Talisman

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1 代数

1.1 FFT-gwx

```
#include <complex>
    int n, m;
int rev[maxn];    //maxn > 2 ^ k
    void fft(Complex *a, int f)
 8
         for(int i = 0; i < m; i++)
         if(i < r[i]) swap(a[i], a[r[i]]);
for(int 1 = 2; 1 <= m; 1 <<= 1)
9
11
12
               int h = 1 >> 1;
               Complex ur = (Complex) {cos(pi / h), f * sin(pi / h)};
for(int i = 0; i < m; i += 1)</pre>
13
14
15
                    Complex w = (Complex) \{1, 0\};
for(int k = 0; k < h; k++, w = w * ur)
16
17
18
19
                          20
21
22
               }
23
24
         if(f == -1)
               for(int i = 0; i < m; i++)
a[i] = a[i] / m;
25
26
27
28
    void multi(Complex *a, Complex *b)
30
         fft(a, 1); fft(b, 1);
for(int i = 0; i < m; i++)
   a[i] *= b[i];</pre>
31
32
33
         fft(a, -1);
34
35
36
37
    void init()
38
         for(m = 1; m <= 2 * n; m <<= 1);
for(int i = 0, j = 0; i < m; i++)
39
40
41
42
43
               for(int x = m >> 1; (j = x) < x; x >>= 1);
44
45
```

1.2 FFT-wrz

```
int len:
   struct comp
 3
 4
5
        double r, i;
        comp operator + (const comp &that) {return (comp){r+that.r, i+that.i};}
comp operator - (const comp &that) {return (comp){r-that.r, i-that.i};}
comp operator * (const comp &that) {return (comp){r*that.r-i*that.i, r*that.i+i*that.r};}
   w[N<<1], a[N<<1], b[N<<1], c[N<<1]; // 数组记得至少开两倍void init()
 9
10
        11
12
13
   void FFT(comp *a, comp *w)
14
15
16
        for(int i = 0, j = 0; i < len; i++)
17
             if(i<j) swap(a[i], a[j]);
for(int 1 = len>>1; (j^=1)<1; 1 >>= 1);
18
19
20
21
        for(int i = 2; i <= len; i <<= 1)
22
23
             int m = i >> 1;
for(int j = 0; j < len; j += i)</pre>
24
25
26
27
                  for(int k = 0; k < m; k++)
28
                       29
30
                       a[j+k] = a[j+k] + tmp;
                  }
31
             }
32
33
   void mul(comp *a, comp *b, comp *c, int 1) // 多项式乘法, c = a * b , c的长度为1
35
36
        for(len = 1; len <= 1; len <<= 1);
init(); FFT(a, w); FFT(b, w);
for(int i = 0; i < len; i++) c[i] = a[i] * b[i];</pre>
37
38
39
        reverse(c+1, c+len); FFT(c, w);
40
        for(int i = 0; i < len; i++) c[i].r /= len; // 转化为int等时应加0.5, 如int(c[i].r+0.5)
42
```

1.3 高精度-wrz

```
1 #include<cmath>
2 #include<cstdio>
3 #include<cstring>
```

```
#include < algorithm >
#define BASE 10000
#define L 20005
  5
  6
7
     using namespace std;
int p;
char s[10*L];
  8
  9
 10
     struct bigint
 11
            int num[L], len;
bigint(int x = 0)
 12
 13
 14
 15
                  memset(num,0,sizeof(num));
                  len = 1;
num[0] = x;
 16
 17
 18
 19
            bigint operator + (bigint b)
 20
                 bigint c;
c.len = max(b.len, len);
for(int i = 0; i < c.len; i++)
21
22
 23
 24
                        c.num[i] += num[i] + b.num[i];
c.num[i+1] = c.num[i] / BASE;
c.num[i] %= BASE;
 25
 26
 27
 28
 29
                  if(c.num[c.len])c.len++;
 30
                  return c;
 31
            bigint operator - (bigint b)
 32
 33
                  bigint c;
c.len = max(len, b.len);
for(int i = 0; i < c.len; i++)</pre>
 34
 35
 36
 37
                        c.num[i] += num[i] - b.num[i];
if(c.num[i] < 0)</pre>
 38
 39
 40
                              c.num[i] += BASE;
c.num[i+1]--;
 41
 42
 43
 44
 45
                  while(!c.num[c.len-1] && c.len > 1)c.len--;
 46
 47
            void operator -= (int b)
 48
 49
 50
                  num[0] -= b;
 51
                  for(int i = 0; i < len; i++)
 52
 53
                        num[i+1] += num[i] / BASE;
                        num[i] %= BASE;
if(num[i] < 0)num[i] += BASE, num[i+1]--;</pre>
 54
 55
 56
                  while(!num[len-1] && len > 1) len--;
 57
 58
 59
            bigint operator * (bigint b)
 60
                  bigint c;
c.len = len + b.len;
 61
 62
 63
                  for(int i = 0; i < len; i++)
 64
 65
                        for(int j = 0; j < b.len; j++)
 66
                              c.num[i+j] += num[i] * b.num[j];
c.num[i+j+1] += c.num[i+j] / BASE;
c.num[i+j] %= BASE;
 67
 68
 69
 70
 71
 72
73
74
                  if(!c.num[c.len-1] && c.len > 1)c.len--;
                  return c;
 75
            bigint operator * (int b)
 76
77
78
                 bigint c;
for(int i = 0; i < len; i++)
    c.num[i] = num[i] * b; // long long
for(int i = 0; i < len; i++)</pre>
 80
 81
                        c.num[i+1] += c.num[i] / BASE;
c.num[i] %= BASE;
 82
 83
 84
 85
                  c.len = len;
                  while(c.num[c.len])c.len++;
 86
 87
                  return c;
 88
 89
            bool substract(bigint b, int pos)
 90
                  if(len < b.len - pos)return false;
else if(len == b.len-pos)
    for(int i = len-1; i>=0; i--)
        if(num[i] < b.num[i+pos])return false;
        else if(num[i] > b.num[i+pos])break;
 91
 92
 93
 94
 95
                  for(int i = 0; i < len; i++)
 97
 98
                        num[i] -= b.num[i+pos];
 99
                        if(num[i] < 0)</pre>
100
101
                              num[i] += BASE;
                              num[i+1] --;
102
103
104
                  while(!num[len-1] && len > 1)len--;
```

```
106
                  return true;
107
108
            // remember to change [BASE] to 10 !!!
// [this] is the remainder
bigint operator / (bigint b)
109
110
111
112
113
                  bigint c;
if(len < b.len)return c;</pre>
114
                  int k = len - b.len;
c.len = k + 1;
for(int i = len-1; i>=0; i--)
115
116
117
118
119
                         if(i>=k)b.num[i] = b.num[i-k];
                         else b.num[i] = 0;
120
121
                  for(int i = 0; i <= k; i++)
    while(this->substract(b,i)) c.num[k-i]++;
for(int i = 0; i < c.len; i++)</pre>
122
123
124
125
126
                         c.num[i+1] += c.num[i] / BASE;
c.num[i] %= BASE;
127
128
129
                  while(!c.num[c.len-1] && c.len > 0) c.len--;
130
131
132
            }
133
134
            // [this] is not the remainder
135
            bigint operator / (int b)
136
                  bigint c;
int tmp = 0;
for(int i = len-1; i>=0; i--)
137
138
139
140
                        tmp = tmp * BASE + num[i];
c.num[i] = tmp / b;
141
142
143
                         tmp %= b;
144
                  for(c.len = len; !c.num[c.len-1] && c.len > 1; c.len--);
145
146
                  return c;
147
            bool scan()
148
149
150
                  int n = -1:
                  char ch = getchar();
while(ch < '0' || ch
151
                  while(ch < '0' || ch > '9') if(ch == EOF)return false; else ch = getchar();
while(ch >= '0' && ch <= '9') s[++n] = ch - '0', ch = getchar();</pre>
152
153
154
                  len = 0;
for(int i = n; i >= 0; i-=4)
155
156
                        num[len] += s[i];
if(i>=1)num[len] += s[i-1] * 10;
if(i>=2)num[len] += s[i-2] * 100;
if(i>=3)num[len] += s[i-3] * 1000;
157
158
159
160
161
                         ++len;
162
163
                  return true;
164
165
            void clr()
166
167
                  memset(num,0,sizeof(num));
168
169
            void print()
170
                  printf("%d",num[len-1]);
for(int i = len-2; i>=0; i--)
   printf("%04d",num[i]);
171
172
173
174
                  printf("\n");
175
            }
```

1.4 线性基-gwx

```
ll solve()
3
         ll res
                   = 0;
         for(int j = 60; j >= 0; j--)

if((a[i] >> j) & 1)
4
5
6
8
9
                          if(!b[j])
10
                               b[j] = a[i];
11
12
                               break;
13
                          a[i] ^= b[j];
14
15
         for(int i = 60; i >= 0; i--)
  res = max(res, res ^ b[i]);
16
17
18
         return res;
19
```

1.5 单纯形

```
1 // max{c * x | Ax <= b, x >= 0}的解, 无解返回空的vector, 否则就是解. 答案在an中
2 template <int MAXN = 100, int MAXM = 100>
3 struct simplex {
   int n, m; double a[MAXM][MAXN], b[MAXM], c[MAXN];
```

```
bool infeasible,
                                  unbounded;
          double v, an[MAXN + MAXM]; int q[MAXN + MAXM];
void pivot (int 1, int e) {
 6
7
               8
10
11
12
13
14
                      t = c[e]; c[e] = 0; v += t * b[1];
for (int j = 0; j < n; ++j) c[j] -= t * a[1][j]; } }</pre>
15
16
17
          bool pre () {
                for (int 1, e; ; ) {
    1 = e = -1;
    for (int i = 0; i < m; ++i) if (b[i] < -EPS && (!~1 || rand () & 1)) l = i;</pre>
18
19
20
21
                      if (!~1) return false;
for (int i = 0; i < n; ++i) if (a[1][i] < -EPS && (!~e || rand () & 1)) e = i;</pre>
22
23
                         (!~e) return infeasible = true;
          pivot (1, e); } }
double solve () {
24
25
                double p; std::fill (q, q + n + m, -1);
for (int i = 0; i < n; ++i) q[i] = i;
v = 0; infeasible = unbounded = false;</pre>
26
28
                if (pre ()) return 0;
for (int 1, e; ; pivot (1, e)) {
    1 = e = -1; for (int i = 0; i < n; ++i) if (c[i] > EPS) { e = i; break; }
29
30
31
                     if (!~e) break; p = INF;
for (int i = 0; i < m; ++i) if (a[i][e] > EPS && p > b[i] / a[i][e])
    p = b[i] / a[i][e], l = i;
if (!~l) return unbounded = true, 0; }
32
33
34
35
36
                             i = n; i < n + m; ++i) if (~q[i]) an[q[i]] = b[i - n];
                return v; } };
```

1.6 NTT-gwx

```
const int G;
   int n. m. inm:
   int rev[maxn], a[maxn], b[maxn];
                                                //maxn > 2 ^ k
 5
   ll power(ll b, int k)
 6
7
         ll res = 1;
 8
         for(; k; k >>= 1, b = b * b % mod)
 9
             if(k & 1)
10
                  res = res * b % mod:
11
        return res;
13
14
   void ntt(ll*a, int f)
15
16
         for(int i = 0; i < m; i++)
        if(rev[i] < i)

swap(a[i], a[rev[i]]);

for(int 1 = 2, h = 1; 1 <= m; h = 1, 1 <<= 1)
17
18
19
21
22
             int ur;
if(f == 1)
                  ur = power(G, (mod - 1) / 1);
23
24
25
26
27
             ur = power(G, mod - 1 - (mod - 1) / 1);
for(int i = 0; i < m; i += 1)
29
30
                  for(int k = i; k < i + h; k++, w = w * ur % mod)
31
                       32
33
34
             }
35
36
37
         if(f == -1)
             for(int i = 0; i < m; i++)
    a[i] = a[i] * inm % mod;</pre>
38
39
40
41
42
   void multi()
43
44
        ntt(a, 1); ntt(b, 1);
for(int i = 0; i < m; i++)
   a[i] = a[i] * b[i] % mod;</pre>
45
46
47
        ntt(a, -1);
48
49
50
    void init()
51
        for(m = 1; m <= 2 * n; m <<= 1) ;
for(int i = 1; i < m; i++)
52
53
             55
56
57
```

2 计算几何

2.1 二维计算几何-wrz

```
#include < bits / stdc++.h>
 3
4
5
   using namespace std;
   const double inf = 1e9;
   const double eps = 1e-9;
   const double pi = acos(-1.0);
 8
10
   /*精度误差下的各种运算*/
11 bool le(double x, double y){return x < y - eps;} // x严格小于y
   bool leq(double x, double y){return x < y + eps;} // x小于等于y
12
   bool equ(double x, double y){return fabs(x - y) < eps;} // x小寸等寸y
double mysqrt(double x) {return x < eps ? 0 : sqrt(x);} // 开根号
double sqr(double x) {return x * x;} // 平方
13
14
15
16
17
   struct point // 点或向量
18
19
        double x, y;
        double operator *
20
                           (point that){return x*that.x + y*that.y;}
       double operator ^ (point that){return x*that, x - y*that.x;}
point operator * (double t){return (point){x*t, y*t};}
21
22
23
        point operator
                           (double t){return (point){x/t, y/t};}
24
       point operator + (point that) {return (point){x + that.x, y + that.y};}
point operator - (point that) {return (point){x - that.x, y - that.y};}
double len(){return mysqrt(x*x+y*y);} // 到原点距离/向量长度
25
26
        point reset_len(double t) // 改变向量长度为t, t为正则方向不变, t为负则方向相反
27
28
            double p = len();
return (point){x*t/p, y*t/p};
29
30
31
       point rot90() {return (point){-y, x};} // 逆时针旋转90度
32
        point rotate(double angle) // 使向量逆时针旋转angle弧度
33
34
            double c = cos(angle), s = sin(angle);
return (point){c * x - s * y, s * x + c * y};
35
36
37
38
   };
39
   struct line // 参数方程表示, p为线上一点, v为方向向量
40
41
42
       point p, v; // p为线上一点, v为方向向量
43
44
        double angle; // 半平面交用, 用atan2计算, 此时v的左侧为表示的半平面。注意有的函数声明一个新的line时没有初始化
45
        bool operator < (const line &that) const {return angle < that.angle;} // 半平面交用,接与x轴夹角排序
   };
46
47
48
   struct circle
49
50
       point c; double r;
51
   };
52
53
54
   double distance(point a, point b) // a, b两点距离
55
56
        return mysqrt(sqr(a.x - b.x) + sqr(a.y - b.y));
57
58
59
   circle make_circle(point a, point b) // 以a, b两点为直径作圆
60
       double d = distance(a, b);
return (circle){(a+b)/2, d/2};
61
62
63
64
   double point_to_line(point a, line b) // 点a到直线b距离
65
66
        return fabs((b.v ^ (a - b.p)) / b.v.len());
67
68
69
70
   point project_to_line(point a, line b) // 点a到直线b的垂足/投影
71
72
73
        return b.v.reset_len((a - b.p) * b.v / b.v.len()) + b.p;
74
   vector<point> circle_inter(circle a, circle b) // 圆a和圆b的交点,需保证两圆不重合,圆的半径必须大于0
75
76
        double d = distance(a.c, b.c);
77
78
        vector < point > ret;
79
        if(le(a.r + b.r, d)
                              || le(a.r + d, b.r) || le(b.r + d, a.r)) return vector<point>(); // 相离或内含
       point r = (b.c - a.c).reset_len(1);
double x = ((sqr(a.r) - sqr(b.r)) / d + d) / 2;
double h = mysqrt(sqr(a.r) - sqr(x));
80
81
82
83
        if(equ(h, 0)) return vector<point>({a.c + r * x}); // 内切或外切
84
        else return vector<point>({a.c + r*x + r.rot90()*h, a.c + r*x - r.rot90()*h}); // 相交两点
85
86
87
   vector<point> line_circle_inter(line a, circle b) // 直线a和 圆b的交点
88
89
        double d = point_to_line(b.c, a);
90
        if(le(b.r, d)) return vector<point>(); // 不交
91
        double x = mysqrt(sqr(b.r) - sqr(d));
        point p = project_to_line(b.c,
92
        if(equ(x, 0)) return vector<point> ({p}); // 相切
93
94
        else return vector<point> ({p + a.v.reset_len(x), p - a.v.reset_len(x)}); // 相交两点
95
96
   point line_inter(line a, line b) // 直线a和直线b的交点,需保证两直线不平行
97
        double s1 = a.v \hat{(b.p - a.p)};
99
```

```
100
        double s2 = a.v \hat{b.p} + b.v
                                        a.p)
        return (b.p * s2 - (b.p + b.v) * s1) / (s2 - s1);
101
102
103
    vector<point> tangent(point p, circle a) // 过点p的圆a的切线的切点, 圆的半径必须大于0
104
105
106
        circle c = make_circle(p, a.c);
107
        return circle_inter(a, c);
108
109
110
    vector<line> intangent(circle a, circle b) // 圆a和圆b的内公切线
111
        point p = (b.c * a.r + a.c * b.r) / (a.r + b.r);
112
        vector<point> va = tangent(p, a), vb = tangent(p, b);
vector<line> ret;
if(va.size() == 2 && vb.size() == 2)
113
114
115
116
            ret.push_back((line){va[0], vb[0] - va[0]});
ret.push_back((line){va[1], vb[1] - va[1]});
117
118
119
120
        else if(va.size() == 1 && vb.size() == 1)
121
            ret.push_back((line){p, (a.c - b.c).rot90()});
122
123
124
        return ret;
125
126
127 // 判断半平面交是否有解,若有解需保证半平面交必须有界,可以通过外加四个大半平面解决
128 // 1cnt为半平面数量,1为需要做的所有半平面的数组,p为存交点的临时数组,h为时刻更新的合法的半平面数组,下标均从1开
129
    bool HP(int lcnt, line *1, line *h, point *p)
130
        sort(l+1, l+1+lcnt);
int head = 1, tail = 1;
h[1] = 1[1];
for(int i = 2; i <= lcnt; i++)</pre>
131
132
133
134
135
136
            line cur = l[i];
            for(; head < tail && le(cur.v ^ (p[tail-1]-cur.p), 0); tail--); // 先删队尾再删队头,顺序不能换for(; head < tail && le(cur.v ^ (p[head]-cur.p), 0); head++);
137
138
            h[++tail] = cur;
139
            if(equ(h[tail].v ^ h[tail-1].v, 0)) // 平行
140
141
                if(le(h[tail].v * h[tail-1].v, 0)) return false; // 方向相反的平行直线,显然已经不可能围出有界半平面
142
143
                tail-
144
                if(le(h[tail+1].v ^ (h[tail].p - h[tail+1].p), 0)) h[tail] = h[tail+1];
145
146
            if(head < tail) p[tail-1] = line_inter(h[tail-1], h[tail]);</pre>
147
        for(; head < tail && le(h[head].v ^ (p[tail-1]-h[head].p), 0); tail--); return tail - head > 1;
148
149
150
151
152 double calc(double X) {return 0;} // 计算给定X坐标上的覆盖的长度,配合辛普森积分使用
153 // 自适应辛普森积分,参数分别为(左端点x坐标,中点x坐标,右端点x坐标,左端点答案,中点答案,右端点答案)154 // 改变计算深度应调整eps
    double simpson(double 1, double mid, double r, double f1, double fm, double fr)
156
        157
158
159
        if(fabs(ansl + ansr - ans) < eps) return ans / 6;</pre>
160
        else return simpson(1,lmid,mid,f1,flm,fm) + simpson(mid,rmid,r,fm,frm,fr);
161
163
    int main(){}
164
```

