

Online Appendix

1 Proof of Proposition 1

Assumption 1. F_1, \dots, F_N are all L -smooth: for all \mathbf{v} and \mathbf{w} , $F_n(\mathbf{v}) \leq F_n(\mathbf{w}) + (\mathbf{v} - \mathbf{w})^T \nabla F_n(\mathbf{w}) + \frac{L}{2} \|\mathbf{v} - \mathbf{w}\|_2^2$

Assumption 2. F_1, \dots, F_N are all μ -strongly convex: for all \mathbf{v} and \mathbf{w} , $F_n(\mathbf{v}) \geq F_n(\mathbf{w}) + (\mathbf{v} - \mathbf{w})^T \nabla F_n(\mathbf{w}) + \frac{\mu}{2} \|\mathbf{v} - \mathbf{w}\|_2^2$.

Assumption 3. For each client $n \in \mathcal{N}$, the stochastic gradient of F_n is unbiased with its variance bounded by σ_n^2 .

Assumption 4. The expected squared norm of stochastic gradients for each client is uniformly bounded by G_n^2 .

Proposition 1 (Convergence Upper Bound with an Arbitrary Sampling Probability \mathbf{q}). *For any given client sampling probability profile \mathbf{q} in Algorithm 1, if we choose the decaying learning rate $\eta_r = \frac{2}{\max\{8L, \mu E\} + \mu r}$, the model parameter after R rounds $\mathbf{w}^R(\mathbf{q})$ has the optimality gap as follows*

$$\mathbb{E}[F(\mathbf{w}^R(\mathbf{q}))] - \min_{\mathbf{w}} F(\mathbf{w}) \leq \frac{1}{R} \left(\sum_{n=1}^N \frac{s_n}{q_n} + \beta \right), \quad (1)$$

where $s_n = \frac{8LEG_n^2}{\mu^2 N^2}$, $\beta = \frac{2L}{\mu^2 E} D + \frac{12L^2}{\mu^2 E} \Gamma + \frac{4L^2}{\mu E} \|\mathbf{w}_0 - \mathbf{w}^*\|^2$, $D = \sum_{n=1}^N (p_n \sigma_n)^2 + 8 \sum_{n=1}^N p_n G_n^2 E^2$, and $\Gamma = F^* - \frac{1}{N} \sum_{n=1}^N \min_{\mathbf{w}} F_n(\mathbf{w})$.

Proof. The proof follows a similar argument of weighted client sampling in [1], where we first show that for any client sampling probabilities \mathbf{q} , the variance between the aggregated model \mathbf{w}^{r+1} and the virtual global model under full participation (i.e., $\bar{\mathbf{w}}^{r+1}$) is bounded as follows:

$$\mathbb{E}_{\mathcal{S}(\mathbf{q})^r} \|\mathbf{w}^{r+1} - \bar{\mathbf{w}}^{r+1}\|^2 \leq 4 \sum_{n=1}^N \frac{(1 - q_n)}{q_n} \left(\frac{\eta^r EG}{N} \right)^2. \quad (2)$$

Note that the main difference of (2) compared to the weighted sampling in is that the client sampling probability q_n is independent among each other. In particular, when $q_n = 1$ for all n , the variance in (2) is tightly bounded by zero, as the aggregated model \mathbf{w}^{r+1} in the left hand side of (2) recovers the aggregated model of full client participation $\bar{\mathbf{w}}^{r+1}$. Then, we use mathematical induction to obtain a non-recursive bound on $\mathbb{E}_{\mathcal{S}(\mathbf{q})^r} \|\mathbf{w}^R - \mathbf{w}^*\|^2$, whose difference compared to the bound of full participation is the variance introduced in (2). After that, we converted the bound of $\mathbb{E}_{\mathcal{S}(\mathbf{q})^r} \|\mathbf{w}^R - \mathbf{w}^*\|^2$ to $\mathbb{E}[F(\mathbf{w}^R(\mathbf{q}))] - F^*$ using L -smoothness, which yields the additional term of $\sum_{n=1}^N \frac{1}{q_n}$ in (1) and concludes the proof. \square

