Modeling Redundance-based Routing in Delay Tolerant Networks

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Abstract

Delay Tolerant Networks (DTNs) are usually intermittently connected mobile wireless networks. Routing in delay tolerant networks is a challenging problem because DTNs have their own unique features and some well-known assumptions of traditional networks are no longer true in DTNs. Various routing schemes have been proposed in literature which adopt message redundance to fasten the message delivery and control the routing overhead. Despite the number of different proposed routing schemes, there lacks a systematical analytic work on modeling the redundance based routing schemes for delay tolerant networks. This motivates our work in thinking the redundance based routing from the analytical perspective of view. In this paper, we model the message delivering in DTNs as a continuous time Markov random process with absorbing state. With this model, we have derived the probability distribution functions of the message delivery delay when different routing schemes are used, including direct transmission routing, source forwarding, and binary forwarding schemes. We have conducted simulations to validate our model and our analytic results. The simulations show that our model can precisely capture the features of those redundance based routing schemes.

1   Introduction

Delay Tolerant Networks (DTNs) are intermittently connected mobile wireless networks in which the connectivity between nodes changes frequently due to nodes movement. Examples of this kind of network include inter-planetary networks [1] [2], wildlife tracking and habitat monitoring sensor networks [3] [4], and etc. In general, the DTN routing works in an extended store-and-forward manner. The messages are sent between two nodes when they encounter and a wireless link is established. The nodes may cache the messages for considerably long time before getting the chance to send them to the next hop nodes.

Because the exact contact information between any two nodes in DTN is usually not known in advance, making the right routing decisions is challenging. One simple and straightforward method is using flooding based epidemic routing. Flooding based routing schemes can cause too much overhead for the network and overflow the buffer of the nodes. In order to control the overhead of flooding based routing schemes, the authors of [5] propose to dispatch a certain number of identical message copies to a fixed number of relay nodes,

instead of sending to every node in the network. When a message is generated by a source node, a "quota" is attached to that message, which represents how many identical copies of the message can be inserted in the network. When a node carrying the message forwards a copy to another node, the remaining quota of the message is distributed between the two copies on the two different nodes. The node will not forward the message to somebody else if the carried message does not have enough quota.

To the best of our knowledge, there is little systematic analytic work on modeling and analyzing redundance based routing, although lots of routing schemes have been proposed. In this paper, we propose a model using continuous time Markov chain with absorbing state to study the performance of redundance based routing in delay tolerant networks. The non-absorbing state of the Markov chain is the number of relays carrying copies of the same message. When the message is successfully delivered to its destination node, the Markov chain enters its absorbing state. The transiting rate between different states depends on the message forwarding schemes and we will discuss details of the state transfer rate for different message forwarding schemes in Section 4.

The rest of this paper is organized as follows. In Section 2 we present a brief overview of the previous works related to this paper. In Section 3 we present our model in detail. We use the Markov model to analyze different redundance based routing schemes in Section 4. The simulation results are presented in Section 5. We conclude this paper in Section 6.

2   Related Work

2.1   Deterministic Routing

Delay tolerant network is currently attracting significant attention. There have been many research activities in the area of DTN routing, which is a fundamental issue in DTNs [6]. The knowledge-based routing [7] formulates the DTN routing problem as moving messages across a graph with time-varying connectivity whose dynamics is predictable. Merugu et al assume the network profile can be accurately predicted over certain period of time and propose the space-time routing scheme in [8]. The dynamic of the networks is modeled as *space-time graph*. Finding the best rout for message is done by looking ahead of time. Both the work in [7] and [8] are deterministic routing schemes.

2.2   Opportunistic Routing

The *direct transmission* routing scheme [9] is the simplest and most resource efficient one, in which only the source node holds the message until it meets the destination node. This scheme, although has minimal overhead, may incur very long delays for message delivery. The utility based routing schemes [10] [11] [12] define parameters to represent the capability of a relay node to deliver the message and use that information to assist making the decision of forwarding message.

Single message copy opportunistic routing may have too long delivery delay. To shorten the delay, message redundance based routing is introduced. The flooding routing is proposed in [13] which forwards the message to all nodes in the networks. Flooding routing has the highest probability of delivery (or least latency) and the most overhead in terms of bandwidth and energy. To bound the overhead of delivering a message, Spyropoulos et al propose to spray only a fixed number of message copies into the network [5] [14]. Wang et

al [15] propose to erasure code the messages into large number of blocks to achieve more diversity in forwarding. The scheme proposed in [16] integrates erasure coding into utility based routing. The schemes proposed in [15] and [16] have fixed overhead as well.

