

# 1 Section 8.6 Exercises

Exercises with solutions from Section 8.6 of [UA].

**Exercise 8.6.1.** (a) Fix  $r \in \mathbf{Q}$ . Show that the set  $C_r = \{t \in \mathbf{Q} : t < r\}$  is a cut.

The temptation to think of all cuts as being of this form should be avoided. Which of the following subsets of  $\mathbf{Q}$  are cuts?

(b)  $S = \{t \in \mathbf{Q} : t \leq 2\}$

(c)  $T = \{t \in \mathbf{Q} : t^2 < 2 \text{ or } t < 0\}$

(d)  $U = \{t \in \mathbf{Q} : t^2 \leq 2 \text{ or } t < 0\}$

*Solution.* (a) It is clear that  $C_r$  satisfies (c1) and (c2). To see that  $C_r$  satisfies (c3), observe that if  $t \in C_r$ , then  $t < \frac{t+r}{2}$  and  $\frac{t+r}{2} \in C_r$ .

(b) This is not a cut, since it has 2 as a maximum element.

(c) This is a cut.  $T$  satisfies (c1) since  $0 \in T$  and  $2 \notin T$ .

Suppose  $t \in T$  and  $r$  is a rational such that  $r < t$ . If  $r < 0$  then certainly  $r \in T$ , so suppose that  $r \geq 0$ , which implies that  $t > 0$ . It follows that  $r^2 < t^2 < 2$  and so  $r \in T$ . Thus  $T$  satisfies (c2).

Suppose  $t \in T$ . If  $t \leq 0$  then let  $r = 1$  and if  $t > 0$  then let  $r = \frac{2t+2}{t+2}$ . In either case, one can verify that  $t < r$  and  $r \in T$ . Thus  $T$  satisfies (c3).

(d) By Theorem 1.1.1 we have  $U = T$  and hence by part (c)  $U$  is a cut.

**Exercise 8.6.2.** Let  $A$  be a cut. Show that if  $r \in A$  and  $s \notin A$ , then  $r < s$ .

*Solution.* Given that  $r \in A$ , the implication  $s \notin A \implies r < s$  is the contrapositive of (c2).

**Exercise 8.6.3.** Using the usual definitions of addition and multiplication, determine which of these properties are possessed by  $\mathbf{N}$ ,  $\mathbf{Z}$ , and  $\mathbf{Q}$ , respectively.

*Solution.*  $\mathbf{N}$  satisfies (f1), (f2), and (f5). It fails (f3) since there is no additive identity and it fails (f4) since no element has an additive inverse and only 1 has a multiplicative inverse (1 is its own inverse).

$\mathbf{Z}$  satisfies (f1), (f2), (f3), and (f5). It fails (f4) since, while each element has an additive inverse, only 1 and  $-1$  have multiplicative inverses (they are their own inverses).

$\mathbf{Q}$  satisfies each property (f1) - (f5).

**Exercise 8.6.4.** Show that this defines an ordering on  $\mathbf{R}$  by verifying properties (o1), (o2), and (o3) from Definition 8.6.5.

*Solution.* Properties (o2) and (o3) are clear, so let us verify property (o1). It will suffice to show that if  $B \not\subseteq A$ , then  $A \subseteq B$ . Since  $B$  is not a subset of  $A$ , there exists some  $s \in B$  such that  $s \notin A$ . Let  $r \in A$  be given. By [Exercise 8.6.2](#) we must have  $r < s$  and so by (c2) we have  $r \in B$ . Thus  $A \subseteq B$ .

**Exercise 8.6.5.** (a) Show that (c1) and (c3) also hold for  $A + B$ . Conclude that  $A + B$  is a cut.

(b) Check that addition in  $\mathbf{R}$  is commutative (f1) and associative (f2).

(c) Show that property (o4) holds.

