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Reactors: A Deterministic Model for Composable Reactive Systems

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Structural Problems in CPS Design

From a embedded software programmer's perspective:

- Writing <u>deterministic concurrent</u> programs is hard
- 2. <u>Guaranteeing timing properties</u> is often impossible



Programming Model

Reactors

Model of Computation

- → Timed
- → Synchronous
- → Deterministic
- Concurrent
- → Event-driven

Establishes an unambigious relationship between <u>physical</u> and <u>logical</u> time.

Lingua Franca

Coordination Language

- → Polyglot
- → Multi-core
- → High-performance
- → Low overhead
- → Compiler + IDE

For specifying <u>deterministic</u> behaviors and imposing <u>timing</u> constraints.



https://github.com/icyphy/lingua-franca



Contributions of our CyPhy Paper

Logical time	t	
Physical time	T	
Set of tags	$\mathbb{T}=\mathbb{N}^2$	
Reactors		
Reactor instance	$r = (I, O, A, S, \mathcal{F}, \mathcal{N}, \mathcal{R}, \mathcal{G}, P, \bullet, \diamond)$	=3
Set of input ports for r	$I(r)\subseteq \Sigma$	
Set of output ports for r	$O(r)\subseteq \Sigma$	
Set of actions for r	$A(r)\subseteq \Sigma$	me)
Initialization action for r	$\bullet(r) \in A(r)$,
Termination action for r	$\diamond(r) \in A(r)$	
Set of state identifiers for r	$S(r)\subseteq {\it \Sigma}$	
Set of transformations contained in r	$\mathcal{F}(r)$	
Set of reactions contained in r	$\mathcal{N}(r)$	C
Set of contained reactors of r	$\mathcal{R}(r)$	Ť
Topology of reactors in $\mathcal{R}(r)$	$\mathcal{G}(r) \subseteq \left(\bigcup_{r' \in \mathcal{R}(r)} O(r')\right) \times \left(\bigcup_{r' \in \mathcal{R}(r)} I(r')\right)$	
Priority function	$P: \mathcal{N} \cup \mathcal{F} ightarrow \mathcal{P}$	
Reactor containing reactor r	C(r)	
Inputs and outputs		3
Input, output instance	$i,o\in {\it \Sigma}$	
Reactions dependent on $i \in I(r)$	$\mathcal{N}(i) = \{ n \in \mathcal{N}(C(i)) \mid i \in D(n) \}$	
Reactions antidependent on $o \in O(r)$	$\mathcal{N}(o) = \{ n \in \mathcal{N}(C(o)) \mid o \in D^{\vee}(n) \}$	-10

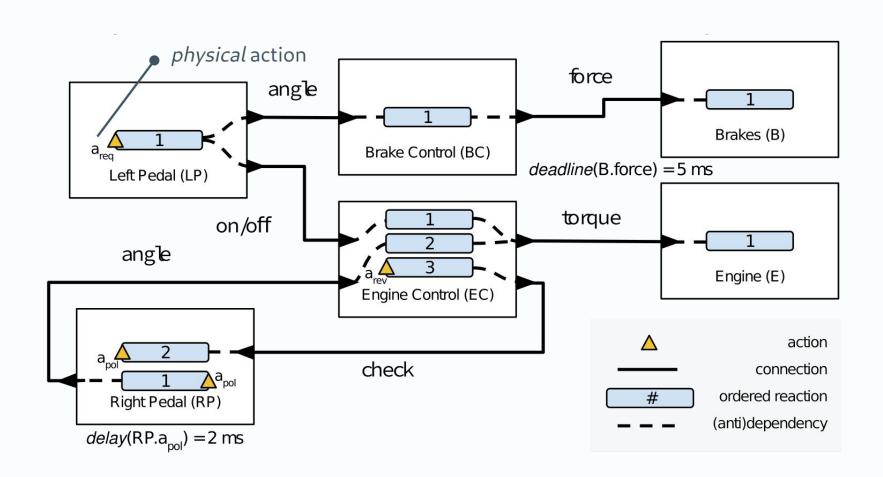


Deterministic Coordination

```
reactor Add {
  input in1:int;
                                                        ComputationC
                                                                          ActuatorA
  input in2:int;
  output out:int;
  reaction(in1, in2) -> out {=
     int result = 0;
     if (in1_is_present) {
                                                   Whether the two triggers are
       result += in1;
                                                   present simultaneously depends
                                                   only on their timestamps, not on
     if (in2_is_present) {
                                                   on the physical time of
       result += in2;
                                                   reception nor on where in the
                                                   network they are sent from.
     set(out, result);
```