2 Proof of Theorem 1

Theorem 1. *A mechanism $m = (q, r)$ is incentive compatible and individual rational if and only if*

- sampling probability $q(\tilde{c})$ is non-increasing in the reported cost \tilde{c} ;
- payment function $r(\tilde{c})$ has the following form:

$$r(\tilde{c}) = \tilde{c} + \frac{1}{q(\tilde{c})} \int_{\tilde{c}}^{c_{\max}} q(z) dz. \quad (3)$$

Proof. We first prove the “if” direction and then prove the “only if” direction.

We plug the payment function (3) into the client i ’s utility function, and obtain the following utility function when the agent i with true cost c_i reports \tilde{c}_i :

$$u(\tilde{c}_i; c_i) = q(\tilde{c}_i)(\tilde{c}_i - c_i) + \int_{\tilde{c}_i}^{c_{\max}} q(z) dz. \quad (4)$$

The derivative of $u(\tilde{c}_i; c_i)$ with respect to reported cost \tilde{c}_i is

$$\frac{\partial u(\tilde{c}_i; c_i)}{\partial \tilde{c}_i} = q'(\tilde{c}_i)(\tilde{c}_i - c_i). \quad (5)$$

Since sampling probability q is non-increasing in reported cost, we have $q'(\tilde{c}_i) \leq 0$. So the derivative $\frac{\partial u(\tilde{c}_i; c_i)}{\partial \tilde{c}_i} \geq 0$ if $\tilde{c}_i \leq c_i$ and $\frac{\partial u(\tilde{c}_i; c_i)}{\partial \tilde{c}_i} \leq 0$ if $\tilde{c}_i \geq c_i$. The agent can maximize his utility when he truthfully reports his cost $\tilde{c}_i = c_i$.

To prove individual rationality, we can verify that

$$\max_{\tilde{c}_i} u(\tilde{c}_i; c_i) = u(c_i; c_i) = \int_{c_i}^{c_{\max}} q(z) dz \geq 0, \forall c_i \leq c_{\max}. \quad (6)$$

Next, we prove the “only if” direction. By incentive compatibility, we have $\max_{\tilde{c}_i} u(\tilde{c}_i; c_i) = u(c_i; c_i)$. By envelope theorem, we have

$$\frac{\partial u(c_i; c_i)}{\partial c_i} = \left. \frac{\partial u(\tilde{c}_i; c_i)}{\partial \tilde{c}_i} \right|_{\tilde{c}_i=c_i} = -q(c_i). \quad (7)$$

Taking the integral from c_i to c_{\max} , we have

$$u(c_{\max}; c_{\max}) - u(c_i; c_i) = - \int_{c_i}^{c_{\max}} q(z) dz. \quad (8)$$

We consider the minimum payment that satisfies the individual rationality such that the client with the maximum cost obtains exactly zero utility, i.e., $u(c_{\max}; c_{\max}) = 0$. Then, we have

$$u(c_i; c_i) = \int_{c_i}^{c_{\max}} q(z) dz \Rightarrow r(c_i) = c_i + \frac{1}{q(c_i)} \int_{c_i}^{c_{\max}} q(z) dz. \quad (9)$$

Recall that

$$\frac{\partial u(\tilde{c}_i; c_i)}{\partial \tilde{c}_i} = q'(\tilde{c}_i)(\tilde{c}_i - c_i). \quad (10)$$

