

# CS 3100, Models of Computation, Spring 2020, Lec 3

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[bit.ly/3100s20Syllabus](https://bit.ly/3100s20Syllabus)



# Lecture 3, covering up to Chapter 4.6

# Introducing DFA

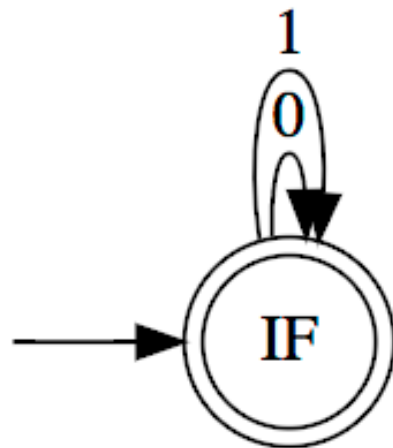
- Machines with a **finite set of states** : “Q”
- Is **defined for an alphabet**  $\Sigma$  (usually  $\{0,1\}$ )
- Has a **single initial state**
- Has **some number of final states** “F”
  - Could be none
  - Could be all of Q
- Has moves defined
  - for every state  $q$  in Q
  - for every member of  $\Sigma$
- **Helps define a language!**
  - Any string that takes the DFA from I to one of the states in “F” is in the DFA’s language
  - The DFA accepts such a string
- **All the strings that go from I to F are in the language recognized by the DFA**

Is this a DFA?  
Which strings are accepted by it?  
What language is recognized by it?



```
In [12]: 1 # Do this in CH4-5-6
          2
          3 dotObj_dfa(md2mc(''
          4
          5 DFA
          6
          7 IF : 0|1 -> IF
          8
          9 '''))
```

Out[12]:

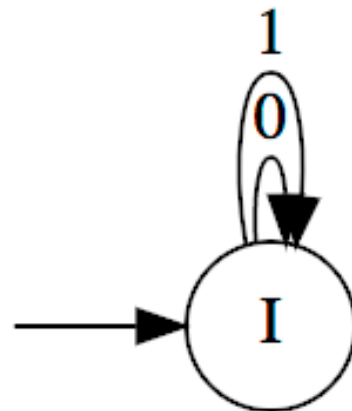


Is this a DFA?  
Which strings are accepted by it?  
What language is recognized by it?



```
In [11]: 1 # Do this in CH4-5-6
          2 |
          3 dotObj_dfa(md2mc(''
          4
          5 DFA
          6
          7 I : 0|1 -> I
          8
          9 '''))
```

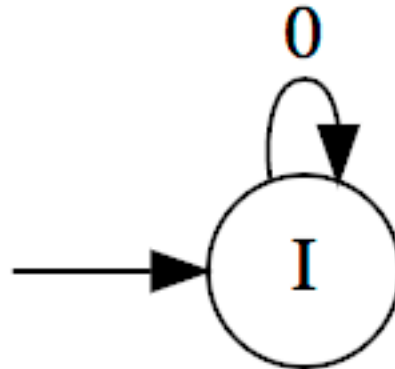
Out[11]:



Is this even a DFA over  $\{0,1\}$  ?? Why or why not?

```
In [13]: 1 # Do this in CH4-5-6
        2
        3 dotObj_dfa(md2mc(''
        4
        5 DFA
        6
        7 I : 0 -> I
        8
        9 '''))
```

Out[13]:

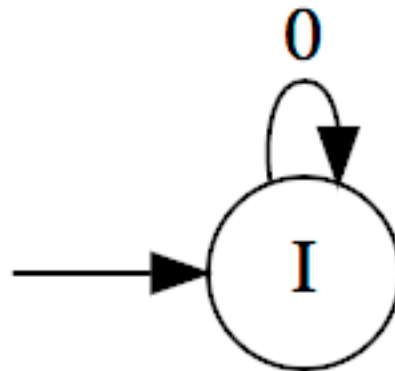


# Is this a DFA over $\{0\}$ ?



```
In [13]: 1 # Do this in CH4-5-6
          2
          3 dotObj_dfa(md2mc(''
          4
          5 DFA
          6
          7 I : 0 -> I
          8
          9 '''))
```

Out[13]:

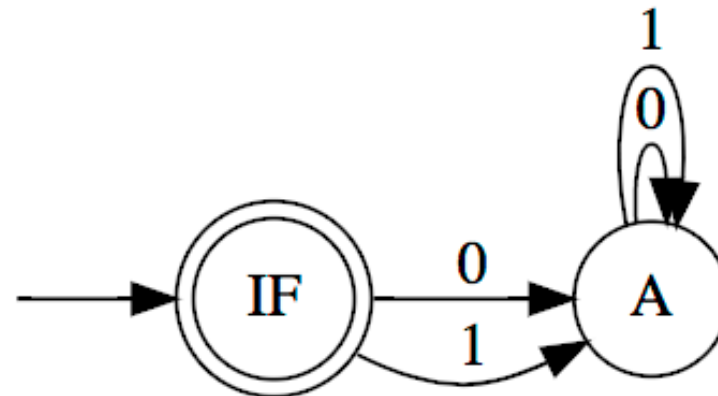


Is this a DFA?  
Which strings are accepted by it?  
What language is recognized by it?



```
In [14]: 1 # Do this in CH4-5-6
        2
        3 dotObj_dfa(md2mc(''
        4
        5 DFA
        6
        7 IF : 0|1 -> A
        8
        9 A : 0|1 -> A
       10
       11 '''))
```

