

Xiaoying Li
CS 325 -Winter 2020
Homework #8

● **Problem 1**

a.

(1) First-Fit

```
int First-Fit (int weight[], int number, int capacity) {  
    // Initialize bin's number to 0.  
    int binNumber = 0;  
    // Allocate an array to store the left room in bins, whose size is the number of items.  
    int* binRoom;  
    binRoom = new int[number];  
  
    // Pack items one by one.  
    for (int i = 0; i < number; i++) {  
        int j;  
        // Put each item into the first bin which it fits.  
        for (j = 0; j < binNumber; j++) {  
            // If such a bin is found, update this bin's left room and move to next item.  
            if (binRoom[j] >= weight[i]) {  
                binRoom[j] = binRoom[j] - weight[i];  
                break;  
            }  
        }  
        // If there is no available bin then open a new bin, set its left room,  
        // and update bins' number, then move to the next item.  
        if (j == binNumber) {  
            binRoom[binNumber] = capacity - weight[i];  
            binNumber++;  
        }  
    }  
  
    delete[] binRoom;  
    // After all items are packed into bins, return the number of bins.  
    return binNumber;  
}
```

The outer loop runs n times, the inner loop runs up to n times. Therefore, the running time of First-Fit algorithm is $O(n^2)$, where n is the number of items.

(2) First-Fit-Decreasing

```
int First-Fit-Decreasing(int weight[], int number, int capacity) {  
    // First sort the items in decreasing order by size.  
    Insertion-Sort(weight, number);  
    // Then use First-Fit on the resulting list and return the number of bins.  
    return First-Fit(weight, number, capacity);  
}
```

I used insertion sort to sort the items, and its running time is $O(n^2)$; the running time of First-Fit is also $O(n^2)$. Therefore, the running time of First-Fit-Decreasing algorithm is $O(n^2)$, where n is the number of items.

(3) Best-Fit

```
int Best-Fit(int weight[], int number, int capacity) {  
    // Initialize bin's number to 0.  
    int binNumber = 0;  
    // Allocate an array to store the left room in bins, whose size is the number of items.  
    int* binRoom;  
    binRoom = new int[number];  
  
    // Pack items one by one.  
    for (int i = 0; i < number; i++) {  
        int j;  
        // Initialize the least left room to bin's capacity plus 1,  
        // and the index of bin with the least left room to 0;  
        int least = capacity + 1;  
        int index = 0;  
  
        // Place the item into the bin which will leave the least room left over after the  
        // item is placed in the bin.  
        for (j = 0; j < binNumber; j++) {  
            // Compare every bin's left room after the item is placed in the bin if the  
            // item fits the bin to find the bin which will leave the least room left over  
            // after the item is placed in the bin.  
            if (binRoom[j] >= weight[i] && binRoom[j] - weight[i] < least) {  
                // Update the least left room and the index of the bin with the least  
                // left room.  
                index = j;  
                least = binRoom[j] - weight[i];  
            }  
        }  
    }  
}
```

```

// If the item does not fit in any bin, start a new bin, set its left room, and
// update bins' number, then move to the next item.
if (least == capacity + 1) {
    binRoom[binNumber] = capacity - weight[i];
    binNumber++;
}

// If such a bin is found, update this bin's left room and move to the next item.
else {
    binRoom[index] -= weight[i];
}
}

delete[] binRoom;
// After all items are packed into bins, return the number of bins.
return binNumber;
}

```

The outer loop runs n times, the inner loop runs up to n times. Therefore, the running time of Best-Fit algorithm is $O(n^2)$, where n is the number of items.

b. The answer is submitted to TEACH.

c.

