

AVL

Definition

An AVL tree is a binary search tree which satisfies:
the heights of the two child sub trees of any node differ by at most one

Remark:

Representation stores the balance factor or the height of the node

Operations over AVL

- search, insert and delete
 - all take $O(\log n)$ time in average and worst cases
 - where n is the number of nodes in the tree prior to the operation.

Consider the next representation:

```
AVLTreeNode =    record
                  info: TComparable
                  left: ^ AVLTreeNode
                  right: ^ AVLTreeNode
                  h: Integer
                end
```

Search

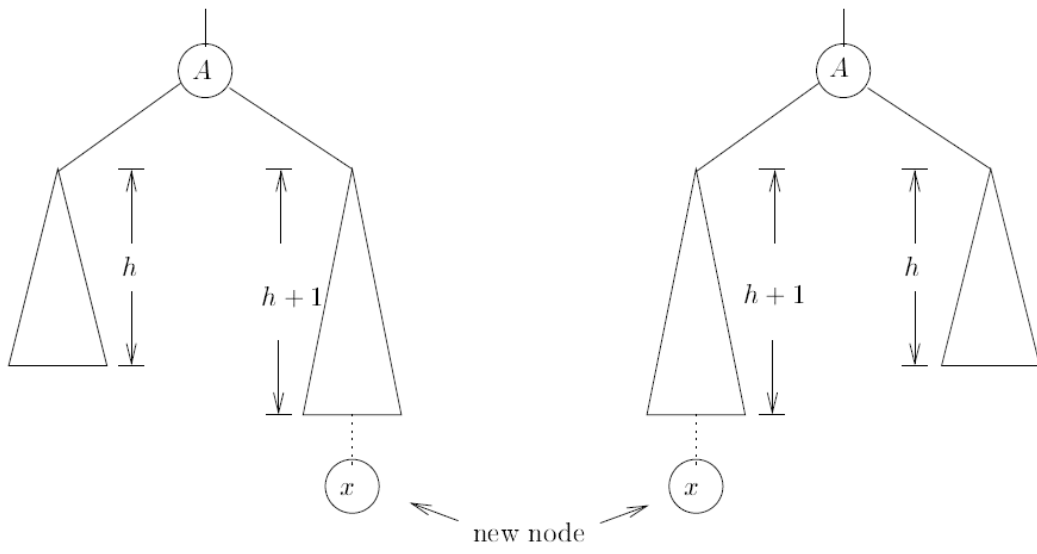
- BST search

Insert

Insertion:

