

Fully Verifying Transformation Contracts for Declarative ATL

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Business Informatics Group

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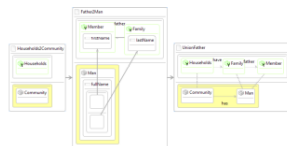
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 - Input independence - verification for all input models
 - Examine combinations of transformation rules

- Translating ATL transformation into DSLTrans language



Translation (ATL HOT)

DSLTrans

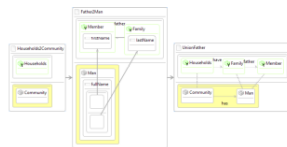


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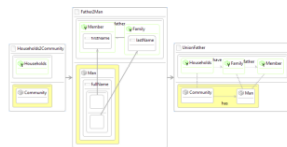


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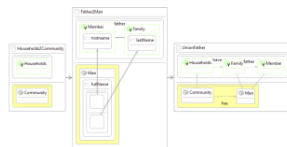
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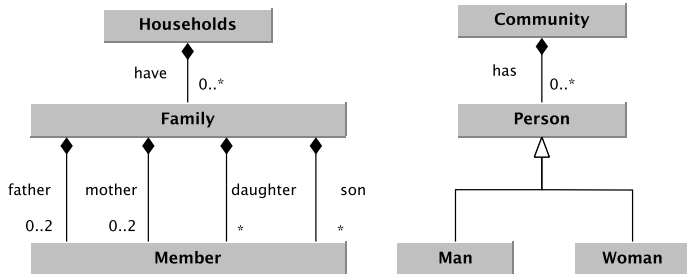
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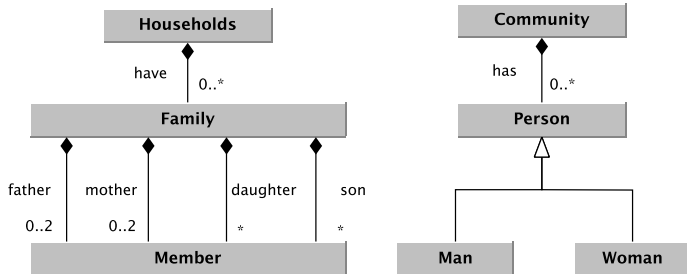
- Performed through a higher-order transformation
 - Specified in ATL

Transformation Metamodels



- Transform *Members* to *Men* and *Women*

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- Transform *Members* to *Men* and *Women*
- NB: Metamodels are not representative of today's society!

ATL Transformation

```
1  module Families2Persons;
2  create OUT : Persons from IN : Families;
3
4
5
6
7
8
9
10
11
12
13
14
15 rule Father2Man { -- R2
16   from
17     mem : Families!Member, fam : Families!Family
18       (fam.father=mem)
19   to
20     m : Persons!Man (
21       fullName <- mem.firstName + fam.lastName --B2
22     )}
```

