

"We use an idea due to
[]."

Bibliography

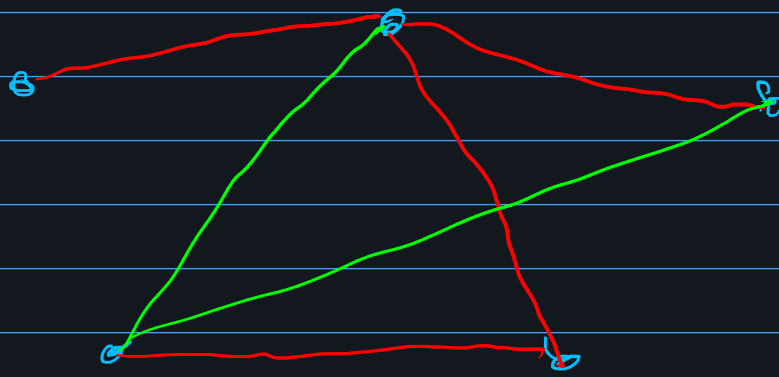
1.
2.
⋮

Bibtex

(a) Show that for every $c \leq 1/1000$ and $p = c/n$, with probability that tends to 1 as n tends to infinity, all the connected components of $G(n, p(n))$ are of size $O(\log n)$. (this is actually true for all $c < 1$, can you prove it?)

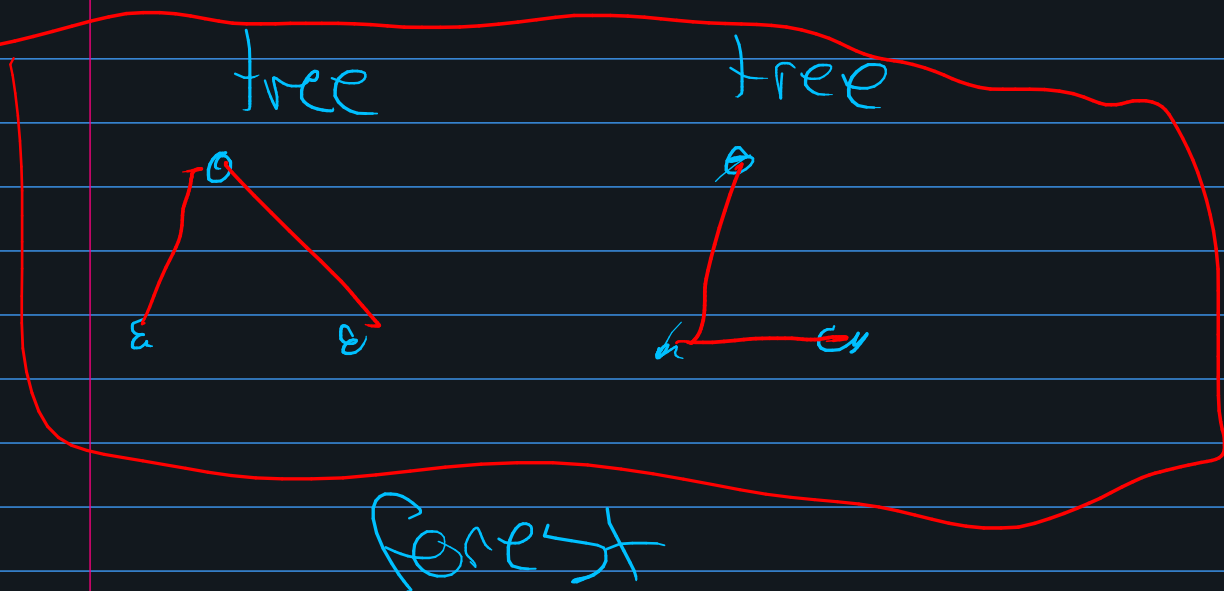
Idea: a set of k vertices forms
a connected component iff it contains
a spanning tree, & there's no edges
into the complement.

(connected)
Tree: graph w/ no cycles



~~all~~ edges
form
a spanning
tree.

In a tree, there is exactly one path between any two vertices



A tree on k vertices has
exactly $k-1$ edges

Let S be a set of k vertices.

$\Pr[S \text{ is a connected component}]$

$= \Pr[S \text{ contains a spanning tree \& has no edges outside of it}]$

\leq

$(\underbrace{\# \text{ of spanning trees on } k \text{ vertices}}_{T_k}) \cdot \Pr[\text{you have some specific spanning tree}]$

$$= p^{k-1} (1-p)^{\overbrace{k(n-k)}^{\substack{\text{\# of possible} \\ \text{edges out of} \\ S}}}$$

$\Pr[\Rightarrow \text{component of size } k]$

$$\leq T_k p^{k-1} (1-p)^{k(n-k)} \binom{n}{k}$$

set $K = C \log n$ & show that
this probability .

$P_r[\exists \text{ component size bigger than } K]$

 $\rightarrow 0$

$$f(n) = o(1) \iff \lim_{n \rightarrow \infty} f(n) = 0$$

$$f(n) = o(g(n)) \iff \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$$

$$\log_a x = \frac{\log_b x}{\log_b a}$$

(b) Show that for every $c \geq 1000$ and $p = c/n$, with probability that tends to 1 as n tends to infinity, there exists a connected component of a linear size. (this is true for all $c > 1$, can you prove something like that?)