To the best of our knowledge, there lacks a systematic analytical work on modeling the redundance based routing, although a number of routing schemes have been proposed. The branching process is used to model the epidemic routing in DTNs [17]. In [14] the authors give upper and lower bound of the message delivery delay in multi-copy message routing scheme. Our work is different from the previous work in that we propose a model which can be used in modeling a wide range of redundance based routing scheme. By this model, we can derive the message delivery delay distribution, which can be a rich source of information for improving the performance of redundance based routing schemes.

3 Modeling Redundance based Routing

3.1 Assumptions and Notations

In order to simplify the analysis, we have a set of assumptions here. We assume that the storage capacity of each node is large enough that there is no cache replacement occurs. For any node, the average arrival rate of contacts is λ. That is to say, each node will meet with another node every $\frac{1}{\lambda}$ seconds on average. We assume that the arrival of contacts is a Poisson random process, i.e. the number of contacts a node meets within t seconds is a random variable $X(t)$ and the distribution of $X(t)$ is $P(X(t) = x) = \frac{(\lambda t)^x}{x!}e^{-\lambda t}$, where λ is the contact arrival rate and we assume all nodes in the network have the same contact arrival rate. The Poisson arrival assumption is reasonable since it is widely used in modeling the arrival of events. Although we assume the contact arrival is Poisson distribution here, we'd like to emphasize that our model does not depend on this assumption, as we will show later.

For the convenience of reading, we summarize notations used in this paper in Table 1.

Table 1: Notations Used in This Paper

λ:	average contact arrival rate
N:	number of nodes in the network
M:	maximum number of relay nodes
a_i, b_i	state transfer rates
$F_i(t)$:	the state probabilities
$f_i(t)$:	the state probability distributions
$\overline{f}(s)$	Laplace transfer of function $f(t)$
$P(D|k)$:	the number of ways to partition D into k integers

3.2 The Model

We consider a DTN with N mobile nodes. M is the maximum number of copies a message can have and $M \leq N$. At the very beginning, the source node generates a message and the source is the only carrier of that message. When the source node meets with other nodes, the source node will forward message copies to them. Other relay nodes may also

forward the message and introduce more relay nodes carrying the message. For the message in concern, we define the number of relay nodes carrying that message as the "state" of the system. Obviously, the number of relay nodes carrying the message can vary from 1 to M. We use S_i to denote the state of the network system in which there are i relay nodes carrying the message in concern. There is another state, S_{end}, in the system, which represents the state that one relay node carrying the message has forwarded the message to the destination node and the message transmission is finished.

When the network system is in state S_i, it can transfer to state S_{i+1}, which means a relay node has forwarded a message copy to another node who does not have that message. The state S_i can also transfer to state S_{end}, in case that one of the i relay nodes carrying the message has met with the destination node and the message transmission is finished. We show the state transferring in Fig. 1. In this figure, the maximum number of relay nodes is M and there are $M+1$ states, including the S_{end} state. The parameter a_i is the transfer rate from state S_i to S_{i+1} and b_i is the transfer rate from state S_i to state S_{end}. The random process shown in Fig. 1 is a continuous time Markov chain with absorbing state [18]. Once the transfer rates between states are known, we can have transient analysis of this Markov chain.
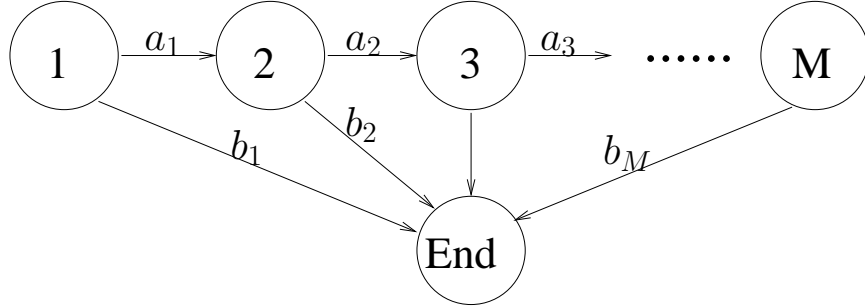

Figure 1: The continuous time Markov model representing the message delivery in redundance-based routing schemes.