2.2 basis

```
int sign(DB x) {
           return (x > eps) - (x < -eps);
 2
 3
 4
    DB msgrt(DB x) {
 5
          return sign(x) > 0 ? sqrt(x) : 0;
 6
    }
    struct Point {
 8
          DB x, y;
          Point rotate(DB ang) const { // 逆时针旋转 ang 弧度 return Point(cos(ang) * x - sin(ang) * y,
10
11
12
                             cos(ang) * y + sin(ang) * x);
13
          Point turn90() const { // 逆时针旋转 90 度 return Point(-y, x);
14
15
16
17
          Point unit() const {
18
                return *this / len();
19
20
    DB dot(const Point& a, const Point& b) {
    return a.x * b.x + a.y * b.y;
22
23
    DB det(const Point& a, const Point& b) {
    return a.x * b.y - a.y * b.x;
24
26
27 #define cross(p1,p2,p3) ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
28 #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
29 | bool isLL(const Line& 11, const Line& 12, Point& p) { // 直线与直线交点
```

```
DB s1 = det(12.b - 12.a, 11.a - 12.a),
    s2 = -det(12.b - 12.a, 11.b - 12.a);
if (!sign(s1 + s2)) return false;
p = (11.a * s2 + 11.b * s1) / (s1 + s2);
return true;
   31
  32
   33
   34
   35
         bool onSeg(const Line& 1, const Point& p) { // 点在线段上
return sign(det(p - 1.a, 1.b - 1.a)) == 0 && sign(dot(p - 1.a, p - 1.b)) <= 0;
  36
  37
  38
         Point projection(const Line & 1, const Point point point projection (const Line & 1, const Point point point point point projection (const Line & 1, const Point point point projection (const Line & 1, const Point point projection (const Line & 1, const Point point projection (const Line & 1, const Point projection (const Lin
  39
   40
  41
         DB disToLine(const Line& 1, const Point& p) { // 点到*直线*距离 return fabs(det(p - l.a, l.b - l.a) / (l.b - l.a).len());
   42
  43
  44
  47
        48
   49
  50
  51
  53
  54
  55
  57
        58
  59
  60
   61
   62
                               DB r = std::min(c1.r, c2.r);
   63
                               return r * r * PI;
   64
                    66
   67
  68
   69
  70
          // 圆与圆交点
          bool isCC(Circle a, Circle b, P& p1, P& p2) {
                    I isCC(Circle a, Circle b, P& p1, P& p2) {
DB s1 = (a.o - b.o).len();
if (sign(s1 - a.r - b.r) > 0 || sign(s1 - std::abs(a.r - b.r)) < 0) return false;
DB s2 = (a.r * a.r - b.r * b.r) / s1;
DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
P o = (b.o - a.o) * (aa / (aa + bb)) + a.o;
P delta = (b.o - a.o).unit().turn90() * msqrt(a.r * a.r - aa * aa);
p1 = o + delta, p2 = o - delta;
return true;
  72
  73
74
   75
   76
   77
   78
  80 }
         81
  82
  83
  84
  85
                    p1 = c.o + p + delta;
p2 = c.o + p - delta;
  87
  88
  89
                     return true;
  90
         91
  92
  93
  94
  95
  96
  97
  98
                    } else
                              Point p = (c1.o * -c2.r + c2.o * c1.r) / (c1.r - c2.r);

Point p1, p2, q1, q2;

if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) {

   if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);

   ret.push_back(Line(p1, q1));
100
101
102
103
104
105
                                         ret.push_back(Line(p2, q2));
                              }
106
107
108
                     return ret;
109 }
if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { // 两圆相切认为没有切线 ret.push_back(Line(p1, q1)); ret.push_back(Line(p2, q2));
115
116
117
118
119
                    return ret;
120 }
         bool contain(vector<Point> polygon, Point p) { // 判断点 p 是否被多边形包含,包括落在边界上int ret = 0, n = polygon.size();
for(int i = 0; i < n; ++ i) {
121
122
123
                               int i = 0; i < n; ff i / i
Point u = polygon[i], v = polygon[(i + 1) % n];
if (onSeg(Line(u, v), p)) return true; // Here I guess.
if (sign(u.y - v.y) <= 0) swap(u, v);
if (sign(p.y - u.y) > 0 || sign(p.y - v.y) <= 0) continue;</pre>
124
125
126
127
                               ret += sign(det(p, v, u)) > 0;
```

```
130
          return ret & 1;
131 }
    // 用半平面 (q1,q2) 的逆时针方向去切凸多边形
std::vector<Point> convexCut(const std::vector<Point>&ps, Point q1, Point q2) {
132 | //
133
          std::vector<Point> qs; int n = ps.size();
for (int i = 0; i < n; ++i) {</pre>
134
135
               136
137
138
139
140
141
          return qs;
142 }
143 // 求凸包
144 std::vector<Point> convexHull(std::vector<Point> ps) {
145
          int n = ps.size(); if (n <= 1) return ps;</pre>
146
          std::sort(ps.begin(), f
std::vector<Point> qs;
for (int i = 0; i < n; qs.push_back(ps[i ++]))
    while (qs.size() > 1 && sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <= 0)</pre>
          std::sort(ps.begin(), ps.end());
147
148
149
          qs.pop_back();
for (int i = n - 2, t = qs.size(); i >= 0; qs.push_back(ps[i --]))
while ((int)qs.size() > t & sign(det(qs[qs.size() - 2], qs.back(), ps[i])) <= 0)
150
151
152
                     qs.pop_back();
154
          return qs;
```

2.3 circle-arc-struct

```
struct circle {
          point o;
double r;
 3
          circle(point o, double r) : o(o), r(r) {}
 5
    };
    struct arcs { // 点1顺时针到点r
          point o, 1, r;
arcs() {}
 8
          arcs(point o, point 1, point r) : o(o), 1(1), r(r) {}
10
    bool isCL(circle a, line 1, point &p1, point &p2) { // 圆与直线的交点
          double x = dot(1.a - a.o, 1.b - 1.a);

double y = (1.b - 1.a).len2();

double d = x * x - y * ((1.a - a.o).len2() - a.r * a.r);

if(sign(d) < 0) return false;
12
13
14
15
16
          d = max(d, 0.0);
          point p = 1.a - ((1.b - 1.a) * (x / y)), delta = (1.b - 1.a) * (sqrt(d) / y); p1 = p + delta, p2 = p - delta;
17
18
          return true;
20
    double ang(const point &d1, const point &d2) { // 向量d1顺时针转到向量d2的角度if (sign(det(d1, d2)) == 0)
21
22
                 return (sign(dot(d1, d2)) > 0) ? 0 : pi;
sign(det(d1, d2)) < 0)
23
24
          if (sign(det(d1, d2)) < 0)
return acos(dot(d1, d2) / d1.len() / d2.len());
25
26
          else return 2 * pi - acos(dot(d1, d2) / d1.len() / d2.len());
27
    7
    bool onArcs(const point &p, const arcs &a) { // 点在圆弧上 return sign(ang(a.1 - a.o, a.r - a.o) - ang(a.1 - a.o, p - a.o)) > 0;
28
29
30
31
32
    /*struct circle {
33
          point o;
double r;
34
          circle(point o, double r) : o(o), r(r) {}
35
    };
36
37
    struct arcs {
39
          //l -> r clockwise
          point o, l, r;
arcs() {}
40
41
42
          arcs(point o, point 1, point r) : o(o), l(1), r(r) {}
43
44
    //Circle intersect with Line
bool isCL(circle a, line l, point &p1, point &p2) {
   double x = dot(l.a - a.o, l.b - l.a);
   double y = (l.b - l.a).len2();
   double d = x * x - y * ((l.a - a.o).len2() - a.i
   if(sign(d) < 0) return false;
}</pre>
45
47
48
                                                        - a.o).len2() - a.r * a.r);
49
50
          d = max(d, 0.0);
point p = 1.a - ((1.b - 1.a) * (x / y)), delta = (1.b - 1.a) * (sqrt(d) / y);
p1 = p + delta, p2 = p - delta;
51
52
53
55
56
    //use acos !precision
57
59
     //angle (d1, d2) clockwise
    double ang(const point &d1, const point &d2) {
   if (sign(det(d1, d2)) == 0)
      return (sign(dot(d1, d2)) > 0) ? 0 : pi;
60
61
62
63
          if (sign(det(d1, d2)) < 0)
          return acos(dot(d1, d2) / d1.len() / d2.len());
else return 2 * pi - acos(dot(d1, d2) / d1.len() / d2.len());
64
65
66
    bool onArcs(const point &p, const arcs &a) {
   return sign(ang(a.1 - a.o, a.r - a.o) - ang(a.1 - a.o, p - a.o)) > 0;
68
69
70
```

2.4 circles-intersections

```
struct Event
              Point p;
double ang;
               int delta
 5
               Event \ (Point \ p = Point(0, \ 0), \ double \ ang = 0, \ double \ delta = 0) \ : \ p(p), \ ang(ang), \ delta(delta) \ \{\} 
 6
      bool operator < (const Event &a, const Event &b) {
 8
              return a.ang < b.ang;
 9

}
void addEvent(const Circle &a, const Circle &b, vector<Event> &evt, int &cnt) {
    double d2 = (a.o - b.o).len2(),
        dRatio = ((a.r - b.r) * (a.r + b.r) / d2 + 1) / 2,
        pRatio = sqrt(-(d2 - sqr(a.r - b.r)) * (d2 - sqr(a.r + b.r)) / (d2 * d2 * 4));

Point d = b.o - a.o, p = d.rotate(PI / 2),
        q0 = a.o + d * dRatio + p * pRatio,
        q1 = a.o + d * dRatio - p * pRatio;

double ang0 = (q0 - a.o).ang(),
        ang1 = (q1 - a.o).ang();
    evt.push_back(Event(q1, ang1, 1));
    evt.push_back(Event(q0, ang0, -1));
    cnt += ang1 > ang0;
}

10
11
12
13
14
15
17
18
19
21
22
      bool issame(const Circle &a, const Circle &b) { return sign((a.o - b.o).len()) == 0 && sign(a.r - b.r) == 0; } bool overlap(const Circle &a, const Circle &b) { return sign(a.r - b.r - (a.o - b.o).len()) >= 0; } bool intersect(const Circle &a, const Circle &b) { return sign((a.o - b.o).len() - a.r - b.r) < 0; }
23
25
     Circle c[N];
double area[N]; /
Point centroid[N];
26
27
                                          // area[k] -> area of intersections >= k.
     bool keep[N];
void add(int cnt, DB a, Point c) {
    area[cnt] += a;
    centroid[cnt] = centroid[cnt] + c * a;
29
30
31
32
34
      void solve(int C) {
              for (int i = 1; i <= C; ++ i) {
    area[i] = 0;
    centroid[i] = Point(0, 0);</pre>
35
36
37
38
39
              for (int i = 0; i < C; ++i) {
                       int cnt = 1;
40
41
                       vector < Event > evt;
                       for (int j = 0; j < i; ++j) if (issame(c[i], c[j])) ++cnt; for (int j = 0; j < C; ++j) {    if (j != i && !issame(c[i], c[j]) && overlap(c[j], c[i])) {
42
43
44
45
                                          +cnt;
46
47
48
                       for (int j = 0; j < C; ++j) {
    if (j != i && !overlap(c[j], c[i]) && !overlap(c[i], c[j]) && intersect(c[i], c[j])) {
49
50
                                        addEvent(c[i], c[j], evt, cnt);
51
52
53
                       if (evt.size() == 0u) {
                               add(cnt, PI * c[i].r * c[i].r, c[i].o);
55
                       } else {
                               sort(evt.begin(), evt.end());
evt.push_back(evt.front());
for (int j = 0; j + 1 < (int)evt.size(); ++j) {
    cnt += evt[j].delta;
    cnt += evt[j].delta;</pre>
56
57
58
59
                                        add(cnt, det(evt[j].p, evt[j + 1].p) / 2, (evt[j].p + evt[j + 1].p) / 3);
double ang = evt[j + 1].ang - evt[j].ang;
if (ang < 0) {
60
61
62
                                                ang += PI * 2;
63
64
                                        f (sign(ang) == 0) continue;
add(cnt, ang * c[i].r * c[i].r / 2, c[i].o +
    Point(sin(ang1) - sin(ang0), -cos(ang1) + cos(ang0)) * (2 / (3 * ang) * c[i].r));
add(cnt, -sin(ang) * c[i].r * c[i].r / 2, (c[i].o + evt[j].p + evt[j + 1].p) / 3);
65
66
68
                               }
69
70
71
                      }
72
73
74
              for (int i = 1; i <= C; ++ i)
if (sign(area[i])) {
                               centroid[i] = centroid[i] / area[i];
75
```

2.5 cirque-area-merge

```
//n^2*logn
   struct point {
        point rotate(const double &ang) {
    return point(cos(ang) * x - sin(ang) * y, cos(ang) * y + sin(ang) * x);
3
5
6
        double ang() {
            return atan2(y, x);
8
10
   struct Circle {
        point o; double r;
11
12
        int tp; // 正圆为1 反向圆为-1
Circle (point o = point(0, 0), double r = 0, int tp = 0) : o(o), r(r), tp(tp) {}
13
14
16
   struct Event
        point p;
double ang;
17
18
19
20
        Event (point p = point(0, 0), double ang = 0, double delta = 0) : p(p), ang(ang), delta(delta) {}
```

```
bool operator < (const Event &a, const Event &b) {
23
                return a.ang < b.ang;
24
25
      void addEvent(const Circle &a, const Circle &b, vector<Event> &evt, int &cnt) {
               d addEvent(const Circle &a, const Circle &b, vector<Event> &evt, int &cnt) {
   double d2 = (a.o - b.o).len2(),
      dRatio = ((a.r - b.r) * (a.r + b.r) / d2 + 1) / 2,
      pRatio = sqrt(-(d2 - sqr(a.r - b.r)) * (d2 - sqr(a.r + b.r))) / d2 / 2;
   point d = b.o - a.o, p = d.rotate(PI / 2),
      q0 = a.o + d * dRatio + p * pRatio,
      q1 = a.o + d * dRatio - p * pRatio;
   double ang0 = (q0 - a.o).ang(),
      ang1 = (q1 - a.o).ang();
   evt.push_back(Event(q1, ang1, b.tp));
   evt.push_back(Event(q0, ang0, -b.tp));
   cnt += (ang1 > ang0) * b.tp;
26
27
28
29
30
31
32
33
35
36
37
      bool issame(const Circle &a, const Circle &b) {
    return sign((a.o - b.o).len()) == 0 && sign(a.r - b.r) == 0;
39
40
      bool overlap(const Circle &a, const Circle &b) {
   return sign(a.r - b.r - (a.o - b.o).len()) >= 0;
41
42
43
      bool intersect(const Circle &a, const Circle &b) {
   return sign((a.o - b.o).len() - a.r - b.r) < 0;
44
45
46
47
      int C;
Circle c[N];
48
49
50
      double area[N];
      double area[N];
void solve() { // area[1]..area[C]
  memset(area, 0, sizeof(double) * (C + 1));
  for (int i = 0; i < C; ++i) {
    int cnt = (c[i].tp > 0);
}
52
53
54
                         int cnt = (c[i].tp > 0);
vector<Event> evt;
for (int j = 0; j < i; ++j) if (issame(c[i], c[j])) cnt += c[j].tp;
for (int j = 0; j < C; ++j)
    if (j != i && !issame(c[i], c[j]) && overlap(c[j], c[i])) cnt += c[j].tp;</pre>
56
57
58
                         for (int j = 0; j < C; ++j)
    if (j != i && !overlap(c[j], c[i]) && !overlap(c[i], c[j]) && intersect(c[i], c[j]))
60
                         addEvent(c[i], c[j], evt, cnt);
if (evt.size() == 0) area[cnt] += c[i].tp * PI * c[i].r * c[i].r;
61
62
63
                          else
64
                                  sort(evt.begin(), evt.end());
65
                                   evt.push_back(evt.front());
                                  evt.pusn_back(evt.front());
for (int j = 0; j + 1 < (int)evt.size(); ++j) {
    cnt += evt[j].delta;
    area[cnt] += c[i].tp * det(evt[j].p, evt[j + 1].p) / 2;
    double ang = evt[j + 1].ang - evt[j].ang;
    if (ang < 0) ang += PI * 2;
    area[cnt] += c[i].tp * (ang * c[i].r * c[i].r / 2 - sin(ang) * c[i].r * c[i].r / 2);</pre>
66
67
68
69
70
71
72
73
74
                         }
                }
      }
75
```

2.6 凸包

```
// 凸包中的点按逆时针方向
   struct Convex
2
3
4
        int n;
std::vector<Point> a, upper, lower;
void make_shell(const std::vector<Point>& p,
5
6
7
                  std::vector<Point>& shell) { // p needs to be sorted.
             8
9
10
                   shell.push_back(p[i]);
11
12
             }
13
14
        void make_convex() {
             std::sort(a.begin(), a.end());
make_shell(a, lower);
std::reverse(a.begin(), a.end());
15
16
17
18
             make_shell(a, upper);
             a.insert(a.end(), upper.begin(), upper.end());
if ((int)a.size() >= 2) a.pop_back();
n = a.size();
19
20
21
22
23
24
25
         void init(const std::vector<Point>& _a) {
             clear(a); a = _a; n = a.size();
make_convex();
26
27
28
29
        void read(int _n) {      // Won't make
      clear(a);       n = _n;       a.resize(n);
      for (int i = 0; i < n; i++)</pre>
                                  // Won't make convex.
30
31
                  a[i].read();
        33
34
35
              assert(r >= 0);
36
             for (; 1 + 1 < r; ) {
   int mid = (1 + r) / 2;
37
38
                   if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
    r = mid.
39
40
                           mid
41
42
43
             return std::max(std::make_pair(det(vec, convex[r]), r),
```

```
std::make_pair(det(vec, convex[0]), 0));
45
46
          int binary_search(Point u, Point v, int 1, int r) {
                int s1 = sign(det(v - u, a[1 % n] - u));
for (; 1 + 1 < r; ) {
   int mid = (1 + r) / 2;</pre>
47
48
49
                       int smid = sign(det(v - u, a[mid % n] - u));
if (smid == s1) l = mid;
50
51
52
                       else r = mid;
54
                 return 1 % n;
55
          // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个int get_tangent(Point vec) {
    std::pair<DB, int> ret = get_tangent(upper, vec);
    ret.second = (ret.second + (int)lower.size() - 1) % n;
56
57
58
59
60
                ret = std::max(ret, get_tangent(lower, vec));
61
                return ret.second;
62
63
          // 求凸包和直线 u, v 的交点, 如果不相交返回 false, 如果有则是和 (i, next(i)) 的交点, 交在点上不确定返回前后两
          bool get_intersection(Point u, Point v, int &i0, int &i1) {
                int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
    if (p0 > p1) std::swap(p0, p1);
    i0 = binary_search(u, v, p0, p1);
    i1 = binary_search(u, v, p1, p0 + n);
65
66
67
68
70
                       return true;
71
72
                 else return false;
73
74
    };
```

