Congruence Closure Algorithm

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If you haven't read the README.md file, I suggest you do so before reading this report.

1 Introduction

This paper presents the research conducted within the scope of the Automatic Reasoning course for the academic year 2022/2023. Specifically, the focus lies on the implementation of the Congruence Closure Algorithm utilizing Directed Acyclic Graphs (DAGs).

The content of the project can be found at the following link: Leonardo Zecchin's project.

2 Implementation

2.1 Project structure

- 1. The mainProgram.py program implements the algorithm by reading an input file that contains formulas in normal form, e.g., f(a,b)=a and f(f(a,b),b)!=a.
- 2. The the Parser.py program implements the algorithm by reading an input file that contains the formulas that need to be brought into DNF form and then brought to normal form, e.g., and (eq(f(a,b),a),dis(f(f(a,b),b),a)) becomes f(a,b)=a and f(f(a,b),b)!=a; or imply(eq(x,g(y,z)),eq(f(x),f(g(y,z)))) becomes and (eq(x,g(y,z)),dis(f(x),f(g(y,z)))) and after x=g(y,z) and f(x)!=f(g(y,z)).
- 3. The mainProgramWithFL.py program implements the algorithm by reading an input file that contains formulas in normal form, e.g., f(a,b)=a and f(f(a,b),b)!=a but it use the Forbidden List.

Within the code folder are the codes that are used by the main programs, in particular cca is used for the Congruence Closure Algorithm, while the other programs were used during implementation.

Inside the classes folder, you will find two important classes: dag and node, which are utilized by the algorithm.

In the input and output folders, there are two types of files:

- 1. input.txt and output.txt: the former contains the formulas in the normal form, and in the latter, you will find the algorithm's resulting outcomes.
- 2. inputToParser.txt and outputToParser.txt: the former contains formulas that must be parsed, and in the latter, you will find the algorithm's resulting outcomes.

2.2 The Algorithm

The structure of the algorithm is the following: Given $\Sigma_{\rm E}$ -formula

$$F: s_1 = t_1 \wedge \cdots \wedge s_m = t_m \wedge s_{m+1} \neq t_{m+1} \wedge \cdots \wedge s_n \neq t_n$$

with subterm set S_F , perform the following steps:

- 1. Construct the initial DAG for the subterm set S_F .
- 2. For $i \in \{1, ..., m\}$, MERGE $s_i t_i$.
- 3. If FIND $s_i = \text{FIND } t_i$ for some $i \in \{m+1, \ldots, n\}$, return unsatisfiable.
- 4. Otherwise (if FIND $s_i \neq$ FIND t_i for all $i \in \{m+1, \ldots, n\}$) return satisfiable.

Inside the dag.py you will find the implementation of the algorithm's functions: *MERGE*, *UNION*, *CONGRUENT*, *CCPAR*, *FIND* and *NODE*, the specific explanation of functions is outside the scope of the paper.

The implementation of the algorithm is within the **cca.py** in particular the function is the **congruenceClosureAlgorith** which takes in input:

- 1. F₋plus: it contains the formulas with equality (=);
- 2. F_minus: it contains the formulas with disequality (\neq) ;
- 3. Sf: the subterm set;
- 4. new_dag: is the DAG that represents the subterm set.

```
\begin{array}{lll} def & congruenceClosureAlgorithm (F\_plus\,,\;\; F\_minus\,, Sf\,, new\_dag\,)\colon\\ & for\;\; f\;\; in\;\; F\_plus\,\colon\\ & \#Step\;\; 1\\ & idx1\,, idx2\,=\,\; getIndex\,(f\,,Sf)\\ & new\_dag\,.MERGE(idx1\,,idx2\,)\\ \#Step\;\; 2\\ & for\;\; f\;\; in\;\; F\_minus\,\colon\\ & idx1\,, idx2\,=\,\; getIndex\,(f\,,Sf)\\ & if\;\; new\_dag\,.FIND(idx1\,)\,\Longrightarrow\,\; new\_dag\,.FIND(idx2\,)\colon\\ & \quad return\;\; False\\ & else\,\colon\\ & \quad return\;\; True \end{array}
```

In the congruenceClosureAlgorithm there are the implementation the second, the third and fourth steps.

