ECE653
Software Testing, Quality Assurance, and Maintenance
Assignment 2 (70 Points), Version 1

Instructor: Werner Dietl
Release Date: June 11, 2024

Due: 22:00, June 28, 2024
Submit: An electronic copy on GitLab

Any source code and test cases for the assignment will be released in the skeleton repository at `https://git.uwaterloo.ca/stqam-1245/skeleton`.

I expect each of you to do the assignment independently. I will follow UW's Policy 71 for all cases of plagiarism.

Submission Instructions:

Please read the following instructions carefully. **If you do not follow the instructions, you may be penalized up to 5 points.** Illegible answers receive no points.

Submit by pushing your changes to the `main` branch of your GitLab repository in directory `a2/`. Make sure to use a web browser to check that your changes have been committed! The submission must contain the following:

- a `user.yml` file with your UWaterloo user information;

- a single pdf file called `a2_sub.pdf`. The first page must include your full name, 8-digit student number and your uwaterloo email address;

- a directory `a2q3` that includes your code for Question 3; and

- a directory `wlang` that includes your code for Question 4.

After submission, **review your submissions on GitLab web interface to make sure you have uploaded the right files/versions.**

You can push changes to the repository before and after the deadline. We will use the latest commit at the time of deadline for marking.

Question 1 (10 points)

Consider the following program `Prog1`:

```
1   havoc x, y;
2   if x + y > 15 then {
3     x = x + 7;
4     y = y - 12 }
5   else {
6     y = y + 10;
7     x = x - 2 };
8
9   x = x + 2;
10
11  if 2 * (x + y) > 21 then {
12    x = x * 3;
13    y = y * 2 }
14  else {
15    x = x * 4;
16    y = y * 3 + x };
17  skip
```

(a) How many execution paths does `Prog1` have? List all the paths as a sequence of line numbers taken on the path.

(b) Symbolically execute each path and provide the resulting path condition. Show the steps of symbolic execution as a table. An example of executing the first line is given below:

Edge	Symbolic State (PV)	Path Condition (PC)
$1 \to 2$	$x \mapsto X_0, y \mapsto Y_0$	true
...	...	...

(c) For each path in part (b), indicate whether it is feasible or not. For each feasible path, give values for X_0 and Y_0 that satisfy the path condition.

path 1:
  1,2, 3,4 9, 11 , 12,13 ,17 ✓

path 2:
  1,25,6,7, 9,11,12,13, 17 ✓   4 paths

path 3: 1,2, 5,6,7,9, 11, 14,15,16,17

path 4: 1, 2, 3.4.9, 11 , 14 15 16 17 ✓

b)

path 1

Edge	Symbolic State (PV)	Path Condition (PC)
$1 \to 2$	$x \mapsto X_0, y \mapsto Y_0$	true
$2 \to 3$	$X = X_0 + 7, \quad Y = Y_0$	$X_0 + Y_0 > 15$
$3 \to 4$	$X = X_0 + 7, \quad Y = Y_0 - 12$	$X_0 + Y_0 > 15$
$4 \to 9$	$X = X_0 + 9, \quad Y = Y_0 - 12$	$X_0 + Y_0 > 15$
$9 \to 11$	$X = X_0 + 9, \quad Y = Y_0 - 12$	$X_0 + Y_0 > 15 \wedge 2*(X+Y-3) > 21$
$11 \to 12$	$X = 3X_0 + 27, \quad Y = Y_0 - 12$	$X + Y > 15 \wedge 2*(X+Y-3) > 21$
$12 \to 13$	$X = 3X_0 + 27, \quad Y = 2Y_0 - 24$	$X + Y > 15 \wedge 2*(X+Y-3) > 21$
$13 \to 17$	$X = 3X_0 + 27, \quad Y = 2Y_0 - 24$	$X + Y > 15 \wedge 2*(X+Y-3) > 21$

