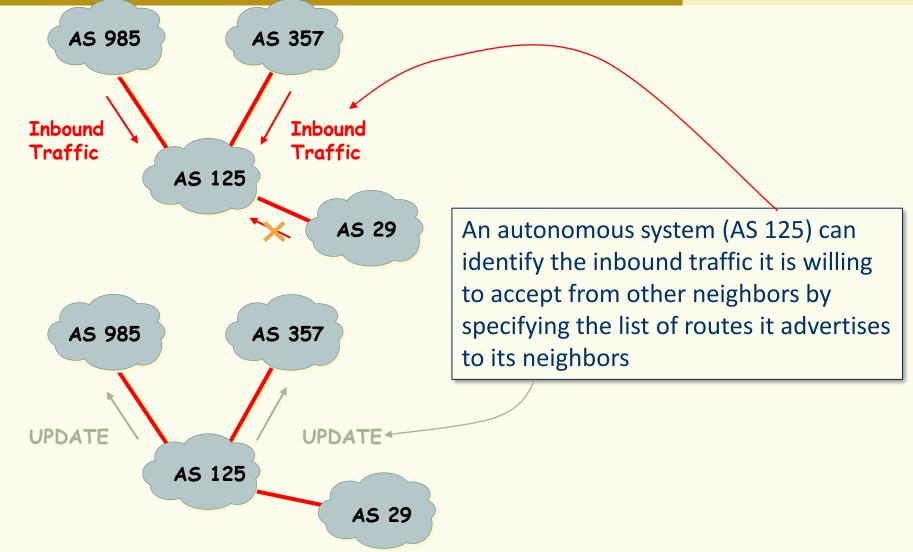
Interdomain routing

BGP-4 decision process

Enzo Mingozzi
Professor @ University of Pisa
enzo.mingozzi@unipi.it

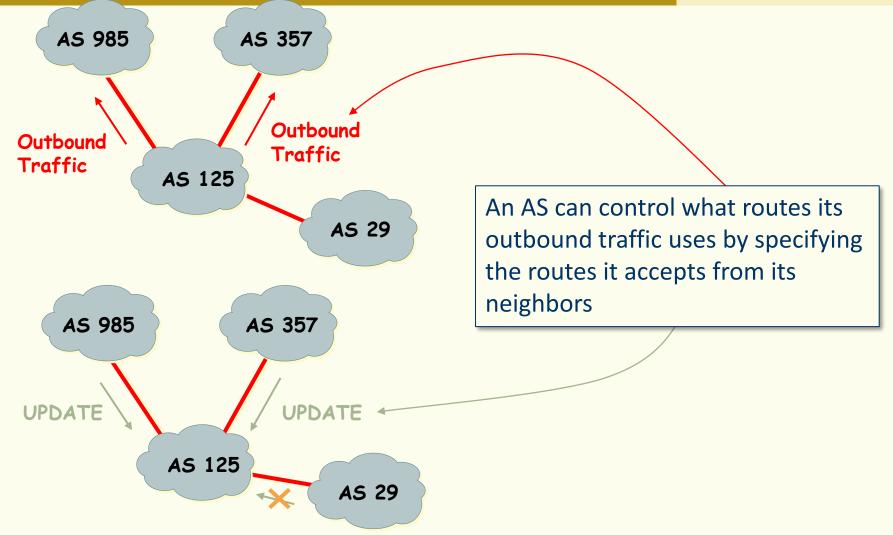
Route Filtering





Route Filtering

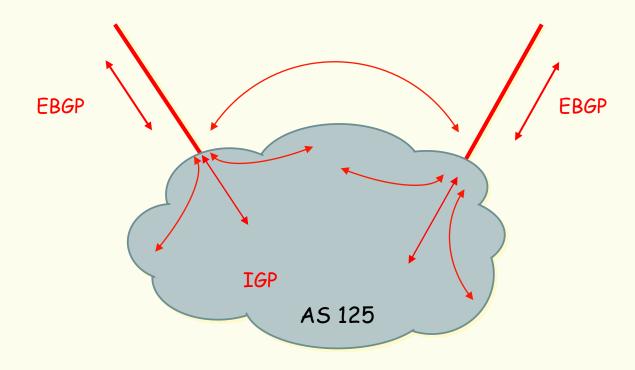




Route Filtering



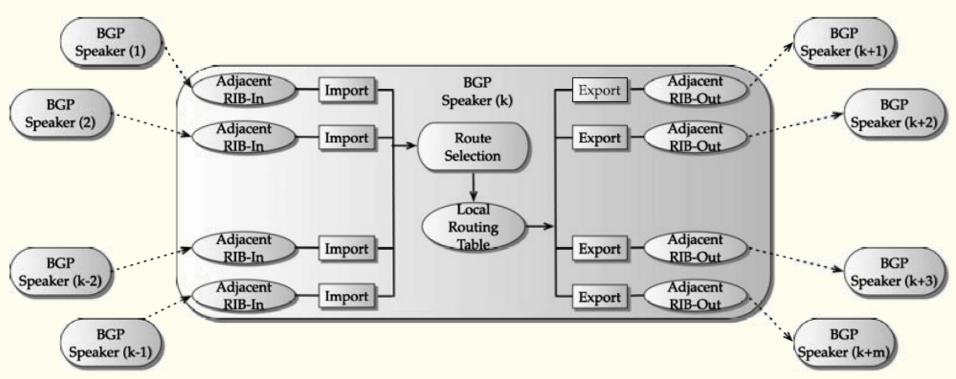
Filtering can also limit routing updates flowing from one protocol to another There is the possibility of injecting BGP routes in the IGP as well as injecting the IGP or static routes into BGP





