第五章 高级发布-订阅模式

第三章和第四章讲述了ZMQ中请求-应答模式的一些高级用法。如果你已经能够彻底理解了,那我要说声恭喜。这一章我们会关注发布-订阅模式,使用上层模式封装,提升ZMQ发布-订阅模式的性能、可靠性、状态同步及安全机制。

本章涉及的内容有:

- 处理慢订阅者(自杀的蜗牛模式)
- 高速订阅者 (黑箱模式)
- 构建一个共享键值缓存(克降模式)

检测慢订阅者(自杀的蜗牛模式)

在使用发布-订阅模式的时候,最常见的问题之一是如何处理响应较慢的订阅者。理想状况下,发布者能以全速 发送消息给订阅者,但现实中,订阅者会需要对消息做较长时间的处理,或者写得不够好,无法跟上发布者的 脚步。

如何处理慢订阅者?最好的方法当然是让订阅者高效起来,不过这需要额外的工作。以下是一些处理慢订阅者的方法:

- **在发布者中贮存消息**。这是Gmail的做法,如果过去的几小时里没有阅读邮件的话,它会把邮件保存起来。但在高吞吐量的应用中,发布者堆积消息往往会导致内存溢出,最终崩溃。特别是当同是有多个订阅者时,或者无法用磁盘来做一个缓冲,情况就会变得更为复杂。
- **在订阅者中贮存消息**。这种做法要好的多,其实ZMQ默认的行为就是这样的。如果非得有一个人会因为内存溢出而崩溃,那也只会是订阅者,而非发布者,这挺公平的。然而,这种做法只对瞬间消息量很大的应用才合理,订阅者只是一时处理不过来,但最终会赶上进度。但是,这还是没有解决订阅者速度过慢的问题。
- **暂停发送消息**。这也是Gmail的做法,当我的邮箱容量超过7.554GB时,新的邮件就会被Gmail拒收或丢弃。这种做法对发布者来说很有益,ZMQ中若设置了阈值(HWM),其默认行为也就是这样的。但是,我们仍不能解决慢订阅者的问题,我们只是让消息变得断断续续而已。
- **断开与满订阅者的连接**。这是hotmail的做法,如果连续两周没有登录,它就会断开,这也是为什么我正在使用第十五个hotmail邮箱。不过这种方案在ZMQ里是行不通的,因为对于发布者而言,订阅者是不可见的,无法做相应处理。

看来没有一种经典的方式可以满足我们的需求,所以我们就要进行创新了。我们可以让订阅者自杀,而不仅仅是断开连接。这就是"自杀的蜗牛"模式。当订阅者发现自身运行得过慢时(对于慢速的定义应该是一个配置项,当达到这个标准时就大声地喊出来吧,让程序员知道),它会哀嚎一声,然后自杀。

订阅者如何检测自身速度过慢呢?一种方式是为消息进行编号,并在发布者端设置阈值。当订阅者发现消息编号不连续时,它就知道事情不对劲了。这里的阈值就是订阅者自杀的值。

这种方案有两个问题:一、如果我们连接的多个发布者,我们要如何为消息进行编号呢?解决方法是为每一个发布者设定一个唯一的编号,作为消息编号的一部分。二、如果订阅者使用ZMQ_SUBSRIBE选项对消息进行了过滤,那么我们精心设计的消息编号机制就毫无用处了。

有些情形不会进行消息的过滤,所以消息编号还是行得通的。不过更为普遍的解决方案是,发布者为消息标注时间戳,当订阅者收到消息时会检测这个时间戳,如果其差别达到某一个值,就发出警报并自杀。

当订阅者有自身的客户端或服务协议,需要保证最大延迟时间时,自杀的蜗牛模式会很合适。撤销一个订阅者也许并不是最周全的方案,但至少不会引发后续的问题。如果订阅者收到了过时的消息,那可能会对数据造成进一步的破坏,而且很难被发现。

以下是自杀的蜗牛模式的最简实现:

suisnail: Suicidal Snail in C

```
#include "czmq.h"
#define MAX_ALLOWED_DELAY 1000 // 毫秒
static void
subscriber (void *args, zctx_t *ctx, void *pipe)
   void *subscriber = zsocket_new (ctx, ZMQ_SUB);
   zsocket_connect (subscriber, "tcp://localhost:5556");
   while (1) {
       char *string = zstr_recv (subscriber);
        int64_t clock;
       int terms = sscanf (string, "%" PRId64, &clock);
        assert (terms == 1);
        free (string);
        if (zclock_time () - clock > MAX_ALLOWED_DELAY) {
           fprintf (stderr, "E: 订阅者无法跟进, 取消中\n");
           break;
```

```
zclock_sleep (1 + randof (2));
   zstr_send (pipe, "订阅者中止");
publisher (void *args, zctx_t *ctx, void *pipe)
   void *publisher = zsocket_new (ctx, ZMQ_PUB);
   zsocket_bind (publisher, "tcp://*:5556");
       char string [20];
       sprintf (string, "%" PRId64, zclock_time ());
       zstr_send (publisher, string);
       char *signal = zstr_recv_nowait (pipe);
       if (signal) {
           free (signal);
           break;
       zclock_sleep (1); // 等待1毫秒
   zctx_t *ctx = zctx_new ();
   void *pubpipe = zthread_fork (ctx, publisher, NULL);
   void *subpipe = zthread_fork (ctx, subscriber, NULL);
   free (zstr_recv (subpipe));
   zstr_send (pubpipe, "break");
   zclock_sleep (100);
   zctx_destroy (&ctx);
```

几点说明:

- 示例程序中的消息包含了系统当前的时间戳(毫秒)。在现实应用中,你应该使用时间戳作为消息头, 并提供消息内容。
- 示例程序中的发布者和订阅者是同一个进程的两个线程。在现实应用中,他们应该是两个不同的进程。 示例中这么做只是为了演示的方便

高速订阅者(黑箱模式)

发布-订阅模式的一个典型应用场景是大规模分布式数据处理。如要处理从证券市场上收集到的数据,可以在证券交易系统上设置一个发布者,获取价格信息,并发送给一组订阅者。如果我们有很多订阅者,我们可以使用TCP。如果订阅者到达一定的量,那我们就应该使用可靠的广播协议,如pgm。

假设我们的发布者每秒产生10万条100个字节的消息。在剔除了不需要的市场信息后,这个比率还是比较合理的。现在我们需要记录一天的数据(8小时约有250GB),再将其传入一个模拟网络,即一组订阅者。虽然10万条数据对ZMQ来说很容易处理,但我们需要更高的速度。

假设我们有多台机器,一台做发布者,其他的做订阅者。这些机器都是8核的,发布者那台有12核。

在我们开始发布消息时,有两点需要注意:

- 1. 即便只是处理很少的数据,订阅者仍有可能跟不上发布者的速度;
- 2. 当处理到6M/s的数据量时,发布者和订阅者都有可能达到极限。

首先,我们需要将订阅者设计为一种多线程的处理程序,这样我们就能在一个线程中读取消息,使用其他线程来处理消息。一般来说,我们对每种消息的处理方式都是不同的。这样一来,订阅者可以对收到的消息进行一次过滤,如根据头信息来判别。当消息满足某些条件,订阅者会将消息交给worker处理。用ZMQ的语言来说,订阅者会将消息转发给worker来处理。

这样一来,订阅者看上去就像是一个队列装置,我们可以用各种方式去连接队列装置和worker。如我们建立单向的通信,每个worker都是相同的,可以使用PUSH和PULL套接字,分发的工作就交给ZMQ吧。这是最简单也是最快速的方式:

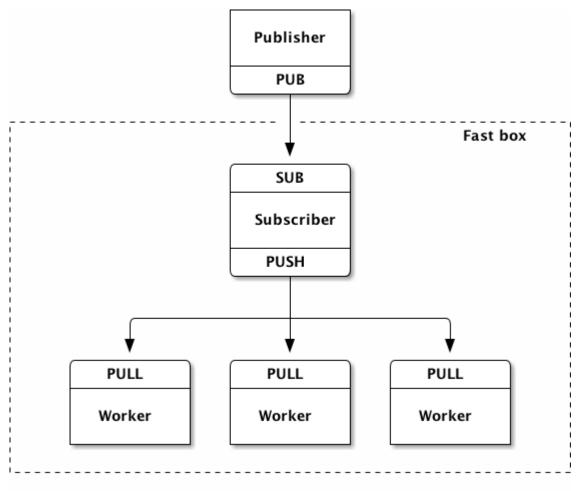


Figure 1 - Simple Black Box Pattern

订阅者和发布者之间的通信使用TCP或PGM协议,订阅者和worker的通信由于是在同一个进程中完成的,所以 使用inproc协议。

下面我们看看如何突破瓶颈。由于订阅者是单线程的,当它的CPU占用率达到100%时,它无法使用其他的核心。单线程程序总是会遇到瓶颈的,不管是2M、6M还是更多。我们需要将工作量分配到不同的线程中去,并发地执行。

很多高性能产品使用的方案是分片,就是将工作量拆分成独立并行的流。如,一半的专题数据由一个流媒体传输,另一半由另一个流媒体传输。我们可以建立更多的流媒体,但如果CPU核心数不变,那就没有必要了。 让我们看看如何将工作量分片为两个流:

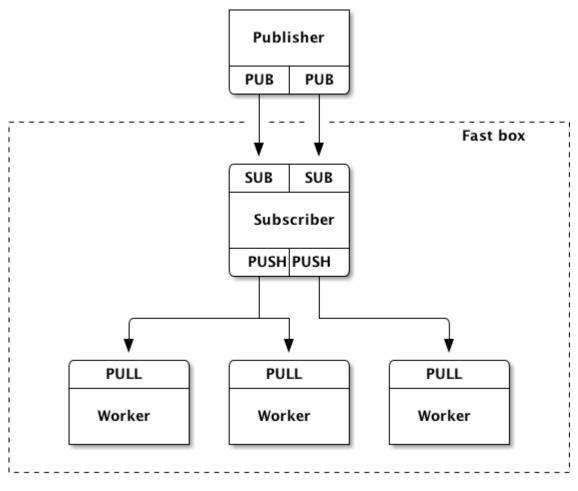


Figure 2 - Mad Black Box Pattern

要让两个流全速工作,需要这样配置ZMQ:

- 使用两个I/O线程,而不是一个;
- 使用两个独立的网络接口;
- 每个I/O线程绑定至一个网络接口;
- 两个订阅者线程,分别绑定至一个核心;
- 使用两个SUB套接字;
- 剩余的核心供worker使用;
- worker线程同时绑定至两个订阅者线程的PUSH套接字。

创建的线程数量应和CPU核心数一致,如果我们建立的线程数量超过核心数,那其处理速度只会减少。另外, 开放多个I/O线程也是没有必要的。

共享键值缓存(克隆模式)

发布-订阅模式和无线电广播有些类似,在你收听之前发送的消息你将无从得知,收到消息的多少又会取决于你的接收能力。让人吃惊的是,对于那些追求完美的工程师来说,这种机器恰恰符合他们的需求,且广为传播,成为现实生活中分发消息的最佳机制。想想非死不可、推特、BBS新闻、体育新闻等应用就知道了。

但是,在很多情形下,可靠的发布-订阅模式同样是有价值的。正如我们讨论请求-应答模式一样,我们会根据"故障"来定义"可靠性",下面几项便是发布-订阅模式中可能发生的故障:

• 订阅者连接太慢,因此没有收到发布者最初发送的消息;

- 订阅者速度太慢,同样会丢失消息;
- 订阅者可能会断开,其间的消息也会丢失。

还有一些情况我们碰到的比较少, 但不是没有:

- 订阅者崩溃、重启,从而丢失了所有已收到的消息;
- 订阅者处理消息的速度过慢,导致消息在队列中堆砌并溢出;
- 因网络过载而丢失消息(特别是PGM协议下的连接);
- 网速过慢,消息在发布者处溢出,从而崩溃。

其实还会有其他出错的情况,只是以上这些在现实应用中是比较典型的。

我们已经有方法解决上面的某些问题了,比如对于慢速订阅者可以使用自杀的蜗牛模式。但是,对于其他的问题,我们最后能有一个可复用的框架来编写可靠的发布-订阅模式。

难点在于,我们并不知道目标应用程序会怎样处理这些数据。它们会进行过滤、只处理一部分消息吗?它们是 否会将消息记录起来供日后使用?它们是否会将消息转发给其下的worker进行处理?需要考虑的情况实在太多 了,每种情况都有其所谓的可靠性。