ATL Transformation

```
1 module Families2Persons;
2 create OUT : Persons from IN : Families;
3
4 rule Households2Community { -- R1
5   from
6     hh: Families!Households
7   to
8     c : Persons!Community (
9       has <- hh.have->collect(f | thisModule.
10         resolveTemp(Tuple{mem=f.father,fam=f}, 'm')), --B11
11       has <- hh.have->collect(f | thisModule.
12         resolveTemp(Tuple{mem=f.mother,fam=f}, 'w')) --B12
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- Implicit resolution mechanism of ATL
- Through collect operation

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- Graph-based, contains rules arranged in layers

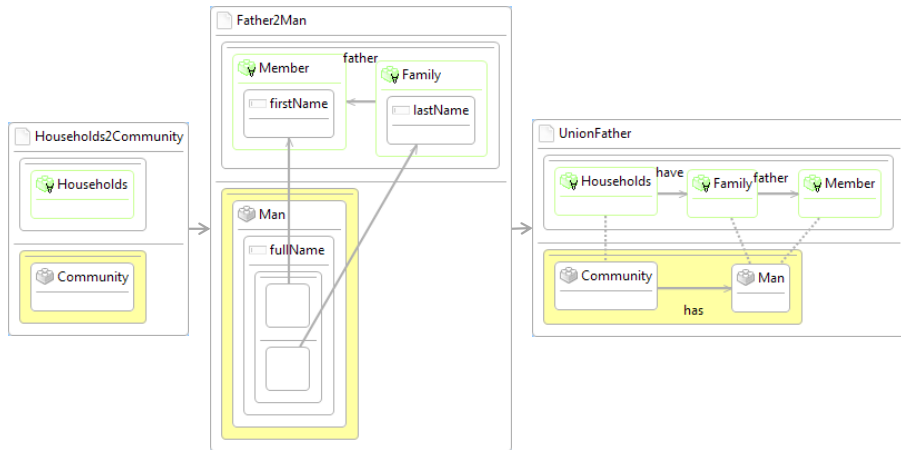
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- All DSLTrans computations are terminating and confluent
- Unbounded loops during execution are not allowed



- Rules arranged in layers
- Match graph on top of rules
- Apply graph on bottom
 - Produced when match graph is found

- Higher-order transformation written in ATL
- Creates a DSLTrans transformation from declarative ATL
 - Informal testing: less than 20 seconds
- Available on our website: <http://msdl.cs.mcgill.ca/people/levi/files/MODELS2015>

TABLE I
FEATURES OF DECLARATIVE ATL CONSIDERED

Matched Rules	✓	Filters	✓
Lazy Rules	✓	OCL Expressions	✓
Several Bindings	✓	Helpers	×
Several <i>InPatternElements</i>	✓	Conditions	×
Several <i>OutPatternElements</i>	✓	Using Block	×

- Covers declarative ATL
- Transformation can be rewritten to avoid missing features