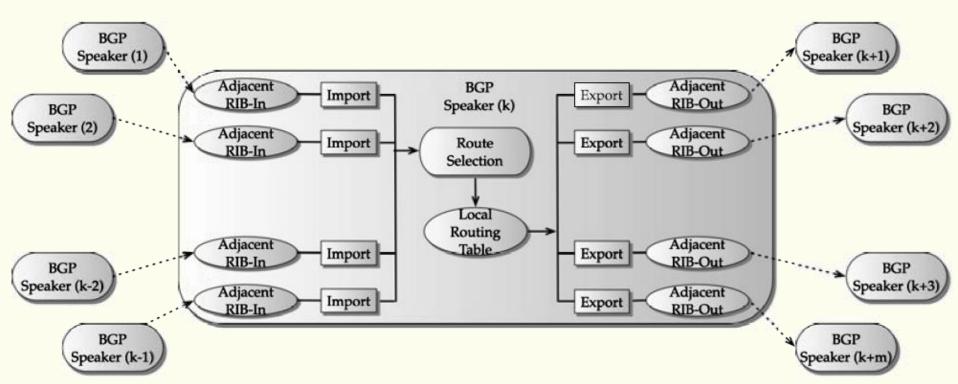
The BGP decision process consists of

- 1) path selection, and
- 2) (aggregation and) dissemination



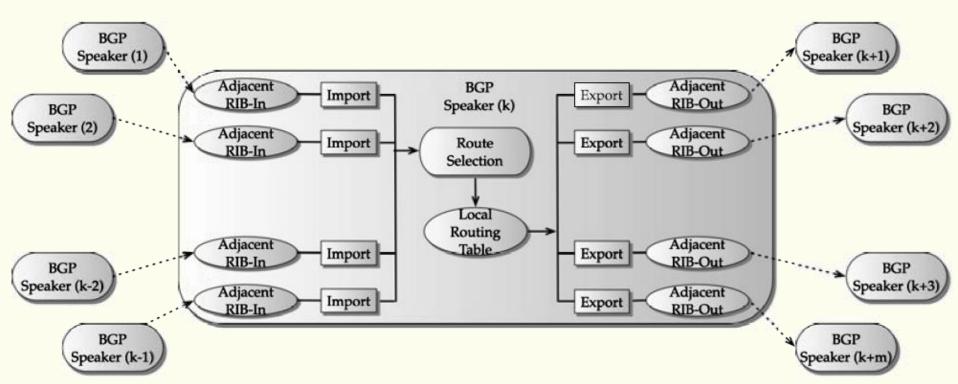


Each BGP speaker maintains several **Routing Information Bases Adjacent RIBs-In (Adj-RIBs-In)** stores AS level routing information for each IP prefix it has learned about from its neighbors through inbound UPDATE messages



