

Semidefinite Relaxations of Truncated Least-Squares in Robust Rotation Search: Tight or Not

Liangzu Peng, Mahyar Fazlyab, René Vidal

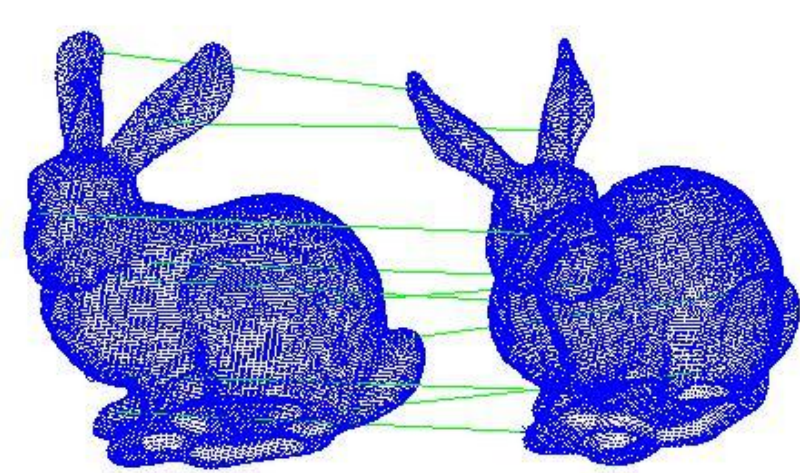
Rotation Search: Problem Setup

Goal:

- Find 3D rotation R_0^* that best aligns 2 point sets

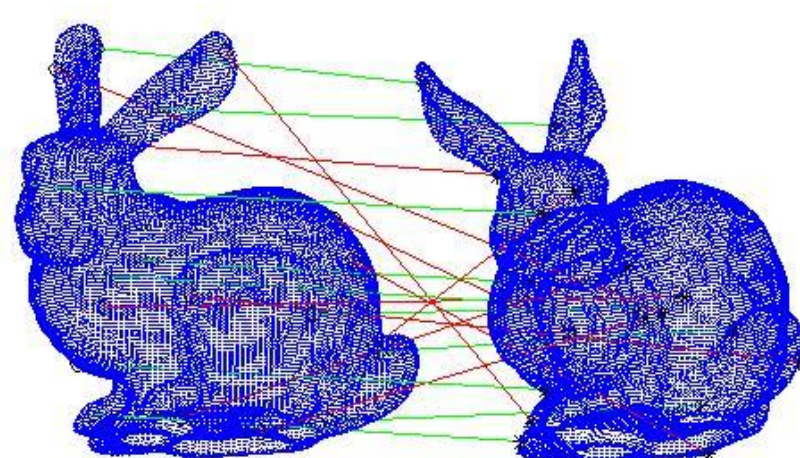
Case without Outliers (Wahba, 1965 [1]):

- $y_i \approx R_0^* x_i, \quad i = 1, \dots, \ell$
 - $R_0^* \in SO(3)$
 - Optimization problem:
- $$\min_{R_0 \in SO(3)} \sum_{i=1}^{\ell} \|y_i - R_0 x_i\|_2^2$$



Case with Outliers:

- Inliers: $y_i \approx R_0^* x_i$
- Outliers: (x_i, y_i) arbitrary
- Optimization problem:

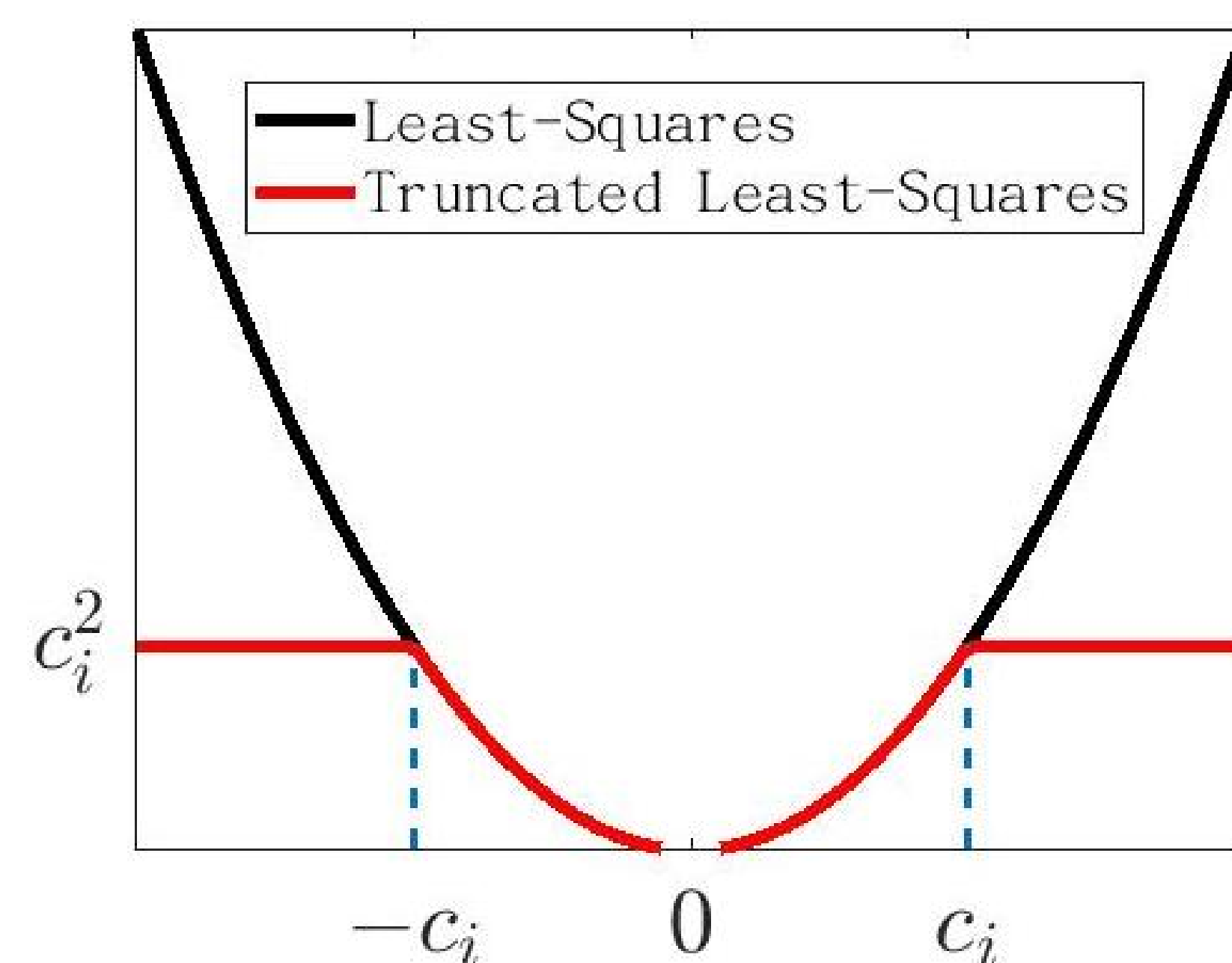


Truncated Least-Squares (TLS-R) [1]:

$$\min_{R_0 \in SO(3)} \sum_{i=1}^{\ell} \min \left\{ \|y_i - R_0 x_i\|_2^2, c_i^2 \right\}$$

(Huber, 1964 [2])

Truncation Parameter



Key challenge:

- Under what conditions can we solve (TLS-R) to global optimality in polynomial time?

From (TLS-R) to Semidefinite Relaxations (SDR)

$$(\text{TLS-R}) \xLeftrightarrow{(1)} (\text{TLS-Q}) \xLeftrightarrow{(2), (3)} (\text{QCQP}) \xRightarrow{(4) \text{ lifting}} (\text{SDR})$$

Step (1): 3D rotations $R_0 \Leftrightarrow$ unit quaternions w_0

$$\min_{R_0 \in SO(3)} \sum_{i=1}^{\ell} \min \left\{ \|y_i - R_0 x_i\|_2^2, c_i^2 \right\}$$

$$\downarrow$$

$$\min_{w_0 \in \mathbb{S}^3} \sum_{i=1}^{\ell} \min \left\{ w_0^\top Q_i w_0, c_i^2 \right\}$$

Step (2): $\min\{a, b\} = \min_{\theta_i \in \{0,1\}} \theta_i a + (1 - \theta_i) b$

$$\min_{w_0 \in \mathbb{S}^3} \sum_{i=1}^{\ell} \min_{\theta_i \in \{0,1\}} \left(\theta_i w_0^\top Q_i w_0 + (1 - \theta_i) c_i^2 \right)$$

Step (3): $w_i := \theta_i w_0$, so $\theta_i = w_i^\top w_0$ [3]

$$\min_{\substack{w_0 \in \mathbb{S}^3 \\ w_i \in \{0, w_0\}}} \sum_{i=1}^{\ell} \left(w_i^\top Q_i w_0 + (1 - w_i^\top w_0) c_i^2 \right)$$

$$\min_{w \in \mathbb{R}^{4(\ell+1)}} \text{trace}(Q w w^\top)$$

$$\text{s.t. } [w w^\top]_{0i} = [w w^\top]_{ii}, \quad \forall i \in \{1, \dots, \ell\}$$

$$\text{trace}([w w^\top]_{00}) = 1$$

Step (4): From (QCQP) to (SDR)

$$\min_{W \succeq 0} \text{trace}(Q W)$$

$$\text{s.t. } [W]_{0i} = [W]_{ii}, \quad \forall i \in \{1, \dots, \ell\}$$

$$\text{trace}([W]_{00}) = 1$$

Tightness of (SDR)

Definition (Tightness):

Let $w \in \mathbb{R}^{4(\ell+1)}$ be a global minimizer of (QCQP). We say (SDR) is tight if its solution is $W^* = w w^\top$.

Main Results:

Positive Result

- Inlier Assumption:** We have $c_i^2 > f(\text{noise})$ for every inlier point pair (y_i, x_i) . Here f is a complicated function of noise with $f(0) = 0$.
- Outlier Assumption:** Each outlier pair (y_j, x_j) is drawn independently and randomly from $\mathbb{R}^3 \times \mathbb{R}^3$ according to some continuous probability distribution, and c_j^2 is set small enough.
- Under the *Inlier* and *Outlier Assumption*, (SDR) is tight.

Negative Result

(SDR) is not tight for “adversarial” outliers.

Remarks:

- For the exact definitions of noise function f and outlier truncation parameter c_j^2 , see our paper.
- “Adversarial” outliers can be point pairs that are defined by a rotation far from the ground-truth.

Acknowledgement: work supported by grants NSF 1704458, NSF 1934979, ONR MURI 503405-78051.

References

- [1] G. Wahba. A least squares estimate of satellite attitude. SIAM Review, 1965.
- [2] P. J. Huber. Robust estimation of a location parameter. Ann. Math. Stat., 1964.
- [3] H. Yang & L. Carlone. A quaternion-based certifiably optimal solution to the Wahba problem with outliers. ICCV 2019.