- Two steps for higher-order transformation

Mapping - Part Three

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- First, each from/to part of an ATL rule is transformed into match/apply graphs in DSLTrans

Mapping - Part Three

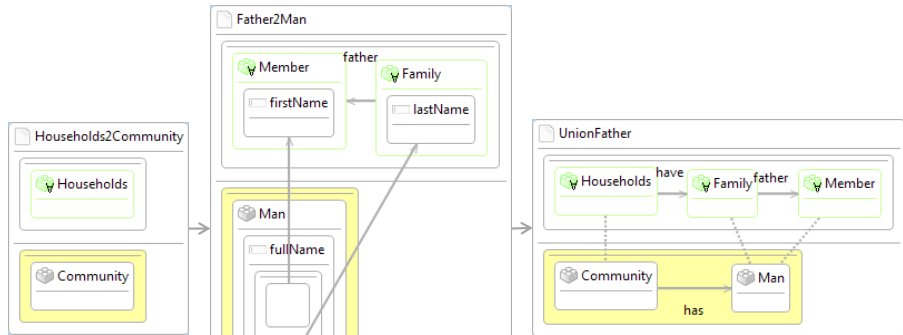
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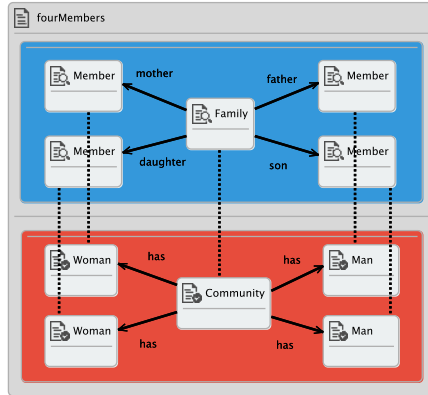
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Mapping - Part Four

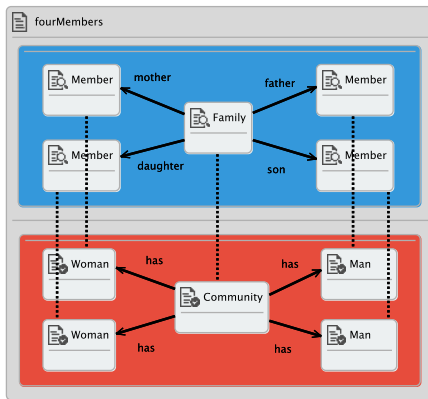
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Contracts

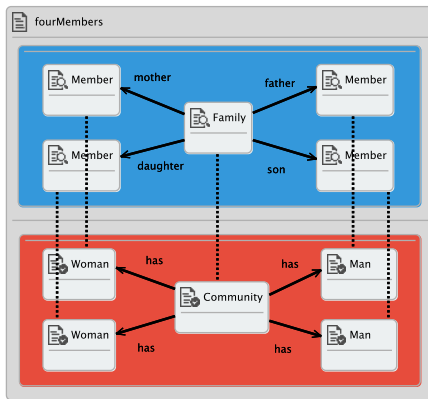


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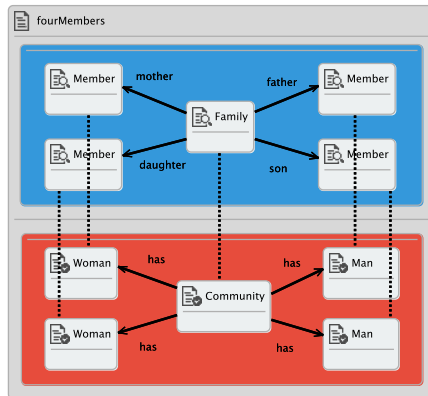
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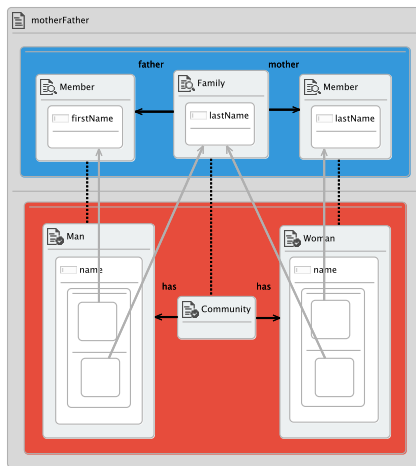
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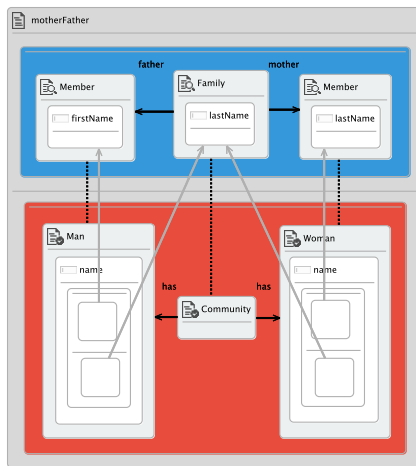


- If blue graph is in input model, then red graph is in output model
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- *A family with a father, mother, son, daughter should always produce two males and two females in the target community*

Contracts

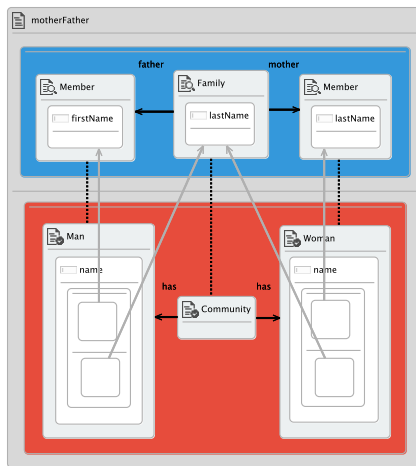


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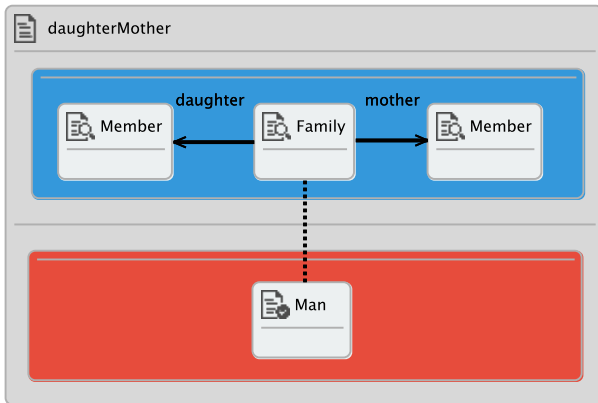
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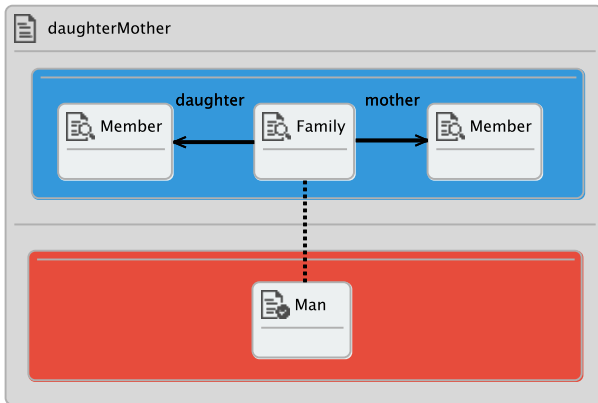
- Reasoning about attributes of elements
- *Is the full name of the produced Person correctly created from the last name of the Family and the first name of the Member?*

