Chapter 10 :: Functional Languages

Evaluation Order

Michael L. Scott



Evaluation Order Revisited

- Applicative order: evaluates all arguments before invoking function
 - what you're used to in imperative languages
 - usually faster
- Normal order: doesn't evaluate arg until you need it
 - sometimes faster
 - terminates if anything will (Church-Rosser theorem)



Evaluation Order (Example)

$$(and (not (= y 0)) (/ x y))$$



Why normal order may be slow

(define double (lambda (x) (+ x x))) (double (* 3 4))

Applicative Order

(double (* 3 4))

- \rightarrow (double 12)
- **→** (+ 12 12)
- **→** 24

Normal Order

(double (* 3 4))

- **→** (+ (* 3 4) (* 3 4))
- **→** (+ 12 (* 3 4))
- **→** (+ 12 12)
- **→** 24



Scheme Evaluation Order

- In Scheme
 - functions use applicative order defined with lambda
 - arguments are evaluated right to left
 - special forms (aka macros) use normal order defined with syntax-rules



Scheme Applicative Order Example

```
(define add (lambda (x) (+ x 20)))
(define min (lambda (x y) (if (\le x y) x y)))
(trace add)
(min (add 5) (add 20))
[Entering #[compound-procedure 4 add] Args: 20]
[40]
      <== #[compound-procedure 4 add] Args: 20]; <==
      means exiting this fct
[Entering #[compound-procedure 4 add] Args: 5]
[25]
      <== #[compound-procedure 4 add] Args: 5]
;Value: 25
```

Strict versus Non-strict Languages

- A *strict* language requires all arguments to be well-defined, so applicative order can be used
- A non-strict language does not require all arguments to be well-defined; it requires normal-order evaluation
- Scheme is strict for functions, but non-strict for special forms
- C is strict, except for boolean expressions



Forcing Normal Order in Scheme

- Use delay and force constructs
 - delay: creates an expression but does not evaluate it
 - force: forces the evaluation of a delayed expression
- Example

```
(define expr (delay (+ a 10)))
(define a 15)
(force expr) → 25
```



Forcing Normal Order in Scheme

```
(define naturals
  (letrec ((next (lambda (n)
     (cons n (delay (next (+ n 1))))))
    (next 1)))
(define head car)
(define tail (lambda (stream) (force (cdr stream))))
(head naturals) \rightarrow 1
```

(head (tail naturals) → 2 (head (tail (tail naturals)) → 3



Memoization

- Memoization: Technique saves an expression's result in some type of fast lookup structure
 - Thereafter references to the expression use this computed value
 - Brings performance of normal order evaluation within a constant factor of applicative order evaluation
- Spreadsheets use memoization

Example:

$$a10 = b10 + c10$$
 $b9 = 5$
 $b10 = 3 * b9$ $c9 = 10$
 $c10 = 8 * c9$

Memoization (Potential Problem)

- May not work properly in the presence of sideeffects
- Example:

```
(define x 5)
(define y 10)
(define (z (* x y))
(set! x 2)
(define (a (* x y))
```