Below are the screen shots of my randomly generated cases. In my code to produce the random inputs, the range of bin's capacity and number of items are both [1, 50], and the weight of each item is set to from 1 to not exceed the capacity of a bin. Totally 20 bin packing instances are generated.

```
Test Case 1:
capacity: 45
item number: 36
weight: 20 43 21 43 11 30 1 43 41 21 35 28 15 22 40 25 12 5 36 41 38 43 10 38 45 3 40 35 43 30 41 29 37 10 40 32
result: First Fit - 27; First Fit Decreasing - 26; Best Fit - 26

Test Case 2:
capacity: 21
item number: 13
weight: 13 6 12 15 20 6 13 7 11 21 4 20 12
result: First Fit - 9; First Fit Decreasing - 9; Best Fit - 9

Test Case 3:
capacity: 35
item number: 38
weight: 26 24 1 19 19 8 10 20 3 26 14 11 2 35 11 16 19 13 33 2 16 7 34 21 20 21 20 4 16 19 3 9 8 34 2 24 23 26
result: First Fit - 20; First Fit Decreasing - 19; Best Fit - 20

Test Case 4:
capacity: 26
item number: 9
weight: 23 8 19 26 21 19 16 5 2
result: First Fit - 6; First Fit Decreasing - 6; Best Fit - 6

Test Case 5:
capacity: 35
item number: 8
weight: 30 19 25 4 3 15 32 15
result: First Fit - 5; First Fit Decreasing - 5; Best Fit - 5

Test Case 6:
capacity: 45
item number: 9
weight: 18 1 4 17 38 42 14 8 15
result: First Fit - 4; First Fit Decreasing - 4; Best Fit - 4

Test Case 7:
capacity: 11
item number: 45
weight: 10 7 5 6 7 8 1 7 1 6 11 2 1 11 10 8 8 1 7 8 9 8 6 10 8 7 3 4 1 8 9 10 5 11 9 6 1 10 8 5 3 8 2 8 10
result: First Fit - 31; First Fit Decreasing - 31; Best Fit - 31

Test Case 8:
capacity: 29
item number: 47
weight: 5 27 19 8 13 19 8 14 22 18 14 7 4 1 23 4 12 28 17 5 6 26 28 8 12 16 24 7 29 1 29 12 29 4 28 15 22 3 11 5 5 7 13 8 27 28 10
result: First Fit - 25; First Fit Decreasing - 24; Best Fit - 25
```

```
Test Case 9:
capacity: 48
item number: 20
weight: 33 17 6 21 35 27 12 15 47 35 31 9 16 11 24 19 47 19 13 29
result: First Fit - 11; First Fit Decreasing - 10; Best Fit - 11

Test Case 10:
capacity: 23
item number: 34
weight: 14 1 18 23 8 16 19 21 23 10 8 3 6 16 15 13 3 2 17 2 16 8 19 8 21 19 7 14 20 8 20 5 22 10
result: First Fit - 21; First Fit Decreasing - 21; Best Fit - 21

Test Case 11:
capacity: 3
item number: 35
weight: 3 1 3 2 3 2 1 1 3 2 2 1 2 2 3 1 3 1 2 1 3 2 1 2 3 2 2 3 3 3 2 3 1 3 1
result: First Fit - 25; First Fit Decreasing - 25; Best Fit - 25

Test Case 12:
capacity: 3
item number: 44
weight: 2 1 2 3 1 1 1 2 2 1 2 2 1 1 3 2 2 1 2 1 1 1 3 3 3 3 3 3 2 2 3 1 3 2 1 1 1 3 3 2 2 1 2
result: First Fit - 29; First Fit Decreasing - 29; Best Fit - 29

Test Case 13:
capacity: 48
item number: 5
weight: 10 6 27 16 3
result: First Fit - 2; First Fit Decreasing - 2; Best Fit - 2

Test Case 14:
capacity: 40
item number: 9
weight: 20 11 20 28 37 30 25 34 4
result: First Fit - 7; First Fit Decreasing - 6; Best Fit - 7

Test Case 15:
capacity: 15
item number: 44
weight: 15 12 14 4 4 7 6 9 5 7 11 4 10 1 15 8 3 2 6 5 13 3 10 12 1 3 3 4 5 14 2 13 15 1 11 10 5 11 8 3 1 11 3 14
result: First Fit - 23; First Fit Decreasing - 22; Best Fit - 23

Test Case 16:
capacity: 46
item number: 26
weight: 13 39 46 21 2 45 18 23 36 38 41 37 10 2 10 3 31 31 20 28 2 36 18 24 10 11
result: First Fit - 15; First Fit Decreasing - 14; Best Fit - 14
```

```

Test Case 17:
capacity: 45
item number: 49
weight: 32 30 20 23 3 39 33 38 26 3 16 7 31 24 31 35 32 40 5 40 9 10 14 22 15 29 6 41 43 32 35 15 18 35 44 1 27 35 24 22 28 19 20 45 24 14 27 21 19
result: First Fit - 31; First Fit Decreasing - 29; Best Fit - 30

Test Case 18:
capacity: 34
item number: 46
weight: 22 28 18 33 20 20 31 19 32 31 7 19 5 15 33 16 19 14 21 31 17 7 6 23 32 20 3 28 11 23 14 32 27 21 22 26 8 1 23 10 6 24 15 28 2 20
result: First Fit - 31; First Fit Decreasing - 30; Best Fit - 30

Test Case 19:
capacity: 18
item number: 35
weight: 11 5 16 2 1 11 6 9 11 7 10 16 10 5 8 13 11 6 17 4 18 4 5 13 4 10 13 4 13 18 15 6 15 7 11
result: First Fit - 21; First Fit Decreasing - 20; Best Fit - 20

Test Case 20:
capacity: 12
item number: 46
weight: 10 12 3 2 5 12 4 11 1 6 3 6 11 7 6 3 8 8 2 4 10 4 6 12 10 2 9 10 2 9 7 8 5 5 1 8 9 7 2 12 8 8 8 5 9 10
result: First Fit - 29; First Fit Decreasing - 28; Best Fit - 29