Incentive compatibility implies that there exists $\epsilon > 0$ such that for any $\tilde{c}_i \in (c_i - \epsilon, c_i)$, $\frac{\partial u(\tilde{c}_i; c_i)}{\partial \tilde{c}_i} \geq 0$ and for any $\tilde{c}_i \in (c_i, c_i + \epsilon)$, $\frac{\partial u(\tilde{c}_i; c_i)}{\partial \tilde{c}_i} \leq 0$. Notice that $\tilde{c}_i - c_i < 0$ for any $\tilde{c}_i \in (c_i - \epsilon, c_i)$, and $\tilde{c}_i - c_i > 0$ for any $\tilde{c}_i \in (c_i, c_i + \epsilon)$. This requires that $q'(\tilde{c}_i) \leq 0$ on $(c_i - \epsilon, c_i)$ and $(c_i, c_i + \epsilon)$. As incentive compatibility holds for all c_i , this in particular implies that $q'(\tilde{c}_i) \leq 0$ for all c_i , which shows that q is non-increasing. \square

3 Proof of Theorem 2

Assumption 5. The virtual cost $\phi(c)$ is non-decreasing in the cost c .

Theorem 2. Under Assumption 5, the optimal sampling probability is as follows:

1. Case 1: If $\bar{B} \leq \sqrt{\phi_{\min}} \mathbf{E}_c[\sqrt{\phi(c)}]$, the optimal sampling probability is

$$q^*(c_n) = \frac{1}{\sqrt{\phi(c_n)}} \cdot \frac{\bar{B}}{\mathbf{E}_c[\sqrt{\phi(c)}]}, \quad (11)$$

for all c_n .

2. Case 2: If $\sqrt{\phi_{\min}} \mathbf{E}_c[\sqrt{\phi(c)}] < \bar{B} < \mathbf{E}_c[\phi(c)]$, the optimal sampling probability is

$$q^*(c_n) = \begin{cases} 1, & c_n \leq \hat{c}; \\ \frac{1}{\sqrt{\phi(c_n)}} \cdot \frac{\bar{B} - \mathbf{E}_c[\phi(c) \cdot \mathbb{1}\{c \leq \hat{c}\}]}{\mathbf{E}_c[\sqrt{\phi(c)} \cdot \mathbb{1}\{c > \hat{c}\}]}, & c_n > \hat{c}. \end{cases} \quad (12)$$

Here, \hat{c} is solution to equation $H(x) \triangleq \mathbf{E}_c[\phi(c) \cdot \mathbb{1}\{c \leq x\}] + \mathbf{E}_c[\sqrt{\phi(c)} \cdot \mathbb{1}\{c > x\}]$, and there exists a unique $\hat{c} \in (c_{\min}, c_{\max})$ satisfying $\bar{B} = H(\hat{c})$. And \hat{c} can be computed by linear grid search over the support of γ .

3. Case 3: If $\bar{B} \geq \mathbf{E}_c[\phi(c)]$, the optimal sampling probability is $q^*(c_n) = 1$ for all c_n .

Proof. We start with the discrete version of P2, in which the cost take discrete values, and obtain discrete solution. Then we transform the discrete solution to continuous solution.

Discrete solution: We consider that the cost takes discrete values in the set $\{c_1, c_2, \dots, c_K\}$, where $c_1 < c_2 < \dots < c_K$. The corresponding virtual cost associated with cost c_k is ϕ_k :

$$\phi_1 = c_1, \quad (13)$$

$$\phi_k = c_k + \frac{c_k - c_{k-1}}{f(c_k)} \cdot F(c_{k-1}), k = 2, \dots, K, \quad (14)$$

$$r_K = c_K. \quad (15)$$

$$r_k = c_k + \sum_{j=k+1}^K \frac{q_j}{q_k} (c_j - c_{j-1}), \quad k = 1, 2, \dots, K-1. \quad (16)$$

where f is the probability of cost (i.e., $f(c_k) = P(c_k)$, $k = 1, 2, \dots, K$) and F is the cumulative density function of cost (i.e., $F(c_k) = \sum_{i=1}^k f(c_i)$, $k = 1, 2, \dots, K$). Then the probability of virtual cost is $P(\phi_k) = f_k$, $1 \leq k \leq K$. Define $\bar{B} = \frac{B}{RN}$.