所以,我们将问题抽象出来,供多种应用程序使用。这种抽象应用我们称之为共享的键值缓存,它的功能是通 过唯一的键名存储二进制数据块。

不要将这个抽象应用和分布式哈希表混淆起来,它是用来解决节点在分布式网络中相连接的问题的;也不要和 分布式键值表混淆,它更像是一个NoSQL数据库。我们要建立的应用是将内存中的状态可靠地传递给一组客户 端,它要做到的是:

- 客户端可以随时加入网络,并获得服务端当前的状态;
- 任何客户端都可以改变键值缓存(插入、更新、删除);
- 将这种变化以最短的延迟可靠地传达给所有的客户端;
- 能够处理大量的客户端,成百上千。

克隆模式的要点在于客户端会反过来和服务端进行通信,这在简单的发布-订阅模式中并不常见。所以我这里使用"服务端"、"客户端"而不是"发布者"、"订阅者"这两个词。我们会使用发布-订阅模式作为核心消息模式,不过还需要夹杂其他模式。

分发键值更新事件

我们会分阶段实施克隆模式。首先,我们看看如何从服务器发送键值更新事件给所有的客户端。我们将第一章中使用的天气服务模型进行改造,以键值对的方式发送信息,并让客户端使用哈希表来保存:

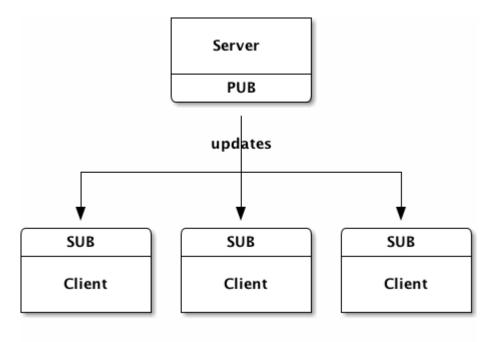


Figure 3 - Simplest Clone Model

以下是服务端代码:

clonesrv1: Clone server, Model One in C

```
#include "kvsimple.c"
   zctx_t *ctx = zctx_new ();
   void *publisher = zsocket_new (ctx, ZMQ_PUB);
   zsocket_bind (publisher, "tcp://*:5556");
   zhash_t *kvmap = zhash_new ();
   int64_t sequence = 0;
   srandom ((unsigned) time (NULL));
   while (!zctx_interrupted) {
        kvmsg_t *kvmsg = kvmsg_new (++sequence);
        kvmsg_fmt_key (kvmsg, "%d", randof (10000));
        kvmsg_fmt_body (kvmsg, "%d", randof (1000000));
        kvmsg_send
                      (kvmsg, publisher);
```

```
kvmsg_store (&kvmsg, kvmap);
}
printf (" 已中止\n已发送 %d 条消息\n", (int) sequence);
zhash_destroy (&kvmap);
zctx_destroy (&ctx);
return 0;
}
```

以下是客户端代码:

clonecli1: Clone client, Model One in C

```
#include "kvsimple.c"
   zctx_t *ctx = zctx_new ();
   void *updates = zsocket_new (ctx, ZMQ_SUB);
   zsocket_connect (updates, "tcp://localhost:5556");
   zhash_t *kvmap = zhash_new ();
   int64_t sequence = 0;
   while (TRUE) {
       kvmsg_t *kvmsg = kvmsg_recv (updates);
       if (!kvmsg)
           break;
       kvmsg_store (&kvmsg, kvmap);
       sequence++;
   printf ("已中断\n收到%d条消息\n", (int) sequence);
   zhash_destroy (&kvmap);
   zctx_destroy (&ctx);
    return 0;
```

几点说明:

• 所有复杂的工作都在kvmsg类中完成了,这个类能够处理键值对类型的消息对象,其实质上是一个ZMQ 多帧消息,共有三帧:键(ZMQ字符串)、编号(64位,按字节顺序排列)、二进制体(保存所有附加

信息)。

- 服务端随机生成消息,使用四位数作为键,这样可以模拟大量而不是过量的哈希表(1万个条目)。
- 服务端绑定套接字后会等待200毫秒,以避免订阅者连接延迟而丢失数据的问题。我们会在后面的模型中解决这一点。
- 我们使用"发布者"和"订阅者"来命名程序中使用的套接字,这样可以避免和后续模型中的其他套接字发生混淆。

以下是kvmsg的代码,已经经过了精简:

kvsimple: Key-value message class in C

```
#include "kvsimple.h"
#include "zlist.h"
#define KVMSG_KEY_MAX 255
#define FRAME_KEY
```

```
#define FRAME_SEQ
#define FRAME_BODY
#define KVMSG_FRAMES
struct _kvmsg {
   int present [KVMSG_FRAMES];
   zmq_msg_t frame [KVMSG_FRAMES];
   char key [KVMSG_KEY_MAX + 1];
kvmsg_t *
kvmsg_new (int64_t sequence)
   kvmsg_t
        *self;
   self = (kvmsg_t *) zmalloc (sizeof (kvmsg_t));
   kvmsg_set_sequence (self, sequence);
   return self;
kvmsg_free (void *ptr)
   if (ptr) {
        kvmsg_t *self = (kvmsg_t *) ptr;
        int frame_nbr;
        for (frame_nbr = 0; frame_nbr < KVMSG_FRAMES; frame_nbr++)</pre>
            if (self->present [frame_nbr])
                zmq_msg_close (&self->frame [frame_nbr]);
```

```
free (self);
kvmsg_destroy (kvmsg_t **self_p)
    assert (self_p);
    if (*self_p) {
       kvmsg_free (*self_p);
       *self_p = NULL;
kvmsg_t *
kvmsg_recv (void *socket)
    assert (socket);
    kvmsg_t *self = kvmsg_new (0);
    int frame_nbr;
    for (frame_nbr = 0; frame_nbr < KVMSG_FRAMES; frame_nbr++) {</pre>
        if (self->present [frame_nbr])
            zmq_msg_close (&self->frame [frame_nbr]);
        zmq_msg_init (&self->frame [frame_nbr]);
        self->present [frame_nbr] = 1;
        if (zmq_recvmsg (socket, &self->frame [frame_nbr], 0) == -1) {
            kvmsg_destroy (&self);
           break;
        int rcvmore = (frame_nbr < KVMSG_FRAMES - 1)? 1: 0;</pre>
        if (zsockopt_rcvmore (socket) != rcvmore) {
            kvmsg_destroy (&self);
            break;
    return self;
```

```
void
kvmsg_send (kvmsg_t *self, void *socket)
    assert (self);
    assert (socket);
    int frame_nbr;
    for (frame_nbr = 0; frame_nbr < KVMSG_FRAMES; frame_nbr++) {</pre>
        zmq_msg_t copy;
        zmq_msg_init (&copy);
        if (self->present [frame_nbr])
            zmq_msg_copy (&copy, &self->frame [frame_nbr]);
        zmq_sendmsg (socket, &copy,
            (frame_nbr < KVMSG_FRAMES - 1)? ZMQ_SNDMORE: 0);</pre>
        zmq_msg_close (&copy);
kvmsg_key (kvmsg_t *self)
    assert (self);
    if (self->present [FRAME_KEY]) {
        if (!*self->key) {
            size_t size = zmq_msg_size (&self->frame [FRAME_KEY]);
            if (size > KVMSG_KEY_MAX)
                size = KVMSG_KEY_MAX;
            memcpy (self->key,
                zmq_msg_data (&self->frame [FRAME_KEY]), size);
            self->key [size] = 0;
        return self->key;
        return NULL;
```

```
int64_t
kvmsg_sequence (kvmsg_t *self)
    assert (self);
    if (self->present [FRAME_SEQ]) {
        assert (zmq_msg_size (&self->frame [FRAME_SEQ]) == 8);
        byte *source = zmq_msg_data (&self->frame [FRAME_SEQ]);
        int64_t sequence = ((int64_t) (source [0]) << 56)</pre>
                         + ((int64_t) (source [1]) << 48)
                         + ((int64_t) (source [2]) << 40)
                         + ((int64_t) (source [3]) << 32)
                         + ((int64_t) (source [4]) << 24)
                         + ((int64_t) (source [5]) << 16)
                         + ((int64_t) (source [6]) << 8)
                         + (int64_t) (source [7]);
        return sequence;
byte *
kvmsg_body (kvmsg_t *self)
   assert (self);
    if (self->present [FRAME_BODY])
        return (byte *) zmq_msg_data (&self->frame [FRAME_BODY]);
size_t
kvmsg_size (kvmsg_t *self)
    assert (self);
    if (self->present [FRAME_BODY])
```

```
return zmq_msg_size (&self->frame [FRAME_BODY]);
    else
        return 0;
void
kvmsg_set_key (kvmsg_t *self, char *key)
   assert (self);
   zmq_msg_t *msg = &self->frame [FRAME_KEY];
   if (self->present [FRAME_KEY])
        zmq_msg_close (msg);
   zmq_msg_init_size (msg, strlen (key));
   memcpy (zmq_msg_data (msg), key, strlen (key));
   self->present [FRAME_KEY] = 1;
void
kvmsg_set_sequence (kvmsg_t *self, int64_t sequence)
   assert (self);
   zmq_msg_t *msg = &self->frame [FRAME_SEQ];
   if (self->present [FRAME_SEQ])
       zmq_msg_close (msg);
   zmq_msg_init_size (msg, 8);
   byte *source = zmq_msg_data (msg);
   source [0] = (byte) ((sequence >> 56) & 255);
   source [1] = (byte) ((sequence >> 48) & 255);
   source [2] = (byte) ((sequence >> 40) & 255);
   source [3] = (byte) ((sequence >> 32) & 255);
   source [4] = (byte) ((sequence >> 24) & 255);
   source [5] = (byte) ((sequence >> 16) & 255);
   source [6] = (byte) ((sequence >> 8) & 255);
   source [7] = (byte) ((sequence) & 255);
   self->present [FRAME_SEQ] = 1;
```

```
kvmsg_set_body (kvmsg_t *self, byte *body, size_t size)
   assert (self);
   zmq_msg_t *msg = &self->frame [FRAME_BODY];
   if (self->present [FRAME_BODY])
        zmq_msg_close (msg);
   self->present [FRAME_BODY] = 1;
   zmq_msg_init_size (msg, size);
   memcpy (zmq_msg_data (msg), body, size);
void
kvmsg_fmt_key (kvmsg_t *self, char *format, ...)
   char value [KVMSG_KEY_MAX + 1];
   va_list args;
   assert (self);
   va_start (args, format);
   vsnprintf (value, KVMSG_KEY_MAX, format, args);
   va_end (args);
   kvmsg_set_key (self, value);
kvmsg_fmt_body (kvmsg_t *self, char *format, ...)
   char value [255 + 1];
   va_list args;
   assert (self);
   va_start (args, format);
```

```
vsnprintf (value, 255, format, args);
    va_end (args);
    kvmsg_set_body (self, (byte *) value, strlen (value));
void
kvmsg_store (kvmsg_t **self_p, zhash_t *hash)
    assert (self_p);
    if (*self_p) {
        kvmsg_t *self = *self_p;
        assert (self);
        if (self->present [FRAME_KEY]
        && self->present [FRAME_BODY]) {
            zhash_update (hash, kvmsg_key (self), self);
            zhash_freefn (hash, kvmsg_key (self), kvmsg_free);
       *self_p = NULL;
void
kvmsg_dump (kvmsg_t *self)
    if (self) {
        if (!self) {
            return;
        size_t size = kvmsg_size (self);
        byte *body = kvmsg_body (self);
        fprintf (stderr, "[seq:%" PRId64 "]", kvmsg_sequence (self));
        fprintf (stderr, "[key:%s]", kvmsg_key (self));
        fprintf (stderr, "[size:%zd] ", size);
        int char_nbr;
        for (char_nbr = 0; char_nbr < size; char_nbr++)</pre>
            fprintf (stderr, "%02X", body [char_nbr]);
```

```
fprintf (stderr, "NULL message\n");
kvmsg_test (int verbose)
   kvmsg_t
        *kvmsg;
   zctx_t *ctx = zctx_new ();
   void *output = zsocket_new (ctx, ZMQ_DEALER);
   int rc = zmq_bind (output, "ipc://kvmsg_selftest.ipc");
   assert (rc == 0);
   void *input = zsocket_new (ctx, ZMQ_DEALER);
    rc = zmq_connect (input, "ipc://kvmsg_selftest.ipc");
   assert (rc == 0);
   zhash_t *kvmap = zhash_new ();
   kvmsg = kvmsg_new (1);
   kvmsg_set_key (kvmsg, "key");
   kvmsg_set_body (kvmsg, (byte *) "body", 4);
   if (verbose)
        kvmsg_dump (kvmsg);
    kvmsg_send (kvmsg, output);
    kvmsg_store (&kvmsg, kvmap);
   kvmsg = kvmsg_recv (input);
   if (verbose)
       kvmsg_dump (kvmsg);
    assert (streq (kvmsg_key (kvmsg), "key"));
    kvmsg_store (&kvmsg, kvmap);
    zhash_destroy (&kvmap);
```

```
zctx_destroy (&ctx);

printf ("OK\n");
return 0;
}
```