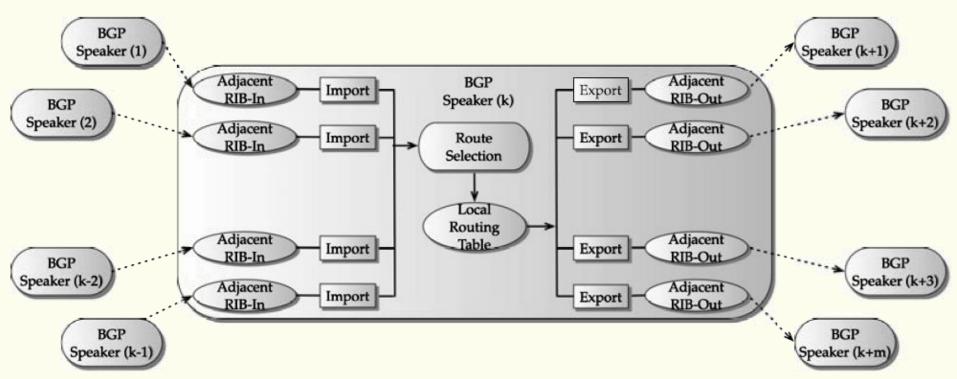


Each BGP speaker maintains several **Routing Information Bases Loc-RIB** stores the routes that have been determined locally by the BGP speaker decision process, used for updating the forwarding table





Each BGP speaker maintains several **Routing Information Bases Adjacent RIBs-Out (Adj-RIBs-Out)** stores the routes for advertisement to its neighboring BGP speakers through outbound UPDATE messages



BGP path selection

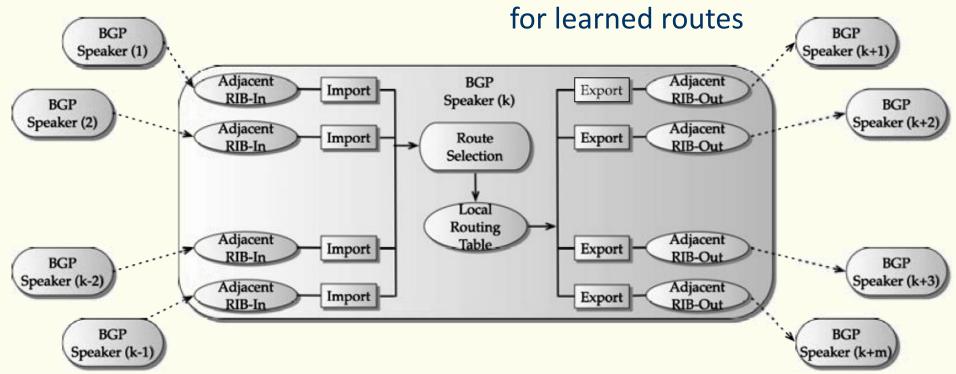


Two phases

- 1) Import policy and filtering
- 2) Best route determination

 Filter out IP prefixes that are not allowed or that should not be reached via that peer

Assess the degree of preference



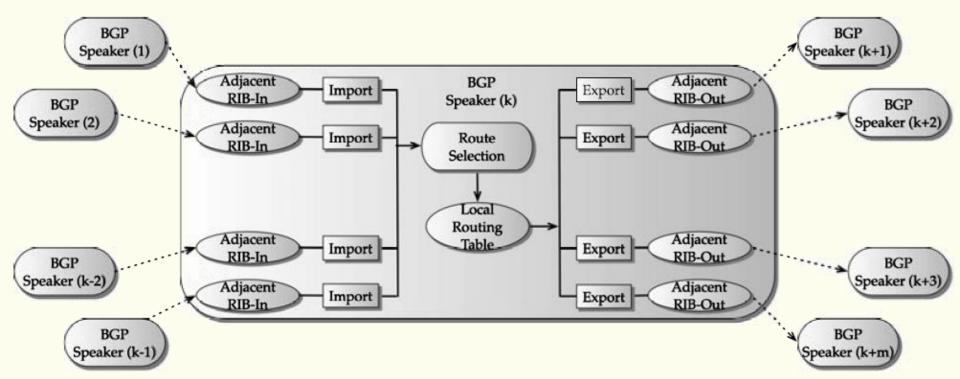
BGP path selection



Two phases

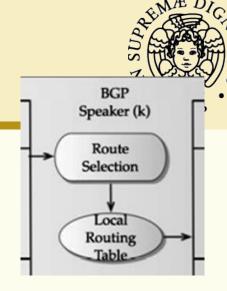
- 1) Import policy and filtering
- 2) Best route determination