2.7 point-in-polygon

```
bool pit_in_polygon(pit q){ // µ
               int cnt = 0;
for(int i = 1; i <= n; ++i){
  3
                        pit p1 = p[i];
pit p2 = p[suc[i]];
  5
                        prt p2 - p[suc[1]];
if(pit_on_seg(q, p1, p2)) return true;
int k = dcmp(det(p2 - p1, q - p1));
int d1 = dcmp(p1.y - q.y);
int d2 = dcmp(p2.y - q.y);
if(k > 0 && d1 <= 0 && d2 > 0) ++cnt;
if(k < 0 && d2 <= 0 && d1 > 0) --cnt;
  6
7
  9
10
11
12
                if(cnt != 0) return true;
14
                else return false;
15
16
      bool seg_in_polygon(pit a, pit b){ //
    vec v = b - a;
               vec v = b - a;
for(int t = 1; t <= 1000; ++t){
   pit c = a + v * (1.00 * (rand() % 10000) / 10000);
   if(pit_in_polygon(c)) continue;
   else return false;</pre>
17
18
19
20
21
22
23
                return true;
```

2.8 point-struct

```
const double eps = 1e-8;
const double PI = acos(-1.0);
   int sign(double x) {
    return (x < -eps) ? -1 : (x > eps);
 4
5
 6
    double sqr(double x) {
 8
         return x * x;
 9
10
11
    struct point {
         double x, y;
point(double x = 0, double y = 0) : x(x), y(y) {}
point(double x = 0, double y = 0) : x(x), y(y) {}
12
13
14
15
              return point(x + rhs.x, y + rhs.y);
16
         point operator - (const point &rhs) const {
   return point(x - rhs.x, y - rhs.y);
17
18
19
20
         point operator * (const double &k) const {
21
              return point(x * k, y * k);
22
23
         point operator / (const double &k) const {
              return point(x / k, y / k);
24
25
26
         double len2() {
              return x * x + y * y;
28
29
         double len() {
30
              return sqrt(len2());
31
         point rotate(const double &ang) { // 逆时针旋转 ang 弧度 return point(cos(ang) * x - sin(ang) * y, cos(ang) * y + sin(ang) * x);
32
33
34
35
         point turn90() { // 逆时针旋转 90 度
36
              return point(-y, x);
37
         double ang() {
```

```
return atan2(y, x);
40
           point operator < (const point &rhs) {
   return (x < rhs.x || x == rhs.x && y < rhs.y);</pre>
41
42
43
    double dot(const point &a, const point &b) {
   return a.x * b.x + a.y * b.y;
45
46
47
    double det(const point &a, const point &b) {
   return a.x * b.y - a.y * b.x;
49
50
51
    struct line {
          point a, b;
line(point a, point b) : a(a), b(b) {}
53
54
55
    bool onSeg(const point &p, const line &l) { // 点在线段上 包含端点 return sign(det(p - l.a, l.b - l.a)) == 0 && sign(dot(p - l.a, p - l.b)) <= 0;
56
58 }
    double disToLine(const point &p, const line &1) { // 点到直线距离 return fabs(det(p - 1.a, 1.b - 1.a) / (1.b - 1.a).len());
59
60
61 l
    double disToSeg(const point &p, const line &1) { // 点到线段距离 return sign(dot(p - 1.a, 1.b - 1.a)) * sign(dot(p - 1.b, 1.a - 1.b)) != 1 ? disToLine(1, p) : min((p - 1.a).len(), (p - 1.b).len());
62
63
65
    point projection(const point &p, const line &1) { // 点到直线投影 return 1.a + (1.b - 1.a) * (dot(p - 1.a, 1.b - 1.a) / (1.b - 1.a).len2());
67
68 }
    point symmetry(const point &a, const point &b) { // 点a关于点b的对称点
return b + b - a;
69
71
    point reflection(const point &p, const line &1) { // 点关于直线的对称点
73
           return symmetry(p, projection(p, 1));
74
    bool isLL(const line &11, const line &12, point &p) { // 直线求交 long long s1 = det(12.b - 12.a, 11.a - 12.a); long long s2 = -det(12.b - 12.a, 11.b - 12.a);
75
76
           if(!sign(s1 + s2)) return false
78
           p = (1\bar{1}.a * s2 + 11.b * s1) / (s1 + s2);
79
80
           return true;
81
    bool p_in_tri(const point &p, const point &a, const point &b, const point &c) { //点在三角形内(包含边界) return sign(abs(det(a - p, b - p)) + abs(det(b - p, c - p)) + abs(det(c - p, a - p)) - abs(det(b - a, c - a))) == 0;
82
84
85
    bool Check(const point &p, const point &d, const point &a, const point &b) {
    return sign(det(d, a - p)) * sign(det(b - p, d)) >= 0;
86
88
    bool isll(const point &p, const point &q, const point &a, const point &b) { // 跨立实验return Check(p, q - p, a, b) && Check(a, b - a, p, q);
89
90
```

2.9 三角形内心,外心,垂心

```
Point inCenter(const Point &A, const Point &B, const Point &C) { // \not | \  \    double a = (B - C).len(), b = (C - A).len(), c = (A - B).len(), s = fabs(det(B - A, C - A)),
3
4
         r = s / p;
return (A * a + B * b + C * c) / (a + b + c);
5
 6
   Point circumCenter(const Point &a, const Point &b, const Point &c) { // 外心
         Point bb = b - a, cc = c - a;
double db = bb.len2(), dc = cc.len2(), d = 2 * det(bb, cc);
return a - Point(bb.y * dc - cc.y * db, cc.x * db - bb.x * dc) / d;
8
9
10
11
   13
14
15
16
17
         return Point(x0, y0);
18
```

2.10 三维计算几何

```
三维绕轴旋转,大拇指指向 axis 向量方向,四指弯曲方向转 w 弧度
 2
       Point rotate(const Point& s, const Point& axis, DB w)
                  DB x = axis.x, y = axis.y, z = axis.z;
DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
                           cosw = cos(w), sinw = sin(w);
                  DB a[4][4];
memset(a, 0, sizeof a);
a[3][3] = 1;
 6
7
                  a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
10
                 a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;

a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;

a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;

a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;

a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;

a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;

a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;

DB ans[4] = {0, 0, 0, 0}, c[4] = {s.x, s.y, s.z, 1};

for (int i = 0; i < 4; ++ i)

for (int j = 0; j < 4; ++ j)

ans[i] += a[j][i] * c[j];
11
12
13
15
16
17
18
19
20
```

```
22 return Point(ans[0], ans[1], ans[2]);
23 }
```

2.11 三维凸包

```
__inline P cross(const P& a, const P& b) {
                        return P(
                                                   a.y * b.z - a.z * b.y,
a.z * b.x - a.x * b.z,
a.x * b.y - a.y * b.x
   3
4
5
                                     );
   7
          }
   8
           __inline DB mix(const P& a, const P& b, const P& c) {
 10
                        return dot(cross(a, b), c);
 11
 12
           __inline DB volume(const P& a, const P& b, const P& c, const P& d) {
 13
 14
                       return mix(b - a, c - a, d - a);
          }
 15
16
          struct Face {
 17
                      uct race {
  int a, b, c;
  __inline Face() {}
  __inline Face(int _a, int _b, int _c):
      a(_a), b(_b), c(_c) {}
  __inline DB area() const {
      return 0 5 * cose(n[b] - n[c] - n[
 19
20
 21
 23
                                 return 0.5 * cross(p[b] - p[a], p[c] - p[a]).len();
                       }
 24
                       __inline P normal() const {
 25
                                    return cross(p[b] - p[a], p[c] - p[a]).unit();
 26
27
28
29
                       __inline DB dis(const P& p0) const {
   return dot(normal(), p0 - p[a]);
 30
31
32
          \label{eq:std:std:std} \mbox{std::vector} < \mbox{Face> face, tmp;} \ // \ \mbox{Should be O(n)} \, .
33
34
          int mark[N][N], Time, n;
36
37
           --inline void add(int v) {
    ++ Time;
                        clear(tmp);
38
                        39
 40
41
 42
 44
                                     }
 45
                                      else {
 46
                                                   tmp.push_back(face[i]);
 47
 48
                       clear(face); face = tmp;
for (int i = 0; i < (int)tmp.size(); ++ i) {
   int a = face[i].a, b = face[i].b, c = face[i].c;
   if (mark[a][b] == Time) face.emplace_back(v, b, a);
   if (mark[b][c] == Time) face.emplace_back(v, c, b);
   if (mark[c][a] == Time) face.emplace_back(v, a, c);
   assert(face.size() < 500u);</pre>
 49
50
51
53
54
55
 56
57
          }
58
          59
60
 61
 62
                                                   std::swap(p[i], p[2]);
for (int j = 3; j < n; ++ j)
    if (sign(volume(p[0], p[1], p[2], p[j]))) {
        std::swap(p[j], p[3]);</pre>
63
64
 65
66
67
                                                                               return:
68
                                                                 }
 69
                                     }
70
71
72
                        }
          }
 73
          void build_convex() {
74
75
76
                        reorder();
                        clear(face);
                        face.emplace_back(0, 1, 2);
face.emplace_back(0, 2, 1);
for (int i = 3; i < n; ++ i)
78
79
80
                                     add(i);
```

3 数据结构

3.1 KD 树

```
1 long long norm(const long long &x) {
2     // For manhattan distance
3     return std::abs(x);
4     // For euclid distance
5     return x * x;
6 }
7
8 struct Point {
9     int x, y, id;
10
```

```
const int& operator [] (int index) const { if (index == 0) {
 12
 13
                      return x;
                } else {
 14
                     return y;
 15
 16
                }
 17
          }
 18
 19
          friend long long dist(const Point &a, const Point &b) {
                long long result = 0;
for (int i = 0; i < 2; ++i) {
    result += norm(a[i] - b[i]);</pre>
 20
21
 22
 23
 24
25
                return result;
 26
     } point[N];
 27
 28
     struct Rectangle {
 29
          int min[2], max[2];
 30
          Rectangle() {
    min[0] = min[1] = INT_MAX; // sometimes int is not enough
    max[0] = max[1] = INT_MIN;
 31
 32
 33
 34
          void add(const Point &p) {
   for (int i = 0; i < 2; ++i) {
      min[i] = std::min(min[i], p[i]);
      max[i] = std::max(max[i], p[i]);</pre>
 36
 37
 38
 39
 40
 41
42
          }
          43
 44
 45
46
 47
 48
                             For maximum distance
                      result += std::max(norm(max[i] - p[i]), norm(min[i] - p[i]));
 49
 50
 51
                return result;
 53
     };
 54
 55
     struct Node {
          Point seperator;
          Rectangle rectangle; int child[2];
 57
 58
 59
 60
           void reset(const Point &p) {
                seperator = p;
rectangle = Rectangle();
rectangle.add(p);
 61
 62
 63
 64
                child[0] = child[1] = 0;
 65
     } tree[N << 1];</pre>
 66
 67
     int size, pivot;
 68
 69
     bool compare(const Point &a, const Point &b) {
   if (a[pivot] != b[pivot]) {
 70
 71
 72
                return a[pivot] < b[pivot];
 73
74
          return a.id < b.id;
 75
 76
 77
     // 左閉右開: build(1, n + 1)
     int build(int 1, int r, int type = 1) {
 78
 79
          pivot = type;
if (1 >= r) {
 80
 81
 82
          int x = ++size;
int mid = 1 + r >> 1;
 83
 84
 85
           std::nth_element(point + 1, point + mid, point + r, compare);
          tree[x].reset(point[mid]);
for (int i = 1; i < r; ++i) {
    tree[x].rectangle.add(point[i]);</pre>
 86
 87
 88
 89
          tree[x].child[0] = build(1, mid, type ^ 1);
tree[x].child[1] = build(mid + 1, r, type ^ 1);
 90
 91
 92
          return x;
 93
 94
     int insert(int x, const Point &p, int type = 1) {
 95
          pivot = type;
if (x == 0) {
 97
                tree[++size].reset(p);
 98
 99
                return size:
100
101
           tree[x].rectangle.add(p);
          if (compare(p, tree[x].seperator)) {
   tree[x].child[0] = insert(tree[x].child[0], p, type ^ 1);
102
103
          } else {
104
105
                tree[x].child[1] = insert(tree[x].child[1], p, type ^ 1);
106
107
          return x:
108
110 // For minimum distance
111 // For maximum:下面递归query时0, 1 换顺序;< and >;min and max
112 void query(int x, const Point &p, std::pair<long long, int> &answer, int type = 1) {
```

```
113
         pivot = type
         if (x == 0] \mid | tree[x].rectangle.dist(p) > answer.first) {
114
115
             return;
116
117
         answer = std::min(answer
         118
119
120
121
         } else {
123
124
             query(tree[x].child[1], p, answer, type
              query(tree[x].child[0], p, answer, type ^ 1);
125
126
127
128 std::priority_queue<std::pair<long long, int> > answer;
129
130
    void query(int x, const Point &p, int k, int type = 1) {
         pivot = type;
if (x == 0 || (int)answer.size() == k && tree[x].rectangle.dist(p) > answer.top().first) {
131
132
133
134
135
         answer.push(std::make\_pair(dist(tree[x].seperator,\ p),\ tree[x].seperator.id));
         if ((int)answer.size() > k) {
    answer.pop();
136
137
138
         if (compare(p, tree[x].seperator)) {
   query(tree[x].child[0], p, k, type ^ 1);
   query(tree[x].child[1], p, k, type ^ 1);
139
140
141
142
          else {
             query(tree[x].child[1], p, k, type ^ 1);
query(tree[x].child[0], p, k, type ^ 1);
143
144
145
146
```

3.2 KD 树-gwx

```
struct Point
 3
        double x, y;
        int id;
       Point operator - (const Point &a) const {
 5
            return (Point){x - a.x, y - a.y, id};
 6
 8
   } b[maxn],
              c[maxn];
 9
   struct node
10
12
       int ch[2];
13
   } a[maxn];
14
   struct rev
15
       int id;
double dis;
16
17
        bool operator < (const rev &a) const{
18
19
            int tmp = sign(dis - a.dis);
20
21
            if(tmp)
            return tmp < 0;
return id < a.id;</pre>
23
24
25
   typedef pr priority_queue <rev>;
   pr p0;
27
   int build(int 1, int r, int f)
28
29
30
       if(1 > r)
       return 0;
int x = (1 + r) >> 1;
if(f == 0)
31
32
33
            nth_element(a + 1, a + x, a + r + 1, cmp0); //按x排序
34
35
       36
37
38
39
       return x:
40
41
42
   void update(pr &a, rev x)
43
44
        if(a.size() < K)
       a.push(x);
else if(x < a.top())</pre>
45
46
47
48
            a.pop();
49
            a.push(x);
50
51
52
   pr merge(pr a, pr b)
53
54
        int s1 = a.size(), s2 = b.size();
55
56
        if(s1 < s2)
57
58
            while(!a.empty())
59
                update(b, a.top());
60
61
                a.pop();
62
63
            return b:
       }
64
```

```
65
66
67
              while(!b.empty())
68
69
                   update(a, b.top());
70
                   b.pop();
71
72
             return a;
73
75
76
77
   pr query(int u, Point x, int f)
78
        if(!u)
        return p0;//empty priority_queue
int d = (dis(a[a[u].ch[0]].p, x) > dis(a[a[u].ch[1]].p, x));
79
80
81
         double dx;
        pr res = query(a[u].ch[d], x, f ^ 1);
update(res, (rev){a[u].p.id, dis(a[u].p, x)});
82
83
84
         if(f == 0)
             dx = abs(x.x - a[u].p.x);
85
        dx = abs(x.y - a[u].p.y);
if(dx > res.top().dis)
87
88
        return res;
res = merge(res, query(a[u].ch[d ^ 1], x, f ^ 1));
89
91
         return res;
92
   }
93
   pr solve(Point p)
                           //离p最近的K个点
94
95
96
         int root = build(1, tot, 0);
97
         return query(root, p, 0);
98
```

3.3 lct-gwx

```
int pa[maxn], st[maxn];
struct node
 3
 4
          int ch[2], pa;
          ll s, w, sw; //s: size of subtree; w: value; sw: sum of value ll m, p; //tags of addition and multiplication bool f; //tag of flip
 5
6
 8
 9
10
    void flip(int u)
11
12
          if(!u) return;
          swap(a[u].ch[0], a[u].ch[1]);
a[u].f ^= 1;
13
14
15
17
    void add(int u, int c)
18
19
          if(!u) return;
          (a[u].p += c) %= mod;
(a[u].w += c) %= mod;
(a[u].sw += a[u].s * c % mod) % mod;
20
21
22
23
25
    void mult(int u, int c)
26
27
          if(!u) return;
28
          if(a[u].m == -1) a[u].m = c;
          lir(a[u].m == -1) a[u].m =
else (a[u].m *= c) %= mod;
(a[u].p *= c) %= mod;
(a[u].w *= c) %= mod;
(a[u].sw *= c) %= mod;
29
30
31
32
34
35
    void pushdown(int u)
36
37
          if(!u) return;
          int l = a[u].ch[0], r = a[u].ch[1];
if(a[u].m != -1)
38
39
40
41
                mult(1, a[u].m); mult(r, a[u].m);
a[u].m = -1;
42
43
          if(a[u].p)
44
45
46
                add(1, a[u].p); add(r, a[u].p);
47
                a[u].p = 0;
48
49
          if(a[u].f)
50
                flip(1); flip(r); a[u].f ^= 1;
51
52
53
54
55
    void maintain(int u)
56
57
58
          pushdown(u);
          int l = a[u].ch[0], r = a[u].ch[1];
a[u].s = a[1].s + a[r].s + 1;
60
          a[u].sw = (a[1].sw + a[r].sw + a[u].w) % mod;
61
62
64 void rotate(int u)
```

```
65
           int x = a[u].pa, y = a[x].pa, d = (a[x].ch[1] == u);
if(!y) pa[u] = pa[x], pa[x] = 0;
else a[y].ch[a[y].ch[1] == x] = u;
a[x].ch[d] = a[u].ch[d ^ 1], a[a[u].ch[d ^ 1]].pa = x;
a[u].ch[d ^ 1] = x; a[x].pa = u; a[u].pa = y;
 66
 67
 68
 69
 70
71
72
           maintain(x); maintain(u);
 73
     void splay(int u)
 75
76
           int t = u:
 77
           while(a[t].pa) st[++top] = t, t = a[t].pa;
 78
           pushdown(t);
 79
           while(top) pushdown(st[top]), top--;
 80
           while(a[u].pa)
 81
                 int x = a[u].pa, y = a[x].pa;
if(!y) {rotate(u); return;}
if(a[x].ch[1] == u ^ a[y].ch[1] == x) rotate(u);
 82
 83
 84
 85
                 else rotate(x);
 86
                 rotate(u);
 87
 88
 89
 90
     void access(int u)
 92
           splay(u);
           if(a[u].ch[1])
    a[a[u].ch[1]].pa = 0, pa[a[u].ch[1]] = u, a[u].ch[1] = 0, maintain(u);
 93
 94
 95
           while(pa[u])
 96
 97
                 int v = pa[u];
                 splay(v);
98
                spray(v),

if (a[v].ch[1])

a[a[v].ch[1]].pa = 0, pa[a[v].ch[1]] = v, a[v].ch[1] = 0;

a[v].ch[1] = u; a[u].pa = v; pa[u] = 0;
 99
100
101
102
                 maintain(v); splay(u);
103
104
105
106
     void sroot(int u)
107
108
           access(u); flip(u);
109
110
     void get(int u, int v)
111
112
113
           sroot(u); access(v);
114
115
116
     void cut(int u, int v)
117
118
           get(u, v); a[v].ch[0] = a[u].pa = 0;
119
           maintain(v);
120
121
122
     void join(int u, int v)
123
           access(u); sroot(v);
a[u].ch[1] = v; a[v].pa = u;
124
125
126
           maintain(u);
127
```