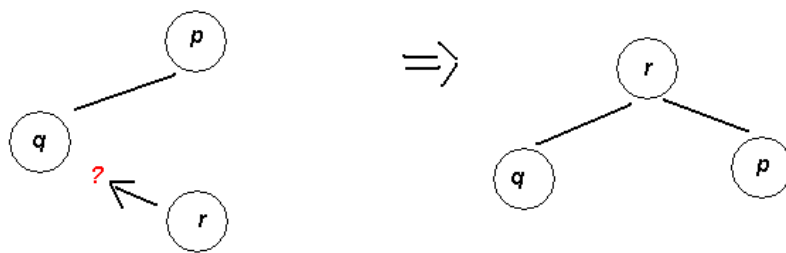
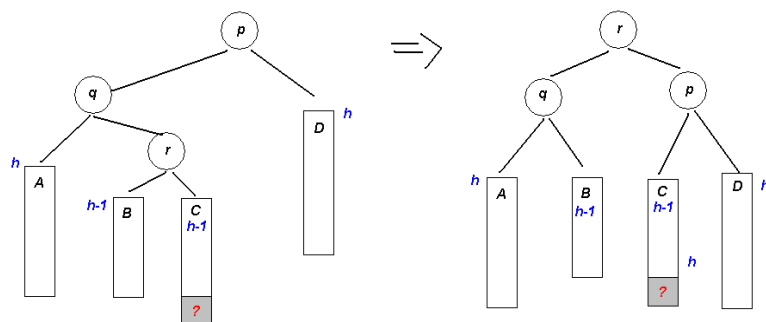
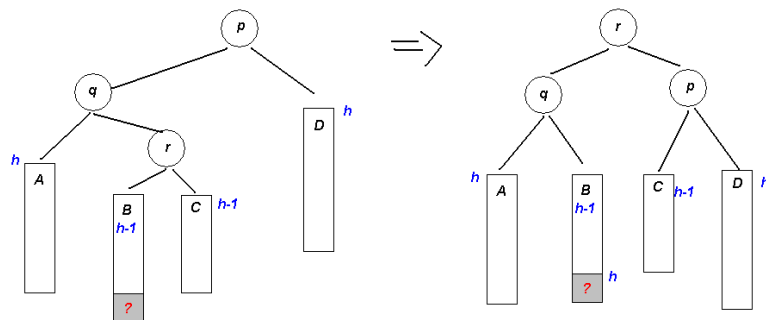
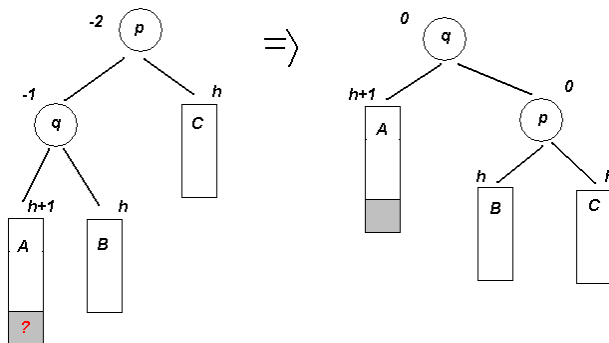
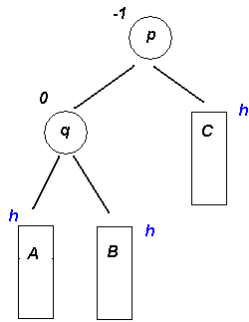
- require the tree to be rebalanced
 - insert an element like in BST case
 - rebalance the tree (if it is the case)
 - consider all the ancestors (to the root)
 - rebalance** → one or more tree rotations.

When to rebalance :

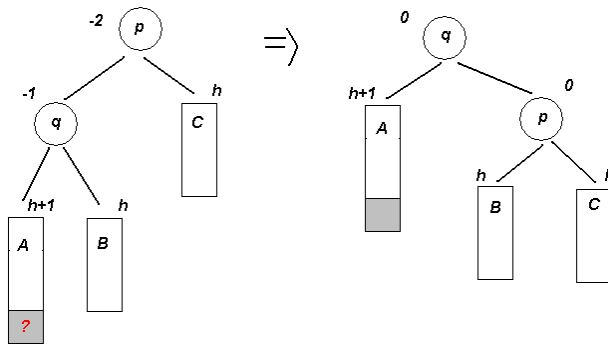


Insert cases - examples

Assume: (next) initial situation & new node on the left subtree



Rotations



/ Representation without link to parent */*

/ Update heights, then return new root */*

Function RotateRight (p)

 q := p^.left

 p^.left := q^.right;

 q^.right := p;

 p^.h := Max(Height(p^.left), Height(p^.right)) + 1;

 q^.h := Max(Height(q^.left), Height(q^.right)) + 1;

 RotateRight := q */* New root */*

end_RotateRight

Function RotateLeft (p)

 q := p^.right;