Out[14]:



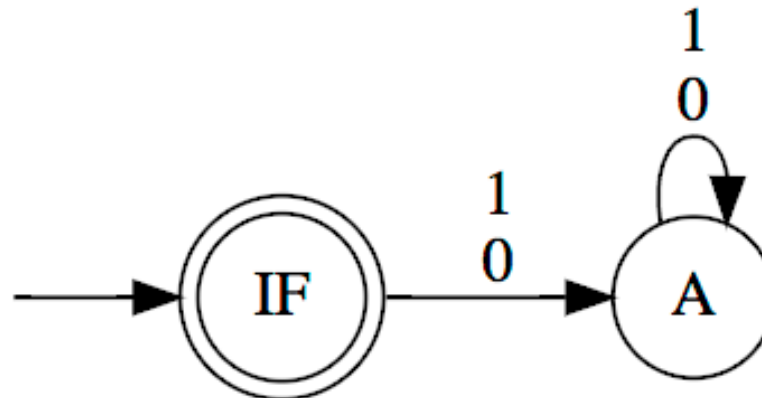


Is this a DFA?  
Which strings are accepted by it?  
What language is recognized by it?



```
In [15]: 1 # Do this in CH4-5-6
          2
          3 dotObj_dfa(md2mc(''
          4
          5 DFA
          6
          7 IF : 0|1 -> A
          8
          9 A : 0|1 -> A
         10
         11 '''), FuseEdges=True)
```

Out[15]:

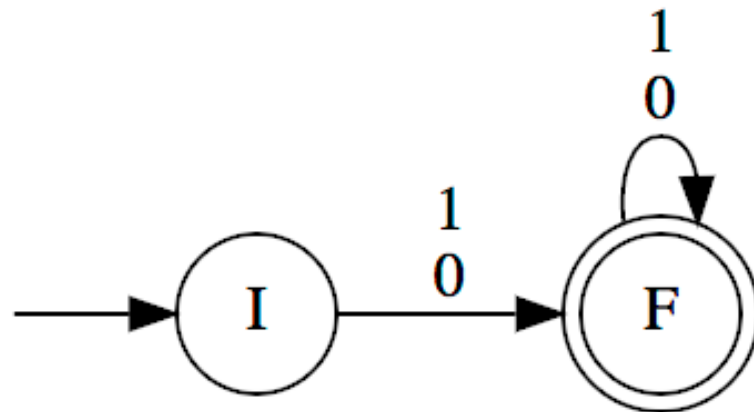


Is this a DFA?  
Which strings are accepted by it?  
What language is recognized by it?



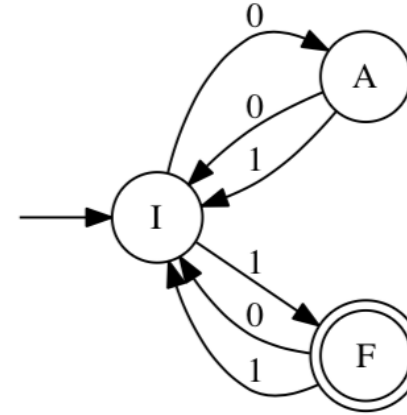
```
In [16]: 1 # Do this in CH4-5-6
          2
          3 dotObj_dfa(md2mc(''
          4
          5 DFA
          6
          7 I : 0|1 -> F
          8
          9 F : 0|1 -> F
         10
         11 '''), FuseEdges=True)
```

Out[16]:



# Formal structure of a DFA, and an example

- $Q$  is a *finite nonempty* set of states,
- $\Sigma$  is a *finite nonempty* alphabet,
- $\delta : Q \times \Sigma \rightarrow Q$  is a *total* transition function,
- $q_0 \in Q$  is the initial state, and
- $F \subseteq Q$  is a *finite, possibly empty* set of final (or *accepting*) states.



```
{'Q': {'A', 'F', 'I'},  
'Sigma': {'0', '1'},  
'Delta': {('A', '0'): 'I',  
          ('A', '1'): 'I',  
          ('F', '0'): 'I',  
          ('F', '1'): 'I',  
          ('I', '0'): 'A',  
          ('I', '1'): 'F'},  
'q0': 'I',  
'F': {'F'}}
```

The language of a DFA is

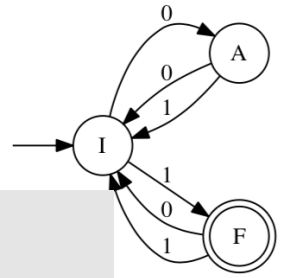
“ the strings that take the DFA from the initial state to one of the final states ”

