

The Code Libraries For XJTU ACM, RpBomb

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Chapter 1

High Precision

1.1 BigInteger

```
//written by chenkun

#include <cstdio>
#include <cstring>

inline __int64 max(__int64 a, __int64 b) { return a>b?a:b; }
struct Bignum {
    static const __int64 MAXUNIT=1000000000;
    //数组中每个元素能表示的最大数字+1
    static const int MAXLEN=9;
    //数组每个元素能表示的最大位数
    static const int MAXDIGITS=180;
    //bignum所能表示的最大位数
    static const int MAXROOMS=MAXDIGITS/MAXLEN; //数组大小
    __int64 v[MAXROOMS];
    int len;
    bool positive;
    Bignum(__int64 init=0) {
        memset(v, 0, sizeof v);
        v[0]=init;
        len=1;
        if(init>=0) positive=true; else positive=false;
    }
    Bignum(const char* str);
    Bignum operator+(const Bignum& rhs);
    Bignum& operator+=(const Bignum& rhs);
    Bignum operator*(const Bignum& rhs);
    bool operator<(const Bignum& rhs);
    bool operator==(const Bignum& rhs);
    bool operator>(const Bignum& rhs);
    void print();
};
```

```

Bignum::Bignum(const char* str) {
    __int64 t;
    char *s=(char*)((void*)str);
    if(s[0]=='-') {
        positive=false;
        s++;
    } else {
        positive=true;
    }
    if(s[0]=='+') s++;
    int l=strlen(s)/MAXLEN;
    if(strlen(s)%MAXLEN!=0) l++;
    len=l;
    t=0;
    for(int i=0;i<strlen(s)%MAXLEN;i++) {
        t=t*10+s[i]-'0';
    }
    v[len-1]=t;
    int start=len-1;
    if(strlen(s)%MAXLEN!=0) start--;
    for(int i=strlen(s)%MAXLEN;i<strlen(s);i+=MAXLEN) {
        t=0;
        for(int j=i;j<i+MAXLEN;j++) {
            t=t*10+s[j]-'0';
        }
        v[start--]=t;
    }
}

Bignum Bignum::operator +(const Bignum &rhs) {
    Bignum ret;
    int maxlen=max(len, rhs.len);
    __int64 r=0;
    for(int i=0;i<maxlen;i++) {
        ret.v[i]=v[i]+rhs.v[i]+r;
        if(ret.v[i]>=MAXUNIT) {
            r=ret.v[i]/MAXUNIT;
            ret.v[i]%=MAXUNIT;
        } else r=0;
    }
    ret.len=maxlen;
    if(r>0) ret.v[ret.len++]=r;
    return ret;
}

Bignum& Bignum::operator +=(const Bignum& rhs) {
    int maxlen=max(len, rhs.len);
    __int64 r=0;
    for(int i=0;i<maxlen;i++) {

```

```

        v[i]=v[i]+rhs.v[i]+r;
        if(v[i]>=MAXUNIT) {
            r=v[i]/MAXUNIT;
            v[i]=v[i]%MAXUNIT;
        } else r=0;
    }
    len=maxlen;
    if(r>0) v[len++]=r;
    return *this;
}

Bignum Bignum::operator *(const Bignum& rhs) {
    Bignum ret;
    __int64 r;
    for(int i=0;i<rhs.len;i++) {
        r=0;
        for(int j=0;j<len;j++) {
            ret.v[i+j]+=(rhs.v[i]*v[j]+r);
            if(ret.v[i+j]>=MAXUNIT) {
                r=ret.v[i+j]/MAXUNIT;
                ret.v[i+j]=ret.v[i+j]%MAXUNIT;
            } else r=0;
        }
        if(r>0) ret.v[len+i]+=r;
    }
    ret.len=rhs.len+len;
    while(ret.v[ret.len-1]==0) ret.len--;
    if(ret.len==0) ret.len=1;
    return ret;
}

bool Bignum::operator <(const Bignum& rhs) {
    if(len>rhs.len) return false;
    if(len<rhs.len) return true;
    for(int i=len-1;i>=0;i--) {
        if(v[i]>rhs.v[i]) return false;
        if(v[i]<rhs.v[i]) return true;
    }
    return false;
}

bool Bignum::operator >(const Bignum& rhs) {
    if(len>rhs.len) return true;
    if(len<rhs.len) return false;
    for(int i=len-1;i>=0;i--) {
        if(v[i]>rhs.v[i]) return true;
        if(v[i]<rhs.v[i]) return false;
    }
    return false;
}

```

```

}

bool Bignum::operator ==(const Bignum& rhs) {
    if(len!=rhs.len) return false;
    for(int i=len-1;i>=0;i--) {
        if(v[i]!=rhs.v[i]) return false;
    }
    return true;
}

void Bignum::print() {
    char buf[15];
    if(!positive) putchar('-');
    printf("%I64d",v[len-1]);
    for(int i=len-2;i>=0;i--) {
        sprintf(buf,"%I64d",v[i]);
        int len=strlen(buf);
        if(len<MAXLEN)
            for(int j=0;j<MAXLEN-len;j++) putchar('0');
        printf("%I64d",v[i]);
    }
    putchar('\n');
}

```

1.2 High Precision Div

Chapter 2

Data Structure

2.1 Disjoint Set

```
int p[MAXN+1], rank[MAXN+1];

inline void makeset(int a) {
    p[a]=a;
    rank[a]=0;
}

int findset(int a) {
    if(a!=p[a]) p[a]=findset(p[a]);
    return p[a];
}

void unionset(int a, int b) {
    if(a==b) return ;
    if(rank[a]>rank[b]) {
        p[b]=a;
    } else {
        p[a]=b;
        if(rank[a]==rank[b])
            rank[b]++;
    }
}
```

2.2 Range Minimum Query)(return the min element)

```
int a[MAXLEN];
int r[20][MAXLEN];
//rmq预处理数组r[j][i]表示a(i...i+2powj-1)中的最小值
int n; //a数组的长度
```

```

void rmq() {
    for (int i=0; i<n; i++)
        r[0][i]=a[i];
    for (int k=1; (1<<k)<=n; k++) {
        for (int i=0; i<n; i++) {
            r[k][i]=r[k-1][i];
            if (i+(1<<k-1)<n&& r[k-1][i+(1<<k-1)]<r[k][i])
                r[k][i]=r[k-1][i+(1<<k-1)];
        }
    }
}

//取得a数组在(x..y)之间的最小值
//int k=0;
//while((1<<k)<=(y-x)) k++;
//k--;
//return min(r[k][x+1], r[k][y-(1<<k)+1]);

```

2.3 Range Minimum Query(return the min element's pos)

```

const int MAX=1<<17;
//MAX is a pow of 2 and just bigger than n
int key[MAX];           //存放数据[0..n)
int n;                  //数据的个数
int pp[MAX*2];          //rmq数组

//i号则返回最小值小标,j则返回最大值小标
inline bool RMQcmp(int i, int j) {
    return key[i]>key[j];
}

//预处理, 建树
void RMQCreate() {
    int f=MAX, t=MAX+n-1, i;
    for (i=0; i<n; i++)
        pp[i+MAX]=i;
    for (; f<t; f/=2, t/=2) {
        for (i=f; i<t; i+=2) {
            if (!RMQcmp(pp[i], pp[i+1]))
                pp[i/2]=pp[i];
            else
                pp[i/2]=pp[i+1];
        }
        if (!(t&1)) pp[t/2]=pp[t];
    }
}

```

```

//返回key[ss,ee]中的最小（大）值下标
int RMQFind(int ss, int ee) {
    int k, f, t;
    ss+=MAX, ee+=MAX;
    k=!RMQcmp(pp[ss], pp[ee])? pp[ss]: pp[ee];
    for (f=ss, t=ee; f<t; f/=2, t/=2) {
        if (!(f&1)&&f+1<t&&RMQcmp(pp[f+1], k))
            k=pp[f+1];
        if ((t&1)&&t-1>f&&RMQcmp(pp[t-1], k))
            k=pp[t-1];
    }
    return k;
}

```

2.4 Leftist Tree

```

//左偏树written by chenkun
int key[MAXN+1]; //节点的键值
int left[MAXN+1], right[MAXN+1], p[MAXN+1], dist[MAXN+1];
//左偏树上每个节点的属性
//left, right表示左儿子和右儿子
//p表示父亲节点
//dist表示距离，使得左偏树能快速合并

```

```

inline void swap(int& a, int& b) {
    int t=a;
    a=b;
    b=t;
}

```

```

//init: 初始化
//读入key[], 初始化key[]=right[]=p[]=dist[]=0

```

```

//合并以a为根，以b为根的2个堆, 返回新堆的根节点

```

```

int merge(int a, int b) {
    if (a==0) return b;
    if (b==0) return a;
    //如果是小堆，则改为if(key[b]>key[a])
    if (key[b]>key[a]) swap(a, b);
    int c=merge(right[a], b);
    right[a]=c; p[c]=a;
    if (dist[right[a]]>dist[left[a]]) {
        int la=left[a];
        int ra=right[a];
        left[a]=ra;
        right[a]=la;
    }
    if (right[a]==0)
        dist[a]=0;
    else

```

```

        dist[a]=dist[right[a]]+1;
    return a;
}

//在以a为根的左偏树中插入节点x,返回根节点
int insert(int x,int a) {
    p[x]=left[x]=right[x]=dist[x]=0;
    return merge(a,x);
}

//返回节点a所在的左偏树的根节点,即返回最小值或最大值
int getroot(int a) {
    if(p[a]==0) return a;
    else return getroot(p[a]);
}

//删除节点a所在的左偏树的根节点, 返回新的堆的根节点
int deletemin(int a) {
    int r=getroot(a);
    p[left[r]]=p[right[r]]=0;
    return merge(left[r],right[r]);
}

/*
//最一般的删除操作, 可以在左偏树中删除任意一个已知的节点x
void delete(int x) {
    int q=p[x];
    int pp=merge(left[x],right[x]);
    p[x]=left[x]=right[x]=dist[x]=0;
    p[pp]=q;
    if(q!=0&&left[q]==x)
        left[q]=p;
    if(q!=0&&right[q]==x)
        right[q]=p;
    while(q) {
        if(dist[left[q]]<dist[right[q]])
            swap(left[q],right[q]);
        if(dist[right[q]]+1==dist[q]) return;
        dist[q]=dist[right[q]]+1;
        pp=q;
        q=p[pp];
    }
}
*/

```

2.5 Interval Tree(common class)

```

//线段树的通用构造方法written by chenkun
//MAXN为线段上点的最大个数
int len=1; //用来记录线段个数的变量

```

```
struct segtree {
    int l,r,sum;      //l,r表示左右端点的值
    segtree *lc,*rc;
    void con( int x, int y); //在(x,y)范围内建树
    void ins( int x,int y);  //在(x,y)范围内进行插入
    void cal( int , int );
}tr[MAXN*2],*rt = &tr[0];
```

```
void segtree::con(int x,int y) {
    l = x;
    r = y;
    sum=0;
    if( x == y ){ lc = rc = 0; return; }
    int mid = ( x + y ) >> 1;
    lc = &tr[ len++];
    rc = &tr[ len++];
    lc -> con( x , mid );
    rc -> con( mid+1 , y );
}
```

```
void segtree::ins(int x,int y) {}
```

```
void segtree::cal(int x,int y) {}
```

2.6 Heap-Dijkstra

```
//堆，另有一个用这个堆实现的Dijkstra
// *push 在某个位置插入
// *pop删除最小值
// *top取最小值
```

```
//堆的大小，也即途中顶点数
const int HEAP_SIZE=100;
template <class COST_TYPE>
class Heap {
public:
    //保存堆，元素值为堆中该位置代表的途中顶点编号
    int data[HEAP_SIZE];
    int size;           //以用堆的大小
    int index[HEAP_SIZE]; //保存图中的顶点对应于堆中的编号
    COST_TYPE cost[HEAP_SIZE];
    //cost[i]表示堆中data[i]位置的值

    void shift_up(int i) {
        int j;
        while(i>0) {
            j=(i-1)/2;
            if(cost[data[i]]<cost[data[j]]) {
                swap(index[data[i]],index[data[j]]);
                swap(data[i],data[j]);
            }
        }
    }
};
```

```

        i=j;
    } else {
        break;
    }
}

}

void shift_down(int i) {
    int j,k;
    while(2*i+1<size) {
        j=2*i+1;
        k=j+1;
        if((k<size)&&
            (cost[data[k]]<cost[data[j]])&&
            (cost[data[k]]<cost[data[i]])) {
            swap(index[data[k]],index[data[i]]);
            swap(data[k],data[i]);
            i=k;
        } else if(cost[data[j]]<cost[data[i]]) {
            swap(index[data[j]],index[data[i]]);
            swap(data[j],data[i]);
        } else { break; }
    }
}

void init() {
    size=0;
    memset(index,-1,sizeof index);
    memset(cost,-1,sizeof cost);
}

bool empty() { return size==0; }
int pop() {
    int res=data[0];
    data[0]=data[size-1];
    index[data[0]]=0;
    size--;
    shift_down(0);
    return res;
}

int top() { return data[0]; }
void push(int x,COST_TYPE c) {
    if(index[x]==-1) {
        cost[x]=c;
        data[size]=x;
        index[x]=size;
        size++;
        shift_up(index[x]);
    } else if(c<cost[x]) {
        cost[x]=c;
        shift_up(index[x]);
    }
}

```

```

                                shift_down(index[x]);
                            }
                        }
};

int Dijkstra(int G[20][20], int n, int s, int t) {
    Heap<int> heap;
    heap.init();
    heap.push(s, 0);
    while(!heap.empty()) {
        int u=heap.pop();
        if(u==t) {
            return heap.cost[t];
        }
        for(int i=0; i<n; i++) {
            if(G[u][i]>=0) {
                heap.push(i, heap.cost[u]+G[u][i]);
            }
        }
    }
    return -1;
}

```

2.7 Binary Indexed Tree

```

//树状数组
//c数组下标从1开始,也可以计算sum(0)
inline int Lowbit(int t) {
    return t & ( t ^ ( t - 1 ) );
}
//在pos位置加上num
void plus(int pos , int num) {
    while(pos <= n) {
        c[pos] += num;
        pos += Lowbit(pos);
    }
}
//返回(1..end)区间所有数的和
int sum(int end) {
    if(end==0) return 0;
    int sum = 0;
    while(end > 0) {
        sum += c[end];
        end -= Lowbit(end);
    }
    return sum;
}

```

2.8 Splay Tree-Dynamic Array

2.9 Interval Tree-Dynamic Ranking

```

//线段树求区间中排序后的第k个数
// *构造函数建树
// *find函数查找出区间中比给定数值小的元素个数
struct IntervalTree {
    int A,B,Mid,N;
    int *Data;
    IntervalTree *Left,*Right;

    IntervalTree(int x,int y) {
        A=x;B=y;Mid=(x+y)>>1;N=y-x+1;
        Left=Right=NULL;
        if(x==y) {
            Data=new int [N];
            Data[0]=List[x];
            return;
        }
        Left=new IntervalTree(A,Mid);
        Right=new IntervalTree(Mid+1,B);
        Data=new int [N];
        memcpy(Data,Left->Data,sizeof(int)*Left->N);
        memcpy(Data+Left->N,Right->Data,sizeof(int)*Right->N);
        sort(Data,Data+N);
    }
    int Rank(int x,int y,int number) {
        if(x>B||y<A) return 0;
        if(x<=A&&B<=y) {
            return lower_bound(Data,Data+N,number)-Data;
        }
        return Left->Rank(x,y,number)+Right->Rank(x,y,number);
    }
};

```

2.10 A Data Structure to Representations hexahedron

```

//前面，后面，上面，下面，左面，右面一次编号为3,1,2,0,4,5
// 前
// —
// —
//左i—i右
// —
// 后
int back[6]={2,3,0,1,5,4}; //记录每个面的对面的编号
int myleft[6][6];
//myleft[i][j]表示当下面，前面的编号分别为i,j是，左边的编号

void pre() {

