

A Genetic Algorithm for Task Offloading problem in Vehicular Edge Computing

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Abstract—In this paper, we consider task offloading in vehicular edge computing systems. The limited battery limits the mileage of electric vehicles. For a longer driving distance, the task of the vehicle can not only be performed locally, but also offload task to save energy. In this way, the energy consumption of the system is reduced and the energy utilization rate of the whole system is improved. The task offloading involves optimization variables which are where to offload tasks and the frequency, And The optimization variables are coupled. So it is hard to decide how to allocate tasks. We propose that we can solve this problem by using genetic algorithm. Specifically, our experiment shows between 20% and 25% energy savings compared to the baseline which the tasks executes workload locally.

Index Terms—Vehicular Edge Computing; Task Offloading; Genetic Algorithm

I. INTRODUCTION

Clean energy is an important trend in the future. Electric energy is representative of clean energy, so more and more people choose electric vehicles. In the era of IoT, almost all devices need to be connected to the Internet, including automobiles.

Electric intelligent vehicles will be an integral part of urban transportation in the future. However, before significant changes in battery technology, battery capacity is the shortcoming of intelligent electric vehicles [1]. Many edge computing systems have battery limitations [2]. Limited battery capacity will cause "range anxiety". Many methods have been proposed to solve the mileage problem, such as using blockchain technology to record the real-time data of vehicle travel [3] and considering complete vehicle energy management [4]. The vehicular coasts are part of driving mode, reducing energy consumption in [5].

Because of the large amount of data processed in driving, the central processing unit(CPU) will consume considerable energy, dramatically impacting the driving mileage.

Some researchers proposed sending the task to mobile edge computing servers to save energy. However, if all data is uploaded to the cloud server for processing, the response time might be too long, and the requirements for low delay will not be met [6]. Cloud computing and edge computing have different focuses [7]. Cloud

computing is characterized by centralization and scale. Edge computing has the advantages of low latency and low power consumption. In paper [8], Liu et al. propose a system to guarantee task-low delay performance. In addition, the current bandwidth and storage cannot satisfy the requirements for transmitting all data.

Yu et al. [9] focus on virtual machine resource allocation in vehicular edge clouds to solve this problem. Liu et al. [10] consider dynamic requirements and resource constraints and classify all tasks into four types of pending lists. Then the tasks in each list will be offloaded to different nodes according to their features.

Task offloading is the main problem in edge computing [11]. The recently popular edge computing technology has been introduced. We can assign tasks not only to the nearby edge server but also to resource-rich users nearby. We propose to assign tasks to nearby vehicles to minimize the system's energy consumption. This offloading process forms a three-tier architecture, including vehicular cloud, edge servers, roadside units, and central cloud processors. The edge servers are the controller of the vehicular cloud. Moreover, the central cloud is responsible for global scheduling.

The problem is non-linear integer programming. We propose a genetic algorithm to solve the task offloading problem. In recently, genetic algorithm is applied to various fields. In [12], Sun et al. use genetic algorithm to design CNN architectures. Zhang et al. [13] use it to minimize energy consumption. Yoon et al. [14] propose an efficient genetic algorithm using a novel normalization method to solve maximum coverage deployment problem in wireless sensor networks. In a complex 3D environment, Roberge et al. [15] use it to compute feasible and quasi-optimal trajectories for fixed wing UAVs.

Our experimental results show that energy saving is apparent when the number of vehicles is less than 30.

The remainder of this paper is organized as follows. Section II introduces the system model and problem formulation. Section III describes the genetic algorithm. Section IV describe our experiment and analysis the result. Section V is conclusion and future work.

TABLE I: Notations in This Paper

Parameter	Interpretation
T	The resource sharing time
C_{it}	The CPU capacity
r_i	The workload of vehicle i
r'_{jt}	The resource requester j need
M_{it}	Millions Instructions Per Second(MIPS)
θ_i	Estimated parameters
E_{it}^{blnc}	Energy consumed by the vehicle i
E_i^{rec}, E_i^{send}	Energy received or send by the vehicle i

II. SYSTEM MODEL AND PROBLEM FORMULATION

We first describe our system, including network architecture and computation models in this section.

A. Network Architecture

Similar to [9], as shown in Figure 1, the hierarchical architecture consisting vehicular cloud, edge cloud, and central cloud.