We are particularly interested in the probability distribution of state S_{end}, $f_{end}(t)$, which is the distribution of the message delivery delay. Let $F_i(t)$ and $F_{end}(t)$ denote the probability that at time t the system is in state S_i and S_{end}, respectively. The initial state of the Markov chain is S_1; that is $F_1(0) = 1$, $F_i(0) = F_{end}(0) = 0$ for $i \in [2, M]$. Following Kolmogorov's equation, we have the following differential equations about the state probabilities $F_{end}(t)$ and $F_i(t)$ for $i \in [1, M]$.

$$
\begin{cases}
\frac{dF_1}{dt} & = -(a_1 + b_1)F_1 \\
\cdots\cdots \\
\frac{dF_i}{dt} & = -(a_i + b_i)F_i + a_{i-1}F_{i-1} \\
\cdots\cdots \\
\frac{dF_M}{dt} & = -b_M F_M + a_{M-1}F_{M-1} \\
\frac{dF_{end}}{dt} & = \sum_{i=1}^{M} b_i F_i
\end{cases}
\tag{1}
$$

Given $F_1(0) = 1$ and $F_i(0) = F_{end}(0) = 0$ for $i \in [2, M]$, we can solve this system using Laplace transforms, as shown in (2), in which $\overline{F}_i(s)$ is the Laplace transfer of $F_i(t)$,

$\overline{F}_i(s) = \mathcal{L}(F_i(t))$.

$$\begin{cases} s\overline{F}_1 - 1 = -(a_1 + b_1)\overline{F}_1 \\ \ldots\ldots \\ s\overline{F}_i = -(a_i + b_i)\overline{F}_i + a_{i-1}\overline{F}_{i-1} \\ \ldots\ldots \\ s\overline{F}_M = -b_M\overline{F}_M + a_{M-1}\overline{F}_{M-1} \\ s\overline{F}_{end} = \sum_{i=1}^{M} b_i\overline{F}_i \end{cases} \qquad (2)$$

If we define $a_0 = 1$ and $a_M = 0$, we can derive $\overline{F}_i$ and $s\overline{F}_{end}$ as the following.

$$\begin{cases} \overline{F}_i = \prod_{j=1}^{i} \frac{a_{j-1}}{s+a_j+b_j} \\ s\overline{F}_{end} = \sum_{i=1}^{M} b_i \prod_{j=1}^{i} \frac{a_{j-1}}{s+a_j+b_j} \end{cases}$$

Since $F_{end}(t)$ is the CDF of message transmission delay, we can get the PDF of the delay as $\frac{dF_{end}(t)}{dt} = f_{end}(t)$. It follows that $\overline{f}_{end}(s) = s\overline{F}_{end}(s) - \overline{F}_{end}(0) = s\overline{F}_{end}(s)$. Considering the initial state $\overline{F}_{end}(0) = 0$, $\overline{f}_{end}(s)$ can be written as:

$$\overline{f}_{end}(s) = \sum_{i=1}^{M} b_i \prod_{j=1}^{i} \frac{a_{j-1}}{s + a_j + b_j} \qquad (3)$$

The inverse Laplace transform of (3) is the convolution of a series of exponential functions, as shown in (4), in which the symbol $\otimes$ represents convolution operation.

$$f_{end}(t) = \sum_{i=1}^{M} b_i \bigotimes_{j=1}^{i} a_{j-1}e^{-(a_j+b_j)t} \qquad (4)$$

Function $f_{end}(t)$ is easy to compute because the convolution of two exponential functions can be computed as the following, when a is not equal to b.

$$\begin{aligned} e^{-at} \otimes e^{-bt} \quad &= \int_0^t e^{-au}e^{-b(t-u)}du \\ &= e^{-bt}\int_0^t e^{(b-a)u}du \\ &= e^{-bt}\frac{1}{b-a}e^{(b-a)u}|_0^t \\ &= \frac{e^{-bt}}{b-a}(e^{(b-a)t} - 1) \\ &= \frac{e^{-at}-e^{-bt}}{b-a} \end{aligned}$$

In case of $a = b$, the convolution of K exponential function is

$$\bigotimes_{i=1}^{K} e^{-at} = \frac{t^{K-1}}{(K-1)!}e^{-at}$$

4 Using the Model to Analyze Different Forwarding Schemes

4.1 Analysis of Direct Transmission

In the direct transmission scheme, the source node will carry the message until it meets with the destination node. That is to say the source will not forward the message to any

other node except the destination node. Because there are $(N-1)$ other nodes in the network but only one of them is the destination of the message, any contact of the source node has a $\frac{1}{N-1}$ chance to be the destination node. Under the assumption that the contact arrival is a Poisson process with arrival rate λ, the probability distribution of the delivery delay is exponential distribution.

$$f_{end}(t) = \frac{\lambda}{N-1} e^{-\frac{\lambda t}{N-1}} \tag{5}$$

We can also derive (5) from the differential equation set in (1). In this case, there are only two states in the Markov chain, state S_1 and S_{end}. The state transfer rate from S_1 to S_{end} is $\frac{\lambda}{N-1}$. We have the following two differential equations:

$$\begin{cases} \frac{dF_1}{dt} = -\frac{\lambda}{N-1} F_1 \\ \frac{dF_{end}}{dt} = \frac{\lambda}{N-1} F_1 \end{cases}$$

By the Laplace transforms, we have $s\overline{F}_{end} = \frac{\lambda}{N-1} \frac{1}{s+\frac{\lambda}{N-1}}$. So the delay distribution is $f_{end}(t) = \frac{\lambda}{N-1} e^{-\frac{\lambda}{N-1}t}$.