我们会在下文编写一个更为完整的kvmsg类,可以用到现实环境中。

客户端和服务端都会维护一个哈希表,但这个模型需要所有的客户端都比服务端启动得早,而且不能崩溃,这显然不能满足可靠性的要求。

创建快照

为了让后续连接的(或从故障中恢复的)客户端能够获取服务器上的状态信息,需要让它在连接时获取一份快照。正如我们将"消息"的概念简化为"已编号的键值对",我们也可以将"状态"简化为"一个哈希表"。为获取服务端状态,客户端会打开一个REQ套接字进行请求:

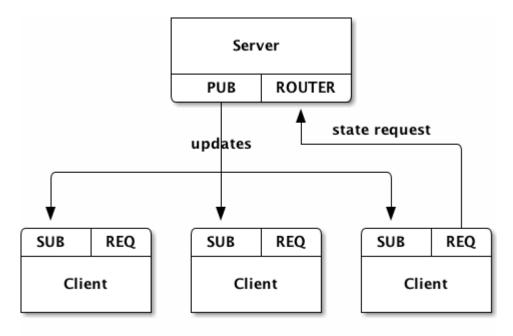


Figure 4 - State Replication

我们需要考虑时间的问题,因为生成快照是需要一定时间的,我们需要知道应从哪个更新事件开始更新快照,服务端是不知道何时有更新事件的。一种方法是先开始订阅消息,收到第一个消息之后向服务端请求"将该条更新之前的所有内容发送给"。这样一来,服务器需要为每一次更新保存一份快照,这显然是不现实的。

所以,我们会在客户端用以下方式进行同步:

- 客户端开始订阅服务器的更新事件,然后请求一份快照。这样就能保证这份快照是在上一次更新事件之后产生的。
- 客户端开始等待服务器的快照,并将更新事件保存在队列中,做法很简单,不要从套接字中读取消息就可以了,ZMQ会自动将这些消息保存起来,这时不应设置阈值(HWM)。

- 当客户端获取到快照后,它将再次开始读取更新事件,但是需要丢弃那些早于快照生成时间的事件。如 快照生成时包含了200次更新,那客户端会从第201次更新开始读取。
- 随后,客户端就会用更新事件去更新自身的状态了。

这是一个比较简单的模型,因为它用到了ZMQ消息队列的机制。服务端代码如下:

clonesrv2: Clone server, Model Two in C

```
#include "kvsimple.c"
static int s_send_single (char *key, void *data, void *args);
static void state_manager (void *args, zctx_t *ctx, void *pipe);
int main (void)
   zctx_t *ctx = zctx_new ();
   void *publisher = zsocket_new (ctx, ZMQ_PUB);
   zsocket_bind (publisher, "tcp://*:5557");
   int64_t sequence = 0;
   srandom ((unsigned) time (NULL));
   void *updates = zthread_fork (ctx, state_manager, NULL);
    free (zstr_recv (updates));
   while (!zctx_interrupted) {
       kvmsg_t *kvmsg = kvmsg_new (++sequence);
        kvmsg_fmt_key (kvmsg, "%d", randof (10000));
       kvmsg_fmt_body (kvmsg, "%d", randof (1000000));
       kvmsg_send (kvmsg, publisher);
       kvmsg_send
                      (kvmsg, updates);
       kvmsg_destroy (&kvmsg);
   printf ("已中断\n已发送 %d 条消息\n", (int) sequence);
    zctx_destroy (&ctx);
```

```
typedef struct {
   void *socket;
   zframe_t *identity; // 请求方的标识
} kvroute_t;
s_send_single (char *key, void *data, void *args)
   kvroute_t *kvroute = (kvroute_t *) args;
   zframe_send (&kvroute->identity,
       kvroute->socket, ZFRAME_MORE + ZFRAME_REUSE);
   kvmsg_t *kvmsg = (kvmsg_t *) data;
   kvmsg_send (kvmsg, kvroute->socket);
   return 0;
state_manager (void *args, zctx_t *ctx, void *pipe)
   zhash_t *kvmap = zhash_new ();
   zstr_send (pipe, "READY");
   void *snapshot = zsocket_new (ctx, ZMQ_ROUTER);
   zsocket_bind (snapshot, "tcp://*:5556");
   zmq_pollitem_t items [] = {
       { pipe, 0, ZMQ_POLLIN, 0 },
       { snapshot, 0, ZMQ_POLLIN, 0 }
   int64_t sequence = 0; // 当前快照版本
   while (!zctx_interrupted) {
       int rc = zmq_poll (items, 2, -1);
       if (rc == -1 && errno == ETERM)
           break;
        if (items [0].revents & ZMQ_POLLIN) {
           kvmsg_t *kvmsg = kvmsg_recv (pipe);
```

```
if (!kvmsg)
           break;
        sequence = kvmsg_sequence (kvmsg);
       kvmsg_store (&kvmsg, kvmap);
    if (items [1].revents & ZMQ_POLLIN) {
        zframe_t *identity = zframe_recv (snapshot);
        if (!identity)
           break;
       char *request = zstr_recv (snapshot);
       if (streq (request, "ICANHAZ?"))
            free (request);
       else {
           printf ("E: 错误的请求,程序中止\n");
           break;
       kvroute_t routing = { snapshot, identity };
       zhash_foreach (kvmap, s_send_single, &routing);
       printf ("正在发送快照,版本号 %d\n", (int) sequence);
       zframe_send (&identity, snapshot, ZFRAME_MORE);
       kvmsg_t *kvmsg = kvmsg_new (sequence);
       kvmsg_set_key (kvmsg, "KTHXBAI");
       kvmsg_set_body (kvmsg, (byte *) "", 0);
       kvmsg_send
                     (kvmsg, snapshot);
       kvmsg_destroy (&kvmsg);
zhash_destroy (&kvmap);
```

以下是客户端代码:

clonecli2: Clone client, Model Two in C

```
//
// 克隆模式 - 客户端 - 模型2
//
```

```
#include "kvsimple.c"
int main (void)
   zctx_t *ctx = zctx_new ();
   void *snapshot = zsocket_new (ctx, ZMQ_DEALER);
    zsocket_connect (snapshot, "tcp://localhost:5556");
   void *subscriber = zsocket_new (ctx, ZMQ_SUB);
    zsocket_connect (subscriber, "tcp://localhost:5557");
    zhash_t *kvmap = zhash_new ();
   int64_t sequence = 0;
    zstr_send (snapshot, "ICANHAZ?");
   while (TRUE) {
        kvmsg_t *kvmsg = kvmsg_recv (snapshot);
       if (!kvmsg)
        if (streq (kvmsg_key (kvmsg), "KTHXBAI")) {
            sequence = kvmsg_sequence (kvmsg);
           printf ("已获取快照,版本号=%d\n", (int) sequence);
           kvmsg_destroy (&kvmsg);
           break;
       kvmsg_store (&kvmsg, kvmap);
   while (!zctx_interrupted) {
        kvmsg_t *kvmsg = kvmsg_recv (subscriber);
        if (!kvmsg)
           break;
        if (kvmsg_sequence (kvmsg) > sequence) {
           sequence = kvmsg_sequence (kvmsg);
           kvmsg_store (&kvmsg, kvmap);
        else
           kvmsg_destroy (&kvmsg);
    zhash_destroy (&kvmap);
    zctx_destroy (&ctx);
```

几点说明:

- 客户端使用两个线程,一个用来生成随机的更新事件,另一个用来管理状态。两者之间使用PAIR套接字通信。可能你会考虑使用SUB套接字,但是"慢连接"的问题会影响到程序运行。PAIR套接字会让两个线程严格同步的。
- 我们在updates套接字上设置了阈值(HWM),避免更新服务内存溢出。在inproc协议的连接中,阈值是两端套接字阈值的加和,所以要分别设置。
- 客户端比较简单,用C语言编写,大约60行代码。大多数工作都在kvmsg类中完成了,不过总的来说, 克隆模式实现起来还是比较简单的。
- 我们没有用特别的方式来序列化状态内容。键值对用kvmsg对象表示,保存在一个哈希表中。在不同的时间请求状态时会得到不同的快照。
- 我们假设客户端只和一个服务进行通信,而且服务必须是正常运行的。我们暂不考虑如何从服务崩溃的情形中恢复过来。

现在,这两段程序都还没有真正地工作起来,但已经能够正确地同步状态了。这是一个多种消息模式的混合体:进程内的PAIR、发布-订阅、ROUTER-DEALER等。

重发键值更新事件

第二个模型中,键值更新事件都来自于服务器,构成了一个中心化的模型。但是我们需要的是一个能够在客户端进行更新的缓存,并能同步到其他客户端中。这时,服务端只是一个无状态的中间件,带来的好处有:

- 我们不用太过关心服务端的可靠性,因为即使它崩溃了,我们仍能从客户端中获取完整的数据。
- 我们可以使用键值缓存在动态节点之间分享数据。

客户端的键值更新事件会通过PUSH-PULL套接字传达给服务端:

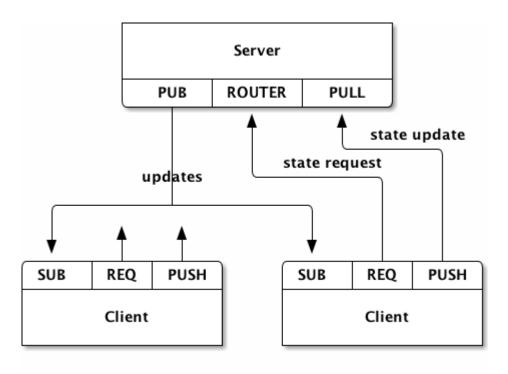


Figure 5 - Republishing Updates

我们为什么不让客户端直接将更新信息发送给其他客户端呢?虽然这样做可以减少延迟,但是就无法为更新事件添加自增的唯一编号了。很多应用程序都需要更新事件以某种方式排序,只有将消息发给服务端,由服务端分发更新消息,才能保证更新事件的顺序。

有了唯一的编号后,客户端还能检测到更多的故障:网络堵塞或队列溢出。如果客户端发现消息输入流有一段空白,它能采取措施。可能你会觉得此时让客户端通知服务端,让它重新发送丢失的信息,可以解决问题。但事实上没有必要这么做。消息流的空挡表示网络状况不好,如果再进行这样的请求,只会让事情变得更糟。所以一般的做法是由客户端发出警告,并停止运行,等到有专人来维护后再继续工作。我们开始创建在客户端进行状态更新的模型。以下是客户端代码:

clonesrv3: Clone server, Model Three in C

```
int main (void)
   zctx_t *ctx = zctx_new ();
   void *snapshot = zsocket_new (ctx, ZMQ_ROUTER);
   zsocket_bind (snapshot, "tcp://*:5556");
   void *publisher = zsocket_new (ctx, ZMQ_PUB);
   zsocket_bind (publisher, "tcp://*:5557");
   void *collector = zsocket_new (ctx, ZMQ_PULL);
   zsocket_bind (collector, "tcp://*:5558");
   int64_t sequence = 0;
   zhash_t *kvmap = zhash_new ();
   zmq_pollitem_t items [] = {
       { collector, 0, ZMQ_POLLIN, 0 },
       { snapshot, 0, ZMQ_POLLIN, 0 }
   while (!zctx_interrupted) {
       int rc = zmq_poll (items, 2, 1000 * ZMQ_POLL_MSEC);
       if (items [0].revents & ZMQ_POLLIN) {
           kvmsg_t *kvmsg = kvmsg_recv (collector);
           if (!kvmsg)
               break;
           kvmsg_set_sequence (kvmsg, ++sequence);
           kvmsg_send (kvmsg, publisher);
           kvmsg_store (&kvmsg, kvmap);
           printf ("I: 发布更新事件 %5d\n", (int) sequence);
       if (items [1].revents & ZMQ_POLLIN) {
           zframe_t *identity = zframe_recv (snapshot);
           if (!identity)
               break;
           char *request = zstr_recv (snapshot);
           if (streq (request, "ICANHAZ?"))
               free (request);
           else {
               printf ("E: 错误的请求,程序中止\n");
               break;
```

```
kvroute_t routing = { snapshot, identity };
           zhash_foreach (kvmap, s_send_single, &routing);
           printf ("I: 正在发送快照,版本号:%d\n", (int) sequence);
           zframe_send (&identity, snapshot, ZFRAME_MORE);
           kvmsg_t *kvmsg = kvmsg_new (sequence);
           kvmsg_set_key (kvmsg, "KTHXBAI");
           kvmsg_set_body (kvmsg, (byte *) "", 0);
                         (kvmsg, snapshot);
           kvmsg_send
           kvmsg_destroy (&kvmsg);
    printf (" 已中断\n已处理 %d 条消息\n", (int) sequence);
    zhash_destroy (&kvmap);
    zctx_destroy (&ctx);
    return 0;
static int
s_send_single (char *key, void *data, void *args)
   kvroute_t *kvroute = (kvroute_t *) args;
   zframe_send (&kvroute->identity,
       kvroute->socket, ZFRAME_MORE + ZFRAME_REUSE);
   kvmsg_t *kvmsg = (kvmsg_t *) data;
   kvmsg_send (kvmsg, kvroute->socket);
    return 0;
```

