

Lecture 0 - Review

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Sets - Notation

A **set** of **elements** a , b , and c is written as

$$\mathcal{S} = \{a, b, c\}.$$

To indicate that a is in set \mathcal{S} , we say ' a is an element of set \mathcal{S} ', and write

$$a \in \mathcal{S}.$$

Sets - Named Sets

- ▶ Natural numbers: $\mathbb{N} = \{1, 2, 3, 4, \dots\}$
- ▶ Integers: $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$
- ▶ Rational numbers: $\mathbb{Q} = \{x \mid x = a/b \text{ where } a, b \in \mathbb{Z}\}$
 - ▶ Notation: the \mid means 'such that'
 - ▶ So, the above says:
' \mathbb{Q} is the set of all numbers x such that x is equal to the quotient a/b where a and b are both integers.'
- ▶ Real numbers: \mathbb{R}
- ▶ Empty (Null) Set: $\emptyset = \{\}$

Sets - Operations and Subsets

Common Set Operations:

- ▶ Union:

$A \cup B$ is the set of all things that are in either A or B .

- ▶ Intersection:

$A \cap B$ is the set of all things that are in both A and B .

- ▶ Complement:

\overline{A} is the set of all things that are not in A .

Subsets:

- ▶ Subset:

$A \subseteq B$, every element in set A is also in set B .

- ▶ Proper Subset:

$A \subset B$, set A is a subset of set B and $A \neq B$.

Boolean Logic

Consider the **literals** P and Q that can have the boolean values True (\top) or False (\perp).

- ▶ Think statements like ‘all pigs can fly,’ or ‘my shirt is blue.’

Boolean Logic - Operations

- ▶ Disjunction (OR): $P \vee Q$ is true if either P **or** Q is true.
- ▶ Conjunction (AND): $P \wedge Q$ is true if both P **and** Q are true.
- ▶ Negation (NOT): $\neg P$ is true if P is **not** true.
- ▶ Exclusive-Or (XOR): $P \oplus Q$, P or Q but not both
- ▶ Implies: $P \Rightarrow Q$ means 'if P is true, then Q is true'
- ▶ Equality: $P \Leftrightarrow Q$ means ' P is true if and only if Q is true'
 - ▶ We write 'iff' as shorthand for 'if and only if':
' P is true iff Q is true'
 - ▶ To prove equality, must prove both directions:
Forward direction $P \Rightarrow Q$ and reverse direction $Q \Rightarrow P$.

Sequences and Tuples

A **sequence** is a list of objects in some order.

- ▶ Ex: $(3,65,2)$ or $(65,3,2)$. Note: these are not the same sequence because order matters.
- ▶ Note: a sequence can be infinite.

A **tuple** is a sequence with a finite length, and a tuple of length k is called a **k-tuple**.

- ▶ Ex: $(3,65,2)$ is a 3-tuple and $(5,4)$ is a 2-tuple.
- ▶ 2-tuples are also called ordered pairs.

Sets, Sequences, and Tuples

Sets and sequences can be elements of other sets and sequences.

- ▶ Ex: the **power set** of set A is the set of all subsets of A .

If $A = \{0, 1\}$, then the power set of A is the set:

$$\{ \emptyset, \{0\}, \{1\}, \{0, 1\} \}.$$

- ▶ Ex: The set of all ordered pairs with elements 0 and 1:

$$\{ (0, 0), (0, 1), (1, 0), (1, 1) \}.$$

The **Cartesian Product** (aka. Cross Product) of sets A and B :

$$A \times B = \{(a, b) | a \in A, b \in B\}.$$

Functions

A function is a mapping of inputs to outputs.

- ▶ The set of all possible inputs to a function is its **domain**.
- ▶ The outputs of a function are all contained in a set called its **range**.
- ▶ We write $f : D \rightarrow R$.
- ▶ Note: a function does not necessarily use all of the elements of the specified range. If it does, we say f maps D **onto** R .

Relations

A **predicate** or **property** is a function whose range is:

$$\{\text{True}, \text{False}\}.$$

A **relation** is a predicate whose domain is a set of k -tuples, and a **binary relation** is a relation whose input is an ordered pair.

- ▶ Infix notation: xRy means $R(x, y) = \text{True}$.
- ▶ Ex: 'less than', with infix notation using the symbol $<$
i.e. $x < y$.