 Select the best route for each separate imported IP prefix



BGP path selection

• **Tie-breaking rules** when multiple routes are available to the same **imported** IP prefix

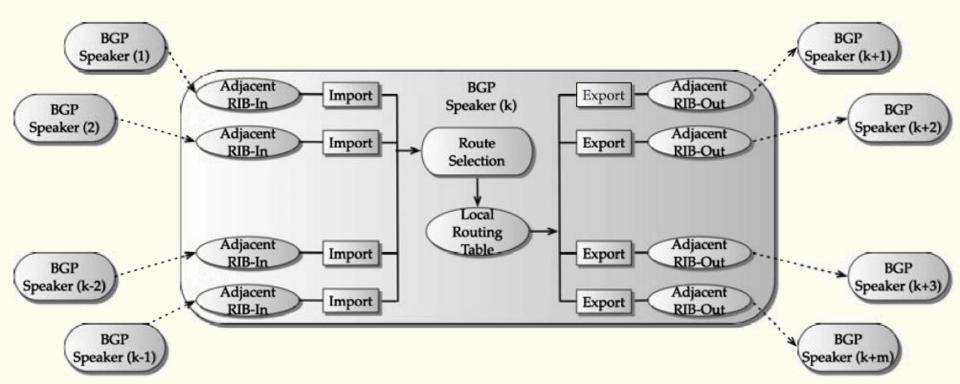


- 1. Ignore routes for which the **NEXT-HOP attribute is not resolvable**
- 2. Apply the **degree of preference** assessed during the *import policy and filtering* phase (either on LOCAL–PREF if the announcement is received from an iBGP speaker, or any locally pre-configured decision)
- 3. Select the route that **originated locally** at the BGP speaker
- 4. Select the route with the **shortest AS path**
- 5. Select the one with the **lowest ORIGIN** attribute (IGP, then EGP, then Incomplete)
- 6. Select the route with the **lowest MED** for eBGP routes (learned from the same AS)
- 7. Select the route received from **eBGP** over iBGP
- 8. Select the route with **shortest (internal) path to the NEXT–HOP router** (as determined by IGP)
- 9. Select the route learned from the eBGP neighbor with the lowest BGP identifier
- 10. Select the route from the iBGP neighbor with the lowest BGP identifier

BGP route aggregation and dissemination



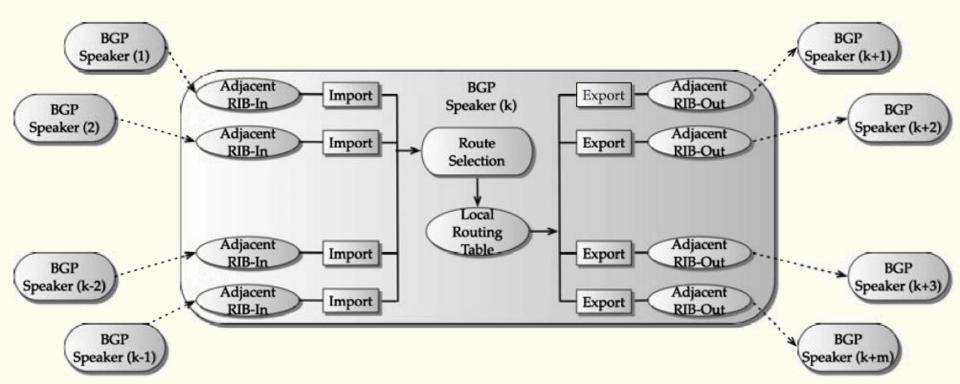
 Optional route aggregation based on CIDR: combine IP prefixes (supernetting) to reduce the number of networks announced to a downstream AS