以下是客户端代码:

clonecli3: Clone client, Model Three in C

```
//
// 克隆模式 - 客户端 - 模型3
//

// 直接编译,不创建类库
#include "kvsimple.c"
```

```
int main (void)
   zctx_t *ctx = zctx_new ();
   void *snapshot = zsocket_new (ctx, ZMQ_DEALER);
   zsocket_connect (snapshot, "tcp://localhost:5556");
   void *subscriber = zsocket_new (ctx, ZMQ_SUB);
   zsocket_connect (subscriber, "tcp://localhost:5557");
   void *publisher = zsocket_new (ctx, ZMQ_PUSH);
   zsocket_connect (publisher, "tcp://localhost:5558");
   zhash_t *kvmap = zhash_new ();
   srandom ((unsigned) time (NULL));
   int64_t sequence = 0;
   zstr_send (snapshot, "ICANHAZ?");
   while (TRUE) {
       kvmsg_t *kvmsg = kvmsg_recv (snapshot);
       if (!kvmsg)
           break;
       if (streq (kvmsg_key (kvmsg), "KTHXBAI")) {
           sequence = kvmsg_sequence (kvmsg);
           printf ("I: 已收到快照,版本号:%d\n", (int) sequence);
           kvmsg_destroy (&kvmsg);
           break;
       kvmsg_store (&kvmsg, kvmap);
   int64_t alarm = zclock_time () + 1000;
   while (!zctx_interrupted) {
       zmq_pollitem_t items [] = { { subscriber, 0, ZMQ_POLLIN, 0 } };
       int tickless = (int) ((alarm - zclock_time ()));
       if (tickless < 0)</pre>
           tickless = 0;
       int rc = zmq_poll (items, 1, tickless * ZMQ_POLL_MSEC);
       if (rc == -1)
           break;
       if (items [0].revents & ZMQ_POLLIN) {
           kvmsg_t *kvmsg = kvmsg_recv (subscriber);
           if (!kvmsg)
           if (kvmsg_sequence (kvmsg) > sequence) {
```

```
sequence = kvmsg_sequence (kvmsg);
           kvmsg_store (&kvmsg, kvmap);
           printf ("I: 收到更新事件:%d\n", (int) sequence);
       else
           kvmsg_destroy (&kvmsg);
    if (zclock_time () >= alarm) {
       kvmsg_t *kvmsg = kvmsg_new (0);
       kvmsg_fmt_key (kvmsg, "%d", randof (10000));
       kvmsg_fmt_body (kvmsg, "%d", randof (1000000));
       kvmsg_send
                     (kvmsg, publisher);
       kvmsg_destroy (&kvmsg);
       alarm = zclock_time () + 1000;
printf (" 已准备\n收到 %d 条消息\n", (int) sequence);
zhash_destroy (&kvmap);
zctx_destroy (&ctx);
```

几点说明:

- 服务端整合为一个线程,负责收集来自客户端的更新事件并转发给其他客户端。它使用PULL套接字获取 更新事件,ROUTER套接字处理快照请求,以及PUB套接字发布更新事件。
- 客户端会每隔1秒左右发送随机的更新事件给服务端,现实中这一动作由应用程序触发。

子树克隆

现实中的键值缓存会越变越多,而客户端可能只会需要部分缓存。我们可以使用子树的方式来实现:客户端在发送快照请求时告诉服务端它需要的子树,在订阅更新事件时也指明子树。

关于子树的语法有很多,一种是"分层路径"结构,另一种是"主题树":

- 分层路径:/some/list/of/paths
 - 主题树:some.list.of.topics

这里我们会使用分层路径结构,以此扩展服务端和客户端,进行子树操作。维护多个子树其实并不太困难,因此我们不在这里演示。

下面是服务端代码,由模型3衍化而来:

clonesrv4: Clone server, Model Four in C

```
#include "kvsimple.c"
static int s_send_single (char *key, void *data, void *args);
typedef struct {
   void *socket;
   zframe_t *identity; // 请求方标识
   char *subtree;
} kvroute_t;
int main (void)
   zctx_t *ctx = zctx_new ();
   void *snapshot = zsocket_new (ctx, ZMQ_ROUTER);
   zsocket_bind (snapshot, "tcp://*:5556");
   void *publisher = zsocket_new (ctx, ZMQ_PUB);
   zsocket_bind (publisher, "tcp://*:5557");
   void *collector = zsocket_new (ctx, ZMQ_PULL);
   zsocket_bind (collector, "tcp://*:5558");
   int64_t sequence = 0;
    zhash_t *kvmap = zhash_new ();
   zmq_pollitem_t items [] = {
        { collector, 0, ZMQ_POLLIN, 0 },
       { snapshot, 0, ZMQ_POLLIN, 0 }
   while (!zctx_interrupted) {
        int rc = zmq_poll (items, 2, 1000 * ZMQ_POLL_MSEC);
        if (items [0].revents & ZMQ_POLLIN) {
           kvmsg_t *kvmsg = kvmsg_recv (collector);
           if (!kvmsg)
               break;
           kvmsg_set_sequence (kvmsg, ++sequence);
           kvmsg_send (kvmsg, publisher);
           kvmsg_store (&kvmsg, kvmap);
```

```
if (items [1].revents & ZMQ_POLLIN) {
           zframe_t *identity = zframe_recv (snapshot);
           if (!identity)
           char *request = zstr_recv (snapshot);
           char *subtree = NULL;
           if (streq (request, "ICANHAZ?")) {
               free (request);
               subtree = zstr_recv (snapshot);
           else {
               printf ("E: 错误的请求,程序中止\n");
               break;
           kvroute_t routing = { snapshot, identity, subtree };
           zhash_foreach (kvmap, s_send_single, &routing);
           printf ("I: 正在发送快照,版本号:%d\n", (int) sequence);
           zframe_send (&identity, snapshot, ZFRAME_MORE);
           kvmsg_t *kvmsg = kvmsg_new (sequence);
           kvmsg_set_key (kvmsg, "KTHXBAI");
           kvmsg_set_body (kvmsg, (byte *) subtree, 0);
           kvmsg_send
                        (kvmsg, snapshot);
           kvmsg_destroy (&kvmsg);
           free (subtree);
   printf (" 已中断\n已处理 %d 条消息\n", (int) sequence);
   zhash_destroy (&kvmap);
   zctx_destroy (&ctx);
   return 0;
s_send_single (char *key, void *data, void *args)
```

printf ("I: 发布更新事件 %5d\n", (int) sequence);

下面是客户端代码:

clonecli4: Clone client, Model Four in C

```
#include "kvsimple.c"
#define SUBTREE "/client/"
   zctx_t *ctx = zctx_new ();
   void *snapshot = zsocket_new (ctx, ZMQ_DEALER);
   zsocket_connect (snapshot, "tcp://localhost:5556");
   void *subscriber = zsocket_new (ctx, ZMQ_SUB);
   zsocket_connect (subscriber, "tcp://localhost:5557");
   zsockopt_set_subscribe (subscriber, SUBTREE);
   void *publisher = zsocket_new (ctx, ZMQ_PUSH);
    zsocket_connect (publisher, "tcp://localhost:5558");
   zhash_t *kvmap = zhash_new ();
   srandom ((unsigned) time (NULL));
    int64_t sequence = 0;
    zstr_sendm (snapshot, "ICANHAZ?");
    zstr_send (snapshot, SUBTREE);
```

```
while (TRUE) {
    kvmsg_t *kvmsg = kvmsg_recv (snapshot);
    if (!kvmsg)
       break;
    if (streq (kvmsg_key (kvmsg), "KTHXBAI")) {
        sequence = kvmsg_sequence (kvmsg);
        printf ("I: 已收到快照,版本号:%d\n", (int) sequence);
       kvmsg_destroy (&kvmsg);
       break;
    kvmsg_store (&kvmsg, kvmap);
int64_t alarm = zclock_time () + 1000;
while (!zctx_interrupted) {
    zmq_pollitem_t items [] = { { subscriber, 0, ZMQ_POLLIN, 0 } };
    int tickless = (int) ((alarm - zclock_time ()));
    if (tickless < 0)</pre>
       tickless = 0;
    int rc = zmq_poll (items, 1, tickless * ZMQ_POLL_MSEC);
    if (rc == -1)
       break;
    if (items [0].revents & ZMQ_POLLIN) {
       kvmsg_t *kvmsg = kvmsg_recv (subscriber);
       if (!kvmsg)
           break; // 中断
       if (kvmsg_sequence (kvmsg) > sequence) {
           sequence = kvmsg_sequence (kvmsg);
           kvmsg_store (&kvmsg, kvmap);
           printf ("I: 收到更新事件:%d\n", (int) sequence);
        else
           kvmsg_destroy (&kvmsg);
    if (zclock_time () >= alarm) {
        kvmsg_t *kvmsg = kvmsg_new (0);
        kvmsg_fmt_key (kvmsg, "%s%d", SUBTREE, randof (10000));
       kvmsg_fmt_body (kvmsg, "%d", randof (1000000));
       kvmsg_send
                     (kvmsg, publisher);
       kvmsg_destroy (&kvmsg);
       alarm = zclock_time () + 1000;
```

```
}
printf (" 已准备\n收到 %d 条消息\n", (int) sequence);
zhash_destroy (&kvmap);
zctx_destroy (&ctx);
return 0;
}
```

瞬间值

瞬间值指的是那些会立刻过期的值。如果你用克隆模式搭建一个类似DNS的服务时,就可以用瞬间值来模拟动态DNS解析了。当节点连接网络,对外发布它的地址,并不断地更新地址。如果节点断开连接,则它的地址也会失效。

瞬间值可以和会话(session)联系起来,当会话结束时,瞬间值也就失效了。克隆模式中,会话是由客户端定义的,并会在客户端断开连接时消亡。

更简单的方法是为每一个瞬间值设定一个过期时间,客户端会不断延长这个时间,当断开连接时这个时间将得 不到更新,服务器就会自动将其删除。

我们会用这种简单的方法来实现瞬间值,因为太过复杂的方法可能不值当,它们的差别仅在性能上体现。如果 客户端有很多瞬间值,那为每个值设定过期时间是恰当的;如果瞬间值到达一定的量,那最好还是将其和会话 相关联,统一进行过期处理。

首先,我们需要设法在键值对消息中加入过期时间。我们可以增加一个消息帧,但这样一来每当我们需要增加消息内容时就需要修改kvmsg类库了,这并不合适。所以,我们一次性增加一个"属性"消息帧,用于添加不同的消息属性。