4.2 Analysis of Source Forward

In source forward scheme, only the source of the message can forward copies of that message to other relay nodes. After receiving a message copy from the source, the relay node will not send that message copy to any node except the destination. Before we begin the analysis source forward routing scheme, we'd like to answer the following question: when there are I nodes carrying the message, how long it takes for the first one to meet with the destination node.

For any relay node carrying the message, the average arrival rate of the destination is $\frac{1}{N-1}\lambda$. Since the contact arrival is Poisson distribution, the inter-arrival time of contact is exponential distribution $f(t) = \frac{\lambda}{N-1} e^{-\frac{\lambda}{N-1}t}$. We consider the message delivery delay when there are I relay nodes carrying the message. The time needed for relay node i to meet with the destination node is an exponential distributed random variable Φ_i. The time needed for the first one of the I relay nodes to encounter the destination node is $\min[\Phi_1, ...\Phi_I]$. If we assume Φ_i for $i \in [1, I]$ are IID (independent and identical distributed) random variables, we can derive the distribution of $\min[\Phi_1, ...\Phi_I]$ as following [1]:

$$f_{\min[\Phi_1, ...\Phi_I]}(t) = \frac{I\lambda}{N-1} e^{-\frac{I\lambda}{N-1}t} \tag{6}$$

From (6) we see that $\min[\Phi_1, ...\Phi_I]$ is an exponential random variable, whose mean is $\frac{I\lambda}{N-1}$. So when there are i relay nodes carrying the message copies, the transfer rate from state S_i to state S_{end} is:

$$r(S_i, S_{end}) = b_i = \frac{i}{N-1}\lambda \tag{7}$$

---

[1]The distribution can be calculated like this: $P(\min[\Phi_1, ...\Phi_I] \leq x) = 1 - P(\min[\Phi_1, ...\Phi_I] > x) = 1 - e^{-\frac{I\lambda}{N-1}x} \Rightarrow f_{\min[\Phi_1, ...\Phi_I]}(t) = \frac{I\lambda}{N-1} e^{-\frac{I\lambda}{N-1}t}$

When there are i relay nodes in the network carrying the message and the $i < M$, the state transfer rate from S_i to S_{end} is $\frac{i}{N-1}\lambda$. The transfer rate from S_i to state S_{i+1} is $\frac{N-i-1}{N-1}\lambda$, because the source node will not forward message copies to those node already having that message and the probability that the source node encounters a node who does not have the message is $\frac{N-i-1}{N-1}$, since there are $(N-i-1)$ nodes in the network can be potential relay nodes of the message.

$$r(S_i, S_{i+1}) = a_i = \frac{N-i-1}{N-1}\lambda \tag{8}$$

Replacing a_i and b_i in (1) with (7) and (8). We can derive the distribution of delay in source forward as shown in (9).

$$
\begin{cases}
f_{end}(t) = \sum_{i=1}^{M} \bigotimes_{j=1}^{i} a_{j-1} e^{-(a_j+b_j)t} \\
a_0 = 1, a_M = 0 \\
a_i = \frac{N-i-1}{N-1}\lambda, \quad i \in [1, M-1] \\
b_i = \frac{i}{N-1}\lambda, \quad i \in [1, M]
\end{cases}
\tag{9}
$$

Note $a_i + b_i = 1$ for $i \in [1, M-1]$, we can write $f_{end}(t)$ as:

$$
\begin{aligned}
f_{end}(t) = \quad & \sum_{i=1}^{M-1} b_i \left(\prod_{j=1}^{i} a_{j-1} \right) \frac{t^{i-1}}{(i-1)!} e^{-\lambda t} \\
& + \frac{b_M \prod_{j=1}^{M} a_{j-1}}{(M-2)!} t^{M-2} e^{-\lambda t} \bigotimes e^{-\frac{M}{N-1}\lambda t}
\end{aligned}
$$

