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CSCE625

Problem 1:

Step 1: first order logic

1. $\forall x (Crow(x) \wedge \neg White(x) \rightarrow \exists y (Jewel(y) \wedge Loves(x, y)))$
2. $\forall x (Talk(x) \rightarrow Intel(x))$
3. $\forall x (Intel(x) \rightarrow \forall y (Book(y) \rightarrow Loves(x, y)))$
4. $\neg \exists x (\forall y (Book(y) \rightarrow Loves(x, y)) \wedge \exists z (Jewel(z) \wedge Loves(x, z)))$
5. $Crow(Jack) \wedge Talk(Jack)$

Step 2: convert to canonical form:

For 1: $\forall x (Crow(x) \wedge \neg White(x) \rightarrow \exists y (Jewel(y) \wedge Loves(x, y)))$

Using Skolemization for the existential quantifier: $f(x)$ is a Skolem function, gives a jewel that x loves.

the formula becomes:

$$\forall x (Crow(x) \wedge \neg White(x) \rightarrow (Jewel(f(x)) \wedge Loves(x, f(x))))$$

To get the CNF, using the formula $\neg P \vee Q$. So, the formula becomes:

$$\forall x (\neg Crow(x) \vee White(x) \vee (Jewel(f(x)) \wedge Loves(x, f(x))))$$

This gives us two clauses:

- a) $\forall x \neg Crow(x) \vee White(x) \vee Jewel(f(x))$
- b) $\forall x \neg Crow(x) \vee White(x) \vee Loves(x, f(x))$

For 2: $\forall x (Talk(x) \rightarrow Intel(x))$

the formula becomes: $\forall x (\neg Talk(x) \vee Intel(x))$

For 3: $\forall x (Intel(x) \rightarrow \forall y (Book(y) \rightarrow Loves(x, y)))$

the formula becomes: $\forall x \forall y (\neg Intel(x) \vee \neg Book(y) \vee Loves(x, y))$

For 4: $\neg \exists x (\forall y (Book(y) \rightarrow Loves(x, y)) \wedge \exists z (Jewel(z) \wedge Loves(x, z)))$

Do the negation, and the formula becomes:

$$\forall x (\exists y (Book(y) \wedge \neg Loves(x, y)) \vee \neg \exists z (Jewel(z) \wedge Loves(x, z)))$$

After negating:

$$\forall x (\exists y (Book(y) \wedge \neg Loves(x, y)) \vee \forall z (\neg Jewel(z) \vee \neg Loves(x, z)))$$

And then do the skolemization for the existential quantifier y , $g(x)$ is a skolem function, specific book that x does not love.

$$\forall x ((Book(g(x)) \wedge \neg Loves(x, g(x))) \vee \forall z (\neg Jewel(z) \vee \neg Loves(x, z)))$$

For 5, $Crow(Jack) \wedge Talk(Jack)$

This gives us two clauses:

- a) $Crow(Jack)$
- b) $Talk(Jack)$

Step 3: Proof by resolution:

List of clauses:

For 1: two clauses:

- a) $\neg \text{Crow}(x) \vee \text{White}(x) \vee \text{Jewel}(f(x))$
- b) $\neg \text{Crow}(x) \vee \text{White}(x) \vee \text{Loves}(x, f(x))$

For 2: $(\neg \text{Talk}(x) \vee \text{Intel}(x))$

For 3: $(\neg \text{Intel}(x) \vee \neg \text{Book}(y) \vee \text{Loves}(x, y))$

For 4: $((\text{Book}(g(x)) \wedge \neg \text{Loves}(x, g(x))) \vee (\neg \text{Jewel}(z) \vee \neg \text{Loves}(x, z)))$

For 5, two clauses:

- a) $\text{Crow}(\text{Jack})$
- b) $\text{Talk}(\text{Jack})$

Conclusion: $\exists x(\text{Crow}(x) \wedge \text{White}(x))$

Negated conclusion: $\neg \exists x(\text{Crow}(x) \wedge \text{White}(x))$

For 6: After negating: $\forall x(\neg \text{Crow}(x) \vee \neg \text{White}(x))$

Resolution step:

For 7: From 5b and 2: $\text{Intel}(\text{Jack})$

For 8: From 5a and 6: $\neg \text{White}(\text{Jack})$

For 9: From 1a and 5a: $\text{White}(\text{Jack}) \vee \text{Loves}(\text{Jack}, f(\text{Jack}))$

For 10: From 8 and 9: $\text{Loves}(x, f(x))$

For 11: From 7 and 3: $\neg \text{Book}(y) \vee \text{Loves}(\text{Jack}, y) \rightarrow \text{false}$, conflict with premise 4

Here is the premise 4: $((\text{Book}(g(x)) \wedge \neg \text{Loves}(x, g(x))) \vee (\neg \text{Jewel}(z) \vee \neg \text{Loves}(x, z)))$

Take the first part: $\text{Book}(g(x)) \wedge \neg \text{Loves}(x, g(x))$

This directly contradicts the derived statement that Jack loves all books, and hence we've reached the required contradiction which shows that our initial negation of the conclusion was false which there must be at least one white crow. Therefore, the proof by resolution shows that there exists a white crow is a logical consequence of the given premises.