其次,我们需要设法删除这条数据。目前为止服务端和客户端会盲目地增改哈希表中的数据,我们可以这样定义:当消息的值是空的,则表示删除这个键的数据。

下面是一个更为完整的kvmsg类代码,它实现了"属性"帧,以及一个UUID帧,我们后面会用到。该类还会负责处理值为空的消息,达到删除的目的:

kvmsg: Key-value message class - full in C

```
#include <uuid/uuid.h>
#include "zlist.h"
#define KVMSG_KEY_MAX 255
#define FRAME_KEY 0
#define FRAME_SEQ
#define FRAME_UUID
#define FRAME_PROPS
#define FRAME_BODY
#define KVMSG_FRAMES 5
struct _kvmsg {
   int present [KVMSG_FRAMES];
   zmq_msg_t frame [KVMSG_FRAMES];
   char key [KVMSG_KEY_MAX + 1];
   zlist_t *props;
   size_t props_size;
```

```
static void
s_encode_props (kvmsg_t *self)
   zmq_msg_t *msg = &self->frame [FRAME_PROPS];
   if (self->present [FRAME_PROPS])
        zmq_msg_close (msg);
   zmq_msg_init_size (msg, self->props_size);
   char *prop = zlist_first (self->props);
   char *dest = (char *) zmq_msg_data (msg);
   while (prop) {
        strcpy (dest, prop);
       dest += strlen (prop);
       *dest++ = '\n';
        prop = zlist_next (self->props);
   self->present [FRAME_PROPS] = 1;
s_decode_props (kvmsg_t *self)
   zmq_msg_t *msg = &self->frame [FRAME_PROPS];
   self->props_size = 0;
   while (zlist_size (self->props))
        free (zlist_pop (self->props));
   size_t remainder = zmq_msg_size (msg);
   char *prop = (char *) zmq_msg_data (msg);
   char *eoln = memchr (prop, '\n', remainder);
   while (eoln) {
        *eoln = 0;
        zlist_append (self->props, strdup (prop));
        self->props_size += strlen (prop) + 1;
        remainder -= strlen (prop) + 1;
        prop = eoln + 1;
        eoln = memchr (prop, '\n', remainder);
```

```
kvmsg_t *
kvmsg_new (int64_t sequence)
    kvmsg_t
        *self;
    self = (kvmsg_t *) zmalloc (sizeof (kvmsg_t));
    self->props = zlist_new ();
    kvmsg_set_sequence (self, sequence);
    return self;
void
kvmsg_free (void *ptr)
   if (ptr) {
        kvmsg_t *self = (kvmsg_t *) ptr;
        int frame_nbr;
        for (frame_nbr = 0; frame_nbr < KVMSG_FRAMES; frame_nbr++)</pre>
            if (self->present [frame_nbr])
                zmq_msg_close (&self->frame [frame_nbr]);
        while (zlist_size (self->props))
            free (zlist_pop (self->props));
        zlist_destroy (&self->props);
        free (self);
kvmsg_destroy (kvmsg_t **self_p)
    assert (self_p);
   if (*self_p) {
       kvmsg_free (*self_p);
       *self_p = NULL;
```

```
kvmsg_t *
kvmsg_dup (kvmsg_t *self)
    kvmsg_t *kvmsg = kvmsg_new (0);
    int frame_nbr;
    for (frame_nbr = 0; frame_nbr < KVMSG_FRAMES; frame_nbr++) {</pre>
        if (self->present [frame_nbr]) {
            zmq_msg_t *src = &self->frame [frame_nbr];
            zmq_msg_t *dst = &kvmsg->frame [frame_nbr];
            zmq_msg_init_size (dst, zmq_msg_size (src));
            memcpy (zmq_msg_data (dst),
                    zmq_msg_data (src), zmq_msg_size (src));
            kvmsg->present [frame_nbr] = 1;
    kvmsg->props = zlist_copy (self->props);
    return kvmsg;
kvmsg_t *
kvmsg_recv (void *socket)
    assert (socket);
    kvmsg_t *self = kvmsg_new (0);
    int frame_nbr;
    for (frame_nbr = 0; frame_nbr < KVMSG_FRAMES; frame_nbr++) {</pre>
        if (self->present [frame_nbr])
            zmq_msg_close (&self->frame [frame_nbr]);
        zmq_msg_init (&self->frame [frame_nbr]);
        self->present [frame_nbr] = 1;
        if (zmq_recvmsg (socket, &self->frame [frame_nbr], 0) == -1) {
            kvmsg_destroy (&self);
            break;
```

```
int rcvmore = (frame_nbr < KVMSG_FRAMES - 1)? 1: 0;</pre>
        if (zsockopt_rcvmore (socket) != rcvmore) {
            kvmsg_destroy (&self);
            break;
    if (self)
        s_decode_props (self);
    return self;
void
kvmsg_send (kvmsg_t *self, void *socket)
   assert (self);
   assert (socket);
    s_encode_props (self);
    int frame_nbr;
    for (frame_nbr = 0; frame_nbr < KVMSG_FRAMES; frame_nbr++) {</pre>
        zmq_msg_t copy;
        zmq_msg_init (&copy);
        if (self->present [frame_nbr])
            zmq_msg_copy (&copy, &self->frame [frame_nbr]);
        zmq_sendmsg (socket, &copy,
            (frame_nbr < KVMSG_FRAMES - 1)? ZMQ_SNDMORE: 0);</pre>
       zmq_msg_close (&copy);
kvmsg_key (kvmsg_t *self)
   assert (self);
   if (self->present [FRAME_KEY]) {
        if (!*self->key) {
            size_t size = zmq_msg_size (&self->frame [FRAME_KEY]);
```

```
if (size > KVMSG_KEY_MAX)
                size = KVMSG_KEY_MAX;
            memcpy (self->key,
                zmq_msg_data (&self->frame [FRAME_KEY]), size);
            self->key [size] = 0;
        return self->key;
       return NULL;
int64_t
kvmsg_sequence (kvmsg_t *self)
    assert (self);
    if (self->present [FRAME_SEQ]) {
        assert (zmq_msg_size (&self->frame [FRAME_SEQ]) == 8);
        byte *source = zmq_msg_data (&self->frame [FRAME_SEQ]);
        int64_t sequence = ((int64_t) (source [0]) << 56)</pre>
                         + ((int64_t) (source [1]) << 48)
                         + ((int64_t) (source [2]) << 40)
                         + ((int64_t) (source [3]) << 32)
                         + ((int64_t) (source [4]) << 24)
                         + ((int64_t) (source [5]) << 16)
                         + ((int64_t) (source [6]) << 8)
                         + (int64_t) (source [7]);
       return sequence;
byte *
kvmsg_uuid (kvmsg_t *self)
    assert (self);
    if (self->present [FRAME_UUID]
```

```
&& zmq_msg_size (&self->frame [FRAME_UUID]) == sizeof (uuid_t))
        return (byte *) zmq_msg_data (&self->frame [FRAME_UUID]);
   else
       return NULL;
byte *
kvmsg_body (kvmsg_t *self)
   assert (self);
   if (self->present [FRAME_BODY])
       return (byte *) zmq_msg_data (&self->frame [FRAME_BODY]);
size_t
kvmsg_size (kvmsg_t *self)
   assert (self);
   if (self->present [FRAME_BODY])
       return zmq_msg_size (&self->frame [FRAME_BODY]);
   else
       return 0;
kvmsg_set_key (kvmsg_t *self, char *key)
   assert (self);
   zmq_msg_t *msg = &self->frame [FRAME_KEY];
   if (self->present [FRAME_KEY])
       zmq_msg_close (msg);
    zmq_msg_init_size (msg, strlen (key));
```

```
memcpy (zmq_msg_data (msg), key, strlen (key));
    self->present [FRAME_KEY] = 1;
void
kvmsg_set_sequence (kvmsg_t *self, int64_t sequence)
   assert (self);
   zmq_msg_t *msg = &self->frame [FRAME_SEQ];
   if (self->present [FRAME_SEQ])
       zmq_msg_close (msg);
    zmq_msg_init_size (msg, 8);
   byte *source = zmq_msg_data (msg);
   source [0] = (byte) ((sequence >> 56) & 255);
   source [1] = (byte) ((sequence >> 48) & 255);
   source [2] = (byte) ((sequence >> 40) & 255);
   source [3] = (byte) ((sequence >> 32) & 255);
   source [4] = (byte) ((sequence >> 24) & 255);
   source [5] = (byte) ((sequence >> 16) & 255);
   source [6] = (byte) ((sequence >> 8) & 255);
   source [7] = (byte) ((sequence) & 255);
   self->present [FRAME_SEQ] = 1;
kvmsg_set_uuid (kvmsg_t *self)
   assert (self);
   zmq_msg_t *msg = &self->frame [FRAME_UUID];
   uuid_t uuid;
   uuid_generate (uuid);
   if (self->present [FRAME_UUID])
       zmq_msg_close (msg);
   zmq_msg_init_size (msg, sizeof (uuid));
   memcpy (zmq_msg_data (msg), uuid, sizeof (uuid));
    self->present [FRAME_UUID] = 1;
```

```
kvmsg_set_body (kvmsg_t *self, byte *body, size_t size)
   assert (self);
   zmq_msg_t *msg = &self->frame [FRAME_BODY];
   if (self->present [FRAME_BODY])
        zmq_msg_close (msg);
   self->present [FRAME_BODY] = 1;
   zmq_msg_init_size (msg, size);
   memcpy (zmq_msg_data (msg), body, size);
void
kvmsg_fmt_key (kvmsg_t *self, char *format, ...)
   char value [KVMSG_KEY_MAX + 1];
   va_list args;
   assert (self);
   va_start (args, format);
   vsnprintf (value, KVMSG_KEY_MAX, format, args);
   va_end (args);
   kvmsg_set_key (self, value);
kvmsg_fmt_body (kvmsg_t *self, char *format, ...)
   char value [255 + 1];
   va_list args;
   assert (self);
   va_start (args, format);
```

```
vsnprintf (value, 255, format, args);
    va_end (args);
   kvmsg_set_body (self, (byte *) value, strlen (value));
kvmsg_get_prop (kvmsg_t *self, char *name)
   assert (strchr (name, '=') == NULL);
   char *prop = zlist_first (self->props);
   size_t namelen = strlen (name);
   while (prop) {
       if (strlen (prop) > namelen
       && memcmp (prop, name, namelen) == 0
       && prop [namelen] == '=')
           return prop + namelen + 1;
        prop = zlist_next (self->props);
void
kvmsg_set_prop (kvmsg_t *self, char *name, char *format, ...)
   assert (strchr (name, '=') == NULL);
   char value [255 + 1];
   va_list args;
   assert (self);
   va_start (args, format);
   vsnprintf (value, 255, format, args);
   va_end (args);
   char *prop = malloc (strlen (name) + strlen (value) + 2);
```

```
sprintf (prop, "%s=", name);
    char *existing = zlist_first (self->props);
   while (existing) {
        if (memcmp (prop, existing, strlen (prop)) == 0) {
            self->props_size -= strlen (existing) + 1;
            zlist_remove (self->props, existing);
            free (existing);
           break;
        existing = zlist_next (self->props);
   strcat (prop, value);
   zlist_append (self->props, prop);
   self->props_size += strlen (prop) + 1;
kvmsg_store (kvmsg_t **self_p, zhash_t *hash)
   assert (self_p);
   if (*self_p) {
       kvmsg_t *self = *self_p;
       assert (self);
        if (kvmsg_size (self)) {
           if (self->present [FRAME_KEY]
            && self->present [FRAME_BODY]) {
               zhash_update (hash, kvmsg_key (self), self);
               zhash_freefn (hash, kvmsg_key (self), kvmsg_free);
        else
            zhash_delete (hash, kvmsg_key (self));
        *self_p = NULL;
```

```
void
kvmsg_dump (kvmsg_t *self)
    if (self) {
        if (!self) {
            return;
        size_t size = kvmsg_size (self);
        byte *body = kvmsg_body (self);
        fprintf (stderr, "[seq:%" PRId64 "]", kvmsg_sequence (self));
        fprintf (stderr, "[key:%s]", kvmsg_key (self));
        fprintf (stderr, "[size:%zd] ", size);
        if (zlist_size (self->props)) {
            char *prop = zlist_first (self->props);
            while (prop) {
                fprintf (stderr, "%s;", prop);
                prop = zlist_next (self->props);
        int char_nbr;
        for (char_nbr = 0; char_nbr < size; char_nbr++)</pre>
            fprintf (stderr, "%02X", body [char_nbr]);
        fprintf (stderr, "NULL message\n");
kvmsg_test (int verbose)
    kvmsg_t
        *kvmsg;
```

```
zctx_t *ctx = zctx_new ();
void *output = zsocket_new (ctx, ZMQ_DEALER);
int rc = zmq_bind (output, "ipc://kvmsg_selftest.ipc");
assert (rc == 0);
void *input = zsocket_new (ctx, ZMQ_DEALER);
rc = zmq_connect (input, "ipc://kvmsg_selftest.ipc");
assert (rc == 0);
zhash_t *kvmap = zhash_new ();
kvmsg = kvmsg_new (1);
kvmsg_set_key (kvmsg, "key");
kvmsg_set_uuid (kvmsg);
kvmsg_set_body (kvmsg, (byte *) "body", 4);
if (verbose)
    kvmsg_dump (kvmsg);
kvmsg_send (kvmsg, output);
kvmsg_store (&kvmsg, kvmap);
kvmsg = kvmsg_recv (input);
if (verbose)
    kvmsg_dump (kvmsg);
assert (streq (kvmsg_key (kvmsg), "key"));
kvmsg_store (&kvmsg, kvmap);
kvmsg = kvmsg_new (2);
kvmsg_set_prop (kvmsg, "prop1", "value1");
kvmsg_set_prop (kvmsg, "prop2", "value1");
kvmsg_set_prop (kvmsg, "prop2", "value2");
kvmsg_set_key (kvmsg, "key");
kvmsg_set_uuid (kvmsg);
kvmsg_set_body (kvmsg, (byte *) "body", 4);
assert (streq (kvmsg_get_prop (kvmsg, "prop2"), "value2"));
if (verbose)
    kvmsg_dump (kvmsg);
kvmsg_send (kvmsg, output);
kvmsg_destroy (&kvmsg);
kvmsg = kvmsg_recv (input);
if (verbose)
    kvmsg_dump (kvmsg);
assert (streq (kvmsg_key (kvmsg), "key"));
assert (streq (kvmsg_get_prop (kvmsg, "prop2"), "value2"));
kvmsg_destroy (&kvmsg);
```

```
// 关闭并销毁所有对象
zhash_destroy (&kvmap);
zctx_destroy (&ctx);

printf ("OK\n");
return 0;
}
```

