

# CVXALGGeo USER MANUAL

MEHDI GHASEMI

## 1. WHAT IS CVXALGGeo?

CVXALGGeo is a package written for SAGE and published under GNU(GPL) to implement a fast method to compute a lower bound for multivariate polynomials over a basic closed semialgebraic set. It is mainly based on methods introduced in [2], [1] and [3]. Moreover, to provide computational tools for usual calculation based on semidefinite programming CVXALGGeo implements some functionalities available on SOSTOOLS and GLOPTIPOLY for MATLAB [4, 8].

At this stage, CVXALGGeo contains implementations of two different methods for polynomial optimization:

**1.1. By Geometric Programming:** Let  $f = \sum_{\alpha} f_{\alpha} \underline{X}^{\alpha}$  be a polynomials on  $X_1, \dots, X_n$  of degree  $d$ , where  $\underline{X}^{\alpha}$  is short for  $X_1^{\alpha_1} \dots X_n^{\alpha_n}$ , and  $\alpha = (\alpha_1, \dots, \alpha_n) \in \mathbb{N}^n$ . The degree of a monomial  $\underline{X}^{\alpha}$  is denoted by  $|\alpha|$  that is defined by  $|\alpha| = \alpha_1 + \dots + \alpha_n$ . Denote  $f_{\alpha}$  where  $\alpha = 2d(\delta_{1i}, \dots, \delta_{ni})$  by  $f_{2d,i}$  and  $f_{(0,\dots,0)} = f(\underline{0})$  by  $f_0$  for simplicity. Degree of a polynomial  $f$  ( $\deg f$ ) is the maximum degree of its monomials. We also denote the set of those exponents  $\alpha$  for which the monomial  $f_{\alpha} \underline{X}^{\alpha}$  is not a square by  $\Delta_f$ , i.e.,

$$\Delta_f := \{\alpha \in \mathbb{N}^n : f_{\alpha} < 0 \text{ or } \alpha_i \text{ is odd for some } 1 \leq i \leq n\}.$$

Let  $\mathbf{g} = (g_1, \dots, g_m)$  be a tuple of polynomials and  $K_{\mathbf{g}} = \{x \in \mathbb{R}^n : g_j(x) \geq 0, j = 1, \dots, m\}$ . Denote the infimum of  $f$  over  $K_{\mathbf{g}}$  by  $f_{*,\mathbf{g}}$ ,  $\Delta = \Delta_f \cup \Delta_{-g_1} \cup \dots \cup \Delta_{-g_m}$  and  $A = (a_{kl})$  be a  $(m+1) \times (m+1)$  matrix such that  $a_{0l} = 1$  if  $l = 1$  and 0 otherwise. Define  $h_k = \sum_{j=0}^m a_{jk} g_j$ ,  $k = 0, \dots, m$  where  $g_0 = -f$  and  $H(\mu) = -\sum_{j=0}^m \mu h_j$ . We can Decompose  $H$  as  $H^+ - H^-$  where all the coefficients of  $H^+$  and  $H^-$  are nonnegative. If  $\rho$  is an optimum value for the following *geometric program*, then  $f_{gp,\mathbf{g}} = -h_0(0) - \rho$

---

*Date:* January 25, 2013.

*Key words and phrases.* sums of squares, semidefinite programming, geometric programming, optimization.

is a lower bound for the global minimum of  $f$  on  $K_{\mathbf{g}}$ .

$$(1) \quad \left\{ \begin{array}{l} \text{Minimize } \sum_{j=1}^m \mu_j h_j(0)^+ + \sum_{\alpha \in \Delta^{<d}} (d - |\alpha|) \left[ \left( \frac{w_\alpha}{d} \right)^d \left( \frac{\alpha}{\mathbf{z}_\alpha} \right)^\alpha \right]^{1/(d-|\alpha|)} \\ \\ \text{s.t. } \sum_{\alpha \in \Delta} z_{\alpha,i} \leq H(\mu)_{d,i}, \quad i = 1, \dots, n \\ \\ \left( \frac{\mathbf{z}_\alpha}{\alpha} \right)^\alpha \geq \left( \frac{w_\alpha}{d} \right)^d, \quad \alpha \in \Delta^=d \\ \\ w_\alpha \geq \max\{H(\mu)_\alpha^+, H(\mu)_\alpha^-\}, \quad \alpha \in \Delta \\ \\ \sum_{k=0}^m a_{jk} \mu_k \geq 0, \quad j = 1, \dots, m \end{array} \right.$$

Here  $h_j^+(0) = \max\{h_j(0), 0\}$  [3, Theorem 4.1].