```


2.10. A DATA STRUCTURE TO REPRESENTATIONS HEXAHEDRON 13

```

    myleft [0][3] = 4;
    myleft [0][1] = 5;
    myleft [0][5] = 3;
    myleft [0][4] = 1;
    myleft [3][0] = 5;
    myleft [3][2] = 4;
    myleft [3][4] = 0;
    myleft [3][5] = 2;
    myleft [2][3] = 5;
    myleft [2][1] = 4;
    myleft [2][4] = 3;
    myleft [2][5] = 1;
    myleft [1][2] = 5;
    myleft [1][0] = 4;
    myleft [1][4] = 2;
    myleft [1][5] = 0;
    myleft [4][3] = 2;
    myleft [4][1] = 0;
    myleft [4][2] = 1;
    myleft [4][0] = 3;
    myleft [5][3] = 0;
    myleft [5][1] = 2;
    myleft [5][2] = 3;
    myleft [5][0] = 1;
}

```


Chapter 3

Strings

3.1 KMP Match

主串: `s[]`, 模板串: `t[]`。

KMP函数返回`t`第一次在`s`中出现的位置(start from 0)。

若`s`不包含`t`, 则返回-1。

调用KMP函数之前要调用一次`getnext`函数

```
#include <stdio.h>
#include <string.h>
#define MAXs 10000
#define MAXt 1000

int next[MAXt];

void getnext(char t[])
{
    int i=0,j=-1,size=strlen(t);
    next[0]=-1;
    while(i<size)
    {
        if(j==-1 || t[i]==t[j]) {i++; j++; next[i]=j;}
        else j=next[j];
    }
}

int KMP(char s[], char t[], int pos)
{
    int i=pos, j=0;
    int s1=strlen(s), s2=strlen(t);
    while(i<s1 && j<s2)
    {
        if(j==-1 || s[i]==t[j]) {i++; j++;}
        else j=next[j];
    }
    if(j>=s2) return i-s2;
}
```

```

        else return -1;
    }

    int main()
    {
        char s[MAXs], t[MAXt];
        int pos;
        while( scanf("%s %s", s, t)==2)
        {
            getnext(t);
            pos=KMP(s, t, 0);
            if(pos>=0)
                printf("t_first_appears_in_s_at_%d\n", pos);
            else printf("t_is_not_included_in_s\n");
        }
        return 0;
    }

```

3.2 Trie tree

```

//字典树
//在危险节点上用邻接表保存引起该危险的单词序号
const int MAXNODE=100000;
int child[MAXNODE][26];
bool danger[MAXNODE];
int nodes;
//邻接表
int head[MAXNODE], q[MAXNODE], next[MAXNODE], tot;
void build_trie() {
    int i, j;
    nodes=1;
    static char words[10];
    for(i=0; i<m; i++) {
        scanf("%s\n", words);
        p=1;
        for(j=0; j<strlen(words); j++) {
            int d=words[j]-'a';
            if(d=='?'-'a') d=26;
            if(d=='*-'a') d=27;
            if(child[p][d]==0) {
                nodes++;
                child[p][d]=nodes;
            }
            p=child[p][d];
        }
        danger[p]=true;
        if(head[p]==0) {
            tot++;
            head[p]=tot;
            q[tot]=i;
        }
    }
}

```

```

        } else {
            int t=head[p];
            while(next[t]) t=next[t];
            tot++;
            next[t]=tot;
            q[tot]=i;
        }
    }
}

```

3.3 Trie Graph

//Trie图，进行多串匹配

```

const int maxn=3000; //节点的最大数目
const int maxchar=27; //字符集中元素的个数

```

```

int child[maxn+1][27];
bool danger[maxn+1];
int suffix[maxn+1],q[maxn+1];
int nodes,f,r,p;

```

//读入模式串，建立trie树

```

void build_trie() {
    int m;
    nodes=1;
    string words;
    cin>>m;
    for(int i=0;i<m;i++) {
        cin>>words;
        p=1;
        for(int j=0;j<words.size();j++) {
            int d=words[j]-'a';
            if(child[p][d]==0) {
                nodes++;
                child[p][d]=nodes;
            }
            p=child[p][d];
            if(danger[p]) break;
        }
        danger[p]=true;
    }
}

```

//根据以建好的trie树建立trie图

```

void build_graph() {
    f=r=0;
    for(int i=0;i<26;i++) {
        if(child[1][i]==0) {
            child[1][i]=1;
        } else {
            r++;
        }
    }
}

```

```

        q[r]=child[1][i];
        suffix[child[1][i]]=1;
    }
}
while(f<r) {
    f++;
    danger[q[f]]=danger[q[f]]||
        danger[suffix[q[f]]];
    if(!danger[q[f]]) {
        for(int i=0;i<26;i++) {
            if(child[q[f]][i]==0)
                child[q[f]][i]=child[suffix[q[f]]][i];
            else {
                r++;
                q[r]=child[q[f]][i];
                suffix[q[r]]=child[suffix[q[f]]][i];
            }
        }
    }
}
}
}

```

3.4 $O(n)$ Suffix Array

```

#define MAXLEN 200100
int n; //字符串长度
int SA[MAXLEN],rank[MAXLEN],h[MAXLEN],height[MAXLEN];
int r[20][MAXLEN]; //rmq预处理数组
int str[MAXLEN]; //字符串

inline bool leq(int a1,int a2,int b1,int b2) {
    return(a1 < b1 || a1 == b1 && a2 <= b2);
}
inline bool leq(int a1,int a2,int a3,int b1,int b2,int b3)
{
    return(a1 < b1 || a1 == b1 && leq(a2,a3, b2,b3));
}

static void radixPass(int* a,int* b,int* r,int n,int K) {
    int* c = new int[K + 1];
    for (int i = 0; i <= K; i++) c[i] = 0;
    for (int i = 0; i < n; i++) c[r[a[i]]]++;
    for (int i = 0, sum = 0; i <= K; i++) {
        int t = c[i]; c[i] = sum; sum += t;
    }
    for (int i = 0; i < n; i++) b[c[r[a[i]]]++] = a[i];
    delete [] c;
}

void suffixArray(int* s, int* SA, int n, int K) {

```

```

int n0=(n+2)/3, n1=(n+1)/3, n2=n/3, n02=n0+n2;
int* s12=new int[n02 + 3];
s12[n02]=s12[n02+1]=s12[n02+2]=0;
int* SA12=new int[n02 + 3];
SA12[n02]=SA12[n02+1]=SA12[n02+2]=0;
int* s0=new int[n0];
int* SA0=new int[n0];

for (int i=0,j=0;i<n+(n0-n1);i++) if(i%3!=0) s12[j++]=i;

radixPass(s12, SA12, s+2, n02, K);
radixPass(SA12, s12, s+1, n02, K);
radixPass(s12, SA12, s, n02, K);

int name = 0, c0 = -1, c1 = -1, c2 = -1;
for(int i=0;i<n02;i++) {
if (s[SA12[i]]!=c0 || s[SA12[i]+1]!=c1 || s[SA12[i]+2]!=c2){
    name++;c0=s[SA12[i]];c1=s[SA12[i]+1];c2=s[SA12[i]+2];
}
if (SA12[i]%3==1){s12[SA12[i]/3]=name;} // left half
else{s12[SA12[i]/3 + n0]=name;} // right half
}

if (name < n02) {
    suffixArray(s12, SA12, n02, name);
    for (int i=0;i<n02;i++) s12[SA12[i]] = i + 1;
} else
    for(int i=0;i<n02;i++) SA12[s12[i]-1] = i;

for (int i=0,j=0;i<n02;i++) if(SA12[i]<n0)
    s0[j++]=3*SA12[i];
radixPass(s0, SA0, s, n0, K);

for(int p=0,t=n0-n1,k=0;k<n;k++) {
#define GetI() (SA12[t]<n0?SA12[t]*3+1:(SA12[t]-n0)*3+2)
    int i = GetI();
    int j = SA0[p];
    if (SA12[t]<n0?
        leq(s[i],s12[SA12[t] + n0],s[j],s12[j/3]):
        leq(s[i],s[i+1],s12[SA12[t]-n0+1],
            s[j],s[j+1],s12[j/3+n0]))
    {
        SA[k]=i;t++;
        if (t == n02) {
            for(k++;p<n0;p++,k++) SA[k]=SA0[p];
        }
    } else {
        SA[k]=j;p++;
        if (p == n0) {
            for(k++;t<n02;t++,k++) SA[k]=GetI();
        }
    }
}

```

```

    }
  }
}
delete [] s12; delete [] SA12; delete [] SA0; delete [] s0;
}

void lcs() {
  for(int i=0; i<n; i++) rank[SA[i]]=i;
  for(int i=0; i<n; i++) {
    if(rank[i]==0) {
      h[i]=0;
      continue;
    }
    int j=rank[i]-1, k=rank[i], s;
    if(i==0 || h[i-1]<=1)
      s=0;
    else
      s=h[i-1]-1;
    for(; SA[k]+s<n && SA[j]+s<n; s++)
      if(str[SA[k]+s]!=str[SA[j]+s]) break;
    h[i]=s;
  }
  for(int i=0; i<n; i++)
    height[rank[i]]=h[i];
}

void rmq() {
  for(int i=0; i<n; i++) r[0][i]=height[i];
  for(int k=1; (1<k)<=n; k++) {
    for(int i=0; i<n; i++) {
      r[k][i]=r[k-1][i];
      if(i+(1<k-1)<n && r[k-1][i+(1<k-1)]<r[k][i])
        r[k][i]=r[k-1][i+(1<k-1)];
    }
  }
}

int asklcs(int x, int y) {
  if(x>y) swap(x, y);
  int k=0;
  while((1<k)<=(y-x)) k++;
  k--;
  return min(r[k][x+1], r[k][y-(1<k)+1]);
}

int main() {
  //读入str, 长度为n, 令SA[i]=str[i]
  str[n]=str[n+1]=str[n+2]=SA[n]=SA[n+1]=SA[n+2]=0;
  suffixArray(str, SA, n, b); //其中b为str数组中最大的数
}

```



```
    lcs ();  
    rmq ();  
    //asklcs(i,j) //求(i..j)的最长公共前缀  
}
```


Chapter 4

Graph Theory

4.1 Strongly Connected Components

```
#include <stdio.h>
#include <string.h>
#define G_size 100000
#define V_size 11000

typedef struct Graph
{
    int id;
    int next;
} Graph;

typedef struct Edge
{
    int s, e;
} Edge;

Edge E[ G_size ];
Graph GA[ G_size ], GT[ G_size ];
int N, M;           //点数, 边数
int G_end;
int order[ V_size ], id[ V_size ], vis[ V_size ];
int cnt;

void Insert( int s, int e ) //建立原图和逆图
{
    int p;
    p = s;
    while (GA[p].next)
        p = GA[p].next;
    GA[ G_end ].id = e;
    GA[p].next = G_end;
```

```

    p = e;
    while (GT[p].next)
        p = GT[p].next;
    GT[G_end].id = s;
    GT[p].next = G_end;

    G_end++;
}

void DFST(int x) //对逆图进行搜索
{
    int p, q;
    vis[x] = 1;
    p = GT[x].next;
    while (p)
    {
        q = GT[p].id;
        if (!vis[q])
            DFST(q);
        p = GT[p].next;
    }
    order[cnt++] = x;
}

void DFSA(int x) //对原图进行搜索
{
    int p, q;
    vis[x] = 1;
    id[x] = cnt;
    p = GA[x].next;
    while (p)
    {
        q = GA[p].id;
        if (!vis[q])
            DFSA(q);
        p = GA[p].next;
    }
}

void Solve() //主要过程
{
    int s, e;
    int i;

    memset(GA, 0, sizeof(GA));
    memset(GT, 0, sizeof(GT));
    memset(E, 0, sizeof(E));
    G_end = N + 1;

    for (i = 0; i < M; i++)

```

```

{
    scanf("%d_%d", &s, &e);
    E[i].s = s - 1;
    E[i].e = e - 1;
    Insert(s - 1, e - 1);
}

memset(vis, 0, sizeof(vis));
cnt = 0;
for (i = 0; i < N; i++)
{
    if (!vis[i])
    {
        DFST(i);
    }
}

memset(vis, 0, sizeof(vis));
cnt = 0;
for (i = N - 1; i >= 0; i--)
{
    if (!vis[order[i]])
    {
        DFSA(order[i]);
        cnt++;
    }
}
//id[v]表示v点所在强连通分量的编号
}

```

4.2 Lowest Common Ancestor(Online,RMQ-LCA)

```

/*rmq-lca方法求最小公共祖先
   writen by chenkun
*/
#include <stdio>
#include <cstring>

//树中节点最大数
const int MAXN=40000;
//使2 pow MAX大于MAXN即可
const int MAX=1<<17;

//下面为rmq数组
int key[MAX], seq[MAX];
int n; //特殊用途，中序遍历的计数
int pp[MAX*2];

int nn,m; //顶点数，边数nn-1==m
//下面为邻接表

```

```

int head [MAXN+1], next [2*MAXN+10],
    q [2*MAXN+10], cost [2*MAXN+10], tot;
int first [MAXN+1];
bool vis [MAXN+1];

inline void swap(int& a, int& b) {
    int t=a; a=b; b=t;
}
//加入边a-b
void insert (int a, int b) {
    int t;
    tot++;
    if (head[a]==0)
        head[a]=tot;
    else {
        t=head[a];
        head[a]=tot;
        next[tot]=t;
    }
    q[tot]=b;
    tot++;
    if (head[b]==0)
        head[b]=tot;
    else {
        t=head[b];
        head[b]=tot;
        next[tot]=t;
    }
    q[tot]=a;
}

void init () {
    int a, b, len;
    char c;
    scanf ("%d_%d", &n, &m);
    for (int i=0; i<m; i++) {
        scanf ("%d_%d", &a, &b);
        insert (a, b);
    }
}

//根据邻接表dfs构造树
void build_tree (int u, int level) {
    vis [u]=true;
    key [u]=level;
    seq [u]=u;
    first [u]=u;
    n++;
    int h=head [u];
    while (h) {

```

```

        if (!vis[q[h]]) {
            build_tree(q[h], level+1);
            key[n]=level;
            seq[n]=u;
            n++;
        }
        h=next[h];
    }

}

inline bool RMQcmp(int i, int j) {
    return key[i]>key[j];
}

void RMQCreate() {
    int f=MAX, t=MAX+n-1, i;
    for (i=0; i<n; i++)
        pp[i+MAX]=i;
    for (; f<t; f/=2, t/=2) {
        for (i=f; i<t; i+=2) {
            if (!RMQcmp(pp[i], pp[i+1]))
                pp[i/2]=pp[i];
            else
                pp[i/2]=pp[i+1];
        }
        if (!(t&1)) pp[t/2]=pp[t];
    }
}

int RMQFind(int ss, int ee) {
    int k, f, t;
    ss+=MAX, ee+=MAX;
    k=!RMQcmp(pp[ss], pp[ee])? pp[ss]: pp[ee];
    for (f=ss, t=ee; f<t; f/=2, t/=2) {
        if (!(f&1)&&f+1<t&&!RMQcmp(pp[f+1], k))
            k=pp[f+1];
        if ((t&1)&&t-1>f&&!RMQcmp(pp[t-1], k))
            k=pp[t-1];
    }
    return k;
}

int main() {
    int cas, s, e, fs, fe, lca;
    init();
    build_tree(1, 0);
    RMQCreate();
    // 查询s点和e点的lca
    scanf("%d %d", &s, &e);
    fs=first[s], fe=first[e];

```

```

        if (fs > fe) swap(fs, fe);
        lca = seq[RMQFind(fs, fe)];
        return 0;
    }

```

4.3 Bipartite Maximum Match(DFS)

```

const int maxm=200; //右图最大顶点数
const int maxn=200; //左图最大顶点数
bool g[maxn][maxm]; //邻接矩阵
int match[maxm];
bool us[maxm];

//dfs寻找增广路
bool go(int v) {
    int i;
    for (i=0; i<m; i++) {
        if (us[i]) continue;
        if (g[v][i]) {
            if (match[i]==-1) {
                us[i]=1;
                match[i]=v;
                return true;
            }
        }
    }
    for (i=0; i<m; i++) {
        if (us[i]) continue;
        if (g[v][i]) {
            us[i]=1;
            if (go(match[i])) {
                match[i]=v;
                return true;
            }
        }
    }
    return false;
}

//返回最大匹配数
int matcher() {
    int i, res=0;
    memset(match, 255, sizeof match);
    for (i=0; i<n; i++) {
        memset(us, 0, sizeof us);
        if (go(i)) ++res;
    }
    return res;
}