(d) Show that the cut

$$O = \{p \in \mathbf{Q} : p < 0\}$$

successfully plays the role of the additive identity (f3). (Showing  $A + O = A$  amounts to proving that these two sets are the same. The standard way to prove such a thing is to show two inclusions:  $A + O \subseteq A$  and  $A \subseteq A + O$ .)

*Solution.* (a) Since  $A$  and  $B$  are non-empty,  $A + B$  must also be non-empty. Since neither  $A$  nor  $B$  contains every rational number, there exist rationals  $r \notin A$  and  $s \notin B$ . It follows from [Exercise 8.6.2](#) that  $a + b < r + s$  for every  $a \in A$  and  $b \in B$ , so that  $r + s \notin A + B$ . Thus  $A + B \neq \mathbf{Q}$  and we have now shown that  $A + B$  satisfies (c1).

Let  $a + b \in A + B$  be given. By (c3), there exist rationals  $r \in A$  and  $s \in B$  such that  $a < r$  and  $b < s$ . It follows that  $a + b < r + s$  and  $r + s \in A + B$ . Thus  $A + B$  satisfies (c3).

(b) Commutativity and associativity of addition in  $\mathbf{R}$  follow immediately from commutativity and associativity of addition in  $\mathbf{Q}$ .

(c) Let  $A, B$ , and  $C$  be cuts such that  $B \subseteq C$ . If  $a + b \in A + B$ , then  $a + b \in A + C$  also since  $B \subseteq C$ . Thus  $A + B \subseteq A + C$ .

(d) Let  $a + p \in A + O$  be given. Then  $p < 0$ , so  $a + p < a$  and it follows from (c2) that  $a + p \in A$ ; thus  $A + O \subseteq A$ .

Now let  $a \in A$  be given. By (c3) there exists some  $b \in A$  such that  $a < b$ . Notice that  $a = b + (a - b) \in A + O$ , since  $a - b < 0$ . It follows that  $A \subseteq A + O$  and we may conclude that  $A + O = A$ .

**Exercise 8.6.6.** (a) Prove that  $-A$  defines a cut.

- (b) What goes wrong if we set  $-A = \{r \in \mathbf{Q} : -r \notin A\}$ ?
- (c) If  $a \in A$  and  $r \in -A$ , show  $a + r \in O$ . This shows  $A + (-A) \subseteq O$ . Now, finish the proof of property (f4) for addition in Definition 8.6.4.

*Solution.* (a) Since  $A \neq \mathbf{Q}$ , there exists a  $t \notin A$ . Then  $-t-1 \in -A$ , since  $t < -(-t-1) = t+1$ . Thus  $-A$  is non-empty. Since  $A$  is non-empty, there exists some  $r \in A$ . Then  $-r \notin -A$ , since if  $t \notin A$  then  $t > -(-r) = r$  by [Exercise 8.6.2](#). Thus  $-A \neq \mathbf{Q}$  and we see that  $-A$  satisfies (c1).

Suppose that  $r \in -A$ , so that there is some  $t \notin A$  such that  $t < -r$ , and suppose that  $s < r$ . Then  $t < -r < -s$ , demonstrating that  $s \in -A$  also. Thus  $-A$  satisfies (c2).

Suppose that  $r \in -A$ , so that there is some  $t \notin A$  such that  $t < -r$ . Define  $s = r - \frac{r+t}{2}$  and notice that  $r < s$  since  $0 < -r - t$ . Furthermore,  $s \in -A$  since

$$t \notin A \quad \text{and} \quad t < \frac{t-r}{2} = -s.$$

Thus  $-A$  satisfies (c3) and we may conclude that  $-A$  is a cut.

- (b) This does not necessarily define a cut. For example, let  $C_2$  be the cut  $\{r \in \mathbf{Q} : r < 2\}$ . Then using this definition, we find that  $-C_2 = \{r \in \mathbf{Q} : r \leq -2\}$ , which fails property (c3).
- (c) There exists a  $t \notin A$  such that  $t < -r$ . By [Exercise 8.6.2](#) it must be the case that  $a < t < -r$  and thus  $a + r < 0$ , i.e.  $a + r \in O$ . Thus  $A + (-A) \subseteq O$ .