Example: Simplified Power Train





Example: Simplified Power Train

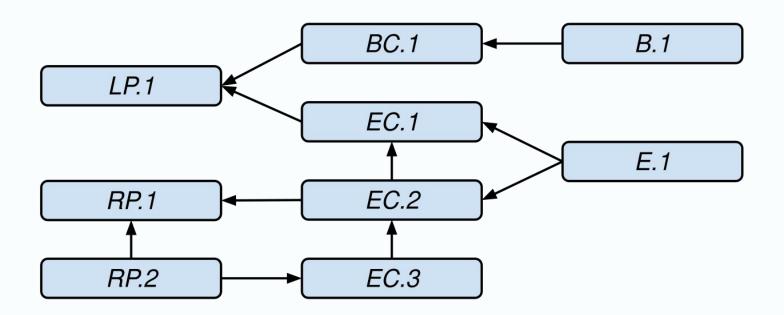
```
target C;
□ reactor EngineControl {
     input on off:bool;
     input angle:int;
     output torque:int;
     output check:bool;
     state braking:bool(true);
     timer pulse(0, 10 msec);
     reaction(on off) -> torque {=
          if(self->braking) {
              self->braking = true;
              set(torque, 0);
          } else {
              self->braking = false;
     =}
```

```
reaction(angle) -> torque {=
    if (!*(self->braking)) {
        set(torque, calc_torq(angle));
    }
=}

reaction(pulse) -> check {=
    set(check, true);
    =}
}
```



Example: Simplified Power Train



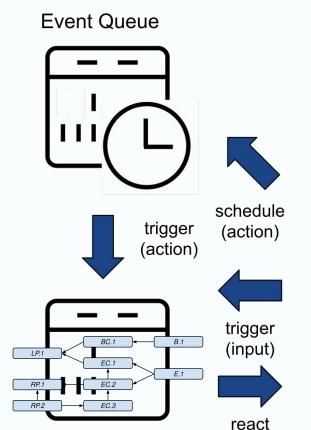


Compilation

- Parse LF code
 - Exclude reaction bodies written in target code
- Compute dependency graph
 - Reject if cyclic
- Generate code that:
 - Instantiates components
 - Connects ports
 - Incorporates verbatim reaction bodies
 - Brings declared ports and actions into reaction scope
 - Encodes dependencies between reactions
- Invoke the target compiler
 - Report possible compile errors

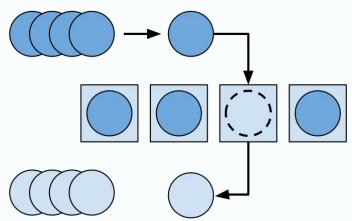


Runtime: Overview



Reaction Queue

- Handle events in tag order
- Observe dependencies between reaction
- Observe deadlines/handle violations





Runtime: Events

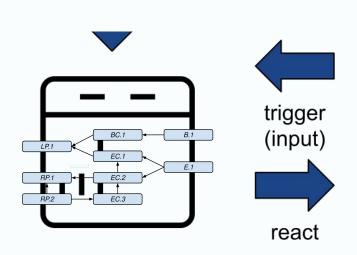
Event Queue trigger (action) schedule (action)

Events are:

- associated with an action a
- associated with a value v
- ordered by their tag g
- Logical time may advance to the next tag g when g <= T (current physical time)
- The runtime provides a schedule() function that can be called from reactions in order to enqueue new events



Runtime: Reactions

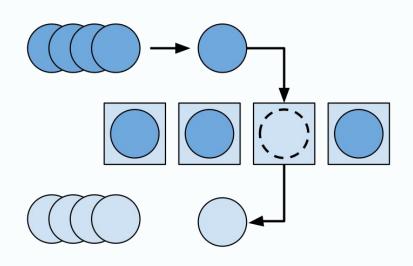


Reaction Queue

- Reactions are:
 - triggered by actions or inputs
 - able to schedule actions, but only with a tag g > t (current logical time)
- Written outputs trigger instantaneous reactions; these are added *directly* to the reaction queue