path 2

Edge	Symbolic State (PV)	Path Condition (PC)
$1 \to 2$	$x \mapsto X_0, y \mapsto Y_0$	true
$2 \to 5$	$x \mapsto X_0,\ y \mapsto Y_0$	$X_0 + Y_0 \le 15$
$5 \to 6$	$x \mapsto X_0,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15$
$6 \to 7$	$x \mapsto X_0 - 2,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15$
$7 \to 9$	$x \mapsto X_0,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15$
$9 \to 11$	$x \mapsto X_0,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15 \ \wedge\ 2*(X_0 + Y_0 + 10) > 21$
$11 \to 12$	$x \mapsto 3X_0,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15 \ \wedge\ 2*(X_0 + Y_0 + 10) > 21$
$12 \to 13$	$x \mapsto 3X_0,\ y \mapsto 2Y_0 + 20$	$X_0 + Y_0 \le 15 \ \wedge\ 2*(X_0 + Y_0 + 10) > 21$
$13 \to 17$	$x \mapsto 3X_0,\ y \mapsto 2Y_0 + 20$	$X_0 + Y_0 \le 15 \ \wedge\ 2(X_0 + Y_0 + 10) > 21$

path 3:

path 3: 1 2, 5, 6, 7, 9, 15 16, 17

Edge	Symbolic State (PV)	Path Condition (PC)
$1 \to 2$	$x \mapsto X_0, y \mapsto Y_0$	true
$2 \to 5$	$x \mapsto X_0, y \mapsto Y_0$	$X_0 + Y_0 \le 15$
$5 \to 6$	$x \mapsto X_0,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15$
$6 \to 7$	$x \mapsto X_0 - 2,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15$
$7 \to 9$	$x \mapsto X_0,\ y \mapsto Y_0 + 10$	$X_0 + Y_0 \le 15$
$9 \to 11$	$x \mapsto X,\ y \mapsto Y + 10$	$X_0 + Y_0 \le 15 \ \wedge\ 2*(X + Y + 10) \le 21$
$11 \to 14$	$x \mapsto X,\ y \mapsto Y + 10$	$X_0 + Y_0 \le 15 \ \wedge\ 2*(X_0 + Y_0 + 10) \le 21$
$14 \to 15$	$x \mapsto 4X,\ y \mapsto Y + 10$	$X_0 + Y_0 \le 15 \ \wedge\ 2*(X + Y + 10) \le 21$
$15 \to 16$	$x \mapsto 4X,\ y \mapsto 3Y_0 + 10 + 4X$	$X + Y \le 15 \ \wedge\ 2(X_0 + Y_0 + 10) \le 21$
$16 \to 17$	$x \mapsto 4X,\ y \mapsto 3Y_0 + 10 + 4X$	$X + Y \le 15 \ \wedge\ 2(X_0 + Y_0 + 10) \le 21$

```
1  havoc x, y;
2  if x + y > 15 then {
3      x = x + 7;
4      y = y - 12 }
5  else {
6      y = y + 10;
7      x = x - 2 };
8
9  x = x + 2;
10
11 if 2 * (x + y) > 21 then {
12 `    x = x * 3;
13     y = y * 2 }
14 else {
15     x = x * 4;
16     y = y * 3 + x };
17 skip
```

path4:

path4: 1, 2, 3.4.9, 11, 14 15 16 17 ✓

Edge	Symbolic State (PV)	Path Condition (PC)
$1 \to 2$	$x \mapsto X_0, y \mapsto Y_0$	true
2->3	$X \mapsto X+7$, $Y \mapsto Y$	$X_0 + Y_0 > 15$
3->4	$X \mapsto X_0+7$, $Y \mapsto Y_0 12$	$X_0 + Y_0 > 15$
4->9	$X \mapsto X_0+9$ $Y \mapsto Y_0 12$	$X_0 + Y_0 > 15$
9->11	$X \mapsto X_0+9$ $Y \mapsto Y_0 12$	$X_0 + Y_0 > 15 \land 2(X_0 + Y_0 - 3) <= 21$
11->14		
14->15		
15->16		
16->17		

15.1 - 3 = 12.1

$\to$ path infeasible

c): path 4 infeasible

path 3: X=0, Y=0

path 2: X=4, Y=10

path 1 : X=100 Y=10|

Question 2 (15 points)

(a) The constraint *at-most-one*$(a_1, \ldots, a_n)$ is satisfied if at most one of the Boolean variables $a_1, \ldots, a_n$ is true. For example, *at-most-one*$(\top, \bot, \bot)$ is true, and *at-most-one*$(\top, \bot, \top)$ is false. Encode the constraint

$$at\text{-}most\text{-}one(a_1, a_2, a_3, a_4)$$

into an equivalent set of clauses (i.e., in CNF).

