

IO Controller

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- **Background Knowledge**
- **What is IO Controller**
- **How to use IO Controller**
- **The Implementation of IO Controller**
- **Problems with IO Controller**
- **Development Status**
- **Other Proposals for IO Bandwidth Controlling**

- CGroup(Control Group)
 - A mechanism to put tasks into groups
 - Implemented as a pseudo filesystem
 - Automatically inherited at fork() by any children
 - Resource control is implemented as “subsystem”

How to use CGroup

- Mount a CGroup subsystem

`#mount -t cgroup -o subsystem name /cgroup`

- Create a group

`#mkdir /cgroup/group1`

- Attach a task

`#echo <PID> > /cgroup/group1/tasks`

- Destroy a group

`#rmdir /cgroup/group1`

■ The goals of IO Scheduler

- Reduce seek and improve global throughput

■ The duty of IO Scheduler

- Merge IO requests
- Sort IO requests
- Dispatch requests in different policies

■ CFQ(Complete Fair Queuing)

- One cfq queue for each task(io context)
- Supports io scheduling priorities and classes
- Round-robin scheduling

■ Deadline

- Minimize read starvation

■ Anticipatory

- Anticipating the next read request

■ Noop

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What is IO Controller



- It's a CGroup subsystem just like memcg cpuset etc.
- It provides proportional bandwidth control
- Control the amount of disk time in terms of group weight
- The use cases
 - Provide io BW isolation among Virtual machines
 - Balance io BW among users
 - Limit io BW for a certain application
 - Reduce latency for concurrently running tasks

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How to use IO Controller



- 1 Apply IO Controller patch and compile kernel

Enable the block layer

--> IO Schedulers

- 2 Mount IO Controller on a mountpoint

- 3 Create a CGroup directory

- 4 Setup IO Controller by using CGroup interface

```
[root@localhost cgroup]# ls
```

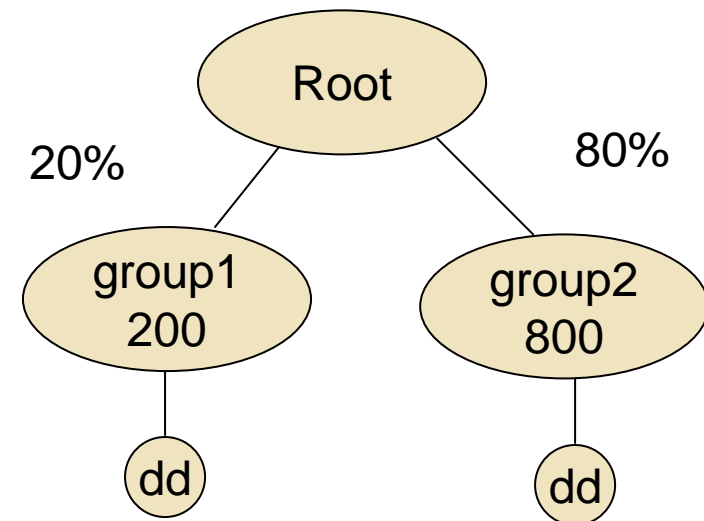
```
blkio.id      io.disk_queue io.disk_time  io.policy  notify_on_release
```

```
io.disk_dequeue io.disk_sectors io.ioprio_class io.weight  tasks
```

- 5 Run tasks in that CGroup

Simple use

- `mkdir /cgroup`
- `mount -t cgroup -o io,blkio io /cgroup`
- `mkdir /cgroup/group1`
- `echo 200 > /cgroup/group1/io.weight`
- `mkdir /cgroup/group2`
- `echo 800 > /cgroup/group2/io.weight`
- `dd if=/sdb2/2000M.1 of=/dev/null &`
- `echo $! > /cgroup/group1/tasks`
- `dd if=/sdb2/2000M.2 of=/dev/null &`
- `echo $! > /cgroup/group2/tasks`



group2: 2097152000 bytes (2.1 GB) copied, 36.3494 seconds, 57.7 MB/s

group1: 2097152000 bytes (2.1 GB) copied, 57.1376 seconds, 36.7 MB/s

sdb disk time usage(**io.disk_time**) at 10th second:

group2: 8:16 **14998** group1: 8:16 **3825**

■ **io.ioprio_class**

- Denotes class of the cgroup (RT, BE, IDLE). This's the default io class of the group for all device. (1 = RT; 2 = BE, 3 = IDLE)