In [1], [2] and [3] the advantages and disadvantages of solving (1) instead of solving the corresponding semidefinite program is discussed.

**1.2. By Semidefinite Programming:** It is well known that a feasible value for the semidefinite program (2) is a lower bound for the polynomial  $f$  on a semialgebraic set  $K = K_S$ , where  $S = \{g_1, \dots, g_m\}$  and  $K_S = \{x \in \mathbb{R}^n : g_i(x) \geq 0, i = 1, \dots, m\}$ .

$$(2) \quad \text{Prg}_k : \left\{ \begin{array}{l} \min_y \quad \sum_{\alpha} f_{\alpha} y_{\alpha} \\ \text{subject to } \quad M_k(y) \geq 0 \\ \quad \quad \quad M_{k-\deg g_i}(g_i y) \geq 0 \quad i = 1, \dots, m \end{array} \right.$$

for  $k \geq \max\{\deg f, \deg g_1, \dots, \deg g_m\}$ . All the notations are as described in [5]. Current implementation is based on the description given on [6, Chapter 4 and 5] and [7].

## 2. CVXALGGeo

CVXALGGeo implements few tools for computations over real polynomial rings in Python for SAGE using CVXOPT. The current version consists of two main classes.

- (1) GLOptGeoPrg
- (2) SosTools

In this section we describe how to use these programs.

### 2.1. GLOptGeoPrg(Prog, Rng [, H=A, Settings]).

**Initialization:** This class implements (1), requires two mandatory arguments, Prog, Rng and two optional arguments H and Settings.

- Prog=[f, Cons] is a list consists of two parts:

- ‘f’ is the polynomial which is subject to minimizations.
- ‘Cons=[g1, ..., gm]’ is a list of polynomials corresponding to the constraints  $g_j \geq 0$ , for  $j = 1, \dots, m$ .

They must be elements of the polynomial ring `Rng, 'PolynomialRing(F, 'x', n)'` where `n` is the number of variables and ‘F’ is the base field.

- `H` is the matrix  $A = (a_{kl})$ .
- `Settings` is a dictionary with the following keys:
  - ‘Iterations’ is the number of iterations the solver tries to find an optimum solution. Its default value is 150.
  - ‘Details’ is by default set to `False`. Set `True` to see the detail output of the CVXOPT solver.
  - ‘tryKKT’ is by default 20 that forces the solver to continue for the given number of steps, even if a singular KKT matrix occurred.
  - ‘AutoMatrix’ accepts a boolean ‘True’ or ‘False’. The value `True` leaves the choice of  $A$  to the program automatically.

**Running:** The method `G1OptGeoPrg.minimize()` solves (1) for `f` and `g1, ..., gm`. The attribute `G1OptGeoPrg.fgp` also contains the output of `minimize`.

**Output:** The attribute `G1OptGeoPrg.Info` is a dictionary which gives detailed information on the results of the method `minimize`.

- ‘gp’ is the lower bound obtained from (1).
- ‘Wall’ contains the wall time consumed by the solver.
- ‘CPU’ contains the CPU time consumed by the solver.
- ‘status’ and ‘Message’ give some informations about the running process.

## 2.2. `SosTools(Prg, Rng [, Settings])`.

This class is on its early stage and aims to provides the same functionality as `SOSTOOLS` does for `MATLAB`.

**Initialization:** This class implements the semidefinite program (2).

- `Prg=[f,Cons]` is a list consists of two parts:
  - ‘f’ is the polynomial which is subject to minimizations.
  - ‘Cons=[g1, ..., gm]’ is a list of polynomials corresponding to the constraints  $g_j \geq 0$ , for  $j = 1, \dots, m$ .

They must be elements of the polynomial ring `Rng, 'PolynomialRing(F, 'x', n)'` where `n` is the number of variables and ‘F’ is the base field.