(b) Show whether the following First Order Logic (FOL) sentence is valid or not. Either give a proof of validity, or show a model in which the sentence is false.

$$(\forall x \cdot \exists y \cdot P(x) \vee Q(y)) \iff (\forall x \cdot P(x)) \vee (\exists y \cdot Q(y))$$

(c) Show whether the following First Order Logic (FOL) sentence is valid or not. Either give a proof of validity, or show a model in which the sentence is false.
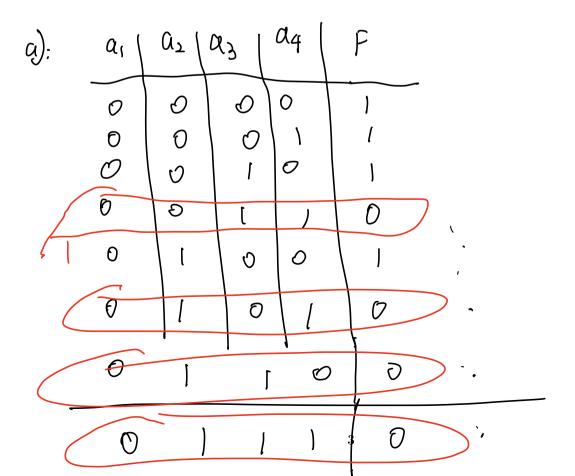
$$(\forall x \cdot \exists y \cdot P(x, y) \vee Q(x, y)) \implies (\forall x \cdot \exists y \cdot P(x, y)) \vee (\forall x \cdot \exists y \cdot Q(x, y))$$
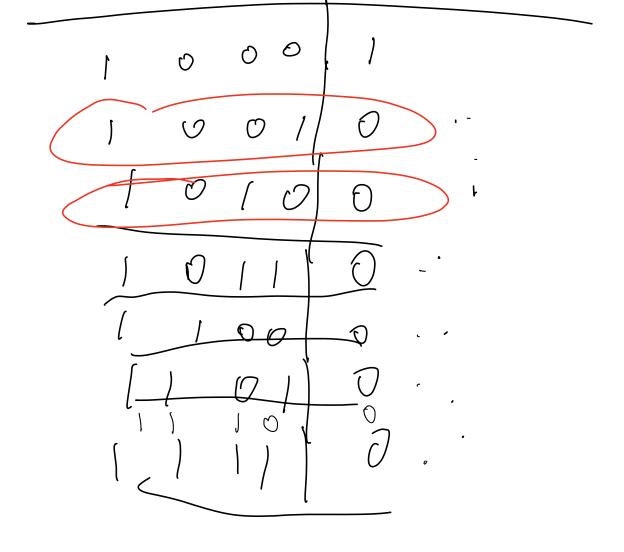
(d) Consider the following FOL formula Φ:

$$\exists x \exists y \exists z \, (P(x, y) \wedge P(z, y) \wedge P(x, z) \wedge \neg P(z, x))$$

For each of the following FOL models, explain whether they satisfy or violate the formula Φ.

   (a) $M = \langle S_1, P_1 \rangle$, where $S_1 = \mathbb{N}$, and $P_1 = \{(x, y) \mid x, y \in \mathbb{N} \wedge x < y\}$. Does $M_1 \models \Phi$?
   (b) $M_2 = \langle S_2, P_2 \rangle$, where $S_2 = \mathbb{N}$ and $P_2 = \{(x, x+1) \mid x \in \mathbb{N}\}$. Does $M_2 \models \Phi$?
   (c) $M_3 = \langle S_3, P_3 \rangle$, where $S_3 = \mathcal{P}(\mathbb{N})$, the powerset of natural numbers, and $P_3 = \{(A, B) \mid A, B \in \mathcal{P}(\mathbb{N}) \wedge A \subseteq B\}$. Does $M_3 \models \Phi$?