```


4.4 Bipartite Optimized Match(KM Algorithm)

```

/*
Hungary 算法求二部图的最优匹配
输入: C-二部图的利润矩阵,C[x][y]表示x和y匹配的利润
      若x和y之间没有边则C[x][y]=0
      nx-二部图的节点集合X中的元素数目
      ny-二部图的节点集合Y中的元素数目
输出: X,Y-最优匹配X,Y集合中节点所匹配的节点的id,-1表示
      该节点没有被匹配,若C[i][X[i]]=0,则最终结果可以删除
      这一匹配,因为有无这一匹配对最大利润没有影响
注意:输入必须保证C[x][y]≥0
*/
typedef int Graph[100][100];
typedef int Path[100];
void optmatch(Graph C,int nx,int ny,Path X,Path Y) {
    Path Lx,Ly,Q,prev;
    int i,j,k,s,head,tail;
    //要保证Y中的节点数目比X中的多
    if(ny<nx) ny=nx;
    for(i=0;i<nx;i++) {
        Lx[i]=Ly[i]=0;
        for(j=0;j<ny;j++)
            Lx[i]=max(Lx[i],C[i][j]);
    }
    memset(X,-1,sizeof Path);
    memset(Y,-1,sizeof Path);
    i=0;
    while(i<nx) {
        memset(prev,-1,sizeof Path);
        for(Q[0]=i,head=0,tail=1;head<tail&&X[i]<0;head++) {
            s=Q[head];
            for(j=0;j<ny&&X[i]<0;j++) {
                if(Lx[s]+Ly[j]>C[s][j]||prev[j]>=0) continue;
                Q[tail++]=Y[j];
                prev[j]=s;
                if(Y[j]<0) {
                    while(j>=0) {
                        s=prev[j];
                        Y[j]=s;
                        k=X[s];
                        X[s]=j;
                        j=k;
                    }
                }
            }
        }
        if(X[i]>=0) {i++;}
        else{

```

```

k=2147483647;
for (head=0;head<tail;head++) {
    s=Q[head];
    for (j=0;j<ny;j++) {
        if (prev[j]==-1) {
            k=min(k,Lx[s]+Ly[j]-C[s][j]);
        }
    }
}
for (j=0;j<tail;j++) Lx[Q[j]]-=k;
for (j=0;j<ny;j++)
    if (prev[j]>=0) Ly[j]+=k;
}
}
}

```

4.5 General Maximum Match

/*
 Method:Maximum Cardinality Matching Problem in General Graph
 By Edmonds Blossom-Contraction Algorithm
 Detail:
 Augmenting Path Theorem: If there is no augmenting path, the
 matching is maximal. An augmenting path is an alternating path
 which is started and ended with unmatched nodes;

Use BFS to find augmenting path.
 If there is an alternating cycle (must have an odd number of nodes),
 this cycle called "blossom" must be contracted as a new node

Notice to maintain the father pointers of the nodes in blossoms.

```

*/
#include <iostream>
#include <algorithm>
using namespace std;
#define SET0(x) memset(x,0,sizeof(x))
#define SET1(x) memset(x,0xff,sizeof(x))

const int MAXN=250;
typedef int Graph[MAXN][MAXN];
typedef int Path[MAXN];

int n,head,tail,start,finish,newbase;
Graph g;Path mat;Path Q,inQ,inP,inB,father,base;
/*Queue,inQueue,inPath,in Blossom,*/
void createGraph() {
    int u,v;
    SET0(g);
    scanf("%d",&n);
    while(2==scanf("%d%d",&u,&v))

```

```

        g[u][v]=g[v][u]=true;
    }
    inline void push(int u) {Q[tail++]=u;inQ[u]=true;}
    inline int pop() {return Q[head++];}

    int findCommonAncestor(int u,int v) {
        SET0(inP);
        while(1) {
            u=base[u];
            inP[u]=true;
            if(u==start) break;
            u=father[mat[u]];
        }
        while(1) {
            v=base[v];
            if(inP[v]) break;
            v=father[mat[v]];
        }
        return v;
    }
    void resetTrace(int u) {
        int v;
        while(base[u]!=newbase) {
            v=mat[u];
            inB[base[u]]=inB[base[v]]=true;
            u=father[v];
            if(base[u]!=newbase) father[u]=v;
        }
    }

    void blossomContract(int u,int v) {
        newbase=findCommonAncestor(u,v);
        SET0(inB);
        resetTrace(u);resetTrace(v);
        if(base[u]!=newbase) father[u]=v;
        if(base[v]!=newbase) father[v]=u;
        for(u=1;u<=n;u++)
            if(inB[base[u]]){
                base[u]=newbase;
                if(!inQ[u]) push(u);
            }
    }
    void findAugmentingPath() {
        int u,v;
        SET0(inQ);SET0(father);
        for(u=1;u<=n;u++) base[u]=u;
        tail=head=1;
        push(start);
        finish=0;
        while(head<tail) {

```

```

        u=pop();
        for (v=1;v<=n;v++)
            if (g[u][v]&&base[u]!=base[v]&&mat[u]!=v)
                if (v==start || mat[v]>0&&father[mat[v]]>0)
                    blossomContract(u,v);
                else if (father[v]==0) {
                    father[v]=u;
                    if (mat[v]>0) push(mat[v]);
                    else { finish=v; return; }
                }
            }
    }
    void augmentPath() {
        int u,v,w;
        u=finish;
        while (u>0) {
            v=father[u];
            w=mat[v];
            mat[v]=u;
            mat[u]=v;
            u=w;
        }
    }
    void edmonds() {
        int u;
        SET0(mat);
        for (u=1;u<=n;u++)
            if (mat[u]==0) {
                start=u;
                findAugmentingPath();
                if (finish>0) augmentPath();
            }
    }
    void printMatch() {
        int u;
        int count=0;
        for (u=1;u<=n;u++)
            if (mat[u]>0) count++;
        printf("%dn",count);
        for (u=1;u<=n;u++)
            if (u<mat[u])
                printf("%d_%d\n",u,mat[u]);
    }
    int main() {
        createGraph();
        edmonds();
        printMatch();
        return 0;
    }

```

4.6 Maximum Flow(Matrix)

```
//方法，将流网络读入矩阵co后(co[i][j]=-co[j][i]),
//while(bfs());即可得到最大流
//最大流的结果存放在flow[]中
//written by chenkun
const int maxn=110; //顶点最大数

struct tnode {
    int pre,now,mincost;
    tnode() {}
    tnode(int P,int N,int C):
        pre(P),now(N),mincost(C) {}
};

int n,m,s,t; //s,t分别为源点，汇点

int flow[maxn][maxn];
int co[maxn][maxn];
tnode queue[maxn];
bool vis[maxn];

//bfs求增广路并进行增流
bool bfs() {
    int f=0,r=0,h,nowf;
    tnode tt;
    memset(vis,false,sizeof vis);
    queue[r++]=tnode(-1,s,9999999);
    vis[s]=true;
    while(f<r) {
        tt=queue[f];
        if(tt.now==t) break;
        for(int i=0;i<=n+1;i++) {
            nowf=co[tt.now][i]-flow[tt.now][i];
            if(!vis[i]&&nowf>0) {
                queue[r++]=tnode(f,i,min(tt.mincost,nowf));
                vis[i]=true;
            }
        }
        f++;
    }
    if(tt.now!=t) return false;
    int minp=tt.mincost;
    nowf=tt.pre;
    while(nowf!=-1) {
        flow[queue[nowf].now][tt.now]+=minp;
        flow[tt.now][queue[nowf].now]
            -=flow[queue[nowf].now][tt.now];
        tt=queue[nowf];
    }
}
```

```

        nowf=tt.pre;
    }
    return true;
}

```

4.7 Maximum Flow(LinkList)

```

//链表最大流
//使用方法:while(Find_Path())Update()
//即可得到最大流的数值在Flow中
#include <cstdio>
#include <vector>
using namespace std;

const int MaxN=100,MaxM=1000,INF=(1<<30);
struct Node{
    int Data,C,F,Op;
};
vector<int> G[MaxN];
Node E[MaxM];

int S,T,M; //S,T为源点和汇点,M边数
int Flow;
int Q[MaxN],vf[MaxN],pre[MaxN];

void Add(int u,int v,int C) {
    E[M].Data=v;
    E[M].C=C;E[M].F=0;E[M].Op=M+1;
    G[u].push_back(M);M++;
    E[M].Data=u;
    E[M].C=0;E[M].F=0;E[M].Op=M-1;
    G[v].push_back(M);
    M++;
}

inline int fmin(int u,int v) {return (u<v)?u:v;}
bool Find_Path() {
    int F=0,R=0,u,v,j;
    memset(pre,-1,sizeof pre);
    Q[0]=S;pre[S]=-2;vf[S]=INF;
    while(F<=R) {
        u=Q[F++];
        for(j=0;j<G[u].size();j++) {
            v=G[u][j];
            if(E[v].C>E[v].F&&pre[E[v].Data]==-1) {
                Q[++R]=E[v].Data;
                pre[E[v].Data]=E[v].Op;
                vf[E[v].Data]=fmin(E[v].C-E[v].F,vf[u]);
                if(E[v].Data==T) return true;
            }
        }
    }
}

```

```

    }
    return false;
}
void Update() {
    int cnt=T;
    while(pre[cnt]!=-2) {
        E[pre[cnt]].F-=vf[T];
        E[E[pre[cnt]].Op].F+=vf[T];
        cnt=E[pre[cnt]].Data;
    }
    Flow+=vf[T];
}

```

4.8 Maximum Flow(LinkList)

```

const int maxn=100; //顶点最大数

struct tnode {
    int pre,now,mincost;
    tnode() {}
    tnode(int P,int N,int C):pre(P),now(N),mincost(C) {}
};

int n,m,s,t; //s,t分别为源点和汇点
int flow[maxn+1][maxn+1]; //流量
tnode queue[maxn];
bool vis[maxn+1];

//邻接表,c[]为容量
int head[maxn],next[maxn*maxn],q[maxn],c[maxn],tot;

//bfs求增广路并进行增流
bool bfs() {
    int f=0,r=0,h,nowf;
    tnode tt;
    memset(vis,false,sizeof vis);
    queue[r++]=tnode(-1,s,9999999);
    vis[s]=true;
    while(f<r) {
        tt=queue[f];
        h=head[tt.now];
        if(tt.now==t) break;
        while(h) {
            nowf=c[h]-flow[tt.now][q[h]];
            if(!vis[q[h]]&&nowf>0) {
                queue[r++]=tnode(f,q[h],min(tt.mincost,nowf));
                vis[q[h]]=true;
            }
            h=next[h];
        }
        f++;
    }
}

```

```

        f++;
    }
    if (tt.now!=t) return false;
    int minp=tt.mincost;
    nowf=tt.pre;
    while (nowf!=-1) {
        flow[queue[nowf].now][tt.now]+=minp;
        flow[tt.now][queue[nowf].now]=-flow[queue[nowf].now][tt.now];
        tt=queue[nowf];
        nowf=tt.pre;
    }
    return true;
}

```

4.9 Minimum Cost Maximum Flow

```

#include <iostream>
using namespace std;
#define MAXN 100
#define inf 1.0e10
int min_cost_max_flow(int n,int mat[][MAXN],int cost[][MAXN],
    int source,int sink,int flow[][MAXN],int& netcost) {
    int pre[MAXN],min[MAXN],d[MAXN],i,j,t,tag;
    if (source==sink) return inf;
    for (i=0;i<n;i++)
        for (j=0;j<n;j+=1) flow[i][j]=0;
    for (netcost=0;;) {
        for (i=0;i<n;i++) pre[i]=0,min[i]=inf;
        for (pre[source]=source+1,min[source]=0,d[source]=inf,tag=1;tag;)
            for (tag=t=0;t<n;t++)
                if (d[t]) for (i=0;i<n;i++)
                    if (j=mat[t][i]-flow[t][i]&&min[t]+cost[t][i]<min[i])
                        tag=1,min[i]=min[t]+cost[t][i],pre[i]=t+1,d[i]=d[t]<j?d[t]:j;
                    else if (j=flow[i][t]&&min[t]-cost[i][t]<min[i])
                        tag=1,min[i]=min[t]-cost[i][t],pre[i]=-t-1,d[i]=d[t]<j?d[t]:j;
        if (!pre[sink]) break;
        for (netcost+=min[sink]*d[i=sink];i!=source;)
            if (pre[i]>0)
                flow[pre[i]-1][i]+=d[sink],i=pre[i]-1;
            else
                flow[i][-pre[i]-1]-=d[sink],i=-pre[i]-1;
        for (j=i=0;j<n;j+=flow[source][i++]);
    }
    return j;
}

```

4.10 Minimum Arborescence

//最小树形图，邻接表written by chenkun


```

#define NOEDGE 999999
const unsigned int maxm=40000;           //最大边数
const int maxn=1000;                     //最大点数
bool vis[maxn];
int N,M;                                 //顶点数, 边数
int head[maxn], next[maxm], q[maxm], tot; //正向边的邻接表存储
int head2[maxn], next2[maxm], q2[maxm], tot2; //反向边的邻接表
存储
int G[maxn][maxn];                       //邻接矩阵
int res;                                 //最优值

void dfs(int v){
    vis[v]=true;
    int i=head[v];
    while(i) {
        if(!vis[q[i]])
            dfs(q[i]);
        i=next[i];
    }
}

//是否存在可行解
bool possible(){
    memset(vis,0,sizeof(vis));
    dfs(0);
    for(int i=1;i<N;++i)
        if(!vis[i])
            return false;
    return true;
}

//求最小树形图
int pre[maxn];
bool del[maxn];

//求最小树形图, 最后结果在res中 (可以输出方案)
void solve(){
    memset(del,0,sizeof(del));
    for(;;){
        int i;
        for(i=1;i<N;i++) {
            if(del[i]) continue;
            pre[i]=i;
            G[i][i]=NOEDGE;
            int j=head2[i];
            while(j) {
                if(!del[q2[j]]) {
                    if(G[pre[i]][i]>G[q2[j]][i])
                        pre[i]=q2[j];
                }
            }
        }
    }
}

```

```

        j=next2[j];
    }
}
for (i=1; i<N; i++) {
    if (del[i]) continue;
    int j=i;
    memset(vis, 0, sizeof vis);
    //寻找环
    while (!vis[j] && j!=0) {
        vis[j]=true;
        j=pre[j];
    }
    //如果没有环, 则退出
    if (j==0) continue;
    i=j;
    res+=G[pre[i]][i];
    //删除环
    for (j=pre[i]; j!=i; j=pre[j]) {
        res+=G[pre[j]][j];
        del[j]=true;
    }
    j=head2[i];
    //更新相应的边
    while (j) {
        if (!del[q2[j]]) {
            G[q2[j]][i] -= G[pre[i]][i];
        }
        j=next2[j];
    }
    for (j=pre[i]; j!=i; j=pre[j]) {
        int k=head[j];
        while (k) {
            if (q[k]!=0 && !del[q[k]] && k!=i) {
                if (G[i][q[k]] > G[j][q[k]]) {
                    //在邻接表中增加边
                    if (G[i][q[k]] == NOEDGE) {
                        tot++; tot2++;
                        int ttt=head[i];
                        head[i]=tot;
                        next[tot]=ttt;
                        q[tot]=q[k];
                        ttt=head2[q[k]];
                        head2[q[k]]=tot2;
                        next2[tot2]=ttt;
                        q2[tot2]=i;
                    }
                    G[i][q[k]] = G[j][q[k]];
                }
            }
            k=next[k];
        }
    }
}

```

```

    }
    k=head2[j];
    while(k) {
        if(!del[q2[k]]) {
            if(G[q2[k]][i]>G[q2[k]][j]-G[pre[j]][j]) {
                if(G[q2[k]][i]==NOEDGE) {
                    tot++; tot2++;
                    int ttt=head[q2[k]];
                    head[q2[k]]=tot;
                    next[tot]=ttt;
                    q[tot]=i;
                    ttt=head2[i];
                    head2[i]=tot2;
                    next2[tot2]=ttt;
                    q2[tot2]=q2[k];
                }
                G[q2[k]][i]=G[q2[k]][j]-G[pre[j]][j];
            }
        }
        k=next2[k];
    }
}
for(j=pre[i]; j!=i; j=pre[j]) {
    del[j]=true;
}
break;
}
if(i>=N) {
    for(int k=1; k<N; k++) {
        if(del[k]) continue;
        res+=G[pre[k]][k];
    }
    break;
}
}
}

void init() {
    tot=tot2=0;
    int a,b,c;
    scanf("%d_%d",&N,&M);
    memset(head,0,sizeof head);
    memset(next,0,sizeof next);
    memset(head2,0,sizeof head2);
    memset(next2,0,sizeof next2);
    for(int i=0; i<N; i++)
        for(int j=i; j<N; j++)
            G[i][j]=G[j][i]=NOEDGE;
    for(int i=0; i<M; i++) {
        scanf("%d_%d_%d",&a,&b,&c);