 p^.right := q^.left

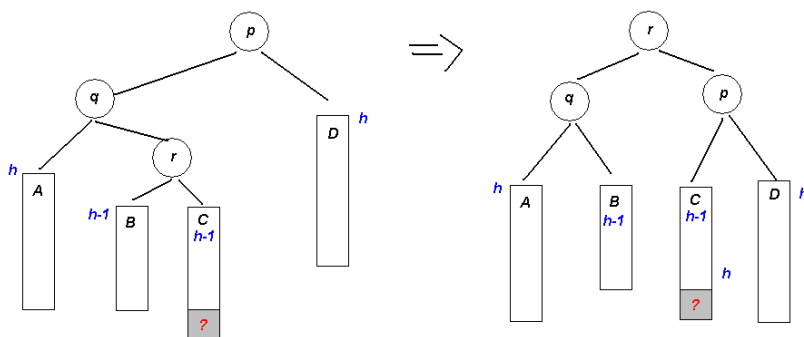
 q^.left := p

 p^.h := Max(Height(p^.left), Height(p^.right)) + 1

 q^.h := Max(Height(q^.left), Height(q^.right)) + 1

 RotateLeft := q

end_RotateLeft



Function DblRotateLeftRight (p)

 p^.left := RotateLeft (p^.left)

 DblRotateLeftRight := RotateRight (p);

end_DblRotateLeftRight

```

Function insert_rec(p , el)
// ElementType el, AvlTreeNode p
// return the new p
  if( p = NIL)
    p := new AvlTreeNode
    p^.info := el
    p^.h := 0;
    p^.left := NIL
    p^.right := NIL
  else
    if( el < p^.info) then
      p^.left := insert_rec(p^.left , el )
      if(Height(p^. right) - Height( p^. left ) = -2 )
        if( el < p^.left^.info)
          p := RotateRight ( p )
        else
          p := DblRotateLeftRight ( p )
        endif
      endif
    else
      // el >= [p].info
      p^.right = insert_rec(p^.right , el )
      if( Height( p^.right ) - Height( p^.left ) = 2 )
        if( el > p^.right^.info ) then
          p := RotateLeft ( p )
        else
          p := DblRotateRightLeft( p );
        endif
      endif
    endif
    p^.h := Max( Height( p^.left ), Height( p^.right ) ) + 1;
  endif
insert_rec := p
End_insert_rec

Subalg. insert(T , el)
  p := getRoot(T)
  np :=insert_rec(p, el)
  setRoot(T, np)
end_insert