For the reverse inclusion, let  $p < 0$  be a given rational number in  $O$ . We claim that there must exist some  $r \in A$  such that  $r - \frac{p}{2} \notin A$ , and we will prove this by contradiction. Suppose that  $r - \frac{p}{2} \in A$  for all  $r \in A$ . Since  $A$  is a cut, there is some  $r_0 \in A$ . An induction argument shows that  $r_0 - \frac{np}{2} \in A$  for all  $n \in \mathbf{N}$ . Let  $q \in \mathbf{Q}$  be given and use the Archimedean property of  $\mathbf{Q}$  to obtain an  $n \in \mathbf{N}$  such that  $r_0 - \frac{np}{2} > q$ ; it follows from (c2) that  $q \in A$ . The conclusion is that  $A = \mathbf{Q}$ , contradicting (c1).

Thus there is some  $r \in A$  such that  $r - \frac{p}{2} \notin A$ . Since  $r - \frac{p}{2} < r - p$ , it follows that  $p - r \in -A$ . Then  $p = r + p - r \in A + (-A)$ , demonstrating that  $O \subseteq A + (-A)$ . We may conclude that  $A + (-A) = O$ .

**Exercise 8.6.7.** (a) Show that  $AB$  is a cut and that property (o5) holds.

- (b) Propose a good candidate for the multiplicative identity (1) on  $\mathbf{R}$  and show that this works for all cuts  $A \geq O$ .
- (c) Show the distributive property (f5) holds for non-negative cuts.

*Solution.* (a) It is clear that  $AB$  is non-empty. If either  $A = O$  or  $B = O$ , then it is straightforward to verify that  $AB = O \neq \mathbf{Q}$ . Suppose that  $A > O$  and  $B > O$ . There exist rationals  $r \notin A$  and  $s \notin B$ ; clearly,  $r, s > 0$ . If  $q \in AB$ , then either  $q < 0$  or  $q = ab$  for  $a \in A, b \in B$  and  $a, b \geq 0$ . By Exercise 8.6.2 we must have  $a < r$  and  $b < s$ , so that  $ab < rs$ . In either case, we have  $q < rs$  and thus  $rs \notin AB$ , demonstrating that  $AB \neq \mathbf{Q}$ . Thus  $AB$  satisfies (c1).

Suppose  $r \in AB$  and  $q < r$ . If  $q < 0$  then  $q \in AB$ , so suppose that  $q \geq 0$ , which implies that  $r > 0$ . We must then have  $r = ab$  for some  $a \in A, b \in B$  with  $a, b > 0$ . Notice that  $\frac{q}{b} < a$ ; (c2) then implies that  $\frac{q}{b} \in A$  and hence  $q = \frac{q}{b} \cdot b \in AB$ . Thus  $AB$  satisfies (c2).

If  $A = O$  or  $B = O$  then  $AB = O$ , which has no maximum element. Suppose that  $A > O$  and  $B > O$  and let  $r \in AB$  be given. If  $r \leq 0$  then let  $q$  be any positive rational in  $AB$ . If  $r > 0$ , then  $r = ab$  for some  $a \in A, b \in B$  with  $a, b > 0$ . By (c3), there exist rationals  $s \in A, t \in B$  such that  $a < s$  and  $b < t$ . Let  $q = st$  and notice that  $q \in AB$  and  $r = ab < st = q$ . In either case, there exists a  $q \in AB$  with  $r < q$ . Thus  $AB$  satisfies (c3) and we may conclude that  $AB$  is a cut.

Property (o5) is clear from the definition of  $AB$ .

- (b) Define  $I = \{p \in \mathbf{Q} : p < 1\}$  and let  $A \geq O$  be given. We claim that  $AI = A$ . Suppose that  $r \in AI$ . If  $r < 0$ , then  $r \in A$ , so suppose that  $r \geq 0$ . Thus  $r = ab$  for some  $a \in A$  such that  $a \geq 0$  and some  $0 \leq b < 1$ . It follows that  $ab < a$  and so by (c2) we have  $r = ab \in A$ . Thus  $AI \subseteq A$ .

