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Q1:

a.) Original: $\exists x P(x) \wedge \exists x Q(x) \wedge A$

$\exists x \exists y (P(x) \wedge Q(y) \wedge A)$ Rewrite in prenex

b.) Original: $\neg(\forall x P(x) \wedge \forall x Q(x))$

De Morgans $\rightarrow \exists x \neg P(x) \vee \exists x \neg Q(x)$

Rewrite in prenex $\exists x \exists y (\neg P(x) \vee \neg Q(y))$

c.) Original: $\exists x P(x) \rightarrow \exists x Q(x)$

~~De Morgans left side and logical of equivalence~~

Logical Equivalence: $\exists x P(x) \rightarrow \exists x Q(x) \equiv \neg \exists x P(x) \vee \exists x Q(x)$

De Morgans Left: $\forall x \neg P(x) \vee \exists x Q(x)$

Rewrite in prenex: $\forall x \exists y (\neg P(x) \vee Q(y))$

Q2:

Let D: Logic is difficult

Let S: Many students like logic

Let E: Mathematics is easy

Assumptions: $D \vee \neg S$ $E \rightarrow \neg D$

a) $S \rightarrow \neg E$

$D \vee \neg S \equiv S \rightarrow D$ (Equivalence) $(p \rightarrow q \equiv \neg p \vee q)$ (commutative) $(p \vee \neg S \equiv \neg S \vee p)$

$E \rightarrow \neg D \equiv \neg(\neg D) \rightarrow \neg E \equiv D \rightarrow \neg E$ ($p \rightarrow q \equiv \neg q \rightarrow \neg p$)

$S \rightarrow D$ Using Hypothetical syllogism
 $D \rightarrow \neg E$
 $\therefore S \rightarrow \neg E$ Valid

b) $\neg E \rightarrow \neg S$

$D \vee \neg S \equiv \neg S \vee D \equiv S \rightarrow D$ (commutative and Equivalence)

If $\neg E$ is true and when $\neg S$ is false

$D \vee \neg S \equiv D \vee F \equiv D$ ✓

$E \rightarrow \neg D \equiv F \rightarrow \neg D$ always true ✓

Counter example

$D \equiv T$

$S \equiv T$

$E \equiv F$

Invalid

c) Prove $\neg E \vee D$

Assumptions:

$$E \rightarrow \neg D \equiv \neg E \vee \neg D$$

Equivalence

$$D \vee \neg S$$

$$\neg E \vee \neg D \quad \text{and} \quad \neg E \vee D$$

Counter example

$$D = F$$

$$(E = F \vee D) \quad (F \vee F = F) \quad (F \vee F = F)$$

d) Prove $\neg D \vee \neg E$

$$E \rightarrow \neg D \equiv \neg E \vee \neg D$$

Equivalence

$$\neg E \vee \neg D \equiv \neg D \vee \neg E$$

Commutative

Valid.

e) Prove $\neg S \rightarrow (\neg E \vee \neg D)$

We know $S \rightarrow D$ from (2a) $\neg E \vee \neg D$
 $S \rightarrow \neg E$ (2a)

$$\neg S \rightarrow (\neg E \vee \neg D)$$

We know $\neg E \vee \neg D$ is true or false

$\neg S \rightarrow \text{True}$ so $\neg S$ must be false ~~if stated~~
to show invalid

$$c) \neg S \rightarrow (\neg E \vee \neg D)$$

$\neg E \vee \neg D$ must be false for it to be invalid

$\neg E = F$ and $\neg D = F$ Must be to be invalid
so $E = T$ $D = T$

BUT we know $E \rightarrow \neg D$

SO $T \rightarrow \neg T$ is false

SO its valid because only scenario
 $\neg S \rightarrow (\neg E \vee \neg D)$ is invalid is when it
breaks assumptions