2.3 Differences with pseudocode

The only difference between my code and the pseudocode is in the UNION function:

```
def UNION(self, id1: int, id2: int) -> None:
print(f"UNION {id1} {id2}")
n1 = self.NODE(id1)
n2 = self.NODE(id2)
n1_ccpar = self.CCPAR(id1)
n2_ccpar = self.CCPAR(id2)
if self.FIND(n1.find) != n1.id:
    self.NODE(self.FIND(n1.find)).find = n2.find
else:
    n1.find = n2.find
n2.ccpar = n2_ccpar+ n1_ccpar
n1.ccpar = []
```

I introduced the if/else statement to address scenarios in which the algorithm needs to modify the find field of the n1 node, but it is connected to another node. In such cases, I must also modify the find field of the node linked to n1. Additionally, a notable distinction is that UNION stores the respective CCPAR values in n1_ccpar and n2_ccpar prior to altering these fields.

3 Forbidden List

I tried to implement the forbidden list. There are some files where I implement the forbidden list:

- 1. cca_fl.py in the congruenceClosureAlgorithm and createDAG functions;
- 2. dag_FL.py in the UNION function;
- 3. node_Fl.py in the MERGE function.

The idea was to create a forbidden list for every node in F_plus, which is the list of the equality formulas. The program checks if a formula in F_plus is in F_minus. This means that if a formula is both equality and disequality formulas the algorithm return UNSAT.

During the CCA before call the MERGE s t recursively the program check if s is in the forbidden list of t or vice versa and if it is true the program return False.

I decide to implement the forbidden list in this way because I think that it is the most efficient way to implement it.

4 Results

In this section we show the results obtain with the algorithm.

4.1 Table Without Forbidden List with Parser

Experiments				
Formulas	DNF	Satisfiability	Time execution	
$-\mathrm{imply}(\mathrm{eq}(\mathrm{x},\mathrm{g}(\mathrm{y},\mathrm{z})),\mathrm{eq}(\mathrm{f}(\mathrm{x}),\mathrm{f}(\mathrm{g}(\mathrm{y},\mathrm{z}))))$	x=g(y,z) and $f(x)!=f(g(y,z))$	UNSAT	0.0004217	
$-\operatorname{and}(\operatorname{eq}(f(a,b),a),\operatorname{dis}(f(f(a,b),b),a))$	f(a,b)=aandf(f(a,b),b)!=a	UNSAT	0.000449	
-and(eq(f(f(f(a))),a),and	f(f(f(a)))=aand	UNSAT	0.0006108	
(eq(f(f(f(f(f(a))))),a),dis(f(a),a)))	f(f(f(f(f(a)))))=a and f(a)!=a			
$-\operatorname{and}(\operatorname{eq}(f(f(f(a))),f(a)),$	f(f(f(a)))=f(a) and $f(f(a))=a$	SAT	0.000388	
and(eq(f(f(a)),a),dis(f(a),a)))	andf(a)!=a			
$-\operatorname{and}(\operatorname{eq}(x,\operatorname{g}(x,z)),\operatorname{dis}(\operatorname{f}(x),\operatorname{f}(\operatorname{g}(y,z))))$	x=g(x,z) and $f(x)!=f(g(y,z))$	SAT	0.00032	

4.2 Table Without Forbidden List

Experiments				
Formulas	Satisfiability	Time execution		
f(f(f(a)))=a and $f(f(f(f(f(a)))))=a$ and $f(a)!=a$	UNSAT	0.001177		
-f(a)=b and $f(a)!=b$	UNSAT	0.000168		
-x=y and $f(x)!=f(y)x=y$ and $f(x)!=f(y)$	UNSAT	0.000255		
-f(x)=f(y) and $x!=y$	SAT	0.00052022		
-f(g(a))=g(f(a)) and $f(g(f(b)))=a$ and $f(b)=a$ and $g(f(a))!=a$	SAT	0.00199		
-f(a,b)=a and $f(f(a,b),b)!=a$	SAT	0.0003209		

4.3 Table With Forbidden List

Experiments		
Formulas	Satisfiability	Time execution
-f(f(f(a)))=a and $f(f(f(f(f(a)))))=a$ and $f(a)!=a$	UNSAT	0.00130
-f(a)=b and $f(a)!=b$	UNSAT	0.000340
-x=y and $f(x)!=f(y)x=y$ and $f(x)!=f(y)$	UNSAT	0.00518
-f(x)=f(y) and $x!=y$	SAT	0.0004017
-f(g(a))=g(f(a)) and $f(g(f(b)))=a$ and $f(b)=a$ and $g(f(a))!=a$	SAT	0.001194
-f(a,b)=a and $f(f(a,b),b)!=a$	SAT	0.0007009

You can find other results in the output folder.

5 Conclusion

The algorithm witout or with Forbidden List with my example are similar so I can't say if the forbidden list version is better or not. It was a very interesting project because I learned a lot of things about the Congruence Closure Algorithm and I spend a lot of time for this project but I think that it was worth it.