Formal Languages Homework 6

Liam Dillingham

March 26, 2019

Problem 4.1.1 and 4.1.2

Provide the CFG for the following languages:

- 1.1 b). The set of strings of balanced parentheses. These are the strings of characters "(" and ")" that can appear in a well-formed arithmetic expression.
- $S \to (S) \mid SS \mid \epsilon$.

 $V = \{S\}.$ $T = \{ "(", ")" \}.$ (The quotations are only to make reading easier).

1.2 d). $\{0^n1^m2^n \mid n \text{ and } m \text{ are abitrary integers }\}.$

$$V = \{S, A\}.$$

 $T = \{0, 1, 2\}.$

$$\begin{array}{ccc} S & \rightarrow & 0S2 \mid 0A2 \mid A \mid \epsilon \\ A & \rightarrow & 1 \mid \epsilon \end{array}$$

1.3 f).
$$\{0^n 1^{2n} \mid n \ge 1\}$$

$$V = \{S, A\}$$

$$T = \{0, 1\}$$

$$\begin{array}{ccc} S & \rightarrow & 0S11 \mid 0A11 \mid \\ A & \rightarrow & \epsilon \end{array}$$

1.4 h). The set of strings of the form $w1^n$, where w is a string of 0's and 1's of length n.

$$V = \{S, a\}$$
$$T = \{0, 1\}$$

$$\begin{array}{ccc} S & \rightarrow & AS1 \mid \epsilon \\ A & \rightarrow & 0 \mid 1 \end{array}$$

2 Problem 5.1.2

The following grammar generates the language of regular expression 0*1(0+1)*.

$$\begin{array}{ccc} S & \rightarrow & A1B \\ A & \rightarrow & 0A \mid \epsilon \\ B & \rightarrow & 0B \mid 1B \mid \epsilon \end{array}$$

Give leftmost and rightmost derivations of the following strings:

2.1 a). 00101.

Leftmost:

 $S\Rightarrow A1B\Rightarrow 0A1B\Rightarrow 00A1B\Rightarrow 00\epsilon 1B\Rightarrow 001B\Rightarrow 0010B\Rightarrow 00101B\Rightarrow 00101\epsilon\Rightarrow 00101.$ Rightmost:

 $S\Rightarrow A1B\Rightarrow A10B\Rightarrow A101B\Rightarrow A101\epsilon\Rightarrow A101\Rightarrow 0A101\Rightarrow 00A101\Rightarrow 00\epsilon 101\Rightarrow 00101.$

2.2 b). 1001.

Leftmost:

 $S \Rightarrow A1B \Rightarrow \epsilon 1B \Rightarrow 10B \Rightarrow 100B \Rightarrow 1001B \Rightarrow 1001\epsilon \Rightarrow 1001.$

Rightmost:

 $S \Rightarrow A1B \Rightarrow A10B \Rightarrow A100B \Rightarrow A1001B \Rightarrow A1001\epsilon \Rightarrow \epsilon 1001 \Rightarrow 1001.$

2.3 c). 00011.

Leftmost:

 $S\Rightarrow A1B\Rightarrow 0A1B\Rightarrow 00A1B\Rightarrow 000A1B\Rightarrow 000\epsilon 1B\Rightarrow 00011B\Rightarrow 00011\epsilon\Rightarrow 00011.$

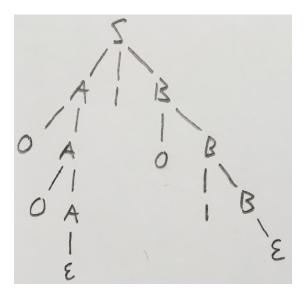
Rightmost:

 $S\Rightarrow A1B\Rightarrow A11B\Rightarrow A11\epsilon\Rightarrow 0A11\Rightarrow 00A11\Rightarrow 000A11\Rightarrow 000\epsilon 11\Rightarrow 00011.$

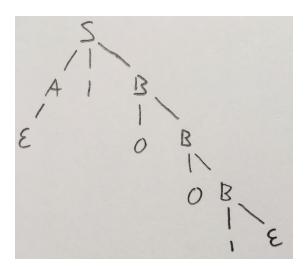
3 Problem 5.2.1

For the grammar and each of the strings in Exercise 5.1.2, give parse trees.

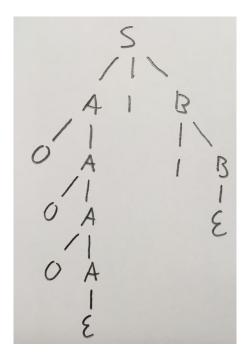
3.1 a). 00101.



3.2 b). 1001.



3.3 c). 00011.



4 Problem 5.4.5

This question concerns the grammar from Exercise 5.1.2

- 4.1 a). Show that this grammar is unambiguous
- 4.2 b). Find a grammar for the same language that *is* ambiguous, and demonstrate its ambiguity

Additional Exercises

5 Problem 1

Design a regular grammar without using a variable goes to ϵ rule $(A \to \epsilon)$ for the set of binary strings beginning with 1, ending with 0, and having even number of 0's and even number of 1's.

6 Problem 2

Design context-free grammars for the following languages

- 6.1 1). $L=\{a^nb^mc^{2n+m}\mid n,m>0\}$ where $\Sigma=\{a,b,c\}$
- 6.2 2). $L=\{a^nb^mc^i\mid m>n+i \text{ and } m,n,i\geq 0\}$, where $\Sigma=a,b,c$
- 6.3 3). $L=\{a^mb^n\mid 0\leq n\leq m\leq 3n\}$, where $\Sigma=\{a,b\}$

7 Problem 3

Let G be the grammar

$$\begin{array}{ccc} S & \rightarrow & ASB \mid ab \mid SS \\ A & \rightarrow & aA \mid \epsilon \\ B & \rightarrow & bB \mid \epsilon \end{array}$$

- 7.1 1). Give a leftmost derivation of aaabb
- 7.2 2). Give a rightmost derivation of aaabb
- 7.3 3). Show that G is ambiguous
- 7.4 4). Construct an unambiguous grammar equivalent to G

.