

# CS 3100, Models of Computation, Fall 2019, Lec 8

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# Designing a DFA (General)

- DFA are simplified C / Java programs with finite states

```
state = I or IF
```

```
While(True) {
```

```
  switch state {
```

```
    case 'IF' : switch input {  
                  case '0' : state = ...update it...  
                  case '1' : state = ...update it...  
                }
```

```
    case 'A' : ...
```

```
  }
```

```
}
```

- More realistic programs are modeled by PDA and all programs are modeled by TM

# Why NFA? Many answers!

## 1. Invented to overcome the limitations of a DFA

- For some regular languages (that have a DFA), the DFA are **exponentially big**
- **In many of those cases, an NFA will be linear / polynomial in size**
- **Therefore reduces tedium (for humans) to specify**
- **This use of NFA also turns into a syntactic approach called Regular Expressions**

## 2. Nondeterminism is a fundamental idea in CS

- **Allows us to classify algorithms into “easy” (P) and “hard” (NP)**

# Features of an NFA

- Finite states
- Multiple initial states
  - Can begin in any initial state
- Transitions on Sigma
- Transitions also on Epsilon
  - Recall that Epsilon is not in Sigma
- Transitions lead to sets of next states
- Has final states (like before)
- Acceptance:
  - Begin at any initial state
  - A journey described by a string (laden perhaps with Epsilon)
  - Ends in a final state

Two NFA designs for “third last is a 1”

## NFA have equivalent DFA - Subset Construction (function nfa2dfa)

- Basically for each NFA that is in a set of states  $\{S1, S2, S3\}$ 
  - For example we assume a set of 3 states an NFA is in.
- First E-close  $\{S1, S2, S3\}$ 
  - Let S1 go to S11, S12 on “
  - Let S2 go to S21 on “
  - Let S3 go on S31, S32, S33 on “
  - E-closure (  $\{S1, S2, S3\}$  ) =  $\{S11, S12, S21, S31, S32, S33\}$
- Fire a symbol, say a in Sigma from each of S11, S12,...S33
- Let the resulting SET OF STATES be
  - $S11', S12', \dots, S33'$  (these are SETS of states)
- Take a set union of S11', S12', ..., S33'
  - E-close that state.
- This is what  $\{S1, S2, S3\}$  transitions to, upon an “a” in Sigma
- Do this for Book77 NFA
- Do this for “third last is a 1” NFA

Book's Figure 7.7 NFA : Write in Jove's markdown in 2 ways,  
then convert to DFA via Subset Construction

