(Not yet Adaptive) Compression of In-Memory Databases

Database Implementation Lab Course

Leon Windheuser

February 28, 2023

Project Introduction

We want to compress the transient part of DuckDB

Project Introduction

We want to compress the transient part of DuckDB

- Open Source SQL OLAP RDBMS in-process developed in Amsterdam research centre CWI (SQLite for OLAP) https://github.com/duckdb/duckdb
- Columnar Storage format
- Vectorized execution engine
- Has already lots of different compression possibilities for persistent data on disk

Project Introduction

We want to compress the transient part of DuckDB

- Open Source SQL OLAP RDBMS in-process developed in Amsterdam research centre CWI (SQLite for OLAP) https://github.com/duckdb/duckdb
- Columnar Storage format
- Vectorized execution engine
- Has already lots of different compression possibilities for persistent data on disk

How do we compress the transient data while having efficient lookups without decompressing everything?

Background: Succinct Data Structures

- Data structures which uses space close to the theoretic lower bound but allows efficient query operations (in-place without needing to decompress)
- Exists for e.g. (bit) vectors, trees, planar graphs, ...

Background: Succinct Data Structures

- Data structures which uses space close to the theoretic lower bound but allows efficient query operations (in-place without needing to decompress)
- Exists for e.g. (bit) vectors, trees, planar graphs, ...

Succinct Integer Vector

Space requirement for integer x is $\ell = \lfloor \log_2(x) \rfloor + 1$ bits

0
1
1
0
0
0
1
0
1
0
0
0
1
0

= 7
= 0
= 5
= 4
= 6

Succinct Integer Vector

Space requirement for integer x is $\ell = \lfloor \log_2(x) \rfloor + 1$ bits



Encode integers with the minimal length of the max integer $3 = |\log_2(7)| + 1$



We already reduce memory by 25%

SDSL: Succinct Data Structure Library

- ► C++11 library and abstraction for succinct data structures
- ▶ Open Source https://github.com/simongog/sdsl-lite
- Contains varierty of different data structures. For now we only used the Integer Vectors.

SDSL: Integer Vectors

SDSL: Integer Vectors

SDSL: Integer Vectors

```
sdsl::int_vector <32> v(10000);
for (size_t i = 0; i < 10000; i++) v[i] = i;</pre>
cout << "Width: " << v.width() << ", size: "</pre>
     << sdsl::size_in_bytes(v) << endl;
sdsl::util::bit_compress(v);
cout << "Width: " << v.width() << ", size: "
     << sdsl::size_in_bytes(v) << endl;
Width: 32, size: 40008
Width: 14, size: 17513
```

Reduces memory by 56.2% (\approx 22.5 KB)

Delta compression of sdsl::int_vector

```
sdsl::int_vector <32> v(10000);
for (size_t i = 0; i < 10000; i++)</pre>
    v[i] = i + 10.000.000;
cout << "Width: " << v.width() << ", size: "</pre>
     << sdsl::size_in_bytes(v) << endl;
sdsl::util::bit_compress(v);
cout << "Width: " << v.width() << ", size: "</pre>
     << sdsl::size_in_bytes(v) << endl;
Width: 32, size: 40008
Width: 24, size: 30008
```

Delta compression of sdsl::int_vector

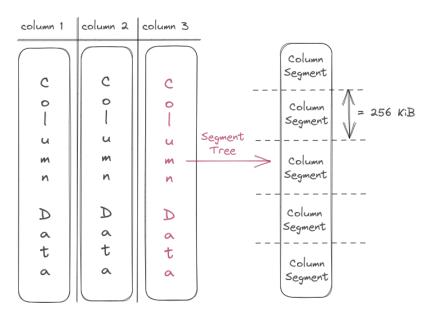
```
sdsl::int_vector <32> v(10000);
for (size_t i = 0; i < 10000; i++)</pre>
    v[i] = i + 10.000.000;
cout << "Width: " << v.width() << ", size: "</pre>
     << sdsl::size_in_bytes(v) << endl;
extractMinFromVector(v);
sdsl::util::bit_compress(v);
cout << "Width: " << v.width() << ", size: "
     << sdsl::size_in_bytes(v) << endl;
Width: 32, size: 40008
Width: 14, size: 17513
```

Delta compression of sdsl::int_vector

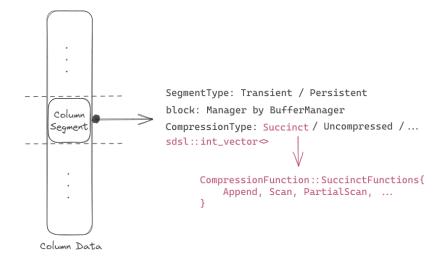
```
sdsl::int_vector <32> v(10000);
for (size_t i = 0; i < 10000; i++)</pre>
    v[i] = i + 10.000.000;
cout << "Width: " << v.width() << ", size: "</pre>
     << sdsl::size_in_bytes(v) << endl;
extractMinFromVector(v);
sdsl::util::bit_compress(v);
cout << "Width: " << v.width() << ", size: "
     << sdsl::size_in_bytes(v) << endl;
Width: 32, size: 40008
Width: 14, size: 17513
```