```

The results for each algorithm are summarized as below, the algorithm with best performance is marked in blue for every test case.

| Case Number | First-Fit | First-Fit- Decreasing | Best-Fit |
|-------------|-----------|-----------------------|----------|
| 1 | 27 | 26 | 26 |
| 2 | 9 | 9 | 9 |
| 3 | 20 | 19 | 20 |
| 4 | 6 | 6 | 6 |
| 5 | 5 | 5 | 5 |
| 6 | 4 | 4 | 4 |
| 7 | 31 | 31 | 31 |
| 8 | 25 | 24 | 25 |
| 9 | 11 | 10 | 11 |
| 10 | 21 | 21 | 21 |
| 11 | 25 | 25 | 25 |
| 12 | 29 | 29 | 29 |
| 13 | 2 | 2 | 2 |
| 14 | 7 | 6 | 7 |
| 15 | 23 | 22 | 23 |
| 16 | 15 | 14 | 14 |
| 17 | 31 | 29 | 30 |
| 18 | 31 | 30 | 30 |
| 19 | 21 | 20 | 20 |
| 20 | 29 | 28 | 29 |

Obviously, First-Fit-Decreasing algorithm performs best since it performs best in 11 cases over 20 test cases. Best-Fit algorithm performs as well as First-Fit-Decreasing in 4 cases, but was never better than First-Fit-Decreasing. And First-Fit algorithm performs worst, which never performs better than the other two algorithms. But there are 9 cases over 20 test cases where the three algorithm's performance is same.

Therefore, First-Fit-Decreasing algorithm performs better, in 11 cases over 20 test cases (55% times).

● Problem 2

a.