■ **io.weight**

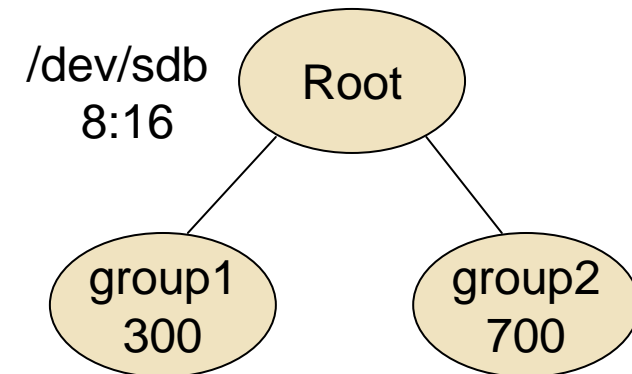
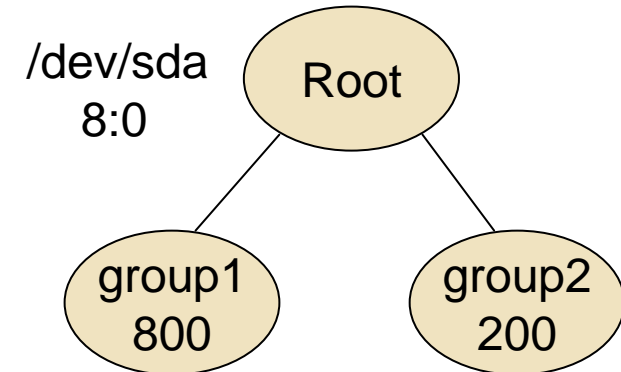
- Denotes per cgroup weight(1 ~ 1000). This's the default weight of the group for all device.

■ **io.policy**

- One can specify per cgroup per device rules using this interface. Syntax:
echo "dev_maj:dev_minor weight ioprio_class" > /path/to/cgroup/io.policy

io.policy example

- **mkdir /cgroup**
- **mount -t cgroup -o io,blkio io /cgroup**
- **mkdir /cgroup/group1**
- **echo 8:0 800 2 > /cgroup/group1/io.policy**
- **echo 8:16 300 2 > /cgroup/group1/io.policy**
- **mkdir /cgroup/group2**
- **echo 8:0 200 2 > /cgroup/group2/io.policy**
- **echo 8:16 700 2 > /cgroup/group2/io.policy**



■ **io.disk_time (read only)**

- disk time allocated to cgroup per device in milliseconds.

```
[root@localhost ~]# cat /cgroup/group1/io.disk_time  
8:16 1930  
8:0 38
```

■ **io.disk_sectors (read only)**

- number of sectors transferred to/from disk by the group.

```
[root@localhost tmp2]# cat /cgroup/group1/io.disk_sectors  
8:16 442520  
8:0 32
```

■ **fairness (0 or 1)**

- Provide better disk time accounting for async IO
- Provide isolation between reads and writes