3.4 LCT-wrz

```
struct node
   node *ch[2], *fa;
uint v, sum, k, b; int rev, siz;
}mem[N], *tot, *null, *pos[N];
 5
   void init()
 6
         null = tot = mem;
null->ch[0] = null->ch[1] = null->fa = null;
null->v = null->sum = null->b = null->rev = null->siz = 0; null->k = 1;
for(int i = 1; i <= n; i++) pos[i] = ++tot, *pos[i] = *null, pos[i]->v = pos[i]->sum = 1;
 8
 9
10
11
   int type(node *x){return x->fa->ch[1]==x?1:0;}
int isroot(node *x){return x->fa->ch[type(x)] != x;}
void mswap(node *&x, node *&y){node *t = x; x = y; y = t;}
13
14
15
    void pushup(node *x)
17
         x->sum = (x->v + x->ch[0]->sum + x->ch[1]->sum) % MOD; x->siz = (x->ch[0]->siz + x->ch[1]->siz + 1) % MOD;
18
19
20
21
    void pushdown(node *x)
22
23
         if(x->rev)
24
               25
26
27
28
29
         for(int i = 0; i <= 1; i++)
30
               x->ch[i]->v = (x->k * x->ch[i]->v % MOD + x->b) % MOD;
31
               32
33
35
```

```
x->k = 1; x->b = 0;
37
38
            void update(node *x){if(!isroot(x))update(x->fa); pushdown(x);}
39
            void rotate(node *x)
 40
                             node *f = x->fa; int d = type(x);
 41
                            factor for the first form f
 42
43
 44
 46
            void splay(node *x)
47
                            update(x);
for(; !isroot(x); )
48
 49
 50
51
                                             if(isroot(x->fa)) rotate(x);
                                            else if(type(x) == type(x->fa)) rotate(x->fa), rotate(x);
else rotate(x),rotate(x);
52
53
55
                            pushup(x);
56
            void access(node *x)
{
57
 58
                            node *tmp = null;
for(; x != null; )
59
60
61
 62
                                            splay(x);
x->ch[1] = tmp;
 63
 64
                                            pushup(x);
 65
                                           tmp = x;
x = x->fa;
 66
 67
 68
69
            void makeroot(node *x)
 70
                             access(x);
                            splay(x);
x->rev ^=
72
73
74
                            swap(x->ch[0], x->ch[1]);
76
            void link(node *x, node *y)
77
 78
                            makeroot(x):
                            x->fa = y;
80
81
            void cut(node *x, node *y)
82
83
                            makeroot(x); access(y);
                            splay(y); y->ch[0] = x->fa = null;
pushup(y);
84
85
86
```

3.5 左偏树-wrz

```
struct heap
 2
            heap *ch[2];
    int dis, siz, v;
}mem[N*2], *h[N], *null, *tot;
 5
 6
7
    heap* newheap()
            heap *p = ++tot;
*p = *null;
 8
 9
           return p;
10
11
     void init()
12
13
           null = tot = mem;
null->ch[0] = null->ch[1] = null;
14
15
16
            null->v = null->dis = null->siz
17
            for(int i = 1; i <= n; i++) h[i] = null;
18
    \label{eq:heap *merge(heap *x, heap *y) // big} \label{eq:heap *merge(heap *x, heap *y) // big}
19
21
            if(x == null) return y;
           if(x == null) return y,
if(x == null) return x;
if(x >> v < y >> v) swap(x, y);
x -> ch[1] = merge(x -> ch[1], y);
if(x -> ch[0] -> dis < x -> ch[1] -> dis) swap(x -> ch[0], x -> ch[1]);
x -> dis = x -> ch[1] -> dis + 1;
x -> siz = x -> ch[0] -> siz + x -> ch[1] -> siz + 1;
22
23
24
25
26
27
28
            return x;
29
30
    heap *pop(heap *x){return merge(x->ch[0], x->ch[1]);}
31
    int main()
32
33
            heap *a = newheap(); a->siz = 1; a->v = 233;
heap *b = newheap(); b->siz = 1; b->v = 233;
34
35
            heap *c = merge(a, b);
36
```

3.6 splay-gwx

```
1 struct node
2 {
3     int pa, ch[2], s, f;
4 }a[maxn];
5     void flip(int u)
```

```
a[u].f ^= 1;
  8
          swap(a[u].ch[0], a[u].ch[1]);
  9
 10
 11
 12
     void pushdown(int k)
 13
          if(!k) return;
int l = a[k].ch[0], r = a[k].ch[1];
if(a[k].flip)
 14
 15
 16
 17
 18
          flip(1);
 19
          flip(r);
a[k].flip ^= 1;
 20
21
22
 23
     int pre(int u)
 25
 26
          u = a[u].ch[0];
 27
          while (a[u].ch[1])
u = a[u].ch[1];
 28
 29
          return u;
 30
 31
 32
     int post(int u)
 33
 34
          u = a[u].ch[1];
          while (a[u].ch[0])
u = a[u].ch[0];
 35
 36
 37
         return u;
 38
 39
 40
     void maintain(int u)
 41
          int l = a[u].ch[0], r = a[u].ch[1];
a[u].s = a[1].s + a[r].s + 1;
 42
 43
 44
 45
 46
     void rotate(int u)
 47
 48
          int x = a[u].pa, y = a[x].pa, d = (a[x].ch[1] == u);
          if(!y)
 49
 50
 51
          else
         else
a[y].ch[a[y].ch[1] == x] = u;
a[x].ch[d] = a[u].ch[d ^ 1];
a[a[u].ch[d ^ 1]].pa = x;
a[u].ch[d ^ 1] = x;
a[x].pa = u;
a[u].pa = y;
maintain(x);
maintain(x);
 52
 54
 55
 56
 57
 58
 59
          maintain(u);
 60
 61
 62
     void splay(int u, int pa) //u的父亲为pa
 63
          int t;
 64
          for(t = u; a[t].pa != pa; t = a[t].pa)
st[++top] = t;
 65
 66
 67
          pushdown(t);
 68
          for(; top; top--)
               pushdown(st[top]);
 69
                                           //pushdown the tags
 70
 71
          while(a[u].pa != pa)
 72
 73
74
          int x = a[u].pa, y = a[x].pa;
if(y == pa)
 75
76
77
78
               return;
 79
          if((a[x].ch[0] == u) ^ (a[y].ch[0] == x))
 80
               rotate(u);
 81
          else
 82
               rotate(x):
 83
          rotate(u);
 84
 85
 86
    void splay2(int u, int &g) //将u旋转到g
87
 88
 89
          while(u != g)
 90
          int x = a[u].pa, y = a[x].pa;
if(x == g)
 91
 92
 94
               rotate(u);
 95
               return;
 96
          if((a[x].ch[0] == u) ^ (a[y].ch[0] == x))
 97
 98
               rotate(u);
          else
99
100
               rotate(x);
101
          rotate(u);
102
103
104
105
    int find_kth(int u, int k)
106
107
          pushdown(u);
```

```
108
         int size = a[a[u].ch[0]].s;
109
         if(k <= size)</pre>
110
         return find_kth(a[u].ch[0], k);
111
         if(k > size + 1)
return find_kth(a[u].ch[1], k - size - 1);
112
113
         return u;
114
115
116
    int get(int 1, int r)
117
         118
119
120
121
         return a[R].ch[0];
122
    }
123
    int new_node() //recycle
124
126
         int res = q.front();
127
         q.pop();
a[res].init();
128
129
         return res;
130
    }
131
    int build(int 1, int r, int pa)
132
133
134
         if(1 > r)
         return 0;
int mid = (1 + r) >> 1, u = new_node();
135
136
137
         a[u].pa = pa;
138
         a[u].ch[0] = build(1, mid - 1, u);
a[u].ch[1] = build(mid + 1, r, u);
139
140
         maintain(u);
141
142
         return u;
143
144
145
    void recycle(int u)
146
         q.push(u);
147
         if (a[u].ch[0])
recycle(a[u].ch[0]);
148
149
150
         if(a[u].ch[1])
151
         recycle(a[u].ch[1]);
152
153
    void del(int 1, int r)
154
155
         int r = get(1, r);
recycle(a[r].ch[0]);
a[a[r].pa].ch[0] = 0;
156
157
158
         maintain(a[r].pa);
159
160
         maintain(root);
161
```

3.7 splay-wrz

```
struct node
   node *ch[2], *fa;
ll key; int siz, tag;
}mem[N*20], *tot, *null, *root;
void init()
 3
 4
 5
 6
7
          root = null = tot = mem;
null->ch[0] = null->ch[1] = null->fa = null;
null->key = null->siz = null->tag = 0;
 8
 9
10
11
    int type(node *x){return x->fa->ch[1]==x;}
node *newnode(11 key)
13
14
          node *p = ++tot; *p = *null;
p->key = key; p->siz = 1;
16
17
          return p;
18
19
    void pushup(node *x)
20
21
          x->siz = x->ch[0]->siz + x->ch[1]->siz + 1;
22
    void rotate(node *x)
24
          25
26
          x \rightarrow ch[!d] = f, f \rightarrow fa = x, pushup(f);
29
30
    void pushdown(node *x)
31
32
          if(x->tag)
33
                int &tag = x->tag;    if(x->ch[0] != null) x->ch[0]->key += tag, x->ch[0]->tag += tag;    if(x->ch[1] != null) x->ch[1]->key += tag, x->ch[1]->tag += tag;
34
35
36
                tag = 0;
37
38
39
40
    void update(node *x)
41
42
43
          if(x==null) return;
update(x->fa);
          pushdown(x);
```

```
46
    void splay(node *x, node *top)
 47
 48
         update(x);
 49
          for(;x->fa!=top;)
 50
              if(x->fa->fa == top) rotate(x);
else if(type(x) == type(x->fa)) rotate(x->fa), rotate(x);
else rotate(x), rotate(x);
 51
 52
 53
 55
         if(top == null) root = x;
         pushup(x);
 56
 57
    void insert(node *x, node *f, node *p, int d)
 59
 60
         if(x == null)
 61
 62
              p->fa = f, f->ch[d] = p;
              return;
 64
 65
         pushdown(x);
         passession(x),
if(p->key < x->key) insert(x->ch[0], x, p ,0);
else insert(x->ch[1], x, p, 1);
 66
 67
 68
         pushup(x);
 69
 70
 71
     void insert(node *p)
 72
         73
 74
 75
    node *findl(node *x){return x->ch[0] ==null?x:findl(x->ch[0]);}
 77
    node *findr(node *x){return x->ch[1]==null?x:findr(x->ch[1]);}
 78
    void insertlr()
 79
 80
          insert(newnode(-INF));
 81
         insert(newnode(INF));
 82
    void delet(node *p)
 83
 84
 85
          splay(p, null);
         spriy(p, half),
node *lp = findr(p->ch[0]), *rp = findl(p->ch[1]);
if(lp == null && rp != null) root = p->ch[1], root->fa = null;
else if(lp != null && rp == null) root = p->ch[0], root->fa = null;
else if(lp == null && rp == null) root = null;
 86
 87
 88
 89
 90
          else
 91
              splay(rp, null); splay(lp,rp);
lp->ch[1] = null; splay(lp,null);
 92
 94
 95
 96
    node* findk(node *p, int k)
 97
 98
99
              100
101
102
103
104
105
106 node* findv(node *p, int v)
107
         node* ret = null;
for(; p!=null; )
108
109
110
              pushdown(p);
if(p->key >= v) ret = p, p = p->ch[0];
111
112
              else p = p \rightarrow ch[1];
113
114
115
         splay(ret, null);
116
         return ret;
117
    void addv(node *p, int v)
118
119
120
         if(p == null) return;
         p->key += v;
p->tag += v;
121
122
123
```

3.8 treap-gwx

```
struct node
 2
 3
         int pri, val, c, s;
int ch[2];
                                          //pri: random value; c: times of showing; s: size of subtree
 5
         int cmp(int x) const {
               if(x == val) return -1;
return x < val ? 0 : 1;</pre>
 6
7
    } a[maxn];
10
    void maintain(int u) {
    a[u].s = a[u].c + a[a[u].ch[0]].s + a[a[u].ch[1]].s;
11
12
13
14
    void rotate(int &u, int d)
15
16
         int tmp = a[u].ch[d ^ 1];
a[u].ch[d ^ 1] = a[tmp].ch[d];
a[tmp].ch[d] = u;
18
19
```

```
maintain(u); maintain(tmp);
21
         u = tmp;
22
 23
 24
    void insert(int &u, int val)
25
26
         if(!u)
27
                = ++cnt;
29
              a[cnt] = (node) {rand(), val, 1, 1};
30
              return;
 31
 32
         a[u].s++;
         33
 34
35
 36
 37
38
39
    int find(int u, int val, int comp, int &res)
40
 41
         int d = a[u].cmp(val);
         if(!u) return -1;
if(d == -1) return u;
if(d == comp)
42
43
 44
 45
              if(d) res = max(res, a[u].val);
else res = min(res, a[u].val);
 46
47
48
 49
         return find(a[u].ch[d], val, comp, res);
50
51
    void remove(int &u)
52
53
         54
55
56
         else
57
              int d = a[a[u].ch[0]].pri < a[a[u].ch[1]].pri ? 0 : 1; rotate(u, d); remove(a[u].ch[d]);
59
60
61
 63
    void del(int &u, int val)
64
65
         if(find(root, val, -2, val) == -1) return;
         a[u].s--;
int d = a[u].cmp(val);
 66
 67
         if(d == -1)
 68
69
 70
              a[u].c--;
 71
             if(!a[u].c) remove(u);
 72
73
74
         else del(a[u].ch[d], val);
76
77
    int find_rank(int u, int val)
 78
         int d = a[u].cmp(val);
         if(d == -1) return 1 + a[a[u].ch[0]].s;
if(d == 0) return find_rank(a[u].ch[0], val);
return a[u].s - a[a[u].ch[1]].s + find_rank(a[u].ch[1], val);
 79
80
81
82
84
    int find_kth(int u, int k)
85
         86
87
         return a[u].val;
88
29
90
91
    int pre(int val)
93
         int ans = -inf;
         int pos = find(root, val, 1, if(pos != -1 && a[pos].ch[0])
94
                                            ans);
95
              pos = a[pos].ch[0];
while(a[pos].ch[1]) pos = a[pos].ch[1];
ans = max(ans, a[pos].val);
97
98
99
100
101
         return ans;
102
    }
103
104
    int post(int val)
105
         int ans = inf;
106
         int pos = find(root, val, 0, ans);
if(pos != -1 && a[pos].ch[1])
107
108
109
              pos = a[pos].ch[1];
while(a[pos].ch[0]) pos = a[pos].ch[0];
ans = min(ans, a[pos].val);
110
111
112
114
         return ans;
115
    //srand()
```

3.9 zkw 线段树

```
int n, M, q;
int d[N << 1];</pre>
 3
     inline void build(int n) {
 4
             for(M = 1; M < n; M <<= 1);
for(int i = M + 1; i <= M + n; i++) t[i] = in();
 8
             for(int i = M - 1; i; --i) d[i] = d[i << 1] + d[i << 1 | 1];
 9
             //max
10
             for(int i = M - 1; i; --i) d[i] = max(d[i << 1], d[i << 1 | 1]);
11
             //min
             for(int i = M - 1; i; --i) d[i] = min(d[i << 1], d[i << 1 | 1]);
12
13
14
     //单点修改
15
     void change(int x, int v) {
t[x = M + x] += v;
16
17
             while(x) d[x >>= 1] = d[x << 1] + d[x << 1 | 1];
18
19
20
21
     //区间查询
     int Sum(int s,int t,int Ans=0){
             for (s = s + M - 1, t = t + M + 1; s ^ t ^ 1; s >>= 1, t >>= 1) {
   if(~ s & 1) Ans += d[s ^ 1];
   if(t & 1) Ans += d[t ^ 1];
23
24
25
26
27
             return Ans;
28
     }
29
             void Sum(int s,int t,int L=0,int R=0){
                    for(s=s+M-1,t=t+M+1;s^t^1;s>>=1,t>>=1){
   L+=d[s],R+=d[t];
   if(~s&1) L=min(L,d[s^1]);
30
31
32
33
                            if(t&1) R=min(R,d[t^1]);
34
35
                    int res=min(L,R); while(s) res+=d[s>>=1];
36
            }
     //单点查询
//差分,当
37
    // 单点查询

// 差分,当前点的值为该点和其父节点的差值

void build(int n) {

    for(M = 1; M <= n + 1; M <<= 1);

    for(int i = M + 1; i <= M + n; i++) d[i] = in();

    for(int i = M - 1; i; --i) {

        d[i] = min(d[i << 1], d[i << 1 | 1]),

        d[i << 1] -= d[i],

        d[i << 1 | 1] -= d[i];
38
39
40
41
42
43
44
45
46
47
     }
48
     void Sum(int x, int res = 0) {
   while(x) res += d[x], x >>= 1;
49
50
51
52
53
     //区间最小(差分)
     Void Sum(int s, int t, int L = 0, int R = 0) {
    for(s = s + M - 1, t = t + M + 1; s ^ t ^ 1; s >>= 1, t >>= 1) {
        L += d[s], R += d[t];
        if(~ s & 1) L = min(L, d[s ^ 1]);
        if(t & 1) R = min(R, d[t ^ 1]);
54
55
57
58
59
             int res = min(L, R);
while(s) res += d[s >>= 1];
60
61
62
    //区间加法,维护最小值(差分)
void Add(int s, int t, int v, int A = 0) {
    for(s = s + M - 1, t = t + M + 1; s ^ t ^ 1; s
        if(~ s & 1) d[s ^ 1] += v;
        if(t & 1) d[t ^ 1] += v;
        A = min(d[s], d[s ^ 1]);
        d[s] -= A, d[s ^ 1] -= A, d[s >> 1] += A;
        A = min(d[t], d[t^1]);
        d[t] -= A, d[t ^ 1] -=A, d[t >> 1] += A;
63
64
                                                                                     t ^ 1;s >>= 1, t >>= 1) {
65
66
67
68
69
70
71
72
73
74
             while(s) {
                    A = min(d[s], d[s ^ 1]),
                    d[s] -= A,
d[s ^ 1] -= A,
d[s >>= 1] += A;
75
76
77
78
             }
79
```