```

```

        tot++;
        tot2++;
        if (head[a]==0) {
            head[a]=tot;
            q[tot]=b;
        } else {
            int j=head[a];
            while(next[j]) j=next[j];
            next[j]=tot;
            q[tot]=b;
        }
        if (head2[b]==0) {
            head2[b]=tot2;
            q2[tot2]=a;
        } else {
            int j=head2[b];
            while(next2[j]) j=next2[j];
            next2[j]=tot2;
            q2[tot2]=a;
        }
        G[a][b]=c;
    }
}

```

4.11 K minimum span tree

```

//K限度生成树written by chenkun
const int maxn=50;
//G原图G2去掉限度节点后的图
int G[maxn][maxn],G2[maxn][maxn];
//label对G2进行编号,统一连同分量的节点标号相同;tot[i]表示标号为i的节点
//个数
int label[maxn],tot[maxn];
//u:dfs编号时的标记;G3:最小生成树
bool u[maxn],G3[maxn][maxn],b[maxn],u2[maxn];
//id表示当前强连同分量的标号值
int k,root,m,minid,maxid,id,maxdel,me1,me2;

//dfs求强连同分量并标号
void dfs(int v) {
    u[v]=true;
    tot[id]++;
    label[v]=id;
    for(int i=minid;i<=maxid;i++)
        if(G2[v][i]>0&&!u[i]) {
            dfs(i);
        }
}

//计算从root-i-v出发的圈中除root-i-v外的边的最大值maxdel

```

```

//me1,me2记录最大边的端点
bool circlemax(int v,int j) {
    u2[v]=true;
    for(int i=minid;i<=maxid;i++) {
        if(!G3[v][i]) continue;
        if(u2[i]) continue;
        if(i==root) {
            maxdel=j;
            return true;
        }
        if(j<G[v][i]) {
            me1=v;me2=i;
        }
        if(circlemax(i,max(j,G[v][i]))) return true;
    }
    return false;
}

//对k限度生成树进行一次换边操作
int addmstedge() {
    int bestdel=-1,bestret=999999,me1b,me2b;
    int addv;
    for(int i=minid;i<=maxid;i++) {
        if(G[root][i]==0) continue;
        if(G3[root][i]==true) continue;
        memset(u2,false,sizeof u2);
        circlemax(i,0);
        if(G[root][i]-maxdel<bestret) {
            bestret=G[root][i]-maxdel;
            addv=i;
            me1b=me1,me2b=me2;
        }
    }
    G3[root][addv]=G3[addv][root]=true;
    G3[me1b][me2b]=G3[me2b][me1b]=false;

    if(bestret==999999) return 0;else return bestret;
}

//求连同分量并标号
void make_cc() {
    for(int i=minid;i<=maxid;i++) {
        for(int j=minid;j<=maxid;j++) {
            G2[i][j]=G[i][j];
            if(i==root || j==root) G2[i][j]=0;
        }
    }
    memset(u,false,sizeof u);
    memset(tot,0,sizeof tot);
    id=0;u[root]=true;

```

```

        for (int i=minid; i<=maxid; i++) {
            if (i!=root)
                if (!u[i]) {
                    id++;
                    dfs(i);
                }
        }
    }

    // 计算标号为id的连同分量的最小生成树
    int calc_mst(int id) {
        int sp, ep, sum=0;
        for (int i=minid; i<=maxid; i++) {
            if (label[i]==id) {
                b[i]=true;
                break;
            }
        }
        for (int i=1; i<tot[id]; i++) {
            int min=999999;
            for (int j=minid; j<=maxid; j++) {
                if (label[j]!=id) continue;
                if (!b[j]) continue;
                for (int kk=minid; kk<=maxid; kk++) {
                    if (label[kk]!=id) continue;
                    if (b[kk]) continue;
                    if (G2[j][kk]==0) continue;
                    if (G2[j][kk]<min) {
                        min=G2[j][kk];
                        sp=j;
                        ep=kk;
                    }
                }
            }
            sum+=min;
            b[ep]=true;
            G3[sp][ep]=G3[ep][sp]=true;
        }
        return sum;
    }

    int solve() {
        int sum=0;
        make_cc();
        memset(b, false, sizeof b);
        memset(G3, false, sizeof G3);
        for (int i=1; i<=id; i++)
            sum+=calc_mst(i);
        for (int i=1; i<=id; i++) {
            int min=999999;

```

```

        int nowid;
        for (int j=minid; j<=maxid; j++) {
            if (label[j]!=i) continue;
            if (G[root][j]==0) continue;
            if (G[root][j]<min) {
                min=G[root][j];
                nowid=j;
            }
        }
        sum+=min;
        G3[root][nowid]=G3[nowid][root]=true;
    }
    int bestans=sum, ret=bestans;
    for (int nowx=id; nowx<k; nowx++) {
        bestans=bestans+addmstedge();
        ret=min(bestans, ret);
    }
    return ret;
}

```

4.12 K minimum span tree

// 最优比率生成树快速迭代 written by chenkun

```

double l[maxn+1][maxn+1], G[maxn+1][maxn+1]; // l 每条边的花费
int c[maxn+1][maxn+1]; // c 每条边的收益

```

```

double prim(double r) {
    double ret=0;
    memset(G, 0, sizeof G);
    for (int i=1; i<=n; i++) {
        for (int j=i+1; j<=n; j++) {
            G[i][j]=G[j][i]=c[i][j]-r*l[i][j];
        }
    }
    lc=0;
    rc=0;
    memset(u, false, sizeof u);
    u[1]=true;
    for (int i=2; i<=n; i++) {
        a[i]=G[1][i];
        cb[i]=c[1][i];
        lb[i]=l[1][i];
    }
    for (int i=2; i<=n; i++) {
        int id;
        double mind=999999999;
        for (int j=1; j<=n; j++) {
            if (!u[j] && a[j]<mind) {
                id=j;
                mind=a[j];
            }
        }
    }
}

```

```

        }
    }
    u[id]=true;
    ret+=mind;
    lc+=cb[id];
    rc+=lb[id];
    for (int j=1;j<=n;j++) {
        if (!u[j]&&G[id][j]<a[j]) {
            a[j]=G[id][j];
            cb[j]=c[id][j];
            lb[j]=l[id][j];
        }
    }
}
return ret;
}

double solve() {
    double r=30,rb=0;
    while (fabs(prim(r))>EPS&&fabs(r-rb)>EPS) {
        rb=r;
        r=lc/rc;           //迭代r=ax/bx
    }
    return r;
}

```


Chapter 5

Computational Geometry

5.1 Geometry 2D

```
//geom2d.cpp
//二维计算几何代码库

//By starfish(2003)
//Enhanced by phoenixinter(2004,2005)

#include<stdio.h>
#include<math.h>
#include<algorithm>
using namespace std;

#define INF 1e10 //无穷大
#define EPS 1e-8 //计算精度
#define PI acos(-1.0)
#define MAXN 100 //多边形的最多顶点数目

/*=====
                        数据结构定义
=====*/

struct Point
{
    double x,y;
    Point(double x0=0,double y0=0):x(x0),y(y0){}
};

struct Line
{
    Point p1,p2;
};
```

```

/*=====
          浮点数处理
=====*/
#define abs(x) ((x)>=0?x:-(x))
#define min(x,y) ((x)<(y)?(x):(y))
#define max(x,y) ((x)>(y)?(x):(y))
#define eq(x,y) (fabs((x)-(y))<EPS)
#define leq(x,y) ((x)<=(y)+EPS)
#define geq(x,y) ((x)+EPS>=(y))
#define zero(x) (((x)>0?(x):-(x))<EPS) //判定x是否为[-EPS,EPS]
#define _sign(x) ((x)>EPS?1:((x)<-EPS?2:0)) //判定x的符号, 返回1,2,0.
/*
    注意:
    如果是一个很小的负的浮点数,
    保留有效位数输出的时候会出现-0.000这样的形式,
    前面多了一个负号, 这就会导致错误!!!!!!
    因此在输出浮点数时, 一定要输出fix(x)!
*/
#define fix(x) (fabs(x)<EPS?0:x)

/*=====
          矢量基本操作
=====*/
bool operator<(Point p1,Point p2)
{
    if(!eq(p1.x,p2.x)) return p1.x<p2.x;
    else return p1.y<p2.y;
}

//计算p1-p2
Point operator-(Point p1,Point p2)
{
    return Point(p1.x-p2.x,p1.y-p2.y);
}

//计算叉积p1*p2
double operator*(Point p1,Point p2)
{
    return (p1.x*p2.y-p1.y*p2.x);
}

//计算叉积(p1-p0)*(p2-p0)
double times(Point p0,Point p1,Point p2)
{
    return (p1.x-p0.x)*(p2.y-p0.y)-(p1.y-p0.y)*(p2.x-p0.x);
}

//计算点积p1.p2
double operator&(Point p1,Point p2)

```

```

{
    return (p1.x*p2.x+p1.y*p2.y);
}

//求矢量u的模
inline double norm(Point u)
{
    return sqrt(u.x*u.x+u.y*u.y);
}

/*
旋转矢量
输入: p 被旋转的矢量
angle 旋转角度, 用弧度表示,
    0表示逆时针旋转,
    0表示顺时针旋转
输出: 旋转后得到的矢量
调用: 无
*/
Point Rotate(Point p, double angle)
{
    Point res;
    res.x=p.x*cos(angle)-p.y*sin(angle);
    res.y=p.x*sin(angle)+p.y*cos(angle);
    return res;
}

/*
矢量V以P为顶点逆时针旋转angle并放大scale倍(By phoenixinter)
*/
Point Rotate(Point v, Point p, double angle, double scale)
{
    Point ret=p;
    v.x-=p.x; v.y-=p.y;
    p.x=scale*cos(angle);
    p.y=scale*sin(angle);
    ret.x+=v.x*p.x-v.y*p.y;
    ret.y+=v.x*p.y+v.y*p.x;
    return ret;
}

/*
求一条直线的倾角, 直线用浮点数表示
(By phoenixinter)
输出: 直线(x1,y1)-(x2,y2)的倾角 (严格来说是向量的)
*/
double GetAngle(Point p1, Point p2)
{
    double dx=p2.x-p1.x, dy=p2.y-p1.y;
    double theta, pi=acos(-1.0);

```

```

        if (dx==0)
        {
            if(dy>0) return 90.0;
            else return 270.0;
        }
        else
        {
            theta=atan((double)dy/(double)dx)*180.0/pi; //弧度化为角度
            if(dx<0) theta+=180;
            else if(dy<0) theta+=360;
            return theta;
        }
    }

//求定比分点的坐标(By phoenixinter)
//输入: 两个点p1,p2,定比k
//输出: 定比分点的坐标
Point dingbipoint(Point p1,Point p2,double k)
{
    Point p;
    p.x=(p1.x+p2.x*k)/(1+k);
    p.y=(p1.y+p2.y*k)/(1+k);
    return p;
}

//二维向量的垂直运算(Perp dot)
//将向量(0,0)-iV逆时针旋转90度, 并保持长度不变。
Point perp(Point v)
{
    Point p;
    p.x=-v.y;p.y=v.x;
    return p;
}

//计算两点之间距离
inline double Dis(Point p1,Point p2)
{
    return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}

/*=====
                        点线关系
=====*/

/*
    判断三点是否共线
    输入: 三个点P1,P2,P3
    输出: true/false;
*/
bool dots_inline(Point p1,Point p2,Point p3)

```

```

{
    return zero(times(p3,p1,p2));
}

/*
P0是否在向量P1P2的左边
;0 说明在左边
=0 说明在P1P2这条直线上
;0 说明在P1P2右边
*/

double isLeft(Point P0,Point P1,Point P2)
{
    return times(P0,P1,P2);
}

/*
计算点p到直线L的距离
调用：无
*/
double Dis2Line(Point p,Line L)
{
    Point a,b;
    a.x=p.x-L.p1.x;a.y=p.y-L.p1.y;
    b.x=L.p2.x-L.p1.x;b.y=L.p2.y-L.p1.y;
    return fabs(a.x*b.y-a.y*b.x)/sqrt(b.x*b.x+b.y*b.y);
}

double ptoline(Point P,Line L)
{
    return fabs(times(L.p2,P,L.p1))/Dis(L.p1,L.p2);
}

/*
计算点p到直线L的最近点
调用：无
*/
Point Npt2Line(Point p,Line L)
{
    Point res;
    double a,b,t;
    a=L.p2.x-L.p1.x;b=L.p2.y-L.p1.y;
    t=((p.x-L.p1.x)*a+(p.y-L.p1.y)*b)/(a*a+b*b);
    res.x=L.p1.x+a*t;
    res.y=L.p1.y+b*t;
    return res;
}

/*
判断点p是否在直线L上

```

调用：无

```

*/
bool OnLine(Point p, Line L)
{
    double res;
    res=(L.p2.x-L.p1.x)*(p.y-L.p1.y)-(L.p2.y-L.p1.y)*(p.x-L.p1.x);
    return fabs(res)<EPS;
}

```

/*
计算点p与直线L的相对关系

输出：

- 0 - 点p在直线L上
- 1 - 点p在直线L左侧
- 2 - 点p在直线L右侧

调用：无

```

*/
int Relation(Point p, Line L)
{
    double res;
    res=(L.p2.x-L.p1.x)*(p.y-L.p1.y)-(L.p2.y-L.p1.y)*(p.x-L.p1.x);
    if (fabs(res)<EPS)
        return 0;
    else
        return (res>0)?1:2;
}

```

/*
求点p关于直线L的对称点
调用：无

```

*/
Point SymPoint(Point p, Line L)
{
    Point res;
    double a, b, t;
    a=L.p2.x-L.p1.x; b=L.p2.y-L.p1.y;
    t=((p.x-L.p1.x)*a+(p.y-L.p1.y)*b)/(a*a + b*b);
    res.x=2*L.p1.x+2*a*t-p.x;
    res.y=2*L.p1.y+2*b*t-p.y;
    return res;
}

```

/*
判断点p是否在线段L上
调用：无

```

*/
bool OnLineSeg(Point p, Line L)
{
    double r;
    r=(L.p2.x-L.p1.x)*(p.y-L.p1.y)-(L.p2.y-L.p1.y)*(p.x-L.p1.x);

```

```

        return (fabs(r)<EPS&&(p.x-L.p1.x)*(p.x-L.p2.x)<=EPS
                &&(p.y-L.p1.y)*(p.y-L.p2.y)<=EPS);
    }

    //判点是否在线段上,包括端点(By phoenixinter)
    int dot_online_in(Point p, Line L)
    {
        return zero(times(L.p2, p, L.p1)) && (L.p1.x-p.x)*(L.p2.x-p.x)<EPS
            && (L.p1.y-p.y)*(L.p2.y-p.y)<EPS;
    }

    //判点是否在线段上,不包括端点(By phoenixinter)
    int dot_online_ex(Point p, Line L)
    {
        return dot_online_in(p, L) && (!zero(p.x-L.p1.x) || !zero(p.y-L.p1.y))
            && (!zero(p.x-L.p2.x) || !zero(p.y-L.p2.y));
    }

    //判两点在线段同侧,点在线段上返回0(By phoenixinter)
    int same_side(Point p1, Point p2, Line L)
    {
        return times(L.p2, L.p1, p1)*times(L.p2, L.p1, p2)>EPS;
    }

    //判两点在线段异侧,点在线段上返回0(By phoenixinter)
    int opposite_side(Point p1, Point p2, Line L)
    {
        return times(L.p2, L.p1, p1)*times(L.p2, L.p1, p2)<-EPS;
    }

    /*
    计算点p到线段L的最近点
    调用: 无
    */
    Point Npt2LineSeg(Point p, Line L)
    {
        Point res;
        double a, b, t, d1, d2;
        a=L.p2.x-L.p1.x; b=L.p2.y-L.p1.y;
        t=((p.x-L.p1.x)*a+(p.y-L.p1.y)*b)/(a*a+b*b);
        if (geq(t, 0) && leq(t, 1))
        {
            res.x=L.p1.x+a*t;
            res.y=L.p1.y+b*t;
        }
        else
        {
            d1=(p.x-L.p1.x)*(p.x-L.p1.x)+(p.y-L.p1.y)*(p.y-L.p1.y);
            d2=(p.x-L.p2.x)*(p.x-L.p2.x)+(p.y-L.p2.y)*(p.y-L.p2.y);
            if (d1<d2)

