

Euler Problem 1 Writeup

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1 Problem Statement:

If we list all the natural numbers below 10 that are multiples of 3 or 5, we get 3, 5, 6 and 9. The sum of these multiples is 23.

Find the sum of all the multiples of 3 or 5 below 1000.

1.1 Coding adaption

Hacker Rank provides the following code adaption:

If we list all the natural numbers below 10 that are multiples of 3 or 5 , We get 3, 5, 6, 9. The sum of these multiples is 23. Find the sum of all the multiples of 3 or 5 below N .

Input Format

First line contains T that denotes the number of test cases. This is followed by T lines, each containing an integer, N .

Constraints

- $1 \leq T \leq 10^5$
- $1 \leq N \leq 10^9$

Output Format

For each test case, print an integer that denotes the sum of all the multiples of 3 or 5 below N .

1.2 Solution

1.2.1 bash

My first attempt was simply to bash this. There are quicker ways and it could all be done in main, but I found this to be the most "human readable" way of doing it.

```
#include <math.h>
#include <stdio.h>

long SumUnderN(long);

int main(){
    int i; //counter
    int T; // # of testcases
    long N; // place to store each instance of N.

    scanf("%d",&T);

    for (i = 0; i < T; i++)
    {
        scanf("%ld",&N);
        printf("%ld\n",SumUnderN(N));
    }
    return 0;
}

long SumUnderN(long N)
{
    long i; //counter
    long sum = 0; // sum
    for(i = 1; i < N; i++)
    {
        sum += i * ((i % 3 == 0) || (i % 5 == 0));
    }
    return sum;
}
```

It should be noted the long data type needs to be used as N can be up to 10^9 .

1.2.2 Optimization / Math

The problem with this method is it's inefficient. For most problems like this with modern computational power this is not a problem, however for the sake of good practice we should do better. As such Hackerank fails two of the automated testcases as it is also testing for this.

The multiples of 3 and 5 are by definition at regular intervals. We could edit our for loops to iterate twice through the given number N , in steps of 3 then 5, but this still isn't optimal. Some basic Number Theory and inclusion/exclusion tell us this can be calculated for a given number with division. If we add up all the multiples of 3, then the multiples of 5 then subtract the multiples of 15 we get our answer. Thus we get:

$$\begin{aligned} \text{answer} &= \left(\sum_{k=1}^{\lfloor \frac{N}{3} \rfloor} (3 \cdot k) \right) + \left(\sum_{k=1}^{\lfloor \frac{N}{5} \rfloor} (5 \cdot k) \right) - \left(\sum_{k=1}^{\lfloor \frac{N}{15} \rfloor} (15 \cdot k) \right) \\ &= 3 \cdot \frac{\lfloor \frac{N}{3} \rfloor (\lfloor \frac{N}{3} \rfloor + 1)}{2} + 5 \cdot \frac{\lfloor \frac{N}{5} \rfloor (\lfloor \frac{N}{5} \rfloor + 1)}{2} - 15 \cdot \frac{\lfloor \frac{N}{15} \rfloor (\lfloor \frac{N}{15} \rfloor + 1)}{2} \end{aligned}$$

We subtract the sum of multiples up to $\lfloor \frac{N}{15} \rfloor$ to remove the double counting of numbers that are both multiples of 3 and 5. We can implement this in code by changing our SumUnderN function using some custom functions from the math.h header file.

```
long SumUnderN(long N)
{
    long M3 = floor(((N-1)/3)); // multiples of 3
    long M5 = floor(((N-1)/5)); // multiples of 5
    long M15 = floor(((N-1)/15)); // multiples of 15
    long sum = (3*(M3*(M3+1))/2) +
               (5*(M5*(M5+1))/2) -
               (15*(M15*(M15+1))/2);
    return sum;
}
```