Formal Languages Homework 6

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1 Problem 6.1.1

Suppose the PDA $P=(\{q,p\},\{0,1\},\{Z_0,X\},\delta,q,Z_0,\{p\})$ has the following transition function

- 1. $\delta(q, 0, Z_0) = \{(q, XZ_0)\}.$
- 2. $\delta(q, 0, X) = \{(q, XX)\}.$
- 3. $\delta(q, 1, X) = \{(q, X)\}.$
- 4. $\delta(q, \epsilon, X) = \{(p, \epsilon)\}.$
- 5. $\delta(p, \epsilon, X) = \{(p, \epsilon)\}.$
- 6. $\delta(p, 1, X) = \{(p, XX)\}.$
- 7. $\delta(p, 1, Z_0) = \{(p, \epsilon)\}.$

2 Problem 6.2.1

Design a PDA to accept each of the following languages. You may accept either by final state or by empty stack, whichever is more convenient.

- 2.1 a). $\{0^n1^n \mid n \ge 1\}$
- 2.2 b). The set of all strings of 0's and 1's such that no prefix has more 1's than 0's.
- 2.3 c). The set of all strings of 0's and 1's with an equal number of 0's and 1's.

3 Problem 6.3.2

Convert the grammar to a PDA that accepts the same language by empty stack

4 Problem 6.4.1

For each of the following PDA's, tell whether or not it is deterministic. Either show that it meets the definition of a DPDA or find a rule or rules that violate it.

$$\begin{array}{ccc} S & \rightarrow & 0S1 \mid A \\ A & \rightarrow & 1A0 \mid S \mid \epsilon \end{array}$$

4.1 The PDA of Example 6.2

The PDA $P = (\{q_0, q_1, q_2\}, \{0, 1\}, \{0, 1, Z_0\}, \delta, q_0, Z_0, \{q_2\})$, where δ is defined by the following rules:

These rules are hard to read

4.2 The PDA of Exercise 6.1.1

4.3 The PDA of Exercise 6.3.3

The PDA $P = (\{q, p\}, \{0, 1\}, \{Z_0, X\}, \delta, q, Z_0)$, where δ is defined to be:

- 1. $\delta(q, 1, Z_0) = \{(q, XZ_0)\}.$
- 2. $\delta(q, 1, X) = \{(q, XX)\}.$
- 3. $\delta(q, 0, X) = \{(p, X)\}.$
- 4. $\delta(q, \epsilon, X) = \{(q, \epsilon)\}.$
- 5. $\delta(p, 1, X) = \{(p, \epsilon)\}.$
- 6. $\delta(p, 0, Z_0) = \{(q, Z_0)\}.$

5 Problem 6.4.2

Give deterministic pushdown automata to accept the following languages:

- 5.1 a). $\{0^n 1^m \mid n \leq m\}$
- 5.2 b). $\{0^n 1^m \mid n \ge m\}$
- 5.3 c). $\{0^n 1^m 0^n \mid n \text{ and } m \text{ are arbitrary }\}$.