 $\hookrightarrow \geq kn$ for some k
 $k < 1$

• Depth-first search:

Claim: If G is such that every two disjoint sets of size k contain an edge between them in G , then G contains a path of length $\geq n - 2k - 1$

Now show any two sets of size $\geq \frac{1}{2}(\frac{1}{2} - \epsilon)n$ in G_n c/n have an edge between them

$$\Rightarrow \exists \text{ path } \geq n - 2\left(\frac{1}{2} - \epsilon\right)n - 1 \\ = \Omega(n)$$

Project 2.

(b) Show that

$$\text{Var}[X] = \log \log n + O(1)$$

$$X_p = \begin{cases} 1 & \text{if } p \mid x \\ 0 & \text{else} \end{cases}$$

$$p \leq M = n^{1/100}$$

$$\text{Var}[X] = \sum_{p \leq m} \text{Var}[X_p]$$

$$+ \sum_{p \neq q} \text{Cov}(X_p, X_q)$$

$$X_p X_q = \begin{cases} 1 & \text{if } p \mid x \text{ \& } q \mid x \Leftrightarrow pq \mid x \\ 0 & \text{otherwise} \end{cases}$$

$$\text{Cov}(X_p, X_q) = E[X_p X_q] - E[X_p] E[X_q]$$

$$\boxed{X \text{ uniform in } [n]} \quad = \Pr[pq | x] - \Pr[p | x] \Pr[q | x]$$

$$E[X_p] = \Pr[p | x] = \frac{\# \text{ multiples of } p \text{ in } [n]}{n}$$

$$= \frac{\lfloor n/p \rfloor}{n}$$

$$\rightarrow = \frac{\lfloor n/pq \rfloor}{n} - \frac{\lfloor n/p \rfloor}{n} \cdot \frac{\lfloor n/q \rfloor}{n}$$

$$x - 1 < \lfloor x \rfloor \leq x \text{ for all } x$$

$$\leq \frac{n/pq}{n} - \frac{n/p - 1}{n} \cdot \frac{n/q - 1}{n}$$

$$= \frac{1}{pq} - \left(\frac{1}{p} - \frac{1}{n} \right) \left(\frac{1}{q} - \frac{1}{n} \right)$$

$$\cancel{\frac{1}{pq}} - \cancel{\frac{1}{pq}} + \frac{1}{nq} + \frac{1}{np} - \cancel{\frac{1}{n^2}}$$

$$\leq \frac{1}{n} \left(\frac{1}{p} + \frac{1}{q} \right)$$

$$\sum_{p \neq q} \text{Cov}(X_p, X_q) \leq \frac{1}{n} \sum_{\substack{p \neq q \\ p, q \leq M}} \left(\frac{1}{p} + \frac{1}{q} \right)$$

$$= \frac{2}{n} \sum_{p < q \leq M} \left(\frac{1}{p} + \frac{1}{q} \right)$$

$$\stackrel{n^{1/60}}{\leq} \frac{2M}{n} \sum_{p \leq M} \frac{1}{p} = o(1)$$

$\log \log n + O(1)$

$$\sum_{p \leq M} \text{Var}[X_p]$$

$$\text{Var}[X_p] = E[X_p^2] - E[X_p]^2$$

$$= \Pr[p|x] - \Pr[p|x]^2$$

$$= \Pr[p|x] \left(1 - \Pr[p|x] \right)$$

$$= \frac{L^n/p}{n} \left(1 - \frac{L^n/p}{n} \right)$$

$$X-1 < |X| \leq X$$

$$\leq \frac{1}{p} \left(1 - \frac{\frac{n}{p} - 1}{n} \right)$$

$$= \frac{1}{p} \left(1 - \frac{1}{p} + \frac{1}{n} \right)$$

$$M \leq n$$

$$= \frac{1}{p} \left(1 - \frac{1}{p} \right) + O\left(\frac{1}{n}\right)$$

$$\sum_{p \leq M} \text{Var}[X_p] \leq \sum_{p \leq M} \left[\frac{1}{p} \left(1 - \frac{1}{p} \right) + O\left(\frac{1}{n}\right) \right]$$

$$1 - \frac{1}{p} \leq 1 \quad \leq \left(\sum_{p \leq M} \frac{1}{p} \right) + O(1)$$

$$= \log \log n + o(1)$$

$$\Rightarrow \text{Var}[X] = \log \log n + O(1) + o(1)$$

$$= \log \log n + O(1).$$

$f(x) = O(g(x))$ means f doesn't grow more quickly than g

$f(x) = o(g(x))$ means f grows more slowly than g .

e.g. $\frac{1}{n^2} = o\left(\frac{1}{n}\right)$

Since $\lim_{n \rightarrow \infty} \frac{\frac{1}{n^2}}{\frac{1}{n}} = \lim_{n \rightarrow \infty} \frac{1}{n} = 0$

$$\Rightarrow \frac{1}{n^2} = o\left(\frac{1}{n}\right)$$

$$\frac{1}{n} + \frac{1}{n^2} = O\left(\frac{1}{n}\right)$$

Since $\frac{\frac{1}{n} + \frac{1}{n^2}}{\frac{1}{n}} = 1 + \frac{1}{n} \rightarrow 1 < \infty$

Big O notation is explained
in the lecture notes that
are on Canvas.