CS1382 Discrete Computational Structures

Lecture 15: Trees

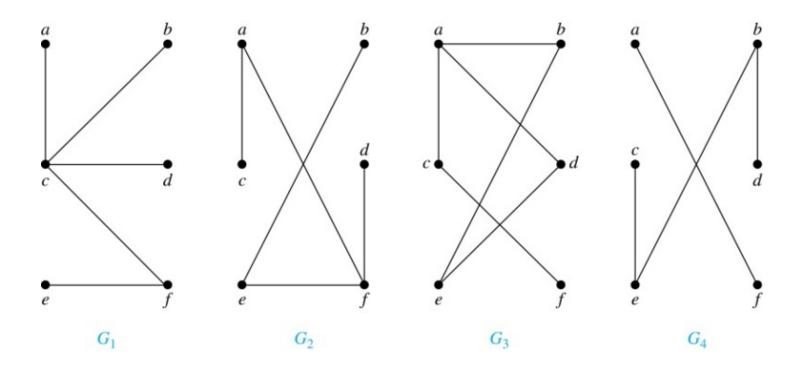
Spring 2019

Richard Matovu



Trees

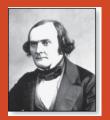
A *tree* is a connected undirected graph with no simple circuits.



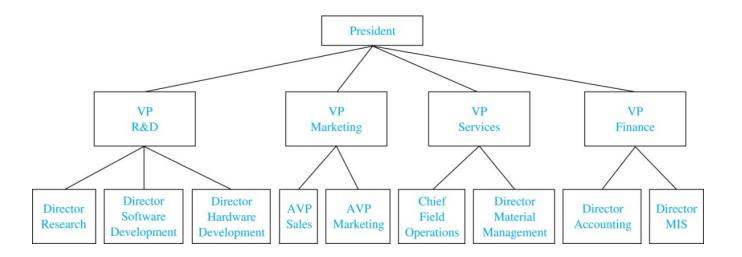
A *forest* is a graph that has no simple circuit, but is not connected.

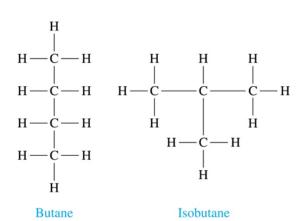
Each of the connected components in a forest is a tree.

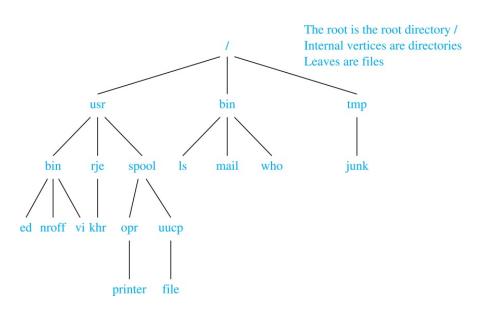
Trees as Models



Used as models in computer science, chemistry, geology, botany, psychology, and many other areas.



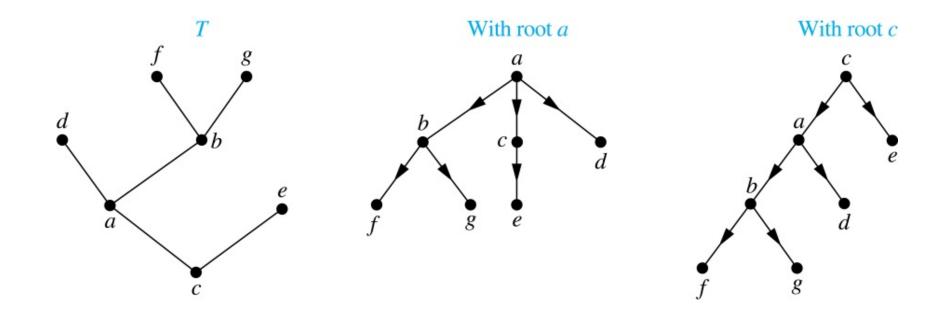




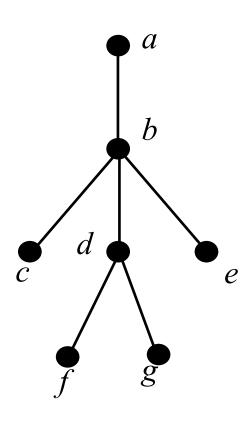
Rooted Trees

A **rooted tree** is a tree in which one vertex has been designated as the **root** and every edge is directed away from the root.

An **unrooted** tree is converted into different rooted trees when different vertices are chosen as the root.