Contracts



- A contract that will not hold

Contracts



- A contract that will not hold
- *A family with a mother and a daughter will always produce a community with a man*

- SyVOLT contract proving tool

L. Lúcio, B. Oakes, and H. Vangheluwe. A technique for symbolically verifying properties of graph-based model transformations. Tech. Report SOCS-TR-2014.1, McGill U, 2014.

Levi Lucio et al. SyVOLT: Full Model Transformation Verification Using Contracts

Contract Proving - Part One

- SyVOLT contract proving tool
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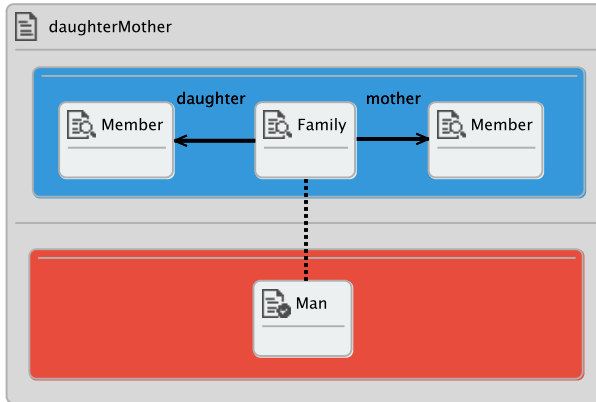
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- A contract holds for a transformation if it holds for all generated path conditions

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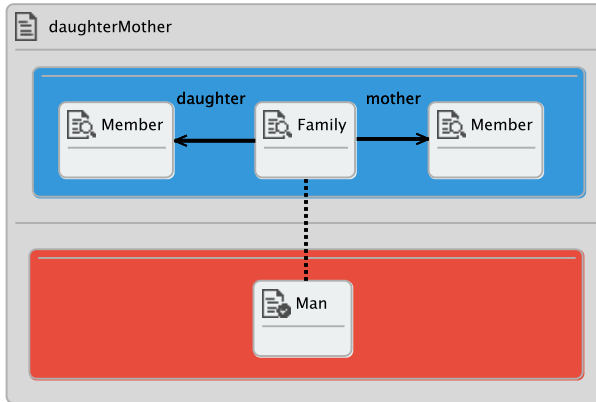
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Contract Proving - Part Two



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Contract Proving - Part Two



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- Fails on path condition:
'HEmpty_HRoot_HMotherRule_HDaughterRule'

Experiments Conducted

- Applicability of the Technique
- Time and Memory Characteristics
- Reducing Contract Proving Time
- Higher-Order Transformation

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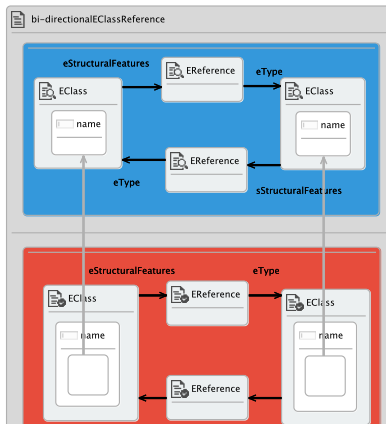
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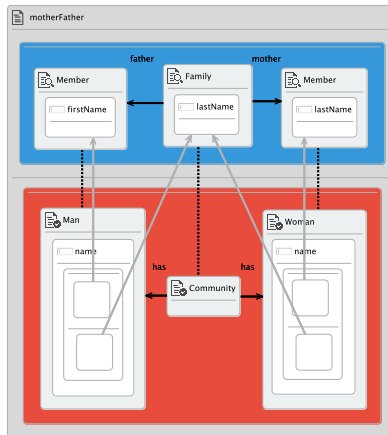