Properties:

- ▶ R is **reflexive** if for every x : xRx .
- ▶ R is **symmetric** if for every x and y : $xRy \Rightarrow yRx$.
- ▶ R is **transitive** if for every x , y , and z : xRy and $yRz \Rightarrow xRz$.

A relation is an **equivalence relation** if it is reflexive, symmetric, and transitive.

Strings

We define an **alphabet** as a nonempty finite set of **symbols**. We typically use the letters Σ and Γ to denote alphabets.

- ▶ $\Sigma = \{0, 1\}$
- ▶ $\Gamma = \{a, b, c, d, e, f, g, h, i, j, k, l, m, n, o, p, q, r, s, t, u, v, w, x, y, z\}$

A **string** over an alphabet is a finite sequence of symbols from that alphabet.

- ▶ $\Sigma = \{0, 1\}$,
Ex: 0011, 01001, 0, 10
- ▶ The **empty string** is a string of length 0, denoted ε .

Properties:

- ▶ The length of a string w is written $|w|$.
- ▶ The reverse of a string w is written $w^{\mathcal{R}}$.
- ▶ A string z is a substring of w if z appears within w .
Ex: 'abba' is a substring of 'bbabba'

Strings

- ▶ Let x be the string $x_1x_2 \dots x_n$ of length n , where x_1, \dots, x_n are all symbols of the given alphabet.
- ▶ Let y be the string $y_1y_2 \dots y_m$ of length m , where y_1, \dots, y_m are all symbols of the given alphabet.
- ▶ The **concatenation** of x and y is written xy and is the string $x_1x_2 \dots x_ny_1y_2 \dots y_m$ of length $n + m$.
- ▶ We can concatenate a string to itself. If we do this k times, we denote the resulting string as x^k .

Languages

A **language** is a set of strings.

Proof Techniques

Direct Proof (Proof by Construction)

Say we want to prove that $P \Rightarrow Q$. A direct proof has 2 steps:

1. Assume that P is true.
2. Use P to show that Q is true.

Ex: Prove that if a and b are consecutive integers, then their sum $a + b$ is odd.

Proof: Assume that a and b are consecutive integers. This means that $b = a + 1$. So we can see that $a + b = 2a + 1$, which is clearly odd.

Proof Techniques

Proof by Contradiction

Say we want to prove that $P \Rightarrow Q$. A proof by contradiction has 2 steps:

1. Assume that P is true but that Q is not true.
2. Use this to demonstrate a contradiction, i.e. an obviously false conclusion.

Ex: Prove that if a and b are consecutive integers, then their sum $a + b$ is odd.

Proof: Assume that a and b are consecutive integers, but that their sum $a + b$ is not odd. Since the sum $a + b$ is not odd, there exists no integer k such that $a + b = 2k + 1$. However, since a and b are consecutive integers, $b = a + 1$ and so $a + b = 2a + 1$. This contradicts the previous sentence, and therefore if a and b are consecutive integers, their sum $a + b$ must be odd.

Proof Techniques

Proof by Induction

Say we want to prove that all of the elements of an ordered infinite set have a given property. Let's call this property $P(x)$, and so our goal is to prove that $P(x)$ is true for all x .

- ▶ Base Case:

We start with the base case, the first (and often simplest) subproblem: prove that $P(1)$ is true.

- ▶ Inductive Hypothesis:

Assume that $P(k)$ is true for some k .

- ▶ Inductive Step:

Using the assumption that $P(k)$ is true, prove that $P(k + 1)$ is true.

Why does this work?

Proof Techniques

Ex: Prove that if a and b are consecutive positive integers, then their sum $a + b$ is odd.

Proof:

- ▶ Base Case:

Let $a = 1$ and $b = 2$. Then $a + b = 3$, which is clearly odd.

- ▶ Inductive Hypothesis (IH):

Consider the integer k . Let $a = k$ and $b = k + 1$. Assume that their sum $a + b = k + (k + 1)$ is odd.

- ▶ Inductive Step:

Consider $k + 1$. Let $a = k + 1$ and $b = k + 2$ so

$$a + b = (k + 1) + (k + 2) = (k + 1) + k + 2 = k + (k + 1) + 2.$$

From the IH, we have assumed that $k + (k + 1)$ is odd.

Adding 2 to a number does not change whether it is odd.

Hence the sum $a + b = k + (k + 1) + 2$ is also odd.