4 图论

4.1 2-SAT

```
1 2-SAT的tarjan做法适用于一类如果A->B,则一定有B'->A'的对称的图。
2 一个强联通分量里的所有点,要么一起选要么一起不选,那就缩起来。
3 一个重要的结论是如果一个强联通分量里同时有A和A',则此图无解,否则一定有解。
4 无解的情况显然正确。
5 有解的情况考虑构造。每次随便从点集里抓一个点A出来,选中A的所有可达点,删去所有可达A'的点。显然这是可以做到的。6 那这样会不会把图弄成无解?考虑如果一个B->A,那么选了A及其可达点,那B选不选是不影响的。对于不可达A的显然也不影响,因此可以这样构造。
7 一个特例是存在A->A'的边,那这样选A就挂了,因此逆拓扑序 来构造才是更一般的做法。
8 因此构造方案只需对于任意一对点A,A',取dfs序大的即可。
```

4.2 上下界网络流

```
1 有源汇上下界费用流:
2 转换为求无源汇上下界最小费用可行循环流,通过T→S连边,流量上下界为(原总流量,INF)。
3
```

```
下界最小费
                                                         可行循环流
                 在原基础上再新增一个超级源点 supS, supT, 构造只有上界的网络。
对于原图的每一条边 (u, v) ,再新图中添加一条 supS→v 流量为 u, v 流量下界的边,一条 u→supT 流量为 u, v 流量下界的边,一条 u→v 流量为 u, v 流量上界-流量下界的边。
做从 supS→supT 的最小费用流,限定到达 supT 的流量为满流 (即 supS 所有出边的流量和)。此即为答案。
HINT: 原图中所有未提及的边费用都应记为 O。新图中的重新构造的边的费用等同原图中对应边的费用。
  5
  6
  7
  8
10
11 | 4.7 上下界网络流
12 | B(u, v) 表示边(u, v) 流量的下界, C(u, v) 表示边(u, v) 流量的上界, F(u, v) 表示边(u, v)
                                         \mathbf{a}, \mathbf{v}) = \mathbf{F} (\mathbf{u}, \mathbf{v}) - \mathbf{B}(\mathbf{u}, \mathbf{v}), \mathbf{L} % \mathbf{f} C(\mathbf{u}, \mathbf{v}) - \mathbf{B}(\mathbf{u}, \mathbf{v})
        的流量。设
                                G(u,
                             v)
                 G(u,
       4.7.1 无源汇的上下界可行流
15
        容量为 B(u, v); u \to T * , 容量为 B(u, v); u \to v , 容量为 C(u, v) - B(u, v) 。 最后求新网络的最大流, 判断从超级源点 S * 出发的边是否都满流即可, 边 (u, v) 的最终解中的实际流量为 G(u, v) + B(u, v) 。
       建立超级源点 S * 和超级汇点 T *
                                                                                       ,对于原图每条边 (u, v) 在新网络中连如下三条边: S * → v,
19
       G(u, v)
       4.7.2 有源汇的上下界可行流
20
       从汇点 T 到源点 S 连一条上界为 \omega,下界为 O 的边。按照无源汇的上下界可行流一样做即可,流量即为 T \to S 边上的流量。
       4.7.3 有源汇的上下界最大流
       1. 在有源汇的上下界可行流中,从汇点 T 到源点 S 的边改为连一条上界为 w,下届为 x 的
       边。x 满足二分性质,找到最大的 x 使得新网络存在无源汇的上下界可行流即为原图的最大
25
       流。
26
26 流。
27 | 2. 从汇点 T 到源点 S 连一条上界为 \omega,下界为 O 的边,变成无源汇的网络。按照无源汇的 28 | 上下界可行流的方法,建立超级源点 S * 和超级汇点 T * ,求一遍 S * T * 的最大流,再将 29 | 从汇点 T 到源点 S 的这条边拆掉,求一次 S \to T 的最大流即可。 30 | 4.7.4 有源汇的上下界可介流 31 | 1. 在有源汇的上下界可介流 T 到源点 T 到源点 T 的边改为连一条上界为 T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T * T 
        边。x 满足二分性质,找到最小的 x 使得新网络存在无源汇的上下界可行流即为原图的最小
33 | 流。
      712. 按照无源汇的上下界可行流的方法,建立超级源点 S* 与超级汇点 T* ,求一遍 S* \to T 最大流,但是注意这一次不加上汇点 T 到源点 S 的这条边,即不使之改为无源汇的网络去 | 求解。求完后,再加上那条汇点 T 到源点 S 上界 \infty 的边。因为这条边下界为 O,所以 S* ,T* 无影响,再直接求一次 S* \to T* 的最大流。若超级源点 S* 出发的边全部满流,则
                                                                                                                                                                                ,求一遍 S * → T * 的
       T→S边上的流量即为原图的最小流,否则无解。
```

4.3 矩阵树定理

```
1 C = 度数矩阵-领接矩阵
2 无向图G的生成树个数 = C的任意n - 1阶主子式(对角线的乘积)
```

4.4 边双联通-gwx

```
//G[i]: 第i个边双联通分量中有哪些点
   void tarjan(int u, int pa)
 3
4
5
        d[u] = 1[u] = ++timer;
for(int i = tail[u]; i; i = e[i].next)
 6
7
8
9
             if(!d[e[i].v])
                 st[++top] = i;
10
                  tarjan(e[i].v, u);
                 l[u] = min(l[u], l[e[
if(l[e[i].v] >= d[u])
11
                                     l[e[i].v]);
12
13
14
15
                      while(true)
16
17
                           int now = st[top--];
if(vst[e[now].u] != bcc)
18
19
20
21
                                vst[e[now].u] = bcc;
                                G[bcc].push_back(e[now].u);
22
                           if(vst[e[now].v] != bcc)
24
25
                                vst[e[now].v] = bcc;
26
                                G[bcc].push_back(e[now].v);
27
28
29
                           if(now == i) break;
                      7
30
                 }
31
             33
        }
34
35
```

4.5 帯花树

```
vector<int> link[maxn];
int n,match[maxn],Queue[maxn],head,tail;
int pred[maxn],base[maxn],start,finish,newbase;
bool InQueue[maxn],InBlossom[maxn];
void push(int u){ Queue[tail++]=u;InQueue[u]=true; }
int pop(){ return Queue[head++]; }
int FindCommonAncestor(int u,int v){
   bool InPath[maxn];
   for(int i=0;i<=,i++) InPath[i]=0;
   while(true){ u=base[u];InPath[u]=true;if(u==start) break;u=pred[match[u]]; }
   while(true){ v=base[v];if(InPath[v]) break;v=pred[match[v]]; }
   return v;
}
void ResetTrace(int u){</pre>
```

```
16
           while (base [u] !=newbase) {
17
                 v=match[u];
                 \label{lossom} \begin{tabular}{l} InBlossom[base[v]] = InBlossom[base[v]] = true; \\ u = pred[v]; \\ \end{tabular}
18
19
20
                 if(base[u]!=newbase) pred[u]=v;
21
22
23
     void BlossomContract(int u,int v){
24
          newbase=FindCommonAncestor(u,v);
25
26
27
           for (int i=0;i<n;i++)
           InBlossom[i]=0;
          ResetTrace(u); ResetTrace(v); if(base[u]!=newbase) pred[u]=v; if(base[v]!=newbase) pred[v]=u;
28
29
30
           for(int i=0;i<n;++i)</pre>
31
32
          if(InBlossom[base[i]]){
   base[i]=newbase;
33
                 if(!InQueue[i]) push(i);
34
35
36
    bool FindAugmentingPath(int u){
   bool found=false;
37
          for(int i=0;i<n;++i) pred[i]=-1,base[i]=i;
for (int i=0;i<n;i++) InQueue[i]=0;
start=u;finish=-1; head=tail=0; push(start);</pre>
38
39
40
41
           while(head<tail){
42
                 int u=pop()
                 for(int i=link[u].size()-1;i>=0;i--){
   int v=link[u][i];
   if(base[u]!=base[v]&&match[u]!=v)
43
44
45
46
                              if (v==start | | (match [v] >=0&&pred [match [v]] >=0))
47
                                   BlossomContract(u,v);
48
                              else if(pred[v]==-1){
49
                                   pred[v]=u;
                                    if(match[v]>=0) push(match[v]);
else{ finish=v; return true; }
50
51
52
                             7
53
                 }
54
55
           return found;
56
57
    void AugmentPath(){
58
           int u=finish, v, w;
59
           while (u>=0) { v=pred[u]; w=match[v]; match[v]=u; match[u]=v; u=w; }
60
    void FindMaxMatching() {
   for(int i=0;i<n;++i) match[i]=-1;</pre>
61
62
           for(int i=0;i<n;++i) if(match[i]==-1) if(FindAugmentingPath(i)) AugmentPath();</pre>
63
64
```

4.6 支配树-gwx

```
/*
 2
          build前设置n为点数, s为源点
树中的i号点对应原图的id[i]号点
 3
 4
 5
 6
7
    struct Dominator_Tree {
          int n, s, cnt;
int dfn[N], id[N], pa[N], semi[N], idom[N], p[N], mn[N];
vector<int>e[N], dom[N], be[N];
void ins(int x, int y) {e[x].push_back(y);}
void dfs(int x) {
 8
 9
10
11
12
                dfn[x] = ++cnt; id[cnt] = x;
13
                 for (int i : e[x])
                       if (!dfn[i])dfs(i), pa[dfn[i]] = dfn[x];
be[dfn[i]].push_back(dfn[x]);
14
15
16
17
          int get(int x) {
   if (p[x] != p[p[x]]) {
      if (semi[mn[x]] > semi[get(p[x])])mn[x] = get(p[x]);
18
19
20
21
                      p[x] = p[p[x]];
22
23
24
                 return mn[x];
25
          void LT() {
26
                 for (int i = cnt; i > 1; i--) {
                       (Int i = cnt; i > i; i --) {
  for (int j : be[i]) semi[i] = min(semi[i], semi[get(j)]);
  dom[semi[i]].push_back(i);
  int x = p[i] = pa[i];
  for (int j : dom[x])idom[j] = (semi[get(j)] < x ? get(j) : x);</pre>
27
28
29
30
31
                       dom[x].clear();
32
33
                for (int i = 2; i <= cnt; i++) {
   if (idom[i] != semi[i])idom[i] = idom[idom[i]];</pre>
34
35
                       dom[id[idom[i]]].push_back(id[i]);
36
37
38
          void build() {
                39
40
41
42
    };
```

4.7 支配树-wrz

```
vector<int> pre[N],bkt[N];
int dom_find(int x)
 3
 4
 5
         if(fo[x]==x) return x;
int r = dom_find(fo[x]);
 6
          if(sdom[vo[fo[x]]] < sdom[vo[x]]) vo[x] = vo[fo[x]];</pre>
 8
          return fo[x] = r;
 9
10
    int dom_eval(int x){dom_find(x); return vo[x];}
    void dom_dfs(int x)
11
12
         redfn[dfn[x]=++dtimer] = x, sdom[x] = dfn[x];
for(int i=last[x];i;i=e[i].next) if(!dfn[e[i].to])
    dom_dfs(e[i].to), fa[e[i].to] = x;
13
14
15
17
    void dom_build(int S)
18
19
          int i,x:
20
          dom_dfs(S);
21
          for(i = dtimer; i >= 2; i--)
22
23
               x = redfn[i];
24
               for(int i = 0, ii = pre[x].size(); i < ii; i++)
25
26
                    int k = pre[x][i];
if(dfn[k]) sdom[x] = min(sdom[x],sdom[dom_eval(k)]);
27
28
               bkt[redfn[sdom[x]]].push_back(x);
int fp = fa[x]; fo[x] = fa[x];
for(int i = 0, ii = bkt[fp].size(); i < ii; i++)</pre>
29
30
31
32
                     int v = bkt[fp][i];
int u = dom_eval(v);
idom[v] = sdom[u] == sdom[v]?fp:u;
34
35
36
37
               bkt[fp].clear();
         for(int i = 2;i <= dtimer; i++) x = redfn[i], idom[x] = idom[x]==redfn[sdom[x]]?idom[x]:idom[idom[x]];
for(int i = 2;i <= dtimer; i++) x = redfn[i], sdom[x] = redfn[sdom[x]];
39
40
41
    void dom_init()
43
44
         dtimer = 0;
for(int i = 1; i <= n; i++)</pre>
45
46
               dfn[i] = 0;
fo[i] = vo[i] = i;
47
48
               pre[i].clear(), bkt[i].clear();
49
51
          for(int x = 1; x <= n; x++) for(int i = last[x]; i; i = e[i].next) pre[e[i].to].push_back(x);
52
    }
53
    /*
          步骤:
54
          1. 建好原图
55
         2.dom_init() // 必须保证原图上所有的边已经连好3.dom_build(S) // S为支配树的根结点标号
56
57
          4.得到idom数组 // idom[x]表示x在支配树上的父结点, 别的数组用处不大
58
59
```

4.8 欧拉回路-wrz

```
#include<cstdio>
    #define N 100005
#define M 200005
    using namespace std;
    int last[N], ecnt = 1, cnt, ans[M], in_deg[N], out_deg[N];
    bool vis[M];
struct edge{int next,to;}e[M<<1];
void addedge(int a, int b)</pre>
 9
          e[++ecnt] = (edge){last[a], b};
10
11
          last[a] = ecnt;
12
    void dfs(int x)
14
          for(int &i = last[x]; i; i = e[i].next)
15
16
17
                int y = e[i].to, j = i;
18
                if(!vis[j>>1])
19
20
                     vis[j>>1] = 1;
21
                     dfs(y);
22
                     ans[++cnt] = j;
23
24
25
               }
         }
26
    int main()
27
28
          int t, n, m, a, b;
scanf("%d%d%d",&t,&n,&m);
for(int i = 1; i <= m; i++)</pre>
29
30
31
                scanf("%d%d",&a,&b);
32
33
               addedge(a,b);
if(t == 1)addedge(b,a), in_deg[a]++, in_deg[b]++;
else ecnt++, in_deg[b]++, out_deg[a]++;
34
36
37
          if(t == 1) // 无向
38
39
```

```
for(int i = 1; i <= n; i++)
                      if((in_deg[i]+out_deg[i]) & 1)
    return !printf("NO\n");
41
42
43
44
          else // 有向
45
                for(int i = 1; i <= n; i++)
    if(in_deg[i] != out_deg[i])
        return !printf("NO\n");</pre>
46
47
48
49
50
          dfs(a);
          if(cnt != m)
51
52
                puts("NO");
53
54
55
56
57
                puts("YES");
58
                for(int i = cnt; i; i--)
60
                      printf("%d", ans[i]&1?-(ans[i]>>1):(ans[i]>>1));
61
62
          }
63
```

4.9 Hopcoft-Karp

```
// O(sqrt(n)m)
        struct hopcoft_karp {
  int mx[MAXN], my[MAXM], lv[MAXN];
  3
  4
                   5
  6
  8
  9
10
                   lv[x] = -1; return false; }
int solve (edge_list <MAXN, MAXM> &e, int n, int m) {
    std::fill (mx, mx + n, -1); std::fill (my, my + m, -1);
    for (int ans = 0; ) {
        std::vector <int> q;
        for (int i = 0; i < n; ++i)
            if (mx[i] == -1) {
                lv[i] = 0; q.push_back (i);
        } else lv[i] = -1;
    for (int head = 0; head < (int) q.size(); ++head) {
        int x = g[head];</pre>
11
12
13
14
16
17
18
19
20
21
22
                                                      int x = q[head];
                                          for (int i = e.begin[x]; ~i; i = e.next[i]) {
    int y = e.dest[i], w = my[y];
    if (~w && lv[w] < 0) { lv[w] = lv[x] + 1; q.push_back (w); } }
int d = 0; for (int i = 0; i < n; ++i) if (!~mx[i] && dfs (e, i)) ++d;
if (d == 0) return ans; else ans += d; } };</pre>
23
```

4.10 KM-truly-n3

```
struct KM {
// Trul
 1 2
               Truly O(n^3)
           // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF, 使用时先对 1 -> n 调用 hungary
 3
          (), 再 get_ans() 求值
int w[N][N];
int v[N]
 4
5
          int lx[N],
                          ly[N], match[N], way[N], slack[N];
 6
          bool used[N];
           void init() {
 8
                for (int i = 1; i <= n; i++) {
                      match[i] = 0;
lx[i] = 0;
ly[i] = 0;
 9
10
11
12
13
                      way[i] = 0;
                }
14
15
          void hungary(int x) {
                i hungary(int x)
match[0] = x;
int j0 = 0;
for (int j = 0; j <= n; j++) {
    slack[j] = INF;
    used[j] = false;</pre>
16
17
18
20
21
22
                }
24
25
                      used[j0] = true;
                      int i0 = match[j0], delta = INF, j1 = 0;
for (int j = 1; j <= n; j++) {
   if (used[j] == false) {</pre>
26
28
29
30
                                  int cur = -w[i0][j] - lx[i0] - ly[j];
if (cur < slack[j]) {
                                        slack[j] = cur;
way[j] = j0;
31
                                  if (slack[j] < delta) {
    delta = slack[j];</pre>
33
34
                                        j1 = j;
35
                            }
37
38
                      for (int j = 0; j <= n; j++) {
    if (used[j]) {
39
                                  lx[match[j]] += delta;
41
42
                                  ly[j] -= delta;
43
                             else slack[j] -= delta;
```

```
46
47
                         j0 = j1;
                  } while (match[j0] != 0);
48
49
                        int j1 = way[j0];
match[j0] = match[j1];
j0 = j1;
50
51
52
                  } while (j0);
53
54
55
           int get_ans() {
   int sum = 0;
   for(int i = 1; i <= n; i++) {</pre>
56
57
58
                        if (w[match[i]][i] == -INF) ; // 无解
if (match[i] > 0) sum += w[match[i]][i];
59
60
61
62
                  return sum;
63
64
    } km;
```