BGP route aggregation and dissemination

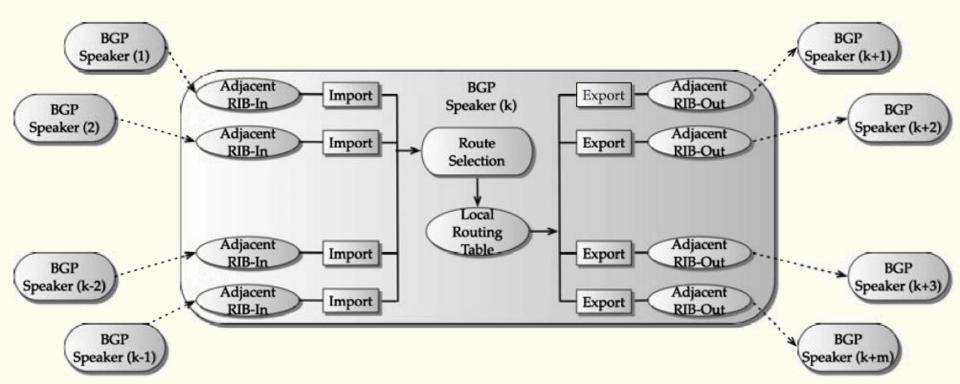


- A BGP speaker applies an export policy before propagating routes to other BGP speakers
- Export policies are separate per neighboring BGP speaker



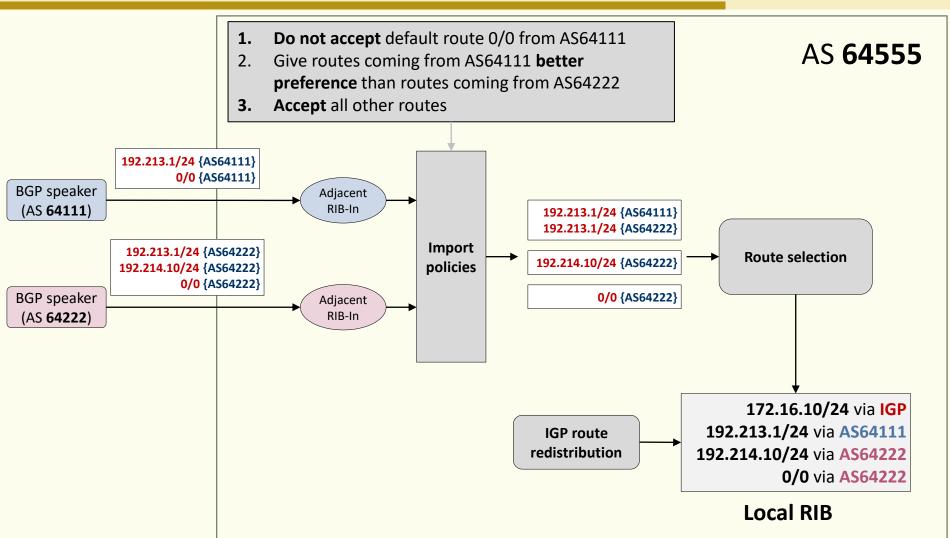


 Policy-based routing: import and export policies are placed at a BGP speaker by a network administrator due to business relations or peering arrangement, i.e., external factors