The central cloud is responsible for global scheduling, and the tasks it performs include complicated computation and global decision.

The edge cloud is the controller of vehicular cloud, and is responsible for creating, maintaining, and deleting vehicular cloud.

The vehicular cloud consists of vehicles which share their computation resources. Each vehicle can access the edge cloud and convey the information to it. The vehicles save the energy through offloading task to other vehicles. It can improve the overall resource utilization.

B. Computation Model

We define resource sharing time as $\mathcal{T} = \{1, \dots, T\}$. And the number of vehicles is defined as N . We use a tuple $J_{it} = \{r_{it}, r'_{it}, d_{it}\}$ to represent the task of vehicle i at time slot t . And d_{it} is the task's data size.

Tasks can be divided into tasks that can be offloaded and tasks that cannot be offloaded and must be executed locally. At time slot t , r_{it} is task required to complete the local by the vehicle i , r'_{it} is task which can offload to others in the vehicle i , and their values are expressed in the number of instructions.

To facilitate the description of the model, we denote $\mathbf{X}_t = \{x_{ij}\} \in \{0, 1\}^{N \times N}$ allocation matrix. If $x_{ij} = 1$, it represents that vehicle i performs the task of offloading vehicle j . Note that if $x_{ii} = 1$, all tasks of vehicle i are executed locally.

We characterize the computing capacity of vehicle i by C_{it} at time t . When a vehicle is selected as a provider (P) at time t , the computing resources of all requester should be less than its computing capacity.

$$\sum_{j=1}^N x_{ij}^t \cdot r'_{jt} \leq C_{it}, i = 1, \dots, N \quad (1)$$

A crucial step is how to quantify the computing resource. The computing capacity is defined based on the

number of instructions executed per unit time (MIPS, M_{it}).

$$C_{it} = M_{it} \cdot \Delta T - r_{it} \quad (2)$$

where ΔT is the duration of resource sharing, and is a smaller time scale than time of establishing vehicular network. After ΔT seconds, the offloading matrix changes.

For a single core CPU, the number of instructions executed per unit time (m_{it}) and CPU frequency (f_{it}) have the following relationship:

$$M_{it} = v_i \cdot f_{it} + \theta_i \quad (3)$$

Where v_i and θ_i are parameters to be estimated.

As a result, the CPU capacity C_{it} in Equation (2) can be calculated following:

$$C_{it} = (v_i \cdot f_{it} + \theta_i) \times \Delta T - r_{it} \quad (4)$$

C. Energy Consumption Model

When we consider the energy consumption, we divide it into two parts: (1)the energy required for calculation, (2)the energy required for transmission.

According to [16], [17], the energy consumption is computed as follows:

$$E = \lambda_i \cdot f_{it}^3 \cdot \Delta T \quad (5)$$

where f_{it} is the frequency of the CPU in vehicle i at time slot t . If a vehicle is selected as the requester(R), its frequency will decrease to f'_{it} . Because its tasks are assigned, so the energy saved is:

$$E_i^{save} = (\lambda_i \cdot f_{it}^3 - \lambda_i \cdot f'^3_{it}) \times \Delta T, \forall i \in R \quad (6)$$

The transmission energy consumption is linear with the transmission time which depends on the ratio between the data size and the data transmission rate(b_{ij}).

$$E = P_0 \cdot \frac{d_{it}}{b_{ij}} \quad (7)$$

where P_0 is the transmission power, although the power can vary [18], [19], [20], for simplicity, we use a fixed value in this article.

Maximum data transmission rate(b) is given by Shannon theorem.

$$b_{ij} = W \log(1 + \text{SNR}) \quad (8)$$

where SNR is Signal-to-noise Ratio, and W is the channel bandwidth. Because they are related to each actual channel, we regard it as a constant in this paper.

For each vehicle, the energy required for receiving information is:

$$E_i^{rec} = \sum_{j=1, j \neq i}^N x_{ij}^t \cdot P_0 \cdot \frac{d_{jt}}{W \log(1 + \text{SNR})}, \quad i = 1, \dots, N \quad (9)$$



Fig. 1: Vehicular networks architecture.