```

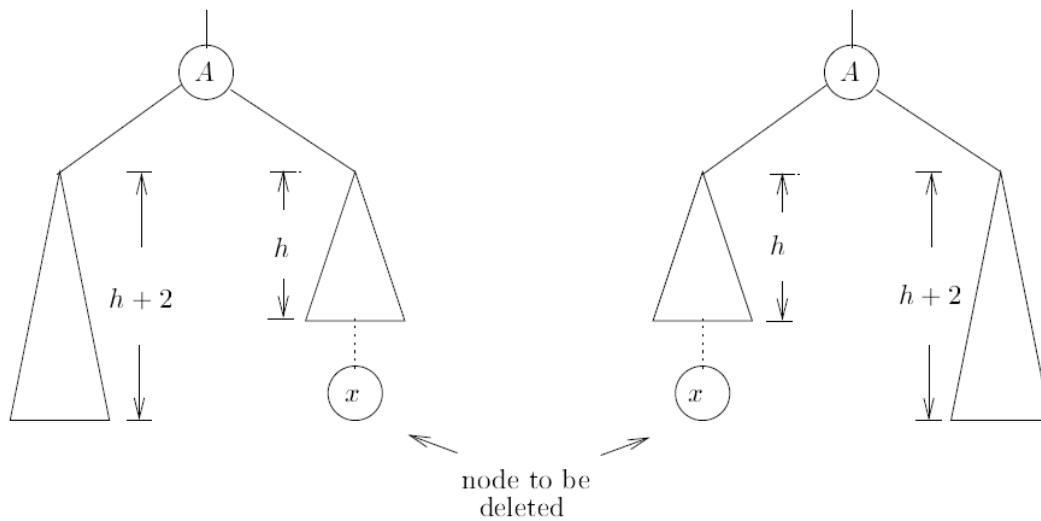
Delete

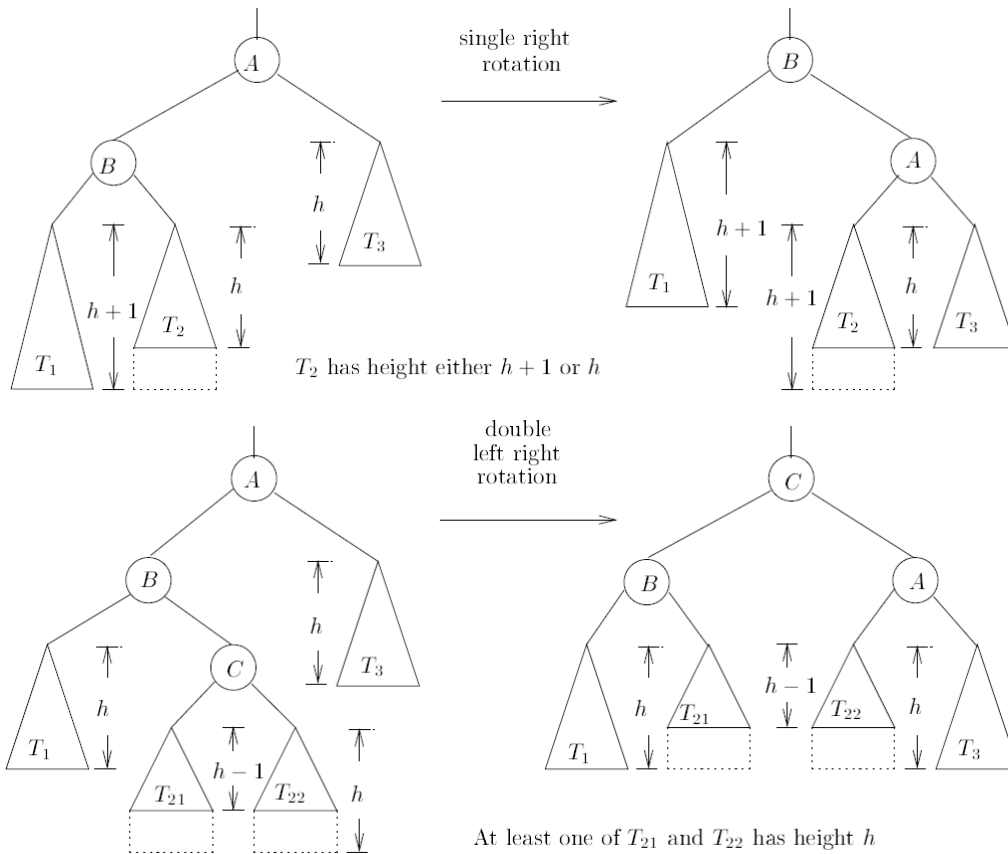
- find the node x where k is stored
- delete the contents of node x ~similar with BST
Deleting a node in an AVL tree can be reduced to deleting a leaf
(next) alg. delete_rec – delete leaves

rebalance

- go from the deleted leaf towards the root
- update the balance factor
- rebalance with rotations if necessary.

Rebalance cases





```

Function delete_rec (p , el)
if p = NIL then return NIL endif
if el = p^.info then
    if p^.left=NIL and p^.right=NIL then
        delete p; p:=NIL
    else
        q:=p
        if p^.left <> NIL      then  p:=p^.left
                               else  p:=p^.right
        endif
        delete q
    endif
else
    if( el > p^.info) then
        p^.right = delete_rec ( p^.right , el );
        if( Height( p^.right ) - Height( p^.left ) = -2 ) then
            if(Height( p^.left ^.left ) = Height( p^.right ) +1) then
                p = RotateRight ( p )
            else // Height( [[p].left ].left = Height( [p].right )
                p = DblRotateLeftRight(p)
            endif
        endif
    else // el( el < [p].info)
        //...
    endif
endif
if p<> NIL then p^.h = Max( Height( p^.left ), Height( p^.right ) ) + 1 endif

```

```
delete_rec := p  
End_delete_rec
```