```

```

        res=L.p1;
    else
        res=L.p2;
    }
    return res;
}

/*
计算点p到线段L的距离(By phoenixinter)
如果p到直线L的垂点在线段上, 那么返回点P到直线L的距离
否则返回P到线段两个端点较近一个的距离
*/
double ptoseg(Point p, Line L)
{
    Point t=p;
    t.x+=L.p1.y-L.p2.y;
    t.y+=L.p2.x-L.p1.x;
    if (times(p,L.p1,t)*times(p,L.p2,t)>EPS)
        return Dis(p,L.p1)<Dis(p,L.p2)?Dis(p,L.p1):Dis(p,L.p2);
    return fabs(times(L.p2,p,L.p1))/Dis(L.p1,L.p2);
}

//点到线段最近距离的平方
double SquaredDistance(Point Y, Line S)
{
    Point D=S.p2-S.p1;
    Point YmP0=Y-S.p1;
    double t=D&YmP0;
    if (t<=0)
        return YmP0&YmP0;
    double DdD=D&D;
    if (t>=DdD)
    {
        Point YmP1=Y-S.p2;
        return YmP1&YmP1;
    }
    return YmP0&YmP0-t*t/DdD;
}

/*=====
                        线线关系
=====*/

/*
直线方程: 两点式->标准式
输入: P1,P2两个点
输出: ax+by+c=0的直线方程的(a,b,c)系数
*/
void convert(Point p1, Point p2, double& a, double& b, double& c)
{
    a=p1.y-p2.y;

```



```

    b=p2.x-p1.x;
    c=p1.x*p2.y-p2.x*p1.y;
}

/*
判断两条直线L1,L2是否相交
调用：无
*/
bool LineIntersect(Line L1,Line L2)
{
    double res;
    res=(L1.p1.x-L1.p2.x)*(L2.p1.y-L2.p2.y)-
        (L1.p1.y-L1.p2.y)*(L2.p1.x-L2.p2.x);
    return fabs(res)>EPS;
}

//判断两条直线是否垂直(By phoenixinter)
bool perpendicular(Line u,Line v)
{
    return zero((u.p1.x-u.p2.x)*(v.p1.x-v.p2.x)+
        (u.p1.y-u.p2.y)*(v.p1.y-v.p2.y));
}

/*
判断两条线段L1,L2是否相交
调用：函数times
说明：快速排斥实验不仅仅是为了提高效率，更是必不可少的!
*/
bool LineSegIntersect(Line L1,Line L2)
{
    return (geq(max(L1.p1.x,L1.p2.x),min(L2.p1.x,L2.p2.x))
        &&geq(max(L2.p1.x,L2.p2.x),min(L1.p1.x,L1.p2.x))
        &&geq(max(L1.p1.y,L1.p2.y),min(L2.p1.y,L2.p2.y))
        &&geq(max(L2.p1.y,L2.p2.y),min(L1.p1.y,L1.p2.y))
        &&times(L1.p1,L2.p1,L1.p2)*times(L1.p1,L2.p2,L1.p2)<=EPS
        &&times(L2.p1,L1.p1,L2.p2)*times(L2.p1,L1.p2,L2.p2)<=EPS);
}

/*
SRbGa书上的线段相交(By phoenixinter)
通过测试:zju 1010,
*/
int dblcmp(double d)
{
    if(fabs(d)<EPS)
        return 0;
    return (d>0)?1:-1;
}

double det(double x1,double y1,double x2,double y2)

```

```

{
    return x1*y2-x2*y1;
}

double cross(Point a, Point b, Point c)
{
    return det(b.x-a.x, b.y-a.y, c.x-a.x, c.y-a.y);
}

double dotdet(double x1, double y1, double x2, double y2)
{
    return x1*x2+y1*y2;
}

double dot(Point a, Point b, Point c)
{
    return dotdet(b.x-a.x, b.y-a.y, c.x-a.x, c.y-a.y);
}

int betweencmp(Point a, Point b, Point c)
{
    return dblcmp(dot(a, b, c));
}

//0 no intersection, 1 proper intersection, 2 improper intersection
//p - point of intersection
//判断线段(a,b),(c,d)是否相交
int segcross(Point a, Point b, Point c, Point d, Point& p)
{
    double s1, s2, s3, s4;
    int d1, d2, d3, d4;
    d1=dblcmp(s1=cross(a, b, c));
    d2=dblcmp(s2=cross(a, b, d));
    d3=dblcmp(s3=cross(c, d, a));
    d4=dblcmp(s4=cross(c, d, b));
    if((d1^d2)==-2&&(d3^d4)==-2)
    {
        p.x=(c.x*s2-d.x*s1)/(s2-s1);
        p.y=(c.y*s2-d.y*s1)/(s2-s1);
        return 1;
    }
    if(d1==0&&betweencmp(c, a, b)<=0||
       d2==0&&betweencmp(d, a, b)<=0||
       d3==0&&betweencmp(a, c, d)<=0||
       d4==0&&betweencmp(b, c, d)<=0)
        return 2;
    return 0;
}

// intersect2D_2Segments(): the intersection of 2 finite 2D segments

```

```

//      Input:  two finite segments S1 and S2
//      Output: *I0 = intersect point (when it exists)
//              *I1 = endpoint of intersect segment [I0,I1] (when it exists)
//      Return: 0=disjoint (no intersect)
//              1=intersect in unique point I0
//              2=overlap in segment from I0 to I1
/*
int intersect2D_Segments( Segment S1, Segment S2, Point* I0, Point* I1 )
{
    Vector    u = S1.P1 - S1.P0;
    Vector    v = S2.P1 - S2.P0;
    Vector    w = S1.P0 - S2.P0;
    float     D = perp(u,v);

    // test if they are parallel (includes either being a point)
    if (fabs(D) < SMALLNUM) { // S1 and S2 are parallel
        if (perp(u,w) != 0 || perp(v,w) != 0) {
            return 0; // they are NOT collinear
        }
        // they are collinear or degenerate
        // check if they are degenerate points
        float du = dot(u,u);
        float dv = dot(v,v);
        if (du==0 && dv==0) { // both segments are points
            if (S1.P0 != S2.P0) // they are distinct points
                return 0;
            *I0 = S1.P0; // they are the same point
            return 1;
        }
        if (du==0) { // S1 is a single point
            if (inSegment(S1.P0, S2) == 0) // but is not in S2
                return 0;
            *I0 = S1.P0;
            return 1;
        }
        if (dv==0) { // S2 a single point
            if (inSegment(S2.P0, S1) == 0) // but is not in S1
                return 0;
            *I0 = S2.P0;
            return 1;
        }
        // they are collinear segments - get overlap (or not)
        float t0, t1; // endpoints of S1 in eqn for S2
        Vector w2 = S1.P1 - S2.P0;
        if (v.x != 0) {
            t0 = w.x / v.x;
            t1 = w2.x / v.x;
        }
        else {
            t0 = w.y / v.y;

```

```

        t1 = w2.y / v.y;
    }
    if (t0 > t1) {
        float t=t0; t0=t1; t1=t;    // must have t0 smaller than t1
        // swap if not
    }
    if (t0 > 1 || t1 < 0) {
        return 0;    // NO overlap
    }
    t0 = t0<0? 0 : t0;    // clip to min 0
    t1 = t1>1? 1 : t1;    // clip to max 1
    if (t0 == t1) {
        // intersect is a point
        *I0 = S2.P0 + t0 * v;
        return 1;
    }

    // they overlap in a valid subsegment
    *I0 = S2.P0 + t0 * v;
    *I1 = S2.P0 + t1 * v;
    return 2;
}

// the segments are skew and may intersect in a point
// get the intersect parameter for S1
float sI = perp(v,w) / D;
if (sI < 0 || sI > 1)    // no intersect with S1
    return 0;

// get the intersect parameter for S2
float tI = perp(u,w) / D;
if (tI < 0 || tI > 1)    // no intersect with S2
    return 0;

*I0 = S1.P0 + sI * u;    // compute S1 intersect point
return 1;
}
//=====
*/

/*
计算两条直线的交点
输入： L1, L2 – 两条直线
输出： P – 两直线的交点
返回值说明了两条直线的位置关系
1 – 共线
2 – 平行
0 – 相交
调用： 宏eq
*/
int CalCrossPoint(Line L1, Line L2, Point& P)
{

```

```

double a1, b1, c1, a2, b2, c2;
a1=L1.p2.y-L1.p1.y; b1=L1.p1.x-L1.p2.x;
c1=L1.p2.x*L1.p1.y-L1.p1.x*L1.p2.y;
a2=L2.p2.y-L2.p1.y; b2=L2.p1.x-L2.p2.x;
c2=L2.p2.x*L2.p1.y-L2.p1.x*L2.p2.y;
if (eq(a1*b2, b1*a2))
{
    if (eq(a1*c2, a2*c1)&&eq(b1*c2, b2*c1))
        return 1;    //共线
    else
        return 2;    //平行
}
else
{
    P.x=(b2*c1-b1*c2)/(a2*b1-a1*b2);
    P.y=(a1*c2-a2*c1)/(a2*b1-a1*b2);
    return 0;        //相交
}
}

/*
求两直线的夹角
输出: 0 PI之间的弧度
调用: 函数norm
*/
double Angle(Line L1, Line L2)
{
    Point u, v;
    u.x=L1.p1.x-L1.p2.x; u.y=L1.p1.y-L1.p2.y;
    v.x=L2.p1.x-L2.p2.x; v.y=L2.p1.y-L2.p2.y;
    return acos((u.x*v.x+u.y*v.y)/(norm(u)*norm(v)));
}

/*
计算线段L1到线段L2的最短距离
调用: 函数LineSegIntersect, Npt2LineSeg, Dis
*/
double MinDis(Line L1, Line L2)
{
    double d1, d2, d3, d4;
    if (LineSegIntersect(L1, L2))
        return 0;
    else
    {
        d1=Dis(Npt2LineSeg(L1.p1, L2), L1.p1);
        d2=Dis(Npt2LineSeg(L1.p2, L2), L1.p2);
        d3=Dis(Npt2LineSeg(L2.p1, L1), L2.p1);
        d4=Dis(Npt2LineSeg(L2.p2, L1), L2.p2);
        return min(min(d1, d2), min(d3, d4));
    }
}

```

```

}

// 计算两个移动物体的最近距离及其时间

struct Track
{
    Point P0;           // 物体的初始位置
    Point v;             // 物体移动的方向向量
};

#define dot1(u,v)      ((u).x*(v).x+(u).y*(v).y)

double cpa_time(Track Tr1,Track Tr2)
{
    Point dv=Tr1.v-Tr2.v;
    double dv2=dot1(dv,dv);
    if(dv2<EPS)
        return 0.0;
    Point w0=Tr1.P0-Tr2.P0;
    double cptime=-dot1(w0,dv)/dv2;
    return cptime;
}

double cpa_distance(Track Tr1,Track Tr2)
{
    double ctime=cpa_time(Tr1,Tr2);
    Point P1,P2;
    P1.x=Tr1.P0.x+(ctime*Tr1.v.x);
    P1.y=Tr1.P0.y+(ctime*Tr1.v.y);
    P2.x=Tr2.P0.x+(ctime*Tr2.v.x);
    P2.y=Tr2.P0.y+(ctime*Tr2.v.y);
    return Dis(P1,P2);
}

/*=====
           三角形
=====*/
struct Triangle
{
    Point P0,P1,P2;
};

/*
    计算三角形面积
    输入: a, b, c是三角形的三个顶点
    输出: 三角形面积, 面积正负按照右手旋规则确定
    调用: 无
*/
double Area(Point a,Point b,Point c)
{

```

```

    return ((b.x-a.x)*(c.y-a.y)-(b.y-a.y)*(c.x-a.x))/2.0;
}

/*
    计算三角形面积
    输入: a, b, c是三角形的三条边
    输出: 三角形面积
    调用: 无
*/
double Area(double a, double b, double c)
{
    double s=(a+b+c)/2.0;
    return sqrt(s*(s-a)*(s-b)*(s-c));
}

/*
    判断点p是否在三角形ABC内
    输出: 0 - 点p在三角形外
           1 - 点p在三角形内
           2 - 点p在三角形边界上
    调用: 函数OnLineSeg, Relation
*/
int InTriangle(Point p, Point A, Point B, Point C)
{
    Point center;
    Line side[3];
    int i, rel;
    center.x=(A.x+B.x+C.x)/3.0; center.y=(A.y+B.y+C.y)/3.0;
    side[0].p1=A; side[0].p2=B;
    side[1].p1=B; side[1].p2=C;
    side[2].p1=C; side[2].p2=A;
    rel=Relation(center, side[0]);
    for(i=0; i<3; i++)
    {
        if(OnLineSeg(p, side[i]))
            return 2; // 点p在三角形边界上
        else if(Relation(p, side[i])!=rel)
            return 0; // 点p在三角形外
    }
    return 1;
}

/*
    计算点到三角形的最近距离
*/
double SquaredDistance(Point Y, Triangle T)
{
    Point D0=T.P1-T.P0, D1=T.P2-T.P0, Delta=Y-T.P0;
    double a00=D0&D0, a01=D0&D1, a11=D1&D1;
    double b0=D0&Delta, b1=D1&Delta;

```

```

double n0=a11*b0-a01*b1;
double n1=a00*b1-a01*b0;
double d=a00*a11-a01*a01;
if (n0+n1<=d)
{
    if (n0>=0)
    {
        if (n1>=0)
        { //点在三角形内,region 0
            return 0;
        }
        else
        { //region 5
            double c=Delta&Delta;
            if (b0>0)
            {
                if (b0<a00)
                    return c-b0*b0/a00;
                else
                    return a00-2*b0+c;
            }
            else
                return c;
        }
    }
    else if (n1>=0)
    { //region 3
        double c=Delta&Delta;
        if (b1>0)
        {
            if (b1<a11)
                return c-b1*b1/a11;
            else
                return a11-2*b1+c;
        }
        else return c;
    }
    else
    { //region 4
        double c=Delta&Delta;
        if (b0<a00)
        {
            if (b0>0)
                return c-b0*b0/a00;
            else
            {
                if (b1<a11)
                {
                    if (b1>0)
                        return c-b1*b1/a11;

```



```

        else
            return c;
    }
    else
        return a11-2*b1+c;
    }
}
else
    return a00-2*b0+c;
}
}
else if (n0<0)
{ //region 2
    double c=Delta&Delta;
    if (b1>0)
    {
        if (b1<a11)
            return c-b1*b1/a11;
        else
        {
            double n=a11-a01+b0-b1, d=a00-2*a01+a11;
            if (n>0)
            {
                if (n<d)
                    return (a11-2*b1+c)-n*n/d;
                else
                    return a00-2*b0+c;
            }
            else
                return a11-2*b1+c;
        }
    }
}
else
    return c;
}
else if (n1<0)
{ //region 6
    double c=Delta&Delta;
    if (b0>0)
    {
        if (b0<a00)
            return c-b0*b0/a00;
        else
        {
            double n=a11-a01+b0-b1, d=a00-2*a01+a11;
            if (n>0)
            {
                if (n<d)
                    return (a11-2*b1+c)-n*n/d;
                else