■ **nr_group_requests**

- maximum number of requests per group

■ **group_idle**

- Prevent from losing the group share

Agenda

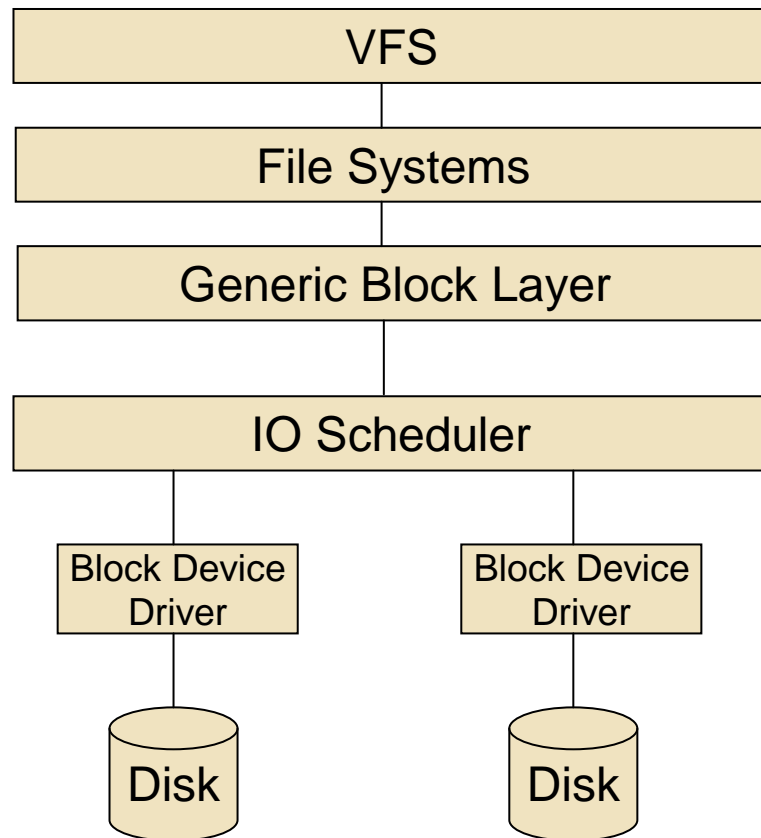


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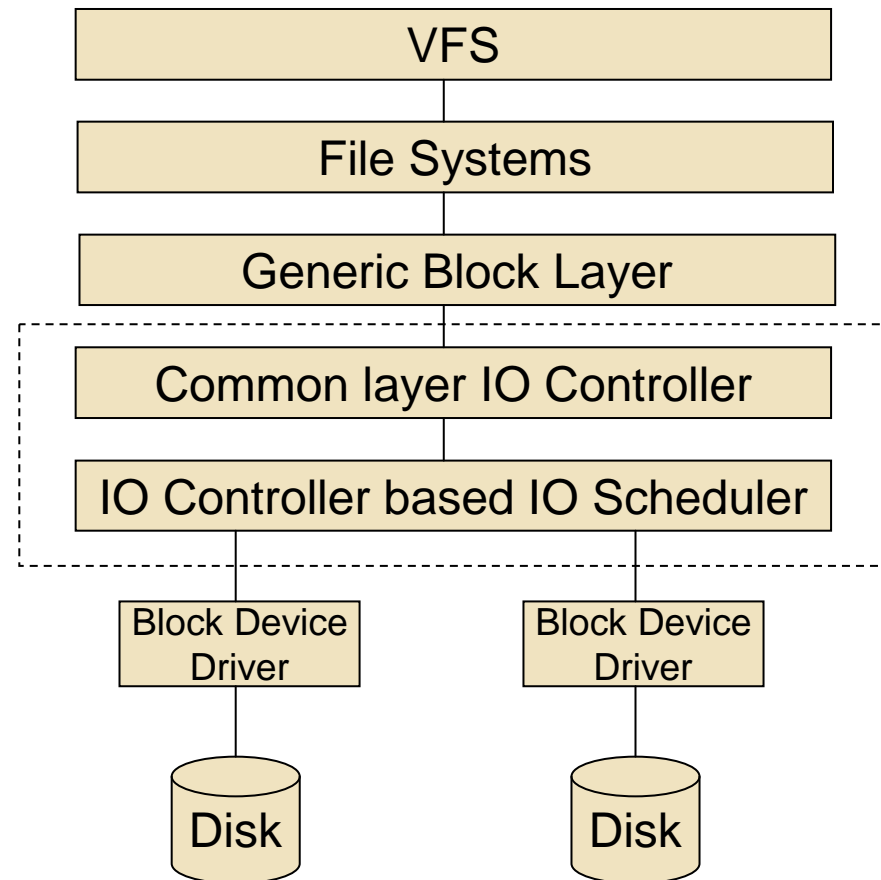
Where does IO Controller locate