For each vehicle, the energy required for sending information is:

$$E_i^{send} = \sum_{i=1, i \neq j}^N x_{ij}^t \cdot P_0 \cdot \frac{d_{it}}{W \log(1 + SNR)}, \quad j = 1, \dots, N \quad (10)$$

Define E_t^{blnc} is all the energy consumption of the vehicle in the time slot t .

$$\begin{aligned} E_{it}^{blnc} &= \sum_{j=1, j \neq i}^N x_{ij}^t \cdot P_0 \cdot \frac{d_{jt}}{W \log(1 + SNR)} \\ &+ \sum_{j=1, j \neq i}^N x_{ji}^t \cdot P_0 \cdot \frac{d_{it}}{W \log(1 + SNR)} \\ &+ \lambda_i \cdot f_{it}^3 \cdot \Delta T, \quad i = 1, \dots, N \end{aligned} \quad (11)$$

The objective of our algorithm is to minimize the quadratic power of energy over all vehicles. The goal is to balance and minimize energy consumption, because the square of the minimum drawing energy consumption.

$$\min \sum_{t=1}^T \sum_{i=1}^N (E_{it}^{blnc})^2 \quad (12)$$

The problem can be formulated as:

$$\min \sum_{t=1}^T \sum_{i=1}^N (E_{it}^{blnc})^2 \quad (13)$$

$$\begin{aligned} \text{s.t. } E_{it}^{blnc} &= \sum_{j=1, j \neq i}^N x_{ij}^t \cdot P_0 \cdot \frac{d_{jt}}{W \log(1 + SNR)} \\ &+ \sum_{j=1, j \neq i}^N x_{ji}^t \cdot P_0 \cdot \frac{d_{it}}{W \log(1 + SNR)} \\ &+ \lambda_i \cdot f_{it}^3 \cdot \Delta T, \quad i = 1, \dots, N \end{aligned} \quad (14)$$

$$\sum_{j=1}^N x_{ij}^t \cdot r'_{jt} \leq C_{it} \quad (15)$$

$$\sum_{i=1}^N x_{ij}^t \geq 1 \quad (16)$$

$$f_{it} \geq 0 \quad (17)$$

$$x_{it}^t \in \{0, 1\} \quad (18)$$

As mentioned above, the Equation(15) shows all the amount of requested cannot exceed available capacity of vehicle i . The Equation(16) shows that the task of vehicle i which can be offloaded must be executed whether by the vehicle i or others. f is the frequency and must be positive number.

D. Example

$$\mathbf{X} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

The \mathbf{X} is a example of allocation matrix. Vehicular set is $\{A, B, C, D, E\}$, the vehicle A, B and D don't assign and load other vehicular tasks, Therefore, they do not generate transmission energy consumption. And vehicle C execute the offloading task from vehicle E. And the sum of each column is greater than or equal to one.

E. Constraints

To ensure the Quality of Service(QoS) between the requester and provider, the following distance constraints need to be met:

$$\min l_{ij}^t < \delta, \quad i = 1, \dots, N, \forall x_{ij}^t = 1 \quad (19)$$

where l_{ij}^t is the distance between vehicle i and j .

If the constraint is not satisfied, the vehicle will exit resource sharing, and the i th row and the i th column of the allocation matrix are all set to zero while $x_{ii}^t = 1$

III. ALGORITHMS

We design a greedy algorithm to minimize the energy consumption of all time slots. and we use genetic algorithm 1 to calculate the minimum value of each moment.

Genetic algorithm is a adaptive heuristic optimization method, which is proposed by John Holland in 1970s [21]. Based on Darwin's theory of natural selection, it simulates the process of natural selection and reproduction to solve optimization problems. Similarly to individuals with favorable variation surviving, the solutions in the algorithm imitate the behavior of chromosomes, such as the mutations of genes and the crossovers of chromosomes.

Algorithm 1 is the pseudo-code of genetic algorithm in our experiment. The algorithm has as input the vector of vehicles with their request size, r'_i , and their local workload $rloc$. The output consists of the allocation matrix \mathbf{X}_t , the frequency f_t , and the energy consumption E_t in the current time.