Now suppose that  $a \in A$ . If  $a \leq 0$ , then (c2) implies that  $2a \in A$  and thus  $a = (2a) \cdot \frac{1}{2} \in AI$ . If  $a > 0$ , then (c3) implies there is some  $r \in A$  with  $a < r$ . Thus  $\frac{a}{r} \in I$  and we see that  $a = r \cdot \frac{a}{r} \in AI$ . Hence  $A \subseteq AI$  and we may conclude that  $AI = A$ .

- (c) Let  $A, B, C \geq O$  be cuts. If  $ABC = O$  then the equality  $A(B + C) = AB + AC$  is clear, so suppose that  $A, B, C > O$  and suppose that  $q \in A(B + C)$ . If  $q < 0$  then  $q = \frac{q}{2} + \frac{q}{2} \in AB + AC$ . Suppose that  $q \geq 0$ . Then  $q = a(b + c) = ab + ac$ , where  $a \in A, b \in B, c \in C$  and  $a, b + c \geq 0$ . There are three cases:  $b, c \geq 0$ ,  $b \geq 0$  and  $c < 0$ , or  $b < 0$  and  $c \geq 0$ . In any of these cases, it is straightforward to verify that  $ab + ac \in AB + AC$ . Thus  $A(B + C) \subseteq AB + AC$ .

Now suppose that  $p + q \in AB + AC$ . If  $p + q < 0$ , then  $p + q \in A(B + C)$ , so suppose that  $p + q \geq 0$ . If  $p, q \geq 0$ , then  $p = a_1b$  and  $q = a_2c$ , for some  $a_1, a_2 \in A, b \in B$ , and  $c \in C$  such that  $a_1, a_2, b, c \geq 0$ . Let  $a = \max\{a_1, a_2\}$  and notice that  $a(b + c) \in A(B + C)$ . Furthermore,  $p + q = a_1b + a_2c \leq ab + ac = a(b + c)$ . It follows from (c2) that  $p + q \in A(B + C)$ .

Next, suppose that  $p < 0$  and  $q \geq 0$ , so that  $q = ac$  for some  $a \in A, c \in C$  with  $a, c \geq 0$ .

Let  $b \in B$  be such that  $b \geq 0$ ; such a  $b$  exists since  $B > O$ . Now notice that

$$p + q = p + ac < ac \leq a(b + c) \in A(B + C).$$

It follows from (c2) that  $p + q \in A(B + C)$ . The case where  $p \geq 0$  and  $q < 0$  is handled similarly. Thus  $AB + AC \subseteq A(B + C)$  and we may conclude that  $A(B + C) = AB + AC$ .

**Exercise 8.6.8.** Let  $\mathcal{A} \subseteq \mathbf{R}$  be nonempty and bounded above, and let  $S$  be the *union* of all  $A \in \mathcal{A}$ .

- (a) First, prove that  $S \in \mathbf{R}$  by showing that it is a cut.
- (b) Now, show that  $S$  is the least upper bound for  $\mathcal{A}$ .

*Solution.* (a) Since  $\mathcal{A}$  is non-empty, it contains some cut  $A$ , so that  $A \subseteq S$ . It follows that  $S$  is non-empty as  $A$  is non-empty. Since  $\mathcal{A}$  is bounded above, there exists some cut  $B$  such that  $A \subseteq B$  for all  $A \in \mathcal{A}$ . It follows that  $S \subseteq B$  and hence that  $S \neq \mathbf{Q}$  since  $B \neq \mathbf{Q}$ . Thus  $S$  satisfies (c1).

Suppose  $r \in S$ , so that  $r \in A$  for some  $A \in \mathcal{A}$ , and suppose  $q < r$ . Since  $A$  is a cut we must have  $q \in A$ , which gives  $q \in S$ . Thus  $S$  satisfies (c2).

