

## Week 13 Accessing and Indexing Data (1)

### Assignment Solutions

1. For a disk of seek time = 5 ms, rotational latency = 3 ms, and transfer time = 0.0001 ms per kb, what's the cost of sequential access to get the next 100 kb of data?

☐ 8 ms

☐ 8.0001 ms

☒ 0.01 ms

**Solution: C**

**Explanation:** Sequential access only costs transfer time, which is  $0.0001\text{ms/kb} * 100\text{kb} = 0.01\text{ms}$ .

2. Consider an ISAM with fanout  $F = 341$  and height  $h = 3$ . What's the maximum number of tuples it can store?

Note: The height of an index tree is the number of pointer hops from its root to reach the data tuples.

☐ 116281

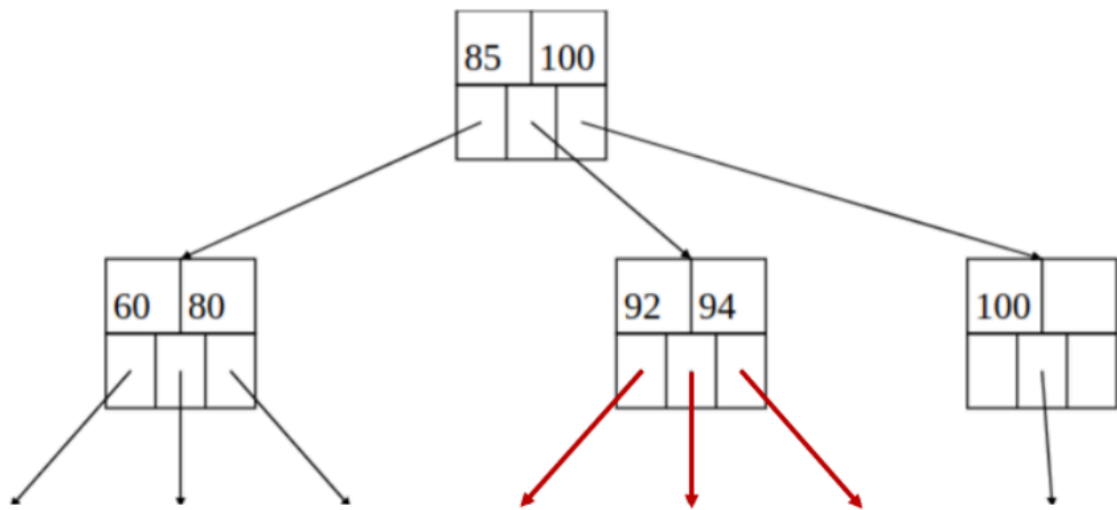
☐ 1000000

☒ 39651821

**Solution: C**

**Explanation:** The total number of tuples available is  $F^h$ , where  $F$  is the fanout and  $h$  is the height, because a node at each of the  $h$  levels can point to  $F$  child nodes. Therefore,  $F^h = 341^3 = 39651821$

3. Considering the following ISAM, what range of values do the pointers in red represent?



☐ (85, 92], (92, 94], (94, 100]

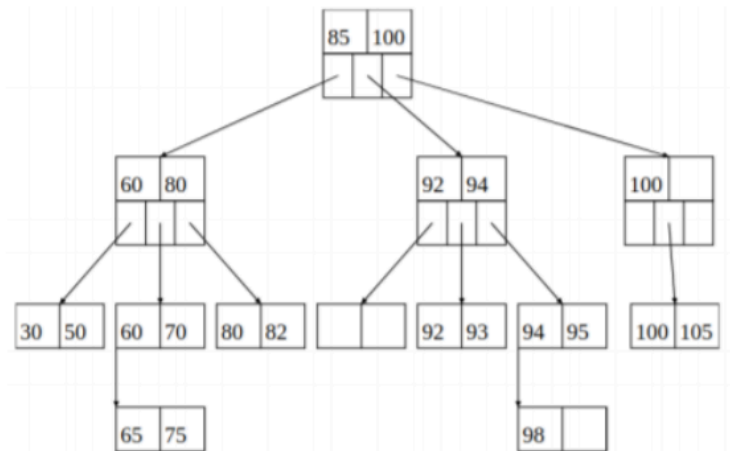
☐  $[-\infty, 92)$ ,  $[92, 94)$ ,  $[94, \infty)$

☒  $[85, 92)$ ,  $[92, 94)$ ,  $[94, 100)$

**Solution: C**

**Explanation:** From the root node, we can identify the range of the middle child node to be  $[85, 100)$ . Thus, the ranges of the three red pointers are further divisions of  $[85, 100)$ , which becomes  $[85, 92)$ ,  $[92, 94)$ , and  $[94, 100)$ .

4. What's the cost of looking up 65 in the following ISAM? Here, we define "cost" as the **number of random accesses** (disk pointers) to reach data tuples.



☐ 3

☐ 2

☐ 5

☒ 4

**Solution: D**

**Explanation:** To reach the data tuple of key 65, we need to go through 3 pointers to reach 65 in ISAM and then another pointer to locate the data tuple.

5. Compared to table scan and binary search, what are the advantages of ISAM?

☐ ISAM takes less time to complete a query.

☐ ISAM automatically sort overflowed pages to balance look up time for any keys.

☐ ISAM requires fewer disk random accesses.

**Solution: AC**

**Explanation:** Table scan takes  $O(n)$ , and binary search takes  $O(\log_2 n)$ . ISAM takes  $O(\log_F n)$ , where  $F$  is the fanout. As fanout is usually much higher than 2, ISAM requires fewer disk pointer traversals (random accesses) -- and thus also takes less time to complete a query.