Lemma 1. The optimal sampling probability is as follows:

1. If $\bar{B} \leq \sqrt{\phi_1} \sum_{k=1}^K f_k \sqrt{\phi_k}$, the optimal sampling probability is

$$q_k^* = \frac{1}{\sqrt{\phi_k}} \cdot \frac{\bar{B}}{\sum_{k=1}^K f_k \sqrt{\phi_k}}. \quad (17)$$

for all k .

2. If $\sqrt{\phi_1} \sum_{k=1}^K f_k \sqrt{\phi_k} < \bar{B} < \sum_{k=1}^K f_k \phi_k$, the optimal sampling probability is

$$q_k^* = \begin{cases} 1, & 1 \leq k \leq \hat{k}; \\ \frac{1}{\sqrt{\phi_k}} \cdot \frac{\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k}{\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}}, & k > \hat{k}. \end{cases} \quad (18)$$

Here, \hat{k} is defined as follows: Let $H(m) = \sum_{k=1}^m f_k \phi_k + \sum_{k=m+1}^K f_k \sqrt{\phi_k} \cdot \sqrt{\phi_{m+1}}$. There is a unique $\hat{k} \in \{1, 2, \dots, K-1\}$ satisfying $H(\hat{k}-1) < \bar{B} < H(\hat{k})$

3. If $\bar{B} \geq \sum_{k=1}^K f_k \phi_k$, the optimal sampling probability is $q_k^* = 1$ for all k .

Proof. Recall that we assume each client's cost is identically and independently distributed according to discrete distribution $f_k, k = 1, 2, \dots, K$. Define $\bar{B} \triangleq B/(NR)$. We write the expectation in objective function and budget constraint explicitly through discrete distribution and obtain the following problem:

$$\min_q \max_{s \in [s_{\min}, s_{\max}]} \sum_{k=1}^K s \cdot \frac{f_k}{q_k} \quad (19a)$$

$$\text{s.t.} \quad \sum_{k=1}^K f_k \cdot q_k \cdot \phi_k \leq \bar{B}; \quad (19b)$$

$$0 < q_k \leq 1, \forall k. \quad (19c)$$

Here, we drop the the monotonic constraint in Problem P2. Latter we will show that the solution indeed satisfies monotonic constraint. Notice that the objective function is an increasing function of s and thus the maximum is obtained when $s = s_{\max}$. Then the objective function of the optimization problem becomes $\sum_{k=1}^K s_{\max} \cdot f_k / q_k$. We find that the optimization problem is a convex problem. Thus, KKT conditions are sufficient and necessary for optimality. The Lagrangian of the optimization problem is

$$L(q, \lambda) = \sum_{k=1}^K s_{\max} \frac{f_k}{q_k} + \lambda \left(\sum_{k=1}^K f_k q_k \phi_k - \bar{B} \right) + \sum_{k=1}^K \lambda_k (q_k - 1). \quad (20)$$

Here we drop the constraint $q_k > 0$, for all k (this is without lose of generality and we will recover a positive solution later). According to KKT conditions, the optimal primal variables q_k^* and dual variables $\lambda^* \geq 0$, $\lambda_k^* \geq 0$ must satisfy

$$\frac{\partial L}{\partial q_k} = -s_{\max} \frac{f_k}{q_k^2} + \lambda^* f_k q_k^* + \lambda_k^* = 0 \quad (21)$$

$$\Rightarrow q_k^* = \sqrt{\frac{s_{\max} f_k}{\lambda^* f_k \phi_k + \lambda_k^*}}, \quad \forall k,$$

$$\lambda^* \left(\sum_{k=1}^K f_k q_k^* \phi_k - \bar{B} \right) = 0, \quad (22)$$

$$\lambda_k^* (q_k^* - 1) = 0, \quad \forall k, \quad (23)$$

$$\sum_{k=1}^K f_k \cdot q_k \cdot \phi_k \leq \bar{B}, \quad (24)$$

$$0 < q_k \leq 1, \forall k, \quad (25)$$

Next, we show that the solution in Theorem 1 exactly satisfies (21)-(25).