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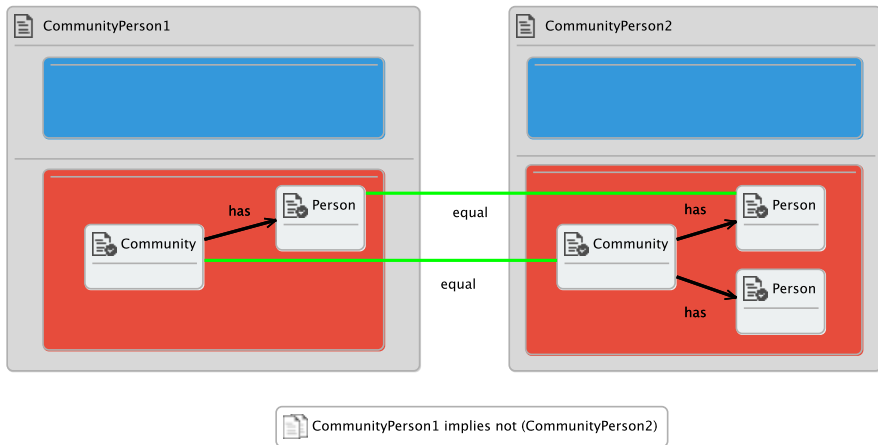
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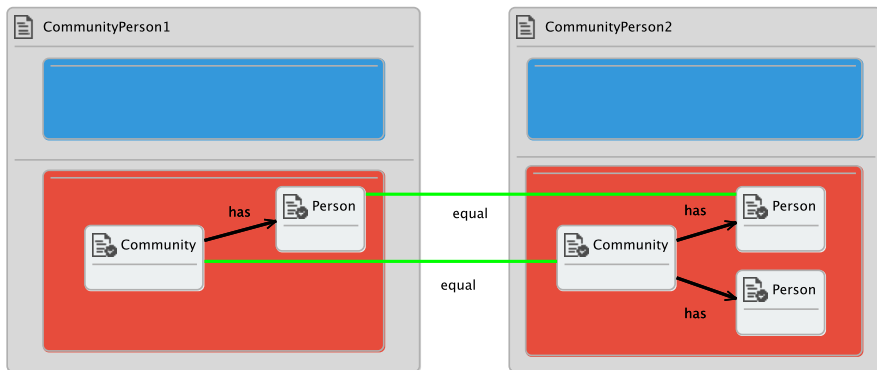
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Applicability of the Technique - Part Three



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 CommunityPerson1 implies not (CommunityPerson2)

- *'If a Community is connected to a Person element, that Community is connected to one and only one Person element'*
- Selim, Gehan. Formal Verification of Graph-Based Model Transformations. PhD Diss. Queen's University, 2015.

Time and Memory Characteristics

	ATL/ DSLTrans Rules	Path Conds. Gen.	Time (s)	Contracts Proved	Time (s)	Memory (MB)
Families-to-Person	5 / 9	52	1.54	4	31.45	45
ER-Copier	5 / 9	70	0.48	1	1.70	43
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Reducing Contract Proving Time

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 - Recursively create rule dependency tree
- Manually performed - slicing has since been automated

- Question: Is a transformation produced by a HOT equivalent to a hand-built one?

G. M. Selim, L. Lúcio, J. R. Cordy, J. Dingel, and B. J. Oakes. Specification and verification of graph-based model transformation properties. In Graph Transformation, pages 113–129. Springer, 2014.

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- Note that number of rules/transformation shape not optimized
- But HOT produces roughly equivalent result