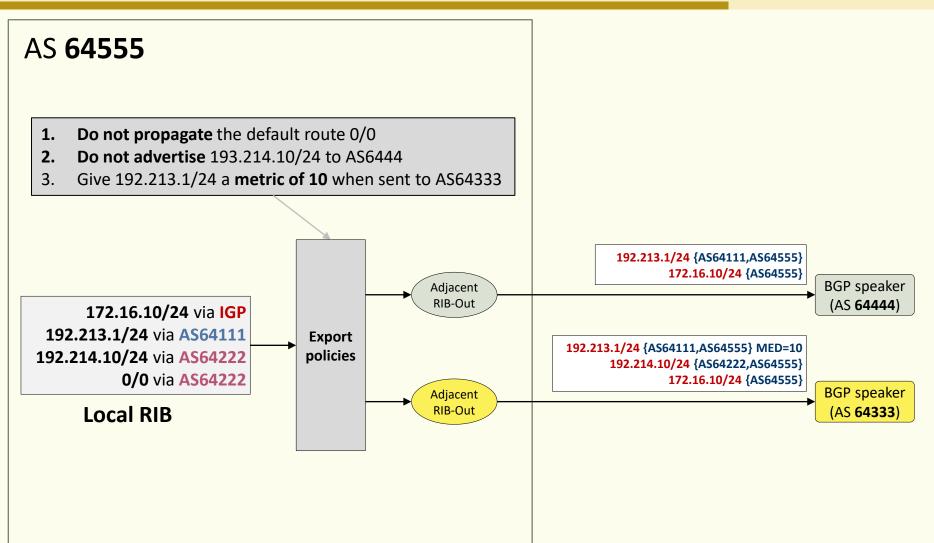
Example





Example





Internal BGP scalability



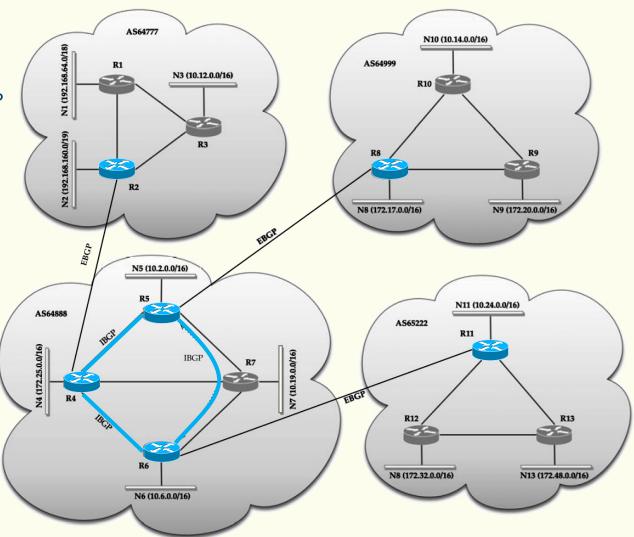
Rule 1 A BGP speaker can advertise IP prefixes it has learned from an eBGP speaker to a neighboring iBGP speaker; similarly, a BGP speaker can advertise IP prefixes it has learned from an iBGP speaker to an eBGP speaker

Rule 2 An iBGP speaker cannot advertise IP prefixes it has learned from an iBGP speaker to another peer iBGP speaker

Two reasons:

- 1. Avoid looping of BGP route updates within the AS
- 2. No need to advertise internal routes





Internal BGP scalability



Rule 1 A BGP speaker can advertise IP prefixes it has learned from an eBGP speaker to a neighboring iBGP speaker; similarly, a BGP speaker can advertise IP prefixes it has learned from an iBGP speaker to an eBGP speaker

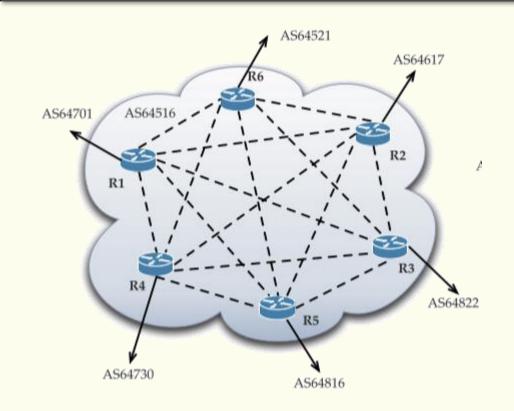
Rule 2 An iBGP speaker cannot advertise IP prefixes it has learned from an iBGP speaker to another peer iBGP speaker

Two reasons:

- 1. Avoid looping of BGP route updates within the AS
- 2. No need to advertise internal routes