The convolution of $t^{M-2} e^{-\lambda t}$ and $e^{-\frac{M}{N-1}\lambda t}$ is:

$$
\begin{aligned}
& t^{M-2} e^{-\lambda t} \bigotimes e^{-\frac{M}{N-1}\lambda t} \\
&= e^{-\frac{M}{N-1}\lambda t} \int_0^t u^{M-2} e^{-\frac{N-M-1}{N-1}\lambda u} du \\
&= e^{-\frac{M}{N-1}\lambda t} \left(\frac{(M-2)!}{(\frac{N-M-1}{N-1}\lambda)^{M-1}} - \frac{e^{-\frac{N-M-1}{N-1}\lambda t}}{(\frac{N-M-1}{N-1}\lambda)^{M-1}} \sum_{i=0}^{M-2} (\frac{N-M-1}{N-1}\lambda t)^i \frac{(M-2)!}{i!} \right) \\
&= \frac{e^{-\lambda t}}{(\frac{N-M-1}{N-1}\lambda)^{M-1}} \sum_{i=0}^{M-2} (\frac{N-M-1}{N-1}\lambda t)^i \frac{(M-2)!}{i!} \\
&= \frac{e^{-\frac{M}{N-1}\lambda t}(M-2)! - e^{-\lambda t} \sum_{i=0}^{M-2} (\frac{N-M-1}{N-1}\lambda t)^i \frac{(M-2)!}{i!}}{(\frac{N-M-1}{N-1}\lambda)^{M-1}} \\
&= \frac{(M-2)!(e^{-\frac{M}{N-1}\lambda t} - e^{-\lambda t} \sum_{i=0}^{M-2} \frac{(\frac{N-M-1}{N-1}\lambda t)^i}{i!})}{(\frac{N-M-1}{N-1}\lambda)^{M-1}}
\end{aligned}
$$

4.3 Analysis of Binary Forward

Binary forward is first proposed in [5]. When a relay node having more them one message quota meets with another node who does not have a copy of that message, they will share the message quotas and each of them will have half of the quotas. Similarly, the transfer rate from state S_i to state S_{end} is still the following:

$$r(S_i, S_{end}) = b_i = \frac{i}{N-1}\lambda$$

In order to derive the transfer rate from state S_i to S_{i+1}, we need to consider the following question: When there are i nodes in the network carrying the message copies (there are i relays for a message), how many of them have the ability to forward message

copy to other nodes. We approximate this problem using the results of integer partition [19], which is a well studied problem in combinatorics. Let $P(D|k)$ denote the number of different partitions when an integer D is partitioned into k parts. We have the following recursive function about $P(D|k)$ [19]:

$$P(D|k) = \sum_{i=1}^{k} P(D-k|i)$$

$P(D|k)$ can be recursively computed by noticing that

$$\begin{cases} P(D|k) = 1, & D = k \\ P(D|k) = 0, & D < k \end{cases}$$

Suppose there are $P(D|k)$ different ways to partition D into k parts, $D = x_1 + x_2 + ... + x_k$, $x_1 \geq x_2 \geq ...x_k \geq 1$. We want to know how many ways we have when we partition D into k parts and s out of the k parts are larger than 1. That is to say $D - k = (x_1 - 1) + (x_2 - 1) + ... + (x_k - 1) = x_1' + x_2' + ... + x_s'$, $x_1' \geq x_2' \geq ...x_s' \geq 1$. To partition $(D-k)$ into s parts, there are $P(D-k|s)$ different ways. Therefore, when we partition D into k parts, on average there are $\frac{\sum_{s=1}^{k} sP(D-k|s)}{P(D|k)}$ parts larger than 1. Back to our problem in binary forwarding, when there are k relay nodes carrying message copies, we can use the $P'(M|k)$ in (10) to approximate the average number of nodes who can forward their message copy to other relay nodes who do not have that message.

$$P'(M|k) = \frac{\sum_{s=1}^{k} sP(M-k|s)}{P(M|k)} \tag{10}$$

The transfer rate from state S_i to state S_{i+1} can be written as the following:

$$r(S_i, S_{i+1}) = a_i = P'(M|i)\frac{N-i-1}{N-1}\lambda \tag{11}$$

For example, $5 = 1 + 1 + 3 = 1 + 2 + 2$, so $P(5|3) = 2$. The average number of parts which are larger than or equal to 2 should be $1 \times 0.5 + 2 \times 0.5 = 1.5$. According to (10), the result is $\frac{1 \times P(2|1) + 2 \times P(2|2) + 3 \times P(2|3)}{P(5|3)} = \frac{1 \times 1 + 2 \times 1 + 3 \times 0}{2} = 1.5$.