3. Part One

Assume true

$$\begin{array}{lll} P_1 \leftrightarrow P_4 & \text{shows} & P_1 \rightarrow P_4 \quad P_4 \rightarrow P_1 \\ P_2 \leftrightarrow P_3 & & P_1 \rightarrow P_3 \quad P_3 \rightarrow P_1 \\ P_1 \leftrightarrow P_3 & & P_3 \rightarrow P_2 \quad P_2 \rightarrow P_3 \end{array}$$

$$\begin{array}{lll} P_1 \rightarrow P_3 & P_1 \rightarrow P_2 & P_1 \rightarrow P_2 \\ \underline{P_3 \rightarrow P_2} & \underline{P_2 \rightarrow P_1} & \underline{P_2 \rightarrow P_1} \\ \therefore P_1 \rightarrow P_2 & \therefore P_2 \rightarrow P_1 & \therefore P_1 \leftrightarrow P_2 \end{array}$$

$$\begin{array}{l} P_4 \rightarrow P_1 \\ P_1 \rightarrow P_2 \\ \hline P_4 \rightarrow P_2 \end{array}$$

$P_4 \rightarrow P_1$ and P_1 implies P_2 and P_3

So P_4 will imply all

Just many Hypo Syllogism

So because $p_1 \leftrightarrow p_4$ and $p_1 \leftrightarrow p_3$ and $p_3 \leftrightarrow p_2$

p_4 is equivalent to p_1, p_2, p_3

Part Two:

$p_1 \rightarrow p_4$ Using Hypo Syllogism you

$p_4 \rightarrow p_2$ can easily show they are

$p_2 \rightarrow p_3$ all equivalent.

$p_3 \rightarrow p_1$

$p_3 \rightarrow p_1$

$p_1 \rightarrow p_4$

$p_4 \rightarrow p_2$

$\therefore p_1 \rightarrow p_2$ etc.

following this order

You show they are all equivalent

4.

Let x be a man

y be a city

z be a state in U.S

Let $P(x, y, z)$: x has visited city y in state z

$\exists x \forall y \forall z P(x, y, z)$

5. Prove no 2 integers k and j satisfy

$$k^2 = 4j + 2$$

$$k^2 = 2(2j + 1)$$

k is even because $k^2 = 2(\text{integer})$ is even

so let $k = 2m$

so

$$(2m)^2 = 4j + 2$$

$$4m^2 = 4j + 2$$

$$2m^2 = 2j + 1$$

$$2m^2 - 2j = 1$$

$$2(m^2 - j) = 1$$

m and j are integers

$$\text{so let } m^2 - j = L$$

$$2(L) = 1$$

$$L = \frac{1}{2}$$

$\frac{1}{2}$ is not an integer

contradicting given k and j are integers

so no integers k and j satisfy $k^2 = 4j + 2$

6.

Alice: C

John: 2)

Carlos: D

Disson: 7D

Test cases

Assume True

Alice did it	John	Carlos	Disson
F	F	T	F
T	F	T	T
F	F	F	T
T	T	T	F
2	①	③	2

a) John did it because that's only case where 1 True

b) Yes then Carlos did it because that's only case one lying

2. Prove $x < x^2 \iff x > 1$

$$x < x^2$$

$$0 < x^2 - x$$

$$0 < x(x-1)$$

only true iff $x > 0$ and $(x-1) > 0$
OR

Both negative $x < 0$ $x-1 < 0$

Given that x is positive real number

$x > 0$ and $(x-1) > 0$ must be true

$x-1 > 0$ is $x > 1$ therefore $x < x^2 \rightarrow x > 1$ is true

Part b

$$x > 1 \rightarrow x < x^2$$

$$x > 1$$

$$x^2 > x$$

$$x < x^2$$

must be true

So $x < x^2 \iff x > 1$

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P: Assume n is an integer

Q: Prove that n^2 ends in digit from set $\{0, 1, 4, 5, 6, 9\}$

$P \rightarrow Q$

$\neg Q \rightarrow \neg P$: If n^2 ends in digit from $\{2, 3, 7, 8\}$
then n is not an integer

a.