- `Settings` is exactly similar to `G1OptGeoPrg` with an extra index ‘Order’ which corresponds to the index  $k$  for `Prgk` in (2). The default value for `Order` is set to be 0 which results in the minimum size for moment matrices.

**Running:**

- The method `SosTools.init_sdp()` initializes the semidefinite program corresponding to (2).
- The method `SosTools.minimize()` returns the an optimum value  $f_{\text{sos}}$ , for (2) if exists.
- The method `SosTools.decompose` returns a vector of polynomials  $[p_1, \dots, p_s]$  such that  $f = \sum_{i=1}^s p_i^2$ .

**Output:** The attribute `SosTools.Info` returns the similar information as for `GlOptGeoPrg.Info`. But instead of ‘gp’ it has the key ‘min’ for  $f_{\text{sos}}$ . The Info key ‘is sos’ is by default `False`, but after executing ‘decompose’ it may become `True`.

### 2.3. Example.

```
## Initial settings ##
n = 3 # Number of variables
d = 4 # Maximum degree of polynomials
m = 2 # Number of constraints
num_monos = 4 # Max number of monomials
R = PolynomialRing(RR,'x',n);
X = R.gens();

#####

Prg=[]
Cns=[]
f=R.random_element(d,randint(1,num_monos))+sum(p^d for p in X)
Prg.append(f)
for i in range(m):
    g=R.random_element(d,num_monos)
    Cns.append(g);
Prg.append(Cns)

conf={'Details':False,'tryKKT':20, 'AutoMatrix':False, 'Order':1}
M=identity_matrix(QQ,m+1)

print Prg

A=GlOptGeoPrg(Prg, R, H=M, Settings=conf)
A.minimize()
```

```

show(A.H)
print A.Info
B=SosTools(Prg, R, Settings=conf)
B.init_sdp()
B.minimize()
print B.Info

```

The code generates random polynomials and computes  $f_{\text{gp}}$  and  $f_{\text{gp,g}}$ : Here is the output:

```

[x0^4 + x1^4 + x2^4 - x2^3 + 122*x0,
[-x0^2*x1 - 4*x2^3 - x1, -x0^2*x1*x2 + 10*x1^3*x2 - 2*x1^2*x2]]

{'status': 'Optimal', 'Message': 'Optimal solution found by solver.',
'Wall': 0.2654759883880615, 'CPU': 0.17999999999999926, 'gp': -285.988059619120}

{'status': 'Optimal', 'min': -285.98744579692226, 'Wall': 1.1604969501495361,
'Message': 'Feasible solution for moments of order 3',
'CPU': 0.84999999999999943, 'Size': [20, 84]}

```

## REFERENCES

- [1] M. Ghasemi, J. B. Lasserre and M. Marshall, *Lower bounds on the global minimum of a polynomial*, to appear.
- [2] M. Ghasemi and M. Marshall, *Lower bounds for polynomials using geometric programming*, SIAM J. Optim. 22(2) (460-473), 2012.
- [3] M. Ghasemi and M. Marshall, *Lower Bounds for a Polynomial on a basic closed semialgebraic set using geometric programming*, in progress.
- [4] D. Henrion and J. B. Lasserre, *Gloptipoly: global optimization over polynomials with matlab and sedumi*, ACM Transactions on Mathematical Software, 29(2):165195, 2003.
- [5] J. B. Lasserre, *Global optimization with polynomials and the problem of moments*, SIAM J. Opt. 11(3)(796817), 2001.
- [6] J. B. Lasserre, *Moments, Positive Polynomials and Their Applications*, Imperial College Press Optimization Series Vol. 1, 2010.
- [7] M. Laurent, *Sums of squares, moment matrices and optimization over polynomials*, 2010.
- [8] S. Prajna, A. Papachristodoulou, P. Seiler, and P. A. Parrilo, *SOSTOOLS: Sum of squares optimization toolbox for MATLAB*, 2004

DEPARTMENT OF MATHEMATICS AND STATISTICS,  
 UNIVERSITY OF SASKATCHEWAN,  
 SASKATOON, SK. S7N 5E6, CANADA  
*E-mail address:* mehdi.ghasemi@usask.ca