1. If $\bar{B} \leq \sqrt{\phi_1} \sum_{k=1}^K f_k \sqrt{\phi_k}$: we show that the primal variables

$$q_k^* = \frac{1}{\sqrt{\phi_k}} \cdot \frac{\bar{B}}{\sum_{k=1}^K f_k \sqrt{\phi_k}}, \quad \forall k, \quad (26)$$

and the dual variables

$$\lambda_k^* = 0, \quad \forall k, \quad (27)$$

$$\lambda^* = \frac{s_{\max}(\sum_{k=1}^K f_k \sqrt{\phi_k})^2}{\bar{B}^2}, \quad (28)$$

satisfy the KKT conditions (21)-(24). To see this, plugging the expressions of λ_k^* and λ^* into (21) yields

$$q_k^* = \sqrt{\frac{s_{\max}}{\lambda^* \phi_k}} = \frac{1}{\sqrt{\phi_k}} \cdot \frac{\bar{B}}{\sum_{k=1}^K f_k \sqrt{\phi_k}}, \quad (29)$$

which is exactly q_k^* in (26). Thus, condition in (21) holds. Meanwhile, we can see that

$$\begin{aligned} \sum_{k=1}^K f_k q_k^* \phi_k &= \sum_{k=1}^K f_k \phi_k \cdot \frac{1}{\sqrt{\phi_k}} \cdot \frac{\bar{B}}{\sum_{k=1}^K f_k \sqrt{\phi_k}} \\ &= \sum_{k=1}^K f_k \sqrt{\phi_k} \cdot \frac{\bar{B}}{\sum_{k=1}^K f_k \sqrt{\phi_k}} = \bar{B}. \end{aligned} \quad (30)$$

Thus, conditions in (22) and (24) holds. As $\lambda_k^* = 0$, condition in (23) holds. Finally, as $\bar{B} \leq \sqrt{\phi_1} \sum_{k=1}^K f_k \sqrt{\phi_k}$, we have

$$q_1^* = \frac{1}{\sqrt{\phi_1}} \cdot \frac{\bar{B}}{\sum_{k=1}^K f_k \sqrt{\phi_k}} \leq 1. \quad (31)$$

We can check that q^* is indeed decreasing ($q_{k+1}^* < q_k^*$, $k = 1, 2, \dots, K-1$) due to increasing virtual cost $\phi_1 < \phi_2 < \dots < \phi_K$. Thus, condition in (25) holds. In conclusion, we have shown the optimality of q_k^* in (26).

2. If $\sqrt{\phi_1} \sum_{k=1}^K f_k \sqrt{\phi_k} < \bar{B} < \sum_{k=1}^K f_k \phi_k$: first of all, we define $H(m)$ as follows:

$$H(m) = \sum_{k=1}^m f_k \phi_k + \sqrt{\phi_{m+1}} \cdot \sum_{k=m+1}^K f_k \sqrt{\phi_k}, \quad (32)$$

where $m = 0, 1, \dots, K$. Notice that we define $H(0) = \sqrt{\phi_1} \sum_{k=1}^K f_k \sqrt{\phi_k}$ and $H(K) = \sum_{k=1}^K f_k \phi_k$. We show the monotonicity of $H(m)$ as follows.

Lemma 2. $H(m+1) > H(m)$, $m = 0, 1, \dots, K-1$. And there exists a unique $\hat{k} \in \{1, 2, \dots, K-1\}$ satisfying $H(\hat{k}-1) < \bar{B} < H(\hat{k})$ for $\bar{B} \in (H(0), H(K))$.