Code and result of the integer program from LINDO:

| | |
|---|--|
| <pre> MIN Y1 + Y2 + Y3 + Y4 + Y5 + Y6 ST A1 + A2 + A3 + A4 + A5 + A6 = 1 B1 + B2 + B3 + B4 + B5 + B6 = 1 C1 + C2 + C3 + C4 + C5 + C6 = 1 D1 + D2 + D3 + D4 + D5 + D6 = 1 E1 + E2 + E3 + E4 + E5 + E6 = 1 F1 + F2 + F3 + F4 + F5 + F6 = 1 4A1 + 4B1 + 4C1 + 8D1 + 8E1 + 8F1 - 10Y1 <= 0 4A2 + 4B2 + 4C2 + 8D2 + 8E2 + 8F2 - 10Y2 <= 0 4A3 + 4B3 + 4C3 + 8D3 + 8E3 + 8F3 - 10Y3 <= 0 4A4 + 4B4 + 4C4 + 8D4 + 8E4 + 8F4 - 10Y4 <= 0 4A5 + 4B5 + 4C5 + 8D5 + 8E5 + 8F5 - 10Y5 <= 0 4A6 + 4B6 + 4C6 + 8D6 + 8E6 + 8F6 - 10Y6 <= 0 END INT Y1 INT Y2 INT Y3 INT Y4 INT Y5 INT Y6 INT A1 INT A2 INT A3 INT A4 INT A5 INT A6 INT B1 INT B2 INT B3 INT B4 INT B5 INT B6 INT C1 INT C2 INT C3 INT C4 INT C5 INT C6 INT D1 INT D2 INT D3 INT D4 INT D5 INT D6 INT E1 INT E2 INT E3 INT E4 INT E5 INT E6 INT F1 INT F2 INT F3 INT F4 INT F5 INT F6 </pre> | <pre> LP OPTIMUM FOUND AT STEP 55 OBJECTIVE VALUE = 3.00000000 NEW INTEGER SOLUTION OF 3.00000000 AT BRANCH 0 PIVOT 55 RE-INSTALLING BEST SOLUTION... OBJECTIVE FUNCTION VALUE 1) 3.000000 VARIABLE VALUE REDUCED COST Y1 1.000000 1.000000 Y2 0.000000 1.000000 Y3 0.000000 1.000000 Y4 0.000000 1.000000 Y5 1.000000 1.000000 Y6 1.000000 1.000000 A1 0.000000 0.000000 A2 0.000000 0.000000 A3 0.000000 0.000000 A4 0.000000 0.000000 A5 1.000000 0.000000 A6 0.000000 0.000000 B1 1.000000 0.000000 B2 0.000000 0.000000 B3 0.000000 0.000000 B4 0.000000 0.000000 B5 0.000000 0.000000 B6 0.000000 0.000000 C1 0.000000 0.000000 C2 0.000000 0.000000 C3 0.000000 0.000000 C4 0.000000 0.000000 C5 0.000000 0.000000 C6 1.000000 0.000000 D1 0.000000 0.000000 D2 0.000000 0.000000 D3 0.000000 0.000000 D4 0.000000 0.000000 D5 1.000000 0.000000 D6 0.000000 0.000000 E1 0.000000 0.000000 E2 0.000000 0.000000 E3 0.000000 0.000000 E4 0.000000 0.000000 E5 0.000000 0.000000 E6 1.000000 0.000000 F1 1.000000 0.000000 F2 0.000000 0.000000 F3 0.000000 0.000000 F4 0.000000 0.000000 F5 0.000000 0.000000 F6 0.000000 0.000000 ROW SLACK OR SURPLUS DUAL PRICES 2) 0.000000 0.000000 3) 0.000000 0.000000 4) 0.000000 0.000000 5) 0.000000 0.000000 6) 0.000000 0.000000 7) 0.000000 0.000000 8) 0.000000 0.000000 9) 0.000000 0.000000 10) 0.000000 0.000000 11) 0.000000 0.000000 12) 0.000000 0.000000 13) 0.000000 0.000000 NO. ITERATIONS= 55 BRANCHES= 0 DETERM.= 1.000E 0 </pre> |
|---|--|

Interpretation:

The result shows that $Y1=1$, $Y5=1$, $Y6=1$, $A5=1$, $B1=1$, $C6=1$, $D5=1$, $E6=1$, and $F1=1$, so the 6 items can be packed into at least 3 bins:

$Y1 - B1$ and $F1$: first bin - items of weight 4 and 6;

$Y5 - A5$ and $D5$: second bin - items of weight 4 and 6;

$Y6 - C6$ and $E6$: third bin - items of weight 4 and 6.

b.

