessarily 1-hop) to the destinations listed in the UPDATE message

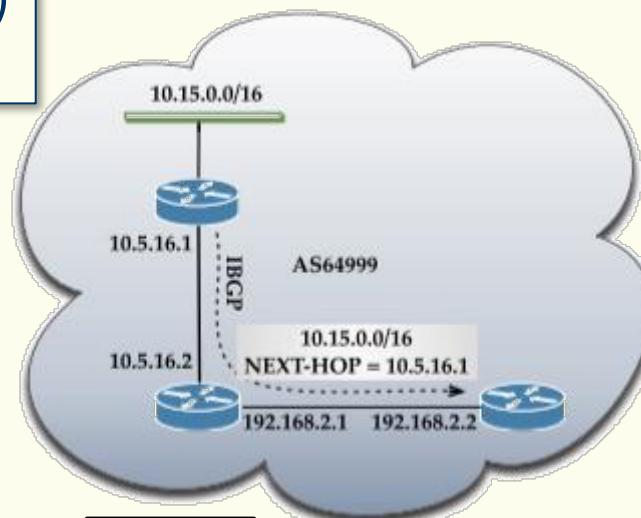


Case 1



Case 3

Well-known mandatory



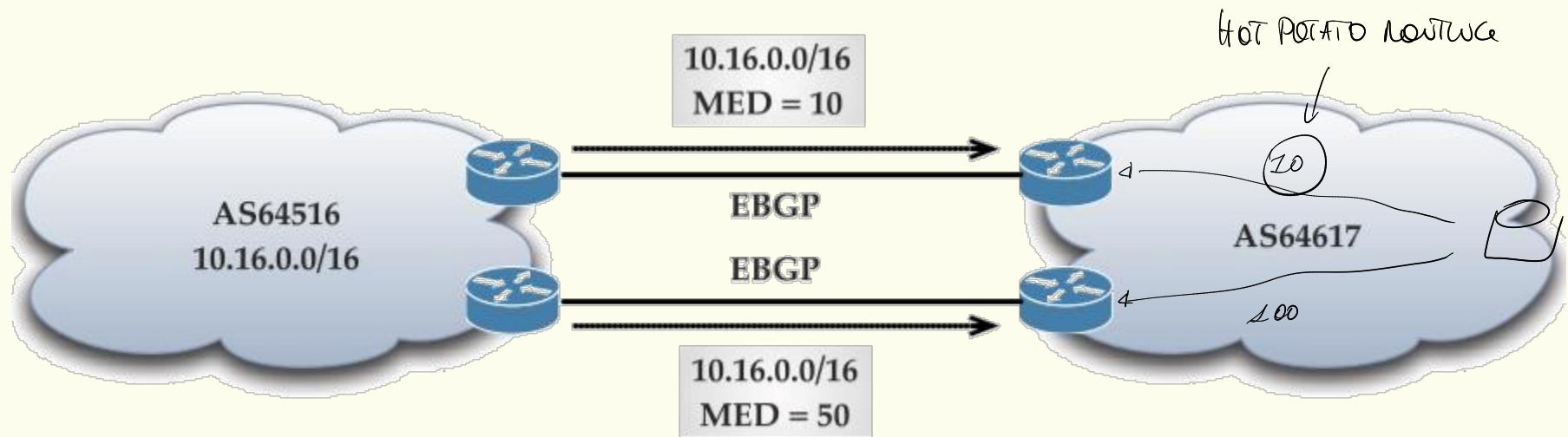
Case 2

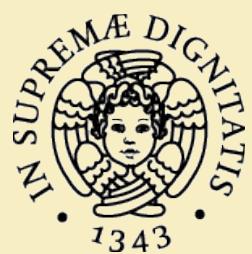
PREM R2 E RA DIRETTAMENTE CONNESSO E STAVET CORE RA CERCA E NELLA RA AFF DO A R2 CHE PASSI L'INTERA RAC

MED (MULTI-EXIT-DISCRIMINATOR)

The attribute is a metric meant for use when there are multiple external links to a neighboring AS

Optional non-transitive





LOCAL_PREF

The attribute defines a metric used internally within an AS between BGP speakers, helpful in selection when the AS has connectivity to multiple ASes

Well-known discretionary