At the initial stage of the algorithm, it will randomly generate a set of feasible solutions which is the first generation of chromosomes(Line2). After that, $\mathbf{X}(t)$ and $f(t)$ is encoded as a string like a chromosome. Then the fitness function is used to calculate the fitness degree for each chromosome(Line5-7), and the probability of each chromosome being selected in the next evolution is calculated according to the fitness degree(Line8). The chromosomes with lower fitness will be eliminated, and only the chromosomes with higher fitness will be retained(Line9). Divide the individual into two parts at random(Line10). A certain position of these two chromosomes is cut off and spliced together to generate a new chromosome. As a result, this new chromosome contains a certain number of genes of both father and mother(Line11-13).

The solution is easier to get the local optimal solution, and there is no way to achieve the global optimal solution. In order to solve this problem, we need to introduce the process of mutation(Line14-16). The next step is to decode the string to the original solution(Line17). Finally we calculate all the energy consumption(Line20-22).

IV. EXPERIMENT

We set the parameters of the experiment and analyzed the result of the experiment in the section.

A. The Setup

As we all known, $U[x, y]$ is the uniform distribution between x and y , and $N(x, y)$ is the normal distribution which mean is x and variance is y . We assume that the time which vehicles share their resource is $T = 10$ seconds and $\Delta T = 1$ second, because of our limited space.

The Cortex-A57 processor is ARM's highest performing processor which is designed to further mobile and enterprise computing applications. The Cortex-A57 has the frequency from 700 MHz to 1900 MHz. The frequency set is $\{700, 800, 900, \dots, 1900\}$, as a result, there is only 13 frequency levels.

As is shown in Equation (3), the number of instructions is $M_{it} = v_i \cdot f_{it} + \theta_i$. And the CPU capacity in Equation (4) is $C_{it} = (v_i \cdot f_{it} + \theta_i) \times \Delta T - r_{it}$,

We estimate the CPU capacity parameters v_i, θ_i is based on the analysis provided in [22], and the value is 7.683 and 4558.52. After calculation, the maximum and minimum M values are respectively 819 and 10039. The required computing resource r_{it} for all tasks to execute workloads is uniformly drawn from $[820, 10000]$, and we assume that the task computing resource value r'_{it} which vehicle i can offload to others is $U[200, r_{it}]$, because some subtask must be execute locally. In our experiment the size of data varies

Algorithm 1: Genetic Algorithm

Input: input parameters $T, M, r', rloc$,
Output: \mathbf{X}, f, E

```

1  $t = 0$  ;
2 Initialize  $\mathbf{X}(0), f(0)$  ;
3 while  $t < T$  do
4   Encoder  $\mathbf{X}(t), f(t)$  ;
5   for  $i = 1$  to  $M$  do
6     Evaluate fitness  $P(ti) = -E_i$ ;
7   end
8   Eliminate fitness smaller individuals from  $P(t)$  ;
9   Replicate the fitness larger individuals from  $P(t)$  ;
10  Divide the array composed of optimization variables into two. ;
11  for  $i = 1$  to  $M/2$  do
12    Crossover operation to  $P(t)$  ;
13  end
14  for  $i = 1$  to  $M$  do
15    Mutation operation ;
16  end
17  Decoder  $\mathbf{X}(t), f(t)$  ;
18   $t \leftarrow t + 1$ ;
19 end
20 for  $i = 1$  to  $N$  do
21    $E \leftarrow E_i + E$ 
22 end
23  $\mathbf{X} = \mathbf{X}(t), f = f(t)$  ;
24 return  $\mathbf{X}, f, E$ 
```

from 1MB to 10 MB. And we set the value of bandwidth d is 27Mb/s.

We have selected seven kinds of vehicle number to participate in resource sharing, which are $\{10, 15, 20, 25, 30, 35, 40\}$. The population number was set to 40 in algorithm. While the number is set to 150 when the number of vehicles is 35 and 40, because it is too small and will not converge.

Our experiment are implemented in python and executed on an Intel Core i7 with 8 GB RAM.

B. Performance Metrics

The performance of our experiment is the percentage of energy saving, which is define as follows:

$$P = \min(100 \cdot (1 - \frac{\sum_{i=1}^N E_{it}^{blnc}}{E_{loc}})), t = 1, \dots, T \quad (20)$$

where E_{loc} is the energy consumed by local task execution.

