

# Breadth-first search

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## Erdős Number - the geek version of Bacon number :)



- Paul Erdős: famous Hungarian mathematician;
- He worked with hundreds of colleagues;
- Published more than 1.400 articles;
- Erdős number is an amused attribute created by his friends;
- Paul Erdős has Erdős number 0;
- Direct collaborators have number 1;
- Collaborators of collaborators have number 2 and so on;

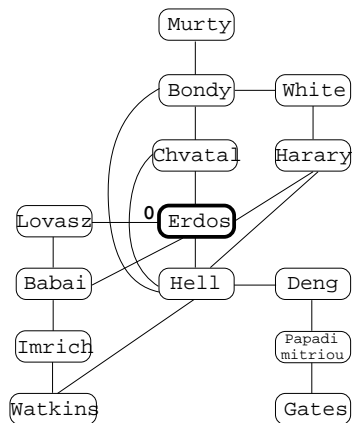
## Definition

- Given a list of people and the relation of collaboration between them, which is Erdős number of each person?

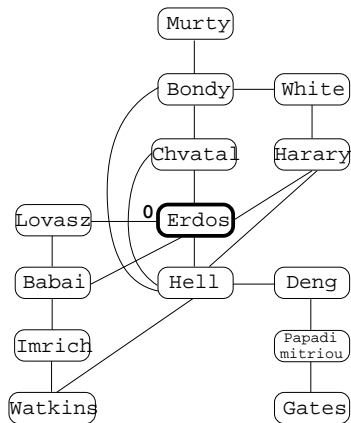
## Definition

- Given a list of people and the relation of collaboration between them, which is Erdős number of each person?
- This problem can be modeled as a graph where:
  - People are vertices;
  - Collaborations are edges.

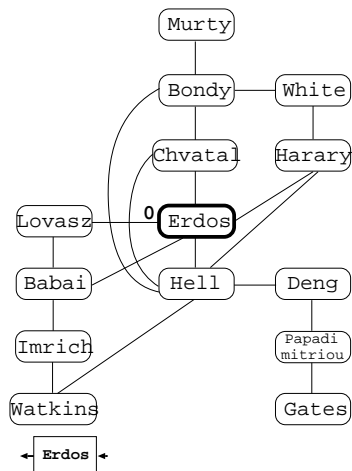
## Example



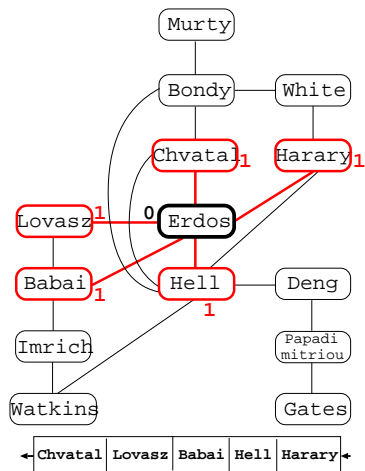
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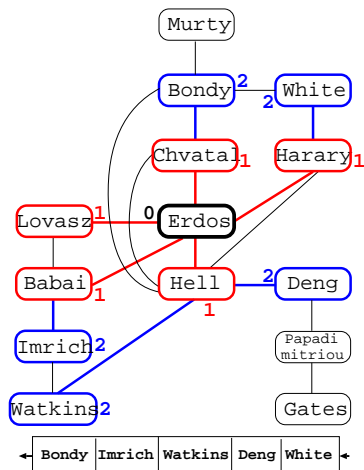


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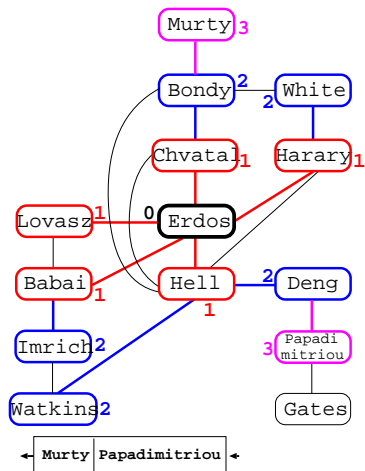




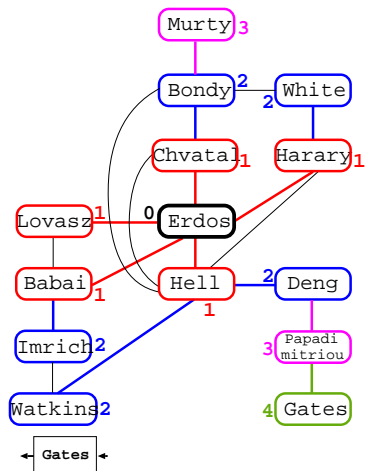
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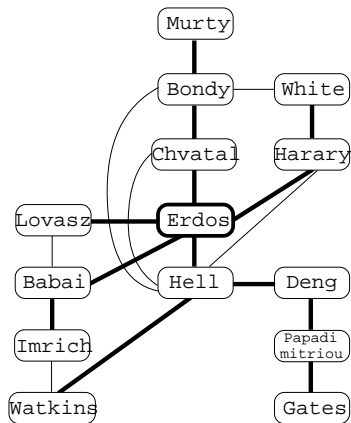
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## Algorithm BFS

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1  bfs(G, s):
2    for u in V(G) do
3        u.visited = False
4        u.d =  $\infty$ 
5        u.p = None
6    s.visited = True
7    s.d = 0
8    Q = Queue()
9    insert(Q, s)
10   while length(Q) > 0 do
11       u = remove(Q)
12       for v in adj(u) do
13           if not v.visited then
14               v.d = u.d + 1
15               v.p = u
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- Each vertex is added once in the queue (line 17):  $O(n)$ ;
- Adjacency list of each vertex is traversed once (line 12):  $O(m)$ ;
- Total complexity:  $O(n + m)$ .

## Bibliography

CORMEN, T. H.; LEISERSON, C. E.; RIVEST, R. L. and STEIN, C.  
*Introduction to Algorithms*, 3rd edition, MIT Press, 2009.