(e) *(Bonus question)* Extend your encoding from part (a) to n variables and use at most $O(n)$ clauses and variables. If your solution is based on external resources, make sure to properly reference them.

a):

a_1	a_2	a_3	a_4	F
0	0	0	0	1
0	0	0	1	1
0	0	1	0	1
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	1
0	1	1	1	0

$$(a_1 \lor a_2 \lor \lnot a_3 \lor \lnot a_4) \land (a_1 \lor \lnot a_2 \lor a_3 \lor \lnot a_4) \land (a_1 \lor \lnot a_2 \lor \lnot a_3 \lor a_4) \land$$

$$(a_1 \lor \lnot a_2 \lor \lnot a_3 \lor \lnot a_4) \land (\lnot a_1 \lor a_2 \lor a_3 \lor \lnot a_4) \land (\lnot a_1 \lor a_2 \lor \lnot a_3 \lor a_4) \land$$

$$(\lnot a_1 \lor a_2 \lor \lnot a_3 \lor \lnot a_4) \land (\lnot a_1 \lor \lnot a_2 \lor a_3 \lor a_4) \land (\lnot a_1 \lor \lnot a_2 \lor a_3 \lor \lnot a_4) \land (\lnot a_1 \lor \lnot a_2 \lor \lnot a_3 \lor a_4)$$

$$(\lnot a_1 \lor \lnot a_2 \lor \lnot a_3 \lor \lnot a_4)$$

b):

$$(\forall x. \exists y . P(x) \lor Q(y))) \Rightarrow (\forall x \cdot P(x)) \lor (\exists y \cdot Q(y))$$

①   ②

↳ For every x there exists some y s.t. $P(x)$ or $Q(y)$ is true

proove ①→②

→ if $P(x)$ is true then ② is true because

                               $(\forall x. P(x))$ is true

→ if $P(x)$ is false, then $Q(y)$ must be true

           So ② is true b.c. $(\exists y \cdot Q(y))$ is true

②→①

if $\forall x \cdot P(x)$ is true then ① is true

     b.c. if for every x $P(x)$ is true

     then $\forall x \cdot \exists y \cdot P(x) \lor Q(y)$ is true because $P(x)$ is true

if $\forall x \cdot P(x)$ is false but $\exists y \cdot Q(y)$ is true

then. there exists some y_0 such that $Q(y_0) =$ true

There fore $\forall x. \exists y \cdot P(x) \lor Q(y)$ is true because there

exists some y s.t. $Q(y) =$ true   So $\forall x \cdot \exists y \cdot P(x) \lor Q(y)$

                                      is true

c)

$S = \{a, b\}$

$P(a, a) = $ false $Q(b, b) = $ false $M(P) = \{\langle a, b\rangle\}$
$P(a, b) = $ true $Q(b, a) = $ true $M(Q) = \{\langle a, b\rangle$
$P(b, a) = $ false $Q(a, b) = $ true
$P(b, b) = $ false $Q(a, a) = $ false $\langle b, a\rangle\}$

RHS

$(\forall x \cdot \exists y \cdot P(x, y) \vee Q(x, y))$

is true b.c. when $x = a$, $P(a, b) \vee Q(a, b)$ is true
when $x = b$, $P(b, b) \vee Q(b, a)$ is true

LHS

$(\forall x \cdot \exists y \cdot P(x, y)) \vee (\forall x \cdot \exists y \cdot Q(x, y))$

$\forall x \cdot \exists y \, P(x, y)$ is always false, because $P(a, a)$, $P(b, a)$ $P(b, b)$ is false

$\forall x \cdot \exists y \cdot Q(x, y)$ is also false because $Q(b, b) = $ false
and $Q(a, a) = $ false

Therefore $(\forall x \cdot \exists y \cdot P(x, y)) \vee (\forall x \cdot \exists y \cdot Q(x, y))$ is false
The FOL sentence is false

(d) Consider the following FOL formula Φ:

$$\exists x \exists y \exists z \, (P(x, y) \wedge P(z, y) \wedge P(x, z) \wedge \neg P(z, x))$$

For each of the following FOL models, explain whether they satisfy or violate the formula Φ.