(if the initial state is also a final state, then the language also contains epsilon)



List five strings in this DFA's language (in numeric order)

# step\_dfa, run\_dfa, accepts\_dfa in code



```
def step_dfa(D, q, a):  
    """Run DFA D from state q on character a.  
    """  
  
    assert(a in D["Sigma"])  
    assert(q in D["Q"])  
    return D["Delta"][(q,a)]
```

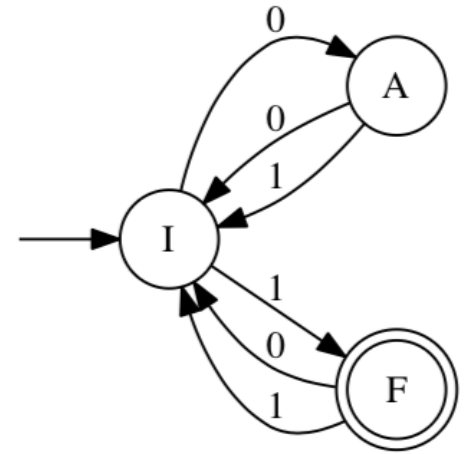
```
def accepts_dfa(D, w):  
    """ Checks for DFA acceptance. Input : DFA D, string w.  
    Output : Boolean (True|False).  
    """  
  
    return run_dfa(D, w) in D["F"]
```

```
def run_dfa(D, w):  
    """In : D (consistent DFA)  
            w (string over D's sigma, including "")  
    Out: next state of D["q0"] via string w.  
    """  
  
    curstate = D["q0"]  
    if w=="":  
        return curstate  
    else:  
        return run_dfa_h(D, w[1:], step_dfa(D,curstate,w[0]))
```

```
def run_dfa_h(D, w, q):  
    """Helper for run_dfa. Compute the next state attained  
    by w running on D starting from state q.  
    """  
  
    if w=="":  
        return q  
    else:  
        return run_dfa_h(D, w[1:], step_dfa(D, q, w[0]))
```

# step\_dfa, run\_dfa, accepts\_dfa in math

- $Q$  is a *finite nonempty* set of states,
- $\Sigma$  is a *finite nonempty* alphabet,
- $\delta : Q \times \Sigma \rightarrow Q$  is a *total* transition function,
- $q_0 \in Q$  is the initial state, and
- $F \subseteq Q$  is a *finite, possibly empty* set of final (or *accepting*) states.

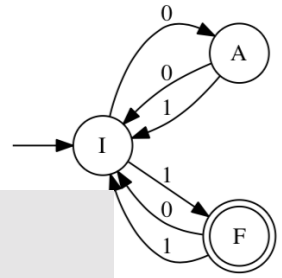


“step\_dfa”       $\delta(D, q, a)$

“run\_dfa”       $\hat{\delta}(D, q, ax) = \hat{\delta}(D, \delta(D, q, a), x)$   
 $\hat{\delta}(D, q, \varepsilon) = q.$

“accepts\_dfa”       $\hat{\delta}(D, q_0, w) \in F.$

# step\_dfa, run\_dfa, accepts\_dfa in both!



```
def step_dfa(D, q, a):  $\delta(D, q, a)$ 
    """Run DFA D from state q on character a.
    """
    assert(a in D["Sigma"])
    assert(q in D["Q"])
    return D["Delta"][(q,a)]
```

```
def accepts_dfa(D, w):
    """ Checks for DFA acceptance. Input : DFA D, string w.
        Output : Boolean (True|False).
    """
    return run_dfa(D, w) in D["F"]
```

$$\hat{\delta}(D, q_0, w) \in F.$$

```
def run_dfa(D, w):
    """In : D (consistent DFA)
           w (string over D's sigma, including "")
       Out: next state of D["q0"] via string w.
    """
    curstate = D["q0"]
    if w=="":
        return curstate
    else:
        return run_dfa_h(D, w[1:], step_dfa(D, curstate, w[0]))
```

$$\hat{\delta}(D, q, \epsilon) = q.$$

$$\hat{\delta}(D, q, ax) = \hat{\delta}(D, \delta(D, q, a), x)$$

```
def run_dfa_h(D, w, q):
    """Helper for run_dfa. Compute the next state attained
       by w running on D starting from state q.
    """
    if w=="":
        return q
    else:
        return run_dfa_h(D, w[1:], step_dfa(D, q, w[0]))
```

# Design a DFA for “ends with 0101”

Express mathematically as a language

Design in Jove syntax, keeping the relevant piece of the input in the state name

i.e. S010 means “seen 010 so far”

# Design a DFA for “contains 0101”



Express mathematically as a language

Design in Jove syntax, keeping the relevant piece of the input in the state name  
i.e. S010 means “seen 010 so far”

Some prefer to draw the DFA out also. But still, name states consistently.



Design a DFA for  $\{0,1\}$  where every block of length 3 contains exactly two 1's

Design a DFA for strings over  $\{0,1\}$  whose numeric value is a multiple of 3