LRC with Multi-Writer Coherence Protocol

P1

Lock(L);

$x \leftarrow$ ⎫ x, y, z pages
$y \leftarrow$ ⎬ modified in c·s
$z \leftarrow$ ⎭

$\Rightarrow x_d, y_d, z_d$ } diffs

Unlock(L);

P2

$\rightarrow$ Lock(L);

$x \leftarrow x$

$Q \leftarrow$ ← fetch at point of access

invalidate x, y, z at lock acquisition (LRC)

P3

Lock(L);

$x \leftarrow \Rightarrow x_d'$

unlock(L);