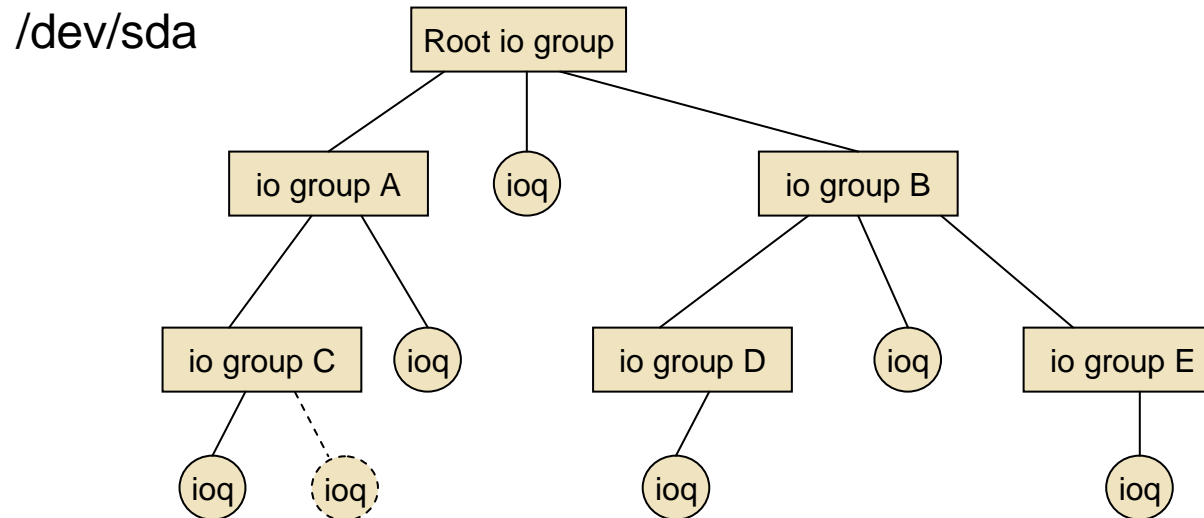
Vanilla kernel



With IO Controller



Scheduling Overview of IO Controller



- Imagine each device as a tree where IO requests stay
- Dynamically create and destroy
- Single ioq for **anticipatory**, **deadline** and **noop** scheduler
- Multi ioq for **cfq** scheduler
- Select one ioq to serve (dispatch requests)

IO Processing in IO Controller



- **bio is submitted into block layer (submit_bio)**
- **Create a new request if merging is not allowed**
- **Find out the CGroup that bio belongs to**
 - for sync io case, tracking CGroup by “current” task
 - for async io case, tracking CGroup by blkio_cgroup
- **Find out corresponding io group**
- **Find out the ioq the request should go into**
- **Insert request into ioq**

- **A CGroup Based Solution**
- **A Proportional and Weight based controller**
- **IO Scheduler Based Solution**
 - Make each IO Scheduler **queue aware**
 - Common layer takes care of queue scheduling
- **Per group requests limitation**
 - `nr_requests` & `nr_group_requests`
 - Make sure that one group won't allocate too many requests and block other groups.

- **A CFS + CFQ Solution**

- Like CFQ: Give time slice to ioq based on its priority
- Like CFS: Decide which ioq should dispatch requests based on the Virtual Disk Time

- **Make difference between sync and async IOs**

- "current" task is used to determine the io group of the sync requests
- Make use of blkio_cgroup to keep track of async requests

- **in-group preemption and in-group merge**

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Problems with IO Controller



- Works only with physical devices (like sda, sdb)
- Doesn't work well if several writers are running
- Doesn't support max bandwidth control
- Which group should be charged if swap writeouts happen
- Extensions of struct bio(bio io_context & bio urgent flag)

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- **The Latest IO Controller Version is V10 (Linux-2.6.31)**

- <http://lkml.org/lkml/2009/9/24/385>

- **Who is working on this project**

- Redhat, Google, Fujitsu

- **The IO Controller mini-summit in Tokyo**

- Only one IO Controller in kernel

- Implement for CFQ only

- Support logical device

- Memcg will implement Per-CGroup dirty ratio and IO Controller will make use of it

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■ **dm-ioband (valinux guys)**

- Implemented as device mapper driver.
- Implemented in higher layers than IO Schedulers (second level)
- Complex setup

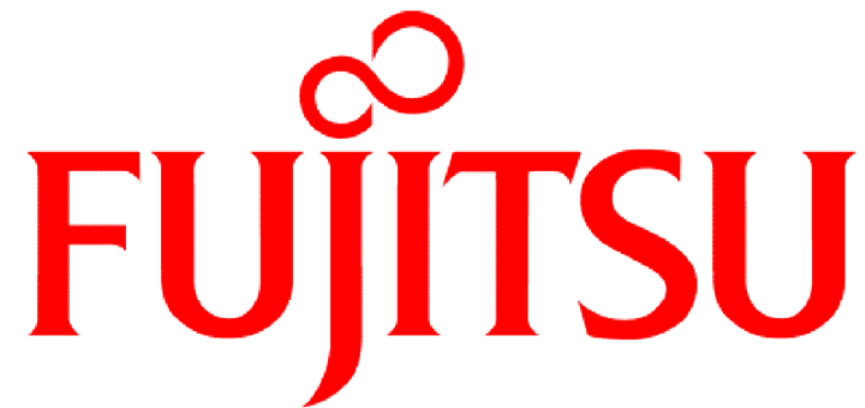
■ **io-throttler (Andrea Righi)**

- Throttle a process if it exceeds the group limit
- Implemented in higher layers than block layer (second level)
- Development has stopped.

- Disk sector vs Disk time accounting
- Bios are buffered in a single queue
 - writes blocks reads issue
 - Doesn't have the notion of io priority
 - More seeks
 - Break the anticipation of underlying IO Scheduler

Thanks!
Questions?





THE POSSIBILITIES ARE INFINITE