

# The LLL Algorithm

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## 1 Motivation

The rows of the following matrix form a basis for a lattice  $L$  in  $\mathbb{R}^4$ :

$$X = \begin{bmatrix} -168 & 602 & 58 \\ 157 & -564 & -57 \\ 594 & -2134 & -219 \end{bmatrix}.$$

One can check that the rows of the following matrix also form a basis for the same lattice:

$$Y = \begin{bmatrix} -6 & 6 & -4 \\ 9 & 4 & 1 \\ -1 & 8 & 6 \end{bmatrix}.$$

Intuitively, the rows of  $X$  seem to be a “worse” basis for  $L$  than those of  $Y$ . Here we make precise the notion of a “nice” basis and introduce a polynomial time algorithm that transforms a “bad” basis into a “good” one.

## 2 Basis Reduction and the LLL Algorithm

A basis is “nice” if the constituent vectors are short and orthogonal to one another. The Gram-Schmidt process transforms a given basis into an orthogonal basis, but when working in a lattice  $L$ , this Gram-Schmidt basis need not live in  $L$ .

**Definition 2.1.** Let  $x_1, \dots, x_n$  be an ordered basis for a lattice  $L$  in  $\mathbb{R}^n$ , and let  $x_1^*, \dots, x_n^*$  be its Gram-Schmidt orthogonalization (GSO). Write  $X = MX^*$  where  $X$  (respectively  $X^*$ ) is the matrix with  $x_i$  (respectively  $x_i^*$ ) as row  $i$  and  $M = (\mu_{ij})$  is the matrix of GSO coefficients. Let  $\alpha$  be a real number with  $\frac{1}{4} < \alpha < 1$ , called the reduction parameter (usually taken to be  $\frac{3}{4}$ ). We say that the basis  $x_1, \dots, x_n$  is  **$\alpha$ -reduced** if it satisfies

1.  $|\mu_{ij}| \leq \frac{1}{2}$  for all  $1 \leq j < i \leq n$ ,
2.  $|x_i^* + \mu_{i,i-1}x_{i-1}^*|^2 \geq \alpha|x_{i-1}^*|^2$  for  $2 \leq i \leq n$ .

Condition (1) says that the  $i$ -th basis vector is “almost orthogonal” to the span of the previous  $i - 1$  vectors. The vector  $x_i^* + \mu_{i,i-1}x_{i-1}^*$  is the vector one obtains when swapping vectors  $x_i$  and  $x_{i-1}$  and then computing the  $(i - 1)$ -st vector of the GSO. Condition (2) then says that this new GSO vector, while potentially shorter than  $x_{i-1}^*$  isn’t “too much” shorter.

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**Algorithm 1** test

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**if**  $i \geq \textit{maxval}$  **then** $i \leftarrow 0$ **else****if**  $i + k \leq \textit{maxval}$  **then** $i \leftarrow i + k$ **end if****end if**

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