Suppose  $r \in S$ , so that  $r \in A$  for some  $A \in \mathcal{A}$ . Since  $A$  is a cut there exists some  $q \in A$  such that  $r < q$ ; note that  $q \in S$  also. Thus  $S$  satisfies (c3). We may conclude that  $S$  is a cut.

- (b) It is clear that  $S$  is an upper bound for  $\mathcal{A}$ . If  $B$  is any upper bound for  $\mathcal{A}$ , then  $B$  contains every  $A \in \mathcal{A}$  and hence must contain the union of all  $A \in \mathcal{A}$ , i.e.  $S \subseteq B$ . It follows that  $S$  is the least upper bound for  $\mathcal{A}$ .

**Exercise 8.6.9.** Consider the collection of so-called “rational” cuts of the form

$$C_r = \{t \in \mathbf{Q} : t < r\}$$

where  $r \in \mathbf{Q}$ . (See [Exercise 8.6.1](#).)

- (a) Show that  $C_r + C_s = C_{r+s}$  for all  $r, s \in \mathbf{Q}$ . Verify  $C_r C_s = C_{rs}$  for the case when  $r, s \geq 0$ .
- (b) Show that  $C_r \leq C_s$  if and only if  $r \leq s$  in  $\mathbf{Q}$ .

*Solution.* (a) Let  $r, s \in \mathbf{Q}$  be given and suppose  $a + b \in C_r + C_s$ , i.e.  $a < r$  and  $b < s$ . It follows that  $a + b < r + s$  and hence that  $a + b \in C_{r+s}$ . Thus  $C_r + C_s \subseteq C_{r+s}$ . Now suppose that  $t \in C_{r+s}$ , so that  $t < r + s$ . Choose a positive integer  $n \in \mathbf{N}$  such that  $t + \frac{1}{n} < r + s$  and note that:

- $s - \frac{1}{n} < s$ , so that  $s - \frac{1}{n} \in C_s$ ;
- $t + \frac{1}{n} - s < r$ , so that  $t + \frac{1}{n} - s \in C_r$ ;
- $t = \left(t + \frac{1}{n} - s\right) + \left(s - \frac{1}{n}\right) \in C_r + C_s$ .

Thus  $C_{r+s} \subseteq C_r + C_s$  and we may conclude that  $C_r + C_s = C_{r+s}$ .

It is clear that  $C_r C_s = C_{rs}$  if  $rs = 0$ , so suppose that  $r, s > 0$  and let  $q \in C_r C_s$  be given. If  $q \leq 0$  then  $q < rs$ , i.e.  $q \in C_{rs}$ . If  $q > 0$  then  $q = ab$  for some  $0 < a < r$  and  $0 < b < s$ . It follows that  $0 < ab < rs$  and thus  $q = ab \in C_{rs}$ . Hence  $C_r C_s \subseteq C_{rs}$ .

Now let  $q \in C_{rs}$  be given. If  $q \leq 0$  then certainly  $q \in C_r C_s$ , so suppose that  $q > 0$  and define  $p = \frac{1}{2}\left(\frac{q}{s} + r\right)$ . Notice that:

- $0 < \frac{q}{s} < p < r$ , so that  $p \in C_r$ ;
- $0 < \frac{q}{p} < s$ , so that  $\frac{q}{p} \in C_s$ ;
- $q = p \cdot \frac{q}{p} \in C_r C_s$ .

Thus  $C_{rs} \subseteq C_r C_s$  and we may conclude that  $C_r C_s = C_{rs}$ .

- (b) If  $r \leq s$  then it is clear that  $C_r \subseteq C_s$ . If  $s < r$ , then it is again clear that  $C_s \subseteq C_r$ . Furthermore, notice that  $C_s \neq C_r$  since  $\frac{s+r}{2}$  belongs to  $C_r$  but not to  $C_s$ . Thus  $C_s \subsetneq C_r$ .