# Z3-Noodler - An Automata based String Solver Theory and Practice of SMT Solving

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- Cross site scripting (XSS)

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  - Some rewriting steps are removed

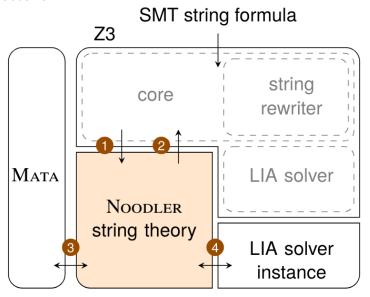


Figure 1: Z3-Noodler architecture

► Axiom Saturation

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- Supoorted String Predicates and Limitations

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- ▶ ¬contains(s, "abc") becomes  $s \notin \Sigma^*$ abc $\Sigma^*$

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- $ightharpoonup zy = z \wedge y \in a^+b^+ \wedge z \in b^+$
- ► Regular constraints enforce *UNSAT*

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- After stability, length constraints are added and fed to the LIA solver

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- Outside of this, the theory is sound but not complete

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- ▶ The same happens in the *Equations* group
- cvc5 is the best one in *Predicates Small*, while Z3-Noodler performed poorly
- Z3-Noodler could, however, be further improved by proper axiom saturation for predicates or lazy predicate evaluation

## The end

► Thanks!