In order to reflect the fairness of our framework, we define the fairness coefficient (FC) based on the standard deviation. The standard deviation is defined as follows.

$$\sigma = \sqrt{\frac{\sum_{i=1}^n (x_i - \mu)^2}{n}} \quad (21)$$

As with the standard deviation, the smaller the value the greater the fairness. FC is define as follows.

$$FC = \max \sqrt{\frac{1}{N} \sum_{i=1}^N (E_{it}^{blnc} - \bar{E})^2}, t = 1, \dots, T \quad (22)$$

C. Experimental Results

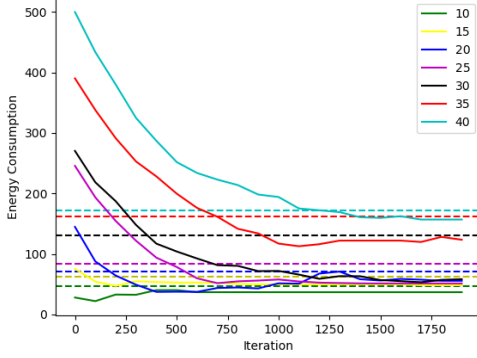


Fig. 2: Energy Consumption.

In Figure 2, the dotted line is the energy consumed by local execution. We can see that when the number of vehicles is 10, the algorithm can get good results by running a few times. While the number of vehicles is 20, it takes 350 iterations to get a better result. And when the number of vehicles is 30, the reduction rate of energy consumption is faster before 750 iterations, and will be significantly slower after 750 iterations. Almost no energy saving when the number of vehicles is 40.

We can conclude that with the increase of the number of vehicles, the number of iterations is also increasing. Therefore, when we run the algorithm on the edge server, we need to adjust the number of iterations reasonably as the number of vehicles changes.

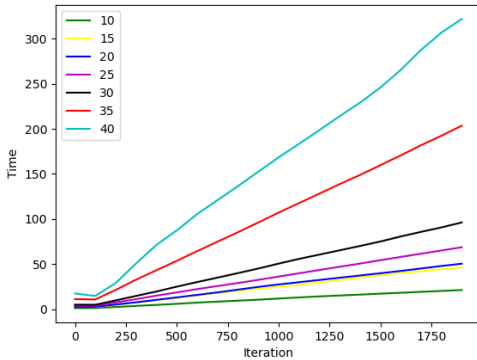
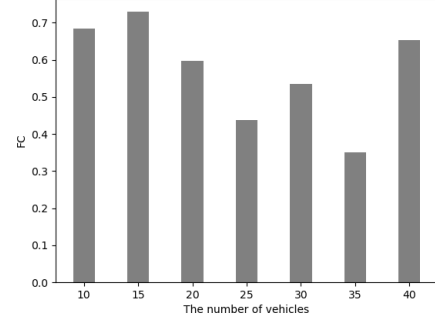
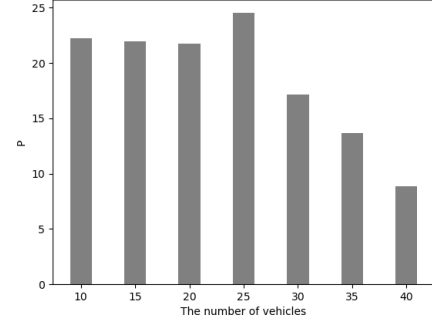


Fig. 3: Experimental running time

In Figure 3, after 100 iterations, the running time and the number of iterations show a linear relationship.



(a) The values of FC



(b) The values of P

Fig. 4: Performance metrics.

However, when the number of iterations is fixed, the relationship between time and vehicle is not linear.

As Figure 4(a) shows, we can see that with the increasing of vehicles, the balance of vehicle energy consumption increases first and then decreases. As Figure 4(b) shows, when the number of vehicles is between 10 and 25, our system can save over 20% the energy consumption, but when the number of vehicles is 30, only 17% can be saved, and when the number of vehicles is 40, almost no savings.

Considering the performance measurement, energy consumption and time, it is appropriate to choose 25 vehicles to participate in resource sharing.

V. CONCLUSIONS AND FUTURE

In this paper, we proposed task offloading based on genetic algorithm for electric smart vehicle. We evaluate the performance of the algorithm through simulation experiments. Experimental results show that our algorithm can achieve the goal of energy conservation.

In the future research, we plan to establish the relationship between the data size and the number of instructions and consider the deadline constraints. And we will also consider the authenticity of the unloading process.

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