

Lecture Notes

Xinyi Li

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We were proving Theorem 5.9 For all $S \in \mathcal{C}$

$$\dim(S) = \liminf_{w \rightarrow S} \dim(w)$$

\leq last time

To see that $\dim(S) \geq \liminf_{w \rightarrow S} \dim(w)$, let s' and s'' be rational numbers such that $s' > s'' > \liminf_{w \in S} \dim(w)$. It suffices to show that there exist only $w \sqsubseteq S$ for which $\dim(w) \leq s'$. Since $s'' > \dim(S)$, there is an algorithmic s'' -supergale d such that $S \in S^\infty[d]$. Define $d' : \mathcal{T} \rightarrow [0, \infty)$ by

$$d'(x) = \begin{cases} d(x) & \text{if } x \in \{0, 1\}^* \\ (2^{s'-s''})d(w) & \text{if } x = w\Box \in \{0, 1\}^*\Box \end{cases}$$

Then d' is an algorithmic s' -supergale, so if for each $s \in [0, \infty)$ we define $\tilde{d}^{(s)} : \mathcal{T} \rightarrow [0, \infty)$ by $\tilde{d}^{(s)}(x) = d'(x)$ for all $x \in \{0, 1\}^*$, $\tilde{d}^{(s)}(x) = 2^{(s-s')}|x|d'(w)$, then the family $\tilde{d} = \{\tilde{d}^{(s)} | s \in [0, \infty)\}$ is an algorithmic termgale. It follows by optimality of \underline{d}_\Box that there is a constant $\alpha > 0$ such that, for all $s \in [0, \infty)$ and all $w \in \{0, 1\}^*$,

$$\underline{d}_\Box^{(s)}(w) > \alpha \tilde{d}^{(s)}(w\Box)$$

Since $S \in S^\infty[d]$, there are infinitely many $w \sqsubseteq S$ such that $\alpha(2^{s'-s''})d(w) > 1$. For all such w , we have

$$\begin{aligned} \underline{d}_\Box^{(s')}(w) &\geq \alpha \tilde{d}^{(s')}(w\Box) \\ &= \alpha d'(w\Box) \\ &= \alpha(2^{s'-s''})d(w) \\ &> 1 \end{aligned}$$

whence $\dim(w) \leq s'$. ■

Theorem (6.1). There is a constant $c \in \mathbb{N}$ such that for all $w \in \{0, 1\}^*$,

$$|K(w) - |w|\dim(w)| \leq c$$

Proof. Let \underline{m} be an optimal algorithmic subprobability measure. ¹ The key fact is that for all $w \in \{0, 1\}^*$ and $s \in [0, \infty)$,

$$\begin{aligned} d[\underline{m}](w\Box) > 1 &\Leftrightarrow 2^{s|w\Box|}\underline{m}(w) > 1 \\ &\Leftrightarrow s > \frac{1}{1-|w|} \log \frac{1}{\underline{m}(w)} \end{aligned}$$

so

$$\dim_{d[\underline{m}]}(w) = \frac{1}{1-|w|} \log \frac{1}{\underline{m}(w)}$$

□

¹ $(K(x) \approx \log \frac{1}{\underline{m}(x)})$