

Assignment 3
CS232 Fall 2021
Xiaoying Li
12/07/2021

Problem 1

$$\lambda = 4 \text{ pkt/s}, \mu = 8 \text{ pkt/s} \rightarrow \rho = \frac{\lambda}{\mu} = \frac{4}{8} = \frac{1}{2}$$

(a)

The probability that the next packet is dropped is the probability that the number of packets in the system is $K=5$. Thus,

$$P[\text{Next Packet Dropped}] = p_K = \frac{(1-\rho)\rho^K}{1-\rho^{K+1}} = \frac{(1-\frac{1}{2})(\frac{1}{2})^5}{1-(\frac{1}{2})^{5+1}} = \frac{1}{63}$$

(b)

$$E[N] = \frac{\rho}{1-\rho} - \frac{(K+1)\rho^{K+1}}{1-\rho^{K+1}} = \frac{\frac{1}{2}}{1-\frac{1}{2}} - \frac{(5+1)(\frac{1}{2})^{5+1}}{1-(\frac{1}{2})^{5+1}} = \frac{19}{21}$$

$$E[T] = \frac{E[N]}{\lambda(1-P_K)} = \frac{\frac{19}{21}}{4(1-\frac{1}{63})} = \frac{57}{248} \text{ second}$$

The expected time a packet needs to go through the system is $\frac{57}{248}$ second.

Problem 2

$$\lambda = 4 \text{ pkt/s}, \mu = 5 \text{ pkt/s} \rightarrow \rho = \frac{\lambda}{\mu} = \frac{4}{5}$$

(a)

$$E[T] = \frac{1/\mu}{1-\rho} = \frac{\frac{1}{5}}{1-\frac{4}{5}} = 1 \text{ second}$$

The expected time a packet spends in the system is 1 second.

(b)

$$E[W] = \frac{(1/\mu)\rho}{(1-\rho)} = \frac{\frac{1}{5} \cdot \frac{4}{5}}{1 - \frac{4}{5}} = \frac{4}{5} \text{ second}$$

The time a packet is expected to wait in the buffer before getting serviced is $\frac{4}{5}$ second.

Problem 3

(a)

When having a dedicated server for each application, we obtain

$$\lambda = 2 \text{ pkt/s}, \mu = 4 \text{ pkt/s} \rightarrow \rho = \frac{\lambda}{\mu} = \frac{2}{4} = \frac{1}{2}$$

$$E[T_{seperate}] = \frac{1/\mu}{1-\rho} = \frac{\frac{1}{4}}{1 - \frac{1}{2}} = \frac{1}{2} \text{ second}$$

When having a single powerful server serving all the five applications simultaneously, we obtain $m = 5$,

$$E[T_{combined}] = \frac{1}{m} E[T_{seperate}] = \frac{1}{5} \cdot \frac{1}{2} = \frac{1}{10} \text{ second}$$

Obviously, $E[T_{seperate}] > E[T_{combined}]$. Therefore, having a single powerful server serving all the five applications simultaneously is better.

(b)

$$n = 2, \rho = \frac{1}{2}$$

$$P\{N(t) = n\} = p_n = (1-\rho)\rho^n \rightarrow P\{N(t) = 2\} = p_2 = \left(1 - \frac{1}{2}\right)\left(\frac{1}{2}\right)^2 = \frac{1}{8}$$

The probability that there are two packets in the system is $\frac{1}{8}$.

Problem 4

(a)

$$\lambda = \frac{1}{2} \text{ call/s}, \mu = 2 \text{ call/s} \rightarrow a = \frac{\lambda}{\mu} = \frac{\frac{1}{2}}{2} = \frac{1}{4}$$

The probability that the next phone call is dropped is the probability that the number of phone calls in the server is $c=5$. Thus,

$$P[N(t) = n] = p_n = \frac{\frac{a^n}{n!}}{\sum_{k=0}^c \frac{a^k}{k!}}, \text{ for } n = 0, 1, \dots, c$$

$$\rightarrow P[N(t) = 2] = p_2 = \frac{\frac{(\frac{1}{4})^2}{2!}}{\sum_{k=0}^2 \frac{a^k}{k!}} = \frac{\frac{1}{32}}{1 + \frac{(\frac{1}{4})^1}{1!} + \frac{(\frac{1}{4})^2}{2!}} = \frac{1}{41}$$

The probability that the next phone call is dropped is $\frac{1}{41}$.

(b)

$$P[\text{second phone call is received before the first one terminates}]$$

$$= P[T(\text{second phone call is received}) < T(\text{first phone call terminates})]$$

$$= \frac{\lambda}{\lambda + \mu} = \frac{\frac{1}{2}}{\frac{1}{2} + 2} = \frac{1}{5}$$

Problem 5

(a)

$$\lambda = 5 \text{ pkts/s}, \mu = 10 \text{ pkts/s}$$

$$\rightarrow \text{traffic load } a = \frac{\lambda}{\mu} = \frac{5}{10} = \frac{1}{2} \text{ Erlangs}$$

The probability of a packet being dropped upon arrival is the probability of the number of packets in the system is $c=1$. Thus,

$$P[N(t) = n] = p_n = \frac{\frac{a^n}{n!}}{\sum_{k=0}^c \frac{a^k}{k!}}, \text{ for } n = 0, 1, \dots, c$$

$$\rightarrow P_b = P[N(t) = 1] = p_1 = \frac{\frac{(\frac{1}{2})^1}{1!}}{\sum_{k=0}^1 \frac{a^k}{k!}} = \frac{\frac{1}{2}}{1 + \frac{1}{2}} = \frac{1}{3}$$

$$\rightarrow \text{carried load} = a(1 - P_b) = \frac{1}{2} \cdot \left(1 - \frac{1}{3}\right) = \frac{1}{3} \text{ Erlangs}$$

(b)

$$\lambda = 5 \text{ pkts/s}, P_b = \frac{1}{3}, E[W] = 2.5 \text{ seconds}$$

$$\rightarrow E[N_q] = \lambda(1 - P_b)E[W] = 5 \cdot \left(1 - \frac{1}{3}\right) \cdot 2.5 = \frac{25}{3}$$

The average number of packets in the buffer is $\frac{25}{3}$.