Runtime: Workers

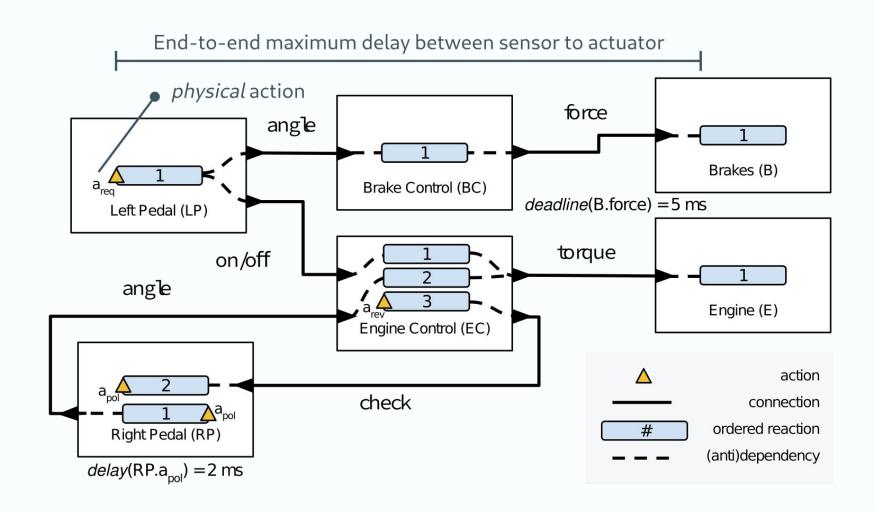


Thread Pool

- Reactions are allowed to be executed if no reactions they depend on are pending or currently executing
- We exploit concurrency transparently by mapping independent reactions to parallel worker threads
- Only a single mutex is required, and locking only occurs to coordinate between next() and schedule()



Runtime: Deadlines

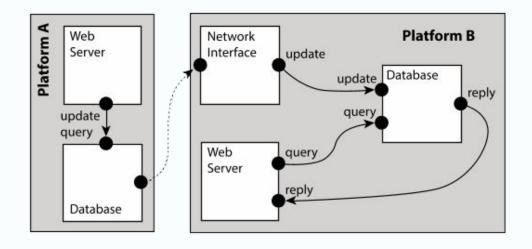




Distributed Deterministic Execution

What does the Network Interface reactor do?

- message receipt raises interrupt
- ISR schedules action at t + E + L
- reaction triggered at t + E + L writes to update port



What about the Database?

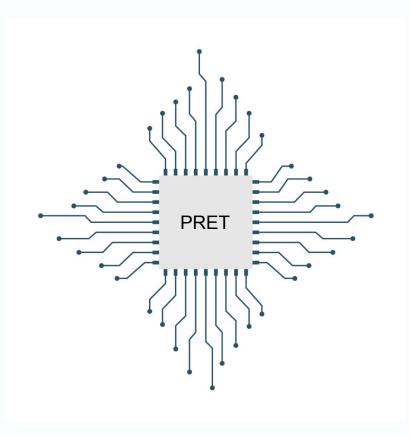
- reaction triggered by query schedules action at t + E + L
- reaction triggered at t + E + L carries out query, eventually producing reply



Pushing the Envelope

Compile-time Timing Guarantees:

- Timing instructions in the ISA
- WCET performed on reaction bodies
- Feasibility of deadlines determined statically





Relaxing synchronization requirements

Precomputation: allow reactions to events with tags greater than the current physical time

- This can be done without violating the semantics:
 - For reactions that do not depend on physical actions and are part of reactors that are free of observable physical side effects



Next Steps

- Operational Semantics
- Proof of Determinism
- Formalize Deadlines
- Implement mutations
- Modularize scheduler + implement EDF
- More efficient encoding of reaction dependencies
- Target Patmos/PRET
- Develop C++ target, TypeScript target
- Precomputation
- Develop syntax and semantics for higher-order reactor
- Integration with PTIDES
- ...



Visit our GitHub:

github.com/icyphy/lingua-franca