Rooted Tree Terminology



a is the **parent** of b, b is the **child** of a,

c, d, e are siblings,

a, b, d are ancestors of f

c, d, e, f, g are descendants of b

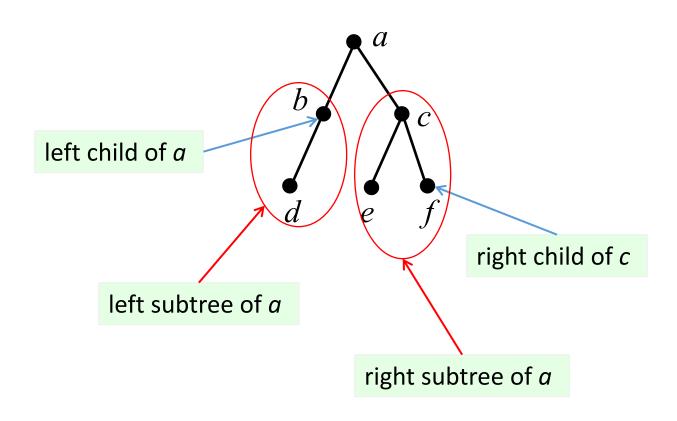
c, e, f, g are leaves of the tree (deg=1)

a, b, d are internal vertices of the tree (at least one child)

subtree with *d* as its root:



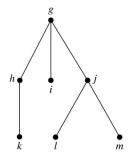
Rooted Tree Terminology

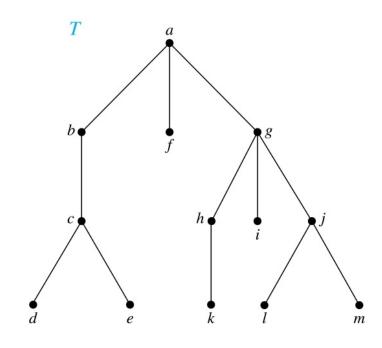


Terminology for Rooted Trees

In the rooted tree T (with root a):

- Find the parent of *c*, the children of *g*, the siblings of *h*, the ancestors of *e*, and the descendants of *b*.
- Find all internal vertices and all leaves.
- What is the subtree rooted at G?

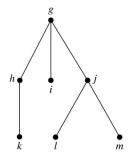


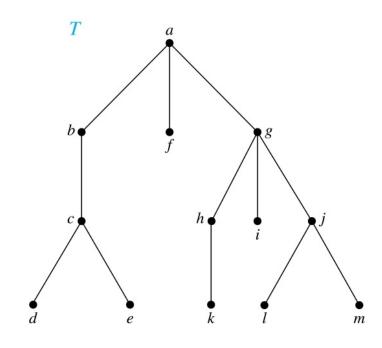


Terminology for Rooted Trees

In the rooted tree T (with root a):

- Find the parent of *c*, the children of *g*, the siblings of *h*, the ancestors of *e*, and the descendants of *b*.
- Find all internal vertices and all leaves.
- What is the subtree rooted at G?

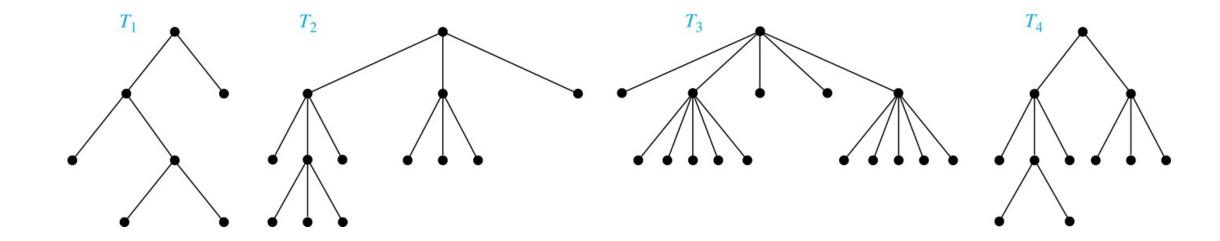




m-ary Rooted Trees

A rooted tree is called an *m-ary tree* if every internal vertex has no more than *m* children.

• An m-ary tree with m = 2 is called a *binary* tree.



Properties of Trees

- An undirected graph is a tree if and only if there is a unique simple path between any two of its vertices.
- A tree with *n* vertices has *n* − 1 edges.

CS1382 Discrete Computational Structures

Tree Traversals

Spring 2019

Richard Matovu



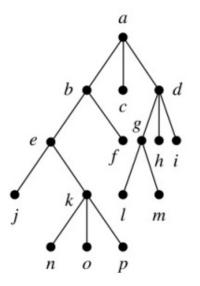
Tree Traversal

Procedures for systematically visiting every vertex of an ordered tree are called *traversals*.