(a) $M_1 = \langle S_1, P_1\rangle$, where $S_1 = \mathbb{N}$ and $P_1 = \{(x, y) \mid x, y \in \mathbb{N} \wedge x < y\}$. Does $M_1 \models \Phi$?

(b) $M_2 = \langle S_2, P_2\rangle$, where $S_2 = \mathbb{N}$ and $P_2 = \{(x, x+1) \mid x \in \mathbb{N}\}$. Does $M_2 \models \Phi$?

(c) $M_3 = \langle S_3, P_3\rangle$, where $S_3 = \mathcal{P}(\mathbb{N})$, the powerset of natural numbers, and $P_3 = \{(A, B) \mid A, B \in \mathcal{P}(\mathbb{N}) \wedge A \subseteq B\}$. Does $M_3 \models \Phi$?

$\exists x \exists y \ \exists z \ (P(x,y) \wedge P(z,y) \wedge P(x,z) \wedge \neg P(z,x))$

there exists x, y, z st. $P(x,y) = $ true and $P(z,y) = $ true

$P(x,z) = $ true

and $(P(z,x) = $ false

a): $x=1$

$y=4$

$z=3$

$P(x,y) = $ true

$P(z,y) = $ true

$P(x,z) = $ true

$P(z,x) = $ false

$\neg P(z,x) = $ true

b): $P_b = \{(x, x+1) \mid x \in \mathbb{N}\}$

for $P(x,y) = $ true $\qquad y = x+1$

for $P(z,y) = $ true, $\qquad y = z+1$

for $P(x,z) = $ true $\qquad z$ must be $x+1$

however $x+1 = z+1$

$x = z$

So $P(x,z)$ can't be true

$N_b \not\models \Phi$

c): $A = \{1\}$ $\qquad$ $P(x,y) = $ true, Let $x = A$, $y = B$

$B = \{1, 4, 5\}$ $\qquad$ $P(z,y) = $ true , Let $z = C$, $y = B$

$C = \{1, 4\}$ $\qquad$ $P(x,z) = $ true if $x = A$, $z = C$

$P(z,x) = $ false if $x = A$, $z = C$ $\qquad$ $C \not\subseteq A$

2e): Use auxillary variable b

where if a_1 is true b is true b_2
b = true all $a_2 a_3$ is false

$$\bigwedge_{i=1}^{n-1} (\neg a_i \vee b) \wedge (\neg b \wedge \neg a_{i+1})$$
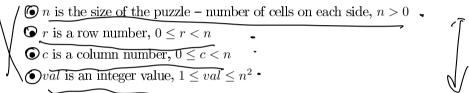
Question 3 (15 points)

In recreational mathematics, a **magic square** is a $n \times n$ square grid filled with distinct positive integers from 1 to n^2 inclusive such that each cell contains a different integer, and the sum of the integers in each row, column, and diagonal is equal.

The following is an example of 3×3 magic square:

(a) Write down quantifier free constraints in First Order Logic to solve the puzzle above for any positive integer n.

(b) Use Z3 Python API to implement a solver for magic square puzzles. Your solver should accept four parameters: n, r, c, val, where:

- n is the size of the puzzle – number of cells on each side, $n > 0$
- r is a row number, $0 \leq r < n$
- c is a column number, $0 \leq c < n$
- val is an integer value, $1 \leq val \leq n^2$

Your program should find a magic square of the given size n with the value val filled at location (r, c), assuming the top left corner corresponds to $(0, 0)$.

Your solver should return a 2D array of integers corresponding to the solution. If there is more than one solution, just return any one of them. If there is no such magic square, your solver should return None. For example,

```
>>> res = solve_magic_square(3, 1, 1, 5)
>>> print_square(res, 3)
4 9 2 =15
3 5 7 =15
8 1 6
```

The skeleton for the solver is provided in `a2q3/magic_square.py`.

You might find it helpful to use Z3 `z3.Distinct(x)` to create a constraint that states that all constants in the list x have distinct values. For example,

```
>>> x, y = z3.Ints ('x y')
>>> z3.Distinct (x, y)
x != y
```

(c) Extend the test suite in `a2q3/puzzle_tests.py` with two additional set of parameters, and one extra set of parameters that does not have a solution (i.e., the solver returns None).