4.11 k 短路 a 星-gwx

```
const int maxn = 1005;
    int n, m;
int S, T, K;
    int dist[maxn], cnt[maxn];
 5
    bool vst[maxn];
    vector<pair<int, int>> G[maxn], H[maxn];
                                                                     //正图&反图
 6
7
    struct node
 9
          11 d;
10
          int id;
          node(){}
11
12
          node(ll d, int id): d(d), id(id) {}
          bool operator< (const node &other) const{
   return d + dist[id] > other.d + dist[other.id];
13
14
15
16
17
    priority_queue <pair<11, int>> q;
priority_queue <node> Q;
18
19
20
21
22
    void init()
23
          for(int i = 1; i <= n; ++i)
24
                G[i].clear(), H[i].clear(), cnt[i] = 0;
25
26
27
    void dijkstra(int S)
          memset(dist, 127, sizeof(dist));
memset(vst, 0, sizeof(vst));
while(!q.empty()) q.pop();
dist[S] = 0;
29
30
31
32
          q.push(make_pair(0, S));
for(int i = 1; i <= n; ++i)</pre>
33
34
35
                if(q.empty()) break;
while(vst[q.top().second]) q.pop();
36
37
                int u = q.top().second; q.pop();
vst[u] = 1:
38
39
40
                for(auto i: H[u])
41
42
                      if(dist[i.first] > dist[u] + i.second)
43
44
                            dist[i.first] = dist[u] + i.second;
45
                            q.push(make_pair(-dist[i.first], i.first));
46
47
                }
48
49
50
51
    int solve()
52
          while(!Q.empty()) Q.pop();
Q.push(node(0, S));
53
54
          while(!Q.empty())
55
56
                auto u = Q.top(); Q.pop();
if(++cnt[u.id] > K) continue;
if(u.d + dist[u.id] > ti) continue;
if(u.id == T && cnt[T] == K)
    return u.d;
for(auto i: G[u.id])
57
59
60
61
62
63
                      Q.push(node(u.d + i.second, i.first));
64
65
          return -1:
66
```

4.12 K 短路可并堆

```
9
  10
        const int maxn=1005, maxe=10005, maxm=maxe*30; //点数, 边数, 左偏树结点数
  11
 12
       //需要用到的结构体定义
        struct A{//用来求最短路
  14
                int x,d;
                 \begin{array}{lll} A ( & \text{int } x, \text{int } d ) : x(x), d(d) \{ \} \\ \text{bool operator} < & \text{(const } A \& a) \text{const} \{ \text{return } d > a.d; \} \\ \end{array} 
  15
  16
  17
 18
 19
        struct node{//左偏树结点
  20
                int w,i,d;//i: 最后一条边的编号 d: 左偏树附加信息
 21
22
                node *lc,*rc;
node(){}
  23
                node(int w,int i):w(w),i(i),d(0){}
void refresh(){d=rc->d+1;}
  24
  25
        }null[maxm],*ptr=null,*root[maxn];
 26
        struct B{//维护答案用
 27
                int x,w;//x是结点编号,w表示之前已经产生的权值
node *rt;//这个答案对应的堆顶,注意可能不等于任何一个结点的堆
B(int x,node *rt,int w):x(x),w(w),rt(rt){}
 28
 29
 30
                 bool operator < (const B &a) const {return w+rt->w>a.w+a.rt->w;}
 31
       }:
 33
 34
       //全局变量和数组定义
 35 | vector < int > G[maxn], W[maxn], id[maxn]; //最开始要存反向图, 然后把G清空作为儿子列表
 36 bool vis[maxn], used[maxe];//used表示边是否在最短路树上37 int u[maxe],v[maxe],w[maxe];//存下每条边,注意是有向边
       int d[maxn],p[maxn];//p表示最短路树上每个点的父边
       int n,m,k,s,t;//s,t分别表示起点和终点
  40
       //以下是主函数中较关键的部分
  42 for(int i=0;i<=n;i++)root[i]=null;//一定要加上!!!
        //(读入&建反向图)
  44
       Dijkstra();
  45
        //(清空G,W,id)
  46
       for(int i=1;i<=n;i++)
                if(p[i]){
 47
                         used[p[i]]=true;//在最短路树上G[v[p[i]]].push_back(i);
 48
  49
 51
       for(int i=1;i<=m;i++){
                w[i]-=d[u[i]]-d[v[i]];//现在的w[i]表示这条边能使路径长度增加多少
  52
 53
                 if(!used[i])
 54
                         root[u[i]]=merge(root[u[i]], newnode(w[i],i));
 55
       dfs(t)
 57
       priority_queue < B > heap;
       heap.push(B(s,root[s],0));//初始状态是找贡献最小的边加进去
 58
       printf("%d\n",d[s]);//第1短路需要特判
while(--k){//其余k-1短路径用二叉堆维护
if(heap.empty())printf("-1\n");
 59
 60
 61
 62
                 else{
                         int x=heap.top().x,w=heap.top().w;
node *rt=heap.top().rt;
 63
  64
                         heap.pop();
printf("%d\n",d[s]+w+rt->w);
  65
  66
                          if (rt->lc!=null||rt->rc!=null)
 67
                          heap.push(B(x,merge(rt->lc,rt->rc),w));//pop掉当前边,换成另一条贡献大一点的边if(root[v[rt->i]]!=null)
 68
  69
  70
                                  heap.push(B(v[rt->i],root[v[rt->i]],w+rt->w));//保留当前边, 往后面再接上另一条边
  71
  72
  73
       //主函数到此结束
  74
        //Dijkstra预处理最短路 O(m\log n)
void Dijkstra(){
 75
76
  77
                memset(d,63,sizeof(d));
  78
                 d[t]=0;
                priority_queue <A>heap;
heap.push(A(t,0));
while(!heap.empty()){
  79
 80
 81
                          int x=heap.top().x;
 83
                         heap.pop();
if(vis[x])continue;
 84
                          vis[x]=true;
  85
                          for(int i=0;i<(int)G[x].size();i++)
  86
                                 if(!rut l-0;1<\lint)G[x]\lint\d[x]\lint\d[x]\lint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d[x]\rint\d
 87
 88
 89
  90
  91
                                  }
 92
                }
 93
 94
       //dfs求出每个点的堆 总计O(m\log n)
 95
       //需要调用merge, 同时递归调用自身
void dfs(int x){
 96
 98
                root[x]=merge(root[x],root[v[p[x]]]);
                for(int i=0;i<(int)G[x].size();i++)
    dfs(G[x][i]);</pre>
 99
100
101
102
       //包装过的new node() O(1)
node *newnode(int w,int i){
103
104
105
                 *++ptr=node(w,i);
                ptr->lc=ptr->rc=null;
106
```

```
107
         return ptr;
108
109
    //带可持久化的左偏树合并 总计O(\log n) //递归调用自身
110
111
112 node *merge(node *x,node *y){
         if(x==null)return y;
113
         if(y==null)return x;
if(x->w>y->w)swap(x,y);
114
115
         node *z=newnode(x->w,x->i);
z->lc=x->lc;
116
117
         z->rc=merge(x->rc,y);
if(z->lc->d>z->rc->d)swap(z->lc,z->rc);
118
119
         z->refresh();
120
121
         return z:
```

4.13 最大团

```
1 /*
   Int g[][]为图的邻接矩阵
 2
 3
        MC(V)表示点集V的最大团
        令Si={vi, vi+1, ..., vn}, mc[i] 录
倒着算mc[i], 那么显然MC(V)=mc[1]
 4
                                 vn}, mc[i]表示MC(Si)
 5
        此外有mc[i]=mc[i+1] or mc[i]=mc[i+1]+1
 6
7
   void init(){
        int i, j;
for (i=1; i<=n; ++i) for (j=1; j<=n; ++j) scanf("%d", &g[i][j]);
 9
10
11
12
   void dfs(int size){
        int i, j, k;
if (len[size] == 0)
    if (size > ans)
13
14
15
16
                  ans=size; found=true;
17
             }
18
             return;
19
20
        for (k=0; k<len[size] && !found; ++k) {
21
             if (size+len[size]-k<=ans) break;</pre>
             i=list[size][k];
if (size+mc[i]<=ans) break;
for (j=k+1, len[size+1]=0; j<len[size]; ++j)
if (g[i][list[size][j]]) list[size+1][len[size+1]++]=list[size][j];</pre>
22
23
24
25
26
27
             dfs(size+1);
        }
28
29
   void work(){
        30
31
32
33
34
             len[1]=0;
             35
36
37
             mc[i]=ans;
38
        }
39
```

4.14 SAP 网络流

```
#include < bits / stdc++.h>
    typedef long long ll;
using std::min;
 4
 5
    void read(int &digit)
 6
           digit=0;
 8
          for (c=getchar();(c<'0' || c>'9') && c!='-';c=getchar());
bool type=false;
if (c=='-')
 9
10
11
          type=true,c=getchar();
for (;c>='0' && c<='9';digit=digit*10+c-'0',c=getchar());</pre>
12
13
          if (type==true)
14
15
                 digit = - digit;
16
17
18
    #define maxn 1010
19
    const int INF=1<<30;</pre>
20
    int n,m;
int S,T;
21
22
    struct Edge
23
24
          int v,flow,next;
e[510010];
25
26
    int g[maxn],tot=1;//tot初值必须赋为1
     void addedge(int x,int y,int flow)
28
           \texttt{e}\, \big[\texttt{++tot}\big] \,.\, \texttt{v=y}\, ; \texttt{e}\, \big[\texttt{tot}\big] \,.\, \texttt{flow=flow}\, ; \texttt{e}\, \big[\texttt{tot}\big] \,.\, \texttt{next=g}\, \big[\texttt{x}\big]\, ; \texttt{g}\, \big[\texttt{x}\big] \, \texttt{=tot}\, ;
29
          e[++tot].v=x;e[tot].flow=0;e[tot].next=g[y];g[y]=tot;
30
31
    int w[maxn], hash[maxn], d[maxn]
33
    int que[maxn],pre1[maxn],pre2[maxn],p[maxn];
34
    bool vis[maxn]:
35
    int maxflow()
36
          for (int i=1;i<=n;i++) hash[i]=0,
for (int i=1;i<=n;i++) p[i]=g[i];
37
                                               hash[i]=0,d[i]=0,vis[i]=false;
38
           //hash[0]=n;
39
```

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```
40
            l=r=1;
que[1]=T; hash[0]=1; vis[T]=true;
 41
 42
 43
            while (1<=r)
 44
 45
                  int u=que[1++];
                  for (int i=g[u];i;i=e[i].next)
if ((i&1) && !vis[e[i].v])
 46
 47
 48
                        que[++r] = e[i].v;
vis[e[i].v] = true;
d[e[i].v] = d[u] + 1;
hash[d[e[i].v]] + +;
 49
 50
 51
 52
 53
 54
            for (int i=1;i<=n;i++)
if (!vis[i]) d[i]=n,hash[n]++;
 55
 56
            int flow=INF;
 58
            int ans=0;
 59
            int u=S;
while (d[S]<n)</pre>
 60
 61
 62
                  w[u]=flow;
                  bool bo=true;
for (int i=p[u];i;i=e[i].next)
if (e[i].flow && d[e[i].v]==d[u]-1)
 63
 64
 65
 66
 67
                        flow=min(flow,e[i].flow);
 68
                        p[u]=i;
                        pre1[e[i].v]=u;
 69
 70
71
72
73
                        pre2[e[i].v]=i;
                        u=e[i].v;
                        bo=false
                        if (u==T)
 74
75
76
                              ans+=flow:
                              while (u!=S)
 77
                                     e[pre2[u]].flow-=flow;
e[pre2[u]^1].flow+=flow;
 79
 80
                                    u=pre1[u];
 81
                              flow=INF;
 83
                        7
 84
                        break:
 85
                  if (!bo)
                                    continue;
 87
                  int minx=n,pos=0;
                  for (int i=g[u];i;i=e[i].next)
if (e[i].flow && d[e[i].v]<minx)
 88
 89
                                                                          minx=d[e[i].v],pos=i;
                  p[u]=pos;
                  hash [d[u]]--
 91
                  if (hash[d[u]]==0) break;
d[u]=minx+1;
 92
 93
                  hash[d[u]]++;
 95
                  if (u!=S) u=pre1[u],flow=w[u];
 96
 97
            return ans:
 98
 99
100
            int n1,n2;
read(n1),read(n2),read(m);
n=n1+n2+2;
101
102
103
            S=n1+n2+1, T=n1+n2+2;
104
105
            tot=1:
            for (int i=1;i<=n1;i++) addedge(S,i,1);
for (int i=1;i<=n2;i++) addedge(i+n1,T,1);</pre>
106
107
            while (m--)
108
109
                  int x,y;
read(x),read(y);
110
111
112
                  addedge(x,y+n1,1);
113
            int mjy=maxflow();
printf("%d\n",mjy);
for (int i=1;i<=n1;i++)</pre>
114
115
116
117
118
                  bool bo=true;
                  \begin{array}{ll} & \text{for (int $j=g[i]$; $j$; $j=e[j].next)} \\ & \text{if (!(j\&1) \&\& e[j].flow==0) } & \text{\{bo=false$; $printf("\%d_{\sqcup}",e[j].v-n1)$; $break$; }\} \\ & \text{if (bo) } & \text{printf("O_{\sqcup}")$;} \end{array}
120
121
122
            printf("\n");
123
124
125
126
127
      //求割的方案:从S开始,沿着非满流边bfs,能遍历到的地方为集合SS,其余为集合TT,横跨两个集合的边为割边
```

4.15 SPFA 判负环-wrz

```
1 int inq[N], inqt[N], dis[N];
bool SPFA()
3 {
    queue<int> q;
    for(int i = 1; i <= n; i++) dis[i] = 0, q.push(i), inq[i] = 1; // 全部入队
6    for(; !q.empty(); )
7    {
        int x = q.front(); q.pop(); inq[x] = 0;
        for(int i = last[x]; i; i = e[i].next)
10    {
```

```
e[i].to;
               if(dis[x] + e[i].val < dis[y])</pre>
12
13
                   dis[y] = dis[x] + e[i].val;
if(!inq[y])
14
15
16
17
                       if(++inqt[y] > n) return false; // 入队n次即有负环
18
                       inq[y] =
                       q.push(y);
19
20
21
22
          }
23
24
      return true:
25
26
27
       步骤:
28
       1.建好原图
29
       2.SPFA() // 若返回为true表示无负环, false表示有负环
30
31
       多次调用时记得清空inqt等数组
32
       有负环时理论复杂度是O(n^2)的
33
```

4.16 斯坦纳树

```
//N\mu , M\pm , P^{\perp} \mu const int inf = 0x3f3f3f3f3f;
    int n, m, p, status, idx[P], f[1 << P][N];
priority_queue<pair<int, int> > q; //int top, h[N];
void_dijkstra(int_dis[]) {}
    8
 9
10
11
12
13
                            \label{eq:continuous_pair} \text{q.pus} \check{\textbf{h}}(\texttt{make\_pair}(-\textbf{f}[\texttt{i}][\texttt{j}],~\texttt{j}));~//\texttt{h}[\texttt{++top}] = \texttt{j},~\texttt{vis}[\texttt{j}] = \texttt{1};
14
15
16
                dijkstra(f[i]); //SPFA(f[i]);
          }
17
18
    int main() {
    scanf("%d%d%d", &n, &m, &p);
19
20
21
22
          tatus = 1 << p;
tot = 0; memset(lst, 0, sizeof(lst));</pre>
23
            ÿ¿
; ª
                                        , Fμ Ι΄ ¼
24
25
                          μ 0
                                  т
          26
27
28
          for (int i = 1;
Steiner_Tree();
                                 i \le p; i++) f[1 << (i - 1)][idx[i]] = 0;
30
31
          int ans = inf;
for (int i = 1; i <= n; i++) ans = min(ans, f[status - 1][i]);</pre>
32
33
34
```

4.17 stoer-wagner 无向图最小割树

```
int cost[maxn] [maxn], seq[maxn], len[maxn], n, m, pop, ans;
    bool used[maxn];
    void Init(){
         int i,j,a,b,c;
for(i=0;i<n;i++) for(j=0;j<n;j++) cost[i][j]=0;
 5
         for(i=0;i<m;i++){
              scanf("%d_{\sqcup}%d_{\sqcup}%d",\&a,\&b,\&c); cost[a][b]+=c; cost[b][a]+=c;
 8
 9
         pop=n; for(i=0;i<n;i++) seq[i]=i;
10
    void Work(){
         ans=inf; int i,j,k,l,mm,sum,pk;
while(pop > 1){
    for(i=1;i<pop;i++) used[seq[i]]=0; used[seq[0]]=1;
    for(i=1;i<pop;i++) len[seq[i]]=cost[seq[0]][seq[i]];</pre>
12
13
14
15
              pk=0; mm=-inf; k=-1;
for(i=1;i<pop;i++) if(len[seq[i]] > mm){ mm=len[seq[i]]; k=i; }
for(i=1;i<pop;i++){</pre>
16
17
18
19
                    used[seq[1=k]]=1;
20
21
                    if (i==pop-2) pk=k;
if (i==pop-1) break;
22
                    mm = -inf;
23
24
25
                    mm=len[seq[j]], k=j;
26
27
              sum=0;
28
29
30
              for(i=0;i<pop;i++) if(i != k) sum+=cost[seq[k]][seq[i]];
              ans=min(ans, sum);
for(i=0;i<pop;i++)
31
                    cost[seq[k]][seq[i]]=cost[seq[i]][seq[k]]+=cost[seq[pk]][seq[i]];
32
              seq[pk]=seq[--pop];
33
         printf("%d\n",ans);
34
```

4.18 tarjan-gwx

```
1 //cut[i]: i是否为割点
    //bridge[i]: e[i]是否为桥
 3
    void dfs(int u, int pa)
 5
           d[u] = 1[u] = ++timer;
           st.push(u); vst[u] = 1;
int child = 0;
for(int i = tail[u]; i; i = e[i].next)
 6
7
 8
 9
                 if(!d[e[i].v])
10
11
                        child++
                        dfs(e[i].v, u);
l[u] = min(l[u], l[e[i].v]);
if(l[e[i].v] >= d[u])
12
13
15
                              cut[u] = 1;
if(1[e[i].v] > d[u])
    bridge[i] = 1;
16
17
18
19
                       }
20
21
           else if(vst[e[i].v]) l[u] = min(l[u], d[e[i].v]); if(!pa && child < 2) cut[u] = 0;
22
23
           if(l[u] == d[u])
24
25
                 int v; scc++;
while(true)
26
                       v = st.top(); st.pop();
id[v] = scc; vst[v] = 0; size[scc]++;
if(u == v) break;
28
29
30
31
           }
33
```

4.19 tarjan-wrz