Now we have all the transfer rates. Using the same technique in Section 4.2 we can get the distribution of the message delay in case of binary forward $\overline{f}_{end}(t)$ as shown in (12).

$$\begin{cases} f_{end}(t) = \sum_{i=1}^{M} b_i \bigotimes_{j=1}^{i} a_{j-1} e^{-(a_j + b_j)t} \\ a_0 = 1, a_M = 0 \\ a_i = P'(M|i)\frac{N-i-1}{N-1}\lambda, & i \in [1, M-1] \\ b_i = \frac{i}{N-1}\lambda, & i \in [1, M] \end{cases} \tag{12}$$

4.4 Analysis of Flooding

In flooding based routing, each node will forward a copy of the message to his contact if the contact does not have that message in his cache. When there are i nodes in the network carrying the message, the state transfer rate from S_i to S_{end} is still $\frac{i}{N-1}\lambda$. When the system is in state S_i, there are i nodes who can forward message copies to other $N - 1 - i$ nodes without the message. Therefore, there transfer rate from state S_i to S_{i+1} is $i\frac{N-1-i}{N-1}\lambda$.
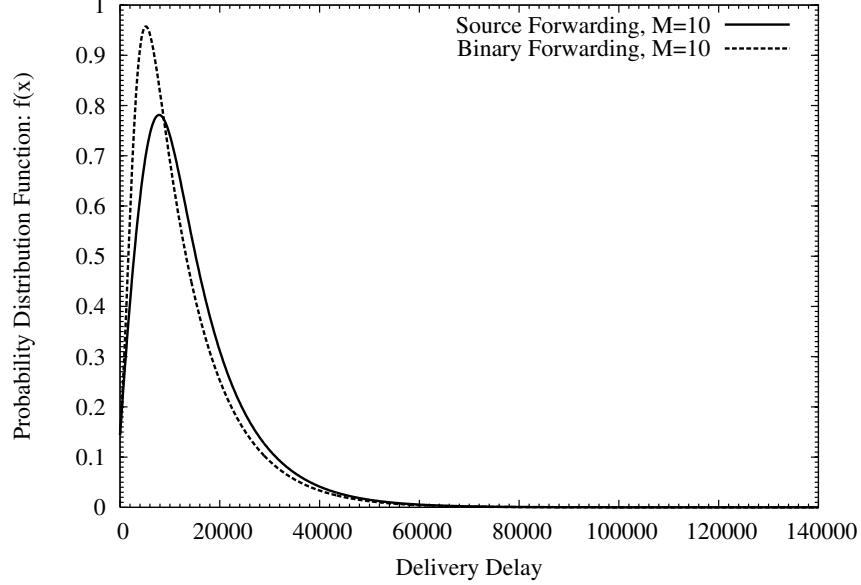
Figure 2: Compare between source forwarding and binary forwarding schemes.

Since there are $N + 1$ states in the system, i.e. state $S_1 \sim S_N$ and S_{end}, we can write the probability distribution function of the message delivery delay as the following:

$$
\begin{cases}
\overline{f}_{end}(s) = \sum_{i=1}^{N} b_i \prod_{j=1}^{i} \frac{a_{j-1}}{s+a_j+b_j} \\
f_{end}(t) = \sum_{i=1}^{N} b_i \bigotimes_{j=1}^{i} a_{j-1} e^{-(a_j+b_j)t} \\
a_0 = 1, a_N = 0 \\
a_i = i \frac{N-1-i}{N-1} \lambda, \quad i \in [1, N-1] \\
b_i = \frac{i}{N-1} \lambda, \quad i \in [1, N]
\end{cases}
\tag{13}
$$

4.5 Discussions

We plot an example of source forward and binary forwarding schemes in Fig. 2 according to (9) and (12), where the maximum number of message copies is set to 6. From the plot in Fig. 2 we can see that the distribution of delivery delay is different from the exponential distribution, although the contact arrival is exponential distribution. It seems that even multiple message copies are generated in the network, the probability that the message being quickly delivered to the destination is still small. For source and binary forward, the probability distribution functions start from a low value and increase rapidly. The probability distribution functions decrease after reaching their maximum values. On the contrary, the exponential distribution is monotonically decreasing. We can also see from Fig. 2 that after reaching the peak value, the source forwarding and binary forwarding schemes do not have too much difference, which means that how the message copies are distributed to relay nodes has more impact on quickly delivering the message to its destination.