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 - Investigate contract-based transformation development

- Developed higher-order transformation to transform ATL into DSLTrans
- Can verify visual contracts on DSLTrans transformations in feasible time
- Contracts verified on all transformation executions
- Future work
 - Integrate HOT into SyVOLT tool
 - Investigate contract-based transformation development
- Thank you for your time!

Fully Verifying Transformation Contracts for Declarative ATL

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<http://msdl.cs.mcgill.ca/people/levi/files/MODELS2015>

Multiplicity Contract

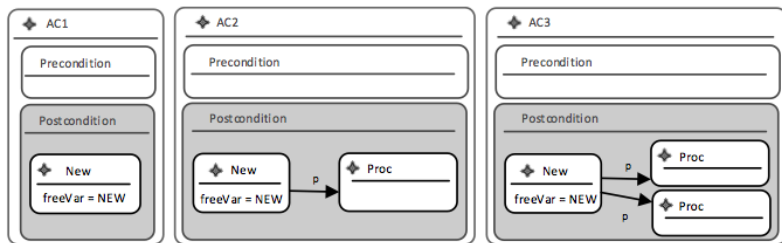


Figure C.1: *AtomicContracts* $AC1$, $AC2$, and $AC3$ that are used to express $MM1$ (Table 6.4) as $AC1 \Rightarrow_{tc} (AC2 \wedge_{tc} \neg_{tc} AC3)$.

“Multiplicity Invariants ensure that the transformation does not produce an output that violates the multiplicities in the Kiltera metamodel”

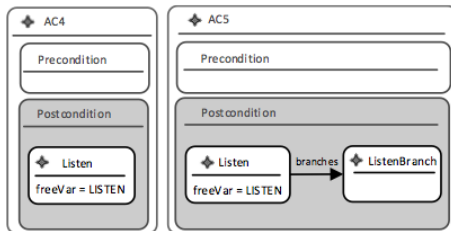


Figure C.5: *AtomicContracts* AC_4 and AC_5 that are used to express MM_5 (Table 6.4) as $AC_4 \Rightarrow_{tc} AC_5$.

“Syntactic Invariants ensure that the generated Kiltera output model is well-formed with respect to Kiltera’s syntax.”

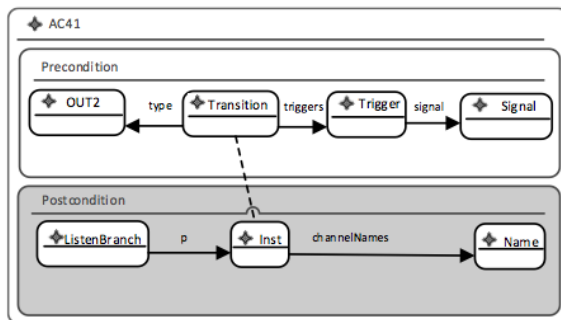


Figure C.16: *AtomicContract AC41* that is used to express *PP2* (Table 6.4).

“Pattern contracts require that if a certain pattern of elements exists in the input model, then a corresponding pattern of elements exists in the output model”