```

```

        return a00-2*b0+c;
    }
    else
        return a11-2*b1+c;
    }
}
else return c;
}
else
{
    //region 1
    double c=Delta&Delta;
    double n=a11-a01+b0-b1,d=a00-2*a01+a11;
    if(n>0)
    {
        if(n<d)
            return (a11-2*b1+c)-n*n/d;
        else
            return a00-2*b0+c;
    }
    else
        return a11-2*b1+c;
}
}

/*
Another Version:
float SquaredDistance (Point Y, Triangle T)
{
    // T has vertices V0, V1, V2
    // t0 = n0/d0 = Dot(Y - V0, V1 - V0) / Dot(V1 - V0, V1 - V0)
    Point D0 = Y - V0, E0 = V1 - V0;
    float n0 = Dot(D0, E0);
    // t1 = n1/d1 = Dot(Y - V1, V2 - V1) / Dot(V2 - V1, V2 - V1)
    Point D1 = Y - V1, E1 = V2 - V1;
    float n1 = Dot(D1, E1);
    if (n0 <= 0 and n1 <= 0) // closest point is V1
        return Dot(D1, D1); // RETURN 0
    // t2 = n2/d2 = Dot(Y - V2, V0 - V2) / Dot(V0 - V2, V0 - V2);
    Point D2 = Y - V2, E2 = V0 - V2;
    float n2 = Dot(D2, E2);
    if (n1 <= 0 and n2 == 0) // closest point is V2
        return Dot(D2, D2); // RETURN 1
    if (n0 <= 0 and n2 <= 0) // closest point is V0
        return Dot(D0, D0); // RETURN 2
    // D0 = Y - V0 = V0 + c1 * (V1 - V0) + c2 * (V2 - V0) = V0 + c1
    // * E1 - c2 * E2 for
    // c0 + c1 + c2 = 1, c0 = m0 / d, c1 = m1 / d, c2 = m2 / d
    float e00 = Dot(E0, E0), e02 = Dot(E0, E2), e22 = Dot(E2, E2);
    float d = e02 * e02 - e00 * e22;
    float a = Dot(D0, E2);

```

```

float m1 = e02 * a - e22 * n0;
float m0, m2;
Point D;
if (d > 0) {
    if (m1 < 0) { // closest point is V2 + t2 * E2
        t2 = n2 / e22;
        D = Y - (V2 + t2 * E2);
        return Dot(D, D); // RETURN 3a
    }
    m2 = e00 * a - e02 * n0;
    if (m2 < 0) { // closest point is V0 + t0 * E0
        t0 = n0 / e00;
        D = Y - (V0 + t0 * E0);
        return Dot(D, D); // RETURN 4a
    }
    m0 = d - m1 - m2;
    if (m0 < 0) { // closest point is V1 + t1 * E1
        t1 = n1 / Dot(E1, E1);
        D = Y - (V1 + t1 * E1);
        return Dot(D, D); // RETURN 5a
    }
} else {
    if (m1 > 0) { // closest point is V2 + t2 * E2
        t2 = n2 / e22;
        D = Y - (V2 + t2 * E2);
        return Dot(D, D); // RETURN 3b
    }
    m2 = e00 * a - e02 * n0;
    if (m2 > 0) { // closest point is V0 + t0 * E0
        t0 = n0 / e00;
        D = Y - (V0 + t0 * E0);
        return Dot(D, D); // RETURN 4b
    }
    m0 = d - m1 - m2;
    if (m0 > 0) { // closest point is V1 + t1 * E1
        t1 = n1 / Dot(E1, E1);
        D = Y - (V1 + t1 * E1);
        return Dot(D, D); // RETURN 5b
    }
}
return 0; // Y is inside triangle, RETURN 6
}
*/

/*=====
                        多边形
=====*/

/*
    计算多边形面积
    输入: poly 多边形顶点数组

```

n 多边形顶点数目

输出: 多边形的面积, 正负按照右手旋规则确定

调用: 无

```

*/
double Area(Point poly[], int n)
{
    double res=0;
    if(n<3) return 0;
    for(int i=0;i<n;i++)
    {
        res+=poly[i].x*poly[(i+1)%n].y;
        res-=poly[i].y*poly[(i+1)%n].x;
    }
    return (res/2.0);
}

```

/*
 计算多边形的重心, 适用于任意简单多边形
 输入的多边形顶点数目必须大于0
 该算法可以一边读入多边形的顶点一边计算重心
 调用: 无

```

*/
Point Orthocenter(Point poly[], int n)
{
    Point p,p0,p1,p2,p3;
    double m,m0;
    p1=poly[0];p2=poly[1];p.x=p.y=m=0;
    for(int i=2;i<n;i++)
    {
        p3=poly[i];
        p0.x=(p1.x+p2.x+p3.x)/3.0;
        p0.y=(p1.y+p2.y+p3.y)/3.0;
        m0=p1.x*p2.y+p2.x*p3.y+p3.x*p1.y-p1.y*p2.x-p2.y*p3.x-p3.y*p1.x;
        if(fabs(m+m0)<EPS)
            m0+=EPS; // 为了防止除0溢出, 对m0做一点点修正
        p.x=(m*p.x+m0*p0.x)/(m+m0);
        p.y=(m*p.y+m0*p0.y)/(m+m0);
        m+=m0;
        p2=p3;
    }
    return p;
}

```

/*
 计算多边形的重心, 采用行列式方法计算
 误差会小一些
 By phoenixinter

```

*/
Point Orthocenter1(Point poly[], int n)
{

```

```

    Point p;
    p.x=p.y=0;
    for (int i=0;i<n;i++)
    {
        p.x+=(poly[i].x+poly[(i+1)%n].x)*(poly[i].x*poly[(i+1)%n].y-
            poly[(i+1)%n].x*poly[i].y);
        p.y+=(poly[i].y+poly[(i+1)%n].y)*(poly[i].x*poly[(i+1)%n].y-
            poly[(i+1)%n].x*poly[i].y);
    }
    p.x/=(6*Area(poly,n));
    p.y/=(6*Area(poly,n));
    return p;
}

/*
判断多边形是否是凸的
调用：函数Relation
*/
bool IsConvex(Point poly[],int n)
{
    int i,rel;
    Line side;
    if(n<3) return false;
    side.p1=poly[0];side.p2=poly[1];
    rel=Relation(poly[2],side);
    for(i=1;i<n;i++)
    {
        side.p1=poly[i];
        side.p2=poly[(i+1)%n];
        if(Relation(poly[(i+2)%n],side)!= rel) return false;
    }
    return true;
}

/*
判定凸多边形,顶点按顺时针或逆时针给出,允许相邻边共线
(By phoenixinter)
*/
int is_convex(int n,Point p[])
{
    int i,s[3]={1,1,1};
    for(i=0;i<n&&s[1]|s[2];i++)
        s[_sign(times(p[i],p[(i+1)%n],p[(i+2)%n]))]=0;
    return s[1]|s[2];
}

/*
判定凸多边形,顶点按顺时针或逆时针给出,不允许相邻边共线
(By phoenixinter)
*/

```

```

int is_convex_v2(int n, Point p[])
{
    int i, s[3]={1,1,1};
    for(i=0; i<n&& s[0]&& s[1]|s[2]; i++)
        s[_sign(times(p[i], p[(i+1)%n], p[(i+2)%n]))]=0;
    return s[0]&& s[1]|s[2];
}

/*
    判点是否在凸多边形内(By phoenixinter)
    顶点按顺时针或逆时针给出,在多边形边上返回0
*/
int inside_convex_v2(Point q, int n, Point p[])
{
    int i, s[3]={1,1,1};
    for(i=0; i<n&& s[0]&& s[1]|s[2]; i++)
        s[_sign(times(p[i], p[(i+1)%n], q))]=0;
    return s[0]&& s[1]|s[2];
}

/*
    判断点p是否在凸多边形poly内
    poly的顶点数目要大于等于3
    输出: 0 - 点p在poly外
           1 - 点p在poly内
           2 - 点p在poly边界上
    调用: 函数OnLineSeg, Relation
*/
int InConvex(Point p, Point poly[], int n)
{
    Point q;
    Line side;
    int i;
    q.x=q.y=0;
    for(i=0; i<n; i++)
    {
        q.x+=poly[i].x;
        q.y+=poly[i].y;
    }
    q.x=1.0*q.x/n; q.y=1.0*q.y/n;
    for(i=0; i<n; i++)
    {
        side.p1=poly[i];
        side.p2=poly[(i+1)%n];
        if(OnLineSeg(p, side))
            return 2; //点p在poly边界上
        else if(Relation(p, side)!=Relation(q, side))
            return 0; //点p在poly外
    }
    return 1; //点p在poly内
}

```

```

}

/*
判断点是否在任意简单多边形内(射线法)
输出: 0 - 点在poly外
      1 - 点在poly内
      2 - 点在poly边界上
调用: 函数eq, OnLineSeg, LineSegIntersect
测试通过: ZJU 1081
*/
int InPolygon(Point p, Point poly[], int n)
{
    int i, c;
    Line ray, side;
    c=0;
    ray.p1=p; ray.p2.y=p.y; ray.p2.x=-INF;
    for (i=0; i<n; i++)
    {
        side.p1=poly[i];
        side.p2=poly[(i+1)%n];
        if (OnLineSeg(p, side))
            return 2; //点在poly边界上
        if (eq(side.p1.y, side.p2.y)) //如果side平行x轴则不作考
            continue;
        if (OnLineSeg(side.p1, ray))
        {
            if (side.p1.y>side.p2.y) c++;
        }
        else if (OnLineSeg(side.p2, ray))
        {
            if (side.p2.y>side.p1.y) c++;
        }
        else if (LineSegIntersect(ray, side))
            c++;
    }
    return ((c%2==1)?1:0); //1:点在poly内;0:点在poly外
}

/*
判断线段是否在任意简单多边形内
调用: 宏eq, 函数InPolygon, OnLineSeg, LineSegIntersect
Point的j 操作符
*/
bool InPolygon(Line L, Point poly[], int n)
{
    bool res;
    int i, m;
    Point p, pts[MAXN];
    Line side;

```

```

    if (!InPolygon(L.p1, poly, n) || !InPolygon(L.p2, poly, n))
        return false;
    m=0;
    for (i=0; i<n; i++)
    {
        side.p1=poly[i]; side.p2=poly[(i+1)%n];
        if (OnLineSeg(L.p1, side))
            pts[m++] = L.p1;
        else if (OnLineSeg(L.p2, side))
            pts[m++] = L.p2;
        else if (OnLineSeg(side.p1, L))
            pts[m++] = side.p1;
        else if (OnLineSeg(side.p2, L))
            pts[m++] = side.p2;
        else if (LineSegIntersect(side, L))
            return false;
    }
    sort(&pts[0], &pts[m]); //对交点进行排序
    for (i=1; i<m; i++)
        if (!eq(pts[i-1].x, pts[i].x) || !eq(pts[i-1].y, pts[i].y))
        {
            p.x=(pts[i-1].x+pts[i].x)/2.0;
            p.y=(pts[i-1].y+pts[i].y)/2.0;
            if (!InPolygon(p, poly, n))
                return false;
        }
    return true;
}

/*
判断线段是否在任意简单多边形外
调用: 宏eq, 函数InPolygon, OnLineSeg, LineSegIntersect
Point的; 操作符
*/
bool OutPolygon(Line L, Point poly[], int n)
{
    bool res;
    int i, m;
    Point p, pts[MAXN];
    Line side;
    if (InPolygon(L.p1, poly, n)==1 || InPolygon(L.p2, poly, n)==1)
        return false;
    m=0;
    for (i=0; i<n; i++)
    {
        side.p1=poly[i]; side.p2=poly[(i+1)%n];
        if (OnLineSeg(L.p1, side) && OnLineSeg(L.p2, side)) return true;
        if (LineSegIntersect(side, L)) return false;
    }
    for (i=0; i<n; i++)

```



```

{
    side.p1=poly[i]; side.p2=poly[(i+1)%n];
    if(OnLineSeg(L.p1, side))
        pts[m++]=L.p1;
    else if(OnLineSeg(L.p2, side))
        pts[m++]=L.p2;
    else if(OnLineSeg(side.p1, L))
        pts[m++]=side.p1;
    else if(OnLineSeg(side.p2, L))
        pts[m++]=side.p2;
}
sort(&pts[0], &pts[m]);
for(i=1; i<m; i++)
{
    if(!eq(pts[i-1].x, pts[i].x) || !eq(pts[i-1].y, pts[i].y))
    {
        p.x=(pts[i-1].x+pts[i].x)/2.0;
        p.y=(pts[i-1].y+pts[i].y)/2.0;
        if(InPolygon(p, poly, n)==1)
            return false;
    }
}
return true;
}

```

/*

用有向直线line切割凸多边形

复杂度: $O(n)$

输入: line 用来切割凸多边形的有向直线

poly 凸多边形顶点数组, 顶点必须按照逆时针排列

n 凸多边形的顶点数目

输出: result[1] 切割后line的左侧部分

result[2] 切割后line的右侧部分

result[0] 没有用到, 只是作为辅助存储空间

m[0..2] m[i]是result[i]中顶点的数目

返回值切口的长度, 如果返回值为0, 说明未做切割

当未作切割时, 如果多边形在该直线的左侧,

则result[1]等于该多边形, 否则result[2]等于该多边形

注意:

1. 被切割的多边形一定要是凸多边形, 顶点按照逆时针排列
2. 可利用这个函数来求多边形的核, 初始的核设为一个很大的矩形, 然后依次用多边形的每条边去割

调用: 函数Relation, CalCrossPoint, Dis

*/

```

double CutConvex(Line line, Point poly[], int n,
                  Point result[3][MAXN], int m[3])
{
    Point pts[3], p;
    Line side;
    int i, cur, pre, npt;

```

```

m[0]=m[1]=m[2]=npt=0;
if (n==0) return 0;
pre=cur=Relation(poly[0], line);
for (i=0; i<n; i++)
{
    cur=Relation(poly[(i+1)%n], line);
    if (cur==pre)
        result[cur][m[cur]++]=poly[(i+1)%n];
    else
    {
        side.p1=poly[i]; side.p2=poly[(i+1)%n];
        CalCrossPoint(side, line, p);
        pts[npt++]=p;
        result[pre][m[pre]++]=p;
        result[cur][m[cur]++]=p;
        result[cur][m[cur]++]=poly[(i+1)%n];
        pre=cur;
    }
}
sort(&pts[0], &pts[npt]);
if (npt<2)
    return 0;
else
    return Dis(pts[0], pts[npt-1]);
}

/*
  寻找一个点到多边形的切点
  输入: 点P(多边形外部的一个点)
        n(多边形的顶点数)
        V(多边形的顶点坐标)
  输出: int rtan, ltan
        分别表示最右边和最左边的切点的坐标的index
*/

// tests for polygon vertex ordering relative to a fixed point P
#define above(P, Vi, Vj) (isLeft(P, Vi, Vj) > 0) // true if Vi is above Vj
#define below(P, Vi, Vj) (isLeft(P, Vi, Vj) < 0) // true if Vi is below Vj

//复杂度O(n)
void tangent_PointPoly(Point P, int n, Point V[], int& rtan, int& ltan)
{
    double eprev, enext;
    int i;
    rtan=ltan=0;
    eprev=isLeft(V[0], V[1], P);
    for (i=1; i<n; i++)
    {
        enext=isLeft(V[i], V[(i+1)%n], P);
        if (eprev<=0&&enext>0)

```

```

        {
            if (!below(P, V[i], V[rtan]))
                rtan=i;
        }
        else if (eprev>0&&enext<=0)
        {
            if (!above(P, V[i], V[ltan]))
                ltan=i;
        }
        eprev=enext;
    }
    return;
}

int Rtangent_PointPolyC(Point P, int n, Point V[]);
int Ltangent_PointPolyC(Point P, int n, Point V[]);

void tangent_PointPolyC(Point P, int n, Point V[], int& rtan, int& ltan)
{
    rtan=Rtangent_PointPolyC(P, n, V);
    ltan=Ltangent_PointPolyC(P, n, V);
}

// 返回：最右边切点坐标的index
int Rtangent_PointPolyC(Point P, int n, Point V[])
{
    int a, b, c, upA, dnC;
    if (below(P, V[1], V[0])&&!above(P, V[n-1], V[0]))
        return 0;
    for (a=0, b=n;;)
    {
        c=(a+b)/2;
        dnC=below(P, V[c+1], V[c]);
        if (dnC&&!above(P, V[c-1], V[c]))
            return c;
        upA=above(P, V[a+1], V[a]);
        if (upA)
        {
            if (dnC)
                b=c;
            else
            {
                if (above(P, V[a], V[c]))
                    b=c;
                else
                    a=c;
            }
        }
    }
    else
    {