5 Experiment Evaluation

5.1 Simulation Setup

We use the random way point mobility model [20] in our simulation experiments. Existing works have shown that the contact inter-arrival time in this mobility model is exponential distribution [21]. We have 100 nodes move in a disc with radius set to 1 unit and there is no node failure. All nodes have the same moving speed and the same communication range. The initial positions of nodes are randomly chosen in the disc. There is no thinking time after nodes arrive at the way points. We set 10000 different random seeds to run the simulation program so we have 10000 simulation instances. We use the message delivery delay in those results to plot the probability distribution of the delay.
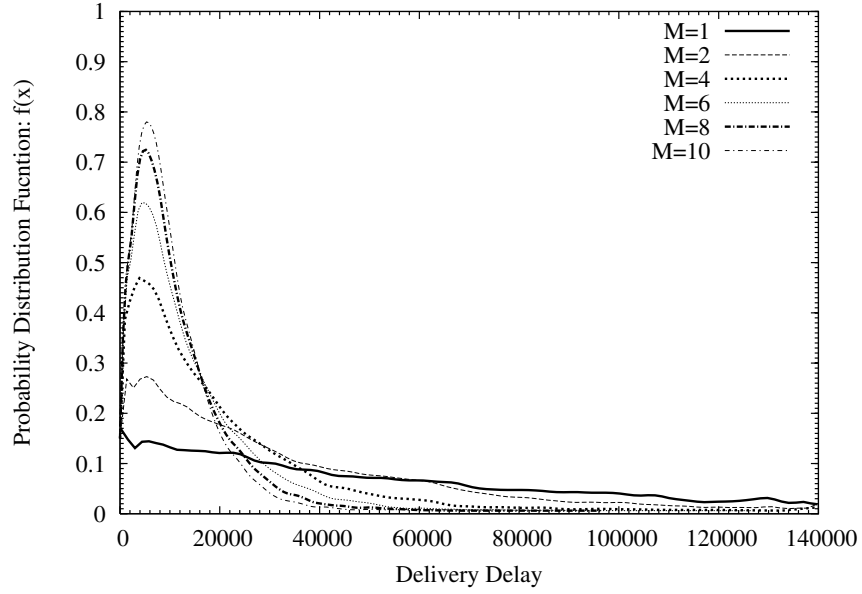
5.2 Simulations of Source Forward

We set the number of maximum copies of a message, M, to different values and run the simulation for each value for 10000 times. The simulation results are plotted in Fig. 3(a). We also plot the probability distributions derived from the analytic result in Fig. 3(b). From the plots in Fig. 3 we can see that the probability distributions derived from our analytical model have almost the same shapes as the simulation results. When the parameter M is 1, the message forwarding is equivalent to the direct transmission scheme and the message delivery delay is an exponentially distributed function. When multiple copies of a message are generated into the network, the message delivery delay function has a hill-like shape.

Both the analytical and simulated results demonstrate that the probability distribution of message delivery delay is not a monotonic function. The probability distribution function starts from a low value and increases to the peak value. After that, the probability distribution function decreases similar to an exponential function. The insight of the probability distribution function can help us in the design of redundance based routing scheme in DTNs, such as setting the appropriate lifetime for a message.