Code and result of the integer program from LINDO:

```

MIN Y1 + Y2 + Y3 + Y4 + Y5
ST
    A1 + A2 + A3 + A4 + A5 = 1
    B1 + B2 + B3 + B4 + B5 = 1
    C1 + C2 + C3 + C4 + C5 = 1
    D1 + D2 + D3 + D4 + D5 = 1
    E1 + E2 + E3 + E4 + E5 = 1
    20A1 + 10B1 + 15C1 + 10D1 + 5E1 - 20Y1 <= 0
    20A2 + 10B2 + 15C2 + 10D2 + 5E2 - 20Y2 <= 0
    20A3 + 10B3 + 15C3 + 10D3 + 5E3 - 20Y3 <= 0
    20A4 + 10B4 + 15C4 + 10D4 + 5E4 - 20Y4 <= 0
    20A5 + 10B5 + 15C5 + 10D5 + 5E5 - 20Y5 <= 0
END
    INT Y1
    INT Y2
    INT Y3
    INT Y4
    INT Y5
    INT A1
    INT A2
    INT A3
    INT A4
    INT A5
    INT B1
    INT B2
    INT B3
    INT B4
    INT B5
    INT C1
    INT C2
    INT C3
    INT C4
    INT C5
    INT D1
    INT D2
    INT D3
    INT D4
    INT D5
    INT E1
    INT E2
    INT E3
    INT E4
    INT E5

```

```

LP OPTIMUM FOUND AT STEP 19
OBJECTIVE VALUE = 3.00000000

NEW INTEGER SOLUTION OF 3.00000000 AT BRANCH 0 PIVOT 19
RE-INSTALLING BEST SOLUTION...

OBJECTIVE FUNCTION VALUE
1) 3.000000

VARIABLE    VALUE    REDUCED COST
Y1          1.000000    1.000000
Y2          1.000000    1.000000
Y3          1.000000    1.000000
Y4          0.000000    1.000000
Y5          0.000000    1.000000
A1          0.000000    0.000000
A2          1.000000    0.000000
A3          0.000000    0.000000
A4          0.000000    0.000000
A5          0.000000    0.000000
B1          1.000000    0.000000
B2          0.000000    0.000000
B3          0.000000    0.000000
B4          0.000000    0.000000
B5          0.000000    0.000000
C1          0.000000    0.000000
C2          0.000000    0.000000
C3          1.000000    0.000000
C4          0.000000    0.000000
C5          0.000000    0.000000
D1          1.000000    0.000000
D2          0.000000    0.000000
D3          0.000000    0.000000
D4          0.000000    0.000000
D5          0.000000    0.000000
E1          0.000000    0.000000
E2          0.000000    0.000000
E3          1.000000    0.000000
E4          0.000000    0.000000
E5          0.000000    0.000000

ROW    SLACK OR SURPLUS    DUAL PRICES
2)     0.000000           0.000000
3)     0.000000           0.000000
4)     0.000000           0.000000
5)     0.000000           0.000000
6)     0.000000           0.000000
7)     0.000000           0.000000
8)     0.000000           0.000000
9)     0.000000           0.000000
10)    0.000000           0.000000
11)    0.000000           0.000000

NO. ITERATIONS= 19
BRANCHES= 0 DETERM.= 1.000E 0

```

Interpretation:

The result shows that $Y1=1$, $Y2=1$, $Y3=1$, $A2=1$, $B1=1$, $C3=1$, $D1=1$, and $E3=1$, so the 5 items can be packed into at least 3 bins:

Y1 – B1 and D1: first bin - items of weight 10 and 10;

Y2 – A2: second bin – item of weight 20;

Y3 – C3 and E3: third bin – items of weight 15 and 5.