Recall that you can execute the test suite using the following command:

```
python -m a2q3.test
```

a):

each cell number range

$$\forall_{i,j} \left({1 \leq i \leq n \, \wedge \atop 1 \leq j \leq n} \right) \rightarrow \left(1 \leq grid[i][j] \leq n^2 \right)$$

distinct

$$\forall_{i,j,p,q} \left({\wedge \; 1 \leq j \leq n \wedge \; 1 \leq p \leq n \atop \wedge \; 1 \leq q \leq n \; \wedge \; p \neq i \; \wedge \; q \neq j} \right) \quad grid[i][j] \neq grid[k][q]$$

Sum for each row, column, diagonal is the same

$$\forall_i \quad {}_{1 \leq i \leq n} \left(\sum_{j=1}^{n} grid[i][j] \right) = M$$

$$\forall_j \quad {}_{1 \leq j \leq n} \quad \sum_{i=1}^{n} grid[i][j] \right) = M$$

$$\forall_i \quad {}_{1 \leq i \leq n} \quad \sum_{i=1}^{n} grid[i][i] = M$$

$$\forall_i \quad {}_{1 \leq i \leq n} \quad \sum_{i=1}^{n} grid[i][n-i+1] = M$$

$$\text{where } M = \frac{n(n^2+1)}{2}$$

$$M = \frac{\frac{n^2 \times (n^2+1)}{2}}{n} = \frac{n \times (n^2+1)}{2}$$

Question 4 (30 points)

Your GitHub repository includes an implementation of a parser and interpreter for the WHILE language from the lecture notes. Your task is to write a symbolic execution engine for it.

The implementation of the interpreter is located in directory `wlang`. You can execute the interpreter using the following command:

```
(venv) $ python3 -m wlang.int wlang/test1.prg
x: 10
```

A sample program is provided for your convinience in `wlang/test1.prg`

You can execute the interpreter using the following command:

```
(venv) $ python3 -m wlang.sym wlang/test1.prg
```

$$\frac{n^3 + n^2 + 2n + 2}{2}$$

$$2+2+2\left(\frac{6}{3}\right)$$
$$1+1+2+2$$
$$1000 + 100 + 20 + 2 = \frac{1122}{2} =$$

A skeleton for a symbolic interpreter is given in `wlang/sym.py`. It includes an implementation of a symbolic state in a class `SymState`. The class is provided for your convenience. You are free to modify it in any way or create your own.

You may find it helpful to look at the implementation of the concrete interpreter in `wlang/int.py`. Note that the concrete interpreter takes a `State` as input, and returns a single `State` as output. On the other hand, symbolic interpreter takes a symbolic state `State` as input, and returns a **list** of symbolic states as output, where each output state corresponds to some execution of the program. In general, the number of output states will be proportional to the number of execution paths in the program.

(a) Implement symbolic execution of straight line code (i.e., programs without `if`- and `while`-statements);

(b) Extend your answer to symbolic execution of programs with `if`-statements;

(c) Extend your answer to symbolic execution of programs with `while`-statements. To handle arbitrary loops, assume that the loop is executed at most **10** times. That is, your symbolic execution engine should explore all feasible program paths in which the body of each loop is executed no more than **10** times.

(d) Extend the test suite `test_sym.py` to achieve 100% branch coverage of your implementation in parts (a), (b), and (c). Recall that you can run the test suite using

```
(venv) $ python3 -m wlang.test
```

and measure coverage of the test suite using

```
(venv) $ coverage run -m wlang.test
(venv) $ coverage html
```

(e) Provide a program on which your symbolic execution engine diverges (i.e., takes longer than a few seconds to run).

e)

```python
        self.assertEquals(len(out),1)
def test_14(self):
    prg1 = 'havoc x; havoc y; havoc z; while x>100 do { while y> 100 do { while z >100 do {z:=z-1}; y:=y-1}; x:=x-1}'
    ast1 = ast.parse_string(prg1)
    engine = sym.SymExec()
    st = sym.SymState()
    out = [s for s in engine.run(ast1, st)]
    self.assertEquals(len(out),1111)
```