客户端模型5和模型4没有太大区别,只是kvmsg类库变了。在更新消息的时候还需要添加一个过期时间的属性:

```
kvmsg_set_prop (kvmsg, "ttl", "%d", randof (30));
```

服务端模型5有较大的变化,我们会用反应堆来代替轮询,这样就能混合处理定时事件和套接字事件了,只是在C语言中是比较麻烦的。下面是代码:

clonesrv5: Clone server, Model Five in C

```
static int s_snapshots (zloop_t *loop, void *socket, void *args);
static int s_collector (zloop_t *loop, void *socket, void *args);
static int s_flush_ttl (zloop_t *loop, void *socket, void *args);
typedef struct {
   zctx_t *ctx;
   zhash_t *kvmap;
   zloop_t *loop;
   int port;
   int64_t sequence;
   void *snapshot;
   void *publisher;
   void *collector;
} clonesrv_t;
```

```
clonesrv_t *self = (clonesrv_t *) zmalloc (sizeof (clonesrv_t));
    self->port = 5556;
    self->ctx = zctx_new ();
    self->kvmap = zhash_new ();
    self->loop = zloop_new ();
    zloop_set_verbose (self->loop, FALSE);
    self->snapshot = zsocket_new (self->ctx, ZMQ_ROUTER);
    self->publisher = zsocket_new (self->ctx, ZMQ_PUB);
    self->collector = zsocket_new (self->ctx, ZMQ_PULL);
    zsocket_bind (self->snapshot, "tcp://*:%d", self->port);
    zsocket_bind (self->publisher, "tcp://*:%d", self->port + 1);
    zsocket_bind (self->collector, "tcp://*:%d", self->port + 2);
    zloop_reader (self->loop, self->snapshot, s_snapshots, self);
    zloop_reader (self->loop, self->collector, s_collector, self);
    zloop_timer (self->loop, 1000, 0, s_flush_ttl, self);
    zloop_start (self->loop);
    zloop_destroy (&self->loop);
    zhash_destroy (&self->kvmap);
    zctx_destroy (&self->ctx);
    free (self);
    return 0;
static int s_send_single (char *key, void *data, void *args);
typedef struct {
   void *socket; // ROUTER套接字 zframe_t *identity; // 请求方标识
   char *subtree;
} kvroute_t;
```

```
s_snapshots (zloop_t *loop, void *snapshot, void *args)
    clonesrv_t *self = (clonesrv_t *) args;
    zframe_t *identity = zframe_recv (snapshot);
    if (identity) {
       char *request = zstr_recv (snapshot);
        char *subtree = NULL;
        if (streq (request, "ICANHAZ?")) {
           free (request);
           subtree = zstr_recv (snapshot);
        else
            printf ("E: 错误的请求,程序中止\n");
        if (subtree) {
           kvroute_t routing = { snapshot, identity, subtree };
           zhash_foreach (self->kvmap, s_send_single, &routing);
           zclock_log ("I: 正在发送快照,版本号:%d", (int) self->sequence);
           zframe_send (&identity, snapshot, ZFRAME_MORE);
           kvmsg_t *kvmsg = kvmsg_new (self->sequence);
           kvmsg_set_key (kvmsg, "KTHXBAI");
           kvmsg_set_body (kvmsg, (byte *) subtree, 0);
           kvmsg_send
                         (kvmsg, snapshot);
           kvmsg_destroy (&kvmsg);
           free (subtree);
   return 0;
s_send_single (char *key, void *data, void *args)
   kvroute_t *kvroute = (kvroute_t *) args;
   kvmsg_t *kvmsg = (kvmsg_t *) data;
   if (strlen (kvroute->subtree) <= strlen (kvmsg_key (kvmsg))</pre>
   && memcmp (kvroute->subtree,
               kvmsg_key (kvmsg), strlen (kvroute->subtree)) == 0) {
```

```
zframe_send (&kvroute->identity,
            kvroute->socket, ZFRAME_MORE + ZFRAME_REUSE);
        kvmsg_send (kvmsg, kvroute->socket);
   return 0;
s_collector (zloop_t *loop, void *collector, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   kvmsg_t *kvmsg = kvmsg_recv (collector);
   if (kvmsg) {
        kvmsg_set_sequence (kvmsg, ++self->sequence);
        kvmsg_send (kvmsg, self->publisher);
        int ttl = atoi (kvmsg_get_prop (kvmsg, "ttl"));
        if (ttl)
           kvmsg_set_prop (kvmsg, "ttl",
               "%" PRId64, zclock_time () + ttl * 1000);
        kvmsg_store (&kvmsg, self->kvmap);
        zclock_log ("I: 正在发布更新事件 %d", (int) self->sequence);
static int s_flush_single (char *key, void *data, void *args);
s_flush_ttl (zloop_t *loop, void *unused, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   zhash_foreach (self->kvmap, s_flush_single, args);
   return 0;
```

```
s_flush_single (char *key, void *data, void *args)
{
    clonesrv_t *self = (clonesrv_t *) args;

    kvmsg_t *kvmsg = (kvmsg_t *) data;
    int64_t ttl;
    sscanf (kvmsg_get_prop (kvmsg, "ttl"), "%" PRId64, &ttl);
    if (ttl && zclock_time () >= ttl) {
        kvmsg_set_sequence (kvmsg, ++self->sequence);
        kvmsg_set_body (kvmsg, (byte *) "", 0);
        kvmsg_send (kvmsg, self->publisher);
        kvmsg_store (&kvmsg, self->kvmap);
        zclock_log ("I: 发布删除事件 %d", (int) self->sequence);
}
    return 0;
}
```

克隆服务器的可靠性

克隆模型1至5相对比较简单,下面我们会探讨一个非常复杂的模型。可以发现,为了构建可靠的消息队列,我们需要花费非常多的精力。所以我们经常会问:有必要这么做吗?如果说你能够接受可靠性不够高的、或者说已经足够好的架构,那恭喜你,你在成本和收益之间找到了平衡。虽然我们会偶尔丢失一些消息,但从经济的角度来说还是合理的。不管怎样,下面我们就来介绍这个复杂的模型。

在模型3中,你会关闭和重启服务,这会导致数据的丢失。任何后续加入的客户端只能得到重启之后的那些数据,而非所有的。下面就让我们想办法让克隆模式能够承担服务器重启的故障。

以下列举我们需要处理的问题:

- 克隆服务器进程崩溃并自动或手工重启。进程丢失了所有数据,所以必须从别处进行恢复。
- 克隆服务器硬件故障,长时间不能恢复。客户端需要切换至另一个可用的服务端。
- 克隆服务器从网络上断开,如交换机发生故障等。它会在某个时点重连,但期间的数据就需要替代的服务器负责处理。

第一步我们需要增加一个服务器。我们可以使用第四章中提到的双子星模式,它是一个反应堆,而我们的程序 经过整理后也是一个反应堆,因此可以互相协作。

我们需要保证更新事件在主服务器崩溃时仍能保留,最简单的机制就是同时发送给两台服务器。

备机就可以当做一台客户端来运行,像其他客户端一样从主机获取更新事件。同时它又能从客户端获取更新事件——虽然不应该以此更新数据,但可以先暂存起来。

所以,相较于模型5,模型6中引入了以下特性:

- 客户端发送更新事件改用PUB-SUB套接字,而非PUSH-PULL。原因是PUSH套接字会在没有接收方时阻塞,且会进行负载均衡——我们需要两台服务器都接收到消息。我们会在服务器端绑定SUB套接字,在客户端连接PUB套接字。
- 我们在服务器发送给客户端的更新事件中加入心跳,这样客户端可以知道主机是否已死,然后切换至备机。
- 我们使用双子星模式的bstar反应堆类来创建主机和备机。双子星模式中需要有一个"投票"套接字,来协助判定对方节点是否已死。这里我们使用快照请求来作为"投票"。
- 我们将为所有的更新事件添加UUID属性,它由客户端生成,服务端会将其发布给所有客户端。
- 备机将维护一个"待处理列表",保存来自客户端、尚未由服务端发布的更新事件;或者反过来,来自服务端、尚未从客户端收到的更新事件。这个列表从旧到新排列,这样就能方便地从顶部删除消息。

我们可以为客户端设计一个有限状态机,它有三种状态:

- 客户端打开并连接了套接字,然后向服务端发送快照请求。为了避免消息风暴,它只会请求两次。
- 客户端等待快照应答,如果获得了则保存它;如果没有获得,则向第二个服务器发送请求。
- 客户端收到快照,便开始等待更新事件。如果在一定时间内没有收到服务端响应,则会连接第二个服务端。

客户端会一直循环下去,可能在程序刚启动时,部分客户端会试图连接主机,部分连接备机,相信双子星模式 会很好地处理这一情况的。

我们可以将客户端状态图绘制出来:

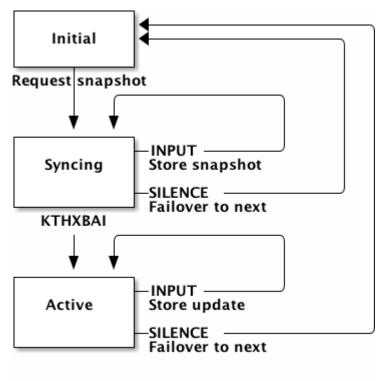


Figure 6 - Clone client FSM

故障恢复的步骤如下:

- 客户端检测到主机不再发送心跳,因此转而连接备机,并请求一份新的快照;
- 备机开始接收快照请求,并检测到主机死亡,于是开始作为主机运行;
- 备机将待处理列表中的更新事件写入自身状态中,然后开始处理快照请求。

当主机恢复连接时:

- 启动为slave状态,并作为克隆模式客户端连接备机;
- 同时,使用SUB套接字从客户端接收更新事件。

我们做两点假设:

- 至少有一台主机会继续运行。如果两台主机都崩溃了,那我们将丢失所有的服务端数据,无法恢复。
- 不同的客户端不会同时更新同一个键值对。客户端的更新事件会先后到达两个服务器,因此更新的顺序可能会不一致。单个客户端的更新事件到达两台服务器的顺序是相同的,所以不用担心。

下面是整体架构图:

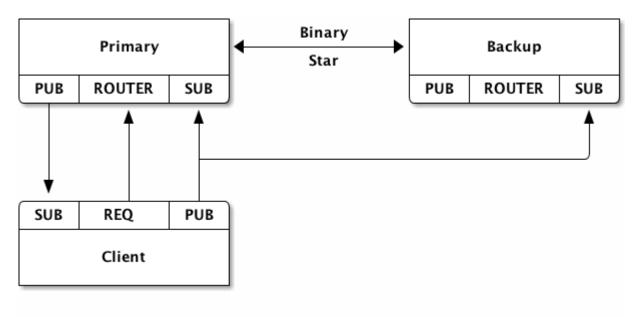


Figure 7 - High availability Clone Server Pair

开始编程之前,我们需要将客户端重构成一个可复用的类。在ZMQ中写异步类有时是为了练习如何写出优雅的代码,但这里我们确实是希望克隆模式可以成为一种易于使用的程序。上述架构的伸缩性来源于客户端的正确行为,因此有必要将其封装成一份API。要在客户端中进行故障恢复还是比较复杂的,试想一下自由者模式和克隆模式结合起来会是什么样的吧。

按照我的习惯,我会先写出一份API的列表,然后加以实现。让我们假想一个名为clone的API,在其基础之上编写克隆模式客户端API。将代码封装为API显然会提升代码的稳定性,就以模型5为例,客户端需要打开三个套接字,端点名称直接写在了代码里。我们可以创建这样一组API:

```
// 为每个套接字指定端点
clone_subscribe (clone, "tcp://localhost:5556");
clone_snapshot (clone, "tcp://localhost:5557");
clone_updates (clone, "tcp://localhost:5558");

// 由于有两个服务端,因此再执行一次
clone_subscribe (clone, "tcp://localhost:5566");
clone_snapshot (clone, "tcp://localhost:5567");
clone_updates (clone, "tcp://localhost:5568");
```

但这种写法还是比较啰嗦的,因为没有必要将API内部的一些设计暴露给编程人员。现在我们会使用三个套接字,而将来可能就会使用两个,或者四个。我们不可能让所有的应用程序都相应地修改吧?让我们把这些信息包装到API中:

```
// 指定主备服务器端点
clone_connect (clone, "tcp://localhost:5551");
clone_connect (clone, "tcp://localhost:5561");
```

这样一来代码就变得非常简洁,不过也会对现有代码的内部就够造成影响。我们需要从一个端点中推算出三个端点。一种方法是假设客户端和服务端使用三个连续的端点通信,并将这个规则写入协议;另一个方法是向服

务器索取缺少的端点信息。我们使用第一种较为简单的方法:

- 服务器状态ROUTER在端点P;
- 服务器更新事件PUB在端点P + 1;
- 服务器更新事件SUB在端点P + 2。

clone类和第四章的flcliapi类很类似,由两部分组成:

- 一个在后台运行的异步克隆模式代理。该代理处理所有的I/O操作,实时地和服务器进行通信;
- 一个在前台应用程序中同步运行的clone类。当你创建了一个clone对象后,它会自动创建后台的clone线程;当你销毁clone对象,该后台线程也会被销毁。

前台的clone类会使用inproc管道和后台的代理进行通信。C语言中,czmq线程会自动为我们创建这个管道。 这也是ZMQ多线程编程的常规方式。

如果没有ZMQ,这种异步的设计将很难处理高压工作,而ZMQ会让其变得简单。编写出来额代码会相对比较复杂。我们可以用反应堆的模式来编写,但这会进一步增加复杂度,且影响应用程序的使用。因此,我们的设计的API将更像是一个能够和服务器进行通信的键值表:

```
clone_t *clone_new (void);
void clone_destroy (clone_t **self_p);
void clone_connect (clone_t *self, char *address, char *service);
void clone_set (clone_t *self, char *key, char *value);
char *clone_get (clone_t *self, char *key);
```

下面就是克隆模式客户端模型6的代码,因为调用了API,所以非常简短:

clonecli6: Clone client, Model Six in C

```
//
// 克隆模式 - 客户端 - 模型6
//
// 直接编译,不建类库
#include "clone.c"

#define SUBTREE "/client/"

int main (void)
{

    // 创建分布式哈希表
    clone_t *clone = clone_new ();

    // 配置
    clone_subtree (clone, SUBTREE);
    clone_connect (clone, "tcp://localhost", "5556");
    clone_connect (clone, "tcp://localhost", "5566");
```

以下是clone类的实现: clone: Clone class in C

```
#include "clone.h"
```

```
#define SERVER_MAX 2
struct _clone_t {
  zctx_t *ctx;
  void *pipe; // 和后台代理间的通信套接字
static void clone_agent (void *args, zctx_t *ctx, void *pipe);
clone_t *
clone_new (void)
   clone_t
      *self;
   self = (clone_t *) zmalloc (sizeof (clone_t));
   self->ctx = zctx_new ();
   self->pipe = zthread_fork (self->ctx, clone_agent, NULL);
   return self;
clone_destroy (clone_t **self_p)
   assert (self_p);
   if (*self_p) {
      clone_t *self = *self_p;
      zctx_destroy (&self->ctx);
```

```
free (self);
        *self_p = NULL;
void clone_subtree (clone_t *self, char *subtree)
   assert (self);
   zmsg_t *msg = zmsg_new ();
   zmsg_addstr (msg, "SUBTREE");
   zmsg_addstr (msg, subtree);
   zmsg_send (&msg, self->pipe);
void
clone_connect (clone_t *self, char *address, char *service)
   assert (self);
   zmsg_t *msg = zmsg_new ();
   zmsg_addstr (msg, "CONNECT");
   zmsg_addstr (msg, address);
   zmsg_addstr (msg, service);
   zmsg_send (&msg, self->pipe);
clone_set (clone_t *self, char *key, char *value, int ttl)
   char ttlstr [10];
   sprintf (ttlstr, "%d", ttl);
   assert (self);
   zmsg_t *msg = zmsg_new ();
    zmsg_addstr (msg, "SET");
```

```
zmsg_addstr (msg, key);
   zmsg_addstr (msg, value);
   zmsg_addstr (msg, ttlstr);
   zmsg_send (&msg, self->pipe);
clone_get (clone_t *self, char *key)
   assert (self);
   assert (key);
   zmsg_t *msg = zmsg_new ();
   zmsg_addstr (msg, "GET");
   zmsg_addstr (msg, key);
   zmsg_send (&msg, self->pipe);
   zmsg_t *reply = zmsg_recv (self->pipe);
   if (reply) {
       char *value = zmsg_popstr (reply);
       zmsg_destroy (&reply);
      return value;
typedef struct {
   char *address;
   int port;
   void *snapshot;
   void *subscriber;
   uint64_t expiry;
   uint requests;
} server_t;
```

```
static server_t *
server_new (zctx_t *ctx, char *address, int port, char *subtree)
   server_t *self = (server_t *) zmalloc (sizeof (server_t));
   zclock_log ("I: adding server %s:%d...", address, port);
   self->address = strdup (address);
   self->port = port;
   self->snapshot = zsocket_new (ctx, ZMQ_DEALER);
   zsocket_connect (self->snapshot, "%s:%d", address, port);
   self->subscriber = zsocket_new (ctx, ZMQ_SUB);
   zsocket_connect (self->subscriber, "%s:%d", address, port + 1);
   zsockopt_set_subscribe (self->subscriber, subtree);
   return self;
static void
server_destroy (server_t **self_p)
   assert (self_p);
   if (*self_p) {
       server_t *self = *self_p;
       free (self->address);
       free (self);
       *self_p = NULL;
#define STATE_INITIAL 0 // 连接之前
#define STATE_SYNCING 1 // 正在同步
#define STATE_ACTIVE
typedef struct {
   zctx_t *ctx;
   void *pipe;
   zhash_t *kvmap;
   char *subtree;
   server_t *server [SERVER_MAX];
   uint nbr_servers;
   uint state;
   uint cur_server;
```

```
int64_t sequence;
   void *publisher;
} agent_t;
static agent_t *
agent_new (zctx_t *ctx, void *pipe)
   agent_t *self = (agent_t *) zmalloc (sizeof (agent_t));
   self->ctx = ctx;
   self->pipe = pipe;
   self->kvmap = zhash_new ();
   self->subtree = strdup ("");
   self->state = STATE_INITIAL;
   self->publisher = zsocket_new (self->ctx, ZMQ_PUB);
   return self;
static void
agent_destroy (agent_t **self_p)
   assert (self_p);
   if (*self_p) {
       agent_t *self = *self_p;
       int server_nbr;
        for (server_nbr = 0; server_nbr < self->nbr_servers; server_nbr++)
            server_destroy (&self->server [server_nbr]);
        zhash_destroy (&self->kvmap);
        free (self->subtree);
        free (self);
       *self_p = NULL;
agent_control_message (agent_t *self)
   zmsg_t *msg = zmsg_recv (self->pipe);
   char *command = zmsg_popstr (msg);
   if (command == NULL)
        return -1;
   if (streq (command, "SUBTREE")) {
        free (self->subtree);
        self->subtree = zmsg_popstr (msg);
```

```
else
    if (streq (command, "CONNECT")) {
        char *address = zmsg_popstr (msg);
        char *service = zmsg_popstr (msg);
        if (self->nbr_servers < SERVER_MAX) {</pre>
            self->server [self->nbr_servers++] = server_new (
                self->ctx, address, atoi (service), self->subtree);
            zsocket_connect (self->publisher, "%s:%d",
                address, atoi (service) + 2);
        else
            zclock_log ("E: too many servers (max. %d)", SERVER_MAX);
        free (address);
        free (service);
   else
   if (streq (command, "SET")) {
        char *key = zmsg_popstr (msg);
        char *value = zmsg_popstr (msg);
        char *ttl = zmsg_popstr (msg);
        zhash_update (self->kvmap, key, (byte *) value);
        zhash_freefn (self->kvmap, key, free);
        kvmsg_t *kvmsg = kvmsg_new (0);
        kvmsg_set_key (kvmsg, key);
        kvmsg_set_uuid (kvmsg);
        kvmsg_fmt_body (kvmsg, "%s", value);
        kvmsg_set_prop (kvmsg, "ttl", ttl);
        kvmsg_send (kvmsg, self->publisher);
        kvmsg_destroy (&kvmsg);
puts (key);
       free (ttl);
        free (key);
    else
   if (streq (command, "GET")) {
        char *key = zmsg_popstr (msg);
        char *value = zhash_lookup (self->kvmap, key);
            zstr_send (self->pipe, value);
        else
            zstr_send (self->pipe, "");
        free (key);
        free (value);
```

```
free (command);
    zmsg_destroy (&msg);
    return 0;
static void
clone_agent (void *args, zctx_t *ctx, void *pipe)
   agent_t *self = agent_new (ctx, pipe);
   while (TRUE) {
        zmq_pollitem_t poll_set [] = {
            { pipe, 0, ZMQ_POLLIN, 0 },
            { 0, 0, ZMQ_POLLIN, 0 }
        int poll_timer = -1;
        int poll_size = 2;
        server_t *server = self->server [self->cur_server];
        switch (self->state) {
            case STATE_INITIAL:
                if (self->nbr_servers > 0) {
                    zclock_log ("I: 正在等待服务器 %s:%d...",
                        server->address, server->port);
                    if (server->requests < 2) {</pre>
                        zstr_sendm (server->snapshot, "ICANHAZ?");
                        zstr_send (server->snapshot, self->subtree);
                        server->requests++;
                    server->expiry = zclock_time () + SERVER_TTL;
                    self->state = STATE_SYNCING;
                    poll_set [1].socket = server->snapshot;
                else
                    poll_size = 1;
            case STATE_SYNCING:
                poll_set [1].socket = server->snapshot;
                break;
            case STATE_ACTIVE:
```

```
poll_set [1].socket = server->subscriber;
       break;
if (server) {
   poll_timer = (server->expiry - zclock_time ())
              * ZMQ_POLL_MSEC;
   if (poll_timer < 0)
       poll_timer = 0;
int rc = zmq_poll (poll_set, poll_size, poll_timer);
if (rc == -1)
   break;
if (poll_set [0].revents & ZMQ_POLLIN) {
   if (agent_control_message (self))
       break; // 中断
else
if (poll_set [1].revents & ZMQ_POLLIN) {
   kvmsg_t *kvmsg = kvmsg_recv (poll_set [1].socket);
   if (!kvmsg)
       break; // 中断
   server->expiry = zclock_time () + SERVER_TTL;
   if (self->state == STATE_SYNCING) {
       server->requests = 0;
       if (streq (kvmsg_key (kvmsg), "KTHXBAI")) {
           self->sequence = kvmsg_sequence (kvmsg);
           self->state = STATE_ACTIVE;
           zclock_log ("I: received from %s:%d snapshot=%d",
               server->address, server->port,
               (int) self->sequence);
           kvmsg_destroy (&kvmsg);
       else
           kvmsg_store (&kvmsg, self->kvmap);
   if (self->state == STATE_ACTIVE) {
       if (kvmsg_sequence (kvmsg) > self->sequence) {
```