Proof. We can see that

$$\begin{aligned}
& H(m+1) - H(m) \\
&= \sum_{k=1}^{m+1} f_k \phi_k + \sqrt{\phi_{m+2}} \cdot \sum_{k=m+2}^K f_k \sqrt{\phi_k} \\
&\quad - \sum_{k=1}^m f_k \phi_k - \sqrt{\phi_{m+1}} \cdot \sum_{k=m+1}^K f_k \sqrt{\phi_k} \\
&= f_{m+1} \phi_{m+1} + \sum_{k=m+2}^K f_k \sqrt{\phi_k} (\sqrt{\phi_{m+2}} - \sqrt{\phi_{m+1}}) \\
&\quad - f_{m+1} \phi_{m+1} \\
&> 0.
\end{aligned} \tag{33}$$

The inequality is due to increasing virtual cost, i.e., ϕ_k is increasing in k . Thus, $H(m)$ is increasing in m . For $\bar{B} \in (H(0), H(K))$, there exists a unique $\hat{k} \in \{1, 2, \dots, K-1\}$ such that $H(\hat{k}-1) < \bar{B} < H(\hat{k})$. \square

Next, we show that the primal variables

$$q_k^* = \begin{cases} 1, & 1 \leq k \leq \hat{k}; \\ \frac{1}{\sqrt{\phi_k}} \cdot \frac{\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k}{\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}}, & k > \hat{k}, \end{cases} \tag{34}$$

and dual variables

$$\lambda^* = \frac{s_{\max} (\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k})^2}{(\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k)^2}, \tag{35}$$

$$\lambda_k^* = \begin{cases} f_k (s_{\max} - \lambda^* \phi_k), & 1 \leq k \leq \hat{k}; \\ 0, & k > \hat{k}, \end{cases} \tag{36}$$

satisfy the KKT conditions (21)-(24).

To see this, for $k \in [1, \hat{k}]$, plugging the expressions of λ^* into (21) yields

$$q_k^* = 1, \tag{37}$$

which is exactly q_k^* in (34). Thus, conditions in (21), (23) and (25) hold. Next, we verify that $\lambda_k^* > 0$.

To see this, plugging the expression of λ^* into λ_k^* , we have

$$\lambda_k^* = s_{\max} f_k \left(1 - \frac{\phi_k (\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k})^2}{(\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k)^2} \right). \tag{38}$$

According to $\bar{B} > H(\hat{k}-1)$, i.e.,

$$\begin{aligned}
\bar{B} &> \sum_{k=1}^{\hat{k}-1} f_k \phi_k + \sqrt{\phi_{\hat{k}}} \cdot \sum_{k=\hat{k}}^K f_k \sqrt{\phi_k} \\
&= \sum_{k=1}^{\hat{k}} f_k \phi_k + \sqrt{\phi_{\hat{k}}} \cdot \sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k},
\end{aligned} \tag{39}$$

we have

$$\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k > \sqrt{\phi_{\hat{k}}} \cdot \sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}. \quad (40)$$

Combining increasing virtual cost ($\phi_k \leq \phi_{\hat{k}}$), we have

$$\frac{\sqrt{\phi_k} \cdot \sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}}{\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k} < 1, k \leq \hat{k}. \quad (41)$$

Thus,

$$\lambda_k^* = s_{\max} f_k \left(1 - \frac{\phi_k (\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k})^2}{(\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k)^2} \right) > 0 \quad (42)$$

As for $k \in [\hat{k} + 1, K]$, plugging the expressions of λ_k^* and λ^* into (21) yields

$$q_k^* = \frac{1}{\sqrt{\phi_k}} \cdot \frac{\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k}{\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}}, \quad (43)$$

which is exactly q_k^* in (34). Thus, condition in (21) holds. According to $\bar{B} < H(\hat{k})$, i.e.,

$$\bar{B} < \sum_{k=1}^{\hat{k}} f_k \phi_k + \sqrt{\phi_{\hat{k}+1}} \cdot \sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}, \quad (44)$$

we have

$$q_{\hat{k}+1}^* = \frac{1}{\sqrt{\phi_{\hat{k}+1}}} \cdot \frac{\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k}{\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}} < 1. \quad (45)$$