```



```

    }
}

// RLtangent_PolyPolyC(): get the RL tangent between two convex polygons
//   Input:  m = number of vertices in polygon 1
//           V = array of vertices for convex polygon 1 with V[m]=V[0]
//           n = number of vertices in polygon 2
//           W = array of vertices for convex polygon 2 with W[n]=W[0]
//   Output: *t1 = index of tangent point V[t1] for polygon 1
//           *t2 = index of tangent point W[t2] for polygon 2
void RLtangent_PolyPolyC( int m, Point* V, int n, Point* W, int* t1, int* t2 )
{
    int ix1, ix2;          // search indices for polygons 1 and 2

    // first get the initial vertex on each polygon
    ix1 = Rtangent_PointPolyC(W[0], m, V); // right tangent from W[0] to V
    ix2 = Ltangent_PointPolyC(V[ix1], n, W); // left tangent from V[ix1] to W

    // ping-pong linear search until it stabilizes
    int done = false;      // flag when done
    while (done==false) {
        done = true;       // assume done until...
        while (isLeft(W[ix2], V[ix1], V[ix1+1]) <= 0){
            ++ix1;          // get Rtangent from W[ix2] to V
        }
        while (isLeft(V[ix1], W[ix2], W[ix2-1]) >= 0){
            --ix2;          // get Ltangent from V[ix1] to W
            done = false;   // not done if had to adjust this
        }
    }
    *t1 = ix1;
    *t2 = ix2;
    return;
}

/*=====
           矩形
=====*/
//表示矩形，左下角坐标是(llx,lly),右上角坐标是(urx,ury)
struct Rect
{
    int llx, lly, urx, ury;
};

/*
判断两个矩形是否相交
相邻不算相交
*/
bool intersect(Rect r1, Rect r2)
{

```

```

        return (max(r1.llx, r2.llx) < min(r1.urx, r2.urx) &&
                max(r1.lly, r2.lly) < min(r1.ury, r2.ury));
    }

    /*
    判断点p是否在矩形内
    调用：函数geq
    通过测试：PKU 1468
    */
    bool inrect(Point p, Rect rect)
    {
        return (geq(p.x, rect.llx) && leq(p.x, rect.urx)
                && geq(p.y, rect.lly) && leq(p.y, rect.ury));
    }

    /*
    用矩形b切割矩形a
    输出：out[4]中保存切割a后得到的新矩形
    a 本身将会被改变
    返回值：如果矩形a,b不相交，返回0
    如果矩形b完全覆盖矩形a,返回-1
    否则返回切割后得到的新矩形的数目
    */
    int CutRect(Rect& a, Rect b, Rect out[4])
    {
        if (b.urx < a.llx || b.llx >= a.urx) return 0;
        if (b.ury < a.lly || b.lly >= a.ury) return 0;
        if (b.llx <= a.llx && b.urx >= a.urx && b.lly <= a.lly && b.ury >= a.ury)
            return -1;
        int n = 0;
        if (b.llx > a.llx)
        {
            out[n] = a; out[n].urx = b.llx;
            n++; a.llx = b.llx;
        }
        if (b.urx < a.urx)
        {
            out[n] = a; out[n].llx = b.urx;
            n++; a.urx = b.urx;
        }
        if (b.lly > a.lly)
        {
            out[n] = a; out[n].ury = b.lly;
            n++; a.lly = b.lly;
        }
        if (b.ury < a.ury)
        {
            out[n] = a; out[n].lly = b.ury;
            n++; a.ury = b.ury;
        }
    }

```

```

    return n;
}

//用长宽表示矩形,w,h分别表示宽度和高度
struct Rect2
{
    double w,h;
};

/*
判断矩形r2是否可以放置在矩形r1内
r1和r2可以任意地旋转
调用：无
*/
bool IsContain(Rect2 r1,Rect2 r2)
{
    double cross,alpha;
    if(r1.h>r1.w) swap(r1.h,r1.w);
    if(r2.h>r2.w) swap(r2.h,r2.w);
    if(leq(r2.h,r1.h)&&leq(r2.w,r1.w))
        return true;
    cross=sqrt(r2.w*r2.w+r2.h*r2.h);
    //注意，现在r1.h肯定大于cross
    alpha=asin(r1.h/cross)-asin(r2.h/cross);
    if(alpha>0&&2.0*alpha<PI)
    {
        if(r2.w*cos(alpha)+r2.h*sin(alpha)<=r1.w+EPS)
            return true;
    }
    swap(r2.h,r2.w);
    alpha=asin(r1.h/cross)-asin(r2.h/cross);
    if(alpha>0&&2.0*alpha<PI)
    {
        if(r2.w*cos(alpha)+r2.h*sin(alpha)<=r1.w+EPS)
            return true;
    }
    return false;
}

/*=====
圆(By phoenixinter)
=====*/
struct Circle
{
    Point c;
    double r;
};

/*
判线段和圆相交，包括端点和相切

```

```

*/
bool intersect(Circle C, Line L)
{
    Point c=C.c;
    double t1=Dis(c,L.p1)-C.r, t2=Dis(c,L.p2)-C.r;
    if (t1<EPS || t2<EPS)
        return t1>-EPS || t2>-EPS;
    Point t=c;
    t.x+=L.p1.y-L.p2.y;
    t.y+=L.p2.x-L.p1.x;
    return times(t,L.p1,c)*times(t,L.p2,c)<EPS&&Dis2Line(c,L)-C.r<EPS;
}

double mario(double a, double b, double c)
{
    return acos(.5*(a*a+b*b-c*c)/(a*b));
}

// 计算两个圆相交部分的面积
// 通过测试: PKU 2546
double commonarea(Circle a, Circle b)
{
    double d=Dis(a.c,b.c), a1,a2,a3;
    if (fabs(a.r-b.r)>=d)
    {
        if (a.r<b.r)
            return PI*a.r*a.r;
        else
            return PI*b.r*b.r;
    }
    else if (a.r+b.r<=d)
        return 0;
    else
    {
        a1=mario(a.r,d,b.r);
        a2=mario(b.r,d,a.r);
        a3=mario(a.r,b.r,d);
        double ans=(a1*a.r*a.r+a2*b.r*b.r-a.r*b.r*sin(a3));
        return ans;
    }
}

/*
    计算三个点所组成的三角形的外接圆
    center是圆心的坐标, r是外接圆的半径
    通过测试: PKU 2242
*/
void outercircle(Point p1, Point p2, Point p3, Point& center, double& r)
{

```



```

    double a,b,c,d,e,f,g;
    a=p2.x-p1.x;b=p2.y-p1.y;
    c=p3.x-p1.x;d=p3.y-p1.y;
    e=a*(p1.x+p2.x)+b*(p1.y+p2.y);
    f=c*(p1.x+p3.x)+d*(p1.y+p3.y);
    g=2.0*(a*(p3.y-p2.y)-b*(p3.x-p2.x));
    center.x=(d*e-b*f)/g;center.y=(a*f-c*e)/g;
    r=sqrt((p1.x-center.x)*(p1.x-center.x)+
           (p1.y-center.y)*(p1.y-center.y));
}

/*
  输入两个点和半径，输出两个可能的圆心
  如果有圆心，返回true
  else return false
*/
bool centre(Point a,Point b,double r,Point& p1,Point& p2)
{
    double rise,run,theta,chordlen,perplen,tantheta,tantheta1;
    Point temp;
    temp.x=a.x-b.x;temp.y=a.y-b.y;
    chordlen=sqrt(temp.x*temp.x+temp.y*temp.y);
    if(chordlen>2*r)
        return false;
    tantheta=sqrt(4*r*r-chordlen*chordlen)/chordlen;
    run=(a.x-b.x)/2;rise=(a.y-b.y)/2;
    p1.x=(a.x+b.x)/2+rise*tantheta;p1.y=(a.y+b.y)/2-run*tantheta;
    p2.x=(a.x+b.x)/2-rise*tantheta;p2.y=(a.y+b.y)/2+run*tantheta;
    return true;
}

/*
  计算两个圆外公切线的交点
  如果不存在，返回false
  通过测试：ZJU 1199
*/
double sqr(double x)
{
    return x*x;
}

bool outertangent(Circle a,Circle b,Point& p)
{
    if(leq(sqr(a.c.x-b.c.x)+sqr(a.c.y-b.c.y),(a.r-b.r)*(a.r-b.r))
        ||eq(a.r,b.r))
        return false;
    p.x=(b.c.x*a.r-a.c.x*b.r)/(a.r-b.r);
    p.y=(b.c.y*a.r-a.c.y*b.r)/(a.r-b.r);
    return true;
}

```

```

/*=====
圆(By freezy)
=====*/

/*求点于圆切线的切点
   poi : 点
   cc : 圆
   输出: result1,result2 为两个切点
*/
void TangentPoint_PC(Point poi, Circle cc, Point& result1, Point& result2){
    double line=sqrt((poi.x-cc.c.x)*(poi.x-cc.c.x)+
                     (poi.y-cc.c.y)*(poi.y-cc.c.y));
    double angel=acos(cc.r/line);
    Point unitvector, lin;
    lin.x=poi.x-cc.c.x;
    lin.y=poi.y-cc.c.y;
    unitvector.x=lin.x/sqrt(lin.x*lin.x+lin.y*lin.y)*cc.r;
    unitvector.y=lin.y/sqrt(lin.x*lin.x+lin.y*lin.y)*cc.r;
    result1=Rotate(unitvector,-angel);
    result2=Rotate(unitvector,angel);
    result1.x+=cc.c.x;
    result1.y+=cc.c.y;
    result2.x+=cc.c.x;
    result2.y+=cc.c.y;
    return;
}

/* 求两圆的外公切线的切点
   两圆: c1,c2;
   输出四个点:
   p1,p2,p3,p4
   其中p1,p2在c1上, p3,p4在c2上。
   而且p1,p3为一条外公切线
   p2,p4为一条外公切线
*/
void TangentPoint_CC_out(Circle c1, Circle c2, Point &p1, Point &p2, Point &p3
    double line=sqrt((c1.c.x-c2.c.x)*(c1.c.x-c2.c.x)+
                     (c1.c.y-c2.c.y)*(c1.c.y-c2.c.y));
    double angell=acos((fabs(c1.r-c2.r)/line));
    Point unitvector1, unitvector2, lin1, lin2;
    if (c1.r>=c2.r){
        lin1.x=c2.c.x-c1.c.x;
        lin1.y=c2.c.y-c1.c.y;
        unitvector1.x=lin1.x/sqrt(lin1.x*lin1.x+lin1.y*lin1.y)*c1.r;
        unitvector1.y=lin1.y/sqrt(lin1.x*lin1.x+lin1.y*lin1.y)*c1.r;
        lin2.x=c1.c.x-c2.c.x;
        lin2.y=c1.c.y-c2.c.y;
        unitvector2.x=lin2.x/sqrt(lin2.x*lin2.x+lin2.y*lin2.y)*c2.r;

```

```

        unitvector2.y=lin2.y/sqrt(lin2.x*lin2.x+lin2.y*lin2.y)*c2.r;
        p1=Rotate(unitvector1, angl1);
        p2=Rotate(unitvector1, -angl1);
        p3=Rotate(unitvector2, -(PI-angl1));
        p4=Rotate(unitvector2, (PI-angl1));
        p1.x+=c1.c.x; p1.y+=c1.c.y;
        p2.x+=c1.c.x; p2.y+=c1.c.y;
        p3.x+=c2.c.x; p3.y+=c2.c.y;
        p4.x+=c2.c.x; p4.y+=c2.c.y;
    }

    else {
        lin2.x=c1.c.x-c2.c.x;
        lin2.y=c1.c.y-c2.c.y;
        unitvector2.x=lin2.x/sqrt(lin2.x*lin2.x+lin2.y*lin2.y)*c2.r;
        unitvector2.y=lin2.y/sqrt(lin2.x*lin2.x+lin2.y*lin2.y)*c2.r;
        p3=Rotate(unitvector2, angl1);
        p4=Rotate(unitvector2, -angl1);
        lin1.x=c2.c.x-c1.c.x;
        lin1.y=c2.c.y-c1.c.y;
        unitvector1.x=lin1.x/sqrt(lin1.x*lin1.x+lin1.y*lin1.y)*c1.r;
        unitvector1.y=lin1.y/sqrt(lin1.x*lin1.x+lin1.y*lin1.y)*c1.r;
        p1=Rotate(unitvector1, -(PI-angl1));
        p2=Rotate(unitvector1, (PI-angl1));
        p1.x+=c1.c.x; p1.y+=c1.c.y;
        p2.x+=c1.c.x; p2.y+=c1.c.y;
        p3.x+=c2.c.x; p3.y+=c2.c.y;
        p4.x+=c2.c.x; p4.y+=c2.c.y;
    }
    return ;
}

/*
求两圆的内公切线的切点
两圆: c1,c2;
输出: p1,p2,p3,p4;
其中p1,p2在c1上, p3,p4在c2上;
且p1,p3为一条内公切线
    p2,p4为一条内公切线
*/
void TangentPoint_CC_in( Circle c1, Circle c2, Point &p1, Point &p2, Point &p3, Point &p4
    double line=sqrt((c1.c.x-c2.c.x)*(c1.c.x-c2.c.x)+
        (c1.c.y-c2.c.y)*(c1.c.y-c2.c.y));
    double angl1=acos((c1.r+c2.r)/line);
    Point unitvector1, unitvector2, lin1, lin2;

    lin1.x=c2.c.x-c1.c.x;
    lin1.y=c2.c.y-c1.c.y;
    unitvector1.x=lin1.x/sqrt(lin1.x*lin1.x+lin1.y*lin1.y)*c1.r;
    unitvector1.y=lin1.y/sqrt(lin1.x*lin1.x+lin1.y*lin1.y)*c1.r;

```

```

        lin2.x=c1.c.x-c2.c.x;
        lin2.y=c1.c.y-c2.c.y;
        unitvector2.x=lin2.x/sqrt(lin2.x*lin2.x+lin2.y*lin2.y)*c2.r;
        unitvector2.y=lin2.y/sqrt(lin2.x*lin2.x+lin2.y*lin2.y)*c2.r;

        p1=Rotate(unitvector1,angell);
        p2=Rotate(unitvector1,-angell);

        p3=Rotate(unitvector2,angell);
        p4=Rotate(unitvector2,-angell);

        p1.x+=c1.c.x;p1.y+=c1.c.y;
        p2.x+=c1.c.x;p2.y+=c1.c.y;
        p3.x+=c2.c.x;p3.y+=c2.c.y;
        p4.x+=c2.c.x;p4.y+=c2.c.y;

        return ;
    }

/*
判断点是否在圆内
返回值：点p在圆内(包括边界)时，返回true
用途：因为圆为凸集，所以判断点集，折线，多边形是否在圆内时，只需要逐一判断点是否在圆内即可。
*/
//(未测试)

bool point_in_circle(Point o,double r,Point p)
{
    double d2=(p.x-o.x)*(p.x-o.x)+(p.y-o.y)*(p.y-o.y);
    double r2=r*r;
    return d2<r2||fabs(d2-r2)<EPS;
}

/*
用途：求不共线的三点确定一个圆
输入：三个点p1,p2,p3
返回值：如果三点共线，返回false；反之，返回true。圆心由q返回，半径由r返回
*/
//(未测试)
bool cocircle(Point p1,Point p2,Point p3,Point &q,double &r)
{
    double a1,a2,a3;
    a1 = Dis(p1,p2);
    a2 = Dis(p2,p3);
    a3 = Dis(p1,p3);
    double s = Area(a1, a2, a3);

```

```

    if (fabs(s) < EPS) return false;

    r = a1 * a2 * a3 / ( 4 * s );

    Point p;
    p.x = p2.x - p1.x;
    p.y = p2.y - p1.y;

    double theda = acos(a1 / 2 / r);
    q = Rotate(p,theda);
    return true;
}

/*
    判断圆于线段是否相交
    输入: C,L
    输出: true表示相交
         flase表示不相交
*/

bool intersect1(Circle C,Line L)
{
    Point c=C.c;
    double t1=Dis(c,L.p1)-C.r,t2=Dis(c,L.p2)-C.r;
    if (t1<EPS||t2<EPS)
        return t1>-EPS||t2>-EPS;
    Point t=c;
    t.x+=L.p1.y-L.p2.y;
    t.y+=L.p2.x-L.p1.x;
    return times(t,L.p1,c)*times(t,L.p2,c)<EPS&&Dis2Line(c,L)-C.r<EPS;
}

/*

/*
    求两圆的相交面积
*/
double commonarea1(Circle a,Circle b)
{
    double d=Dis(a.c,b.c),a1,a2,a3;
    if (fabs(a.r-b.r)>=d)
    {
        if(a.r<b.r)
            return PI*a.r*a.r;
        else
            return PI*b.r*b.r;
    }
    else if(a.r+b.r<=d)
        return 0;
}