最后是克隆服务器的模型6代码:

clonesrv6: Clone server, Model Six in C

```
bstar_t *bstar;
    int64_t sequence;
   int port;
   int peer;
   void *publisher;
   void *collector;
   void *subscriber;
   zlist_t *pending;
   Bool primary;
   Bool master;
   Bool slave;
} clonesrv_t;
int main (int argc, char *argv [])
   clonesrv_t *self = (clonesrv_t *) zmalloc (sizeof (clonesrv_t));
   if (argc == 2 && streq (argv [1], "-p")) {
       zclock_log ("I: 作为主机master运行,正在等待备机slave连接。");
       self->bstar = bstar_new (BSTAR_PRIMARY, "tcp://*:5003",
                               "tcp://localhost:5004");
       bstar_voter (self->bstar, "tcp://*:5556", ZMQ_ROUTER,
                    s_snapshots, self);
       self->port = 5556;
       self->peer = 5566;
       self->primary = TRUE;
   else
   if (argc == 2 && streq (argv [1], "-b")) {
       zclock_log ("I: 作为备机slave运行,正在等待主机master连接。");
       self->bstar = bstar_new (BSTAR_BACKUP, "tcp://*:5004",
                                "tcp://localhost:5003");
       bstar_voter (self->bstar, "tcp://*:5566", ZMQ_ROUTER,
                    s_snapshots, self);
       self->port = 5566;
       self->peer = 5556;
       self->primary = FALSE;
       printf ("Usage: clonesrv4 { -p | -b }\n");
       free (self);
       exit (0);
   if (self->primary)
       self->kvmap = zhash_new ();
```

```
self->ctx = zctx_new ();
    self->pending = zlist_new ();
   bstar_set_verbose (self->bstar, TRUE);
    self->publisher = zsocket_new (self->ctx, ZMQ_PUB);
    self->collector = zsocket_new (self->ctx, ZMQ_SUB);
    zsocket_bind (self->publisher, "tcp://*:%d", self->port + 1);
    zsocket_bind (self->collector, "tcp://*:%d", self->port + 2);
   self->subscriber = zsocket_new (self->ctx, ZMQ_SUB);
    zsocket_connect (self->subscriber, "tcp://localhost:%d", self->peer + 1);
   bstar_new_master (self->bstar, s_new_master, self);
   bstar_new_slave (self->bstar, s_new_slave, self);
    zloop_reader (bstar_zloop (self->bstar), self->collector, s_collector, self);
    zloop_timer (bstar_zloop (self->bstar), 1000, 0, s_flush_ttl, self);
    zloop_timer (bstar_zloop (self->bstar), 1000, 0, s_send_hugz, self);
   bstar_start (self->bstar);
   while (zlist_size (self->pending)) {
       kvmsg_t *kvmsg = (kvmsg_t *) zlist_pop (self->pending);
       kvmsg_destroy (&kvmsg);
    zlist_destroy (&self->pending);
   bstar_destroy (&self->bstar);
   zhash_destroy (&self->kvmap);
   zctx_destroy (&self->ctx);
    free (self);
    return 0;
static int s_send_single (char *key, void *data, void *args);
```

```
typedef struct {
   void *socket;
   zframe_t *identity; // 请求放标识
   char *subtree;
} kvroute_t;
s_snapshots (zloop_t *loop, void *snapshot, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   zframe_t *identity = zframe_recv (snapshot);
   if (identity) {
       char *request = zstr_recv (snapshot);
       char *subtree = NULL;
       if (streq (request, "ICANHAZ?")) {
           free (request);
           subtree = zstr_recv (snapshot);
       else
           printf ("E: 错误的请求,正在退出.....\n");
       if (subtree) {
           kvroute_t routing = { snapshot, identity, subtree };
           zhash_foreach (self->kvmap, s_send_single, &routing);
           zclock_log ("I: 正在发送快照,版本号:%d", (int) self->sequence);
           zframe_send (&identity, snapshot, ZFRAME_MORE);
           kvmsg_t *kvmsg = kvmsg_new (self->sequence);
           kvmsg_set_key (kvmsg, "KTHXBAI");
           kvmsg_set_body (kvmsg, (byte *) subtree, 0);
           kvmsg_send
                         (kvmsg, snapshot);
           kvmsg_destroy (&kvmsg);
           free (subtree);
   return 0;
```

```
static int
s_send_single (char *key, void *data, void *args)
   kvroute_t *kvroute = (kvroute_t *) args;
   kvmsg_t *kvmsg = (kvmsg_t *) data;
   if (strlen (kvroute->subtree) <= strlen (kvmsg_key (kvmsg))</pre>
    && memcmp (kvroute->subtree,
                kvmsg_key (kvmsg), strlen (kvroute->subtree)) == 0) {
        zframe_send (&kvroute->identity,
            kvroute->socket, ZFRAME_MORE + ZFRAME_REUSE);
        kvmsg_send (kvmsg, kvroute->socket);
static int s_was_pending (clonesrv_t *self, kvmsg_t *kvmsg);
static int
s_collector (zloop_t *loop, void *collector, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   kvmsg_t *kvmsg = kvmsg_recv (collector);
   kvmsg_dump (kvmsg);
   if (kvmsg) {
        if (self->master) {
            kvmsg_set_sequence (kvmsg, ++self->sequence);
            kvmsg_send (kvmsg, self->publisher);
            int ttl = atoi (kvmsg_get_prop (kvmsg, "ttl"));
                kvmsg_set_prop (kvmsg, "ttl",
                    "%" PRId64, zclock_time () + ttl * 1000);
            kvmsg_store (&kvmsg, self->kvmap);
            zclock_log ("I: 正在发布更新事件:%d", (int) self->sequence);
        else {
            if (s_was_pending (self, kvmsg))
                kvmsg_destroy (&kvmsg);
```

```
else
                zlist_append (self->pending, kvmsg);
   return 0;
s_was_pending (clonesrv_t *self, kvmsg_t *kvmsg)
   kvmsg_t *held = (kvmsg_t *) zlist_first (self->pending);
   while (held) {
        if (memcmp (kvmsg_uuid (kvmsg),
                   kvmsg_uuid (held), sizeof (uuid_t)) == 0) {
            zlist_remove (self->pending, held);
           return TRUE;
       held = (kvmsg_t *) zlist_next (self->pending);
   return FALSE;
static int s_flush_single (char *key, void *data, void *args);
static int
s_flush_ttl (zloop_t *loop, void *unused, void *args)
    clonesrv_t *self = (clonesrv_t *) args;
   zhash_foreach (self->kvmap, s_flush_single, args);
   return 0;
s_flush_single (char *key, void *data, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   kvmsg_t *kvmsg = (kvmsg_t *) data;
    int64_t ttl;
```

```
sscanf (kvmsg_get_prop (kvmsg, "ttl"), "%" PRId64, &ttl);
    if (ttl && zclock_time () >= ttl) {
        kvmsg_set_sequence (kvmsg, ++self->sequence);
       kvmsg_set_body (kvmsg, (byte *) "", 0);
       kvmsg_send (kvmsg, self->publisher);
       kvmsg_store (&kvmsg, self->kvmap);
       zclock_log ("I: 正在发布删除事件:%d", (int) self->sequence);
s_send_hugz (zloop_t *loop, void *unused, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   kvmsg_t *kvmsg = kvmsg_new (self->sequence);
   kvmsg_set_key (kvmsg, "HUGZ");
   kvmsg_set_body (kvmsg, (byte *) "", 0);
   kvmsg_send
                (kvmsg, self->publisher);
   kvmsg_destroy (&kvmsg);
   return 0;
s_new_master (zloop_t *loop, void *unused, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   self->master = TRUE;
   self->slave = FALSE;
    zloop_cancel (bstar_zloop (self->bstar), self->subscriber);
```

```
while (zlist_size (self->pending)) {
        kvmsg_t *kvmsg = (kvmsg_t *) zlist_pop (self->pending);
        kvmsg_set_sequence (kvmsg, ++self->sequence);
        kvmsg_send (kvmsg, self->publisher);
       kvmsg_store (&kvmsg, self->kvmap);
       zclock_log("I: 正在发布延迟列表中的更新事件:%d",(int) self->sequence);
    return 0;
s_new_slave (zloop_t *loop, void *unused, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   zhash_destroy (&self->kvmap);
   self->master = FALSE;
   self->slave = TRUE;
    zloop_reader (bstar_zloop (self->bstar), self->subscriber,
                 s_subscriber, self);
   return 0;
static int
s_subscriber (zloop_t *loop, void *subscriber, void *args)
   clonesrv_t *self = (clonesrv_t *) args;
   if (self->kvmap == NULL) {
       self->kvmap = zhash_new ();
       void *snapshot = zsocket_new (self->ctx, ZMQ_DEALER);
        zsocket_connect (snapshot, "tcp://localhost:%d", self->peer);
        zclock_log ("I: 正在请求快照:tcp://localhost:%d",
                   self->peer);
        zstr_send (snapshot, "ICANHAZ?");
        while (TRUE) {
            kvmsg_t *kvmsg = kvmsg_recv (snapshot);
```

```
if (!kvmsg)
            break;
        if (streq (kvmsg_key (kvmsg), "KTHXBAI")) {
            self->sequence = kvmsg_sequence (kvmsg);
            kvmsg_destroy (&kvmsg);
            break;
       kvmsg_store (&kvmsg, self->kvmap);
    zclock_log ("I: 收到快照,版本号:%d", (int) self->sequence);
    zsocket_destroy (self->ctx, snapshot);
kvmsg_t *kvmsg = kvmsg_recv (subscriber);
if (!kvmsg)
if (strneq (kvmsg_key (kvmsg), "HUGZ")) {
    if (!s_was_pending (self, kvmsg)) {
       zlist_append (self->pending, kvmsg_dup (kvmsg));
    if (kvmsg_sequence (kvmsg) > self->sequence) {
       self->sequence = kvmsg_sequence (kvmsg);
       kvmsg_store (&kvmsg, self->kvmap);
       zclock_log ("I: 收到更新事件:%d", (int) self->sequence);
    else
       kvmsg_destroy (&kvmsg);
else
    kvmsg_destroy (&kvmsg);
return 0;
```

这段程序只有几百行,但还是花了一些时间来进行调通的。这个模型中包含了故障恢复,瞬间值,子树等等。 虽然我们前期设计得很完备,但要在多个套接字之间进行调试还是很困难的。以下是我的工作方式:

• 由于使用了反应堆(bstar,建立在zloop之上),我们节省了大量的代码,让程序变得简洁明了。整个服务以一个线程运行,因此不会出现跨线程的问题。只需将结构指针(self)传递给所有的处理器即可。此外,使用发应堆后可以让代码更为模块化,易于重用。

- 我们逐个模块进行调试,只有某个模块能够正常运行时才会进入下一步。由于使用了四五个套接字,因此调试的工作量是很大的。我直接将调试信息输出到了屏幕上,因为实在没有必要专门开一个调试器来工作。
- 因为一直在使用valgrind工具进行测试,因此我能确定程序没有内存泄漏的问题。在C语言中,内存泄漏 是我们非常关心的问题,因为没有什么垃圾回收机制可以帮你完成。正确地使用像kvmsg、czmq之类的 抽象层可以很好地避免内存泄漏。

这段程序肯定还会存在一些BUG,部分读者可能会帮助我调试和修复,我在此表示感谢。

测试模型6时,先开启主机和备机,再打开一组客户端,顺序随意。随机地中止某个服务进程,如果程序设计 得是正确的,那客户端获得的数据应该都是一致的。

克隆模式协议

花费了那么多精力来开发一套可靠的发布-订阅模式机制,我们当然希望将来能够方便地在其基础之上进行扩展。较好的方法是将其编写为一个协议,这样就能让各种语言来实现它了。

我们将其称为"集群化哈希表协议",这是一个能够跨集群地进行键值哈希表管理,提供了多客户端的通信机制;客户端可以只操作一个子树的数据,包括更新和定义瞬间值。

http://rfc.zeromq.org/spec:12