```
void tarjan(int x) // 找割点
 2
 3
          low[x] = dfn[x] = ++timer;
          int siz = 0;
for(int i = last[x]; i; i = e[i].next)
 5
6
7
                int y = e[i].to;
if(!dfn[y])
 8
 9
                     tarjan(y); siz++;
cmin(low[x], low[y]);
if(x != 1 && low[y] >= dfn[x]) cut[x] = 1;
10
11
12
13
14
                else cmin(low[x], dfn[y]);
15
          if(x == 1 && siz > 1) cut[1] = 1;
16
18
19
    void tarjan(int x) // 有向图 缩
20
21
          dfn[x] = low[x] = ++timer; sta[++stacnt] = x; insta[x] = 1;
for(int i = last[x]; i; i = e[i].next)
22
23
24
               if(!dfn[y]) tarjan(y), low[x] = min(low[x], low[y]); // 根据不同需求适当修改else if(insta[y])low[x] = min(low[x], dfn[y]);
25
26
27
28
          if(low[x] == dfn[x])
29
               bel[x] = ++bcnt; insta[x] = 0;
for(; sta[stacnt] != x; stacnt--)
   bel[sta[stacnt]] = bcnt, insta[sta[stacnt]] = 0;
30
31
32
33
34
35
```

4.20 朱刘算法-gwx

```
//时间复杂度: O(nm)
2 int N, m;
3 int pre[maxn], in[maxn], f[maxn], id[maxn];
    struct node {int u, v, w;} a[maxm * 2]; //边表
 6
    int find(int x)
 8
          return f[x] == x ? x : f[x] = find(f[x]);
 9
10
11
    int mst()
12
          long long res = 0;
          int root = 1;
int n = N;
14
15
16
          while(true)
17
                for(int i = 1; i <= n; i++) in[i] = INT_MAX, pre[i] = 0;
for(int i = 1; i <= m; i++)
    if(a[i].u != a[i].v && in[a[i].v] > a[i].w)
        in[a[i].v] = a[i].w, pre[a[i].v] = a[i].u;
18
19
20
21
22
                for(int i = 1; i <= n; i++)
    if(in[i] == INT_MAX && i != root) return 0;</pre>
23
24
                 int cnt = 0:
```

```
for(int i = 1; i <= n; i++) f[i] = i, id[i] = 0;
                 for(int i = 1; i <= n; i++)
26
27
                      if(i == root) continue;
res += in[i];
28
29
                       if(find(i) != find(pre[i])) f[f[i]] = f[pre[i]];
30
31
                       else
32
                             cnt++;
                            for(int j = i; j && !id[j]; j = pre[j])
    id[j] = cnt;
34
35
36
                      }
37
                fif(!cnt) break;
for(int i = 1; i <= n; i++)
    if(!id[i]) id[i] = ++cnt;
for(int i = 1; i <= m; i++)</pre>
39
40
41
42
                      if(id[a[i].u] != id[a[i].v]) a[i].w -= in[a[i].v];
a[i].u = id[a[i].u];
a[i].v = id[a[i].v];
43
44
45
46
47
                root = id[root];
48
49
50
          return res;
51
```

4.21 zkw 费用流

```
//稠密图、二分图中较快,稀疏图中不如SPFA int flow, cost, price;
 4
     int dfs(int u, int f)
 5
 6
           if(u == t)
                 flow += f;
cost += price * f;
return f;
 8
 9
10
11
12
           vst[u] = 1;
           vst[u] - 1,
int used = 0;
for(int i = tail[u]; i; i = e[i].next)
    if(!vst[e[i].v] && e[i].c > 0 && e[i].w == 0)
13
14
16
                        int w = dfs(e[i].v, min(e[i].c, f - used));
e[i].c -= w; e[i ^ 1].c += w; used += w;
if(used == f) return f;
17
18
19
20
21
           return used:
22
23
     bool modlabel()
24
           int d = inf;
25
           for(int u = s; u <= t; u++)
    if(vst[u])</pre>
26
27
28
                        for(int i = tail[u]; i; i = e[i].next)
           if(e[i].c > 0 && !vst[e[i].v]) d = min(d, e[i].w);
if(d == inf) return 0;
for(int u = s; u <= t; u++)
   if(vst[u])</pre>
29
30
31
                        for(int i = tail[u]; i; i = e[i].next)
e[i].w -= d, e[i ^ 1].w += d;
33
34
           price += d;
return 1;
35
37
38
    void zkw()
39
40
           do memset(vst, 0, sizeof(vst));
  while(dfs(s, inf) > 0);
while(modlabel());
41
42
43
```

5 数论

5.1 杜教筛

```
#define N 1000005 // (10^9)^(2/3)
#define M 3333331 // hash siz
    int prime[N], notprime[N], pcnt, mu[N], pre[N];
int hash[M], nocnt; struct node{int id, f, next;}no[1000000];
int F(int n) // calculate mu[1]+mu[2]+...+mu[n]
 6
7
          if(n<N) return pre[n];</pre>
          int h = n%N; for(int i = hash[h]; i; i = no[i].next) if(no[i].id == n) return no[i].f; int ret = 1;
 8
          for(int i = 2, j; i <= n; i = j + 1)
10
11
                j = n/(n/i);
ret -= F(n/i) * (j-i+1);
12
13
14
          no[++nocnt] = (node) {n, ret, hash[h]};
15
          hash[h] = nocnt;
return ret;
16
17
19
    void init()
20
          mu[1] = 1;
21
```

```
for(int i = 2; i < N; i++)
23
                  if(!notprime[i]) prime[++pcnt] = i, mu[i] = -1;
for(int j = 1; j <= pcnt && prime[j] * i < N; j++)</pre>
24
25
26
                        notprime[prime[j] * i] = 1;
if(i % prime[j]) mu[prime[j] * i] = -mu[i];
else {mu[prime[j] * i] = 0; break;}
27
28
29
30
31
32
           for(int i = 1; i < N; i++) pre[i] = pre[i-1] + mu[i];
33
34
35
           用之前必须先init()
如果n很大,求和记得开long long
如果有取模,求和记得改取模
36
37
38
```

5.2 求逆元

```
void ex_gcd(long long a, long long b, long long &x, long long &y) {
    if (b == 0) {
        x = 1;
        y = 0;
        return;
}
long long xx, yy;
ex_gcd(b, a % b, xx, yy);
y = xx - a / b * yy;
x = yy;
}
long long inv(long long x, long long MODN) {
    long long inv_x, y;
ex_gcd(x, MODN, inv_x, y);
return (inv_x % MODN + MODN) % MODN;
}
```

5.3 直线下整点

5.4 拉格朗日插值

```
#define MOD 1000000007
int inv[N], invf[N], f[N];
int fpow(int a, int b)
 4
 5
           int r = 1:
           for(; b; b >>= 1)
                if(b & 1) r = 1ll*r*a%MOD;
a = 1ll*a*a%MOD;
 8
10
           return r;
12
13
    int la(int x, int k) // k次, 求f(x)
14
          int lim = k+2, ff = 1;
for(int i = 1; i <= lim; i++)
   ff = 111 * ff * (x-i) % N</pre>
15
16
17
18
           for(int i = 1; i <= lim; i++)
          f[i] = (f[i-1] + fpow(i, k)) % MOD; // 预处理 <math>f(1), f(2), \dots, f(lim), 注意修改 if(x \le lim) \ return \ f[x];
19
20
          int ret = 0;
for(int i = 1; i <= lim; i++)</pre>
21
23
24
                 (ret += 1ll * f[i]
                             * ff % MOD *
                             * ff % MOD * (x-i < N ? inv[x-i] : fpow(x-i, MOD-2)) % MOD // 复杂度
* invf[i-1] % MOD * invf[lim-i] % MOD * ((lim-i) % 2 ? MOD-1 : 1) % MOD
25
26
27
                ) \%= MOD;
29
          return ret;
30
31
    void init()
          inv[1] = 1;
for(int i = 2; i < N; i++) inv[i] = 111 * (MOD - MOD / i) * inv[MOD % i] % MOD;
invf[0] = 1;
for(int i = 1; i < N; i++) invf[i] = 111 * invf[i-1] * inv[i] % MOD;</pre>
33
34
35
37
38
           用之前必须先init()
39
           如果所有的逆元都能预处理就是O(n)的, 否则是O(nlogn)的
40
41
```

5.5 线性回归

```
1 // O(m^2logn)
2 // Given a[0], a[1], ..., a[m - 1]
3 // a[n] = c[0] * a[n - m] + ... + c[m - 1] * a[n - 1]
```

```
// Solve for a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1]
 5
     void linear_recurrence(long long n, int m, int a[], int c[], int p) {
   long long v[M] = {1 % p}, u[M << 1], msk = !!n;
   for(long long i(n); i > 1; i >>= 1) {
      msk <<= 1;
   }
}</pre>
 6
7
 8
 9
10
           for(long long x(0); msk; msk >>= 1, x <<= 1) {
   fill_n(u, m << 1, 0);
   int b(!!(n & msk));</pre>
11
12
13
14
                      |= b;
                  if(x < m) {
 u[x] = 1 % p;
15
16
                  }else {
17
                         for(int i(0); i < m; i++) {</pre>
18
                               for(int j(0), t(i + b); j < m; j++, t++) {
    u[t] = (u[t] + v[i] * v[j]) % p;
19
20
21
22
                         for(int i((m << 1) - 1); i >= m; i--) {
   for(int j(0), t(i - m); j < m; j++, t++) {
      u[t] = (u[t] + c[j] * u[i]) % p;</pre>
23
24
25
26
27
                         }
28
29
                  copy(u, u + m, v);
30
31
            //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
           for(int i(m); i < 2 * m; i++) {
    a[i] = 0;
32
33
                  for(int j(0); j < m; j++) {
    a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
34
35
36
37
           for(int j(0); j < m; j++) {
   b[j] = 0;
38
39
                  for(int i(0); i < m; i++) {
    b[j] = (b[j] + v[i] * a[i + j]) % p;
40
41
42
43
           for(int j(0); j < m; j++) {
    a[j] = b[j];
44
45
46
```

5.6 素数测试-gwx

```
11 multi(ll x, ll y, ll M) {
    ll res = 0;
    for(; y; y >>= 1, x = (x + x) % M)
        if(y & 1) res = (res + x) % M;
 3
 5
         return res;
 6
   11 power(11 x, 11 y, 11 p)
 8
        ll res = 1;
for(; y; y >>= 1, x = multi(x, x, p))
    if(y & 1) res = multi(res, x, p);
 9
10
11
         return res;
13
14
   int primetest(ll n, int base)
15
16
         11 n2 = n - 1, res;
17
         int s = 0;
        int s - 0,
while(!(n2 & 1)) n2 >>= 1, s++;
res = power(base, n2, n);
if(res == 1 || res == n - 1) return 1;
18
19
20
21
22
         while(s >= 0)
23
              res = multi(res, res, n);
if(res == n - 1) return 1;
25
26
              s--:
28
         return 0;
                      // n is not a strong pseudo prime
29
30
   int isprime(ll n)
31
        33
34
35
36
37
              if(n < lim[i]) return 1;</pre>
              if(!primetest(n, testNum[i])) return 0;
38
39
         return 1;
41
   ll pollard(ll n)
42
43
         11 i, x, y, p;
if(isprime(n)) return n;
45
         if(!(n & 1)) return 2;
for(i = 1; i < 20; i++)
46
47
49
                      y = func(x, n), p = gcd(y - x, n);
              while(p == 1)
50
51
52
                   x = func(x, n);
                   y = func(func(y, n), n);
```

5.7 原根-gwx

```
bool check_force(int g, int p)
 2
         int cnt = 0, prod = g;
for(int i = 1; i <= p - 1; ++i, prod = prod * g % p)
    if(prod == 1)
    if(int == 1)</pre>
 4
 5
                    if(++cnt > 1) return 0;
 6
7
         return 1;
 8
 9
   //d[]: prime divisor of (p - 1)
bool check_fast(int g, int p)
10
11
12
         13
14
15
16
         return 1;
17
18
19
    int primitive_root(int p)
20
         for(int i = 2; i < p; ++i)
    if(check(i, p)) return i;</pre>
21
22
23
```

5.8 勾股数

```
1 a=m^2-n^2, b=2mn, c=m^2+n^2
2 其中m和n中有一个是偶数,则(a,b,c)是素勾股数
```

6 字符串

6.1 AC 自动机-gwx

```
void add(int now)
 2
3
4
           int k = 0;
           for(int i = 1; i <= n; i++)
 5
 6
                int c = s[i] - 'A';
if(!ch[k][c])
    ch[k][c] = ++cnt;
 8
                k = ch[k][c];
 9
10
          ed[k] = 1;
                                 //或vector全部记录
11
12
          id[now] = k;
13
14
15
    void build()
16
17
           q.push(0);
18
           while(!q.empty())
19
                int u = q.front(), v;
q.pop();
for(int i = 0; i < m; i++)
    if(v = ch[u][i])
20
22
23
24
                             int k = pa[u];
while(k && !ch[k][i])
    k = pa[k];
if(u)
25
26
27
28
                             pa[v] = ch[k][i];
q.push(v);
29
30
31
32
                       else
33
                             ch[u][i] = ch[pa[u]][i];
```

6.2 AC 自动机-wrz

```
struct ACAM
 2
          ACAM *next[S], *fail;
   int ban;

}mem[N], *tot,

ACAM *newACAM()
                         *null, *root, *q[N];
 6
7
          ACAM *p = ++tot;
 9
          *p = *null; return p;
10
11
    void init()
12
         null = tot = mem;
for(int i = 0; i < alpha; i++) null->next[i] = null;
null->fail = null; null->ban = 0;
13
14
15
16
         root = newACAM();
17
   void inser(char *s) {
18
```

```
20
          ACAM *p = root;
          for(int i = 0; s[i]; i++)
21
22
23
               int w = s[i] - 'a';
if(p->next[w] == null) p->next[w] = newACAM();
24
25
               p = p->next[w];
26
27
          p->ban = 1;
28
    void build()
30
          31
32
33
               if(root->next[i] == null) root->next[i] = root;
else root->next[i]->fail = root, q[tail++] = root->next[i];
34
35
36
37
          for(; head < tail; head++)</pre>
38
               ACAM *p = q[head];
p->ban |= p->fail->ban;
for(int i = 0; i < alpha; i++)</pre>
39
40
41
42
                     if(p->next[i] == null) p->next[i] = p->fail->next[i];
else p->next[i]->fail = p->fail->next[i], q[tail++] = p->next[i];
43
44
45
46
          }
```

6.3 exKMP-gwx

```
void get_next()
 2
           int a = 0, p = 0;
nxt[0] = m;
 4
 5
            for(int i = 1; i < m; i++)
 6
7
                   if(i >= p || i + nxt[i - a] >= p)
 8
                         \begin{array}{lll} \mbox{if} (\mbox{i} >= \mbox{p}) & p = \mbox{i}; \\ \mbox{while} (\mbox{p} < \mbox{m} & \&\& \mbox{t} [\mbox{p}] == \mbox{t} [\mbox{p} - \mbox{i}]) & p++; \\ \mbox{nxt} [\mbox{i}] & = \mbox{p} - \mbox{i}; \end{array}
10
11
12
                         a = i;
13
14
                   else nxt[i] = nxt[i - a];
15
16
17
18
     void exkmp()
19
           int a = 0, p = 0;
get_next();
for(int i = 0; i < n; i++)</pre>
20
21
22
23
24
25
                   if(i >= p || i + nxt[i - a] >= p) // i >= p 的作用: 举个典型例子, s 和 t 无一字符相同
                         if(i >= p) p = i;
while(p < n && p</pre>
26
                                                      -i < m && s[p] == t[p - i]) p++;
28
29
                         ext[i] = p - i;
                         a = i;
30
31
                   else ext[i] = nxt[i - a];
           }
32
33
```

6.4 最小表示-wrz

```
int min_represent(char *s, int len) // 当s不是字符串时应该将char改成int等,len是s的长度,下标从0开始到n-1结束
         int i = 0, j = 1;
for(; i < len && j < len; )
3
4
5
6
7
8
9
              for (; s[(i+k)\%len] == s[(j+k)\%len] && k < len; k++); if (k == len) break; if (s[(i+k)\%len] > s[(j+k)\%len])
10
                   i += k+1;
if(i <= j) i = j + 1;
12
13
14
              }
              else
16
                     += k+1;
                   if(j \le i) j = i + 1;
17
18
19
         return i < j ? i : j;
20
21
```

6.5 最小表示-gwx

6.6 马拉车-gwx

```
//maxn = 2 * n
    void manacher(int n)
 3
 4
           int p = 0, r = 0;
for(int i = 1; i <= n; i++)
 6
7
                  if(i <= r) len[i] = min(len[2 * p - i], r - i + 1);
 8
                 else len[i] = 1;
while(b[i + len[i]] == b[i - len[i]]) len[i]++;
 9
10
                 if(i + len[i] - 1 >= r)
r = i + len[i] - 1, p = i;
11
12
13
15
    int main()
16
17
          scanf("%d\n%s", &n, a + 1);
b[++tot] = '@'; b[++tot] = '#';
for(int i = 1; i < n; i++)
    b[++tot] = a[i], b[++tot] = '#';</pre>
19
20
          b[++tot] = a[n];
b[++tot] = '#'; b[++tot] = '$';
21
22
           manacher(tot);
23
```

6.7 回文树-wrz

```
char s[N], out[N];
    struct PT
 3
 4
          PT *fail, *next[A];
    int len;
}mem[N], *tot, *null, *root1, *root0, *last;
 6
    PT *newPT()
 8
          PT *p = ++tot;
*p = *null; return p;
10
11
    void init()
12
14
          null = tot = mem;
          null->fail = null;
for(int i = 0; i < A; i++) null->next[i] = null;
null->len = 0;
15
16
18
          root1 = newPT(); root1->fail = root1; root1->len = -1;
          root0 = newPT(); root0->fail = root1; last = root1;
19
20
    int extend(int c, int i) // 返回这一次是否多了一个回文子串
21
         PT *p = last;
for(; s[i-p->len-1] != c+'a'; p = p->fail);
if(p->next[c] != null) {last = p->next[c]; return 0;}
PT *np = p->next[c] = last = newPT(); np->len = p->len + 2;
23
24
25
26
          if (p->len == -1) np->fail = root0;
28
29
               for (p=p->fail; s[i-p->len-1] != c+'a'; p = p->fail); np->fail = p->next[c];
30
31
32
33
          return 1:
34
    void main()
36
          scanf("%s",s+1); init();
for(int i = 1, ii = strlen(s+1); i <= ii; i++)
  out[i] = extend(s[i]-'a', i)?'1':'0';</pre>
37
38
39
40
          puts(out+1);
41
```

6.8 后缀数组-gwx

```
//sa[i]: 排第i的串的开头位置 rank[i]: 开头位置为i的串的排名//maxn = 2 ^{\circ} k
 3
 4
    void trans(int*s1, int*s2, int*r1, int*r2)
 5
 6
7
           for(int i = 1; i <= n; i++)
                 v[r1[s1[i]]] = i;
          V[r1[s1[1]]] = 1;
for(int i = n; i >= 1; i--)
    if(s1[i] > k)
        s2[v[r1[s1[i] - k]]--] = s1[i] - k;
for(int i = n - k + 1; i <= n; i++)
        s2[v[r1[i]]--] = i;
for(int i = 1; i <= n; i++)
    r2[s2[i]] = r2[s2[i - 1]] + (r1[s2[i]] != r1[s2[i - 1]] || r1[s2[i] + k] != r1[s2[i - 1] + k]);</pre>
 8
 9
10
12
13
14
15
16
17
    int lcp(int s, int t)
18
19
           s = rank[p][s], t = rank[p][t];
           if(s > t) swap(s, t);
21
           int k = Log[t - s + 1];
```