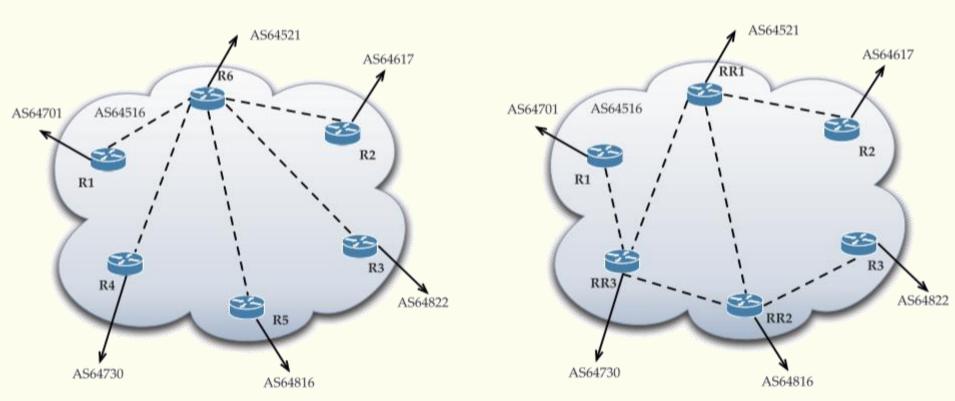
n iBGP speakers $\rightarrow n(n-1)/2$ iBGP sessions each speaker handling n-1 sessions



Route reflector



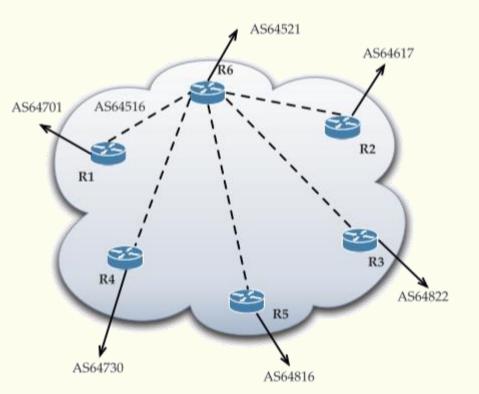
- One or more iBGP speakers act as concentration routers
- The other iBGP speakers establish only one BGP session to a route reflector (route reflector clients)
- Each route reflector with its clients form a cluster, identified by a CLUSTER-ID

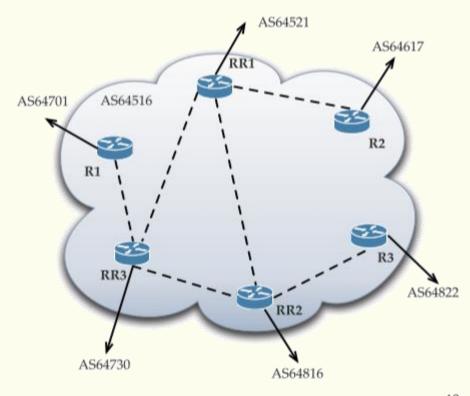


Route reflector



- Announcement received **from another route reflector** → reflect/pass it to its clients
- Announcement received from a route reflector client → reflect to another route reflector
- Announcement received from an eBGP speaker → reflect to all other route reflectors and clients





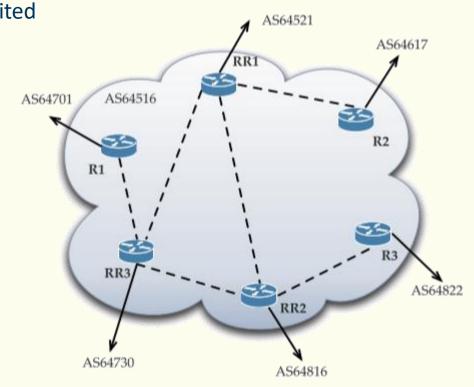
Route reflector



- Route reflectors must form a full mesh connectivity among themselves!
- How to avoid routing loops? Two additional attributes
- ORIGINATOR-ID: identifies a route reflector through its 4-byte router ID, added only by the originating route reflector

2. CLUSTER-LIST: stores a sequence of 4-byte CLUSTER-ID values to indicate the path of

clusters that an advertised IP prefix has visited AS64617 AS64701 AS645 R1 AS64822 AS64730 AS64816



References



- D. Medhi, K. Ramasamy, Network Routing: Algorithms, Protocols, and Architectures, 2nd/ed. Morgan Kaufmann, ©2018
- RFC
 - RFC4271, A Border Gateway Protocol 4 (BGP-4),
 Jan. 2006
 - RFC4360, BGP Extended Communities Attribute,
 Feb. 2006