- Preorder traversal
- Inorder traversal
- Postorder traversal.

Preorder Traversal

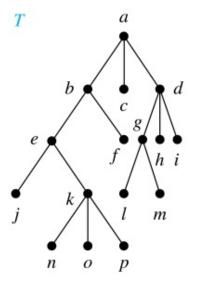
```
procedure preorder (T: ordered rooted tree)
r := root of T
list r
for each child c of r from left to right
    T(c) := subtree with c as root
    preorder(T(c))
```



Preorder traversal: Visit root, visit subtrees left to right

Inorder Traversal (continued)

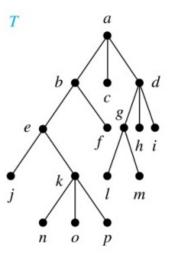
```
procedure inorder (T: ordered rooted tree)
r := \text{root of } T
if r is a leaf then list r
else
   l := first child of r from left to right
  T(I) := subtree with I as its root
  inorder(T(I))
  list(r)
  for each child c of r from left to right
     T(c) := subtree with c as root
     inorder(T(c))
```



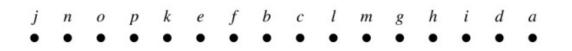
Inorder traversal: Visit leftmost subtree, visit root, visit other subtrees left to right

Postorder Traversal (continued)

```
procedure postordered (T: ordered rooted tree)
r := root of T
for each child c of r from left to right
    T(c) := subtree with c as root
    postorder(T(c))
list r
```



Postorder traversal: Visit subtrees left to right; visit root



CS1382 Discrete Computational Structures

Spanning Trees

Spring 2019

Richard Matovu



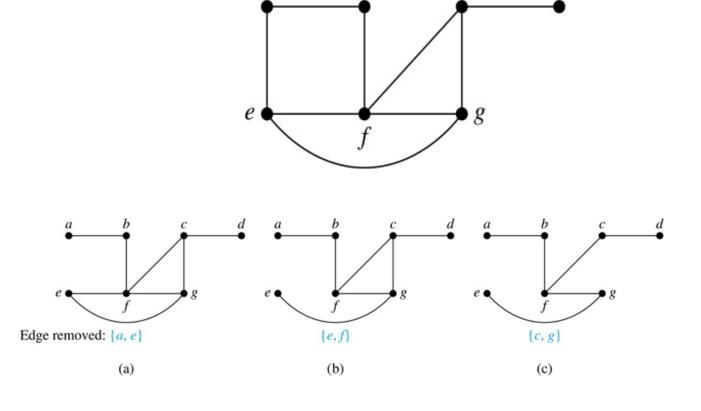
Spanning Trees

A **spanning tree** of simple graph *G* is a subgraph of *G* that is a tree containing every vertex of *G*.

Applications:

- Communication/PowerGrid Networks
- Data Clustering
- Maze Generation

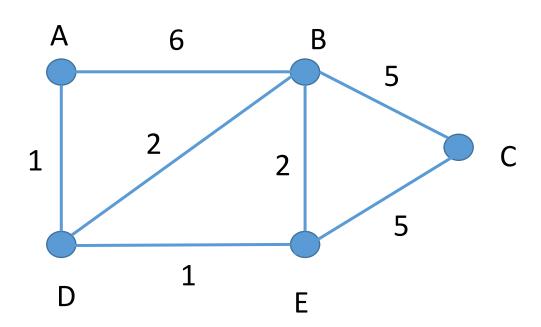
Find the spanning tree of the simple graph below:

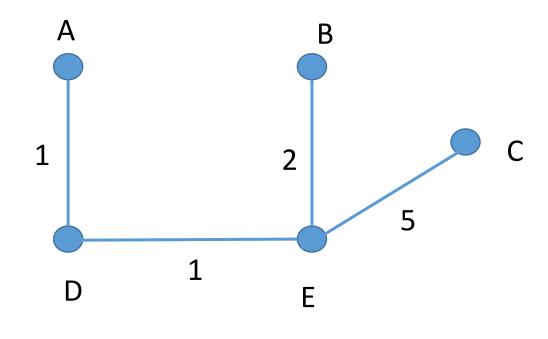


Minimum Spanning Trees (MST)

MST is a spanning tree whose sum of the weights of edges is minimum

- Kruskal's Algorithm
- Prim's Algorithm





Questions?

Thank You!