```
return min(f[s][k], f[t + 1 - (1 << k)][k]);
24
25
26
    void work()
27
         for(int k = 0; k <= maxk; k++)
  for(int i = 1 << k; i < (1 << k + 1) && i <= n; i++)
        Log[i] = k;
int n = 0</pre>
28
29
30
         int p = 0, q = 1;
for(int i = 1; i <= n; i++)
31
32
         v[a[i]]++;
for(int i = 1; i <= S; i++) //S:alphabet_size
v[i] += v[i - 1];</pre>
33
34
         35
37
38
39
40
41
         while(k < n)
42
43
               45
46
47
         for(int i = 1: i \le n: i++)
48
               \begin{array}{ll} h[i] &= max(h[i-1]-1,\ 0)\,;\\ int\ j &= sa[p][rank[p][i]-1]\,;\\ while(a[i+h[i]] &== a[j+h[i]]) \end{array}
49
50
51
52
                    h[i]++;
         for(int i = 2; i <= n; i++)
f[i][0] = h[i];
54
55
         for(int k = 1; k <= maxk; i++)
  for(int i = 2; i + (1 << k) - 1 <= n; i++)
    f[i][k] = min(f[i][k - 1], f[i + (1 << k - 1)][k - 1]);</pre>
56
58
59
```

6.9 后缀数组-wrz

```
1 | // 对都是数字的数组做SA时要保证数组中没有0,否则height等可能由于s[0]=s[n+1]=0出问题
           使用要保证s[0]=s[n+1]=0
   char s[N];
          t1[N],
   int n,
                   t2[N], sa[N], rank[N], sum[N], height[N], lef, rig; // 数组开两倍
   void SA_build()
6
7
        int *x = t1, *y = t2, m = 30; for(int i = 1; i <= n; i++) sum[x[i] = s[i] - 'a' + 1]++; for(int i = 1; i <= m; i++) sum[i] += sum[i-1];
9
10
        for(int i = n; i >= 1; i--) sa[sum[x[i]]--] = i;
        for(int k = 1; k <= n; k <<= 1)
11
12
            13
14
15
16
17
            for(int i = 1; i <= m; i++) sum[i] =
            for(int i = 1; i <= n; i++) sum[x[i]]++;
for(int i = 1; i <= m; i++) sum[i] += sum[i-1];
for(int i = n; i >= 1; i--) sa[sum[x[y[i]]]--] = y[i];
18
19
20
21
22
23
            24
              = x[sa[n]];
26
27
            if(m == n) break;
        for(int i = 1; i <= n; i++) rank[sa[i]] = i;
for(int i = 1, k = 0; i <= n; height[rank[i++]] = k?k--:k)
    for(; s[i+k] == s[sa[ran[i]-1]+k] && i+k <= n && sa[ran[i]-1]+k <= n; k++);</pre>
28
29
30
31
```

6.10 后缀数组 SAIS

```
// string is 0-based
                 // sa[] is 1-based
// s[n] < s[i] i = 0...n-1
     3
                  namespace SA
                  int sa[MAXN], rk[MAXN], ht[MAXN], s[MAXN << 1], t[MAX << 1], p[MAXN], cnt[MAXN], cur[MAXN];
#define pushS(x) sa[cur[s[x]]--] = x
#define pushL(x) sa[cur[s[x]]++] = x</pre>
      6
                                            fine inducedSort(v) std::fill_n(sa, n, -1)
for (int i = 0; i < n; i++) cnt[s[i]]++;\</pre>
                                                                                                                                                                                                                                                                                          -1); std::fill_n(cnt, m, 0);\
                  #define
                                            for (int i = 1; i < m; i++) cnt[i] += cnt[i-1];\
for (int i = 0; i < m; i++) cur[i] = cnt[i]-1;\
for (int i = n1-1; ~i; i--) pushS(v[i]);\
for (int i = 1; i < m; i++) cur[i] = cnt[i-1];\
for (int i = 0; i < n; i++) cur[i] = cnt[i-1];\
for (int i = 0; i < n; i++) if (sa[i] > 0 && t[sa[i]-1]) pushL(sa[i]-1);\
for (int i = 0; i < n; i++) if (sa[i] > 0 && t[sa[i]-1]) pushL(sa[i]-1);\
for (int i = 0; i < n; i++) if (sa[i] > 0 && t[sa[i]-1]) pushL(sa[i]-1);\
for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1);\
for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) if (sa[i]-1) for (int i = 0; i < n; i++) i
 10
 11
 12
 13
 14
                                           for (int i = 0; i < n; i++) if (sa[i] > 0 && t[sa[i]-1]) pushL(sa[i]-1);\
for (int i = 0; i < m; i++) cur[i] = cnt[i]-1;\
for (int i = n-1; ~i; i--) if (sa[i] > 0 && !t[sa[i]-1]) pushS(sa[i]-1)
void sais(int n, int m, int *s, int *t, int *p) {
   int n1 = t[n-1] = 0, ch = rk[0] = -1, *s1 = s+n;
   for (int i = n-2; ~i; i--) t[i] = s[i] == s[i+1] ? t[i+1] : s[i] > s[i+1]
   for (int i = 1; i < n; i++) rk[i] = t[i-1] && !t[i] ? (p[n1] = i, n1++) :
   inducedSext(n);</pre>
 15
 16
 17
 18
                                                                                                                                                                                                                                                                                                                                                                                                                                                                  s[i] > s[i+1];
20
 21
                                                                    inducedSort(p); for (int i = 0, x, y; i < n; i++) if (~(x = rk[sa[1]])), if (ch < 1 || p[x+1] - p[x] != p[y+1] - p[y]) ch++; else for (int j = p[x], k = p[y]; j <= p[x+1]; j++, k++) if ((s[j]<<1|t[j]) != (s[k]<<1|t[k])) {ch++; break;}
                                                                       inducedSort(p);
24
25
```

```
28
29
                   inducedSort(s1); }
int mapCharToInt(int n,
30
31
                                                                                      const T *str) {
32
                               int m = *std::max_element(str, str+n);
                  int m = *std::max_element(str, str+n);
    std::fill_n(rk, m+1, 0);
    for (int i = 0; i < n; i++) rk[str[i]] = 1;
    for (int i = 0; i < n; i++) rk[i+1] += rk[i];
    for (int i = 0; i < n; i++) s[i] = rk[str[i]] - 1;
    return rk[m]; }

void suffixArray(int n, const T *str) {
    int m = mapCharToInt(++n, str);
    sais(n, m, s, t, p);
    for (int i = 0; i < n; i++) rk[sa[i]] = i;
    for (int i = 0, h = ht[0] = 0; i < n-1; i++) {
        int j = sa[rk[i]-1];
    }
}</pre>
33
34
35
37
38
39
40
41
42
                                         int j = sa[rk[i]-1];
while (i+h < n && j+h < n && s[i+h] == s[j+h]) h++;
if (ht[rk[i]] = h) h--; } };</pre>
43
44
```

6.11 后缀自动机-gwx

```
int root = 1, cnt = 1, last = 1;
int pa[maxn], l[maxn], ch[maxn][maxs];
     void add(int c) //c : 0 ~ alpha_size
 5
6
7
            int np = ++cnt, p = last; last = cnt;
l[np] = x; r[np] = 1;
while(p && !ch[p][c])
 9
                   ch[p][c] = np, p = pa[p];
             if(!p)
10
11
                   pa[np] = root;
return;
12
13
14
            int q = ch[p][c];
if(1[q] == 1[p] + 1)
pa[np] = q;
15
16
17
18
             else
19
                   20
21
22
23
24
25
26
27
            }
28
30
     void get_right()
31
            for(int i = 1; i <= n; i++) add(i);
for(int i = 1; i <= cnt; i++) v[1[i]]++;
for(int i = 1; i <= n; i++) v[i] += v[i - 1];
for(int i = cnt; i; i--) t[v[1[i]]--] = i;
for(int i = cnt; i; i--) if(pa[t[i]]) r[pa[t[i]]] += r[t[i]];</pre>
32
33
34
35
36
            \mathbf{r}[1] = 0;
```

6.12 后缀自动机-wrz

```
1
    struct SAM
 2
    SAM *next[A], *fail;
int len, mi, mx;
}mem[N], *tot, *null, *root, *last, *q[N];
SAM *newSAM(int len)
 6
7
           SAM *p = ++tot;
*p = *null;
p->len = p->mi = len;
p->mx = 0;
 8
 9
10
11
12
            return p;
13
14
     void init()
15
            null = tot = mem;
for(int i = 0; i < A; i++) null->next[i] = null;
16
17
            null->fail = null;
null->len = null->mi = null->mx = 0;
18
20
            root = last = newSAM(0);
21
22
    void extend(int v)
            24
25
26
27
28
                  \begin{array}{lll} \mathtt{SAM} & *q &=& \mathtt{p->next[v]}; \\ & \mathtt{if}(\mathtt{q->len} &==& \mathtt{p->len+1}) & \mathtt{np->fail} &=& \mathtt{q}; \end{array}
29
30
31
                   else
32
33
                          SAM *nq = newSAM(p->len+1);
                         San *Inq = newSan(p->len+1),
memcpy(nq->next, q->next, sizeof(nq->next));
nq->fail = q->fail;
q->fail = np->fail = nq;
for(; p->next[v] == q && p != null; p = p->fail) p->next[v] = nq;
34
35
36
37
```

```
39 | }
40 | }
```

6.13 ex 后缀自动机-wrz

```
struct sam
 3
        sam *fail, *next[A];
   int len;
}mem[N<<1], *tot, *null, *root;</pre>
 5
   sam* newsam()
 8
         *++tot = *null;
        return tot;
10
11
   void init()
12
13
        null = tot = mem; null->fail = null; null->len = 0;
        for(int i = 0; i < A; i++) null->next[i] = null; root = newsam();
14
15
16
17
   sam* extend(sam *p, int v)
18
19
         if(p->next[v] != null)
20
21
              sam *q = p -> next[v];
              if(p->len + 1 == q->len) return q;
22
23
24
25
                  sam *nq = newsam(); *nq = *q; nq->len = p->len + 1; q->fail = nq;
26
27
                  for(; p->next[v] == q && p != null; p = p->fail) p->next[v] = nq;
28
29
                  return nq;
             }
30
31
32
33
             sam *np = newsam(); np->len = p->len + 1;
for(; p->next[v] == null && p != null; p = p->fail) p->next[v] = np;
if(p == null) np->fail = root;
34
35
36
37
              else
38
                  sam *q = p->next[v];
39
                  if(p\rightarrow len + 1 == q\rightarrow len) np\rightarrow fail = q;
40
41
                  else
                  {
                       42
43
44
45
                        for(; p->next[v] == q && p != null; p = p->fail) p->next[v] = nq;
46
47
             return np;
48
49
50
   void build_tree()
51
        for(sam *i = tot; i != mem; i--)
   addedge(i->fail - mem, i - mem);
52
53
54
```

7 其他

7.1 蔡勒公式

```
int zeller(int y,int m,int d) {
    if (m<=2) y--,m+=12; int c=y/100; y%=100;
    int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
    if (w<0) w+=7; return(w);
}</pre>
```

7.2 dancing-links

```
struct Node {
            Node *1, *r, *u, *o
int size, line_no;
                                  *u, *d, *col;
 3
 4
5
            Node() {
                  size = 0; line_no = -1;
l = r = u = d = col = NULL;
 8
    } *root;
10
     void cover(Node *c) {
            c ->l ->r = c ->r; c ->r ->l = c ->l;
for (Node *u = c ->d; u != c; u = u ->d)
11
12
                  for (Node *v = u->r; v != u; v = v->r) {
    v->d->u = v->u;
    v->u->d = v->d;
13
14
15
16
                          -- v->col->size;
17
18
19
20
     void uncover(Node *c) {
           for (Node *v = c->u; u != c; u = u->u) {
    for (Node *v = u->1; v != u; v = v->1) {
        ++ v->col->size;
        v->u->d = v;
21
22
23
24
25
                          v \rightarrow d \rightarrow u = v;
26
                  }
27
            c->1->r = c; c->r->1 = c;
```

```
29
 30
       std::vector<int> answer;
bool search(int k) {
   if (root->r == root) return true;
   Node *r = NULL;
 31
 32
 33
 34
               for (Node *u = root->r; u != root; u = u->r)
    if (r == NULL || u->size < r->size)
        r = u;
 35
 36
               if (r == NULL || r->size == 0) return false;
 38
                else {
 39
 40
                        cover(r);
                        bool succ = false;
for (Node *u = r->d; u != r && !succ; u = u->d) {
 41
 42
                               answer.push_back(u->line_no);
for (Node *v = u->r; v != u; v = v->r) // Cover row
 43
 44
 45
                                       cover(v->col);
 46
47
                                succ |= search(k + 1);
for (Node *v = u->1; v != u; v = v->1)
 48
                                        uncover(v->col);
 49
                                if (!succ) answer.pop_back();
 51
                       uncover(r):
 52
                       return succ;
 53
 54
 55
       bool entry[CR][CC];
Node *who[CR][CC];
 56
 57
       int cr, cc;
 59
       void construct() {
   root = new Node();
   Node *last = root;
   for (int i = 0; i < cc; ++ i) {
       Node *u = new Node();
       last->r = u; u->l = last;
       Node *u = new Node();
 60
 61
 62
 63
 64
 65
                       last > r = u; u -> l = last;
Node *v = u; u -> line_no = i;
last = u;
for (int j = 0; j < cr; ++ j
    if (entry[j][i]) {</pre>
 66
67
                                                           < cr; ++ j)
 68
 69
 70
71
72
73
74
75
                                        ++ u->size;
                                        Node *cur = new Node();
who[j][i] = cur;
                                        cur->line_no = j;
cur->col = u;
                                        cur - > u = v; v - > d = cur;
 76
                                        v = cur;
 77
78
                       v \rightarrow d = u; u \rightarrow u = v;
 79
                last->r = root; root->l = last;
               fast->r = root; root->l = last;
for (int j = 0; j < cr; ++ j) {
  Node *last = NULL;
  for (int i = cc - 1; i >= 0; -- i)
        if (entry[j][i]) {
            last = who[j][i];
            break;
}
 81
 82
 83
 84
 85
 86
 87
                        for (int i = 0; i < cc; ++ i)
                               (int 1 = 0; 1 \ cc, ++ 1
if (entry[j][i]) {
    last -> r = who[j][i];
    who[j][i]->1 = last;
    last = who[j][i];
 89
 90
 91
 92
 93
 94
 95
 96
       void destruct() {
    for (Node *u = root->r; u != root; ) {
        for (Node *v = u->d; v != u; ) {
            Node *nxt = v->d;
            root; }
 98
 99
100
101
                                delete(v);
102
                                v = nxt;
103
104
                       Node *nxt = u->r;
105
                       delete(u); u = nxt;
106
107
               delete root;
108
```

7.3 枚举子集

7.4 梅森旋转

```
#include <random>

int main() {
    std::mt19937 g(seed); // std::mt19937_64
    std::cout << g() << std::endl;
}</pre>
```

7.5 乘法取模

```
2 long long mult(long long x, long long y, long long MODN) {
3 long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % MODN;
4 return t < 0 ? t + MODN : t;
5 }</pre>
```

8 提示

8.1 make 支持 c++11

```
1 export CXXFLAGS='-std=c++11u-Wall' source .bashrc
```

8.2 Java

```
import java.util.*;
import java.math.*;
3
4
   public class javaNote
       static BigInteger q[] = new BigInteger [5000000]; // 定义数组的正确姿势, 记得分配内存
5
6
7
       public static void main(String[] args)
8
           long currentTime = System.currentTimeMillis(); // 获取时间, 单位是ms
10
11
12
            Scanner sc = new Scanner(System.in); // 定义输入
13
            int a = sc.nextInt(), b;
           System.out.println("integer□=□" + a); // 输出
14
           BigInteger x = new BigInteger("233"), y = new BigInteger("666");
16
           BigInteger.valueOf(1); // 将指定的表达式转化成BigInteger类型
17
18
           x.add(y); //x+y
x.subtract(y); //x-y
x.multiply(y); //x*y
19
20
21
22
           x.divide(y);
23
           x.pow(233); // x**233
24
           x.compareTo(y); // 比較x和y, x < y : -1, x = y : 0, x > y : 1
26
           BigDecimal n = new BigDecimal("233"), m = new BigDecimal("666");
27
           n.divide(m,a,RoundingMode.DOWN); //n/m并精确到小数点后第a位, a=0表示精确到个位, a为负数表示精确到小数点前
                -a+1位, 可能变成科学计数法
28
29
                取整方式
                RoundingMode.CEILING: 取右边最近的整数,即向正无穷取RoundingMode.FLOOR: 取左边最近的整数,即向负无穷取整
30
                                                           即向正无穷取整
31
                RoundingMode.DOWN: 向0取整
32
                RoundingMode.UP: 远离O取整
33
                RoundingMode.HALF_UP:上取整的四舍五入, >=0.5会进位, <0.5会舍去, 负数会先取绝对值再四舍五入再变回负数RoundingMode.HALF_DOWN:下取整的四舍五入, >0.5会进位, <=0.5会舍去, 负数原理同上RoundingMode.HALF_EVEN:分奇偶的四舍五入, >0.5会进位, <0.5会舍去, =0.5会向最近的偶数取整, 如2.5->2,
34
35
36
                     (-2.5) -> (-2)
38
39
           Math.max(a, b);//取大
40
            Math.min(a, b);//取小
41
           Math.PI;//pi
42
43
           HashSet<BigInteger> hash = new HashSet<BigInteger>(); // hash table
           hash.contains(x); // hash table中是否有a, 有则返回true, 反之返回false hash.add(x); // 把x加进hash table
44
45
46
           hash.remove(x); // 从hash table 中删去x
47
48
            Arrays.sort(arr, 1, n+1); // arr 是需要排序的数组,后两个参数分别是排序的起始位置和结束位置+1,还可以有第
                四个参数是比较函数
49
            // Arrays.sort(arr, a, b, cmp) = sort(arr+a, arr+b, cmp)
50
       }
51
```

8.3 cout 输出小数

```
1 std::cout << std::fixed << std::setprecision(5);
```

8.4 释放容器内存

```
template <typename T>
2 __inline void clear(T& container) {
3     container.clear(); // 或者删除了一堆元素
4     T(container).swap(container);
5 }
```

8.5 tuple

8.6 读入优化

```
1 // getchar()读入优化 << 关同步cin << 此优化 
2 |// 用isdigit()会小幅变慢 
3 |// 返回 false 表示读到文件尾 
4 namespace Reader {
```

9 附录-数学公式