In DuckDB we know the minimum of the vector directly without searching (column statistics)

DuckDB Storage Architecture 100 meter view



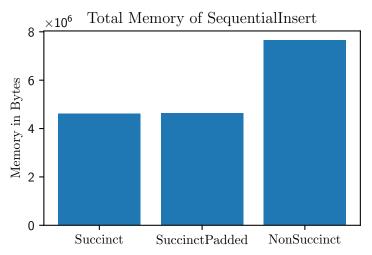
DuckDB Storage Architecture 10 meter view



Benchmarks: Sequential Insert and total Scan

Scanning 10⁶ rows.

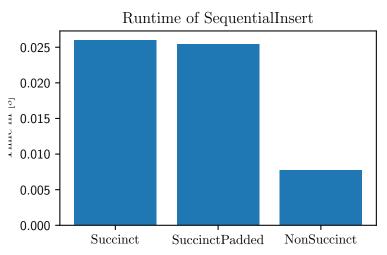
```
SELECT * FROM t1;
```



Benchmarks: Sequential Insert and total Scan

Scanning 10⁶ rows.

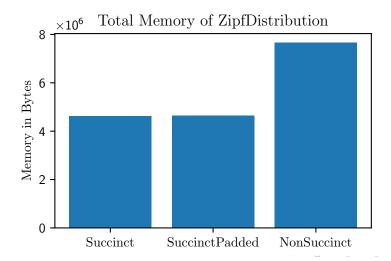
```
SELECT * FROM t1;
```



Benchmarks: Zipf Selection

10.000 selections with Zipf Distribution of 10⁶ total rows.

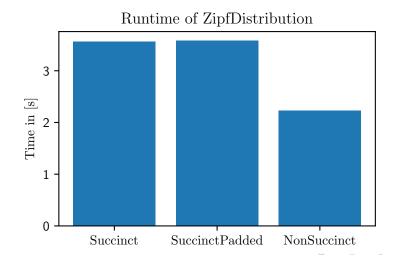
```
SELECT i FROM t1
WHERE i == {ZIPF_DISTRIBUTED_NUMBER};
```



Benchmarks: Zipf Selection

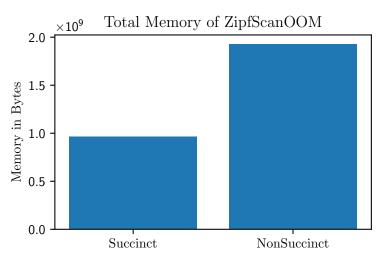
10.000 selections with Zipf Distribution of 10⁶ total rows.

```
SELECT i FROM t1
WHERE i == {ZIPF_DISTRIBUTED_NUMBER};
```



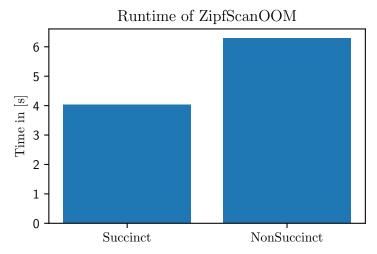
Benchmarks: Zipf Out Of Memory (Limit 1GB)

```
SELECT i FROM t1
WHERE i == {ZIPF_DISTRIBUTED_NUMBER};
```



Benchmarks: Zipf Out Of Memory (Limit 1GB)

```
SELECT i FROM t1
WHERE i == {ZIPF_DISTRIBUTED_NUMBER};
```



Conclusion and Future Work

- ► For OLAPish queries it is not (yet?) worth it, large overhead.
- ▶ For OLTP transactions it might be worth it. Reduces memory by $\approx 40\%$ but increases runtime by $\approx 35\%$.
- Huge benefit if succinct representation fits in memory vs spilling to disk.

Conclusion and Future Work

- ► For OLAPish queries it is not (yet?) worth it, large overhead.
- ▶ For OLTP transactions it might be worth it. Reduces memory by $\approx 40\%$ but increases runtime by $\approx 35\%$.
- Huge benefit if succinct representation fits in memory vs spilling to disk.

Future Work:

- Need to investigate in access statistics to only compress rarely used column segments.
- Currently copying a lot of data as execution engine expects an uncompressed flat vector. Adapt execution engine to allow succinct vectors as well since they allow random access.