Thus, for $k \in [\hat{k} + 1, K]$, combining increasing virtual cost ($\phi_{k+1} > \phi_k$), we have $q_k^* \leq q_{\hat{k}+1}^* < 1$, which means condition in (25) holds. As $\lambda_k^* = 0$, condition in (23) holds.

Finally, for $k \in [1, K]$,

$$\begin{aligned} & \sum_{k=1}^K f_k \cdot q_k^* \cdot \phi_k \\ &= \sum_{k=1}^{\hat{k}} f_k \cdot \phi_k + \sum_{k=\hat{k}+1}^K \frac{f_k \phi_k}{\sqrt{\phi_k}} \cdot \frac{\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k}{\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}} \\ &= \sum_{k=1}^{\hat{k}} f_k \cdot \phi_k + \frac{\bar{B} - \sum_{k=1}^{\hat{k}} f_k \phi_k}{\sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k}} \cdot \sum_{k=\hat{k}+1}^K f_k \sqrt{\phi_k} \\ &= \bar{B}. \end{aligned} \quad (46)$$

That is, conditions in (22) and (24) hold. In conclusion, we have shown the optimality of q_k^* in (34).

3. If $\bar{B} \geq \sum_{k=1}^K f_k \phi_k$, we show that the primal variables

$$q_k^* = 1, \quad \forall k, \quad (47)$$

and the dual variables

$$\lambda^* = 0 \quad (48)$$

$$\lambda_k^* = s_{\max} f_k, \quad \forall k. \quad (49)$$

To see this, plugging the expressions of λ_k^* and λ^* into (21) yields

$$q_k^* = 1, \quad (50)$$

which is exactly q_k^* in (34). Thus, conditions in (21), (23) and (25) hold. As $\lambda^* = 0$, condition in (22) holds. Finally, we have

$$\sum_{k=1}^K f_k \cdot q_k^* \cdot \phi_k = \sum_{k=1}^K f_k \phi_k \leq \bar{B}, \quad (51)$$

which means condition in (24) holds. In conclusion, we have shown the optimality of q_k^* in (47). \square

From discrete to continuous costs: The continuous cost can be considered as a special case of discrete cost by setting the number of discrete costs K to be infinity. To transform discrete solution to continuous solution, we replace summations with integrals considering K goes to infinity as follows:

$$\sum_{k=1}^K f_k \sqrt{\phi_k} \Rightarrow \int_{c_{\min}}^{c_{\max}} f(c) \sqrt{\phi(c)} dc = \mathbf{E}_c[\sqrt{\phi(c)}], \quad (52)$$

$$\sum_{k=1}^K f_k \phi_k \Rightarrow \int_{c_{\min}}^{c_{\max}} f(c) \phi(c) dc = \mathbf{E}_c[\phi(c)], \quad (53)$$

$$\begin{aligned} H(m) &= \sum_{k=1}^m f_k \phi_k + \sum_{k=m+1}^K f_k \sqrt{\phi_k} \cdot \sqrt{\phi_{m+1}} \\ &\Rightarrow H(x) = \mathbf{E}_c[\phi(c) \cdot \mathbb{1}\{c \leq x\}] + \mathbf{E}_c[\sqrt{\phi(c)} \cdot \mathbb{1}\{c > x\}] \end{aligned} \quad (54)$$

Finally, we are able to derive the continuous solution in Theorem 2. \square

Reference

[1] Xiang Li, Kaixuan Huang, Wenhao Yang, Shusen Wang, and Zhihua Zhang. On the convergence of fedavg on non-iid data. In *International Conference on Learning Representation*, 2019.