```

```

    else
    {
        a1=mario(a.r,d,b.r);
        a2=mario(b.r,d,a.r);
        a3=mario(a.r,b.r,d);
        double ans=(a1*a.r*a.r+a2*b.r*b.r-a.r*b.r*sin(a3));
        return ans;
    }
}

/*
    求三角形内切圆圆心
    输入: p1,p2,p3;
    输出: q,r;
*/

void Inscribed_circle(Point p1, Point p2, Point p3, Point &q , double &r)
{
    double a1,a2,a3;
    a1 = Dis(p1,p2);
    a2 = Dis(p1,p3);
    a3 = Dis(p2,p3);
    r = 2 * fabs(Area(p1,p2,p3)) / (a1 + a2 + a3);
    double theda = (mario(a1,a2,a3)) * 0.5;
    double Len = r / sin(theda);
    Line l;
    l.p1 = p1;
    l.p2 = p2;
    if (Relation(p3,l) == 2) theda = - theda;
    Point temp;
    temp.x = p2.x - p1.x;
    temp.y = p2.y - p1.y;
    temp = Rotate(temp,theda);
    temp.x = temp.x / a1 * Len;
    temp.y = temp.y / a1 * Len;
    temp.x+=p1.x;temp.y+=p1.y;
    q = temp;
};

```

5.2 Geometry 3D

See [hyh's_paper](#)

5.3 Convex Hull(Graham Algorithm)

```

//Graham-Scan
//计算平面点集的凸包
///通过测试:
//PKU 2178,ZJU 1453,PKU 2595

```

```

#include <iostream>
#include <cmath>
#include <algorithm>
using namespace std;

#define EPS 1.0e-8
#define MAXN 100
#define eq(x,y) (fabs((x)-(y))<EPS)
#define geq(x,y) (x+EPS>=y)

struct Point {
    double x,y,angle,dis;
};
bool operator<(const Point& p1,const Point& p2) {
    if(eq(p1.angle,p2.angle))
        return p1.dis<p2.dis;
    return p1.angle<p2.angle;
}

//计算叉积(p1-p0)*(p2-p0)
double times(Point p0,Point p1,Point p2) {
    return ((p1.x-p0.x)*(p2.y-p0.y)-(p1.y-p0.y)*(p2.x-p0.x));
}

/*
Graham扫描求点集凸包
输入:pts 点集数组
n 点集大小
输出: con 凸包上的顶点数组, 按逆时针排序
m 凸包上的顶点数目
*/
void Graham(Point pts[],int n,Point con[],int& m) {
    int i,k;
    m=0;
    if(n<3) {
        con[0]=pts[0];
        con[1]=pts[1];
        return;
    }
    /*选取pts中y坐标最小的pts[k],
    如果这样的点有多个, 则取最左边的一个
    */
    k=0;
    for(i=1;i<n;i++) {
        if(eq(pts[i].y,pts[k].y)) {
            if(pts[i].x<=pts[k].x) k=i;
        } else if(pts[i].y<pts[k].y) k=i;
    }
    swap(pts[0],pts[k]);
    for(i=1;i<n;i++) {

```

```

        pts[i].angle=atan2(pts[i].y-pts[0].y,pts[i].x-pts[0].x);
        pts[i].dis=(pts[i].x-pts[0].x)*(pts[i].x-pts[0].x)+
            (pts[i].y-pts[0].y)*(pts[i].y-pts[0].y);
    }
    sort(pts+1,pts+n);
    con[m++]=pts[0];
    for(i=1;i<n;i++) {
        //如果有极角相同的点，只去相对于pts[0]最远的一个
        if(i+1<n&&eq(pts[i].angle,pts[i+1].angle))
            continue;
        if(m>=3)
            while(geq(times(con[m-2],pts[i],con[m-1]),0)) m--;
        con[m++]=pts[i];
    }
}

int main() {
    int i,n,m;
    Point pts[MAXN],con[MAXN];
    scanf("%d",&n);
    for(i=0;i<n;i++)
        scanf("%lf_%lf",&pts[i].x,&pts[i].y);
    Graham(pts,n,con,m);
    for(i=0;i<m;i++)
        printf("%lf_%lf\n",con[i].x,con[i].y);
    return 0;
}

```

5.4 binary div to determin whether all points is lay the same side of a line

```

//二分快速判断平面上的点集是否在一条直线的同一侧
//writen by chenkun
//通过测试:pku 1912
//输入p: 平面点集
// 每次查询的直线(p1,p2)
// 调用solve(0,m-1)返回结果,m为点集凸包顶点数
// true:点集在直线的同一侧
// false:否则
#include "stdio.h"
#include "stdlib.h"
#include "math.h"
#define MAXN 100000
#define EPS 1.0e-6

struct point
{
    double x, y;
} p[MAXN+1],h[MAXN+1],h2[MAXN+1],p1,p2;

```



```

typedef struct point point;
typedef int (*compfn)(const void*,const void*);
int compare(point *a, point *b) {
    if ((a->x-b->x<-EPS) || ((fabs(a->x-b->x)<=EPS)&&a->y-b->y<-EPS))
        return -1;
    if ((a->x-b->x>EPS) || ((fabs(a->x-b->x)<=EPS)&&a->y-b->y>EPS))
        return 1;
    return 0;
}
double distance(const point& p1,const point& p2) {
    return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y));
}
double multiply(const point& sp,const point& ep,const point& op) {
    return ((sp.x-op.x)*(ep.y-op.y)-(ep.x-op.x)*(sp.y-op.y));
}

int partition(point a[],int p,int r) {
    int i=p,j=r+1,k;
    double ang,dis;
    point R,S;
    k=(p+r)/2; //防止快排退化
    R=a[p];
    a[p]=a[k];
    a[k]=R;
    R=a[p];
    dis=distance(R,a[0]);
    while(1) {
        while(1) {
            ++i;
            if(i>r) {
                i=r;
                break;
            }
            ang=multiply(R,a[i],a[0]);
            if(ang>0)
                break;
            else if(ang==0) {
                if(distance(a[i],a[0])>dis)
                    break;
            }
        }
        while(1) {
            --j;
            if(j<p) {
                j=p;
                break;
            }
            ang=multiply(R,a[j],a[0]);
            if(ang<0)

```

```

        break;
    else if (ang==0) {
        if (distance(a[j],a[0])<dis)
            break;
    }
    }
    if (i>=j) break;
    S=a[i];
    a[i]=a[j];
    a[j]=S;
}
a[p]=a[j];
a[j]=R;
return j;
}

void anglesort(point a[],int p,int r) {
    if(p<r) {
        int q=partition(a,p,r);
        anglesort(a,p,q-1);
        anglesort(a,q+1,r);
    }
}

void Graham_scan(point PointSet[],point ch[],int n,int &len) {
    int i,k=0,top=2;
    point tmp;
    //选取PointSet中y坐标最小的点PointSet[k],
    //如果这样的点有多个.则取最左边的一个
    for(i=1;i<n;i++)
        if(PointSet[i].x<PointSet[k].x||
            (PointSet[i].x==PointSet[k].x)&&(PointSet[i].y<PointSet[k].y))
            k=i;
    tmp=PointSet[0];
    PointSet[0]=PointSet[k];
    PointSet[k]=tmp; //现在PointSet中y坐标最小的点在PointSet[0]
    /* 对顶点按照相对PointSet[0]的极角从小到大进行排序, 极角相同
       的按照距离PointSet[0]从近到远进行排序 */
    anglesort(PointSet,1,n-1);
    if(n<3) {
        len=n;
        for(int i=0;i<n;i++) ch[i]=PointSet[i];
        return ;
    }

    ch[0]=PointSet[0];
    ch[1]=PointSet[1];
    ch[2]=PointSet[2];
    for(i=3;i<n;i++){
        while (multiply(PointSet[i],ch[top],ch[top-1])>=0) top--;
    }
}

```

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```

        ch[++top]=PointSet[i];
    }
    len=top+1;
}

double cross(point a, point b, point c) {
    return ((b.x-a.x)*(c.y-b.y)-(c.x-b.x)*(b.y-a.y));
}

//计算直线(rp1,rp2)与(p1,p2)的交点res
void cross_point(point rp1, point rp2, point& res) {
    if (fabs(rp1.x-rp2.x)<=EPS) {
        res.x=rp1.x;
        res.y=p1.y+(p2.y-p1.y)*(res.x-p1.x)/(p2.x-p1.x);
        return;
    }
    if (fabs(p1.x-p2.x)<=EPS) {
        res.x=p1.x;
        res.y=rp1.y+(rp2.y-rp1.y)*(res.x-rp1.x)/(rp2.x-rp1.x);
        return;
    }
    double a=p1.y;
    double b=(p2.y-p1.y)/(p2.x-p1.x);
    double c=rp1.y;
    double d=(rp2.y-rp1.y)/(rp2.x-rp1.x);
    res.x=(a-c-b*p1.x+d*rp1.x)/(d-b);
    res.y=a+b*(res.x-p1.x);
}

bool solve(int l, int r) {
    point cro, p3;
    double e;
    int mid=(l+r)>>1;
    if (l>r) return false;
    if (fabs((h[mid].x-h[1].x)*(p2.y-p1.y)-
        (p2.x-p1.x)*(h[mid].y-h[1].y))<=EPS) {
        e=cross(p1, h[1], h[mid]);
        if (l==r&&fabs(e)>EPS) return false;
        if (fabs(e)<=EPS) {
            return true;
        } else if (e<0) {
            return solve(l, mid-1);
        } else {
            return solve(mid+1, r);
        }
    }
    cross_point(h[1], h[mid], cro);
    if (cro.x>=h[1].x&&cro.x<=h[mid].x&&
        ((cro.y>=h[1].y&&cro.y<=h[mid].y)||
        (cro.y>=h[mid].y&&cro.y<=h[1].y))) {

```

```

        return true;
    }
    if (cro.x > h[mid].x) {
        p3.x = h[mid].x + p2.x - cro.x;
        p3.y = h[mid].y + p2.y - cro.y;
        double c1 = cross(h[1], h[mid], p3);
        double c2 = cross(p3, h[mid], h[mid+1]);
        if (c1*c2 >= 0) {
            return solve(mid+1, r);
        }
        c1 = cross(h[mid-1], h[mid], p3);
        c2 = cross(p3, h[mid], h[1]);
        if (c1*c2 >= 0) {
            return solve(1, mid-1);
        }
        return false;
    } else {
        p3.x = h[1].x + p2.x - cro.x;
        p3.y = h[1].y + p2.y - cro.y;
        double c1 = cross(h[mid], h[1], p3);
        double c2 = cross(p3, h[1], h[(int)h[0].x-1]);
        if (c1*c2 >= 0) {
            return solve(mid+1, r);
        }
        c1 = cross(h[2], h[1], p3);
        c2 = cross(p3, h[1], h[mid]);
        if (c1*c2 >= 0) {
            return solve(1, mid-1);
        }
        return false;
    }
}

int main() {
    int n;
    int i, len;
    scanf("%d", &n);
    for (i=0; i<n; i++) {
        scanf("%lf %lf", &p[i].x, &p[i].y);
    }
    Graham_scan(p, h2, n, len);
    h[0].x = len+1;
    h[1].x = h2[0].x;
    h[1].y = h2[0].y;
    for (int i=len-1; i>=1; i--) h[len+1-i].x = h2[i].x, h[len+1-i].y = h2[i].y;
    while (scanf("%lf %lf %lf %lf", &p1.x, &p1.y, &p2.x, &p2.y) != EOF) {
        if (solve(2, (int)h[0].x-1)) {
            printf("BAD\n");
        } else {
            printf("GOOD\n");
        }
    }
}

```

5.4. *BINARY DIV TO DETERMIN WHETHER ALL POINTS IS LAY THE SAME SIDE OF A LINE*89

```
        }  
    }  
    return 0;  
}
```


Chapter 6

Mathematics

6.1 Basic Number Theory

```
//扩展欧几里德
//使得m*x+n*y=1,返回m,n的最大公约数
long long exgcd(long long m, long long & x, long long n, long long & y)
{
    long long x1, y1, x0, y0;
    x0=1; y0=0;
    x1=0; y1=1;
    long long r=(m%n+n)%n;
    long long q=(m-r)/n;
    x=0; y=1;
    while(r)
    {
        x=x0-q*x1; y=y0-q*y1; x0=x1; y0=y1;
        x1=x; y1=y;
        m=n; n=r; r=m%n;
        q=(m-r)/n;
    }
    return n;
}

//快速取模mod exp, calculates a pow b mod n
int ModExp(int a, int b, int n) {
    int c=1, d=a;
    while(b) {
        if(b&1) c=(c*d)%n;
        d=(d*d)%n;
        b>>1;
    }
    return c;
}

//求解模线性方程ax=b(mod n), n!=0
void modular_linear_equation(long long a, long long b, long long n) {
    long long e, d, x, y;
```

```

    int i;
    d=exgcd(a,x,n,y);
    if(b%d!=0) printf("No answer!\n");
    else {
        e=x*(b/d)%n;
        for(i=0;i<d;i++)
            printf("The %dth answer is %lld\n",i+1,(e+i*(n/d))%n);
    }
}

/*中国余数订立, 求解模线性方程组
a=B[1](mod W[1])
a=B[2](mod W[2])
.....
a=B[n](mod W[n])
其中W,B已知,W[i]≠0且W[i]与W[j]互质, 求a
*/
int linear_modular_equation_system(long long B[], long long W[], int k)
{
    int i;
    long long d,x,y,a=0,m,n=1;
    for(i=0;i<k;i++) n*=W[i];
    for(i=0;i<k;i++) {
        m=n/W[i];
        d=exgcd(W[i],x,m,y);
        a=(a+y*m*B[i])%n;
    }
    if(a>0) return a;
    else return a+n;
}

/*=====
二分求矩阵幂
=====*/
//list[0][][]: 单位矩阵
//list[1][][]: 结果矩阵,初始为单位矩阵
void matrixmul(int mode) { //mode 0:乘单位矩阵1:自乘
    for(int i=1;i<=nodes;i++)
        for(int j=1;j<=nodes;j++)
            temp[i][j]=0;
    for(int i=1;i<=nodes;i++)
        for(int j=1;j<=nodes;j++)
            for(int k=1;k<=nodes;k++)
                temp[i][j]+=graph[1][i][k]*graph[mode][k][j];
    for(int i=1;i<=nodes;i++)
        for(int j=1;j<=nodes;j++)
            graph[1][i][j]=temp[i][j]%100000;
}

void binarymul(int n) {

```



```

        if (n==1) return;
        binarymul(n/2);
        matrixmul(1);
        if (n%2==1) matrixmul(0);
    }

// 求矩阵的n次方:binarymul(n);

/*=====
                        素数测试随机化算法
=====*/
// 调用isprim(n,s)返回结果, true为素数
// n为要判断的数,s为随机测试次数
long long a,N;

long long mimod(long long n) {
    if (n==1) return a%N;
    long long t=mimod(n/2);
    long long r=(t*t)%N;
    if (n%2==1) r=(r*a)%N;
    return r;
}

bool isprim(long long n,int s) {
    N=n;
    for (int i=0;i<s;i++) {
        a=rand()%(n-1)+1;
        if (mimod(n-1)!=1)
            return false;
    }
    return true;
}

/*=====
                        数的因子分解
=====*/
void PrimeFactor(long long n) {
    int i,j;
    long long res=1;

    if (n < 2) {
        return n;
    }
    i = 2;
    while (i <= (int) (sqrt ((double) (n)))) {
        j = 0;
        while (n % i == 0) {
            n /= i;
            j ++;
        }
    }
}

```

```

        if(j){
                                //因子i出现j次
        }
        i++;
    }
    if (n!= 1) {
        //最后还有个因子n
    }
}

```

6.2 Gaussian Elimination

see hyh's_paper

6.3 Romberg Numerical Integration

```

double F(double t) {
    return (R/(f*cos(t)+h));
}

double romberg(double a, double b, double eps){
    vector<double> R;
    int k=-1;
    double r=0.5*(b-a)*(F(a)+F(b));
    R.push_back(r);
    do{
        k+=1;
        r=0.0;
        for (int i=0; i<pow(2.0, k); i++)
            r+=F(a+(b-a)*(i+0.5)/pow(2.0, k));
        r*=(b-a)/pow(2.0, k+1);
        r+=0.5*R[k];
        R.push_back(r);
        for (int m=0; m<=k; m++)
            R[k-m]=(pow(4.0, m+1)*R[k+1-m]-R[k-m])/(pow(4.0, m+1)-1);
    } while (fabs(R[0]-R[1])>eps);
    return R[0];
}

```