5.3 Simulations of Binary Forward

Fig. 4(a) shows the simulation results of the delay probability distribution when binary forwarding is used. The analytical results are shown in Fig. 4(b) as well. From the plots in Fig. 4 we can see that the probability distributions in case of binary forwarding still have hill-like shapes when multiple copies are used. But the curves increase more rapidly before reaching their highest points. After comparing Fig. 3 and 4 we can see that how the source distributes the message copies has significant impact on message delivery, especially when the maximum number of message copies, M, is set to a relatively high value. In order to achieve fast message delivery, we should carefully design how to distribute the message copies. In those networks where all nodes are identical, binary forwarding is optimal. But for networks where nodes have different attributions, such as moving velocity and buffer capacity, we may need more sophisticated schemes. We leave this issue as the future work of this paper.
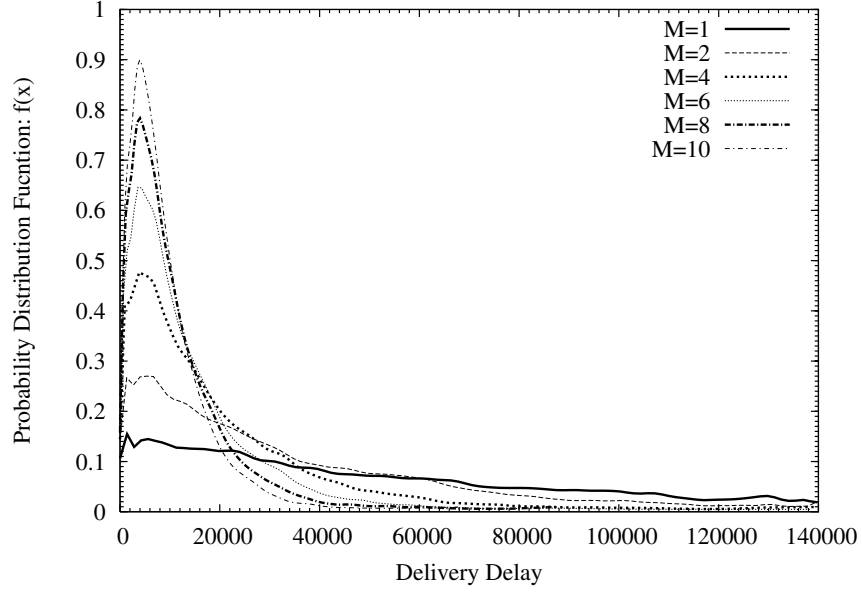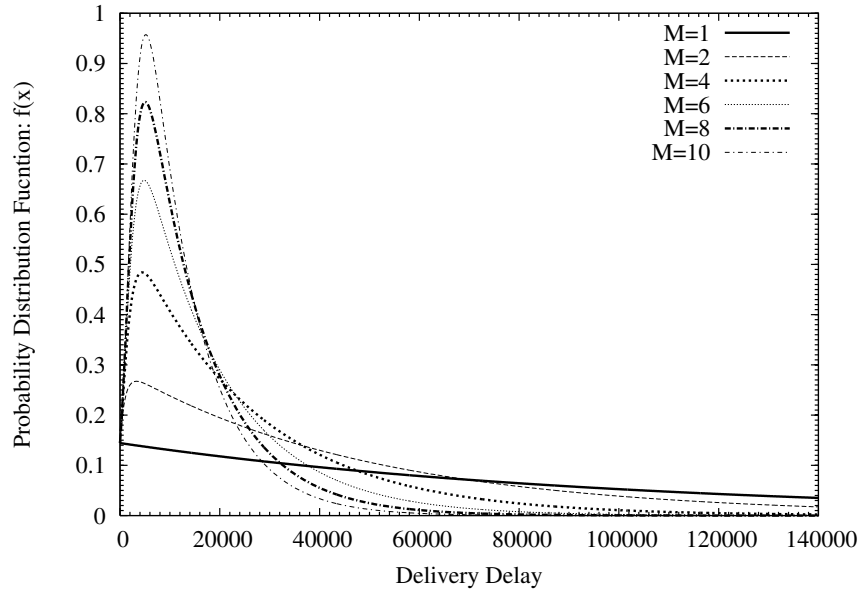
(a) simulation result
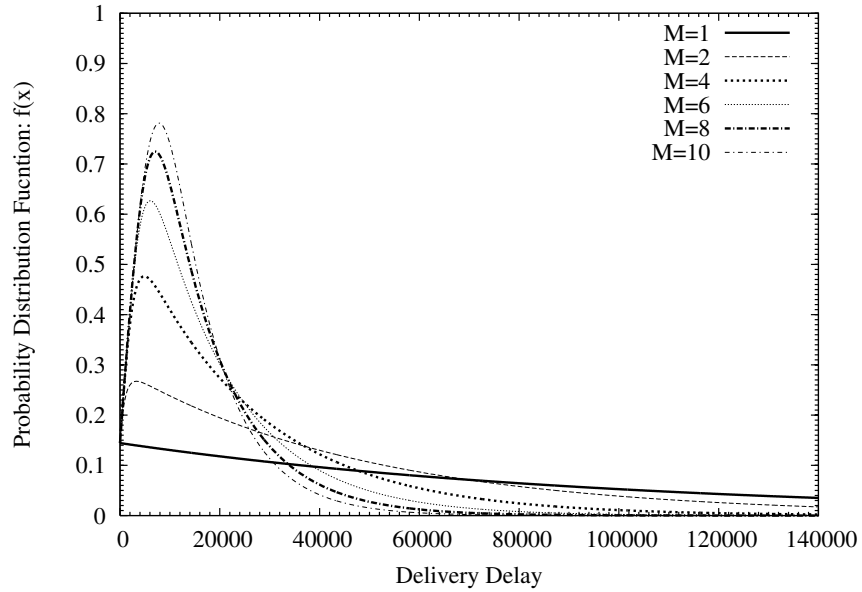
(b) analytic result

Figure 3: Probability distribution of message delivery delay in case of source forwarding.

(a) simulation result

(b) analytic result

Figure 4: Probability distribution of message delivery delay in case of binary forwarding.

6 Conclusion

In this paper, we have presented a continuous time Markov random process model for modeling the message transmission in redundance based routing schemes for delay tolerant networks. With the help of this model, we derived the probability distribution functions of various DTN routing schemes. We used simulation experiments to verify our model. The results show that our analytical model can precisely capture the features of message delivery in redundance based message forwarding schemes.

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